```
makra.cpp
#ifndef LOCAL
#pragma GCC optimize ("03")
#endif
#include <bits/stdc++.h>
using namespace std:
//st, nd, mp, pb, eb
#define siz(c) ((int)(c).size())
#define all(c) (c).begin(), (c).end()
#define sim template < class c</pre>
#define ris return * this
#define mor > muu & operator << (
define R22(r) sim>typename enable if<1 r sizeof dud<c>(0),muu&>::type operator<<(c g)#
sim > struct rae { c b. e: }:
sim > rge<c> range(c b, c e) { return {b, e}; }
sim > auto dud(c* r) -> decltype(cerr << *r);</pre>
sim > char dud(...):
struct muu {
#ifdef LOCAL
stringstream a;
~muu() {    cerr << a.str() << endl;    }
R22(<) a << boolalpha << q; ris; }
R22(==) ris << range(begin(g), end(g)); }
sim, class b mor pair < b, c > r) { ris << "(" << r.st << ", " << r.nd << ")"; }
sim mor rae<c> u) {
 a << "[";
 for (c i = u.b; i != u.e; ++i) *this << ", " + 2 * (i == u.b) << *i;
 ris << "]";
template <class...c mor tuple<c...> x) {
 int q = 0;
 a << "(":
 apply([&](c...y){
   ((*this << ", " + 2 * !q++ << v), ...);
 }, x);
 ris << ")";
#define qel(t) sim, class d, class...e mor t<c,d,e...> x){ris << *(d*)&x;}
qel(stack) qel(queue) qel(priority queue)
sim mor const c&) { ris: }
#endif
#define imie(r...) "[" #r ": " << (r) << "] "
#define range(b, e) "[[" #b ", " #e "): " << range(b, e) << "] "
#define debug muu() << FUNCTION << "#" << LINE << ": '
#define endl '\n'
//using ll,ld,pii,vi,vpii,ull,pdd
sim > void mini(c \&a, const c \&b) {if (a > b) a = b;}
sim > void maxi(c \&a, const c \&b) {if (a < b) a = b;}
int32 t main() {ios base::sync with stdio(0); cin.tie(0);}
                                       geany-config
Interface:
 Message window: right
Keybindings:
 Build: Compile: F8. Build: F9
 Focus: Switch to Editor: F1, Switch to VTE: F2
Set Build Commands:
 compile: q++ -Wall -Wshadow -o "%e" "%f" -g -O3 -std=c++1z
 build: a++ -Wall -Wshadow -o "%e" "%f" -fsanitize=undefined -fsanitize=address
   -DLOCAL -D GLIBCXX DEBUG -g -std=c++1z
                                Rzeczy Na Dzien Probny.txt
1. int128/ uint128 t
2. Czy rand() działa tak samo na sprawdzaczce.
3. while (clock() < 1.9 * CLOCKS PER SEC)</pre>
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    Czy time(NULL) daje różne wartości między uruchomieniami

5. Czy kompresja (wklejanie binarnych rzeczy do kodu) działa
6. float128
7. \overline{\text{ordered}} set
8. getchar unlocked/ getchar nolock
9. Porównać czas działania sprawdzaczki i kompa
10. qp hash table / cc hash table
  #include <ext/pb ds/assoc container.hpp>
  using namespace gnu pbds;
                               Data Structures/OrderedSet.cpp
#include<ext/pb ds/assoc container.hpp> // ordered set
#include<ext/pb ds/tree policy.hpp>
using namespace qnu pbds; template <typename T> using ordered set =
    tree<T, null type, less<T>, rb tree tag, tree order statistics node update>:
ordered set<int> s; s.insert(1); s.insert(2);
 s.order of key(1); // Out: 0.
*s.find by order(1); // Out: 2.
                                 Data Structures/Treap.cpp
struct node {
 node *L, *R;
  int ind, prior, sub, lazy; // sub opcjonalne
 node(int ind) : L(0), R(0), ind(ind), prior(rand()), sub(1), lazy(0) {}
void rev(node* v) { // przykladowy update, odwraca kolejnosc w poddrzewie
 v \rightarrow lazy = 1;
  swap(v \rightarrow L, v \rightarrow R);
void push(node* v) { // opcjonalne
 if (v->lazy) {
    if (v->L) rev(v->L):
    if (v->R) rev(v->R);
    v -> lazv = 0:
node* attach(node* v, node* l, node* r) {
 v->L = l; // jesli chcemy trzymac ojca to update w tej funkcji
  v \rightarrow R = r;
 v->sub = 1 + (l ? l->sub : 0) + (r ? r->sub : 0); // opcjonalne
  return v:
node* merge(node* v, node* u) {
 if (!u) return v:
  if (!v) return u;
  push(v);
  push(u);
  if (v->prior > u->prior) return attach(v, v->L, merge(v->R, u));
  else return attach(u, merge(v, u->L), u->R);
pair<node*, node*> split size(node* v, int k) { // (prefiks rozmiaru k, reszta)
 if (!v) return mp(v, v):
  int lewo = v->L ? v->L->sub : 0;
  push(v):
  if (lewo >= k) {
    auto s = split size(v->L, k);
    return mp(s.st. attach(v. s.nd. v->R)):
    auto s = split size(v->R, k - lewo - 1);
    return mp(attach(v, v->L, s.st), s.nd);
pair<node*, node*> split lex(node *v, int k) { // (ind <= k, reszta)</pre>
 if (!v) return mp(v, v);
  if (k < v->ind) {
    auto s = split lex(v->L, k):
    return mp(s.st, attach(v, s.nd, v->R));
  } else {
    auto s = split lex(v->R, k);
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return mp(attach(v, v->L, s.st), s.nd);
int find pos(node *v, int val) { // -1 jesli nie ma
 if (!v) return -1:
 int lewo = v->L ? v->L->sub : 0;
 if (v->ind == val) return 1 + lewo:
 if (val < v->ind) return find pos(v->L, val);
 else return 1 + lewo + find pos(v->R, val);
                              Flows And Matchings/Dinic.cpp
using T = long long;
struct Flow {
 struct E {
   int dest:
  T orig, *lim, *rev;
 int zr, uj, n = 0;
 vector<unique ptr<T>> ts;
 vector<vector<E>> graf:
 vi ptr, odl;
 void vert(int v) {
   n = \max(n, v + 1);
   graf.resize(n):
   ptr.resize(n);
   odl.resize(n);
 bool iszero(T v) {
   return !v; // Zmienić dla doubli.
 void bfs() {
   fill(all(odl), 0);
   vi kol = \{zr\}:
   odl[zr] = 1;
   for (int i = 0; i < siz(kol); i++) {</pre>
     for (E& e : graf[kol[i]]) {
       if (!odl[e.dest] and !iszero(*e.lim)) {
         odl[e.dest] = odl[kol[i]] + 1;
         kol.pb(e.dest);
     }
 T dfs(int v, T lim) {
   if (v == uj) return lim;
   T ret = 0, wez:
   for (int& i = ptr[v]; i < siz(graf[v]); i++) {</pre>
     E\& e = qraf[v][i]:
     if (odl[e.dest] == odl[v] + 1 and !iszero(*e.lim) and
          !iszero(wez = dfs(e.dest, min(*e.lim, lim)))) {
       ret += wez:
       *e.lim -= wez:
       *e.rev += wez:
       lim -= wez:
       if (iszero(lim)) break:
     }
   return ret;
 void add edge(int u, int v, T lim, bool bi = false /* bidirectional? */) {
   vert(max(u, v));
   T *a = new T(lim), *b = new T(lim * bi);
   ts.eb(a):
   ts.eb(b);
   graf[u].pb(E{v, lim,
                              a, b});
   graf[v].pb(E{u, lim * bi, b, a});
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T dinic(int zr , int uj ) {
   zr = zr_; uj = uj_;
   vert(ma\overline{x}(zr, uj));
   T ret = 0;
   while (true) {
     bfs();
     fill(all(ptr), 0);
     const T sta = dfs(zr, numeric limits<T>::max()); // Dla doubli można dać
     if (iszero(sta)) break;
                                                        // infinity() zamiast
     ret += sta:
                                                         // max().
   return ret;
 vi cut() {
   vi ret;
   bfs():
   for (int i = 0; i < n; i++)
     if (odl[i])
        ret.pb(i);
   return ret;
 map<pii, T> get flowing() { // Tam gdzie plynie 0 może nie być
                              // krawedzi.
   map<pii, T> ret;
   for (int i = 0; i < n; i++)
     for (E& e : graf[i])
       if (*e.lim < e.oria)</pre>
          ret[mp(i, e.dest)] += e.orig - *e.lim;
   for (auto& i : ret) {
     const pii rev{i.st.nd, i.st.st};
     const T x = min(i.nd, ret[rev]);
     i.nd -= x:
     ret[rev] -= x;
   return ret;
 }
                            Flows And Matchings/Gomory Hu.cpp
#include "Dinic.cpp"
struct GomoryHu { // #define int ll//jeśli long longi potrzebne
 vector<vpii> graph, tree:
 vector<vi> nodes;
 vector<bool> visited;
                                   //wymaga naszego dinica
 vi groupId, contrId;
                                   //numeracia od zera
 int wnode. n:
 GomoryHu(int N) : graph(N), visited(N), groupId(N), contrId(N), tree(N), n(N) {}
 void addEdge(int u, int v, int cap) {
   graph[u].eb(v, cap);
   graph[v].eb(u, cap);
 void dfs(int v, int type) {
   visited[v] = true; contrId[v] = type;
   for (auto P : tree[v]) { if (!visited[P.st]) { dfs(P.st, type); } }
 vector <pair<pii,int>> run() {
   vi allNodes(n):
   iota(all(allNodes), 0);
   nodes = vector<vi>{allNodes};
   tree = vector<vpii>(n);
   fill(all(groupId), 0);
   for (int step = 1; step < n; step++) {</pre>
     Flow flow:
     for (int i = 0; i < siz(nodes); i++) {
       if (siz(nodes[i]) > 1) { wnode = i; break; }
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C D F G M M N S
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visited[wnode] = true;
     for (auto P : tree[wnode]) { dfs(P.st, nodes[P.st][0]); }
     for (int v = 0; v < n; v++) {
       int a = groupId[v] == wnode ? v : contrId[groupId[v]];
       for (auto& P : graph[v]) {
         int b = groupId[P.st] == wnode ? P.st : contrId[groupId[P.st]];
         if (a != b) { flow.add edge(a, b, P.nd); }
     int a = nodes[wnode][0], b = nodes[wnode][1], f = flow.dinic(a, b);
     auto pom = flow.cut();
     vector <bool> cut(n, false);
     for (int i : pom)
       cut[i]=1;
     for (int v = 0; v < step; v++) {
       if (v == wnode) { continue; }
       for (auto& P : tree[v]) {
         if (P.st == wnode \&\& !cut[contrId[v]]) \{ P.st = step; \}
     vpii PA, PB;
     for (auto& P : tree[wnode]) { (cut[contrId[P.st]] ? PA : PB).pb(P); }
     tree[wnode] = PA; tree[step] = PB;
     tree[wnode].eb(step, f);
     tree[step].eb(wnode, f);
     vi A. B:
     for (int v : nodes[wnode]) {
       (cut[v] ? A : B).pb(v);
       if (!cut[v]) { groupId[v] = step; }
     nodes[wnode] = A:
     nodes.pb(B);
   vector <pair<pii,int>> res;
   for (int i = 0; i < n; i++)
     for (auto P : tree[i])
       if (nodes[i][0]<nodes[P.st][0])
         res.pb({{nodes[i][0], nodes[P.st][0]}, P.nd});
   return res;
}; // #undef int
                    Flows And Matchings/Highest Label Push Relabel.cpp
class LinkedList {
nublic:
 LinkedList(int N) : N(N), next(N) { clear(); }
 void clear() { head.assign(N, -1); }
 int front(int h) { return head[h]; }
 void pop(int h) { head[h] = next[head[h]]; }
 void push(int h, int u) { next[u] = head[h], head[h] = u; }
private:
 int N:
 vi next, head;
class DoublyLinkedList {
private:
 struct Node { int prev, next; };
public:
 DoublyLinkedList(int N) : N(N), nodes(N) { clear(); }
 void clear() { head.assign(N, -1); }
 void erase(int h, int u) {
   int pv = nodes[u].prev, nx = nodes[u].next;
   if (nx >= 0) nodes[nx] prev = pv;
   if (pv >= 0) nodes[pv].next = nx;
   else head[h] = nx;
```

fill(all(visited), false);

