

*Operating
Systems:
Internals
and Design
Principles*

Chapter 3 Process Description and Control

Seventh Edition
By William Stallings

Operating Systems: Internals and Design Principles

The concept of process is fundamental to the structure of modern computer operating systems. Its evolution in analyzing problems of synchronization, deadlock, and scheduling in operating systems has been a major intellectual contribution of computer science.



WHAT CAN BE AUTOMATED?:
THE COMPUTER SCIENCE AND
ENGINEERING RESEARCH STUDY,
MIT Press, 1980

Summary of Earlier Concepts

- A computer platform consists of a collection of hardware resources
- Computer applications are developed to perform some task
- It is inefficient for applications to be written directly for a given hardware platform
- The OS was developed to provide a convenient, feature-rich, secure, and consistent interface for applications to use
- We can think of the OS as providing a uniform, abstract representation of resources that can be requested and accessed by applications

OS Management of Application Execution

- Resources are made available to multiple applications
- The processor is switched among multiple applications so all will appear to be progressing
- The processor and I/O devices can be used efficiently

Process Elements

- Two essential elements of a process are:

Program code

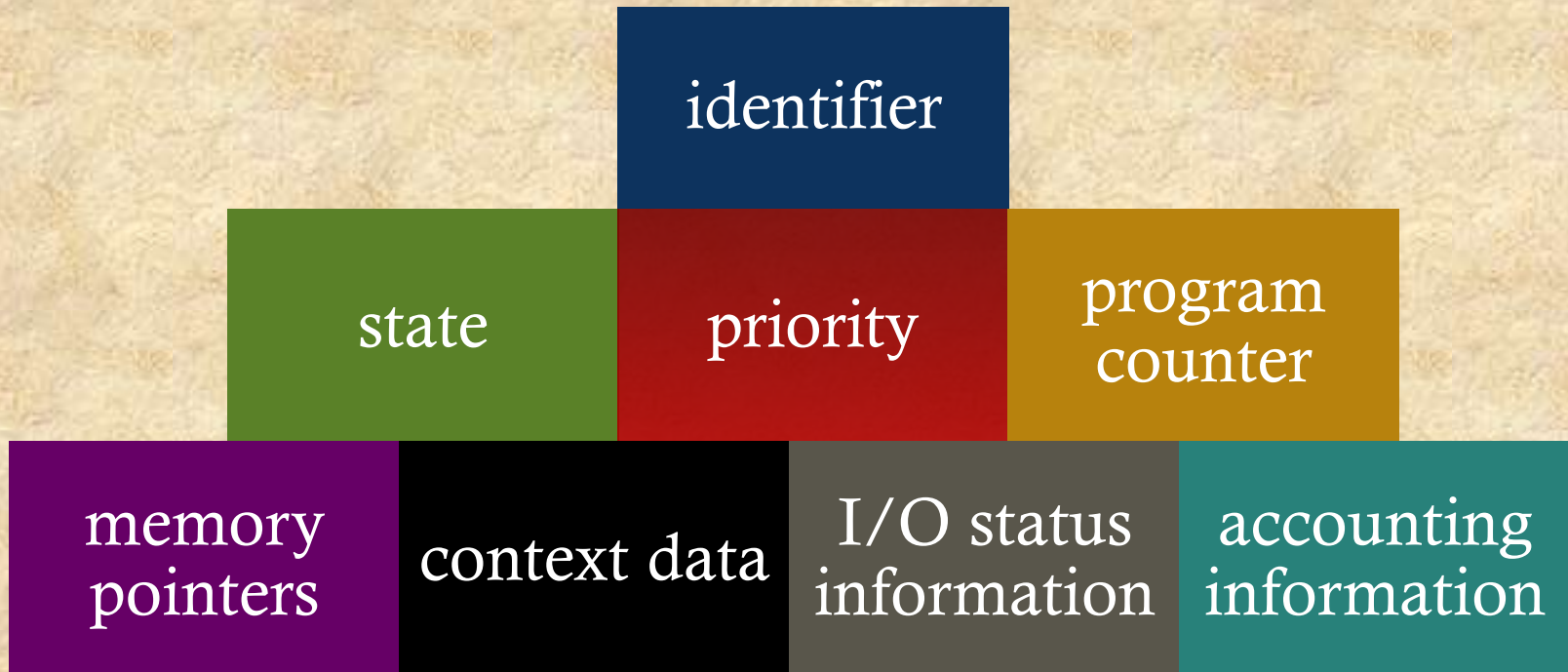
- ⑩ which may be shared with other processes that are executing the same program

A set of data associated with that code

- When the processor begins to execute the program code, we refer to this executing entity as a *process*

Process Elements

- While the program is executing, this process can be uniquely characterized by a number of elements, including:



Process Control Block

- Contains the process elements
- Makes it possible to interrupt a running process and later resume execution as if the interruption had not occurred
- Created and managed by the operating system
- Key tool that allows support for multiple processes (multiprogramming)

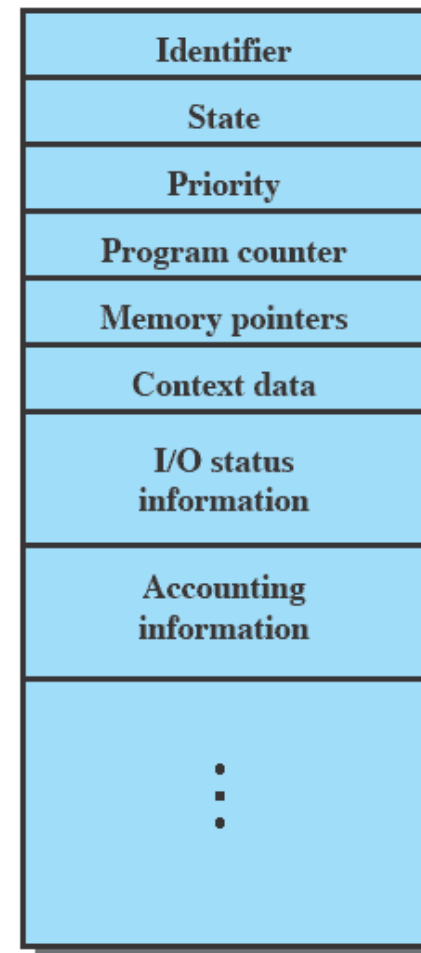


Figure 3.1 Simplified Process Control Block

Process States

Trace

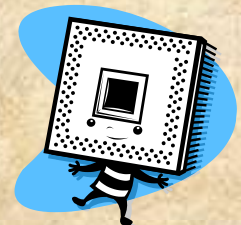
the behavior of an individual *process* can be characterized by listing the sequence of instructions that execute for that process

the behavior of the *processor* can be characterized by showing how the traces of the various processes are interleaved

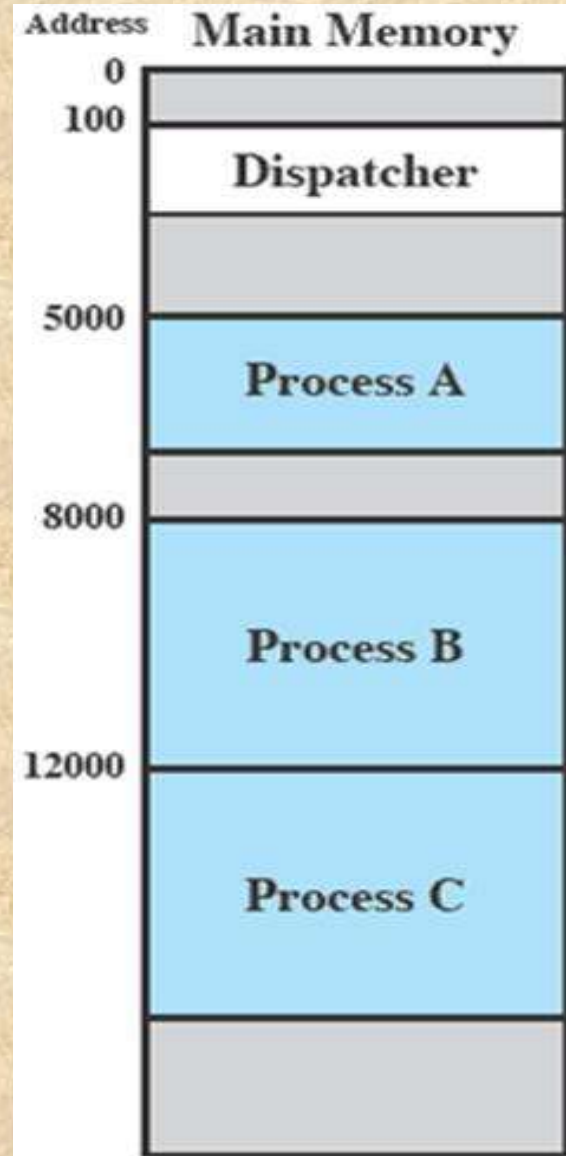


Dispatcher

small program that switches the processor from one process to another



Process Execution



Traces of Processes of Figure 3.2

5000	8000	12000
5001	8001	12001
5002	8002	12002
5003	8003	12003
5004		12004
5005		12005
5006		12006
5007		12007
5008		12008
5009		12009
5010		12010
5011		12011

(a) Trace of Process A

(b) Trace of Process B

(c) Trace of Process C

5000 = Starting address of program of Process A

8000 = Starting address of program of Process B

12000 = Starting address of program of Process C

Figure 3.3

Combined Trace of Processes of Figure 3.2

1	5000				
2	5001				
3	5002				
4	5003				
5	5004				
6	5005				
----- Timeout					
7	100				
8	101				
9	102				
10	103				
11	104				
12	105				
13	8000				
14	8001				
15	8002				
16	8003				
----- I/O Request					
17	100				
18	101				
19	102				
20	103				
21	104				
22	105				
23	12000				
24	12001				
25	12002				
26	12003				
27	12004				
28	12005				
----- Timeout					
29	100				
30	101				
31	102				
32	103				
33	104				
34	105				
35	5006				
36	5007				
37	5008				
38	5009				
39	5010				
40	5011				
----- Timeout					
41	100				
42	101				
43	102				
44	103				
45	104				
46	105				
47	12006				
48	12007				
49	12008				
50	12009				
51	12010				
52	12011				
----- Timeout					

