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```
#ifndef LOCAL
#pragma GCC optimize ("03")
#endif
#include <bits/stdc++.h>
using namespace std:
//st, nd, mp, pb, eb
#define siz(c) ((int)(c).size())
#define all(c) (c).begin(), (c).end()
#define sim template < class c
#define ris return * this
#define mor > muu & operator << (</pre>
#define R22(r) sim>typename enable if<1 r sizeof dud<c>(0),muu&>::type operator<<(c g){
sim > struct rae { c b. e: }:
sim > rge<c> range(c b, c e) { return {b, e}; }
sim > auto dud(c* r) -> decltype(cerr << *r);
sim > char dud(...);
struct muu {
#ifdef LOCAL
stringstream a;
~muu() {    cerr << a.str() << endl;    }
R22(<) a << boolalpha << q: ris: }
R22(==) ris << range(begin(g), end(g)); }
sim, class b mor pair < b, c > r) { ris << "(" << r.st << ", " << r.nd << ")"; }
sim mor rge<c> u) {
 a << "[":
 for (c i = u.b: i != u.e: ++i) *this << ". " + 2 * (i == u.b) << *i:
 ris << "]";
template <class...c mor tuple<c...> x) {
 int a = 0:
 a << "(":
 apply([&](c...y){
  ((*this << ", " + 2 * !q++ << y), ...);
 }, x);
 ris << ")";
#define gel(t) sim. class d. class...e mor t < c.d.e... > x) {ris} << *(d*)&x:}
gel(stack) gel(queue) gel(priority queue)
#else
sim mor const c&) { ris; }
#endif
#define imie(r...) "[" #r ": " << (r) << "] "
#define range(b, e) "[[" #b ", " #e "): " << range(b, e) << "] "
#define debug muu() << FUNCTION << "#" << LINE << ":
#define endl '\n'
//using ll,ld,pii,vi,vpii,ull,pdd
sim > void mini(c \&a, const c \&b) {if (a > b) a = b;}
sim > void maxi(c \&a. const c \&b) {if (a < b) a = b:}
int32 t main() {ios base::sync with stdio(0); cin.tie(0);}
                                       geany-config
Interface:
 Message window: right
Keybindīnas:
 Build: Compile: F8, Build: F9
 Focus: Switch to Editor: F1, Switch to VTE: F2
Set Build Commands:
 compile: g++ -Wall -Wshadow -o "%e" "%f" -g -03 -std=c++1z
 build: q++ -Wall -Wshadow -o "%e" "%f" -fsanitize=undefined -fsanitize=address -DLOCAL
-D GLIBCXX DEBUG -q -std=c++1z
                                Rzeczy Na Dzien Probny.txt

    int128/ uint128 t

2. Czy rand() działa tak samo na sprawdzaczce.
3. while (clock() < 1.9 * CLOCKS PER SEC)</pre>
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```
    Czy time(NULL) daje różne wartości między uruchomieniami

5. Czy kompresia (wklejanie binarnych rzeczy do kodu) działa
6. float128
7. ordered set
8. getchar unlocked/ getchar nolock
9. Porównać czas działania sprawdzaczki i kompa
10. gp hash table / cc hash table
  #include <ext/pb ds/assoc container.hpp>
  using namespace gnu pbds;
                               Data Structures/OrderedSet.cpp
#include<ext/pb ds/assoc container.hpp> // ordered set
#include<ext/pb ds/tree policy.hpp>
using namespace gnu pbds; template <typename T> using ordered set =
    tree<T, null type, less<T>, rb tree tag, tree order statistics node update>:
ordered set<int> s; s.insert(1); s.insert(2);
 s.order of key(1); // Out: 0.
*s.find by \overline{o}rder(1); // Out: 2.
                                 Data Structures/Treap.cpp
struct node {
  node *L, *R;
  int ind, prior, sub, lazy; // sub opcionalne
  node(int ind) : L(0), R(0), ind(ind), prior(rand()), sub(1), lazv(0) {}
void rev(node* v) { // przykladowy update, odwraca kolejnosc w poddrzewie
 v->lazy ^= 1;
  swap(v->L, v->R):
void push(node* v) { // opcionalne
 if (v->lazy) {
    if (v->L) rev(v->L):
    if (v->R) rev(v->R):
    v \rightarrow lazy = 0;
node* attach(node* v, node* l, node* r) {
 v->L = l; // jesli chcemy trzymac ojca to update w tej funkcji
  v \rightarrow R = r:
  v->sub = 1 + (l ? l->sub : 0) + (r ? r->sub : 0); // opcjonalne
  return v:
node* merge(node* v, node* u) {
  if (!u) return v;
  if (!v) return u:
  push(v):
  push(u);
  if (v->prior > u->prior) return attach(v, v->L, merge(v->R, u));
  else return attach(u, merge(v, u->L), u->R);
pair<node*, node*> split size(node* v, int k) { // (prefiks rozmiaru k, reszta)
  if (!v) return mp(v, v):
  int lewo = v->L ? v->L->sub : 0:
  push(v):
  if (lewo >= k) {
    auto s = split size(v->L, k);
    return mp(s.st. attach(v. s.nd. v->R)):
 } else {
    auto s = split size(v->R, k - lewo - 1);
    return mp(attach(v, v->L, s.st), s.nd);
pair<node*, node*> split lex(node *v, int k) { // (ind <= k, reszta)</pre>
 if (!v) return mp(v, v);
  if (k < v->ind) {
    auto s = split lex(v->L, k):
    return mp(s.st, attach(v, s.nd, v->R));
  } else {
    auto s = split lex(v->R, k);
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return mp(attach(v, v->L, s.st), s.nd);
int find pos(node *v, int val) { // -1 jesli nie ma
 if (!v) return -1;
 int lewo = v->L ? v->L->sub : 0;
 if (v->ind == val) return 1 + lewo:
 if (val < v->ind) return find pos(v->L, val);
 else return 1 + lewo + find pos(v->R, val);
                              Flows And Matchings/Dinic.cpp
using T = long long;
struct Flow {
 struct E {
   int dest:
  T orig, *lim, *rev;
 int zr, uj, n = 0;
 vector<unique ptr<T>> ts:
 vector<vector<E>> graf:
 vi ptr, odl;
 void vert(int v) {
   n = max(n, v + 1):
   graf.resize(n);
   ptr.resize(n);
   odl.resize(n);
 bool iszero(T v) {
   return !v; // Zmienić dla doubli.
 void bfs() {
   fill(all(odl), 0):
   vi kol = {zr};
   odl[zr] = 1:
   for (int i = 0: i < siz(kol): i++) {
     for (E& e : graf[kol[i]]) {
       if (!odl[e.dest] and !iszero(*e.lim)) {
         odl[e.dest] = odl[kol[i]] + 1;
         kol.pb(e.dest):
     }
 }
T dfs(int v, T lim) {
   if (v == uj) return lim;
   T ret = 0. wez:
   for (int\& i = ptr[v]: i < siz(graf[v]): i++) {
     E\& e = graf[v][i];
     if (odl[e.dest] == odl[v] + 1 and !iszero(*e.lim) and
         !iszero(wez = dfs(e.dest, min(*e.lim, lim)))) {
       ret += wez:
       *e.lim -= wez:
       *e.rev += wez;
       lim -= wez:
       if (iszero(lim)) break:
   }
   return ret;
 void add edge(int u, int v, T lim, bool bi = false /* bidirectional? */) {
   vert(max(u, v));
   T *a = new T(lim), *b = new T(lim * bi);
   ts.eb(a):
   ts.eb(b):
   graf[u].pb(E{v, lim,
                             a, b});
   graf[v].pb(E{u, lim * bi, b, a});
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```
T dinic(int zr_, int uj_) {
    zr = zr ; uj = uj ;
   vert(max(zr, ui));
   T ret = 0;
    while (true) {
     hfs():
     fill(all(ptr), 0);
     const T sta = dfs(zr, numeric limits<T>::max()); // Dla doubli można dać
     if (iszero(sta)) break:
                                                        // infinitv() zamiast
     ret += sta:
                                                        // max().
   return ret;
 vi cut() {
   vi ret;
   bfs();
    for (int i = 0; i < n; i++)
     if (odl[i])
       ret.pb(i);
   return ret;
 map<pii, T> get flowing() { // Tam gdzie plynie 0 może nie być
   map<pii, T> ret:
                              // krawędzi.
   for (int i = 0; i < n; i++)
     for (E& e : graf[i])
       if (*e.lim < e.orig)</pre>
          ret[mp(i, e.dest)] += e.orig - *e.lim;
    for (auto& i : ret) {
     const pii rev{i.st.nd. i.st.st}:
     const T x = min(i.nd, ret[rev]);
     i.nd -= x;
     ret[rev] -= x;
   return ret:
 }
                            Flows And Matchings/Gomory Hu.cpp
#include "Dinic.cop"
struct GomoryHu { // #define int ll//jeśli long longi potrzebne
 vector<vpii> graph, tree;
 vector<vi> nodes:
 vector<bool> visited:
                                   //wymaga naszego dinica
 vi groupId, contrId;
                                   //numeracia od zera
 int wnode. n:
 GomorvHu(int N): graph(N), visited(N), groupId(N), contrId(N), tree(N), n(N) {}
 void addEdge(int u, int v, int cap) {
   graph[u].eb(v, cap);
   graph[v] eb(u, cap);
 void dfs(int v, int type) {
   visited[v] = true; contrId[v] = type;
    for (auto P : tree[v]) { if (!visited[P.st]) { dfs(P.st, type); } }
 vector <pair<pii,int>> run() {
   vi allNodes(n);
   iota(all(allNodes), 0);
   nodes = vector<vi>{allNodes}:
   tree = vector<vpii>(n):
    fill(all(groupId), 0);
    for (int step = 1; step < n; step++) {</pre>
     Flow flow:
     for (int i = 0; i < siz(nodes); i++) {
       if (siz(nodes[i]) > 1) { wnode = i; break; }
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visited[wnodel = true:
     for (auto P : tree[wnode]) { dfs(P.st, nodes[P.st][0]); }
     for (int v = 0; v < n; v++) {
       int a = groupId[v] == wnode ? v : contrId[groupId[v]];
       for (auto& P : graph[v]) {
         int b = groupId[P.st] == wnode ? P.st : contrId[groupId[P.st]];
         if (a != b) { flow.add edge(a, b, P.nd); }
     int a = nodes[wnode][0], b = nodes[wnode][1], f = flow.dinic(a, b);
     auto pom = flow.cut();
     vector <bool> cut(n, false);
     for (int i : pom)
       cut[i]=1:
     for (int v = \theta; v < step; v++) {
       if (v == wnode) { continue; }
       for (auto& P : tree[v]) {
         if (P.st == wnode \&\& !cut[contrId[v]]) \{ P.st = step; \}
     vpii PA. PB:
     for (auto& P : tree[wnode]) { (cut[contrId[P.st]] ? PA : PB).pb(P); }
     tree[wnode] = PA; tree[step] = PB;
     tree[wnode].eb(step, f);
     tree[step].eb(wnode, f):
     vi A. B:
     for (int v : nodes[wnode]) {
       (cut[v] ? A : B).pb(v);
       if (!cut[v]) { groupId[v] = step; }
     nodes[wnode] = A;
     nodes.pb(B);
   vector <pair<pii,int>> res;
   for (int i = 0; i < n; i++)
     for (auto P : tree[i])
       if (nodes[i][0]<nodes[P.st][0])</pre>
          res.pb({{nodes[i][0], nodes[P.st][0]}, P.nd});
   return res;
}; // #undef int
                    Flows And Matchings/Highest Label Push Relabel.cpp
class LinkedList {
public:
 LinkedList(int N) : N(N). next(N) { clear(); }
 void clear() { head.assign(N, -1); }
 int front(int h) { return head[h]; }
 void pop(int h) { head[h] = next[head[h]]; }
 void push(int h. int u) { next[u] = head[h]. head[h] = u: }
private:
 int N;
 vi next, head;
class DoublyLinkedList {
private:
 struct Node { int prev, next; };
public:
 DoublyLinkedList(int N) : N(N), nodes(N) { clear(); }
 void clear() { head.assign(N, -1); }
 void erase(int h, int u) {
   int pv = nodes[u].prev, nx = nodes[u].next;
   if (nx >= 0) nodes[nx].prev = pv;
   if (pv >= 0) nodes[pv].next = nx;
   else head[h] = nx;
```