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void insert(int h, int u) {
    nodes[u] = \{ -1, head[h] \};
    if (head[h] >= 0) nodes[head[h]].prev = u;
    head[h] = u;
  template <typename Func>
  void erase all(int first, int last, Func f) {
    for (int i = first; i <= last; ++i) {</pre>
      for (int h = head[i]; h >= 0; h = nodes[h].next) f(h);
      head[i] = -1;
 }
private:
 int N:
 vi head;
 vector<Node> nodes:
template <
    typename CapType, typename TotalCapType,
    bool UseGlobal = true, bool UseGap = true
class HighestLabelPushRelabel {
private:
 TotalCapType inf = pow(10., sizeof(TotalCapType) / 4 * 9);
  struct Edge { int to, rev; CapType cap; };
public:
  HighestLabelPushRelabel(int N) : N(N), E(0), G(N), que(N), list(N), dlist(N) {}
  TotalCapType maximum flow(int s, int t) {
    if (s == t) return 0:
    highest active = 0; // highest label (active)
    highest = 0; // highest label (active and inactive)
    height.assign(N, 0); height[s] = N;
    for (int i = 0; i < N; ++i) if (i != s) dlist.insert(height[i], i);
    count.assign(N, 0); count[0] = N - 1;
    excess.assign(N, 0); excess[s] = inf; excess[t] = -inf;
    for (auto\& e : G[s]) push(s, e);
    global relabel(t):
    for (int u = -1, rest = N; highest active >= 0; ) {
      if ((u = list.front(highest active)) < 0) { --highest active; continue; }</pre>
      list.pop(highest active):
      discharge(u);
      if (--rest == 0) rest = N, global relabel(t);
    return excess[t] + inf;
  inline void add directed edge(int u, int v, CapType cap) {
    G[u].pb({v, siz(G[v]), cap});
    G[v].pb(\{u, siz(G[u]) - 1, 0\});
  inline void add undirected edge(int u, int v, CapType cap) {
    G[u].pb(\{v, \overline{siz}(G[v]), \overline{cap}\});
    G[v].pb({u, siz(G[u]) - 1, cap});
private:
  void global relabel(int t) {
    if (!UseGlobal) return;
    height.assign(N, N); height[t] = 0;
    count.assign(N. 0):
    int qh = 0, qt = 0;
    for (que[qt++] = t; qh < qt; ) {
      int u = que[qh++], h = height[u] + 1;
      for (auto \& e : G[u]) if (height[e.to] == N && G[e.to][e.rev].cap > 0) {
          count[height[e.to] = h] += 1;
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que[\overline{qt++}] = e.to;
       }
   list.clear(): dlist.clear():
   for (int u = 0; u < N; ++u) if (height[u] < N) {
        dlist.insert(height[u], u);
        if (excess[u] > 0) list.push(height[u], u);
   highest = highest active = height[que[qt - 1]];
 void push(int u, Edge& e) {
   int v = e.to:
   CapType df = min(excess[u], TotalCapType(e.cap));
   e.cap -= df. G[v][e.rev].cap += df:
   excess[u] -= df, excess[v] += df;
   if (0 < excess[v] && excess[v] <= df) list.push(height[v], v);</pre>
 void discharge(int u) {
   int nh = N:
   for (auto\& e : G[u]) if (e.cap > 0) {
        if (height[u] == height[e.to] + 1) {
          push(u, e);
          if (excess[u] == 0) return;
        } else nh = min(nh, height[e.to] + 1);
   int h = height[u];
   if (UseGap && count[h] == 1) {
     auto f = [\&] (int u) { count[height[u]]--, height[u] = N; };
     dlist.erase all(h, highest, f);
     highest = h - 1;
   } else {
     count[h]--; dlist.erase(h, u);
     height[u] = nh;
     if (nh == N) return;
     count[nh]++; dlist.insert(nh, u);
     highest = max(highest, highest active = nh);
     list.push(nh, u);
 int N, E, highest active, highest;
 vector< vector<Edge> > G:
 vi height, count, que;
 vector<TotalCapType> excess;
 LinkedList list;
 DoublvLinkedList dlist:
using HIPR GP = HighestLabelPushRelabel<int. int. true. true>;
                            Flows And Matchings/Hungarian.cpp
//a[1..n][1..m] - wagi, n \le m, O(n^2 * m), znajduje minimalny koszt (masymalny: *-1)
//u[0..n], v[0..m] - funkcja potencjalu, p[0..m], ans[0..n] - skojarzenie
struct hungarian { // cost = -v[0]
 int n, m;
 vector<vi> a;
 vi u, v, p, way, ans;
 hungarian(int n, int m) : n( n), m( m) {
   #define re(x, y) (x).resize(\overline{y} + 1)
   re(a, n); re(u, n); re(v, m); re(p, m); re(way, m), re(ans, n);
   for (auto \&x: a) re(x, m);
 void addcost(int v1, int v2, int c) {
   a[v1][v2] = c;
 void solve() {
   for(int i = 1; i <= n; ++i) {</pre>
     p[0] = i;
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int j0 = 0;
      vi minv(m + 1, INT MAX); // uwaga, moze byc potrzebny LL!
     vector<char> used(m + 1, false);
       used[i0] = true;
       int i0 = p[j0], delta = INT MAX, j1 = 0;
       for (int j = 1; j \le m; ++j) {
          if (!used[j]) {
            int cur = a[i0][j] - u[i0] - v[j];
            if (cur < minv[j]) {</pre>
              minv[j] = cur;
              way[j] = j0;
            if (minv[i] < delta) {</pre>
              delta = minv[j];
              j1 = j;
       } } }
        for (int j = 0; j <= m; ++j) {
         if (used[j]) {
            u[p[i]] += delta;
            v[j] -= delta;
          else {
            minv[j] -= delta;
       i0 = i1;
     } while (p[j0] != 0);
       int i1 = way[i0];
       p[j0] = p[j1];
       i0 = i1;
     } while (j0);
   for (int j = 1; j \le m; ++j)
     ans[p[i]] = j; // odzyskiwanie wyniku
                             Flows And Matchings/Matching.cpp
struct Matching {
 int n, tim = 0, top = 0;
 vi mat, fa, s, q, pre, vis, head:
 vpii e;
 Matching(int N) : n(N + 1), mat(n, -1), fa(n), s(n), q(n), pre(n), vis(n),
  void edge impl(int x, int y) {e.eb(y, head[x]);head[x] = siz(e) - 1;}
  void add edge(int x, int y) {edge impl(x, y), edge impl(y, x);}
  int find(int x) {return x == fa[x] ? x : fa[x] = find(fa[x]):}
  int lca(int x, int y) {
   for (++tim, x = find(x), y = find(y); ; swap(x, y)) {
     if (~x) {
       if (vis[x] == tim) return x;
       vis[x] = tim:
       x = \text{-mat}[x] ? find(pre[mat[x]]) : -1;
 void blossom(int x, int y, int l) {
   while (find(x) != l) {
     pre[x] = y;
     if (s[mat[x]] == 1) s[q[top++] = mat[x]] = 0;
     if (fa[x] == x) fa[x] = l;
     if (fa[mat[x]] == mat[x]) fa[mat[x]] = l;
     y = mat[x];
     x = pre[y];
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bool match(int x) {
   iota(all(fa), 0);
   fill(all(s), -1);
   s[q[0] = x] = 0;
   top = 1:
   for (int i = 0; i < top; ++i) {
     for (int t = head[q[i]]; ~t; t = e[t].nd) {
        int y = e[t].st;
       if (s[y] == -1) {
          pre[y] = q[i];
          s[y] = 1;
          if (mat[y] == -1) {
           for (int u = y, v = q[i], lst; \sim v; u = lst, v = \sim u? pre[u] : -1)
             lst = mat[v], mat[v] = u, mat[u] = v;
            return true;
         s[q[top++] = mat[y]] = 0;
        else if (!s[y] \&\& find(y) != find(g[i])) {
         int l = lca(y, q[i]);
          blossom(v, q[i], l);
          blossom(q[i], y, l);
     }
   }
   return false;
 int run() {
   int size = 0:
   for (int i = 0; i < n; ++i)
     if (mat[i] == -1 \&\& match(i))
        size++:
   return size;
                             Flows And Matchings/Min Cost.cpp
struct MinCost {
 struct kra {
   int cel, *prze1, *prze2;
   ll koszt:
 };
 int n=0, zr, uj;
 const ll inf=1e9;
 vector <vector <kra>> graf:
 vi bylo, aktu;
 vector <ll> odl:
 void vert(int v) {
   if (v>n) {
     graf.resize(n):
     bylo.resize(n);
     aktu.resize(n);
     odl.resize(n):
 }
 void add edge(int v, int u, int prze, ll koszt) {
   vert(v+1); vert(u+1);
   kra ret1{u, new int(prze), new int(0), koszt};
   kra ret2{v, ret1.prze2, ret1.prze1, -koszt};
   graf[v].pb(ret1);
   graf[u].pb(ret2);
 }
 void spfa() {
   for (int i=0; i<n; i++) {</pre>
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aktu[i]=1;
     odl[i]=inf;
    aktu[zr]=odl[zr]=0:
    queue <int> kol;
    kol.push(zr):
    while(!kol.empty()) {
     int v=kol.front();
     kol.pop();
     if (aktu[v])
        continue:
     aktu[v]=1:
     for (kra i : graf[v]) {
       if (*i.przel && odl[v]+i.koszt<odl[i.cel]) {</pre>
          odl[i.cel]=odl[v]+i.koszt;
          aktu[i.cel]=0;
          kol.push(i.cel);
   }
 int dfs(int v) {
   if (v==ui)
     return 1;
    bvlo[v]=1;
    for (int i=0; i < siz(graf[v]); i++) {</pre>
     if (!bylo[graf[v][i].cel] && (*graf[v][i].przel) &&
     odl[v]+graf[v][i].koszt==odl[graf[v][i].cel] && dfs(graf[v][i].cel)) {
        (*graf[v][i].prze1)--;
        (*graf[v][i].prze2)++;
        return 1;
   }
    return 0;
  pair <int,ll> flow(int zrzr, int ujuj) {
   zr=zrzr; ui=uiui;
    vert(zr+1); vert(uj+1);
    pair <int,ll> ret{0, 0};
    while(1) {
     spfa():
     for (int i=0; i<n; i++)</pre>
       bylo[i]=0;
     if (!dfs(zr))
       break:
      ret.st++;
     ret.nd+=odl[uil:
    return ret;
 }
                          Flows And Matchings/Turbo Matching.cpp
struct matching {
 int n:
 vector<vi> V;
 vector<bool> odw:
 vi para, strona;
 matching(int n) : n( n) { // zakladam numeracje od 1, nie moze byc 0
   V.resize(n + 1);
   odw.resize(n + 1):
   para.resize(n + 1);
 void addedge(int a, int b) { // zakladam ze a jest z lewej, b z prawej
    V[a].pb(b);
    strona.pb(a);
```

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bool skojarz(int x) { // x nalezy do strona
   odw[x] = 1;
   for (auto v : V[x]) {
     if (!para[v] || (!odw[para[v]] && skojarz(para[v]))) {
        para[v] = x;
        para[x] = v;
        return 1;
   } }
   return 0;
 }
 vpii go() {
   sort(all(strona));
   strona.resize(unique(all(strona)) - strona.begin());
   for (int i = 1; i \le n; ++i)
     shuffle(all(V[i]), default random engine(2137));
   sort(all(strona), [∆](int a, int b){return siz(V[a]) < siz(V[b]);});
   bool ok = 0:
   do {
     ok = 0;
     for (auto i : strona) odw[i] = 0;
     for (auto i : strona)
        if (!para[i] && skojarz(i))
          ok = 1;
   } while (ok);
   vpii res:
   for (auto i : strona)
     if (para[i])
        res.pb(mp(i, para[i]));
   return res:
} };
                        Flows And Matchings/Weighted Matching.cpp
static const int INF = INT MAX;
static const int N = 1017:
struct edge {
 int u, v, w;
 edge() {} edge(int ui, int vi, int wi) : u(ui), v(vi), w(wi) {}
int n, n_x;
edge q[N^* 2][N * 2];
int lab[N * 2], match[N * 2], slack[N * 2], st[N * 2], pa[N * 2];
int flo from[N * 2][N + 1], S[N * 2], vis[N * 2];
vi flo[\overline{N} * 2];
queue<int> q;
int e delta(const edge \deltae) { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
void update slack(int u, int x) {
 if (!slack[x] \mid | e delta(g[u][x]) < e delta(g[slack[x]][x])) slack[x] = u;
void set slack(int x) {
 slack[x] = 0;
 for (int u = 1; u <= n; u++)
   if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
     update slack(u, x);
void q push(int x) {
 if (x \le n) q.push(x):
 else for (int i = 0; i < siz(flo[x]); i++) q push(flo[x][i]);
void set st(int x, int b) {
 st[x] = b;
 if (x > n) for (int i = 0; i < siz(flo[x]); i++) set st(flo[x][i], b);
int get pr(int b, int xr) {
 int pr = find(all(flo[b]), xr) - flo[b].begin();
 if (pr % 2 == 1) {
```

```
reverse(flo[b].begin() + 1, flo[b].end());
    return siz(flo[b]) - pr;
 } else return pr;
void set match(int u, int v) {
 match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = q[u][v];
  int xr = flo from[u][e.u], pr = get pr(u, xr);
  for (int i = 0; i < pr; i++) set match(flo[u][i], flo[u][i ^{\wedge} 1]);
  set match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
void augment(int u. int v) {
 for (;;) {
    int xnv = st[match[u]];
    set match(u, v):
    if (!xnv) return;
    set match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u:
    vis[u] = t; u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add blossom(int u, int lca, int v) {
  int b = n + 1; while (b \le n \times \&\& st[b]) ++b;
  if (b > n x) ++n x;
  lab[b] = \overline{0}, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear(); flo[b].pb(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].pb(x), flo[b].pb(y = st[match[x]]), q push(y);
  reverse(flo[b].begin() + 1, flo[b].end()):
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].pb(x), flo[b].pb(y = st[match[x]]), q_push(y);
  set st(b, b);
  for (int x = 1; x \le n x; ++x) g[b][x].w = g[x][b].w = 0;
  for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
  for (int i = 0; i < siz(flo[b]); i++) {
    int xs = flo[b][i];
    for (int x = 1; x <= n x; ++x)
      if (q[b][x].w == 0 \mid | e_delta(g[xs][x]) < e_delta(g[b][x]))
        q[b][x] = q[xs][x], q[x][b] = q[x][xs];
    for (int x = 1; x \le n; ++x) if (flo from[xs][x]) flo from[b][x] = xs;
  set slack(b);
void expand blossom(int b) {
  for (int \bar{i} = 0; i < siz(flo[b]); ++i)
    set st(flo[b][i], flo[b][i]);
  int xr = flo from[b][g[b][pa[b]].u], pr = get pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set slack(xns);
    q push(xns);
```

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S[xr] = 1, pa[xr] = pa[b];
 for (int i = pr + 1; i < siz(flo[b]); i++) {</pre>
   int xs = flo[b][i];
   S[xs] = -1, set slack(xs);
 st[b] = 0;
bool on found edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
   pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
   slack[v] = slack[nu] = 0; S[nu] = 0, q push(nu);
 } else if (S[v] == 0) {
   int lca = get lca(u, v);
   if (!lca) return augment(u, v), augment(v, u), true;
   else add blossom(u, lca, v);
 return false:
bool matching() {
 memset(S + 1, -1, sizeof(int) * n x);
 memset(slack + 1, 0, sizeof(int) * n x);
 q = queue<int>();
 for (int x = 1; x \le n x; x++) {
   if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0, q_push(x);
 if (q.empty()) return false;
 for (;;) {
   while (siz(q)) {
     int u = q.front(); q.pop();
     if (S[st[u]] == 1) continue;
     for (int v = 1; v \le n; v++) {
       if (g[u][v].w > 0 \&\& st[u] != st[v]) {
         if (e delta(q[u][v]) == 0) {
           if (on found edge(g[u][v])) return true;
         } else update slack(u, st[v]);
   int d = INF:
   for (int b = n + 1; b \le n \times b + k ) {
     if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2);
   for (int x = 1; x \le n x; x++) {
     if (st[x] == x \&\& slack[x]) {
       if (S[x] == -1) d = min(d, e delta(g[slack[x]][x]));
       else if (S[x] == 0) d = min(d, e delta(g[slack[x]][x]) / 2);
   for (int u = 1; u <= n; u++) {
     if (S[st[u]] == 0) {
       if (lab[u] <= d) return 0;</pre>
       lab[u] -= d:
     } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n \times b + k ) {
     if (st[b] == b) {
       if (S[st[b]] == 0) lab[b] += d * 2;
       else if (S[st[b]] == 1) lab[b] -= d * 2;
   }
   q = queue<int>();
   for (int x = 1; x \le n x; x++)
     if (st[x] == x \&\& slack[x] \&\& st[slack[x]] != x \&\& e delta(g[slack[x]][x]) == 0)
```