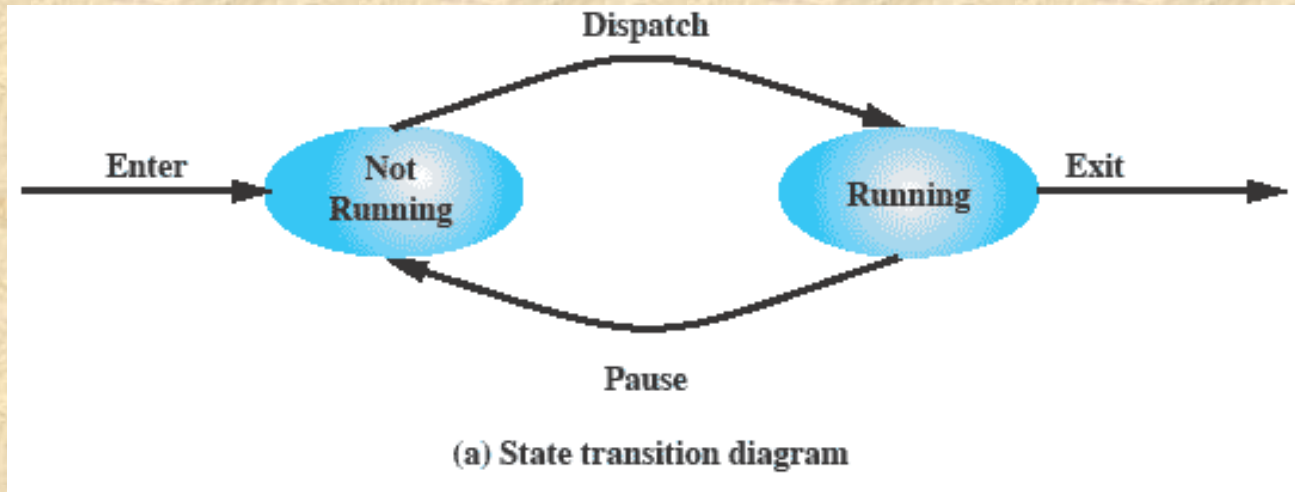
100 = Starting address of dispatcher program

Shaded areas indicate execution of dispatcher process;
first and third columns count instruction cycles;
second and fourth columns show address of instruction being executed

Figure 3.4

Two-State Process Model

- A process may be in one of two states:
 - running
 - not-running



Queuing Diagram

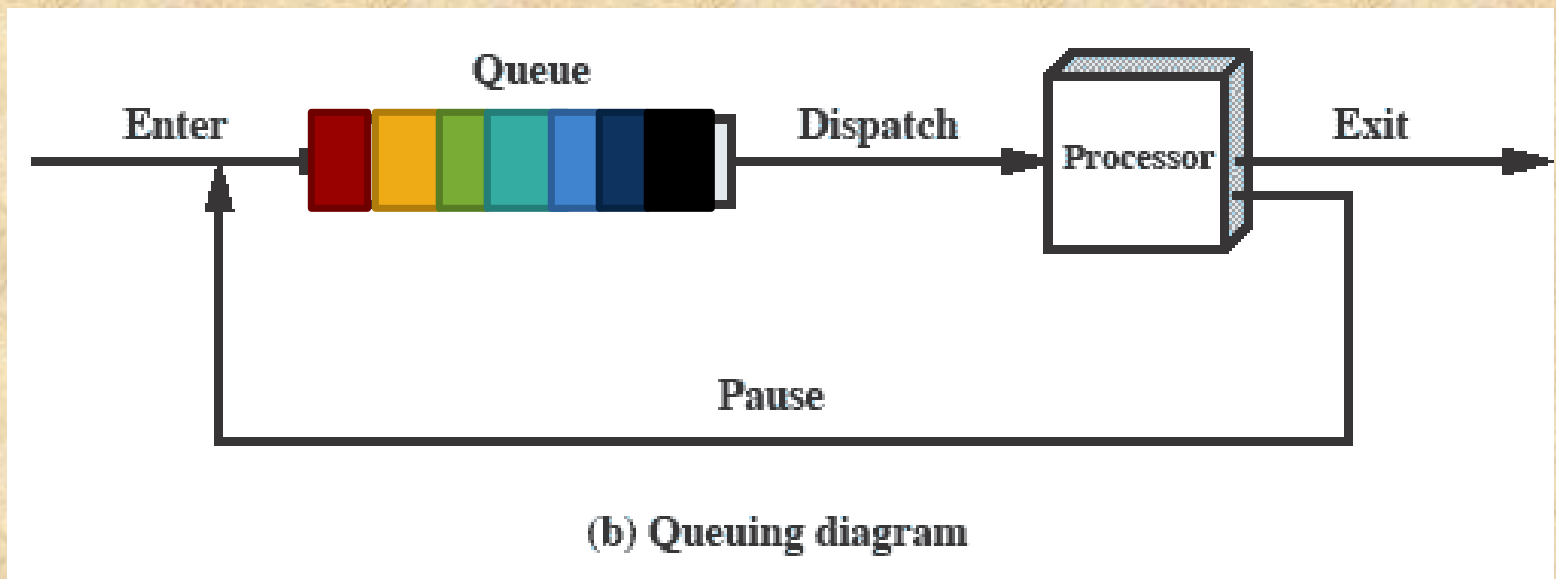
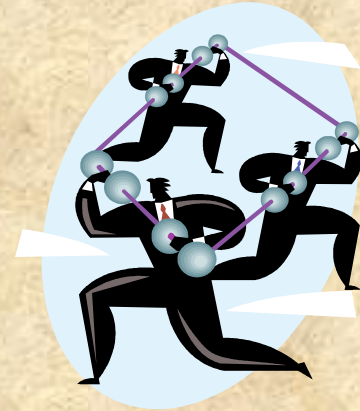


Table 3.1 Reasons for Process Creation

New batch job	The OS is provided with a batch job control stream, usually on tape or disk. When the OS is prepared to take on new work, it will read the next sequence of job control commands.
Interactive logon	A user at a terminal logs on to the system.
Created by OS to provide a service	The OS can create a process to perform a function on behalf of a user program, without the user having to wait (e.g., a process to control printing).
Spawned by existing process	For purposes of modularity or to exploit parallelism, a user program can dictate the creation of a number of processes.

Process Creation

Process spawning

- when the OS creates a process at the explicit request of another process

Parent process

- is the original, creating, process

Child process

- is the new process

Process Termination

- There must be a means for a process to indicate its completion
- A batch job should include a HALT instruction or an explicit OS service call for termination
- For an interactive application, the action of the user will indicate when the process is completed (e.g. log off, quitting an application)



Table 3.2

Reasons for Process Termination



Normal completion	The process executes an OS service call to indicate that it has completed running.
Time limit exceeded	The process has run longer than the specified total time limit. There are a number of possibilities for the type of time that is measured. These include total elapsed time ("wall clock time"), amount of time spent executing, and, in the case of an interactive process, the amount of time since the user last provided any input.
Memory unavailable	The process requires more memory than the system can provide.
Bounds violation	The process tries to access a memory location that it is not allowed to access.
Protection error	The process attempts to use a resource such as a file that it is not allowed to use, or it tries to use it in an improper fashion, such as writing to a read-only file.
Arithmetic error	The process tries a prohibited computation, such as division by zero, or tries to store numbers larger than the hardware can accommodate.
Time overrun	The process has waited longer than a specified maximum for a certain event to occur.
I/O failure	An error occurs during input or output, such as inability to find a file, failure to read or write after a specified maximum number of tries (when, for example, a defective area is encountered on a tape), or invalid operation (such as reading from the line printer).
Invalid instruction	The process attempts to execute a nonexistent instruction (often a result of branching into a data area and attempting to execute the data).
Privileged instruction	The process attempts to use an instruction reserved for the operating system.
Data misuse	A piece of data is of the wrong type or is not initialized.
Operator or OS intervention	For some reason, the operator or the operating system has terminated the process (e.g., if a deadlock exists).
Parent termination	When a parent terminates, the operating system may automatically terminate all of the offspring of that parent.
Parent request	A parent process typically has the authority to terminate any of its offspring.