```
void insert(int h, int u) {
    nodes[u] = \{ -1, head[h] \};
    if (head[h] >= 0) nodes[head[h]].prev = u;
    head[h] = u;
  template <typename Func>
  void erase all(int first, int last, Func f) {
    for (int i = first; i <= last; ++i) {</pre>
      for (int h = head[i]: h >= 0: h = nodes[h].next) f(h):
      head[i] = -1:
 }
private:
 int N:
 vi head;
 vector<Node> nodes;
template <
    typename CapType, typename TotalCapType,
    bool UseGlobal = true, bool UseGap = true
class HighestLabelPushRelabel {
private:
 TotalCapType inf = pow(10., sizeof(TotalCapType) / 4 * 9);
  struct Edge { int to, rev: CapType cap: }:
public:
  HighestLabelPushRelabel(\underbrace{int}\ N)\ :\ N(N),\ E(\theta),\ G(N),\ que(N),\ list(N),\ dlist(N)\ \{\}
  TotalCapType maximum flow(int s, int t) {
    if (s == t) return 0:
    highest active = 0; // highest label (active)
    highest = 0; // highest label (active and inactive)
    height.assign(N, 0); height[s] = N;
    for (int i = 0: i < N: ++i) if (i != s) dlist.insert(height[i]. i):
    count.assign(N, 0); count[0] = N - 1;
    excess.assign(N, \theta); excess[s] = inf; excess[t] = -inf;
    for (auto\& e : G[s]) push(s, e);
    global relabel(t):
    for (int u = -1, rest = N; highest_active >= 0; ) {
      if ((u = list.front(highest active)) < 0) { --highest active; continue; }
      list.pop(highest active);
      discharge(u):
      if (--rest == 0) rest = N, global relabel(t);
    return excess[t] + inf;
  inline void add directed edge(int u, int v, CapType cap) {
    G[u].pb({v, siz(G[v]), cap});
    G[v].pb(\{u. siz(G[u]) - 1. 0\}):
  inline void add undirected edge(int u, int v, CapType cap) {
    G[u].pb({v, siz(G[v]), cap});
    G[v].pb(\{u, siz(G[u]) - 1, cap\}):
private:
 void global relabel(int t) {
    if (!UseGlobal) return:
    height.assign(N, N); height[t] = 0;
    count.assign(N, 0);
    int gh = 0, gt = 0;
    for (que[qt++] = t; qh < qt; ) {
      int u = que[qh++], h = height[u] + 1;
      for (auto\& e : G[u]) if (height[e.to] == N \&\& G[e.to][e.rev].cap > 0) {
          count[height[e.to] = h] += 1;
```

fill(all(visited), false);

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que[qt++] = e.to;
   list.clear(); dlist.clear();
   for (int u = 0; u < N; ++u) if (height[u] < N) {
       dlist.insert(height[u], u):
       if (excess[u] > 0) list.push(height[u], u);
   highest = highest active = height[que[qt - 1]];
 void push(int u, Edge& e) {
   int v = e.to;
   CapType df = min(excess[u], TotalCapType(e.cap));
   e.cap -= df. G[v][e.rev].cap += df:
   excess[u] -= df, excess[v] += df;
   if (0 < excess[v] && excess[v] <= df) list.push(height[v], v);</pre>
 void discharge(int u) {
   int nh = N:
   for (auto& e : G[u]) if (e.cap > 0) {
       if (height[u] == height[e.to] + 1) {
         push(u, e):
         if (excess[u] == 0) return;
       } else nh = min(nh, height[e.to] + 1);
   int h = height[u]:
   if (UseGap && count[h] == 1) {
     auto f = [\&] (int u) { count[height[u]]--, height[u] = N; };
     dlist.erase all(h, highest, f);
     highest = h - 1:
   } else {
     count[h]--; dlist.erase(h, u);
     height[u] = nh;
     if (nh == N) return:
     count[nh]++; dlist.insert(nh, u);
     highest = max(highest, highest active = nh);
     list.push(nh, u);
 int N, E, highest active, highest;
 vector< vector<Edge> > G;
 vi height, count, que:
 vector<TotalCapType> excess;
 LinkedList list;
 DoublyLinkedList dlist;
using HIPR GP = HighestLabelPushRelabel<int, int, true, true>;
                           Flows And Matchings/Hungarian.cpp
//a[1..n][1..m] - wagi, n<=m, O(n^2*m), znajduje minimalny koszt (masymalny: *-1)
//u[0..n], v[0..m] - funkcja potencjalu, p[0..m], ans[0..n] - skojarzenie
struct hungarian { // cost = -v[0]
 int n, m;
 vector<vi> a:
 vi u, v, p, way, ans;
 hungarian(int n, int m) : n( n), m( m) {
   #define re(x, y) (x).resize(y + 1)
   re(a, n); re(u, n); re(v, m); re(p, m); re(way, m), re(ans, n);
   for (auto \&x: a) re(x, m):
 void addcost(int v1, int v2, int c) {
   a[v1][v2] = c;
 void solve() {
   for(int i = 1; i <= n; ++i) {</pre>
     p[0] = i;
```

```
int i0 = 0;
      vi minv(m + 1, INT MAX); // uwaga, moze byc potrzebny LL!
      vector < char > used(m + 1, false):
        used[i0] = true;
        int i0 = p[i0], delta = INT MAX, i1 = 0;
        for (int j = 1; j <= m; ++j\bar{j}) {
          if (!used[i]) {
            int cur = a[i0][i] - u[i0] - v[i];
            if (cur < minv[i]) {</pre>
              minv[j] = cur;
              way[j] = j0;
            if (minv[i] < delta) {</pre>
              delta = minv[j];
              j1 = j;
        for (int j = 0; j \le m; ++j) {
          if (used[j]) {
            u[p[j]] += delta;
            v[i] -= delta;
          else {
            minv[j] -= delta;
        i0 = i1:
      } while (p[j0] != 0);
      do {
        int i1 = way[i0];
        p[j0] = p[j1];
        j0 = j1;
     } while (j0);
    for (int j = 1; j \le m; ++j)
     ans[p[j]] = j; // odzyskiwanie wyniku
 }
                              Flows And Matchings/Matching.cpp
struct Matching {
 int n, tim = 0, top = 0;
 vi mat, fa, s, q, pre, vis, head;
 vpii e:
 Matching(int N): n(N + 1), mat(n, -1), fa(n), s(n), g(n), pre(n), vis(n), head(n, -1)
<mark>.</mark>{}
 void edge impl(int x, int y) {e.eb(y, head[x]);head[x] = siz(e) - 1;}
 void add edge(int x, int y) {edge impl(x, y), edge impl(y, x);}
  int find(int x) {return x == fa[x] ? x : fa[x] = find(fa[x]);}
 int lca(int x, int y) {
    for (++tim, x = find(x), y = find(y); ; swap(x, y)) {
     if (~x) {
        if (vis[x] == tim) return x;
        vis[x] = tim;
        x = \text{-mat}[x] ? find(pre[mat[x]]) : -1;
 void blossom(int x, int y, int l) {
   while (find(x) != l) {
     pre[x] = y;
      if (s[mat[x]] == 1) s[q[top++] = mat[x]] = 0;
      if (fa[x] == x) fa[x] = 1;
     if (fa[mat[x]] == mat[x]) fa[mat[x]] = l;
     y = mat[x];
      x = pre[y];
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for (int i = 0; i < top; ++i) {
     for (int t = head[q[i]]; ~t; t = e[t].nd) {
       int y = e[t].st;
       if (s[v] == -1) {
          pre[y] = q[i];
         s[y] = 1;
          if (mat[y] == -1) {
            for (int u = v, v = a[i], lst: \sim v: u = lst, v = \sim u? pre[u]: -1)
             lst = mat[v], mat[v] = u, mat[u] = v;
           return true:
         s[q[top++] = mat[y]] = 0;
       else if (!s[y] \&\& find(y) != find(q[i])) {
         int l = lca(v, q[i]);
          blossom(v. a[i], l):
          blossom(q[i], y, l);
     }
   return false;
 }
 int run() {
   int size = 0:
   for (int i = 0: i < n: ++i)
     if (mat[i] == -1 \&\& match(i))
       size++;
   return size:
 }
                             Flows And Matchings/Min Cost.cpp
struct MinCost {
 struct kra {
   int cel, *prze1, *prze2;
   ll koszt;
 };
 int n=0. zr. ui:
 const ll inf=1e9;
 vector <vector <kra>> graf;
 vi bvlo. aktu:
 vector <ll> odl:
 void vert(int v) {
   if (v>n) {
     n=v:
     graf.resize(n);
     bylo.resize(n);
     aktu.resize(n);
     odl.resize(n):
   }
 }
 void add edge(int v, int u, int prze, ll koszt) {
   vert(v+1): vert(u+1):
   kra ret1{u, new int(prze), new int(0), koszt};
   kra ret2{v, ret1.prze2, ret1.prze1, -koszt};
   graf[v].pb(ret1);
   graf[u].pb(ret2):
 void spfa() {
   for (int i=0; i<n; i++) {</pre>
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```
aktu[i]=1;
      odl[i]=inf:
   aktu[zr]=odl[zr]=0;
    queue <int> kol;
    kol.push(zr):
    while(!kol.empty()) {
     int v=kol.front();
      kol.pop();
     if (aktu[v])
       continue:
      aktu[v]=1;
      for (kra i : graf[v]) {
       if (*i.przel && odl[v]+i.koszt<odl[i.cel]) {</pre>
          odl[i.cel]=odl[v]+i.koszt;
          aktu[i.cel]=0;
          kol.push(i.cel);
   }
 int dfs(int v) {
   if (v==ui)
      return 1:
   bylo[v]=1;
    for (int i=0; i < siz(graf[v]); i++) {</pre>
     if (!bylo[graf[v][i].cel] && (*graf[v][i].przel) &&
      odl[v]+graf[v][i].koszt==odl[graf[v][i].cel] && dfs(graf[v][i].cel)) {
       (*graf[v][i].prze1)--;
       (*graf[v][i].prze2)++:
       return 1:
   }
   return 0:
 pair <int,ll> flow(int zrzr, int ujuj) {
   zr=zrzr; ui=uiui;
   vert(zr+1): vert(ui+1):
   pair <int, ll> ret{0, 0};
   while(1) {
     spfa();
      for (int i=0: i<n: i++)</pre>
       bylo[i]=0;
      if (!dfs(zr))
       break:
      ret.st++:
     ret.nd+=odl[uil:
    return ret;
                          Flows And Matchings/Turbo Matching.cpp
struct matching {
 int n:
 vector<vi> V:
 vector<bool> odw;
 vi para, strona;
 matching(int n): n(n) { // zakladam numeracie od 1. nie moze bvc 0
   V.resize(n + 1);
   odw.resize(n + 1);
   para.resize(n + 1);
 void addedge(int a, int b) { // zakladam ze a jest z lewej, b z prawej
   V[a].pb(b);
    strona.pb(a);
```