```
if (on found edge(g[slack[x]][x])) return true;
    for (int b = n + 1; b \le n \times b + 1)
      if (st[b] == b \&\& S[b] == 1 \&\& lab[b] == 0) expand_blossom(b);
 return false;
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
  n x = n;
  int n matches = 0;
  long \overline{l}ong tot weight = 0;
  for (int u = 0; u \le n; u++) st[u] = u, flo[u].clear();
  int w max = 0;
  for (int u = 1: u \le n: u++)
    for (int v = 1; v \le n; v++) {
      flo from[u][v] = (u == v ? u : 0);
      w \max = \max(w \max, g[u][v].w);
  for (int u = 1; u \le n; u + +) lab[u] = w max;
  while (matching()) ++n matches;
  for (int u = 1; u <= n; u++) {
    if (match[u] \&\& match[u] < u) {
      tot weight += g[u][match[u]].w;
 return mp(tot weight, n matches);
void add edge(int u, int v, int w) {
 g[u][v].w = g[v][u].w = max(g[u][v].w, w);
void init(int n) {
 n = n:
  for (int u = 1; u <= n; u++)
    for (int v = 1; v <= n; v++) {</pre>
      g[u][v] = edge(u, v, 0);
                                    Geometry/Dewolaj.cpp
typedef long long T;
struct P {
 T x, y;
  int id;
 P operator-(P b) { return P{x - b.x, y - b.y}; }
 T cross(P b) { return x * b.y - y * b.x; }
 T cross(P b, P c) const { return (b - *this).cross(c - *this); }
 T dot(P b) { return x * b.x + y * b.y; }
  bool inTriangle(const P &a, const P &b, const P &c) const {
#define tmp(a, b) (cross(a, b) > 0)
    return tmp(a, b) == tmp(b, c) \&\& tmp(b, c) == tmp(c, a);
#undef tmp
 }
int cmpCircle(P a, P b, P c, P d) {
 P v1 = b - a, v2 = d - a, v3 = b - c, v4 = d - c;
 ld tmp=(ld)abs(v1.cross(v2))*v3.dot(v4)+(ld)v1.dot(v2)*abs(v3.cross(v4));
 if (abs(tmp) < 1e-8) return 0; if (tmp == 0) return 0;
 if (tmp > 0) return 1; return -1;
struct pair_hash {
 template <class T1, class T2>
  std::size t operator()(const std::pair<T1, T2> &p) const {
    return p.first * 10000 + p.second:
 }
unordered map<pair<int, int>, pair<int, int>, pair hash> mt;
```

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set<pair<int, int>> edges;
vector<vector<int>>> triangles;
void rec(int a, int c, const vector<P> &points);
void trim(int a, int b) {
 assert(a < b);</pre>
 auto it = mt.find({a, b});
 if ((it != mt.end()) \&\& (it->second == make pair(-1, -1))) mt.erase(it);
/oid change(int a, int b, int from, int to) {
 if (a > b) swap(a, b);
 if (!mt.count(\{a, b\})) mt[\{a, b\}] = \{-1, -1\};
 for (int *x : vector<int *>{&mt[{a, b}].first, &mt[{a, b}].second})
   if (*x == from) {
     *x = to:
     trim(a, b);
     return;
 assert(false);
/oid rec(int a, int c, const vector<P> &points) {
 if (a > c) swap(a, c);
 if (!mt.count({a, c})) return;
 int b = mt[{a, c}].first;
 int d = mt[{a, c}].second;
 if (b > d) swap(b, d);
 if (b == -1 | | d == -1) return;
 for (int rep = 0; rep < 2; ++rep) {
   if (points[a].inTriangle(points[b], points[c], points[d])) return;
 if (cmpCircle(points[a], points[b], points[c], points[d]) != 1) return;
 assert((!mt.count(\{b, d\})) \mid | (mt[\{b, d\}] == make pair(-1, -1)));
 mt[{b, d}] = {a, c};
 trim(b, d);
 mt.erase(make pair(a, c));
 change(a, b, c, d);change(b, c, a, d);change(a, d, c, b);change(c, d, a, b);
 rec(a, b, points);rec(b, c, points);rec(c, d, points);rec(d, a, points);
/oid addTriangle(int a, int b, int c) {
 change(a, b, -1, c); change(a, c, -1, b); change(b, c, -1, a);
/oid anyTriangulation(vector<P> points) {
 sort(points.begin(), points.end(), [](const P &a, const P &b) {
   return make pair(a.x, a.y) < make pair(b.x, b.y);});</pre>
 vector<P> upper, lower;
 for (P C : points) {
#define backback(w) w[(int)w.size() - 21
   while ((int)upper.size() \geq 2 \&\& backback(upper).cross(upper.back(), C)>0)
     addTriangle(C.id, backback(upper).id, upper.back().id);
     upper.pop back();
   upper.push back(C);
   while ((int)lower.size() >= 2 \&\&
       lower[(int)lower.size() - 2].cross(lower.back(), C) < 0) {</pre>
     addTriangle(C.id, backback(lower).id, lower.back().id);
     lower.pop back():
   lower.push back(C);
#undef backback
 }//all collinear
 if(lower.size()==upper.size()&&lower.size()==points.size())assert(false);
const int nax = 1e6 + 5;
int memo x[nax], memo y[nax];
void dewolaj() {
```

```
mt.clear();edges.clear();triangles.clear();
  scanf("%d", &n);
  vector<P> points(n);
  for (int i = 0; i < n; ++i) {
    scanf("%d%d". &memo x[i]. &memo v[i]):
    points[i] = P\{\text{memo } x[i], \text{ memo } y[i], i\};
  mt.reserve(4123123);
  anyTriangulation(points);
  vector<pair<int, int>> init;
  for (auto ppp : mt) init.push back(ppp.first);
  for (pair<int, int> p : init)
    if (mt.count(p) \&\& mt[p] != make pair(-1, -1)) rec(p.first, p.second, points);
  for (auto ppp : mt)
    if (ppp second != make pair(-1, -1)) {
     int i = ppp.first.first, j = ppp.first.second;
      assert(i != i);
      edges.insert({i, j});
      if (mt[{i, j}].first > j)
        triangles.push back(vector<int>{i, j, mt[{i, j}].first});
      if (mt[{i, j}].second > j)
        triangles.push back(vector<int>{i, j, mt[{i, j}].second});
  // edges zawiera krawędzie triangulacji
                                  Geometry/Halfplanes.cpp
pii operator-(pii a, pii b) {return {a.first - b.first, a.second - b.second};}
pii rot(pii x) { return {x.second, -x.first}; }
using hpl = pair <pii, ll>; //(v, m) = \{x : sc(x, v) >= m\}
ll sc(pii a, pii b) {return a.first*1ll*b.first + a.second*1ll*b.second;}
ll ve(pii a, pii b) {return a.first*1ll*b.second - a.second*1ll*b.first:}
//wartości rzędu współrzędne ^ 4
 int128 det(hpl a, hpl b, hpl c) {
 return a.second * ( int128)ve(b.first, c.first) +
   b.second * (\underline{\text{int}}\overline{128})ve(c.first, a.first) +
    c.second * ( int128)ve(a.first, b.first);
bool subset(hpl a, hpl b) {
 if (ve(a.first, b.first) || sc(a.first, b.first) < 0) return false;</pre>
  return a.second * ( int128) abs(b.first.first) >=
     b.second * (\underline{\text{int}}128) abs(a.first.first) &&
      a.second * ( int128) abs(b.first.second) >=
     b.second * ( int128) abs(a.first.second);
bool disjoint(hpl a, hpl b) {
 return subset(a, {{-b.first.first, -b.first.second}, -b.second}); }
bool around(hpl a, hpl b, hpl c) {
 ll ab = ve(a.first, b.first):
 ll bc = ve(b.first, c.first);
 ll ca = ve(c.first, a.first);
 assert(ab > 0 || bc > 0 || ca > 0);
  return (ab \geq 0 && bc \geq 0 && ca \geq 0) || (ab \leq 0 && bc \leq 0 && ca \leq 0);
void ang sort(vector <hpl> &a) {
 if (a.emptv()) return:
  hpl mid = a.back();
  a.pop back();
  vector <hpl> left, right;
  for (hpl c : a) {
   ll v = ve(c.first, mid.first), s = sc(c.first, mid.first);
    if (make pair(v, s) > pair < ll. ll > (0, 0))
      left.push back(c);
    else
      right.push back(c);
```

```
sort(right.begin(), right.end(), [](hpl x, hpl y){return ve(x.first, y.first) > 0;});
 sort(left.begin(), left.end(), [](hpl x, hpl y){return ve(x.first, y.first) > 0;});
 left.push back(mid):
 left.insert(left.end(), right.begin(), right.end());
 swap(a, left):
//Jeśli przecięcie jest nieograniczone, to półproste idące do nieskończoności
//sa wvznaczone
//przez pierwszy i ostatni element zwróconego wektora
//Jeśli przecięcie jest puste, zwraca pusty wektor
vector <hpl> find hull(vector <hpl> vec) {
 vector <hpl> hull;
 int first = 0:
 ang sort(vec);
 for (hpl curr : vec) {
   if (!hull.empty() && disjoint(curr, hull.back()))
     return {}; //Przecięcie jest puste
   if (!hull.empty() && disjoint(curr, hull[first]))
     return {}; //Przeciecie jest puste
   if (!hull.empty() && subset(hull.back(), curr)) //Case kiedy jedna półpłaszczyzna
     continue; //zawiera się w drugiej, na ogół można wywalić
   if (!hull.empty() && subset(curr, hull.back())) //J.w.
     hull.pop back();
   while (hull.size() - first >= 2u &&
        det(hull.back(), *(hull.end() - 2), curr) <= 0) {
     if (around(hull[hull.size() - 2], hull.back(), curr))
        return {}; //Przeciecie jest puste
       hull.pop back();
   while (hull.size() - first \geq 2u && det(curr, hull[first], hull[first + 1]) \geq 0) {
     if (around(curr, hull[first], hull[first + 1]))
        return {}; //Przeciecie jest puste
     else
       first++:
   if (hull.size() - first < 2u || det(hull.back(), curr, hull[first]) < 0)</pre>
     hull.push back(curr);
 return vector <hpl>(hull.begin() + first, hull.end());
//Półpłaszczyzna wyznaczona przez prostą xy
hpl make hpl(pii x, pii y) {
 pii v = rot(y - x);
 assert(sc(v, x) == sc(v, y));
 return \{v, sc(x, v)\};
                              Geometry/Halfplanes Online.cpp
#define X real()
#define Y imag()
typedef complex<ll> P;
struct line {
 ll a.b.c:
 line(ll a = 0, ll b = 0, ll c = 0): a(a), b(b), c(c)
   {} // <= 10^9
 line (P const &A, P const &B): a(A.Y-B.Y), b(B.X-A.X),
   c(A.X*B.Y-A.Y*B.X) {} //pts <= 10^6
 line operator - () const {return line(-a, -b, -c); }
 bool up() const { return a?(a<0):(b>0);}
inline ll wek(line const &a. line const &b) {return a.a*b.b-a.b*b.a:}
inline bool rown(line a, line b) {return wek(a,b) == 0;}
inline bool pokr(line a, line b)
 {return rown(a,b) && a.a*b.c == b.a*a.c && a.b*b.c == b.b*a.c;}
```

```
inline bool podobne(line a, line b) {return rown(a,b) && a.up() == b.up();}
inline complex<ld> prosta prosta(line a, line b) {
 11 det = wek(a,b);
 ll x = -a.c*b.b+b.c*a.b:
 ll y = -a.a*b.c+a.c*b.a;
  return complex<ld>(x,v)/(ld)det:
inline ll weaker (line a, line b) { // czy a jest slabsze niz b
  assert(rown(a,b));
 if (abs(a.a) > abs(a.b)) return a.c*abs(b.a) - b.c*abs(a.a);
 else return a.c*abs(b.b) - b.c*abs(a.b);
struct Comp {
  bool operator()(const line& a, const line& b) const {
    if (a.up() != b.up()) return a.up() > b.up();
    return wek(a,b) > 0;
|};
const ld EPS = 1e-12:
struct przeciecie polplaszczyzn {
  bool empty, pek;
  set<line, Comp> S;
  typedef set<line. Comp>::iterator iter:
  przeciecie polplaszczyzn() : empty(false), pek(false) {};
  iter next(iter it){return (++it == S.end() ? S.begin() : it);}
  iter prev(iter it){return (it == S.begin() ? --S.end() : --it);}
  bool hide(line a, line b, line c) {
    if (rown(a,b)) {
      if (weaker(a, -b) < 0) empty = true;
      return false:
    if (wek(a,b) < 0) swap(a,b):
    complex<ld> r = prosta prosta(a,b);
    ld v = r.X * c.a + r.Y * c.b + c.c;
    if (wek(a,c) >= 0 \&\& wek(c,b) >= 0 \&\& v > -EPS) return true;
    if (wek(a,c) < 0 \&\& wek(c,b) < 0) {
      if (v < -EPS) empty = true;</pre>
      else if (v < EPS) pek = true;</pre>
    return false;
  void add(line l) {
    if (empty) return;
    if (l.a == 0 && l.b == 0) {
      if (l.c < 0) empty = true;
      return;
    iter it = S.lower bound(l);
    //rownoleale
    if(it != S.end() && podobne(*it, l)) {
      if (weaker(l, *it)>=0) return;
      iter del = it:
      it = next(it);
      S.erase(del):
    //*it>p
    if(siz(S) >= 2 \&\& it == S.end()) it = S.begin();
    while(siz(S) \Rightarrow 2 && hide(l, *next(it), *it)) {
      iter del = it;
      it = next(it):
      S.erase(del);
    //*it<p
    if(siz(S) >= 2) it = prev(it);
    while(siz(S) \geq 2 && hide(l, *prev(it), *it)) {
```

```
iter del = it;
     it = prev(it);
     S.erase(del);
   if(S.size() < 2 || !hide(*it, *next(it), l)) S.insert(l);</pre>
     0 - puste 1 - punkt 2 - odcinek 3 - półprosta 4 - prosta
    5 - dodatnie (może nieskończone) pole (S.size() daje wowczas liczbę boków) */
 int type() {
   if(empty) return 0;
   if(siz(S) \ll 4)
     vector<line> res(all(S)):
     if (siz(res) == 2 \&\& rown(res[0], res[1]) \&\& weaker(res[0], -res[1])<0)
     for (int i = 0; i < siz(res); ++i) for (int j = 0; j < i; ++j)
       if(pokr(res[i], res[j])) {
         if(siz(res) == 2) return 4:
         if(siz(res) == 3) return 3;
         if(siz(res) == 4 \&\& pokr(res[0], res[2]) \&\& pokr(res[1], res[3])) return 1;
          return 2;
     if(siz(res) == 3 \&\& pek) return 1;
   return 5;
                                    Geometry/Rotor.cpp
pii in[nax];
int wh[nax];
pair <pii, pii> dir[nax * nax / 2];
pii operator-(pii a, pii b) {return {a.first - b.first, a.second - b.second;}}
ll pro(pii a. pii b) {return a.first*1ll*b.second - a.second*1ll*b.first;}
bool cmp(pair <pii, pii> a, pair <pii, pii> b) {
 ll p = pro(a.first, b.first):
 if (p > 0) return 1; if (p < 0) return 0; return a.second < b.second;
 //Jak nie ma trzech współliniowych to po prostu: return pro(a.first, b.first) > 0;
int main() {
 scanf("%d", \&n);
 for (int i = 0; i < n; ++i)
   scanf("%d%d", &in[i].first, &in[i].second);
 sort(in, in + n);
 for (int i = 0; i < n; ++i)
   wh[i] = i;
 int cou = 0:
 for (int i = 0; i < n; ++i)
   for (int i = i + 1; i < n; ++i)
     dir[cou++] = {max(in[i] - in[j], in[j] - in[i]), {i, j}};
 sort(dir, dir + cou, cmp);
 for (int i = 0; i < cou; ++i) {
   debug << imie(i) imie(dir[i]);</pre>
   auto c = dir[i].second;
   pii x = in[wh[c.first]], y = in[wh[c.second]];
   swap(wh[c.first], wh[c.second]);
   swap(in[wh[c.first]], in[wh[c.second]]);
   //Policz wynik dla posortowanych prostopadle do prostei xy
 }
                                     Graphs/2sat.cpp
const int nax = 100000:
vector <int> implies[2*nax]; //wymuszenia, 2*x to zmienna 2*x+1 to jej negacja
int sat val[2*nax].sat vis[2*nax].sat sort[2*nax].sat ile:
inline void sat or(int a,int b){
 implies[a^1].push back(b);
 implies[b^1] push back(a);
```

```
void sat dfs mark(int x){
 sat vis[x]=1;
 sat val[x]=sat val[x^1]==-1:
 for (int i : implies[x]) if (!sat_vis[i]) sat_dfs_mark(i);
void sat dfs(int x){
 sat vis[x]=1:
 for (int i : implies[x]) if (!sat vis[i]) sat dfs(i);
 sat sort[--sat ile]=x^1;
bool sat2(int n) {//n - liczba zmiennych, zmienne numerujemy od 0
#define REP for (int i = 0; i < 2 * n; ++i)
 sat ile=2 * n:
 REP sat vis[i]=0,sat val[i]=-1;
 REP if (!sat vis[i]) sat dfs(i);
 REP sat vis[i]=0:
  REP if (!sat vis[sat sort[i]]) sat dfs mark(sat sort[i]);
 REP if (sat val[i]) for(int j : implies[i]) if (!sat val[j]) return 0;
  return 1;
                                   Graphs/Centroid.cpp
struct centro { // indeksowane od 0, par to drzewo centroidow
 vector<vi> edges;
 vector<bool> vis:
 vi par, sz;
  int n:
  centro(int n) : n(n) {
   edges.resize(n);
   vis.resize(n):
   par.resize(n);
   sz.resize(n):
 void edge(int a, int b) {
   edges[a].pb(b);
   edges[b].pb(a);
 int find size(int v. int p = -1) {
   if (vis[v]) return 0;
   sz[v] = 1;
   for (int x: edges[v]) {
     if (x != p) {
       sz[v] += find size(x, v);
   return sz[v];
 int find centroid(int v, int p, int n) {
   for (int x: edges[v]) {
     if (x != p) {
       if (!vis[x] && sz[x] > n / 2) {
          return find centroid(x, v, n);
   return v:
  void init centroid(int v = 0, int p = -1) {
   find size(v);
   \overline{c} = find centroid(v, -1, sz[v]);
   vis[c] = true;
   par[c] = p:
   for (int x: edges[c]) {
     if (!vis[x]) {
       init centroid(x, c);
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}
                                     Graphs/DMST.cpp
struct RollbackUF {
 vi e;
 vpii st:
 RollbackUF(int n) : e(n, -1) {}
 int size(int x) { return -e[find(x)]; }
 int find(int x) { return e[x] < 0 ? x : find(e[x]); }
 int time() { return siz(st); }
 void rollback(int t) {
   for (int i = time(): i-- > t:) e[st[i].first] = st[i].nd:
   st.resize(t);
 bool unite(int a. int b) {
   a = find(a), b = find(b);
   if (a == b) return false:
   if (e[a] > e[b]) swap(a, b);
   st.pb({a, e[a]});
   st.pb({b, e[b]});
   e[a] += e[b];
   e[b] = a;
   return true:
 }
struct Edge {
 int a, b;
 ll w:
struct Node { /// lazv skew heap node
 Edge key;
 Node *l, *r;
 ll delta;
 void prop() {
   key.w += delta;
   if (l) l->delta += delta:
   if (r) r->delta += delta;
   delta = 0;
 Edge top() {
   prop();
   return key;
Node* merge(Node* a, Node* b) {
 if (!a || !b) return a ?: b;
 a->prop(), b->prop():
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->l, (a->r = merge(b, a->r)));
 return a:
void pop(Node*& a) {
 a->prop();
 a = merge(a->l, a->r):
} // wierzcholki numerujemy od 0, r to korzen dmst
pair<ll, vi> dmst(int n, int r, vector<Edge>& g) {
 RollbackUF uf(n);
 vector<Node*> heap(n):
 for (Edge e : g) heap[e.b] = merge(heap[e.b], new Node{e});
 ll res = 0:
 vector<int> seen(n, -1), path(n), par(n);
 seen[r] = r;
 vector<Edge> Q(n), in(n, \{-1, -1\}), comp;
```

```
deque<tuple<int, int, vector<Edge>>> cycs;
  for (int s = 0; s < n; ++s) {
   int u = s, qi = 0, w:
   while (seen[u] < 0) {
     if (!heap[u]) return {-1, {}};
     Edge e = heap[u]->top();
     heap[u]->delta -= e.w, pop(heap[u]);
     Q[qi] = e, path[qi++] = u, seen[u] = s;
     res += e.w, u = uf.find(e.a);
     if (seen[u] == s) { /// found cycle, contract
       Node* cvc = 0:
       int end = qi, time = uf.time();
       do cyc = merge(cyc, heap[w = path[--qi]]);
       while (uf.unite(u, w)):
       u = uf.find(u), heap[u] = cyc, seen[u] = -1;
       cycs push front(\{u, time, \{\&Q[qi], \&Q[end]\}\}\);
   for (int i = 0: i < ai: ++i) in[uf.find(0[i].b)] = 0[i]:
 for (auto& [u, t, comp] : cycs) { // restore sol (optional)
   uf.rollback(t);
   Edge inEdge = in[u]:
   for (auto& e : comp) in[uf.find(e.b)] = e;
   in[uf.find(inEdge.b)] = inEdge;
 for (int i = 0; i < n; ++i) par[i] = in[i].a;
 return {res, par};
                                  Graphs/Dominatory.cpp
struct Dominators {
 int n oria, n:
 vector<int> parent, semi, vertex, dom, ancestor, label;
 vector<vector<int>> succ, pred, bucket;
 Dominators(int n): n orig(n), n(n + 1), parent(n), semi(n), vertex(n), dom(n),
     ancestor(n), label(n), succ(n), pred(n), bucket(n) {
   n = n \text{ orig};
 void add edge(int a, int b){
   a++; b++; succ[a].push back(b);
 void COMPRESS(int v) {
   if (ancestor[ancestor[v]] != 0) {
     COMPRESS(ancestor[v]);
     if (semi[label[ancestor[v]]] < semi[label[v]]) label[v] = label[ancestor[v]];</pre>
     ancestor[v]=ancestor[ancestor[v]]; } }
 void LINK(int v. int w) { ancestor[w]=v: }
 int EVAL(int v) {
   if(ancestor[v] == 0) return v:
   else {
     COMPRESS(v):
     return label[v]:} }
 void DFS(int v) {
   semi[v] = ++n: vertex[n] = v:
   for(auto ng : succ[v]) {
     if(semi[na] == 0) {
       parent[ng]=v;
       DFS(ng);
     pred[ng].push back(v); } }
 // Zwraca vector dominatorow (-1 dla 0). // Wszystko numerowane od 0,
 vector<int> doit() { // wszystko musi być osiągalne z 0.
   iota(label.begin(), label.end(), 0);
   n = 0; DFS(1);
   for (int i = n; i >= 2; --i) {
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int w = vertex[i]:
     for (auto ng : pred[w]) {
       int u = EVAL(ng);
       if (semi[u] < semi[w]) { semi[w] = semi[u]: }</pre>
     bucket[vertex[semi[w]]].push back(w):
     LINK(parent[w],w);
     while (!bucket[parent[w]].empty()) {
        int v = bucket[parent[w]].back();
       bucket[parent[w]].pop back();
       int u = EVAL(v):
       if (semi[u] < semi[v]) dom[v] = u;</pre>
       else dom[v] = parent[w];
   for (int i = 2; i <= n; ++i) {
     int w = vertex[i]:
     if (dom[w] != vertex[semi[w]]) { dom[w] = dom[dom[w]]; }
   dom[1] = 0;
   vector<int> res(n orig);
   for (int i = 0; i < n orig; ++i) res[i] = dom[i + 1] - 1;
   return res; } };
                                  Graphs/Euler_Cycle.cpp
vi ans:
set<int> S[N]:
void rec(int v) {
 while(!S[v].empty()) {
   auto u = *S[v].begin();
   S[v].erase(u):
   S[u].erase(v);
   rec(u):
   ans.pb(v);
void eul() {
 vi deq;
 for(int i = 1; i \le n; ++i) if(S[i].size() % 2) deg.pb(i);
 if (siz(deg) == 0) {
   ans.pb(1);
   rec(1):
 } else if (siz(deg) == 2) {
   sort(all(deg));
   rec(deg[0]);
   reverse(all(ans));
   ans.pb(deg[1]);
 } else assert(0); //Cycle doesn't exist
                                      Graphs/HLD.cpp
vi drz[nax];
int roz[nax], jump[nax], pre[nax], post[nax], fad[nax], czas;
void dfs roz(int v) {
 roz[v]=1;
                                   // drz[] ma nie zawierać krawedzi
 for (int &i : drz[v]) {
                                   // do oica.
   fad[i]=v;
                                 // Init:
   dfs roz(i):
                                 // dfs roz(root):
   roz[v]+=roz[i];
                                 // dfs pre(root);
   if (roz[i]>roz[drz[v][0]]) // Użycie get path(v, u) zwróci
     swap(i, drz[v][0]);
                               // przedziały odpowiadające ścieżce
                                   // z v do u. Przedziały odpowiadające
                                     // ścieżce z v do lca mają
void dfs pre(int v)
                                     // first>=second, zaś te dla ścieżki
                                     // z lca do u maja first<=second.
 if (!jump[v])
                                   // Przedziały są po kolei.
   jump[v]=v;
                                 // Lca występuje w nich dwa razy,
```