Five-State Process Model

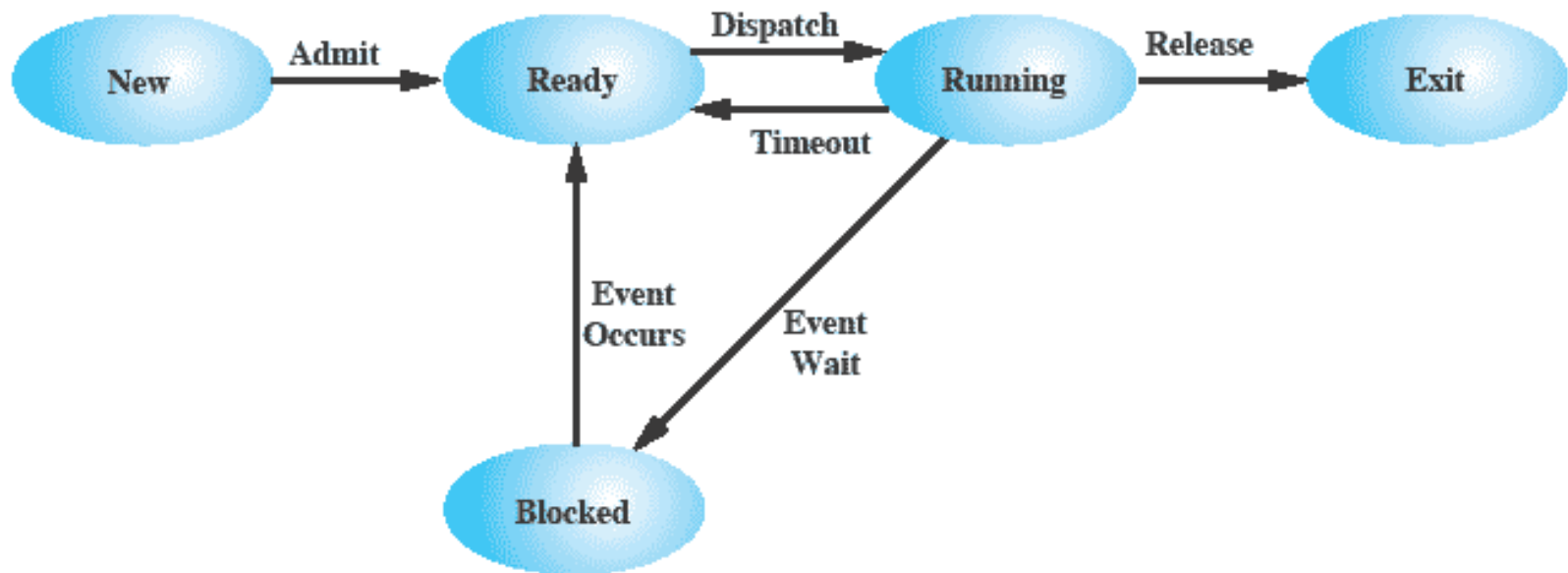


Figure 3.6 Five-State Process Model

Process States for Trace of Figure 3.4

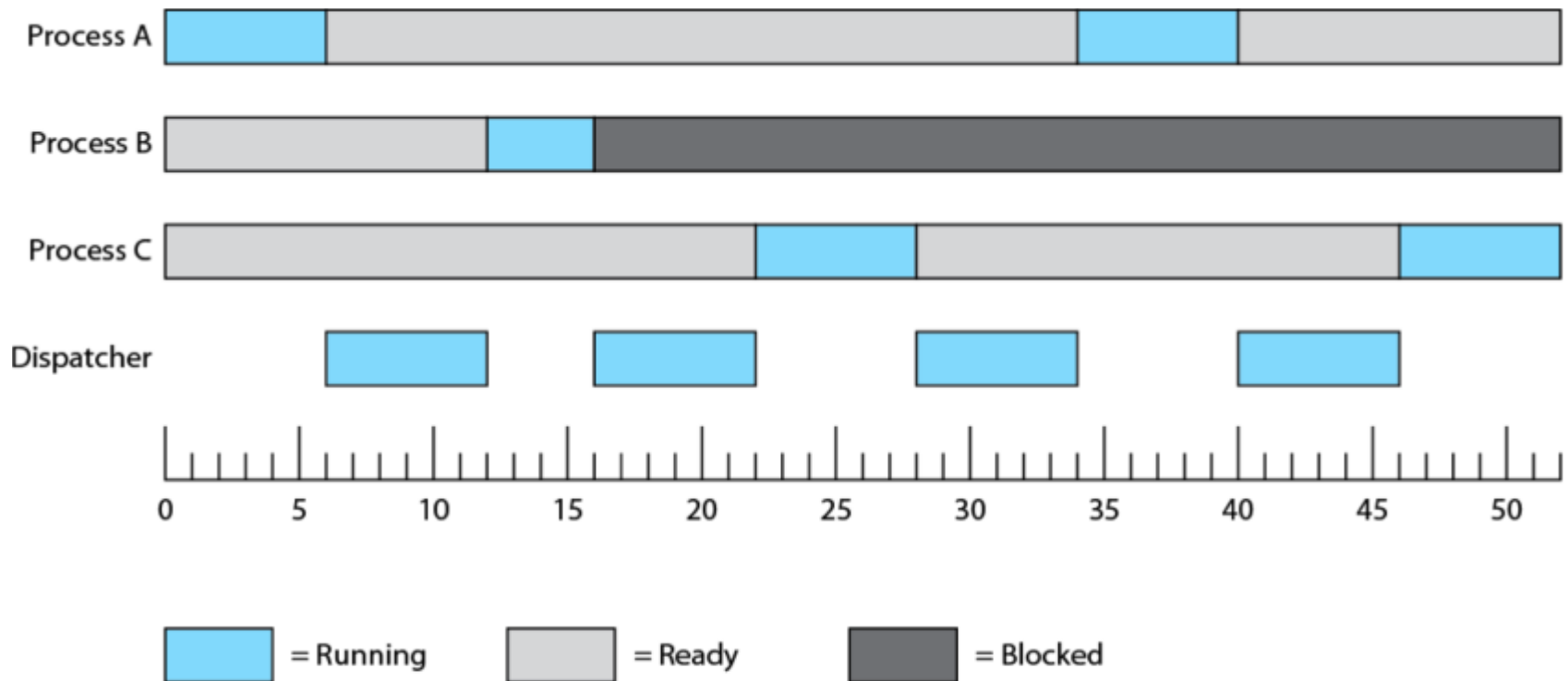
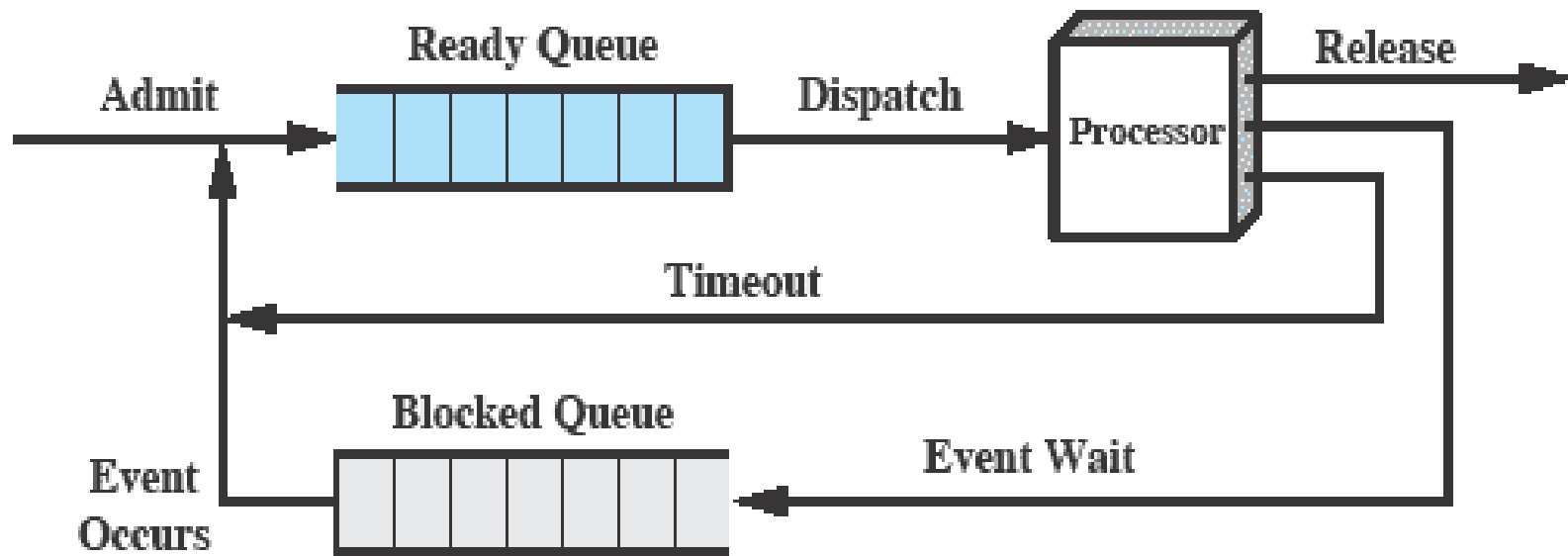


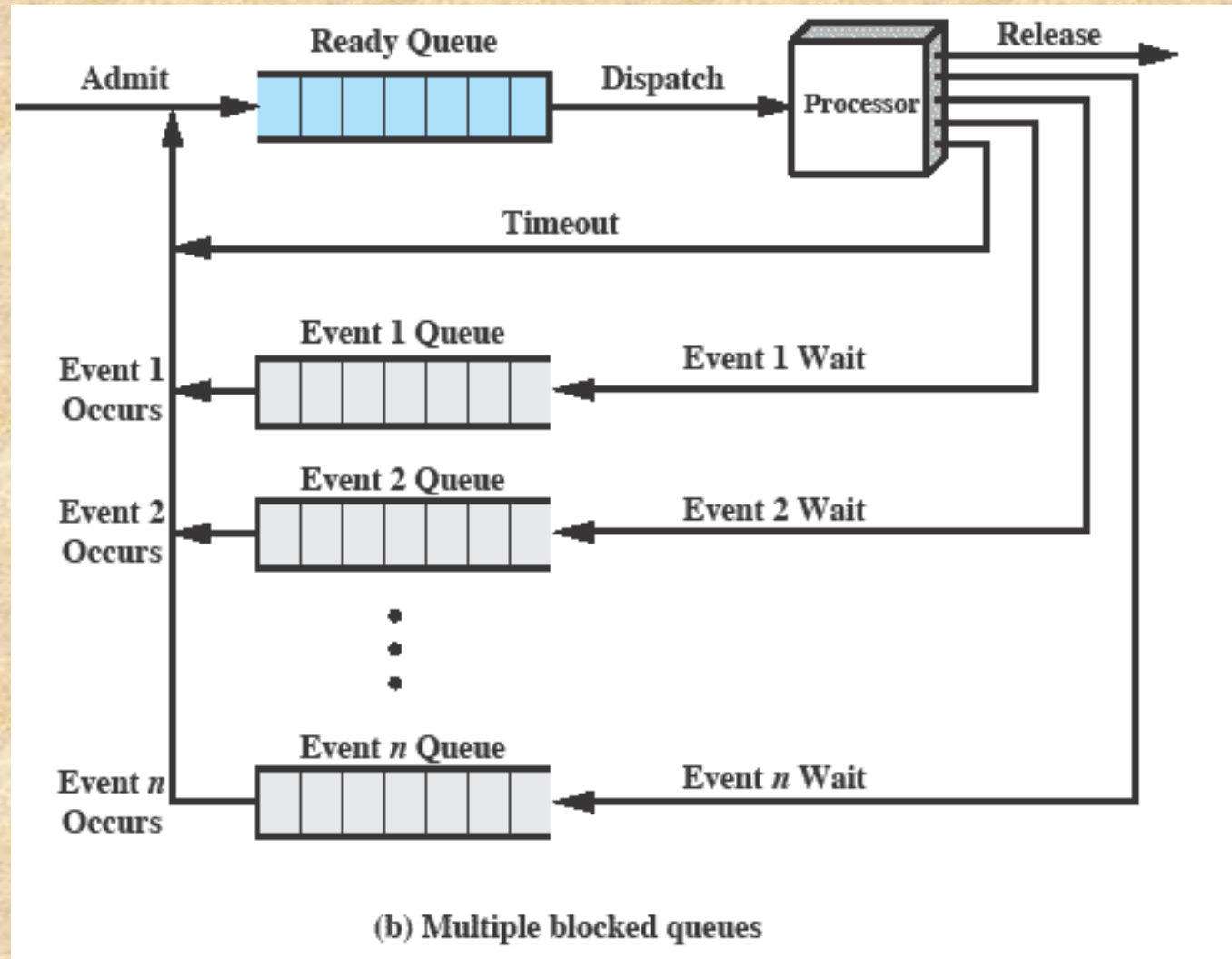
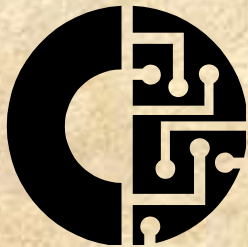
Figure 3.7 Process States for Trace of Figure 3.4

