Computer Science
Spring 2021 Lecture 14:
Dynamic Programming:
Longest Common Subsequence

## Let LCS(n,m) be the length of the longest common subsequence of X[1...n] and Y[1...m]. If O = n or O = m: return OIf X[n] = Y[m]: return I + L(S(n-1, m-1))else return I + L(S(n-1, m-1)) I(S(n-1, m-1))

```
LCS: Iterative Solution

LCS(n,m): // recursive for reference

if n = 0 or m = 0 then

return 0

else if X[n] = Y[m] then

return 1 + LCS(n-1,m-1)

else

return max( LCS(n-1,m), LCS(n,m-1))

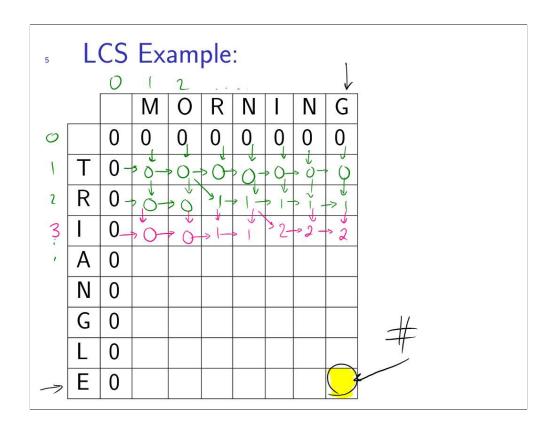
Declare L(S[0...n], 0...m]

for i = 0...n L(S[i,0] = 0] base

for j = 0...m L(S[i,0] = 0] (Mn)

for i = 1...m //fill in L(S[i,j])

O(i) \begin{cases} if & X[i] = Y[j] : L(S[i,j] = 1 + L(S[i,j-1]) \\ else & L(S[i,j-1]) \end{cases}
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Finding the LCS itself

We have LCS[] filled in for all n, m

i = n, j = m, S = empty stack

while i>0 and j>0 do

{ // is x[i] and y[j] part of output?

if x[i] == y[j]

push x[i] to stack. } e guiv:

else if L(S[i,j] == L(S[i-1,j])

i --

else j --

}

While S not empty

Output S.top()

S.pop()
```