bool match(int x) {

iota(all(fa), 0):

fill(all(s), -1);

s[q[0] = x] = 0;

top = 1:

```
D
```

```
Ν
```

```
bool skojarz(int x) { // x nalezv do strona
   odw[x] = 1:
    for (auto v : V[x]) {
     if (!para[v] || (!odw[para[v]] && skojarz(para[v]))) {
        para[v] = x;
        para[x] = v;
        return 1;
   } }
   return 0:
 vpii go() {
   sort(all(strona));
   strona.resize(unique(all(strona)) - strona.begin());
   for (int i = 1: i \le n: ++i)
     shuffle(all(V[i]), default random engine(2137));
    sort(all(strona), [&](int a, int b){return siz(V[a]) < siz(V[b]);});</pre>
    bool ok = 0:
    do {
     ok = 0:
      for (auto i : strona) odw[i] = 0;
      for (auto i : strona)
        if (!para[i] && skojarz(i))
          ok = 1:
   } while (ok);
    vpii res:
   for (auto i : strona)
     if (para[i])
        res.pb(mp(i, para[i]));
   return res:
} };
                        Flows And Matchings/Weighted Matching.cpp
static const int INF = INT MAX;
static const int N = 1017:
struct edge {
 int u, v, w;
 edge() {} edge(int ui, int vi, int wi) : u(ui), v(vi), w(wi) {}
int n, n_x;
edge q[N^* 2][N * 2];
int lab[N * 2], match[N * 2], slack[N * 2], st[N * 2], pa[N * 2];
int flo from[N * 2][N + 1], S[N * 2], vis[N * 2];
vi flo(\overline{N} * 2\overline{1}):
queue<int> q;
int e delta(const edge &e) { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
void update slack(int u, int x) {
 if (!slack[x] \mid | e delta(q[u][x]) < e delta(q[slack[x]][x])) slack[x] = u;
void set slack(int x) {
 slack[x] = 0:
 for (int u = 1: u \le n: u++)
   if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
      update slack(u, x);
void q push(int x) {
 if (x \le n) q.push(x);
 else for (int i = 0; i < siz(flo[x]); i++) q push(flo[x][i]);
void set st(int x, int b) {
 st[x] = b;
 if (x > n) for (int i = 0; i < siz(flo[x]); i++) set st(flo[x][i], b);
int get pr(int b, int xr) {
 int pr = find(all(flo[b]), xr) - flo[b].begin();
 if (pr % 2 == 1) {
```

```
reverse(flo[b].begin() + 1, flo[b].end());
    return siz(flo[b]) - pr;
  } else return pr:
void set match(int u, int v) {
  match[\overline{u}] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo from[u][e.u], pr = get pr(u, xr);
  for (int i = 0; i < pr; i++) set match(flo[u][i], flo[u][i ^ 1]);
  set match(xr, v):
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
void augment(int u. int v) {
  for (::) {
    int xnv = st[match[u]];
    set match(u, v);
    if (!xnv) return:
    set match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
   if (u == 0) continue:
    if (vis[u] == t) return u:
    vis[u] = t; u = st[match[u]];
   if (u) u = st[pa[u]];
  return 0:
void add blossom(int u, int lca, int v) {
  int b = n + 1; while (b \le n \times \&\& st[b]) ++b;
  if (b > n x) ++n x:
  lab[b] = \overline{0}, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear(); flo[b].pb(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].pb(x), flo[b].pb(y = st[match[x]]), q push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].pb(x), flo[b].pb(y = st[match[x]]), q push(y);
  set st(b, b);
  for (int x = 1; x \le n x; ++x) q[b][x].w = q[x][b].w = 0;
  for (int x = 1: x <= n: ++x) flo from[b][x] = 0:
  for (int i = 0; i < siz(flo[b]); i++) {
    int xs = flo[b][i];
    for (int x = 1; x \le n x; ++x)
      if (g[b][x].w == 0 \mid [e_delta(g[xs][x]) < e_delta(g[b][x]))
        q[b][x] = q[xs][x], q[x][b] = q[x][xs];
    for (int x = 1; x \le n; ++x) if (flo from[xs][x]) flo from[b][x] = xs;
 set slack(b):
void expand blossom(int b) {
 for (int \bar{i} = 0; i < siz(flo[b]); ++i)
    set st(flo[b][i]. flo[b][i]):
  int xr = flo from[b][g[b][pa[b]].u], pr = get pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = q[xns][xs].u:
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set slack(xns);
    a push(xns);
```

```
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M
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S[xr] = 1, pa[xr] = pa[b]:
 for (int i = pr + 1; i < siz(flo[b]); i++) {</pre>
   int xs = flo[b][i];
   S[xs] = -1, set slack(xs);
 st[b] = 0:
bool on found edge(const edge &e) {
 int u = st[e.u], v = st[e.v];
 if (S[v] == -1) {
   pa[v] = e.u, S[v] = 1; int nu = st[match[v]];
   slack[v] = slack[nu] = 0; S[nu] = 0, gpush(nu);
 } else if (S[v] == 0) {
   int lca = get lca(u, v);
   if (!lca) return augment(u, v), augment(v, u), true;
   else add blossom(u, lca, v);
 return false;
bool matching() {
 memset(S + 1, -1, sizeof(int) * n x):
 memset(slack + 1, 0, sizeof(int) \overline{*} n x);
 q = queue<int>();
 for (int x = 1; x \le n x; x++) {
   if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0, q push(x);
 if (q.empty()) return false;
 for (;;) {
   while (siz(a)) {
     int u = q.front(); q.pop();
     if (S[st[u]] == 1) continue;
     for (int v = 1; v \le n; v++) {
       if (a[u][v], w > 0 \&\& st[u] != st[v]) {
          if (e delta(q[u][v]) == 0) {
            if (on found edge(g[u][v])) return true;
          } else update slack(u, st[v]);
   int d = INF;
   for (int b = n + 1; b \le n \times (b++) {
     if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2):
   for (int x = 1; x \le n x; x++) {
     if (st[x] == x \&\& slack[x]) {
       if (S[x] == -1) d = min(d, e delta(g[slack[x]][x]));
       else if (S[x] == 0) d = min(d, e delta(g[slack[x]][x]) / 2);
   for (int u = 1; u <= n; u++) {
     if (S[st[u]] == 0) {
       if (lab[u] <= d) return 0;</pre>
       lab[u] -= d:
     } else if (S[st[u]] == 1) lab[u] += d;
   for (int b = n + 1; b \le n \times (b++) {
     if (st[b] == b) {
       if (S[st[b]] == 0) lab[b] += d * 2;
       else if (S[st[b]] == 1) lab[b] -= d * 2;
   q = queue<int>();
   for (int x = 1; x <= n x; x++)
     if (st[x] == x \& \& slack[x] \& \& st[slack[x]] != x \& \& e delta(g[slack[x]][x]) == 0)
```

```
if (on found edge(g[slack[x]][x])) return true;
    for (int b = n + 1: b \le n \times (b++)
      if (st[b] == b \&\& S[b] == 1 \&\& lab[b] == 0) expand blossom(b);
 return false;
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
 n x = n;
 int n matches = 0:
  long \overline{l}ong tot weight = 0:
  for (int u = \overline{0}; u \le n; u++) st[u] = u, flo[u].clear();
  int w max = 0;
  for (\overline{int} \ u = 1; \ u <= n; \ u++)
    for (int v = 1; v \le n; v++) {
      flo from[u][v] = (u == v ? u : 0);
      w \max = \max(w \max, q[u][v].w);
  for (int u = 1; u \le n; u ++) lab[u] = w max;
  while (matching()) ++n matches;
  for (int u = 1; u <= n; u++) {
   if (match[u] && match[u] < u) {</pre>
      tot weight += g[u][match[u]].w;
 return mp(tot weight, n matches):
void add edge(int u, int v, int w) {
 q[u][v].w = q[v][u].w = max(q[u][v].w, w);
void init(int n) {
  for \overline{(int u = 1; u \leq n; u++)}
   for (int v = 1: v \le n: v++) {
      g[u][v] = edge(u, v, 0);
                                     Geometry/Dewolai.cpp
typedef long long T;
struct P {
 T x, y;
 int id:
 P operator-(P b) { return P{x - b.x, y - b.y}; }
 T cross(P b) { return x * b.y - y * b.x; }
 T cross(P b, P c) const { return (b - *this).cross(c - *this); }
 T dot(P b) \{ return x * b.x + v * b.v. \}
 bool inTriangle(const P &a, const P &b, const P &c) const {
#define tmp(a, b) (cross(a, b) > 0)
    return tmp(a, b) == tmp(b, c) && tmp(b, c) == tmp(c, a);
#undef tmp
 }
int cmpCircle(P a, P b, P c, P d) {
 P v1 = b - a, v2 = d - a, v3 = b - c, v4 = d - c;
 ld tmp=(ld)abs(v1.cross(v2))*v3.dot(v4)+(ld)v1.dot(v2)*abs(v3.cross(v4));
 if (abs(tmp) < 1e-8) return 0;if (tmp == 0) return 0;</pre>
 if (tmp > 0) return 1; return -1;
struct pair hash {
  template <class T1, class T2>
  std::size t operator()(const std::pair<T1, T2> &p) const {
    return p.first * 10000 + p.second:
 }
unordered map<pair<int, int>, pair<int, int>, pair hash> mt;
```

```
set<pair<int, int>> edges;
vector<vector<int>>> triangles:
void rec(int a, int c, const vector<P> &points);
void trim(int a, int b) {
 assert(a < b);
 auto it = mt.find({a, b});
 if ((it != mt.end()) && (it->second == make pair(-1, -1))) mt.erase(it):
void change(int a, int b, int from, int to) {
 if (a > b) swap(a, b):
 if (!mt.count(\{a, b\})) mt[\{a, b\}] = \{-1, -1\};
 for (int *x : vector<int *>{&mt[{a, b}].first, &mt[{a, b}].second})
   if (*x == from) {
     *x = to:
     trim(a, b):
     return:
 assert(false):
void rec(int a, int c, const vector<P> &points) {
 if (a > c) swap(a, c);
 if (!mt.count({a, c})) return:
 int b = mt[{a, c}].first;
 int d = mt[{a, c}].second;
 if (b > d) swap(b, d);
 if (b == -1 || d == -1) return:
 for (int rep = 0; rep < 2; ++rep) {</pre>
   if (points[a].inTriangle(points[b], points[c], points[d])) return;
   swap(a, c);
 if (cmpCircle(points[a], points[b], points[c], points[d]) != 1) return;
 assert((!mt.count({b, d})) || (mt[{b, d}] == make pair(-1, -1)));
 mt[{b, d}] = {a, c};
 trim(b. d):
 mt.erase(make pair(a, c));
 change(a, b, c, d);change(b, c, a, d);change(a, d, c, b);change(c, d, a, b);
 rec(a, b, points);rec(b, c, points);rec(c, d, points);rec(d, a, points);
void addTriangle(int a, int b, int c) {
 change(a, b, -1, c); change(a, c, -1, b); change(b, c, -1, a);
void anvTriangulation(vector<P> points) {
 sort(points.begin(), points.end(), [](const P &a, const P &b) {
  return make pair(a.x, a.y) < make pair(b.x, b.y);});</pre>
 vector<P> upper, lower;
 for (P C : points) {
#define backback(w) w[(int)w.size() - 21
   while ((int)upper.size() >= 2 \&\& backback(upper).cross(upper.back(), C)>0){}
     addTriangle(C.id, backback(upper).id, upper.back().id);
     upper.pop back():
   upper.push back(C);
   while ((int)lower.size() >= 2 \& lower[(int)lower.size() - 2].cross(lower.back(), C)
     addTriangle(C.id. backback(lower).id. lower.back().id):
     lower.pop back();
   lower.push back(C):
#undef backback
 }//all collinear
 if(lower.size() == upper.size() &&lower.size() == points.size()) assert(false);
const int max = 1e6 + 5:
int memo x[nax], memo y[nax];
void dewolaj() {
```

```
mt.clear();edges.clear();triangles.clear();
  int n:
  scanf("%d". &n):
  vector<P> points(n);
  for (int i = 0; i < n; ++i) {
    scanf("%d%d". &memo x[i]. &memo v[i]):
    points[i] = P\{\text{memo } x[i], \text{ memo } y[i], i\};
 mt.reserve(4123123);
  anvTriangulation(points):
  vector<pair<int. int>> init:
  for (auto ppp : mt) init.push back(ppp.first);
  for (pair<int, int> p : init)
   if (mt.count(p) \&\& mt[p] != make pair(-1, -1)) rec(p.first, p.second, points):
  for (auto ppp : mt)
    if (ppp.second != make pair(-1, -1)) {
      int i = ppp.first.first, j = ppp.first.second;
      assert(i != i):
      edges.insert({i, j});
      if (mt[{i, j}].first > j) triangles.push back(vector<int>{i, j, mt[{i, j}].first});
      if (mt[{i, j}].second > j) triangles.push back(vector<int>{i, j, mt[{i, j}].second}.
<u>,</u>);
  // edges zawiera krawędzie triangulacji
                                  Geometry/Halfplanes.cpp
pii operator-(pii a, pii b) {return {a.first - b.first, a.second - b.second};}
pii rot(pii x) { return {x.second, -x.first}; }
using hpl = pair <pii, ll>; //(v, m) = \{x : sc(x, v) >= m\}
|ll sc(pii a, pii b) {return a.first*1ll*b.first + a.second*1ll*b.second;}
ll ve(pii a, pii b) {return a.first*1ll*b.second - a.second*1ll*b.first;}
//wartości rzędu współrzędne ^ 4
 int128 det(hpl a, hpl b, hpl c) {
 return a.second * ( int128)ve(b.first, c.first) +
    b.second * ( int128)ve(c.first, a.first) +
    c.second * ( int128)ve(a.first, b.first);
bool subset(hpl a, hpl b) {
 if (ve(a.first, b.first) || sc(a.first, b.first) < 0) return false;</pre>
 return a.second * ( int128) abs(b.first.first) >= b.second * ( int128) abs(a.first.fi
<mark>⊾</mark>rst) &&
   a.second * ( int128) abs(b.first.second) >= b.second * ( int128) abs(a.first.second
bool disjoint(hpl a, hpl b) {
 return subset(a, {{-b.first.first.-b.first.second}, -b.second}): }
bool around(hpl a, hpl b, hpl c) {
 ll ab = ve(a.first, b.first);
 ll bc = ve(b.first, c.first);
 ll ca = ve(c.first, a.first):
  assert(ab > 0 || bc > 0 || ca > 0);
  return (ab \geq 0 && bc \geq 0 && ca \geq 0) | (ab \leq 0 && bc \leq 0 && ca \leq 0);
void and sort(vector <hpl> &a) {
  if (a.empty()) return;
  hpl mid = a.back();
  a.pop back();
  vector <hpl> left, right;
  for (hpl c : a) {
   ll v = ve(c.first, mid.first), s = sc(c.first, mid.first);
    if (make pair(v, s) > pair < ll, ll > (0, 0))
      left.push back(c):
    else
      right.push back(c);
```