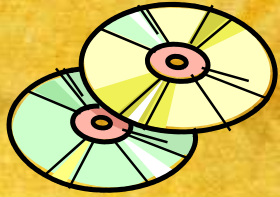
Using Two Queues



(a) Single blocked queue

Multiple Blocked Queues

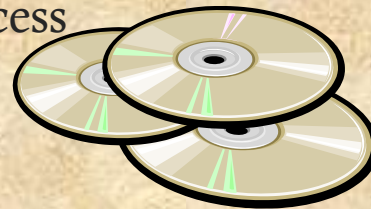




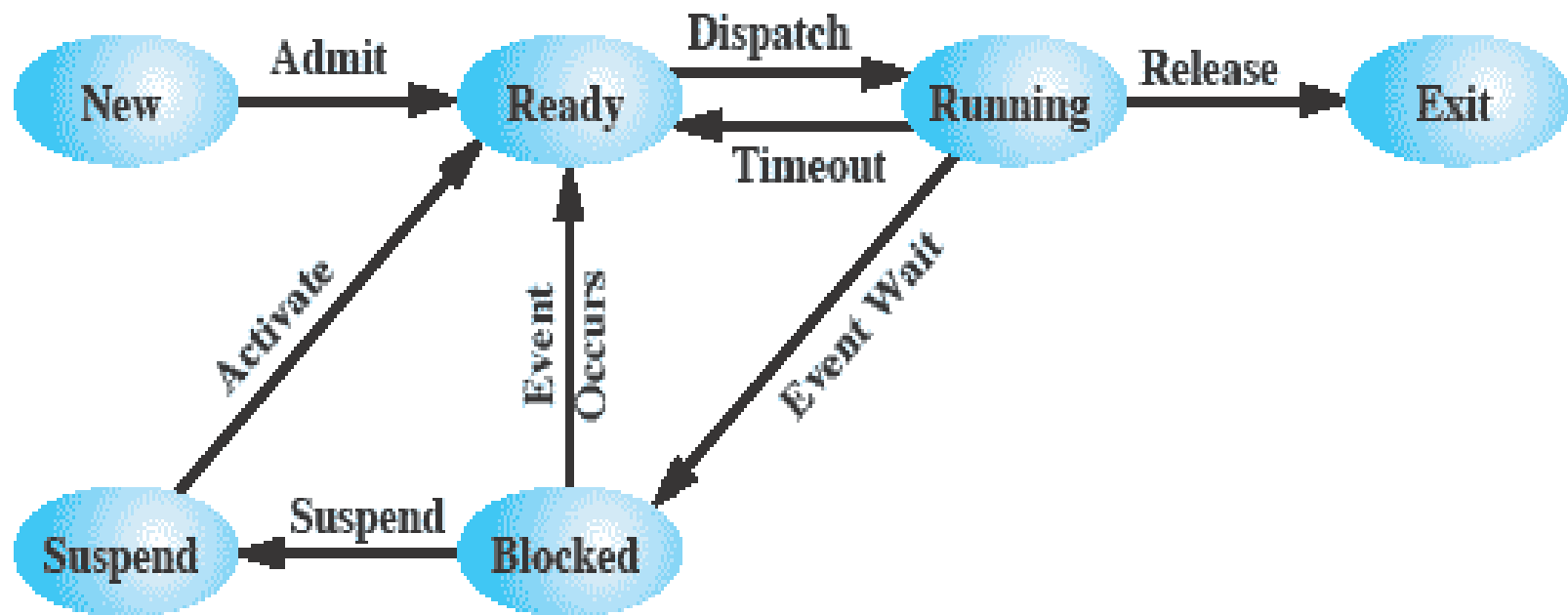
Suspended Processes

■ Swapping

- involves moving part or all of a process from main memory to disk
- when none of the processes in main memory is in the Ready state, the OS swaps one of the blocked processes out on to disk into a suspend queue to make room for a new process, or a previously suspended process that is now ready to execute
- Commonly used in systems that had no virtual memory; less likely to be used with virtual memory since process size can be controlled through the paging mechanism.

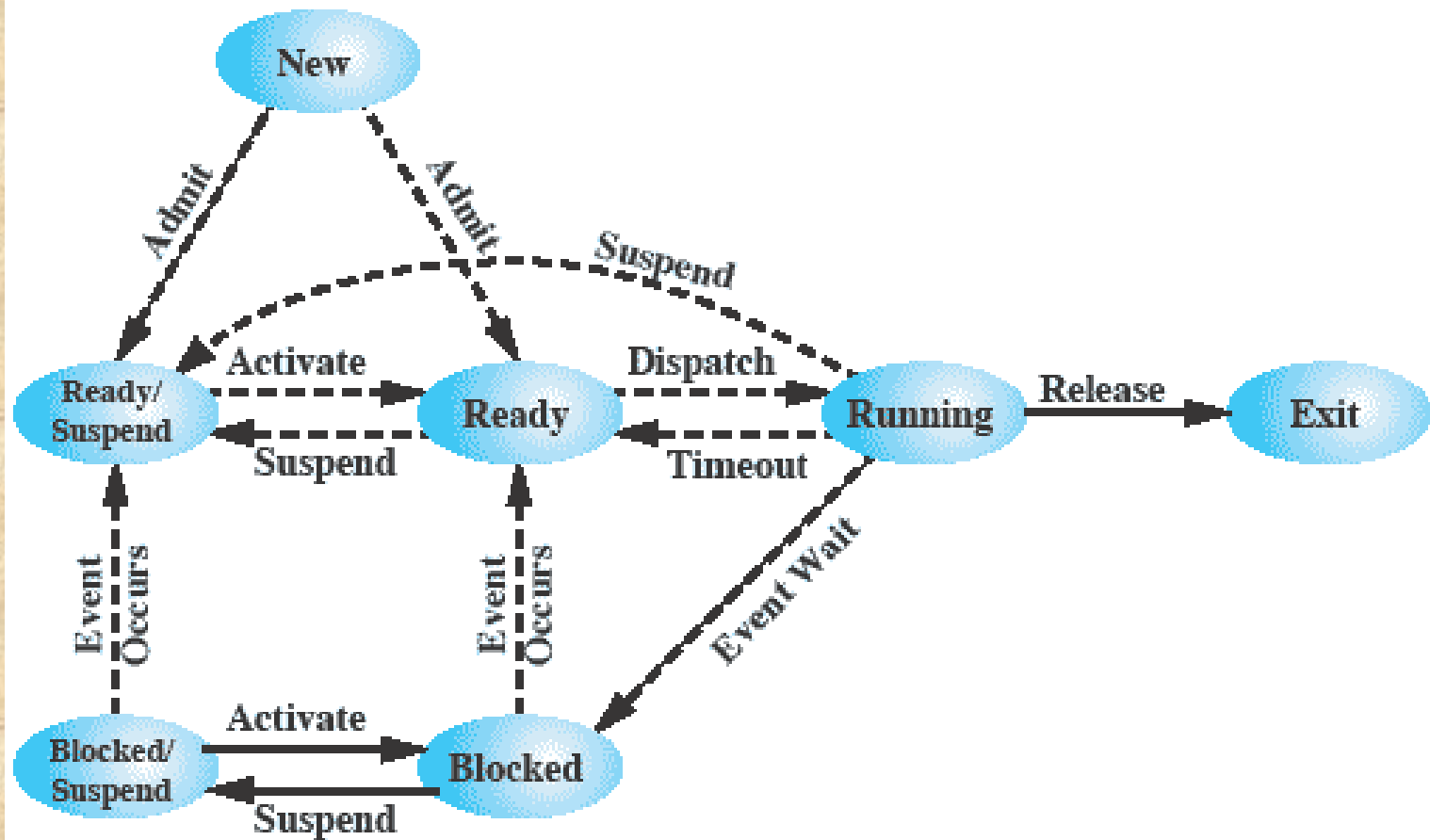


One Suspend State



(a) With One Suspend State

Two Suspend States



(b) With Two Suspend States

Characteristics of a Suspended Process

- The process is not immediately available for execution
- The process may or may not be waiting on an event
- The process was placed in a suspended state by an agent: either itself, a parent process, or the OS, for the purpose of preventing its execution
- The process may not be removed from this state until the agent explicitly orders the removal

Reasons for Process Suspension

Swapping	The OS needs to release sufficient main memory to bring in a process that is ready to execute.
Other OS reason	The OS may suspend a background or utility process or a process that is suspected of causing a problem.
Interactive user request	A user may wish to suspend execution of a program for purposes of debugging or in connection with the use of a resource.
Timing	A process may be executed periodically (e.g., an accounting or system monitoring process) and may be suspended while waiting for the next time interval.
Parent process request	A parent process may wish to suspend execution of a descendent to examine or modify the suspended process, or to coordinate the activity of various descendants.

Table 3.3 Reasons for Process Suspension

Processes and Resources

What does the OS need to know to manage processes & resources?

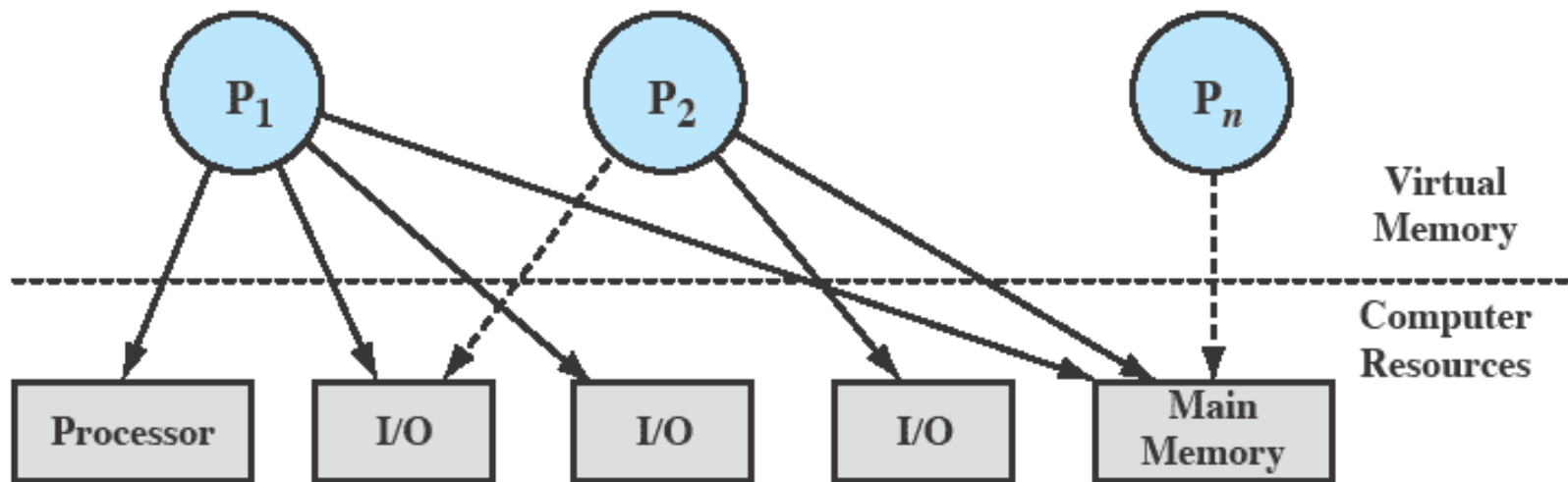


Figure 3.10 Processes and Resources (resource allocation at one snapshot in time)