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pre[v]=(++czas);
                                   // najpierw jako second,
  if (!drz[v].empty())
                                   // a zaraz potem jako first.
   jump[drz[v][0]]=jump[v];
  for (int i : drz[v])
   dfs pre(i);
  post[v]=czas:
int lca(int v. int u) {
 while(jump[v]!=jump[u]) {
   if (pre[v]<pre[u])</pre>
      swap(v, u);
   v=fad[jump[v]];
 return (pre[v]<pre[u] ? v : u):
vpii path up(int v, int u) {
 vpii ret:
  while(jump[v]!=jump[u]) {
    ret.pb({pre[jump[v]], pre[v]});
   v=fad[jump[v]];
  ret.pb({pre[u], pre[v]});
  return ret:
vpii get path(int v, int u) {
  int w=lca(v. u):
  auto ret=path up(v, w):
  auto pom=path up(u, w);
  for (auto &i : ret)
   swap(i.st. i.nd):
  while(!pom.empty()) {
   ret.pb(pom.back()):
   pom.pop back();
  return ret;
                                   Graphs/Link Cut.cpp
struct Splay {
  Splay *l = 0, *r = 0, *p = 0;
  bool flip = false; // Wywal jak nie używasz make root.
  void update() {
   //update anything stored for nodes
    assert(!flip and (!l or !l->flip) and (!r or !r->flip)); }
  void touch() { //Do any lazy prop here
   if (flip) {
      swap(l, r);
      if (l) l->flip = !l->flip:
      if (r) r->flip = !r->flip;
      flip = false;
  bool sroot() { return !p or (p->l != this and p->r != this); }
  void connect(Splay* c, bool left) { (left ? l : r) = c; if (c) c->p = this; }
  void rotate() {
   Splay* f = p; Splay* t = f - p;
    const bool isr = f->sroot():
    const bool left = (this == f->l);
    f->connect(left ? r : l, left);
    connect(f, !left);
    if (isr) p = t;
   else t->connect(this, f == t->l);
   f->update():
  void push(){sroot()?touch():p->push();if (l) l->touch();if (r) r->touch();}
  void splay() {
```

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push();
   while (!sroot()) {
     Splay* x = p->p;
     if ([p->sroot()) (([p->l == this) == (x->l == p)) ? p : this)->rotate():
   update();
 Splay* expose(){//v będzie korzeniem splaya zawierającego ścieżkę do korzenia
   Splay *q = this, *x = 0;//prawe dziecko bedzie nullem. Jak zejdziemy w dół,
   while (g) { // to potem trzeba zrobić splay().
     q->splay();// LCA(u, v): u->expose(); return v->expose();
     q->r = x; q->update();
     x = q; q = q -> p;
   splay();
   return x:
 Splay* root() { // Zwraca roota drzewowego (nie spleiowego!).
   Splay* s = this;
   while (s->touch(), s->l) s = s->l;
   s->splav():
   return s;
 void cut() { // Usuwa krawedź znajdującą sie nad danym wierzchołkiem.
   expose(): assert(l /* Nie iest rootem. */):
   Splay* s = l;
   while (s->touch(), s->r) s = s->r;
   s->splay(); s->r->p = 0; s->r = 0;
 void link(Splay* to) {
   expose(); assert(!l /* Jest rootem. */);
   p = to:
 // Sprawia, że wierzchołek jest rootem w logicznym i w splayowym drzewie.
 void make root() { expose(); flip = !flip; touch(); }
                                     Graphs/Low.cpp
struct Low { // dziala dla multikrawedzi, petli, niespojnego grafu
 int n. cnt. edges:
 vector<vpii> G;
 vi low, pre, par, par nr, used, most, root, vis s, add s;
 Low(int n) : n(n), edges(0) {
   G.resize(n + 1);
   low.resize(n + 1);
   pre.resize(n + 1):
   par.resize(n + 1);
   par nr.resize(n + 1):
   root.resize(n + 1);
 void edge(int a, int b) {
   ++edges;
   G[a].eb(b. edges):
   G[b].eb(a, edges);
 void dfs(int v) {
   pre[v] = ++cnt:
   low[v] = pre[v];
   for (auto it : G[v]) {
     int u = it.st;
     int nr = it.nd:
     if (used[nr]) continue;
     used[nr] = 1;
     if (!pre[u]) {
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par[u] = v;
      par nr[u] = nr;
      dfs(u):
      mini(low[v], low[u]);
      mini(low[v], pre[u]);
void go() { // trzeba wywolac na poczatku!
  used.resize(edges + 1):
  for(int i = 1; i \le n; ++i) if (!pre[i]) root[i] = 1, dfs(i);
vpii mosty() {
 most.resize(edges + 1):
  vpii ans;
  for(int i = 1; i <= n; ++i) {</pre>
   if (!root[i] && low[i] == pre[i]) {
     ans.eb(i, par[i]);
      most[par nr[i]] = 1;
  return ans;
vi pkt art() { // tylko jesli potrzebujemy
 vi ans, take(n + 1), root sons(n + 1);
  for(int i = 1; i <= n; ++i) {
   if (par[i] && root[par[i]]) {
      ++root sons[par[i]];
  for(int i = 1; i <= n; ++i) {</pre>
   if (root[i] \&\& root sons[i] >= 2) take[i] = 1;
   if (!root[i] \&\& !root[par[i]] \&\& low[i] >= pre[par[i]]) take[par[i]] = 1;
  for(int i = 1; i <= n; ++i) if (take[i]) ans.pb(i);</pre>
  return ans:
// kod nizei tylko jesli potrzebujemy 2spojnych
using comps = vector<vpii>;
void dfs s(int v, vpii &moja) {
  vis s[v] = 1:
  for (auto it : G[v]) {
   int u = it.st:
   int nr = it.nd;
   if (most[nr]) continue;
   if (!add s[nr]) {
     moia.eb(v. u):
     add s[nr] = 1;
   if (!vis s[u]) dfs s(u, moja);
comps bico() {
 vis s.resize(n + 1):
  add s.resize(edges + 1);
  comps ans:
  for (auto it : mosty()) ans.pb({it});
  for(int i = 1; i \le n; ++i) {
   if (!vis s[i]) {
     vpii curr:
      dfs s(i, curr);
     if (!curr.emptv()) ans.pb(curr):
 }
  return ans;
```