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C D F G M M N S
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sort(right.begin(), right.end(), [](hpl x, hpl y){return ve(x.first, y.first) > 0;});
 sort(left.begin(), left.end(), [1(hpl x, hpl v){return ve(x,first, v,first) > 0:});
 left.push back(mid):
 left.insert(left.end(), right.begin(), right.end());
 swap(a, left);
//Jeśli przeciecie jest nieograniczone, to półproste idace do nieskończoności sa wyznaczoł
//przez pierwszy i ostatni element zwróconego wektora
//Jeśli przeciecie iest puste, zwraca pusty wektor
vector <hpl> find hull(vector <hpl> vec) {
 vector <hpl> hull;
 int first = 0;
 and sort(vec):
 for (hpl curr : vec) {
   if (!hull.empty() && disjoint(curr, hull.back()))
     return {}; //Przeciecie jest puste
   if (!hull.emptv() && disjoint(curr, hull[first]))
     return {}: //Przeciecie iest puste
   if (!hull.empty() && subset(hull.back(), curr)) //Case kiedy jedna półpłaszczyzna zaw
liera się w drugiej, na ogół można wywalić
     continue:
   if (!hull.emptv() && subset(curr, hull.back())) //J.w.
     hull.pop back():
   while (hull.size() - first >= 2u \&\& det(hull.back(), *(hull.end() - 2), curr) <= 0) {
     if (around(hull[hull.size() - 2], hull.back(), curr))
        return {}: //Przeciecie iest puste
     else
       hull.pop back();
   while (hull.size() - first \geq 2u \&\& det(curr, hull[first], hull[first + 1]) \geq 0) {
     if (around(curr, hull[first], hull[first + 1]))
       return {}; //Przeciecie jest puste
     else
       first++:
   if (hull.size() - first < 2u || det(hull.back(), curr, hull[first]) < 0)</pre>
     hull.push back(curr);
 return vector <hpl>(hull.begin() + first, hull.end());
//Półpłaszczyzna wyznaczona przez prostą xy
hpl make hpl(pii x, pii y) {
 pii v = rot(y - x);
 assert(sc(v, x) == sc(v, y));
 return \{v, sc(x, v)\}:
                              Geometry/Halfplanes Online.cpp
#define X real()
#define Y imag()
typedef complex<LL> P:
struct line {
   LL a.b.c:
   line(LL a = 0, LL b = 0, LL c = 0): a(a), b(b), c(c) {} // <= 10^{\circ}9
   line (P const &A, P const &B): \overline{a}(A.Y-B.Y), \overline{b}(B.X-\overline{A.X}), \overline{c}(A.X*B.Y-A.Y*B.X) {} //pts <=\overline{a}
   line operator - () const {return line(-a, -b, -c); }
   bool up() const { return a?(a<0):(b>0):}
inline LL wek(line const \&a, line const \&b) {return a.a*b.b-a.b*b.a;}
inline bool rown(line a, line b) {return wek(a,b) == 0;}
inline bool pokr(line a. line b) {return rown(a.b) && a.a*b.c == b.a*a.c && a.b*b.c == b.
inline bool podobne(line a, line b) {return rown(a,b) && a.up() == b.up();}
inline complex<LD> prosta prosta(line a, line b) {
```

```
LL det = wek(a,b);
    LL x = -a.c*b.b+b.c*a.b:
    LL v = -a.a*b.c+a.c*b.a:
    return complex<LD>(x,y)/(LD)det;
inline LL weaker (line a, line b) { // czy a jest slabsze niz b
    assert(rown(a.b)):
    if (abs(a.a) > abs(a.b)) return a.c*abs(b.a) - b.c*abs(a.a);
    else return a.c*abs(b.b) - b.c*abs(a.b);
struct Comp {
    bool operator()(const line& a, const line& b) const {
        if (a.up() != b.up()) return a.up() > b.up();
        return wek(a.b) > 0:
    }
const LD EPS = 1e-12;
struct przeciecie polplaszczyzn {
    bool empty, pek:
    set<line. Comp> S:
    typedef set<line, Comp>::iterator iter;
    przeciecie polplaszczvzn() : emptv(false), pek(false) {};
    iter next(iter it){return (++it == S.end() ? S.begin() : it);}
    iter prev(iter it){return (it == S.begin() ? --S.end() : --it);}
    bool hide(line a, line b, line c) {
        if (rown(a.b)) {
            if (weaker(a, -b) < 0) empty = true;</pre>
            return false:
        if (wek(a.b) < 0) swap(a.b):
        complex<LD> r = prosta prosta(a.b);
        LD v = r.X * c.a + r.Y * c.b + c.c;
        if (wek(a,c) >= 0 \&\& wek(c,b) >= 0 \&\& v > -EPS) return true;
        if (wek(a,c) < 0 \& wek(c,b) < 0) {
            if (v < -EPS) emptv = true:</pre>
            else if (v < EPS) pek = true;</pre>
        return false:
    void add(line l) {
        if (empty) return;
        if (l.a == 0 && l.b == 0) {
            if (l.c < 0) emptv = true:</pre>
            return:
        iter it = S.lower bound(l):
        //rownolegle
        if(it != S.end() && podobne(*it, l)) {
            if (weaker(l, *it)>=0) return;
            iter del = it:
            it = next(it):
            S.erase(del);
        if(SZ(S) \ge 2 \&\& it == S.end()) it = S.begin():
        while(SZ(S) \geq 2 && hide(l, *next(it), *it)) {
            iter del = it:
            it = next(it):
            S.erase(del):
        //*it<p
        if(SZ(S) >= 2) it = prev(it):
        while(SZ(S) \geq 2 && hide(l, *prev(it), *it)) {
            iter del = it;
            it = prev(it):
```