OS Control Tables

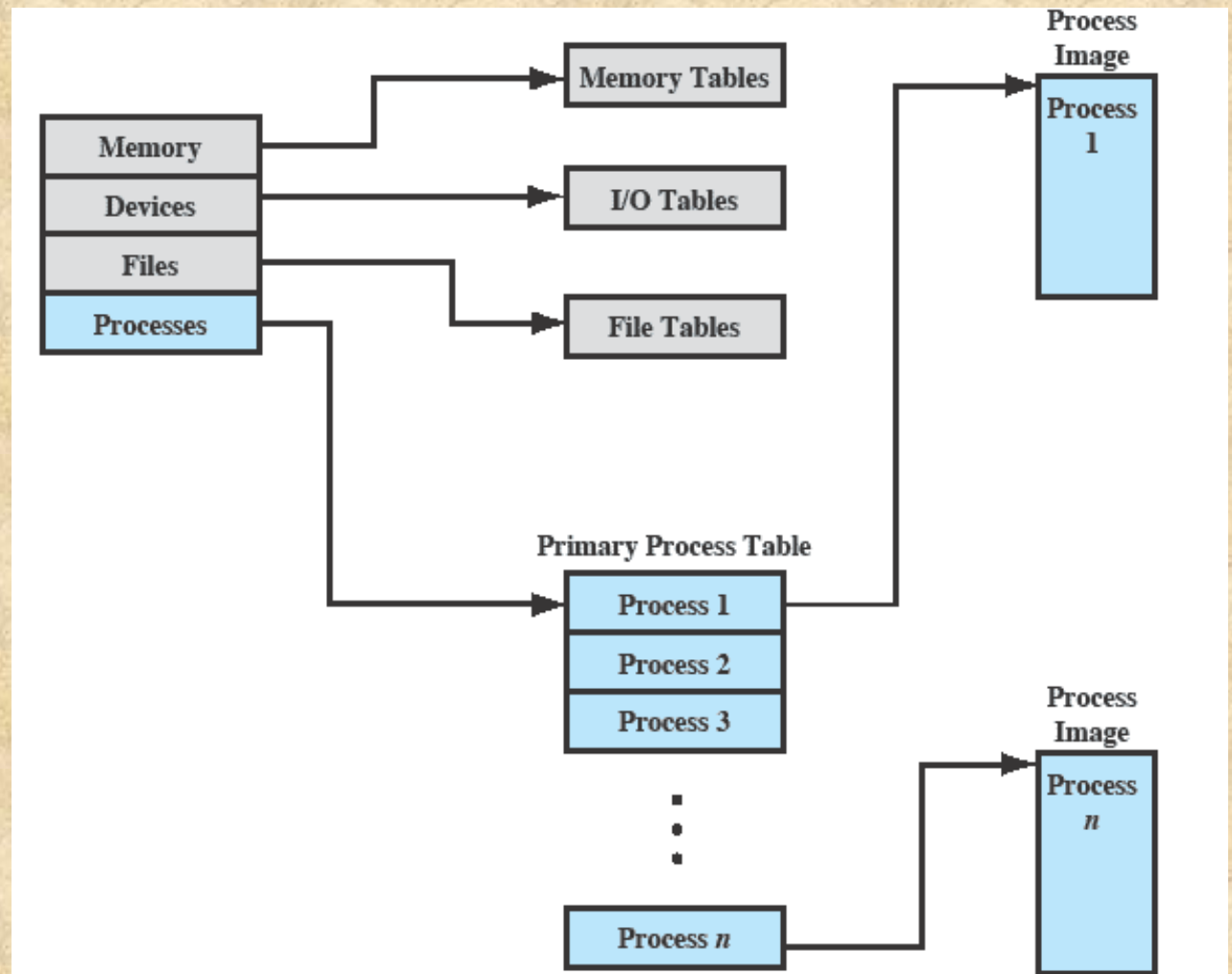


Figure 3.11 General Structure of Operating System Control Tables

Memory Tables

- Used to keep track of both main (real) and secondary (virtual) memory
- Processes are maintained on secondary memory using some sort of virtual memory or simple swapping mechanism

Must include:

allocation of main memory to processes

allocation of secondary memory to processes

protection attributes of blocks of main or virtual memory

information needed to manage virtual memory

I/O Tables

- Used by the OS to manage the I/O devices and channels of the computer system
- At any given time, an I/O device may be available or assigned to a particular process

If an I/O operation is in progress, the OS needs to know:

the status of the I/O operation

the location in main memory being used as the source or destination of the I/O transfer

File Tables

These tables provide information about:

- Information may be maintained and used by a file management system, or the OS itself
 - existence of files
 - location on secondary memory
 - current status
 - other attributes

Process Tables

- Must be maintained to manage processes
- Process tables must have some reference to memory, I/O, and file tables
 - e.g., pointers from each individual process entry to its data in the other three sets of tables.
- The tables themselves must be accessible by the OS and therefore are subject to memory management
- Processes are represented by a *process image*

Process Control Structures - PCB

Process Components

- A process must include a program or set of programs to be executed
- A process will consist of at least sufficient memory to hold the program(s) and data of that process
- The execution of a program typically involves a stack that is used to keep track of procedure calls and parameter passing between procedures

Process Attributes

- Each process has associated with it a number of attributes that are used by the OS for process control –maintained in a Process Control Block (PCB)
- The collection of program, data, stack, and attributes is referred to as the process image
- Process image location will depend on the memory management scheme being used

Process Attributes

Process Identification

Identifiers

Numeric identifiers that may be stored with the process control block include

- Identifier of this process
- Identifier of the process that created this process (parent process)
- User identifier

Processor State Information

User-Visible Registers

A user-visible register is one that may be referenced by means of the machine language that the processor executes while in user mode. Typically, there are from 8 to 32 of these registers, although some RISC implementations have over 100.

Control and Status Registers

These are a variety of processor registers that are employed to control the operation of the processor. These include

- **Program counter:** Contains the address of the next instruction to be fetched
- **Condition codes:** Result of the most recent arithmetic or logical operation (e.g., sign, zero, carry, equal, overflow)
- **Status information:** Includes interrupt enabled/disabled flags, execution mode

Stack Pointers

Each process has one or more last-in-first-out (LIFO) system stacks associated with it. A stack is used to store parameters and calling addresses for procedure and system calls. The stack pointer points to the top of the stack.

Process Identification

- Each process is assigned a unique numeric identifier
 - otherwise there must be a mapping that allows the OS to locate the appropriate tables based on the process identifier
- Many of the tables controlled by the OS may use process identifiers to cross-reference process tables
- Memory tables may be organized to provide a map of main memory with an indication of which process is assigned to each region
 - similar references will appear in I/O and file tables
- When processes communicate with one another, the process identifier informs the OS of the destination of a particular communication
- When processes are allowed to create other processes, identifiers indicate the parent and descendents of each process

Processor State Information

Consists of the contents of processor registers

- user-visible registers
- control and status registers
- stack pointers

Program status word (PSW)

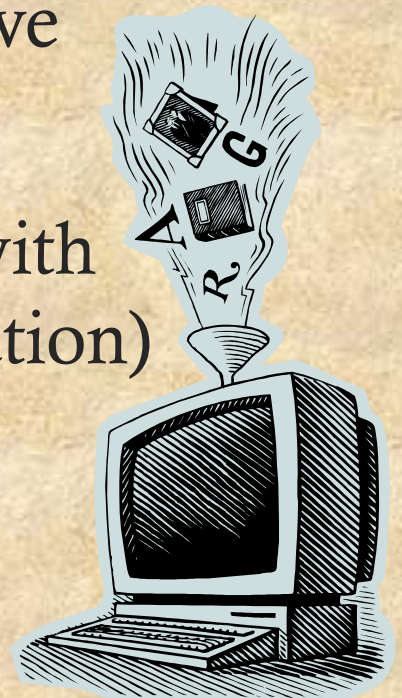
- contains condition codes plus other status information
- EFLAGS register is an example of a PSW used by any OS running on an x86 processor

31	/21															16 /15										0
	I	V	V	A	V	R		N	IO	O	D	I	T	S	Z		A		P		C					
	D	P	F	C	M	F		T	PL	F	F	F	F	F	F		F		F		F					

DF = Direction flag
IF = Interrupt enable flag
TF = Trap flag
SF = Sign flag
ZF = Zero flag
AF = Auxiliary carry flag
PF = Parity flag
CF = Carry flag

Process Control Information

- The additional information needed by the OS to control and coordinate the various active processes
- Forms the 3rd part of the PCB (along with process ID and processor state information)



Typical Elements of a Process Control Block



Process Control Information

Scheduling and State Information

This is information that is needed by the operating system to perform its scheduling function. Typical items of information:

- **Process state:** Defines the readiness of the process to be scheduled for execution (e.g., running, ready, waiting, halted).
- **Priority:** One or more fields may be used to describe the scheduling priority of the process. In some systems, several values are required (e.g., default, current, highest-allowable)
- **Scheduling-related information:** This will depend on the scheduling algorithm used. Examples are the amount of time that the process has been waiting and the amount of time that the process executed the last time it was running.
- **Event:** Identity of event the process is awaiting before it can be resumed.

Data Structuring

A process may be linked to other process in a queue, ring, or some other structure. For example, all processes in a waiting state for a particular priority level may be linked in a queue. A process may exhibit a parent-child (creator-created) relationship with another process. The process control block may contain pointers to other processes to support these structures.

Interprocess Communication

Various flags, signals, and messages may be associated with communication between two independent processes. Some or all of this information may be maintained in the process control block.

Process Privileges

Processes are granted privileges in terms of the memory that may be accessed and the types of instructions that may be executed. In addition, privileges may apply to the use of system utilities and services.

Memory Management

This section may include pointers to segment and/or page tables that describe the virtual memory assigned to this process.

Resource Ownership and Utilization

Resources controlled by the process may be indicated, such as opened files. A history of utilization of the processor or other resources may also be included; this information may be needed by the scheduler.