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Graphs/Planarity Check.cpp
struct FU {
                                     //numeracia od zera
 vi p; vector<bool> b;
                                     //nie dawać multikrawedzi ani petli
 FU() {}
                                     //nie odpalać IsPlanar() ponad raz
 FU(int n) : p(n), b(n) { iota(p.begin(), p.end(), 0); }
 pair<int, bool> Find(int v) {
   if (p[v] == v) \{ return \{v, \theta\}; \}
   auto res = Find(p[v]); res.second ^= b[v];
   p[v] = res.first; b[v] = res.second; return res;
 bool Union(int x, int y, bool flip) {
   bool xb, yb; tie(x, xb) = Find(x); tie(y, yb) = Find(y);
   if (x == y) { return !(xb ^ yb ^ flip); }
   p[x] = y; b[x] = xb ^ yb ^ flip; return true;
struct PlanarityTest {
 int N, M, tm;
 vector<vi> adj, dn, up; vector<pair<int,int>> e up;
 vector<bool> vis; vi low, pre;
 FU fu:
 void DfsLow(int v, int p) {
   vis[v] = true; low[v] = pre[v] = tm++;
   for (int s : adj[v]) {
     if (s == p) { continue; }
     if (!vis[s]) {
       dn[v].pb(s); DfsLow(s, v); low[v] = min(low[v], low[s]);
     } else if (pre[s] < pre[v]) {</pre>
        up[v].pb(siz(e up)); e up.pb({v, s});
        low[v] = min(low[v], pre[s]);
   }
 vi Interlace(const vi &ids, int lo) {
   for (int e : ids) if (pre[e up[e].second] > lo) ans.pb(e);
   return ans;
 bool AddFU(const vi &a. const vi &b) {
   for (int k = 1; k \le siz(a) - 1; ++k) if (!fu.Union(a[k - 1], a[k], 0)) { return 0; }
   for (int k = 1; k \le siz(b) - 1; ++k) if (!fu.Union(b[k - 1], b[k], 0)) { return 0; }
   if (siz(a) \&\& siz(b) \&\& !fu.Union(a[0], b[0], 1)) { return 0; }
   return 1:
 bool DfsPlanar(int v, int p) {
   for (int s : dn[v]) if (!DfsPlanar(s, v)) { return false; }
   auto sz = siz(dn[v]):
   for (int i = 0; i < sz; ++i) {
     for (int j = i + 1; j < sz; ++j)
        auto a = Interlace(up[dn[v][i]], low[dn[v][j]]);
       auto b = Interlace(up[dn[v][j]], low[dn[v][i]]);
       if (!AddFU(a, b)) { return false; }
     for (int j : up[v]) {
       if (e up[j].first != v) { continue; }
       auto a = Interlace(up[dn[v][i]], pre[e up[j].second]);
       auto b = Interlace({i}, low[dn[v][i]]);
       if (!AddFU(a, b)) { return false; }
   for (int s : dn[v]) {
     for (int idx : up[s]) {
       if (pre[e up[idx].second] < pre[p]) { up[v].pb(idx); }</pre>
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```
up[s].clear(); up[s].shrink to fit();
   return true:
 PlanarityTest(int n) : N(n), M(0), adj(n) {}
 void AddEdge(int u, int v) { adj[u].pb(v); adj[v].pb(u); ++M; }
 bool IsPlanar() {
   if (N <= 3) { return true; }
   if (M > 3 * N - 6) { return false; }
   vis = vector<bool>(N);
   up = dn = vector<vi>(N):
   low = pre = vi(N);
   for (int v = 0; v < N; ++v) if (!vis[v]) {
     e up.clear(); tm = 0; DfsLow(v, -1);
     fu = FU(siz(e up));
     if (!DfsPlanar(v, -1)) { return false; }
   return true:
 }
                                   Graphs/Smulewicz.cpp
struct SPFA{ // #define int ll//jeśli long longi potrzebne
 int n; vi odl, oj, czok;
 vector<vpii> d; vector<vi> d2;
 const int inf = 1e9;
 SPFA(int n):n(n+1){
   odl.resize(n, inf); oj.resize(n); czok.resize(n);
   d.resize(n); d2.resize(n);
 vi cykl; int root;
 bool us(int nr){
   if(nr == root) return 1;
   czok[nr] = 0:
   for(int ak:d2[nr]){
     if(oj[ak] == nr){
       if(us(ak)){
          cykl.pb(nr);
                                        //numeracja od zera
         return 1;
   d2[nr].clear();
   return 0;
 bool licz sciezki(int s){ // false, gdy z s da sie dojsc do ujemnego cyklu
   vi st, st2; // znaleziony cykl jest w wektorze cykl
   odl[s] = 0; czok[s] = 1; st.pb(s);
   while(siz(st)){
     for(int i=0; i<siz(st); i++){</pre>
       int ak = st[i];
       if(czok[ak]) for(pii x:d[ak]){
         int nei, cost; tie(nei, cost) = x;
         if(odl[ak] + cost < odl[nei]){</pre>
            root = ak;
            if(us(nei)){
             cykl.pb(ak); reverse(all(cykl));
             return 0;
            odl[nei] = odl[ak] + cost; oj[nei] = ak;
           d2[ak].pb(nei); czok[nei] = 1;
            st2.pb(nei):
         }
```

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st.clear(); swap(st, st2);
   return 1;
 vi ujemny cykl(){
   for (int i=0; i<n-1; i++) add edge(n-1, i, 0);</pre>
   if (licz sciezki(n-1)){
     return {};
   } else {
      return cykl;
 }
 void add edge(int a, int b, int cost){
   d[a].pb({b, cost});
}; // #undef int
                                      Graphs/SCC.cpp
struct Scc {
 using graph = vector<vi>;
 int n, cnt; // spoine w porzadku topo 1 ... cnt
 graph G, T;
 vi vis, comp, post; // comp - numer spojnej wierzcholka
 Scc(int n) : n(n), cnt(0) {
   G.resize(n + 1);
   T.resize(n + 1);
   vis.resize(n + 1);
   comp.resize(n + 1):
 void edge(int a, int b) {
   G[a].pb(b);
   T[b].pb(a);
 void dfs(int v, graph &gra) {
   vis[v] = 1:
   if (\&gra == \&T) comp[v] = cnt;
   for (auto it : gra[v]) if (!vis[it]) dfs(it, gra);
   if (\&gra == \&G) post.pb(v);
 void go() {
   for (int i = 1; i \le n; ++i) if (!vis[i]) dfs(i, G);
   reverse(all(post)):
   fill(all(vis), 0);
   for (auto it : post) if (!vis[it]) ++cnt, dfs(it, T);
 }
                                      Graphs/SSS.cpp
const int nax=100*1007:
vector <int> graf[nax], farg[nax];
int ost[nax], bylo[nax], post[nax], spo[nax], counter, coudfs;
vector <vector<pair<int,int>>> mer;
void dfs1(int v) {
 if (bylo[v])
   return;
 bvlo[v]=1:
 for (int i : graf[v]) dfs1(i);
 coudfs--:
 post[coudfs]=v;
void dfs2(int v, int s) {
 if (spo[v] >= 0)
  return;
 spo[v]=s:
 for (int i : farg[v]) dfs2(i, s);
void rek(int l, int r, vector <pair<pair<int,int>,int>> &kra) {
```

```
if (l>r) return;
  counter++:
  vector <int> ver:
  for (auto i : kra) {
   if (ost[i.first.first]<counter) {</pre>
     ver.push back(i.first.first):
     ost[i.first.first]=counter;
    if (ost[i.first.second]<counter) {</pre>
     ver.push back(i.first.second);
     ost[i.first.second]=counter;
 for (int i : ver) {
   bylo[i]=0; spo[i]=-1;
   graf[i].clear(); farg[i].clear();
  int s=(l+r)>>1;
 for (auto i : kra) {
   if (i.second<=s) {</pre>
     graf[i.first.first].push back(i.first.second);
      farg[i.first.second].push back(i.first.first);
 }
 coudfs=ver.size():
 for (int i : ver)
   dfs1(i):
 for (int i=0; i<(int)ver.size(); i++)</pre>
   dfs2(post[i], post[i]);
  for (int i : ver)
   if (i!=spo[i])
      mer[s].push back({i, spo[i]});
  vector <pair<pair<int,int>,int>> lew, pra;
  for (auto i : kra) {
   if (spo[i.first.first] == spo[i.first.second])
     lew.push back(i);
      pra.push back({{spo[i.first.first],spo[i.first.second]}, i.second});
 rek(l, s-1, lew); rek(s+1, r, pra);
void sss(vector <pair<int,int>> kra) {
 mer.clear();
 mer.resize(kra.size());
  vector <pair<int,int>,int>>daj;
  for (int i=0; i<(int)kra.size(); i++) {</pre>
   dai.push back({kra[i]. i}):
    ost[kra[i].first]=-1;
   ost[kra[i].second]=-1:
 counter=0:
 rek(0, (int)kra.size()-1, daj);
                                     Graphs/Vizing.cpp
#define int ll//jeśli long longi potrzebne
struct Vizina {
 vector<vector<pair<int,int>>> adj;
 map<pair<int,int>, int> edges;
 vector<int> edge colors;
 vector<vector<int>> color to:
 vector<vector<int>>> color queue;
 vector<int> unused color:
 int N, M, K;
 Vizing(int size) : adj(size), N(size), M(0) {}
  void AddEdge(int u, int v) {
```

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adj[u].emplace back(u, M);
  adi[v].emplace back(v, M);
  edges[make pair(min(u, v), max(u, v))] = M;
  ++M:
}
                                       //numeracia od zera
                                       //multikrawedzie niebezpieczne
int MaxDegree() const {
                                       //ogólnie nietestowane
 int answer = 0:
 for (int i = 0; i < N; ++i)
    answer=max(answer, (int)adj[i].size());
  return answer:
void FindFree(int v) {
  while (!color queue[v].emptv()) {
    const int c = color queue[v].back();
    if (color to[v][c] = -1) {
      unused color[v] = c;
      return;
    } else {
      color queue[v].pop back();
  assert(false);
void FlipPath(int from, int a, int b) {
  const int to = color to[from][a];
  color queue[from].push back(a);
  color queue[from] push back(b);
  color to[from][b] = -1;
  if (to == -1) { return; }
  FlipPath(to, b, a);
  color to[from][b] = to;
  color to[to][b] = from;
  color to[from][a] = -1;
  FindFree(from);
 FindFree(to);
void ColorEdge(int x, int y) {
  for (int col = 0; col <= K; ++col)</pre>
    if (color to[x][col] == -1 && color to[y][col] == -1) {
      color to[x][col] = y;
      color to[y][col] = x;
      FindFree(x); FindFree(y);
  vector<int> fan{y}, fan colors{-1};
  vector<bool> has on(N):
  const int c = unused_color[x];
  int d:
  while (true) {
    d = unused color[fan.back()];
    if (color \overline{to}[x][d] == -1 \mid \mid d == c) \{ break; \}
    if (has on[d]) {
      FlipPath(x, d, c);
      FindFree(x);
      return ColorEdge(x, v):
    fan.push back(color to[x][d]);
    fan colors.push bac\overline{k}(d);
    has on[d] = true;
  fan colors.push back(-1):
  for (int i = 0; i < (int)fan.size(); ++i) {</pre>
    const int from = fan colors[i], to = fan colors[i + 1];
    if (from != -1) {
```

```
color queue[fan[i]].push back(from);
        color queue[x].push back(from);
        color to[fan[i]][from] = -1;
     if (to != -1) {
        color to[fan[i]][to] = x;
        color to[x][to] = fan[i];
    assert(color to[x][d] == -1);
    color to[x][\overline{d}] = fan.back();
    color to[fan.back()][d] = x;
    for (int v : fan) { FindFree(v); }
   FindFree(x):
  vector<int> ColorGraph() {
    K = MaxDegree():
    edge colors = vector<int>(M, -1);
    color to = vector<vector<int>>>(N, vector<int>(K + 1, -1));
    vector<int> all colors(K + 1);
    iota(all_colors.begin(), all_colors.end(), 0);
    color queue = vector<vector<int>>>(N, all colors);
    unused color = vector<int>(N);
    for (auto &edge : edges)
     ColorEdge(edge.first.first, edge.first.second);
    for (int v = 0; v < N; ++v)
     for (int c = 0; c \le K; ++c)
        if (color to[v][c] > v) {
          assert(color to[color to[v][c]][c] == v);
          edge colors[edges[make pair(v, color_to[v][c])]] = c;
    return edge colors;
 }
#undef int
                                    Math/BitoweFFT.cpp
void xor fft(ll* p, int n, bool inv) {
 for (int len = 1; 2 * len <= n; len <<= 1) {
    for (int i = 0; i < n; i += 2 * len) {
     for (int j = 0; j < len; ++j) {
       ll u = p[i + i]:
       ll v = p[i + len + j];
        p[i + j] = u + v;
        p[i + len + j] = u - v;
 } } }
 if (inv) {
   for (int i = 0; i < n; ++i) {
     p[i] /= n; // uwaga gdy liczymy modulo!!
void and fft(ll* p, int n, bool inv) {
 for (int len = 1; 2 * len <= n; len <<= 1) {
    for (int i = 0; i < n; i += 2 * len) {
     for (int j = 0; j < len; ++j) {
       ll u = p[i + i]:
       ll v = p[i + len + j];
        if (!inv) {
          p[i + j] = v;
          p[i + len + j] = u + v;
        } else {
          p[i + j] = -u + v;
          p[i + len + j] = u;
void or fft(ll* p, int n, bool inv) {
 for (int len = 1; 2 * len <= n; len <<= 1) {
    for (int i = 0; i < n; i += 2 * len) {
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for (int j = 0; j < len; ++j) {
       ll u = p[i + j];
       ll v = p[i + len + j];
       if (!inv) {
         p[i + j] = u + v;
         p[i + len + j] = u;
       } else {
         p[i + j] = v;
         p[i + len + j] = u - v;
Math/FFT.cpp
#define REP(i,n) for(int i = 0; i < int(n); ++i)
/*Precision error max ans/1e15 (2.5e18) for (long) doubles.
So integer rounding works for doubles with answers 0.5*1e15,
e.g. for sizes 2^20 and RANDOM positive integers up to 45k.
Those values assume DBL MANT DIG=53 and LDBL MANT DIG=64.
For input in [0, M], you can decrease everything by M/2.
If there are many small vectors, uncomment "BRUTE FORCE".*/
typedef double ld: // 'long double' is 2.2 times slower
struct C {
 ld real, imag;
 C operator * (const C & he) const {
   return C{real * he.real - imag * he.imag, real * he.imag + imag * he.real};
 void operator += (const C & he) {real += he.real; imag += he.imag;} };
/oid dft(vector<C> & a. bool rev) {
 const int n = a.size():
 for(int i = 1, k = 0; i < n; ++i) {
   for(int bit = n / 2; (k ^= bit) < bit; bit /= 2);;;</pre>
   if(i < k) swap(a[i], a[k]):
 for(int len = 1, who = 0; len < n; len *= 2, ++who) {
   static vector<C> t[30]:
   vector<C> & om = t[who]:
   if(om.empty()) {
     om.resize(len):
     const ld ang = 2 * acosl(0) / len;
     REP(i, len) om[i] = i\%2 || !who?C{cos(i*ang), sin(i*ang)}:t[who-1][i/2];
   for(int i = 0; i < n; i += 2 * len)
     REP(k. len) {
        const C x = a[i+k], y = a[i+k+len]
           * C{om[k] real, om[k] imag * (rev ? -1 : 1)};
       a[i+k+len] = C{x.real - y.real, x.imag - y.imag};
 if(rev) REP(i, n) a[i].real /= n;
template<typename T>vector<T> multiply(const vector<T> &a, const vector<T> &b,
   bool split = false) {
 if(a.empty() || b.empty()) return {};
 int n = a.size() + b.size();
 vector<T> ans(n - 1):
 /* if(min(a.size(),b.size()) < 190) { // BRUTE FORCE
   REP(i. a.size()) REP(i. b.size()) ans[i+i] += a[i]*b[i]:
   return ans; } */
 while(n\&(n-1)) ++n;
 auto speed = [\&] (const vector<C> \& w, int i, int k) {
   int j = i ? n - i : 0, r = k ? -1 : 1;
   return C{w[i].real + w[j].real * r, w[i].imag
        - w[j].imag * r * (k ? C{0, -0.5} : C{0.5, 0});
 if(!split) { // standard fast version
   vector<C> in(n), done(n);
```