```
S.erase(del);
       if(S.size() < 2 || !hide(*it, *next(it), l)) S.insert(l);
   }
        0 - puste 1 - punkt 2 - odcinek 3 - półprosta 4 - prosta
         5 - dodatnie (może nieskończone) pole (S.size() daje wowczas liczbę boków) */
    int type() {
       if(empty) return 0;
       if(SZ(S) \ll 4){
            vector<line> res(ALL(S)):
            if (SZ(res) == 2 \&\& rown(res[0], res[1]) \&\& weaker(res[0], -res[1])<0) return
L 0;
            REP(i, SZ(res)) REP(j, i) if(pokr(res[i], res[j])) {
                if(SZ(res) == 2) return 4:
                if(SZ(res) == 3) return 3:
               if(SZ(res) == 4 \&\& pokr(res[0], res[2]) \&\& pokr(res[1], res[3])) return 1;
                return 2;
            if(SZ(res) == 3 \&\& pek) return 1:
        return 5;
   }
                                    Geometry/Rotor.cpp
pii in[nax];
int wh[nax]:
pair <pii, pii> dir[nax * nax / 2];
pii operator-(pii a, pii b) {return {a.first - b.first, a.second - b.second;}}
ll pro(pii a, pii b) {return a.first*1ll*b.second - a.second*1ll*b.first;}
bool cmp(pair <pii, pii> a, pair <pii, pii> b) {
 ll p = pro(a.first, b.first):
 if (p > 0) return 1; if (p < 0) return 0; return a.second < b.second;
 //Jak nie ma trzech współliniowych to po prostu: return pro(a.first, b.first) > 0;
int main() {
 scanf("%d", &n);
 for (int i = 0; i < n; ++i)
   scanf("%d%d", &in[i].first, &in[i].second);
  sort(in. in + n):
 for (int i = 0; i < n; ++i)
   wh[i] = i;
  int cou = 0;
 for (int i = 0: i < n: ++i)
   for (int j = i + 1; j < n; ++j)
     dir[cou++] = {max(in[i] - in[i], in[i] - in[i]), {i, i}};
  sort(dir. dir + cou. cmp):
  for (int i = 0: i < cou: ++i) {
   debug << imie(i) imie(dir[i]);</pre>
   auto c = dir[i].second;
   pii x = in[wh[c.first]], y = in[wh[c.second]];
    swap(wh[c.first], wh[c.second]);
   swap(in[wh[c.first]], in[wh[c.second]]);
    //Policz wynik dla posortowanych prostopadle do prostej xy
 }
                                     Graphs/2sat.cpp
const int nax = 100000:
vector <int> implies[2*nax]: //wvmuszenia, 2*x to zmienna 2*x+1 to iei negacia
int sat val[2*nax],sat vis[2*nax],sat sort[2*nax],sat ile;
inline void sat or(int a,int b){
 implies[a^1].push back(b);
 implies[b^1].push back(a):
void sat dfs mark(int x){
 sat vis[x]=1;
```

```
sat val[x]=sat val[x^1]==-1;
  for (int i : implies[x]) if (!sat vis[i]) sat dfs mark(i):
void sat dfs(int x){
  sat vis[x]=1;
  for (int i : implies[x]) if (!sat vis[i]) sat dfs(i):
  sat sort[--sat ile]=x^1;
bool sat2(int n) {//n - liczba zmiennych, zmienne numerujemy od 0
#define REP for (int i = 0: i < 2 * n: ++i)
  sat ile=2 * n:
  REP sat vis[i]=0,sat_val[i]=-1;
  REP if (!sat vis[i]) sat dfs(i);
 REP sat vis[i]=0:
  REP if (!sat vis[sat sort[i]]) sat dfs mark(sat sort[i]);
  REP if (sat val[i]) for(int j : implies[i]) if (!sat val[j]) return 0;
  return 1;
                                     Graphs/Centroid.cpp
struct centro { // indeksowane od 0, par to drzewo centroidow
 vector<vi> edges;
  vector<bool> vis:
  vi par. sz:
  int n:
  centro(int n) : n(n) {
    edges.resize(n):
    vis.resize(n):
    par.resize(n);
    sz.resize(n);
  void edge(int a, int b) {
    edges[a].pb(b);
    edges[b].pb(a);
  int find size(int v, int p = -1) {
    if (vis[v]) return 0;
    sz[v] = 1;
    for (int x: edges[v]) {
      if (x != p)
        sz[v] += find size(x, v);
    return sz[v];
 int find centroid(int v, int p, int n) {
    for (int x: edges[v]) {
      if (x != p) {
        if (!vis[x] && sz[x] > n / 2) {
          return find centroid(x, v, n);
    return v;
  void init centroid(int v = 0, int p = -1) {
    find size(v):
    \overline{\mathsf{int}} \ \overline{\mathsf{c}} = \mathsf{find} \ \mathsf{centroid}(\mathsf{v}, -1, \, \mathsf{sz}[\mathsf{v}]);
    vis[c] = true:
    par[c] = p;
    for (int x: edges[c]) {
      if (!vis[x]) {
        init centroid(x, c):
    }
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Graphs/DMST.cpp
struct RollbackUF {
 vi e;
 vpii st;
 RollbackUF(int n) : e(n. -1) {}
 int size(int x) { return -e[find(x)]; }
 int find(int x) { return e[x] < 0 ? x : find(e[x]); }
 int time() { return siz(st); }
 void rollback(int t) {
   for (int i = time(): i-- > t:) e[st[i].first] = st[i].nd:
   st.resize(t):
 bool unite(int a. int b) {
   a = find(a), b = find(b):
   if (a == b) return false;
   if (e[a] > e[b]) swap(a, b);
   st.pb({a. e[a]}):
   st.pb({b, e[b]});
   e[a] += e[b];
   e[b] = a;
   return true:
struct Edge {
 int a. b:
 ll w:
struct Node { /// lazy skew heap node
 Edae kev:
 Node *l. *r:
 ll delta;
 void prop() {
   kev.w += delta:
   if (l) l->delta += delta:
   if (r) r->delta += delta;
   delta = 0;
 Edge top() {
   prop();
   return key;
Node* merge(Node* a, Node* b) {
 if (!a | | !b) return a ?: b;
 a->prop(), b->prop():
 if (a->key.w > b->key.w) swap(a, b);
 swap(a->l, (a->r = merge(b, a->r)));
 return a:
void pop(Node*& a) {
 a->prop();
 a = merge(a->l, a->r);
} // wierzcholki numeruiemv od 0. r to korzen dmst
pair<ll, vi> dmst(int n, int r, vector<Edge>& g) {
 RollbackUF uf(n):
 vector<Node*> heap(n);
 for (Edge e : q) heap[e,b] = merge(heap[e,b], new Node{e});
 11 res = 0:
 vector<int> seen(n, -1), path(n), par(n);
 seen[r] = r;
 vector<Edge> O(n), in(n, \{-1, -1\}), comp;
 deque<tuple<int, int, vector<Edge>>> cycs;
 for (int s = 0; s < n; ++s) {
   int u = s, qi = 0, w;
```