Structure of Process

Images in Virtual Memory

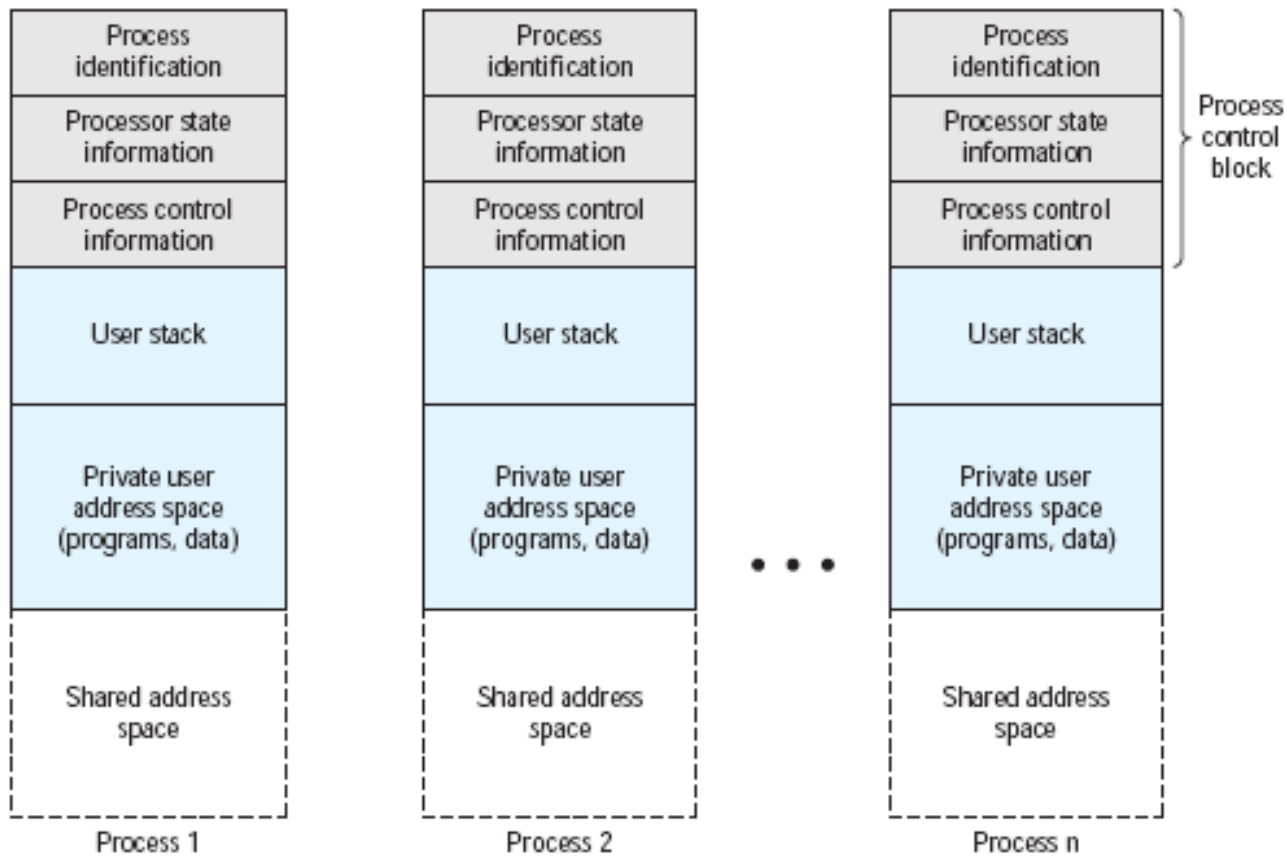
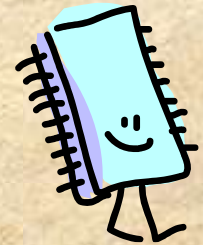


Figure 3.13 User Processes in Virtual Memory

Process List Structures

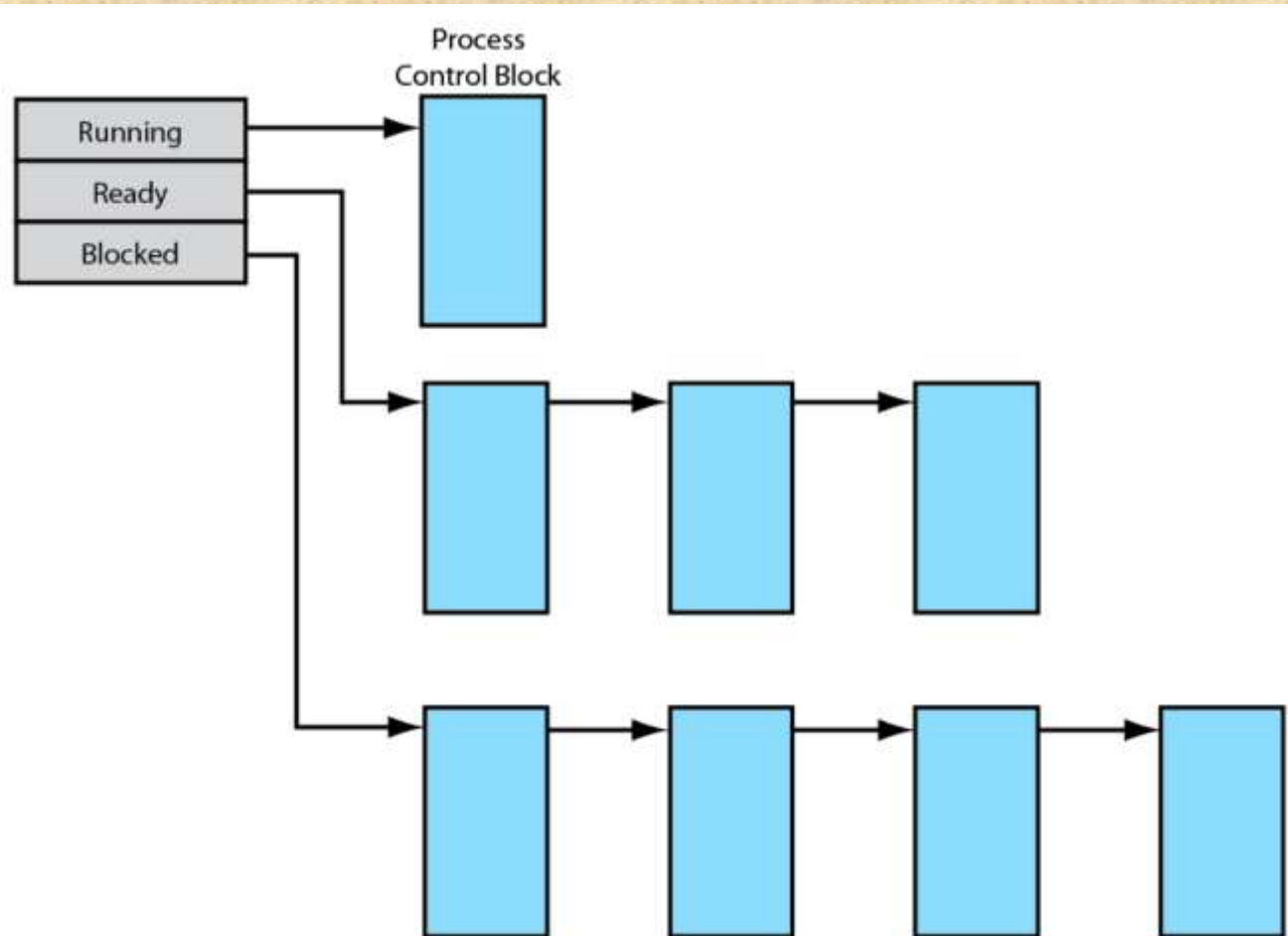


Figure 3.14 Process List Structures

Role of the Process Control Block

- The most important data structure in an OS
 - contains all of the information about a process that is needed by the OS
 - blocks are read and/or modified by virtually every module in the OS
 - defines the state of the OS
- Difficulty is not access, but protection
 - a bug in a single routine could damage process control blocks, which could destroy the system's ability to manage the affected processes
 - a design change in the structure or semantics of the process control block could affect a number of modules in the OS

Modes of Execution

User Mode

- less-privileged mode
- user programs typically execute in this mode



System Mode

- more-privileged mode
- also referred to as control mode or *kernel mode*
- kernel of the operating system

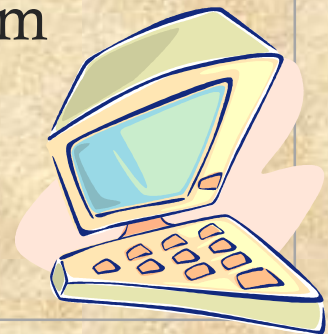


Table 3.7

Typical

Functions

of an

Operating

System

Kernel

Process Management

- Process creation and termination
- Process scheduling and dispatching
- Process switching
- Process synchronization and support for interprocess communication
- Management of process control blocks

Memory Management

- Allocation of address space to processes
- Swapping
- Page and segment management

I/O Management

- Buffer management
- Allocation of I/O channels and devices to processes

Support Functions

- Interrupt handling
- Accounting
- Monitoring

Process Creation

- Once the OS decides to create a new process it:

assigns a unique process identifier
to the new process



allocates space for the process



initializes the process control
block



sets the appropriate linkages



creates or expands other data
structures

Process (Context) Switching

A process switch may occur any time that the OS has gained control from the currently running process. Possible events giving OS control are:

Mechanism	Cause	Use
Interrupt	External to the execution of the current instruction	Reaction to an asynchronous external event
Trap	Associated with the execution of the current instruction	Handling of an error or an exception condition
Supervisor call	Explicit request	Call to an operating system function

System Interrupts

Interrupt

- Due to some sort of event that is external to and independent of the currently running process
 - clock interrupt
 - I/O interrupt
 - memory fault
- Time slice
 - the maximum amount of time that a process can execute before being interrupted

Trap

- An error or exception condition generated within the currently running process
- OS determines if the condition is fatal
 - moved to the Exit state and a process switch occurs
 - action will depend on the nature of the error

Mode Switching

If no interrupts are pending the processor:



proceeds to the fetch stage and fetches the next instruction of the current program in the current process

If an interrupt is pending the processor:



sets the program counter to the starting address of an interrupt handler program



switches from user mode to kernel mode so that the interrupt processing code may include privileged instructions

Mode Switching

- A mode switch is distinct from a process switch.
- A mode switch can take place without changing the state of the executing process
- However, while processing a mode switch the OS may decide a process switch is needed; e.g., if the mode switch was caused by a timer interrupt or a process request for I/O

Change of Process State

- The steps in a full process switch are:

save the context of the processor



update the process control block of the process currently in the Running state



move the process control block of this process to the appropriate queue



select another process for execution



update the process control block of the process selected



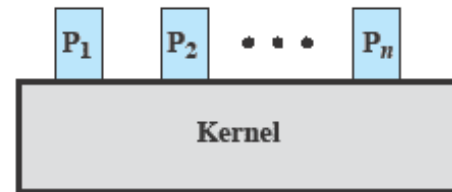
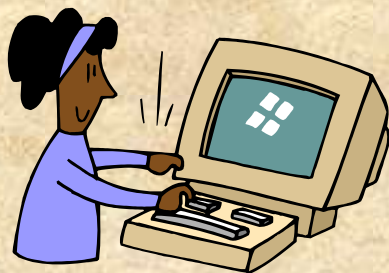
update memory management data structures