```
REP(i, a.size()) in[i].real = a[i];
    REP(i, b.size()) in[i].imag = b[i];
    dft(in, false);
    REP(i, n) done[i] = speed(in, i, 0) * speed(in, i, 1);
    dft(done, true);
    REP(i, ans.size()) ans[i] = is integral<T>::value ?
        llround(done[i].real) : done[i].real;
 else {
    const int M = 1 \ll 15:
    vector <C> t[2];
    for (int x = 0; x < 2; ++x) {
     t[x].resize(n);
      const vector \langle T \rangle \& in = (x ? b : a):
      for (int i = 0; i < (int) in.size(); ++i)</pre>
        t[x][i] = C\{ld(in[i] % M), ld(in[i] / M)\};
      dft(t[x]. false):
    vector \langle C \rangle d1(n), d2(n);
    for (int i = 0; i < n; ++i) {
      d1[i] += speed(t[0], i, 0) * speed(t[1], i, 0);
      d1[i] += speed(t[0], i, 1) * speed(t[1], i, 1) * C{0, 1};
      d2[i] += speed(t[0], i, 0) * speed(t[1], i, 1);
      d2[i] += speed(t[0], i, 1) * speed(t[1], i, 0);
    dft(d1, true);
    dft(d2, true):
    for (int i = 0; i < n; ++i)
      d1[i].imag /= n;
    for (int i = 0; i < (int) ans.size(); ++i)
      ans[i] = (llround(d1[i].real) + llround(d2[i].real) % mod * M +
        llround(d1[i].imag) % \mod * (M * M)) % \mod:
  return ans;
                            Math/Linear Function Convex Hull.cpp
const ll is query = -(1LL << 62):</pre>
struct Line {
 ll m. b:
  mutable function<const Line *()> succ;
  bool operator<(const Line &rhs) const {
    if (rhs.b != is query) return m < rhs.m;</pre>
    const Line *s = succ();
    if (!s) return 0;
    ll x = rhs.m:
    return b - s->b < (s->m - m) * x; //ld?
struct HullDvnamic : public multiset<Line> {
 bool bad(iterator y) {
    auto z = next(v):
    if (y == begin()) {
      if (z == end()) return 0;
      return v->m == z->m \&\& v->b <= z->b:
    auto x = prev(v):
    if (z == end()) return y->m == x->m \&\& y->b <= x->b;
    return (x-b-y-b) * (z-m-y-m) >= (y-b-z-b) * (y-m-x-m);//ld?
  void insert line(ll m, ll b) {
    auto y = insert({m, b});
    y->succ = [=] { return next(y) == end() ? 0 : \&*next(y); };
    if (bad(y)) {
      erase(y);
      return:
```

```
while (next(y) != end() \&\& bad(next(y))) erase(next(y));
   while (y \mid = begin() \&\& bad(prev(y))) erase(prev(y));
 ll eval(ll x) {
   auto l = *lower bound((Line) {x, is query});
   return l.m * x + l.b;
                                      Math/Massey.cpp
const int mod = 1e9 + 7; // const ll mod = (ll)1e18 + 3;
template <class c> void add self(c \& a, c b) \{ a += b; if(a >= mod) a -= mod; \}
template <class c> void sub self(c & a, c b) { a -= b; if(a < 0) a += mod; }
int mul(int a, int b) { return (ll) a * b % mod; }
ll mul(ll a, ll b) {return ( int128) a * b % mod; }
template <class c> c my pow(\overline{c} a, c b) \{/*...*/\}
template <class c> c my inv(c a) { return my pow<c>(a, mod - 2); }
template <class c> c negative (c a) {return (mod - a) % mod;}
template <class F> struct Massev {
 vector<F> start, coef; // 3 optional lines
 vector<vector<F>>> powers;
 F memo inv;
 // Start here and write the next ~25 lines until "STOP"
 int L; // L == coef.size() <= start.size()</pre>
 Massey(vector<F> in) { // O(N^2)
   const int N = in.size();
   vector<F> C{1}, B{1};
   for(int n = 0; n < N; ++n) {
     // assert(0 <= in[n] && in[n] < mod); // invalid input
     B.insert(B.begin(), 0);
     F d = 0:
     for(int i = 0; i \le L; ++i)
        add self(d, mul(C[i], in[n-i]));
     if(d == 0) continue;
     vector<F> T = C:
     C.resize(max(B.size(), C.size()));
     for(int i = 0; i < (int) B.size(); ++i)
        sub self(C[i], mul(d, B[i]));
     if(2 \overline{*} L <= n) {
       L = n + 1 - L:
       B = T;
        d = my inv(d);
        for (F \overline{\&} x : B) x = mul(x, d);
   }
   cerr << "L = " << L << "\n":
   assert(2 * L <= N - 2); // NO RELATION FOUND :(
   // === STOP ===
   for(int i = 1; i < (int) C.size(); ++i)</pre>
     coef.push back(negative(C[i]));
   assert((int) coef.size() == L);
   for(int i = 0; i < L; ++i)
     start.push back(in[i]):
   while(!coef.empty() && coef.back() == 0) { coef.pop back(); --L; }
   if(!coef.emptv()) memo inv = mv inv(coef.back());
   powers.push back(coef);
 vector<F> mul cut(vector<F> a, vector<F> b) {
   vector<F> r(2 * L - 1);
   for(int i = 0; i < L; ++i)
     for(int i = 0: i < L: ++i)
        add_self(r[i+j], mul(a[i], b[j]));
   while((\overline{int}) r.size() > L) {
     F value = mul(r.back(), memo inv); // div(r.back(), coef.back());
```

```
const int X = r.size():
      add self(r[X-L-1], value);
      for(int i = 0; i < L; ++i)
       sub self(r[X-L+i], mul(value, coef[i]));
      assert(r.back() == 0);
      r.pop back();
    return r;
  F get(ll k) { // O(L^2 * log(k))
    if(k < (int) start.size()) return start[k];</pre>
    if(L == 0) return 0:
    k -= start.size();
    vector<F> vec = coef:
    for(int i = 0; (1LL << i) <= k; ++i) {
      if(i == (int) powers.size())
        powers.push back(mul_cut(powers.back(), powers.back()));
      if(k & (1LL << i))
       vec = mul cut(vec. powers[i]):
    F total = 0;
    for(int i = 0; i < L; ++i)
      add self(total, mul(vec[i], start[(int)start.size()-1-i]));
    return total; } };
                              Math/Multipoint Evaluation.cpp
ll dziel(ll a, ll b) {...}
vll inv(vll p. int n) {
  assert(p[0]);
 ll odw=dziel(1, p[0]);
  for (ll &i : p) i=(i*odw)%mod;
  vll q{1};
  for (int i=1: i<n: i<<=1) {</pre>
   vll r=multiply(vll(p.begin(), p.begin()+min(2*i, (int)p.size())), q, true);
    r.resize(2*i):
    r.erase(r.begin(), r.begin()+i);
    for (ll &j : r) j=(mod-j)%mod;
    r=multiply(r, q, true);
    a.resize(2*i):
    for (int j=0; j<min(i, (int)r.size()); j++) q[i+j]=r[j];</pre>
 a.resize(n):
  while(!q.empty() && !q.back()) q.pop back();
  for (ll &i : p) i=(i*odw)%mod;
  return q;
vll div(vll a, vll b) {
  int s=(int)a.size()-(int)b.size()+1;
  if (s<=0) return {};
  reverse(a.begin(), a.end()); reverse(b.begin(), b.end());
  b=inv(b, s);
  a=multiply(a, b, true);
  a.resize(s);
  reverse(a.begin(), a.end());
  return a:
vll rem(vll a, vll b) {
 vll d=div(a, b);
  d=multiply(d, b, true);
  for (int i=0; i<(int)d.size(); i++) a[i]=(a[i]-d[i]+mod)%mod;</pre>
  while(!a.empty() \&\& !a.back()) a.pop back();
 return a:
void precalc(int v, vector <vll> &help, vll query) {
 if ((int)help.size()<=v) help.resize(v+1);</pre>
  int s=query.size();
```

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if (s==1)
   help[v]={(mod-query[0])%mod, 1};
   return:
 vll a, b;
 for (int i=0; i<s/2; i++) a.push back(query[i]);</pre>
 for (int i=s/2; i<s; i++) b.push back(query[i]);</pre>
 precalc(v^*2, help, a); precalc(v^*2+1, help, b);
 help[v]=multiply(help[v*2], help[v*2+1], true);
/oid calc(int v, vll &res, vector <vll> &help, vll wek) {
 wek=rem(wek. help[v]):
 if ((int)help[v].size()==2) {
   res.push back(wek.emptv() ? 0 : wek[0]):
   return;
 calc(v*2, res. help, wek):
 calc(v*2+1, res, help, wek);
/ll multi eva(vll wek, vll guery) {
 vector <vll> help;
 precalc(1, help, query);
 vll res:
 calc(1, res, help, wek);
 return res:
                                     Math/Simplex.cpp
struct Simplex {
                       // Maximize c*x subject to Ax <= b.
 using T = double;
                      // Initialize the structure, set A, b, c and then run
 vector<vector<T>> A: // solve(), Max objective is stored in res, To recover
 vector<T> b, c;
                       // the best result, use getVars().
 int V. E:
 vector<int> eqIds, varIds, cols;
 T res:
 static constexpr T kEps = 1e-9;
 Simplex(int vars, int eqs) : A(eqs, vector<T>(vars)), b(eqs), c(vars),
     V(vars), E(eqs), eqIds(eqs), varIds(vars), res(0) {
   iota(varIds.begin(), varIds.end(), 0);
   iota(eqIds.begin(), eqIds.end(), vars);
 void pivot(int eq. int var) {
   T coef = 1 / A[eq][var];
   cols.clear();
   for (int i = 0; i < V; i++) {
     if (abs(A[eq][i]) > kEps) { cols.push back(i); A[eq][i] *= coef; }
   A[eq][var] *= coef: b[eq] *= coef:
   for (int row = 0; row < E; row++) {
     if (row == eq || abs(A[row][var]) < kEps) { continue; }</pre>
     T k = -A[row][var];
     A[row][var] = 0:
     for (int i : cols) { A[row][i] += k * A[eq][i]; }
     b[row] += k * b[eq];
   T q = c[var]; c[var] = 0;
   for (int i : cols) { c[i] -= q * A[eq][i]; }
   res += q * b[eq];
   swap(varIds[var], eqIds[eq]);
 bool solve() {
   while (true) {
     int eq = -1, var = -1:
     for (int i = 0; i < E; i++) { if (b[i] < -kEps) { eq = i; break; } }
     if (eq == -1) { break; }
     for (int i = 0; i < V; i++) { if (A[eq][i] < -kEps) { var = i; break; } }
```

```
if (var == -1) { res = -le9; return false; /* No solution */ }
      pivot(eq, var);
    while (true) {
      int var = -1, eq = -1;
      for (int i = 0: i < V: i++) { if (c[i] > kEps) { var = i: break: } }
      if (var == -1) { break; }
      for (int i = 0: i < E: i++) {
        if (A[i][var] < kEps) { continue; }</pre>
        if (eq \ge 0 \& \& b[i] / A[i][var] \ge b[eq] / A[eq][var]) { continue; }
      if (eg == -1) { res = 1e9; return false; /* Unbounded */ }
      pivot(eq. var):
    return true;
  vector<T> getVars() { // Optimal assignment of variables.
   vector<T> result(V):
    for (int i = 0; i < E; i++) if (eqIds[i] < V) result[eqIds[i]] = b[i];
    return result:
 }
};
                                          Math/Simpson.cpp
V/ Either run integral(A, B) once or split the interval [A, B] into up to ~1000
// smaller intervals -- if the function f behaves oddly or the interval is long.
ld simp(ld low, ld high, const ld * old, vector<ld> & nowe) {
  const int n = 500; // n must be even!!! Try n = 2 and n = 10.
  nowe.resize(n + 1);
  ld\ total = 0, jump = (high - low) / n;
  for(int i = 0; i \le n; ++i) {
    int mul = i == 0 || i == n ? 1 : 2 + i % 2 * 2: // 1 4 2 4 2 ... 4 1
    nowe[i] = !old | | i % 2 ? f(low + i * jump) : old[i/2];
    total += nowe[i] * mul; // uses a global function ld f(ld x)
 return total * (high - low) / n / 3;
ld rec(ld low, ld high, ld prv, const vector<ld> <math>\& old) {
 ld mid = (low + high) / 2:
 vector<ld> left, right;
 ld L = simp(low, mid, old.data(), left);
 ld R = simp(mid, high, old.data() + old.size() / 2, right);
 ld del = L + R - prv;
                                                         // try without del/15
 if(abs(del) < 15 * eps) return L + R + del / 15; // eps ~ required abs precision</pre>
  return rec(low, mid, L, left) + rec(mid, high, R, right);
ld integral(ld low, ld high) {
 vector<ld> old;
 ld prv = simp(low, high, 0, old);
  return rec(low, high, prv, old);
                                           Math/Wzory.tex
 \int \sqrt{x^2 + 1} \, dx = \frac{1}{2} \left( x \sqrt{x^2 + 1} + \operatorname{arcsinh} x \right) + c
 \int \sqrt{1-x^2} \, dx = \frac{1}{2} \left( x \sqrt{1-x^2} + \arcsin x \right) + c
 \int \frac{1}{ax^2 + bx + c} dx = \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}}
                                                               (\Delta < 0)
 \int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln |ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}
 \int \tan x \, dx = -\ln|\cos x| + c
(\arcsin x)' = \frac{1}{\sqrt{1 - x^2}}, \qquad (\arccos x)' = -\frac{1}{\sqrt{1 - x^2}}
                                          Math/Wzory-2.tex
```

G

 $tgamma(t) = \Gamma(t) = \int_0^\infty e^{t-1}e^{-x} dx$  $\frac{1}{\pi} = 0.31831, \quad \pi^2 = 9.86960, \quad \frac{1}{\pi^2} = 0.10132, \quad \frac{1}{e} = 0.36788, \quad \gamma = 0.10132, \quad \frac{1}{e} = 0.10132, \quad \frac{1}{e} = 0.36788, \quad \gamma = 0.10132, \quad \frac{1}{e} = 0.36788, \quad \gamma = 0.10132, \quad \frac{1}{e} = 0.10132, \quad \frac{1}{e} = 0.36788, \quad \gamma = 0.10132, \quad \frac{1}{e} = 0.10132, \quad$ 0.577215664901532

 $H_n = \ln n + \gamma + \frac{1}{2n} - \frac{1}{12n^2} + O(n^{-4})$   $\ln n! = n \ln n - n + \frac{1}{2} \ln(2\pi n) + \frac{1}{12n} - \frac{1}{360n^3} + \frac{1}{1260n^5} - O(n^{-7})$ 

Jeśli  $n = 2^{a_0} \cdot p_1^{2 \cdot a_1} \cdot \dots \cdot p_r^{2 \cdot a_r} \cdot q_1^{b_1} \cdot \dots \cdot q_s^{b_s}$  i  $B = \prod (b_i + 1)$ , gdzie  $p_i$ to liczby postaci  $4 \cdot k + 3$ , a  $q_i$  to liczby postaci  $4 \cdot k + 1$ , to liczba sposobów na zapisanie n jako sumy dwóch kwadratów liczb naturalnych wynosi:

- 0. ieśli któraś z liczb  $a_i$  nie jest całkowita
- $\frac{B}{2}$ , jeśli B jest parzyste
- $\frac{B-(-1)^{a_0}}{2}$ , jeśli B jest nieparzyste

$$S_j = \sum_{1 \leqslant i_1 < i_2 < \dots < i_j \leqslant n} | A_{i_1} \cap A_{i_2} \cap \dots A_{i_j} |$$

 $\sum\limits_{j\geqslant k}{j\choose k}(-1)^{j+k}S_j$  - liczba elementów należących do dokładnie kzbiorów

 $\sum\limits_{j\geqslant k}{j-1\choose k-1}(-1)^{j+k}S_j$  - liczba elementów należących do co najmniej kzbiorów.