```
while (seen[u] < 0) {
     if (!heap[u]) return {-1, {}};
     Edge e = heap[u]->top();
     heap[u]->delta -= e.w, pop(heap[u]);
     Q[qi] = e, path[qi++] = u, seen[u] = s;
     res += e.w. u = uf.find(e.a):
     if (seen[u] == s) { /// found cycle, contract
       Node* cvc = 0:
       int end = gi, time = uf.time();
       do cyc = merge(cyc, heap[w = path[--qi]]);
       while (uf.unite(u, w));
       u = uf.find(u), heap[u] = cyc, seen[u] = -1;
       cycs.push front(\{u, \text{ time, } \{\&Q[qi], \&Q[end]\}\}\);
    for (int i = 0; i < qi; ++i) in[uf.find(Q[i].b)] = Q[i];
 for (auto& [u, t, comp] : cycs) { // restore sol (optional)
   uf.rollback(t):
   Edge inEdge = in[u];
   for (auto\& e : comp) in[uf.find(e.b)] = e;
   in[uf.find(inEdge.b)] = inEdge:
 for (int i = 0; i < n; ++i) par[i] = in[i].a;
 return {res, par};
                                  Graphs/Dominatory.cpp
struct Dominators {
 int n orig, n;
 vector<int> parent, semi, vertex, dom, ancestor, label:
 vector<vector<int>> succ. pred. bucket:
 Dominators(int n): n orig(n), n(n + 1), parent(n), semi(n), vertex(n), dom(n), ances
tor(n), label(n), succ(n), pred(n), bucket(n) {
   n = n \text{ oria}:
 void add edge(int a, int b){
   a++; b++; succ[a].push back(b);
 void COMPRESS(int v) {
   if (ancestor[ancestor[v]] != 0) {
     COMPRESS(ancestor[v]);
     if (semi[label[ancestor[v]]] < semi[label[v]]) label[v] = label[ancestor[v]];</pre>
     ancestor[v]=ancestor[ancestor[v]]; } }
 void LINK(int v, int w) { ancestor[w]=v; }
 int EVAL(int v) {
   if(ancestor[v] == 0) return v:
   else {
     COMPRESS(v);
     return label[v];} }
 void DFS(int v) {
   semi[v] = ++n; vertex[n] = v;
    for(auto ng : succ[v]) {
     if(semi[nq] == 0) {
       parent[ngl=v:
       DFS(na):
     pred[ng].push back(v); } }
 // Zwraca vector dominatorow (-1 dla 0). // Wszystko numerowane od 0,
 vector<int> doit() { // wszystko musi być osiągalne z 0.
   iota(label.begin(), label.end(), 0);
   n = 0; DFS(1);
    for (int i = n; i >= 2; --i) {
     int w = vertex[i]:
     for (auto ng : pred[w]) {
       int u = EVAL(ng):
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if (semi[u] < semi[w]) { semi[w] = semi[u]; }</pre>
     bucket[vertex[semi[w]]].push back(w);
     LINK(parent[w],w);
     while (!bucket[parent[w]].empty()) {
       int v = bucket[parent[w]].back();
       bucket[parent[w]].pop back();
       int u = EVAL(v):
       if (semi[u] < semi[v]) dom[v] = u;</pre>
       else dom[v] = parent[w]:
   for (int i = 2; i \le n; ++i) {
     int w = vertex[i]:
     if (dom[w] != vertex[semi[w]]) { dom[w] = dom[dom[w]]: }
   dom[1] = 0;
   vector<int> res(n orig);
   for (int i = 0; i < n orig; ++i) res[i] = dom[i + 1] - 1;
   return res; } };
                                      Graphs/HLD.cpp
vi drz[nax]:
int roz[nax], jump[nax], pre[nax], post[nax], fad[nax], czas;
void dfs roz(int v) {
 roz[v]=1;
                                   // drz[] ma nie zawierać krawędzi
 for (int &i : drz[v]) {
                                   // do oica.
   fad[i]=v:
                                 // Init:
   dfs roz(i);
                                 // dfs roz(root);
   roz[v]+=roz[i];
                                 // dfs pre(root);
   if (roz[i]>roz[drz[v][0]]) // Użvcie get path(v, u) zwróci
     swap(i, drz[v][0]);
                               // przedziały odpowiadające ścieżce
 }
                                   // z v do u. Przedziały odpowiadające
                                     // ścieżce z v do lca maja
                                     // first>=second, zaś te dla ścieżki
void dfs pre(int v)
                                     // z lca do u maja first<=second.</pre>
 if (!jump[v])
                                   // Przedziały sa po kolei.
   jump[v]=v;
                                 // Lca występuje w nich dwa razy,
 pre[v]=(++czas):
                                   // naipierw iako second.
 if (!drz[v].empty())
                                   // a zaraz potem jako first.
   jump[drz[v][0]]=jump[v];
 for (int i : drz[v])
   dfs pre(i);
 post[v]=czas;
int lca(int v, int u) {
 while(iump[v]!=iump[u]) {
   if (pre[v]<pre[u])</pre>
     swap(v, u);
   v=fad[jump[v]];
 return (pre[v]<pre[u] ? v : u);</pre>
vpii path up(int v, int u) {
 vpii ret:
 while(jump[v]!=jump[u]) {
   ret.pb({pre[jump[v]], pre[v]});
   v=fad[jump[v]];
 ret.pb({pre[u], pre[v]});
 return ret;
vpii get path(int v. int u) {
 int w=lca(v, u);
 auto ret=path up(v, w);
 auto pom=path up(u, w);
```

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for (auto &i : ret)
    swap(i.st, i.nd):
  while(!pom empty()) {
   ret.pb(pom.back());
   pom.pop back();
 return ret:
                                   Graphs/Link Cut.cpp
struct Splay {
 Splay *l = 0, *r = 0, *p = 0;
 bool flip = false; // Wywal jak nie używasz make root.
 void update() {
   //update anything stored for nodes
   assert(!flip and (!l or !l->flip) and (!r or !r->flip)); }
  void touch() { //Do any lazy prop here
   if (flip) {
     swap(l. r):
     if (l) l->flip = !l->flip;
     if (r) r->flip = !r->flip;
     flip = false:
 bool sroot() { return !p or (p->l != this and p->r != this); }
 void connect(Splay* c, bool left) { (left ? l : r) = c; if (c) c->p = this; }
  void rotate() {
   Splay* f = p; Splay* t = f->p;
   const bool isr = f->sroot();
    const bool left = (this == f->l);
    f->connect(left ? r : l, left):
    connect(f, !left):
    if (isr) p = t;
    else t->connect(this, f == t->l);
    f->update():
 void push(){sroot()?touch():p->push();if (l) l->touch();if (r) r->touch();}
 void splay() {
   push():
   while (!sroot()) {
     Splay* x = p -> p;
     if ((p->sroot())) (((p->l == this) == (x->l == p)) ? p : this)->rotate();
     rotate():
   update();
 Splay* expose(){//v bedzie korzeniem splaya zawierającego ścieżke do korzenia
   Splay *q = this, *x = 0;//prawe dziecko będzie nullem. Jak zejdziemy w dół,
   while (g) { // to potem trzeba zrobić splay().
     q->splay();// LCA(u, v): u->expose(); return v->expose();
     a -> r = x: a -> update():
     x = q; q = q->p;
    splay();
    return x:
 Splay* root() { // Zwraca roota drzewowego (nie splejowego!).
   expose():
   Splav* s = this:
   while (s->touch(), s->l) s = s->l;
   s->splay();
    return s;
 void cut() { // Usuwa krawedź znajdującą się nad danym wierzchołkiem.
    expose(); assert(l /* Nie jest rootem. */);
    Splay* s = l;
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while (s->touch(), s->r) s = s->r;
   s->splay(): s->r->p = 0: s->r = 0:
 void link(Splay* to) {
   expose(); assert(!l /* Jest rootem. */);
   p = to;
 // Sprawia, że wierzchołek jest rootem w logicznym i w splayowym drzewie.
 void make root() { expose(); flip = !flip; touch(); }
                                      Graphs/Low.cpp
struct Low { // dziala dla multikrawedzi, petli, niespojnego grafu
 int n, cnt, edges;
 vector<vpii> G:
 vi low, pre, par, par nr, used, most, root, vis s, add s;
 Low(int n) : n(n), edges(0) {
   G.resize(n + 1);
   low.resize(n + 1):
   pre.resize(n + 1):
   par.resize(n + 1);
   par nr.resize(n + 1);
   root.resize(n + 1):
 void edge(int a, int b) {
   ++edges;
   G[a].eb(b. edges):
   G[b].eb(a, edges);
 void dfs(int v) {
   pre[v] = ++cnt:
   low[v] = pre[v];
   for (auto it : G[v]) {
     int u = it.st;
     int nr = it.nd:
     if (used[nr]) continue;
     used[nr] = 1;
     if (!pre[u]) {
       par[u] = v:
       par nr[u] = nr;
       dfs(u);
       mini(low[v], low[u]);
       mini(low[v], pre[u]);
   }
 }
 void qo() { // trzeba wywolac na poczatku!
   used.resize(edges + 1):
   for(int i = 1; i <= n; ++i) if (!pre[i]) root[i] = 1, dfs(i);</pre>
 vpii mostv() {
   most.resize(edges + 1);
   vpii ans;
   for(int i = 1; i \le n; ++i) {
     if (!root[i] && low[i] == pre[i]) {
       ans.eb(i, par[i]);
       most[par nr[i]] = 1;
   return ans;
 vi pkt art() { // tylko jesli potrzebujemy
   vi ans. take(n + 1), root sons(n + 1):
   for(int i = 1; i \le n; ++\overline{i}) {
     if (par[i] && root[par[i]]) {
       ++root sons[par[i]];
```

```
for(int i = 1: i <= n: ++i) {
     if (root[i] \&\& root sons[i] >= 2) take[i] = 1;
     if (!root[i] \&\& !root[par[i]] \&\& low[i] >= pre[par[i]]) take[par[i]] = 1;
    for(int i = 1; i <= n; ++i) if (take[i]) ans.pb(i);</pre>
    return ans;
 // kod nizei tvlko iesli potrzebujemy 2spojnych
  using comps = vector<vpii>:
  void dfs s(int v, vpii &moja) {
   vis s[v] = 1;
    for (auto it : G[v]) {
     int u = it.st:
     int nr = it.nd;
     if (most[nr]) continue;
     if (!add s[nrl) {
       moia.e\overline{b}(v, u);
       add s[nr] = 1;
     if (!vis s[u]) dfs s(u, moja);
 }
 comps bico() {
   vis s.resize(n + 1):
   add s.resize(edges + 1);
   comps ans;
    for (auto it : mosty()) ans.pb({it});
    for(int i = 1: i \le n: ++i) {
     if (!vis s[i]) {
       vpii curr;
       dfs s(i, curr);
       if (!curr.emptv()) ans.pb(curr):
    return ans;
 }
                                Graphs/Planarity Check.cpp
#define PB push back
#define SZ(x) ((int)(x).size())
#define FOR(i.a.b) for (int i = (a); i \le (b); ++i)
#define VI vector<int>
                                     //numeracia od zera
struct FU {
                                     //nie dawać multikrawędzi ani petli
 VI p: vector<bool> b:
                                     //nie odpalać IsPlanar() ponad raz
 FU() {}
 FU(int n) : p(n), b(n) { iota(p.begin(), p.end(), 0); }
 pair<int, bool> Find(int v) {
   if (p[v] == v) { return {v. 0}: }
   auto res = Find(p[v]); res.second ^= b[v];
   p[v] = res.first; b[v] = res.second; return res;
 bool Union(int x, int v, bool flip) {
   bool xb, yb; tie(x, xb) = Find(x); tie(y, yb) = Find(y);
   if (x == y) { return !(xb ^ yb ^ flip); }
   p[x] = y; b[x] = xb ^ yb ^ flip; return true;
struct PlanarityTest {
 int N, M, tm;
 vector<VI> adi. dn. up: vector<pair<int.int>> e up:
 vector<bool> vis; VI low, pre;
 FU fu;
 void DfsLow(int v, int p) {
```