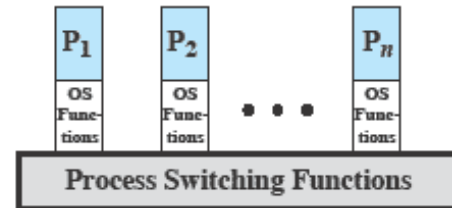
restore the context of the processor to that which existed at the time the selected process was last switched out

If the currently running process is to be moved to another state (Ready, Blocked, etc.), then the OS must make substantial changes in its environment

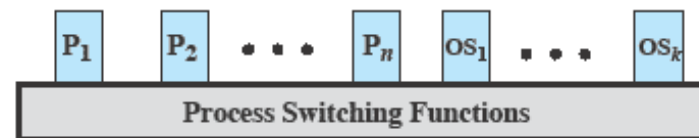
Execution of the Operating System



(a) Separate kernel



(b) OS functions execute within user processes



(c) OS functions execute as separate processes

Figure 3.15 Relationship Between Operating System and User Processes

Execution *Within* User Processes

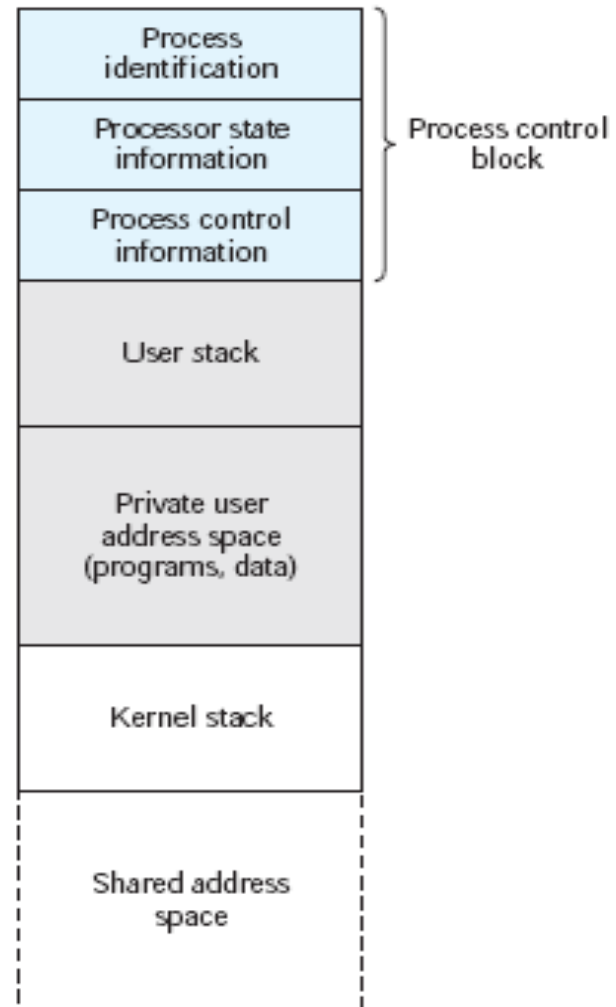
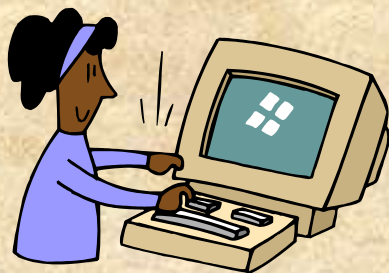


Figure 3.16 Process Image: Operating System Executes within User Space

Security Issues

- An OS associates a set of privileges with each process
- Typically a process that executes on behalf of a user has the privileges that the OS recognizes for that user
- Highest level of privilege is referred to as administrator, supervisor, or root access
- A key security issue in the design of any OS is to prevent, or at least detect, attempts by a user or malware from gaining unauthorized privileges on the system and from gaining root access



System Access Threats

Intruders

- Often referred to as a hacker or cracker

Classes:

- Masquerader
- Mifeasor
- Clandestine user



- Objective is to gain access to a system or to increase the range of privileges accessible on a system
- Attempts to acquire information that should have been protected

Malicious Software

- Most sophisticated types of threats to computer systems

Categories:

- those that need a host program (parasitic)
 - viruses, logic bombs, backdoors
- those that are independent
 - worms, bots

- Can be relatively harmless or very damaging



Countermeasures: Intrusion Detection

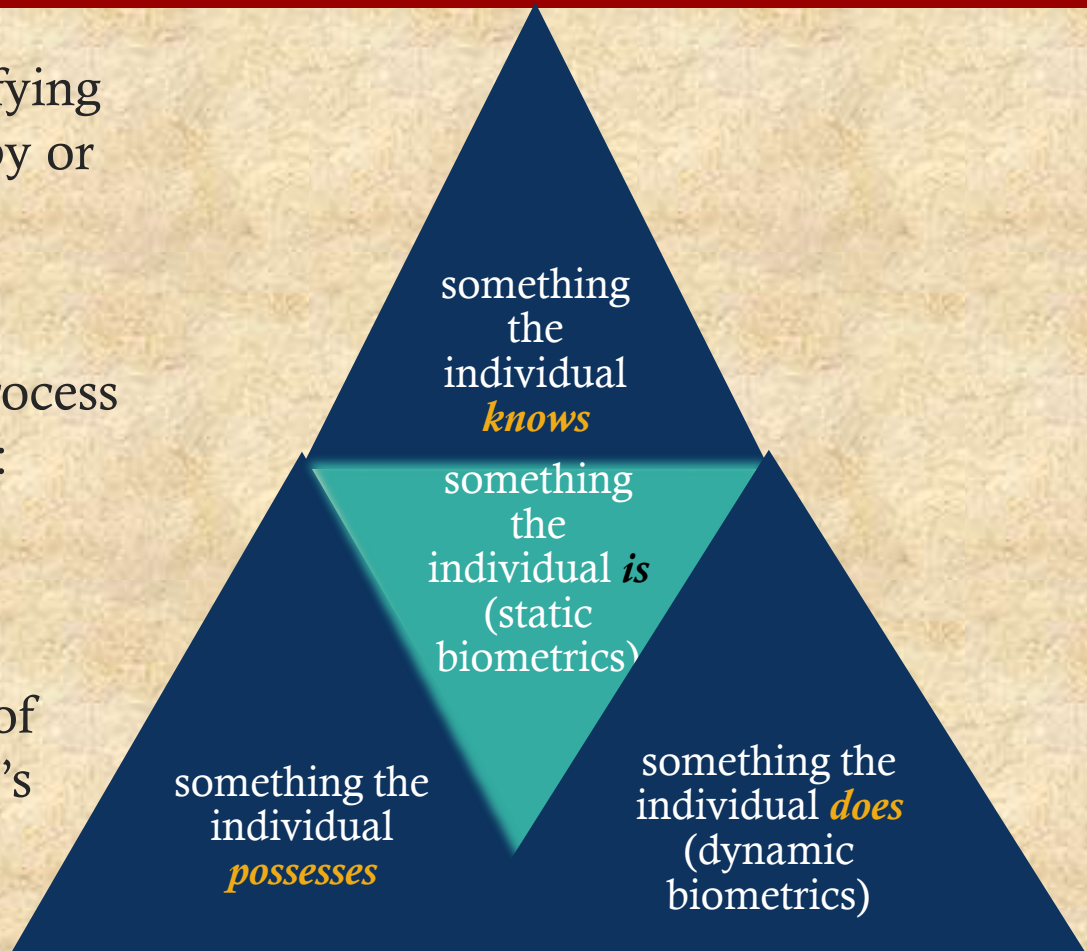
- “A security service that monitors and analyzes system events for the purpose of finding, and providing real-time or near real-time warning of, attempts to access system resources in an unauthorized manner” (RFC 2828)
- May be host or network based
- An intrusion detection system (IDS) comprises three logical components:



- IDSs are typically designed to detect human intruder behavior as well as malicious software behavior

Countermeasures: Authentication

- “The process of verifying an identity claimed by or for a system entity.” (RFC2828)
- An authentication process consists of two steps:
 - Identification
 - Verification
- Four general means of authenticating a user’s identity:



Countermeasures: Access Control

- Implements a security policy that specifies who or what may have access to each specific system resource and the type of access that is permitted in each instance
- Mediates between a user and system resources
- A security administrator maintains an authorization database
- An auditing function monitors and keeps a record of user accesses to system resources



Countermeasures: Firewalls

A dedicated
computer
that:

- interfaces with computers outside a network
- has special security precautions built into it to protect sensitive files on computers within the network

Design
goals of a
firewall:

- all traffic must pass through the firewall
- only authorized traffic will be allowed to pass
- immune to penetration



Unix SVR4



- Uses the model where most of the OS executes within the environment of a user process
- Two process categories: system processes and user processes
- System processes run in kernel mode
 - executes operating system code to perform administrative and housekeeping functions independent of any specific user process.
- User Processes
 - operate in user mode to execute user programs and utilities
 - operate in kernel mode to execute instructions that belong to the kernel
 - enter kernel mode by issuing a system call, when an exception is generated, or when an interrupt occurs

UNIX Process States

User Running	Executing in user mode.
Kernel Running	Executing in kernel mode.
Ready to Run, in Memory	Ready to run as soon as the kernel schedules it.
Asleep in Memory	Unable to execute until an event occurs; process is in main memory (a blocked state).
Ready to Run, Swapped	Process is ready to run, but the swapper must swap the process into main memory before the kernel can schedule it to execute.
Sleeping, Swapped	The process is awaiting an event and has been swapped to secondary storage (a blocked state).
Preempted	Process is returning from kernel to user mode, but the kernel preempts it and does a process switch to schedule another process.
Created	Process is newly created and not yet ready to run.
Zombie	Process no longer exists, but it leaves a record for its parent process to collect.