 $\sum\limits_{j\geqslant 1} (-2)^{j-1} S_j$  - liczba elementów należących do nieparzyście wiele zbiorów

Zaczynajac w punkcie 0, idac w każdym ruchu o -1 z prawdopodobieństwem p i o +1 z 1-p, kończac w -A lub B prawdopodobieństwo dojścia do B wynosi  $\frac{1+r+r^2+\ldots+r^{A-1}}{1+r+r^2+\ldots+r^{A+B-1}} = \frac{1-r^A}{1-r^{A+B}}^* = 1 - \frac{1+r'+r'^2+\ldots+r'^{B-1}}{1+r'+r'^2+\ldots+r'^{A+B-1}} = 1 - \frac{1-r'^B}{1-r'^{A+B}}^*,$ (\*działa o ile  $p \neq \frac{1}{2}$ ) gdzie  $r = \frac{p}{1-p}, r' = \frac{1}{r}$ , dla  $p = \frac{1}{2}$  równe  $\frac{A}{A+B}$ 

Dla  $x_1, x_2, \ldots, x_n$  macierz  $a_{i,j} = x_i^{j-1}$  ma wyznacznik  $\prod (x_i - x_i)$  i odwrotność

$$b_{i,j} = (-1)^{i \frac{\sum\limits_{A \subset [n] \setminus \{j\}, |A| = n - i} \prod\limits_{k \in A} x_k}{\prod\limits_{k \neq i} (x_k - x_j)}}$$

Miscellaneous/De Bruijn.cpp

```
//Generuje słowo cykliczne, Działa dla alph >= 2
vi de bruijn(int len, int alph){
 vi res, lyn{0};
 while (lyn[0] != alph - 1){
   int r = siz(lyn);
   if (len % r == 0)
     for (int c : lyn)
       res.pb(c);
   for (int i = r; i <= len - 1; ++i)
     lyn.pb(lyn[i - r]);
   while (lyn.back() == alph - 1)
     lvn.pop back():
   lyn.back()++;
 res.pb(alph - 1);
```

```
return res;
                                 Miscellaneous/Loops.cpp
sim> bool next inc(c beg, c end, int bou) { //[0, bou)
  for (c i = end; i != beq;) {
    int v = *--i:
   if (v + 1 < bou) {
      iota(i, end, v + 1);
      return true;
   bou = *i;
  return false;
sim> bool next nondec(c beg, c end, int bou) {//[0, bou)
  for (c i = end; i != beg;) {
   int v = *--i:
   if (v + 1 < bou) {
     fill(i. end. v + 1):
      return true;
  return false;
sim> bool next any(c beg, c end, int bou) {
 if (!bou) return false:
  for (c i = end: i != bea:) {
   if (++*--i == bou) *i = 0;
    else return true;
  return false;
sim>bool next diff(c beg,c end,int bou){//starts with 0,1,2...
  return next permutation(beg, end) || next inc(beg, end, bou);
sim>bool next split(c beg, c end) {//starts with all 0s, ends with 0,1,2...
 vector <int> seen(distance(beg, end));
  for (c i = beg; i != end; ++i) seen[*i]++;
  for (c i = end; i != beg;) {
   if (seen[*--i] > 1) {
      (*i)++:
      fill(++i, end, 0);
      return true:
  return false;
V/for (int m = (1 << k) - 1; m < (1 << n); m = nux(m)) - k-podzbiory [0,n)
sim>inline c nux(c m){
 if (!m) return numeric limits<c>::max();
  c A = m \& -m;
  c B = \sim ((A - 1) ^ m);
  c C = B \& -B;
  c D = (C >> (1 + ctz(A))) - 1;
  return C | (m & ~(C - 1)) | D;
V/for (unt i = 0; i < (1u << n); i = rrh(i)) maska bez sasiednich zapalonych, rosnaco
sim> inline c rrh(c m) {
 c a = \sim (m \mid m >> 1);
  cb = a \& -a:
 return (m | b) &~ (b - 1);
V/for (unt i = (1u << (n + 1)) / 3; ~i; i = lyl(i)) maska bez sąsiednich, malejąco
sim> inline c lyl(c m) {
 if (!m) return -1;
```

```
c x = m \& -m:
 return (m ^ x) | ((x << 1) / 3);
for (ll val = 1, low = n / 2, high = n; ; val = n / low, high = low, low = n / (val + 1).
 //include n / i = val for i in (low, high)
 if (val == n) break;
                               Number Theory/Cornacchia.cpp
usina T = int:
                       // Typ, w którym mieści się modulo.
using T2 = long long; // Typ, w którym mieści sie kwadrat modulo.
struct Pierwiastek {// Dla pierwszego p > 2, szuka pierwiastków modulo p.
 Pierwiastek(T p) : s(\theta), p(p), m(p-1) {//template <int p> jeśli stałe p
   assert(p > 2):
   while (m \% 2 == 0) \{ m /= 2; s++; \}
   c = 2;
   while (Pot(c, p / 2) == 1) c = rand() % (p - 1) + 1; //rand() musi być duży
 T Pot(T a, T pot) const {
   T r = 1;
   while (pot) {
     if (pot & 1) r = r * (T2) a % p;
     a = a * (T2) a % p;
     pot >>= 1;
   return r;
 } // Zwraca a**pot % p.
 T Licz(T a) const { // Znajduje pierwiastek z a modulo p.
   if (a == 0) return 0;
                                      // Sprawdza, czy a jest resztą
   if (Pot(a, p / 2) != 1) return -1; // kwadratowa. Jeśli nie, zwraca -1.
   T z = Pot(c, m);
   T v = Pot(a. m / 2):
   T u = (T2) v * v % p;
   v = (T2) v * a % p:
   u = (T2) u * a % p;
   for (int i = s - 1; i >= 1; i --) {
     if (Pot(u, T(1) << (i - 1)) != 1) {
       u = (T2) u * z % p:
       u = (T2) u * z % p;
       v = (T2) v * z % p;
     z = (T2) z * z % p;
   return v; // Pierwiastkami liczby a są: {v, (p - v) % p}.
private: int s;
   T p. m. c:
template<class num t>
inline num t isgrt(num t k) {
 num t r = sqrt(k) + 1;
 while (r * r > k) r--;
 return r:
pii cornacchia(int p, int d) {//dla p pierwszego, znajduje a, b: p = a^2 + db^2
 int x0 = Pierwiastek(p).Licz(p - d):
 if (x0 == -1) return \{-1, -1\};
 x\theta = max(x\theta, p - x\theta);
 int a = p, b = x0;
 int l = isart(p):
 while (b > l) {
   int r = a \% b:
   a = b;
   b = r;
```

```
int c = p - b * b;
  if (c % d) return {-1, -1};
  c /= d;
  int cc = isart(c):
  if (cc * cc != c) return {-1, -1};
  return {b, cc}:
                                  Number Theory/CRT.cpp
void eukl(ll &x, ll &y, ll a, ll b) {
  if (!a) { x = 0; y = 1; return; }
  eukl(y, x, b % a, a);
 x -= y * (b / a);
ll mno(ll a, ll b, ll mod) {/*,..*/}
pll crt2(ll p, ll a, ll g, ll b) { //x==p \pmod{a}, x==q \pmod{b}
 if (a==-1) return {-1, -1};// Zarówno wynik, jak i argumenty są
  ll x, y;
                          // postaci x = first (mod second).
  eukl(x, y, a, b);
                           // Jeśli kongruencja jest niespełnialna
  ll nwd=x*a+y*b;
                           // to zwraca (-1, -1).
  if ((p%nwd)!=(g%nwd))
    return {-1, -1};
  a/=nwd; b/=nwd;
  ll nww=a*b:
  ll ret=mno(x*a, q/nwd, nww)+mno(y*b, p/nwd, nww);
  if ((ret%=nww)<0) ret+=nww:</pre>
  return {ret*nwd+(p%nwd), nww*nwd};
                                Number Theory/Euklides.cpp
ll trzesienie(ll a, ll b, ll c) {
 if (c<0) return 0:
 if (a>b) swap(a, b);
 ll p=c/b:
 ll k=b/a:
 ll d=(c-p*b)/a;
  return trzesienie(b-k*a, a, c-a*(k*p+d+1))+(p+1)*(d+1)+k*p*(p+1)/2;
\} //counts pairs of nonnegative integers (x, y) such that ax + by <= c
                            Number Theory/Fibonacci Cycle.txt
Znaleziony cykl niekoniecznie jest najmniejszym cyklem, wiadomo jednak, że cykl(m)<=6m
Dla m, które nie są postaci 2 * 5^r, zachodzi ograniczenie: cykl(m) <= 4m.
Jeśli nie zrobi się Nww, to wiadomo, ze cykl(m) <= m * 2**(1 + #dz.pierw.m).</pre>
|cvcle(2)| = 3, cvcle(5)| = 20, cvcle(6)| = 1 or 9 \mod 10, |cvcle(6)| = 1
cycle(p == 3 \text{ or } 7 \text{ mod } 10) = 2(p+1) cycle(p^k) = p^{k} - 1 * cycle(k)
cycle(product p i^a i) = lcm(cycle(p i^a^i))
                           Number Theory/number theory data.txt
2: 1 prime divisors, largest=2
6: 2 prime divisors, largest=3
30: 3 prime divisors, largest=5
210: 4 prime divisors, largest=7
2.310: 5 prime divisors, largest=11
30.030: 6 prime divisors, largest=13
510.510: 7 prime divisors, largest=17
9.699.690: 8 prime divisors, largest=19
223.092.870: 9 prime divisors, largest=23
6.469.693.230: 10 prime divisors, largest=29
200.560.490.130: 11 prime divisors, largest=31
7.420.738.134.810: 12 prime divisors. largest=37
304.250.263.527.210: 13 prime divisors, largest=41
13.082.761.331.670.030: 14 prime divisors, largest=43
614.889.782.588.491.410: 15 prime divisors, largest=47
pi(10)=4
                     pi(20)=8
                                            pi(40)=12
                                                                    pi(75)=21
pi(100)=25
                     pi(200)=46
                                            pi(400)=78
                                                                    pi(750)=132
pi(1000)=168
                     pi(2000)=303
                                            pi(4000)=550
                                                                    pi(7.500)=950
pi(10000)=1.229
                     pi(20000)=2.262
                                            pi(40000)=4.203
                                                                    pi(75000)=7.393
pi(100000)=9.592
                                                                    pi(750000)=60.238
                     pi(200000)=17.984
                                            pi(400000)=33.860
pi(1000000)=78.498
                                                                    pi(7.500000)=508.261
                     pi(2000000)=148.933
                                            pi(4000000)=283.146
```

F

G

G

M

```
pi(10000000)=664.579 pi(20000000)=1.270.607 pi(40000000)=2.433.654
pi(75000000)=4.394.304
pi(100000000)=5.761.455
                                            pi(200000000)=11.078.937
pi(400000000)=21.336.326
                                            pi(750000000)=38.703.181
pi(10000000000)=50.847.534
                                            pi(2000000000)=98.222.287
pi(40000000000)=189.961.812
                                            pi(7.500000000)=345.826.612
pi(10000000000)=455.052.511
                                            pi(20000000000)=882.206.716
pi(40000000000)=1.711.955.433
                                            pi(75000000000)=3.125.641.477
pi(1000000000000)=4.118.054.813
                                            pi(200000000000)=8.007.105.059
pi(4000000000000)=15.581.005.657
                                            pi(750000000000)=28.514.968.374
Largest possible number of divisors leg 10^k for k = 3..18
840=2^3*3*5*7
                                         32 divisors
7.560=2^3*3^3*5*7
                                         64 divisors
720.720=2^4*3^2*5*7*11*13
                                         240 divisors
8.648.640=2^6*3^3*5*7*11*13
                                         448 divisors
73.513.440=2^5*3^3*5*7*11*13*17
                                         768 divisors
735.134.400=2^6*3^3*5^2*7*11*13*17
                                         1.344 divisors
                                                             2.304 divisors
6.983.776.800=2^5*3^3*5^2*7*11*13*17*19
                                                             4.032 divisors
97.772.875.200=2^6*3^3*5^2*7^2*11*13*17*19
963.761.198.400=2^6*3^4*5^2*7*11*13*17*19*23
                                                             6.720 divisors
9.316.358.251.200=2^6*3^3*5^2*7*11*13*17*19*23*29
                                                              10.752 divisors
97.821.761.637.600=2^5*3^4*5^2*7^2*11*13*17*19*23*29
                                                              17.280 divisors
866.421.317.361.600=2^6*3^4*5^2*7*11*13*17*19*23*29*31
                                                              26.880 divisors
8.086.598.962.041.600=2^8*3^3*5^2*7^2*11*13*17*19*23*29*31
                                                                        41.472 divisors
74.801.040.398.884.800=2^6*3^3*5^2*7^2*11*13*17*19*23*29*31*37
                                                                        64.512 divisors
897.612.484.786.617.600=2^8*3^4*5^2*7^2*11*13*17*19*23*29*31*37
                                                                        103.680 divisors
                                  Number Theory/Phi.cpp
struct Coprimes {
 vector <ll> w, dp;
 int gdz(ll v) {
   if (v \le w.back()/v)
     return v-1:
   return w.size()-w.back()/v;
 ll phi(ll n) \{ //sum [k = 1 to n] phi(k) \}
   for (ll i=1; i*i<=n; i++) {
     w.push back(i);
     if ((n/i)!=i)
        w.push back(n/i);
   sort(w.begin(), w.end());
   for (ll i : w) {
     dp.push back(i*(i+1)/2);
     for (ll j=1; j*j<=i; j++) {
       if (i>1) dp.back()-=dp[gdz(j)]*(i/j-i/(j+1));
       if (j>1 \&\& j<=i/(j+1)) dp.back()-=dp[gdz(i/j)];
   return dp.back();
 ll ask(ll v) { return dp[gdz(v)]; }//v==n/u for some u
: //phi(10 ** 10) = 2726238982321888942
                                   Number Theory/Pi.cpp
struct Primes {//Pi(10^10) = 455052511
 vector <ll> w, dp;
 int adz(ll v) {
   if (v <= w.back() / v) return v - 1;</pre>
   return w.size() - w.back() / v;
 ll pi(ll n) {
   for (ll i = 1; i * i <= n; i++) {
     w.pb(i):
     if (n / i != i) w.push back(n/i);
   sort(all(w));
```

```
for (ll i : w) dp.push back(i-1);
    for (ll i = 1; (i + 1)^{-*} (i + 1) <= n; i++) {
      if (dp[i] == dp[i - 1]) //Might not work when calculating modulo or sum of
        continue; //some weird function over primes, meant to skip if i + 1 is not prime
      for (int j = siz(w) - 1; w[j] >= (i + 1) * (i + 1); j--)
        dp[i] = dp[adz(w[i] / (i + 1))] - dp[i - 1];
    return dp.back();
  ll ask(ll v) { return dp[gdz(v)]; }//Działa tylko jak v==n/u for some u
1: //Sum upto(10^10) mod 1e9+7 = 35972319
                                   Number Theory/Rho.cpp
// szybkie, faktoryzuje liczby z przedzialu [1e18, 1e18 + 1e5] w 9s na cfie
namespace rho {
const int maxv = 50;
const int maxp = 1e6 + 7; // preprocesujemy rozklad do maxp - 1
int ptot:
int d[maxp]:
int pr[maxp]:
vector<ll> ans; // rozklad, czyscimy zazwyczaj po kazdym wywolaniu decompose()
inline ll mod(ll a, ll n) {
 if (a >= n) a -= n;
  return a:
inline ll add(ll a, ll b, ll n) {
  a += b:
  mod(a, n);
  return a:
inline ll mul(ll x, ll y, ll p) { // uwaga, ta funkcja rzuca overflow, tak ma byc
 ll ret = x * y - (ll)((ld)x * y / p + 0.5) * p;
  return ret < 0 ? ret + p : ret;</pre>
ll fast(ll x, ll k, ll p) {
 ll ret = 1 \% p;
  for (; k > 0; k >>= 1, x = mul(x, x, p)) (k \& 1) \&\& (ret = mul(ret, x, p));
  return ret;
bool rabin(ll n) {
 if (n == 2) return 1;
  if (n < 2 \mid | !(n \& 1)) return 0:
 ll s = 0, r = n - 1;
  for (; !(r \& 1); r >>= 1, ++s)
  for (int i = 0; pr[i] < n && pr[i] < maxv; ++i) {
    ll cur = fast(pr[i], r, n), nxt;
    for (int j = 0; j < s; ++j) {
      nxt = mul(cur, cur, n);
      if (nxt == 1 \&\& cur != 1 \&\& cur != n - 1) return false:
      cur = nxt;
    if (cur != 1) return false;
  return true:
ll factor(ll n) {
  static ll seg[maxp];
  while (true) {
    ll x = rand() % n, y = x, c = rand() % n;
    ll *px = seq, *py = seq, tim = 0, prd = 1;
    while (true) {
      *py++ = y = add(mul(y, y, n), c, n);
      *py++ = y = add(mul(y, y, n), c, n);
      if ((x = *px++) == y) break;
      ll tmp = prd;
```