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M M

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}
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```
vector<vpii> d; vector<vi> d2;
  const int inf = 1e9:
  SPFA(int n):n(n+1){
    odl.resize(n, inf); oj.resize(n); czok.resize(n);
    d.resize(n); d2.resize(n);
  vi cvkl: int root:
  bool us(int nr){
   if(nr == root) return 1;
    czok[nr] = 0:
    for(int ak:d2[nrl){
      if(oj[ak] == nr){
        if(us(ak)){
          cvkl.pb(nr):
                                        //numeracia od zera
          return 1:
    d2[nrl.clear():
    return 0;
  bool licz sciezki(int s){ // false, adv z s da sie doisc do uiemnego cyklu
    vi st, st2; // znaleziony cykl jest w wektorze cykl
    odl[s] = 0; czok[s] = 1; st.pb(s);
    while(siz(st)){
      for(int i=0: i<siz(st): i++){</pre>
        int ak = st[i]:
        if(czok[ak]) for(pii x:d[ak]){
          int nei, cost; tie(nei, cost) = x;
          if(odl[ak] + cost < odl[nei]){</pre>
            root = ak:
            if(us(nei)){
              cykl.pb(ak); reverse(all(cykl));
              return 0:
            odl[nei] = odl[ak] + cost; oj[nei] = ak;
            d2[ak].pb(nei); czok[nei] = 1;
            st2.pb(nei):
      st.clear(): swap(st. st2):
    return 1;
 vi uiemnv cvkl(){
    for (int i=0; i< n-1; i++) add edge(n-1, i, 0);
    if (licz sciezki(n-1)){
      return {};
    } else {
      return cykl;
 void add edge(int a, int b, int cost){
    d[a].p\overline{b}(\{b, cost\});
}; // #undef int
                                       Graphs/SSS.cpp
const int nax=100*1007:
vector <int> graf[nax], farg[nax];
int ost[nax], bylo[nax], post[nax], spo[nax], counter, coudfs;
vector <vector<pair<int.int>>> mer:
void dfs1(int v) {
 if (bylo[v])
    return:
```

vis[v] = true; low[v] = pre[v] = tm++;

 $up[v].PB(SZ(e up)); e up.PB({v, s});$ 

dn[v].PB(s); DfsLow(s, v); low[v] = min(low[v], low[s]);

for (int s : adi[v]) {

**if** (!vis[s]) {

if (s == p) { continue; }

} else if (pre[s] < pre[v]) {</pre>

low[v] = min(low[v], pre[s]);

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bylo[v]=1;
 for (int i : graf[v]) dfs1(i);
 coudfs--:
 post[coudfs]=v;
void dfs2(int v. int s) {
 if (spo[v]>=0)
   return:
 spo[v]=s;
 for (int i : farg[v]) dfs2(i, s);
void rek(int l, int r, vector <pair<pair<int,int>,int>> &kra) {
 if (l>r) return;
 counter++:
 vector <int> ver:
 for (auto i : kra) {
   if (ost[i.first.first]<counter) {</pre>
     ver.push back(i.first.first):
     ost[i.first.first]=counter:
   if (ost[i.first.second]<counter) {</pre>
     ver.push back(i.first.second):
     ost[i.first.second]=counter:
 for (int i : ver) {
   bvlo[i]=0: spo[i]=-1:
   graf[i].clear(); farg[i].clear();
 int s=(l+r)>>1:
 for (auto i : kra) {
   if (i.second<=s) {</pre>
     graf[i.first.first].push back(i.first.second);
      farg[i.first.second].push back(i.first.first):
 }
 coudfs=ver.size();
 for (int i : ver)
   dfs1(i):
 for (int i=0; i<(int)ver.size(); i++)</pre>
   dfs2(post[i], post[i]);
 for (int i : ver)
   if (i!=spo[i])
     mer[s].push back({i, spo[i]});
 vector <pair<pair<int,int>,int>> lew, pra;
 for (auto i : kra) {
   if (spo[i.first.first] == spo[i.first.second])
     lew.push back(i);
   else
     pra.push back({{spo[i.first.first].spo[i.first.second]}. i.second}):
 rek(l, s-1, lew); rek(s+1, r, pra);
void sss(vector <pair<int.int>> kra) {
 mer.clear():
 mer.resize(kra.size());
 vector <pair<int,int>,int>>daj;
 for (int i=0; i<(int)kra.size(); i++) {</pre>
   daj.push back({kra[i], i});
   ost[kra[i].first]=-1;
   ost[kra[i].second]=-1;
 counter=0:
 rek(0, (int)kra.size()-1, daj);
```

```
Graphs/Vizing.cpp
#define int ll//ieśli lona lonai potrzebne
struct Vizina {
 vector<vector<pair<int,int>>> adj;
 map<pair<int,int>, int> edges;
 vector<int> edge colors:
 vector<vector<int>>> color to;
 vector<vector<int>>> color queue;
 vector<int> unused color;
 int N. M. K:
 Vizing(int size) : adj(size), N(size), M(0) {}
 void AddEdge(int u, int v) {
   adj[u].emplace back(u, M);
   adi[v].emplace back(v, M):
   edges[make pair(min(u, v), max(u, v))] = M;
   ++M;
 }
                                        //numeracia od zera
                                        //multikrawedzie niebezpieczne
 int MaxDegree() const {
                                        //ogólnie nietestowane
   int answer = 0;
   for (int i = 0; i < N; ++i)
     answer=max(answer, (int)adi[i].size());
   return answer:
 void FindFree(int v) {
   while (!color queue[v].emptv()) {
     const int c = color queue[v].back();
     if (color to[v][c] == -1) {
       unused color[v] = c;
       return:
     } else {
       color queue[v].pop back();
   assert(false):
 void FlipPath(int from, int a, int b) {
   const int to = color to[from][a]:
   color queue[from].push back(a);
   color queue[from].push back(b);
   color to[from][b] = -1;
   if (to == -1) { return: }
   FlipPath(to, b, a);
   color to[from][b] = to;
   color to[to][b] = from;
   color to [from][a] = -1:
   FindFree(from):
   FindFree(to);
 void ColorEdge(int x, int v) {
   for (int col = 0; col <= K; ++col)</pre>
     if (color to[x][col] == -1 \&\& color to[y][col] == -1) {
       color to[x][col] = y;
       color to[y][col] = x;
       FindFree(x); FindFree(y);
       return;
   vector<int> fan{v}, fan colors{-1};
   vector<bool> has on(N);
   const int c = unused color[x];
   int d:
   while (true) {
     d = unused color[fan.back()];
     if (color to[x][d] == -1 || d == c) { break; }
     if (has on[d]) {
```

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FlipPath(x, d, c);
       FindFree(x):
       return ColorEdge(x, v):
     fan.push back(color to[x][d]);
     fan colors.push back(d):
     has on[d] = true:
   fan colors.push back(-1);
   for (int i = 0; i < (int)fan.size(); ++i) {
     const int from = fan colors[i], to = fan colors[i + 1];
     if (from != -1) {
       color gueue[fan[i]].push back(from);
       color queue(x).push back(from):
       color to[fan[i]][from] = -1;
     if (to != -1) {
       color to[fan[i]][to] = x:
       color to[x][to] = fan[i];
   assert(color to[x][d] == -1):
   color to[x][\overline{d}] = fan.back();
   color to[fan.back()][d] = x;
   for (int v : fan) { FindFree(v); }
   FindFree(x):
 vector<int> ColorGraph() {
   K = MaxDegree();
   edge colors = vector<int>(M, -1):
   color to = vector<vector<int>>(N, vector<int>(K + 1, -1));
   vector<int> all colors(K + 1);
   iota(all colors.begin(), all colors.end(), 0);
   color queue = vector<vector<int>>>(N, all colors);
   unused color = vector<int>(N);
   for (auto &edge : edges)
     ColorEdge(edge.first.first, edge.first.second);
   for (int v = 0: v < N: ++v)
     for (int c = 0; c \le K; ++c)
       if (color to[v][c] > v) {
         assert(color to[color to[v][c]][c] == v);
         edge colors[edges[make pair(v, color to[v][c])]] = c;
   return edge colors;
 }
#undef int
                                    Math/BitoweFFT.cpp
void xor fft(ll* p, int n, bool inv) {
 for (int len = 1; 2 * len <= n; len <<= 1) {
   for (int i = 0; i < n; i += 2 * len) {
     for (int j = 0; j < len; ++j) {
       ll u = p[i + j];
       ll v = p[i + len + i]:
       p[i + j] = u + v;
       p[i + len + j] = u - v;
 } } }
 if (inv) {
   for (int i = 0: i < n: ++i) {
     p[i] /= n; // uwaga qdy liczymy modulo!!
} } }
void and fft(ll* p, int n, bool inv) {
 for (int len = 1; 2 * len <= n; len <<= 1) {
   for (int i = 0; i < n; i += 2 * len) {
     for (int j = 0; j < len; ++j) {
```

```
ll u = p[i + j];
        ll v = p[i + len + i]:
        if (!inv) {
          p[i + j] = v;
          p[i + len + j] = u + v;
        } else {
          p[i + j] = -u + v;
          p[i + len + j] = u;
void or fft(ll* p. int n. bool inv) {
 for (int len = 1, 2 * len <= n, len <<= 1) {
    for (int i = 0; i < n; i += 2 * len) {</pre>
      for (int j = 0; j < len; ++j) {
        ll u = p[i + i]:
        ll v = p[i + len + j];
        if (!inv) {
          p[i + j] = u + v;
          p[i + len + i] = u:
        } else {
          p[i + j] = v;
          p[i + len + j] = u - v;
Math/FFT.cpp
#define REP(i,n) for(int i = 0; i < int(n); ++i)
/*Precision error max ans/le15 (2.5e18) for (long) doubles.
So integer rounding works for doubles with answers 0.5*1e15.
e.g. for sizes 2^20 and RANDOM positive integers up to 45k.
Those values assume DBL MANT DIG=53 and LDBL MANT DIG=64.
For input in [0, M], you can decrease everything by M/2.
If there are many small vectors, uncomment "BRUTE FORCE".*/
typedef double ld: // 'long double' is 2.2 times slower
struct C {
 ld real, imag;
 C operator * (const C & he) const {
    return C{real * he.real - imag * he.imag. real * he.imag + imag * he.real};
  void operator += (const C & he) {real += he.real; imag += he.imag;} };
void dft(vector<C> & a, bool rev) {
  const int n = a.size():
  for(int i = 1, k = 0; i < n; ++i) {
    for(int bit = n / 2; (k ^= bit) < bit; bit /= 2);;;</pre>
    if(i < k) swap(a[i], a[k]):
  for(int len = 1, who = 0; len < n; len *= 2, ++who) {
    static vector<C> t[30];
    vector<C> & om = t[who]:
    if(om.emptv()) {
      om.resize(len):
      const ld ang = 2 * acosl(0) / len;
      REP(i, len) om[i] = i%2 || !who?C{cos(i*ang), sin(i*ang)}:t[who-1][i/2];
    for(int i = 0; i < n; i += 2 * len)
     REP(k, len) {
         const C x = a[i+k], y = a[i+k+len]
            * C{om[k].real, om[k].imag * (rev ? -1 : 1)};
        a[i+k] += y;
        a[i+k+len] = C{x.real - y.real, x.imag - y.imag};
 if(rev) REP(i, n) a[i].real /= n;
template<typename T>vector<T> multiply(const vector<T> &a. const vector<T> &b. bool split
= false) {
 if(a.empty() || b.empty()) return {};
  int n = a.size() + b.size();
```

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return y->m == z->m \&\& y->b <= z->b;
    auto x = prev(v):
    if (z == end()) return y->m == x->m \&\& y->b <= x->b;
    return (x-b-y-b) * (z-m-y-m) >= (y-b-z-b) * (y-m-x-m);//ld?
  void insert line(ll m, ll b) {
    auto y = insert({m, b});
    y->succ = [=] { return next(y) == end() ? 0 : \delta*next(y); };
    if (bad(v)) {
      erase(y);
      return;
    while (\text{next}(v) != \text{end}() \&\& \text{bad}(\text{next}(v))) = \text{rase}(\text{next}(v));
    while (y \mid = begin() \&\& bad(prev(y))) erase(prev(y));
 ll eval(ll x) {
    auto l = *lower bound((Line) {x, is_query});
    return l.m * x + l.b:
 }
                                      Math/Massev.cpp
const int mod = 1e9 + 7; // const ll mod = (ll)1e18 + 3;
template <class c> void add self(c \& a, c b) { a += b; if(a >= mod) a -= mod; }
template <class c> void sub_self(c & a, c b) { a -= b; if(a < 0) a += mod; }
int mul(int a, int b) { return (ll) a * b % mod: }
ll mul(ll a, ll b) {return (__int128) a * b % mod; }
template <class c> c my pow(\overline{c} a, c b) \{/*...*/\}
template <class c> c my inv(c a) { return my pow<c>(a, mod - 2); }
template <class c> c negative (c a) {return (mod - a) % mod;}
template <class F> struct Massey {
 vector<F> start, coef; // 3 optional lines
 vector<vector<F>>> powers:
 F memo inv:
 // Start here and write the next ~25 lines until "STOP"
  int L; // L == coef.size() <= start.size()</pre>
 Massey(vector<F> in) { // O(N^2)
   L = 0:
    const int N = in.size():
    vector<F> C{1}, B{1};
    for(int n = 0; n < N; ++n) {
      // assert(0 <= in[n] && in[n] < mod): // invalid input
      B.insert(B.begin(), 0):
      F d = 0:
      for(int i = 0; i \le L; ++i)
        add self(d, mul(C[i], in[n-i])):
      if(d == 0) continue:
      vector<F> T = C;
      C.resize(max(B.size(), C.size()));
      for(int i = 0; i < (int) B.size(); ++i)
        sub self(C[i], mul(d, B[i]));
      if (2 \times L \le n) {
       L = n + 1 - L;
        B = T:
        d = my inv(d);
        for (F \& x : B) x = mul(x, d);
    cerr << "L = " << L << "\n":
    assert(2 * L <= N - 2); // NO RELATION FOUND :(
    // === STOP ===
    for(int i = 1: i < (int) C.size(): ++i)
      coef.push back(negative(C[i]));
```