UNIX Process State Transition Diagram

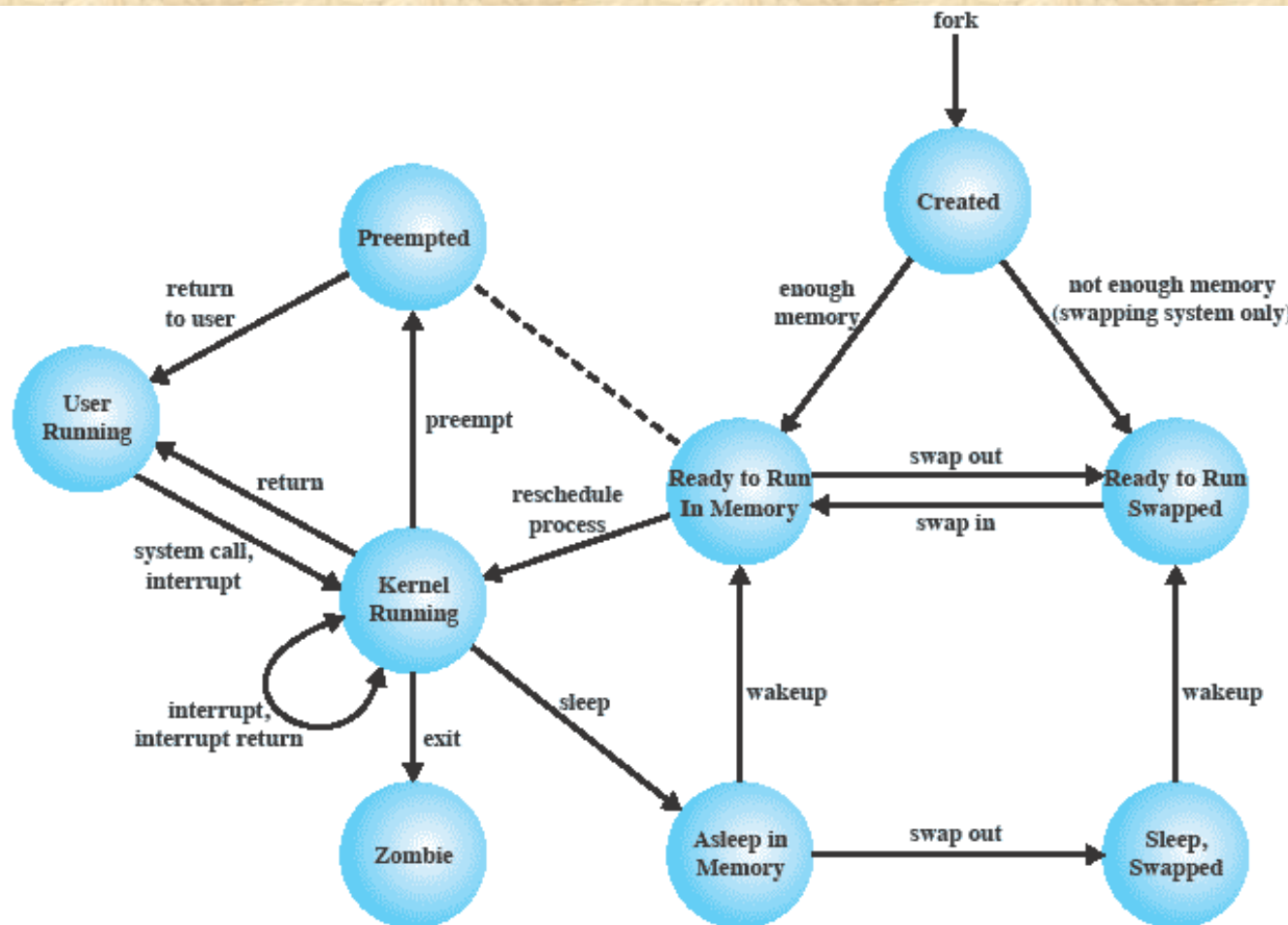


Figure 3.17 UNIX Process State Transition Diagram

A Unix Process



User-Level Context

Process text	Executable machine instructions of the program
Process data	Data accessible by the program of this process
User stack	Contains the arguments, local variables, and pointers for functions executing in user mode
Shared memory	Memory shared with other processes, used for interprocess communication

Register Context

Program counter	Address of next instruction to be executed; may be in kernel or user memory space of this process
Processor status register	Contains the hardware status at the time of preemption; contents and format are hardware dependent
Stack pointer	Points to the top of the kernel or user stack, depending on the mode of operation at the time or preemption
General-purpose registers	Hardware dependent

System-Level Context

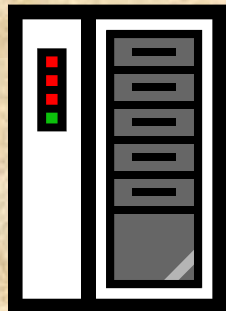
Process table entry	Defines state of a process; this information is always accessible to the operating system
U (user) area	Process control information that needs to be accessed only in the context of the process
Per process region table	Defines the mapping from virtual to physical addresses; also contains a permission field that indicates the type of access allowed the process: read-only, read-write, or read-execute
Kernel stack	Contains the stack frame of kernel procedures as the process executes in kernel mode

Table 3.11

UNIX

Process

Table Entry



Process status	Current state of process.
Pointers	To U area and process memory area (text, data, stack).
Process size	Enables the operating system to know how much space to allocate the process.
User identifiers	The real user ID identifies the user who is responsible for the running process. The effective user ID may be used by a process to gain temporary privileges associated with a particular program; while that program is being executed as part of the process, the process operates with the effective user ID.
Process identifiers	ID of this process; ID of parent process. These are set up when the process enters the Created state during the fork system call.
Event descriptor	Valid when a process is in a sleeping state; when the event occurs, the process is transferred to a ready-to-run state.
Priority	Used for process scheduling.
Signal	Enumerates signals sent to a process but not yet handled.
Timers	Include process execution time, kernel resource utilization, and user-set timer used to send alarm signal to a process.
P_link	Pointer to the next link in the ready queue (valid if process is ready to execute).
Memory status	Indicates whether process image is in main memory or swapped out. If it is in memory, this field also indicates whether it may be swapped out or is temporarily locked into main memory.

Table 3.12

UNIX U

Area

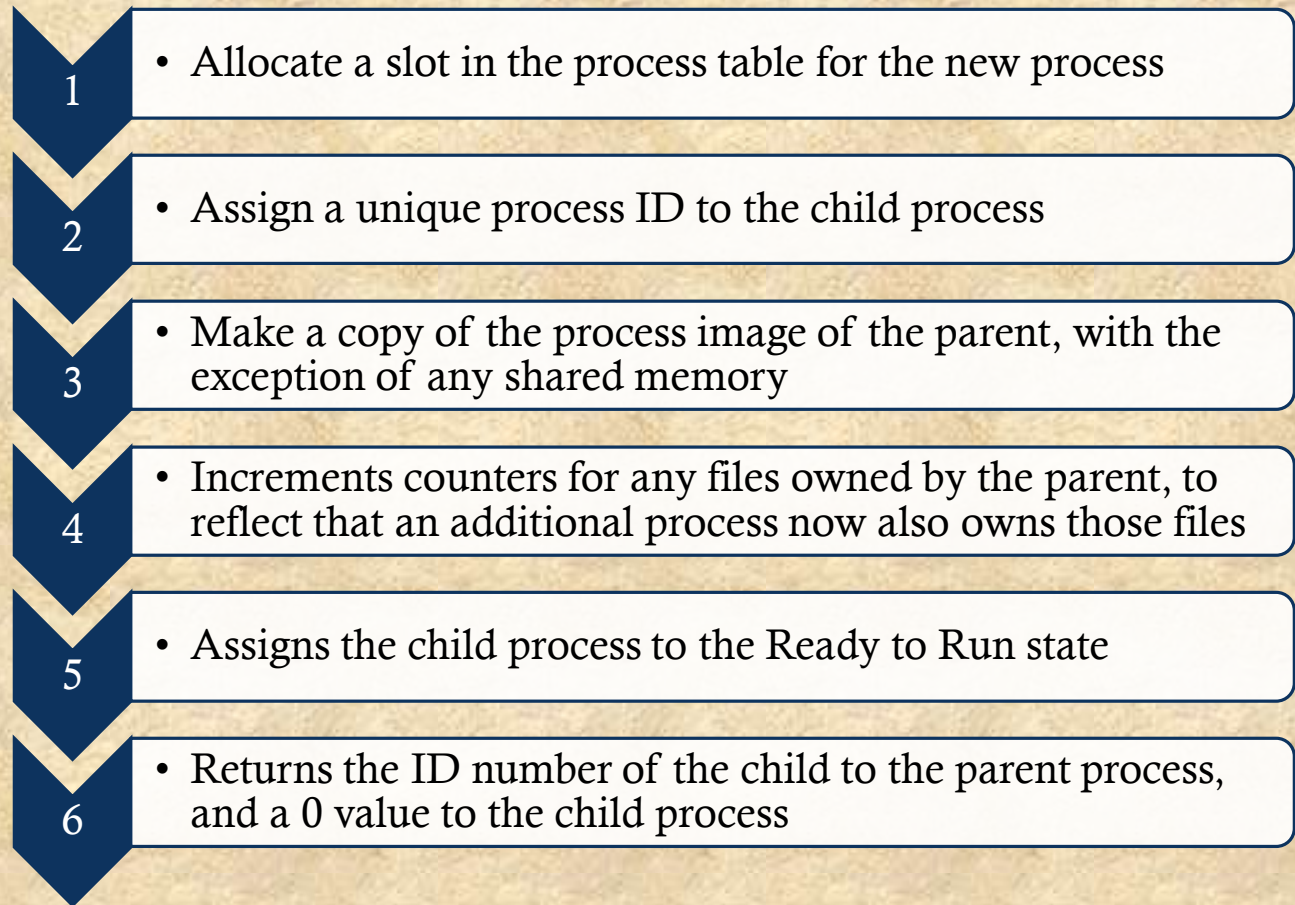


Process table pointer	Indicates entry that corresponds to the U area.
User identifiers	Real and effective user IDs. Used to determine user privileges.
Timers	Record time that the process (and its descendants) spent executing in user mode and in kernel mode.
Signal-handler array	For each type of signal defined in the system, indicates how the process will react to receipt of that signal (exit, ignore, execute specified user function).
Control terminal	Indicates login terminal for this process, if one exists.
Error field	Records errors encountered during a system call.
Return value	Contains the result of system calls.
I/O parameters	Describe the amount of data to transfer, the address of the source (or target) data array in user space, and file offsets for I/O.
File parameters	Current directory and current root describe the file system environment of the process.
User file descriptor table	Records the files the process has opened.
Limit fields	Restrict the size of the process and the size of a file it can write.
Permission modes fields	Mask mode settings on files the process creates.

Process Creation

- Process creation is by means of the kernel system call, `fork()`

- This causes the OS, in Kernel Mode, to:



After Creation

- After creating the process the Kernel can do one of the following, as part of the dispatcher routine:
 - stay in the parent process
 - transfer control to the child process
 - transfer control to another process



Summary

- The most fundamental concept in a modern OS is the process
- The principal function of the OS is to create, manage, and terminate processes
- Process control block contains all of the information that is required for the OS to manage the process, including its current state, resources allocated to it, priority, and other relevant data
- The most important states are Ready, Running and Blocked
- The running process is the one that is currently being executed by the processor
- A blocked process is waiting for the completion of some event
- A running process is interrupted either by an interrupt or by executing a supervisor call to the OS