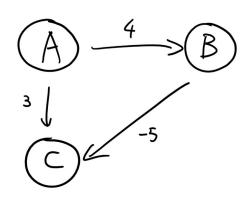
No, consider the following graph



Try using Dijkstra's Algorithm

Vertex	k	d	p
A	False	0	None
В	False	Infinity	None
С	False	Infinity	None

We will start from (A) and find the distance to (B) and (C)

Vertex	k	d	p
A	True	0	None
В	False	4	A
С	False	3	A

Then we mark k value for (C) to true, there is no outgoing edge in (C)

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	Vertex	k	d	p
	A	True	0	None
	В	False	4	A
	С	True	3	A

Update k value for (B) and find its outgoing edges, however (C) is already marked true, we cannot update distance to (C)

apatite distance to (G)				
	Vertex	k	d	p
	A	True	0	None
	В	True	4	A
	С	True	3	A

Here, the shortest distance from (A) to (C) is wrong, the actual shortest path is A->B->C with distance of 4 + (-5) = -1