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prd = mul(prd, abs(y - x), n);
     if (!prd) return gcd(tmp, n);
     if ((++tim) == maxv) {
       if ((prd = gcd(prd, n)) > 1 \&\& prd < n) return prd;
       tim = 0;
   if (tim \&\& (prd = gcd(prd, n)) > 1 \&\& prd < n) return prd;
void decompose(ll n) { // glowna funkcja
 if (n < maxp) {
   while (n > 1) ans.pb(d[n]), n /= d[n];
 } else if (rabin(n))
   ans.pb(n);
 else {
   ll fact = factor(n):
   decompose(fact), decompose(n / fact);
void init() {  // wywolujemy przed pierwsza faktoryzacja, raz na caly program
 d[1] = 1;
 for (int i = 2; i * i < maxp; ++i)</pre>
   if (!d[i])
     for (int j = i * i; j < maxp; j += i) d[j] = i;
 for (int i = 2; i < maxp; ++i)
   if (!d[i]) {
     d[i] = i;
     pr[ptot++] = i;
  // namespace rho
                                     Strings/Aho.cpp
const int MAXN = 404, sigma = 26;
int term[MAXN], len[MAXN], to[MAXN][sigma], link[MAXN], sz = 1;
void add str(string s) {
 int cur = 0;
 for(auto c: s) {
   if(!to[cur][c - 'a']) {
     to[cur][c - 'a'] = sz++;
     len[to[cur][c - 'a']] = len[cur] + 1;
   cur = to[cur][c - 'a'];
 term[cur] = cur;
void push links() {
 int que[sz];
 int st = 0. fi = 1:
 que[0] = 0;
 while(st < fi)</pre>
   int V = que[st++];
   int U = link(V):
   if(!term[V]) term[V] = term[U];
   for(int c = 0: c < sigma: c++)
     if(to[V][c]) {
       link[to[V][c]] = V ? to[U][c] : 0;
       que[fi++] = to[V][c];
     } else {
       to[V][c] = to[U][c];
 }
                                   Strings/Eertree.cpp
```

```
const int maxn=1000*1000+7, alfa=26;
int len[maxn], link[maxn], to[maxn][alfa];
int sz, last, n;
char s[maxn]:
void init() {
 s[n++]=len[1]=-1:
 link[0]=1;
 sz=2:
int get link(int v) {
 while (s[n-len[v]-2]!=s[n-1]) v=link[v];
 return v:
void add letter(char c) {
  s[n++]=c-='a';
  last=get link(last);
  if(!to[last][c]) {
    len[sz]=len[last]+2;
    link[sz]=to[get link(link[last])][c];
    to[last][c]=sz++;
 last=to[last][c];
int main() {
 init():
 //\sim for(int i=1: i<=nn: i++)
   //~ add letter(tekst[i]);
                                  Strings/Graf Podslow.cpp
struct suffix automaton {
  vector<map<char,int>> edges;
  vi link. lenath:
  int last:
  suffix automaton(string s) {
    edges.pb(map<char,int>());
    link.pb(-1):
    length.pb(0);
    last = 0:
    for (int i=0; i<siz(s); i++) {</pre>
      edges.pb(map<char,int>());
      length.pb(i+1):
      link.pb(0);
      int r = siz(edges) - 1;
      int p = last;
      while (p \ge 0 \&\& !edges[p].count(s[i]))  {
        edges[p][s[i]] = r;
        p = link[p]:
      if (p != -1) {
        int q = edges[p][s[i]];
        if (length[p] + 1 == length[q]) {
          link[r] = q;
        else {
          edges.pb(edges[q]);
          length.pb(length[p] + 1):
          link.pb(link[q]);
          int qq = siz(edges)-1;
          link[q] = link[r] = qq;
          while (p \ge 0 \&\& edges[p][s[i]] == q) {
            edges[p][s[i]] = qq;
            p = link[p];
```

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last = r;
  }
 }
                                   Strings/KMP.cpp
int l = 0:
for (int i = 1; i < n; i++) {
 while (l > 0 \&\& s[l] != s[i]) l = P[l-1];
 l += (s[l] == s[i]);
 P[i] = l;
} //Słowo indeksowane od 0, tablica P też
                                  Strings/Lyndon.cpp
// 1) Przyjmuje słowo s (wypełnione na pozycjach 0, 1, ..., n-1).
     Dzieli słowo s na pewną liczbę słów Lyndona p_1, ... p_k tak, że:
     p \ 1 \ge p \ 2 \ge \dots \ge p \ k \ (leksykograficznie)
     Podział jest zapisywany w tablicy b - na i-tej pozycji jest true,
     jeśli nastąpiło cięcie przed i-tą literką.
// 2) Znajduje minimalne leksykograficznie przesunięcie cykliczne słowa.
// 3) Znajduje minimalny leksykograficznie sufiks słowa.
void lyndon(char * s,  // Słowo zaczynające sie na pozycji 0:
                      // 2) s powinno być sklejone: xx.
                      // Długość słowa s (licząc ew. podwojenie).
           int& suf, // 3) pozycja minimalnego leksykograficznie sufiksu.
           int& cyk, // 2) pozycja minimalnego leksykograficznie przes. cykl.
           bool* b) { // Tablica ciecia b.
 for (int i = 0; i < n; i++) b[i] = false; // wykomentuj, jeśli nie 1)
 int p = 0, k = 0, m = 1;
 while (p < n) {
   if (m == n \text{ or } s[m] < s[k]) {
     if (p < n / 2) cyk = p; // wykomentuj, jeśli nie 2)
     while (p <= k) {
       p += m - k:
       if (p < n) {
         suf = p; // wykomentuj, jeśli nie 3)
         b[p] = true; // wykomentuj, jeśli nie 1)
     m = (k = p) + 1;
  } else if (s[m++] != s[k++]) k = p;
                                 Strings/Manacher.cpp
// @s[0..n-1]  - napis długości @n.
// @r[0..2n-2] - tablica promieni palindromów.
// s: a b a a b a a c a a b b b b a a c
void Manacher(const char* s, int n, int* r) {
 for (int i = 0, m = 0, k = 0, p = 0; i < 2 * n - 1; m = i++ - 1) {
   while (p < k \text{ and } i / 2 + r[m] != k)
     r[i++] = min(r[m--], (k + 1 - p++) / 2);
   while (k + 1 < n \text{ and } p > 0 \text{ and } s[k + 1] == s[p - 1])
     k++, p--;
   r[i] = (k + 1 - p++) / 2;
bool is pal(int p, int k) {    //Przedział [p, k] obustronnie domknięty
 return r[p + k] >= (k - p + 1) / 2;
                                    Strings/Pref.cpp
void Pref(const char* s, int n, int* p) {
 p[0] = n;
 int i = 1. m = 0:
 while (i < n) {
   while (m + i < n \text{ and } s[m + i] == s[m]) m++;
   p[i++] = m;
```

```
m = max(m - 1, 0);
   for (int k = 1; p[k] < m; m--) p[i++] = p[k++];
 }
                                     Strings/Suf.cpp
void Sufar(const int* s, int n, int alpha, int* sa, int* lcp = nullptr) {
 if (n > 0) sa[0] = 0;
 if (n <= 1) return;</pre>
 vi roz(alpha + 1), wsk(alpha), typ(n + 1), ids(n, -1), news, pos;
 auto star = [&](int i) -> bool { return typ[i] == 3; };
 auto Indukuj = [&]() -> void {
   copy(roz.begin(), roz.end() - 1, wsk.begin());
   sa[wsk[s[n - 1]]++] = n - 1;
   for(int i = 0; i < n; ++i)
     if (sa[i] > 0 \&\& !typ[sa[i] - 1])
       sa[wsk[s[sa[i] - 1]]++] = sa[i] - 1;
   copy(roz.begin() + 1, roz.end(), wsk.begin());
   for(int i = n-1; i >= 0; --i)
     if (sa[i] > 0 \&\& typ[sa[i] - 1])
       sa[--wsk[s[sa[i] - 1]]] = sa[i] - 1;
 typ[n] = 3;
 for(int i = n-1; i >= 0; --i) {
   sa[i] = -1;
   roz[s[i] + 1]++;
   if (i != n - 1 \&\& s[i] < s[i + 1] + !!typ[i + 1]) {
     typ[i] = 1;
   else if (typ[i + 1]) {
     typ[i + 1] = 3;
 partial sum(all(roz), roz.begin());
 copy(roz.begin() + 1, roz.end(), wsk.begin());
 for(int i = 0; i < n; ++i) if (star(i)) sa[--wsk[s[i]]] = i;
 Indukuj();
 int nast id = 0, b = -1;
 for(int \bar{i} = 0; i < n; ++i) {
   int a = sa[i]:
   if (!star(a)) continue;
   if (b >= 0) while (a == sa[i] || !star(a) || !star(b)) {
     if (star(a) != star(b) || s[a++] != s[b++]) {
       nast id++;
       break;
   ids[b = sa[i]] = nast id;
 for (int i = 0: i < n: i++) {
   if (ids[i] == -1) continue;
   news.pb(ids[i]);
   pos.pb(i);
 vi new sa(siz(news));
 Sufar(news.data(), siz(news), nast id + 1, new sa.data());
 fill(sa, sa + n, -1);
 copy(roz.begin() + 1, roz.end(), wsk.begin());
 reverse(all(new sa)):
 for (int j : new sa) sa[--wsk[s[pos[j]]]] = pos[j];
 Indukuj();
 if (lcp) {
   for (int i = 0; i < n; ++i) ids[sa[i]] = i;
   for (int i = 0, k = 0; i < n; i++, k = max(0, k - 1)) {
     if (ids[i] == n - 1) { k = 0; continue; }
     const int j = sa[ids[i] + 1];
     while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
     lcp[ids[i]] = k;
```

```
Strings/Ukkonen.cpp
template <tvpename Char>
struct Ukkonen {
 // Musi być ściśle większe niż jakakolwiek długość słowa.
 static constexpr int kInfinity = numeric limits<int>::max();
 struct Node {
   map<Char, pair<Node*, pii>> transition;
   Node* suflink;
 // Ta metoda jest wywoływana zawsze gdy tworzona jest krawedź {node}[a, +oo).
 void CreateLeafCallback(Node* node. int a) {
  // printf("CreateLeafCallback({%p}[%d, +oo))\n", node, a);
 // Ta metoda jest wywoływana zawsze gdy krawedź {node}[a, b] zamienia sie
 // w dwie krawędzie: {node}[a, c-1], {middle}[c, b].
 void SplitEdgeCallback(Node* node, int a, int b, Node* middle, int c) {
  // printf("SplitEdgeCallback({%p}[%d, %d] -> {%p}[%d, %d] + {%p}[%d, %d])\n",
            node, a, b, node, a, c - 1, middle, c, b);
   //
 }
 // vector<unique ptr<Node>> nodes to delete; // Odkomentować w celu usuwania.
 Node* NewNode() {
   Node* node = new Node():
   // nodes to delete.insert(node);
                                               // Odkomentować w celu usuwania.
   return node:
 vector<Char> text: // Słowo powinno zaimować indeksy [0.,n-1],
 Node* root;
 Node* pin;
 Node* last explicit node; // Ostatni wierzchołek ,,explicit''.
 int last length; // Liczba literek do ostatniego wierzchołka ,,implicit''.
 // "reserve" warto ustawić na maksymalna dlugość słowa, ale wcale nie trzeba,
 Ukkonen(const int reserve = 0) : root(nullptr), pin(nullptr) {
   text.reserve(reserve):
   root = NewNode();
   pin = NewNode();
   root->suflink = pin;
   last explicit node = root:
   last length = 0;
 void Canonize(Node** s, int* a, int b) {
   if (b < *a) return;</pre>
   pair<Node*, pii> t = (*s)->transition[text[*a]];
   Node* sp = t.st;
   int ap = t.nd.st:
   int bp = t.nd.nd;
   while (bp - ap <= b - *a) {
     *a = *a + bp - ap + 1;
     *s = sp:
     if (*a <= b) {
       t = (*s)->transition[text[*a]];
       sp = t.st:
       ap = t.nd.st;
       bp = t.nd.nd:
   }
 bool TestAndSplit(Node* s, int a, int b, Char c, Node** ret) {
   if (a <= b) {
     pair<Node*. pii>& t = s->transition[text[a]]:
     Node* sp = t.st;
     int ap = t.nd.st:
     int bp = t.nd.nd;
     if (c == text[ap + b - a + 1]) {
       *ret = s;
```

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return true:
      *ret = NewNode();
     t.nd.nd = ap + b - a:
     t.st = *ret;
     (*ret)->transition[text[ap + b - a + 1]] =
          mp(sp, mp(ap + b - a + 1, bp));
     SplitEdgeCallback(s, ap, bp, *ret, ap + b - a + 1);
     return false:
    *ret = s:
    return s->transition.find(c) != s->transition.end();
  void Update(Node** s. int* a. int i) {
   Node* oldr = root;
    Node* r;
    bool end = TestAndSplit(*s, *a, i - 1, text[i], &r);
    while (!end) {
     CreateLeafCallback(r. i):
     r->transition[text[i]] = mp(nullptr, mp(i, kInfinity));
     if (oldr != root) oldr->suflink = r;
     oldr = r;
     *s = (*s)->suflink:
     Canonize(s, a, i - 1);
     end = TestAndSplit(*s, *a, i - 1, text[i], &r);
   if (oldr != root) oldr->suflink = *s:
  // Dodaje kolejną literę do drzewa.
  void AddLetter(Char z) {
   const int i = siz(text);
    text.pb(z):
    auto it = pin->transition.find(z);
   if (it == pin->transition.end())
      pin->transition[z] = mp(root, mp(i, i));
    Update(&last explicit node, &last length, i);
   Canonize(&last explicit node, &last length, i);
  // Zamienia wszystkie krawedzie: [x, +oo) -> [x, text.size()-1].
  void ClearInfinities(Node* node = nullptr) {
   if (node == nullptr) node = root:
   for (auto& it : node->transition) {
     if (it.nd.nd.nd == kInfinity)
       it.nd.nd.nd = siz(text) - 1;
     else ClearInfinities(it.nd.st);
 }
template <tvpename Char>
constexpr int Ukkonen<Char>::kInfinity;
int main() { // Przykład użycia.
 string s = "abcdefgh#";
 Ukkonen<char> u(s.size() /* reserve */);
 for (char c : s) u.AddLetter(c):
 u.ClearInfinities();
  return 0:
```

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