assert((int) coef.size() == L);

for(int i = 0; i < L; ++i)

```
vector<T> ans(n - 1);
 /* if(min(a.size(),b.size()) < 190) { // BRUTE FORCE</pre>
   REP(i, a.size()) REP(i, b.size()) ans[i+i] += a[i]*b[i];
   return ans: } */
 while (n\&(n-1)) ++n;
 auto speed = [&](const vector<C> & w, int i, int k) {
   int j = i ? n - i : 0, r = k ? -1 : 1;
   return C{w[i].real + w[j].real * r, w[i].imag
        - w[i].imag * r * (k ? C{0, -0.5} : C{0.5, 0});
 if(!split) { // standard fast version
   vector<C> in(n), done(n);
   REP(i, a.size()) in[i].real = a[i];
   REP(i, b.size()) in[i].imag = b[i]:
   dft(in. false):
   REP(i, n) done[i] = speed(in, i, \theta) * speed(in, i, 1);
   dft(done, true);
   REP(i, ans.size()) ans[i] = is integral<T>::value ?
        llround(done[i].real) : done[i].real;
 else {
   const int M = 1 << 15:</pre>
   vector <C> t[2]:
   for (int x = 0; x < 2; ++x) {
     t[x].resize(n);
     const vector \langle T \rangle \& in = (x ? b : a):
     for (int i = 0: i < (int) in.size(): ++i)
       t[x][i] = C\{ld(in[i] % M), ld(in[i] / M)\};
     dft(t[x], false);
   vector \langle C \rangle d1(n), d2(n);
   for (int i = 0; i < n; ++i) {
     d1[i] += speed(t[0], i, 0) * speed(t[1], i, 0);
     d1[i] += speed(t[0], i, 1) * speed(t[1], i, 1) * C(0, 1);
     d2[i] += speed(t[0], i, 0) * speed(t[1], i, 1);
     d2[i] += speed(t[0], i, 1) * speed(t[1], i, 0);
   dft(d1. true):
   dft(d2, true):
   for (int i = 0; i < n; ++i)
     dl[i].imag /= n;
   for (int i = 0; i < (int) ans.size(); ++i)
     ans[i] = (llround(d1[i].real) + llround(d2[i].real) % mod * M + llround(d1[i].imag).
 % mod * (M * M)) % mod;
 return ans:
                           Math/Linear Function Convex Hull.cpp
const ll is query = -(1LL \ll 62);
struct Line {
 ll m. b:
 mutable function<const Line *()> succ;
 bool operator<(const Line &rhs) const {</pre>
   if (rhs.b != is query) return m < rhs.m;</pre>
   const Line *s = succ();
   if (!s) return 0;
   ll x = rhs.m:
   return b - s->b < (s->m - m) * x://ld?
struct HullDynamic : public multiset<Line> {
 bool bad(iterator v) {
   auto z = next(v):
   if (y == begin()) {
     if (z == end()) return 0;
```

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start.push back(in[i]);
   while(!coef.empty() && coef.back() == 0) { coef.pop back(); --L; }
   if(!coef.empty()) memo inv = my inv(coef.back());
   powers.push back(coef);
 vector<F> mul cut(vector<F> a, vector<F> b) {
   vector<F> r(2 * L - 1):
   for(int i = 0; i < L; ++i)
     for(int j = 0; j < L; ++j)
       add self(r[i+i], mul(a[i], b[i])):
   while((int) r.size() > L) {
     F value = mul(r.back(), memo inv); // div(r.back(), coef.back());
     const int X = r.size();
     add self(r[X-L-1], value):
     for(int i = 0; i < L; ++i)
       sub self(r[X-L+i], mul(value, coef[i]));
     assert(r.back() == 0);
     r.pop back();
   return r;
 F get(ll k) \{ // O(L^2 * log(k)) \}
   if(k < (int) start.size()) return start[k]:</pre>
   if(L == 0) return 0;
   k -= start.size();
   vector<F> vec = coef:
   for(int i = 0: (1LL << i) <= k: ++i) {
     if(i == (int) powers.size())
       powers.push back(mul cut(powers.back(), powers.back()));
     if(k & (1LL << i))
       vec = mul cut(vec. powers[i]);
   F total = 0;
   for(int i = 0: i < L: ++i)
     add self(total, mul(vec[i], start[(int)start.size()-1-i]));
   return total; } };
                              Math/Multipoint Evaluation.cpp
ll dziel(ll a, ll b) {...}
vll inv(vll p, int n) {
 assert(p[0]);
 ll odw=dziel(1, p[0]);
 for (ll \&i : p) i=(i*odw)%mod:
 vll a{1}:
 for (int i=1; i<n; i<<=1) {</pre>
   vll r=multiply(vll(p.begin(), p.begin()+min(2*i, (int)p.size())), q, true);
   r.resize(2*i):
   r.erase(r.begin(), r.begin()+i);
   for (ll &j : r) j=(mod-j)%mod;
   r=multiply(r, q, true);
   q.resize(2*i):
   for (int j=0; j<min(i, (int)r.size()); j++) q[i+j]=r[j];</pre>
 q.resize(n);
 while(!g.emptv() && !g.back()) g.pop back():
 for (ll &i : p) i=(i*odw)%mod;
 return q;
vll div(vll a. vll b) {
 int s=(int)a.size()-(int)b.size()+1;
 if (s<=0) return {};
 reverse(a.begin(), a.end()); reverse(b.begin(), b.end());
 b=inv(b, s):
 a=multiply(a, b, true);
 a.resize(s);
 reverse(a.begin(), a.end());
```

```
return a;
vll rem(vll a, vll b) {
 vll d=div(a, b);
  d=multiply(d, b, true);
 for (int i=0: i<(int)d.size(): i++) a[i]=(a[i]-d[i]+mod)%mod:</pre>
  while(!a.empty() && !a.back()) a.pop back();
 return a;
void precalc(int v. vector <vll> &help. vll querv) {
 if ((int)help.size()<=v) help.resize(v+1);</pre>
 int s=query.size();
 if (s==1) {
   help[v]={(mod-querv[0])}mod. 1}:
    return:
 vll a, b;
  for (int i=0; i<s/2; i++) a.push back(query[i]);</pre>
  for (int i=s/2: i<s: i++) b.push back(querv[i]);</pre>
 precalc(v^*2, help, a); precalc(v^{\overline{*}}2+1, help, b);
 help[v]=multiply(help[v*2], help[v*2+1], true);
void calc(int v, vll &res, vector <vll> &help, vll wek) {
 wek=rem(wek, help[v]);
 if ((int)help[v].size()==2) {
    res.push back(wek.emptv() ? 0 : wek[0]):
    return:
 calc(v*2, res, help, wek);
 calc(v*2+1, res. help, wek):
vll multi eva(vll wek, vll query) {
 vector <vll> help;
 precalc(1, help, query);
  vll res:
  calc(1, res, help, wek);
  return res;
                                      Math/Simplex.cpp
struct Simplex {
                        // Maximize c*x subject to Ax <= b.
                        // Initialize the structure, set A, b, c and then run
  using T = double;
  vector<vector<T>> A: // solve(). Max objective is stored in res. To recover
  vector<T> b. c:
                        // the best result, use getVars().
  int V, E;
  vector<int> eqIds, varIds, cols;
  static constexpr T kEps = 1e-9:
  Simplex(int vars, int eqs) : A(eqs, vector<T>(vars)), b(eqs), c(vars),
     V(vars), E(egs), egIds(egs), varIds(vars), res(0) {
    iota(varIds.begin(), varIds.end(), 0);
    iota(egIds.begin(), egIds.end(), vars);
  void pivot(int eq, int var) {
   T coef = 1 / A[eq][var];
    cols.clear():
    for (int i = 0; i < V; i++) {
      if (abs(A[eq][i]) > kEps) { cols.push back(i); A[eq][i] *= coef; }
    A[eq][var] *= coef: b[eq] *= coef:
    for (int row = 0; row < E; row++) {
     if (row == eq || abs(A[row][var]) < kEps) { continue; }</pre>
     T k = -A[row][var]:
      A[row][var] = 0:
      for (int i : cols) { A[row][i] += k * A[eq][i]; }
      b[row] += k * b[ea]:
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T q = c[var]: c[var] = 0:
   for (int i : cols) { c[i] -= q * A[eq][i]; }
   res += q * b[eq];
   swap(varIds[var], eqIds[eq]);
 bool solve() {
   while (true) {
     int eq = -1, var = -1;
     for (int i = 0: i < E: i++) { if (b[i] < -kEps) { eq = i: break: } }
     if (eq == -1) { break; }
     for (int i = 0; i < V; i++) { if (A[eq][i] < -kEps) { var = i; break; } }
     if (var == -1) { res = -1e9; return false; /* No solution */ }
     pivot(eq. var):
   while (true) {
     int var = -1, eq = -1;
     for (int i = 0: i < V: i++) { if (c[i] > kEps) { var = i: break: } }
     if (var == -1) { break; }
     for (int i = 0; i < E; i++) {
       if (A[i][var] < kEps) { continue; }</pre>
       if (eq >= 0 \&\& b[i] / A[i][var] >= b[eq] / A[eq][var]) { continue: }
       eq = i;
     if (eg == -1) { res = 1e9; return false; /* Unbounded */ }
     pivot(eq. var):
   return true;
 vector<T> getVars() { // Optimal assignment of variables.
   vector<T> result(V);
   for (int i = 0; i < E; i++) if (eqIds[i] < V) result[eqIds[i]] = b[i];
   return result;
 }
                                     Math/Simpson.cpp
// Either run integral(A, B) once or split the interval [A, B] into up to ~1000
\prime/ smaller intervals -- if the function f behaves oddly or the interval is long.
ld simp(ld low, ld high, const ld * old, vector<ld> & nowe) {
 const int n = 500; // n must be even!!! Try n = 2 and n = 10.
 nowe.resize(n + 1);
 ld total = 0. iump = (high - low) / n:
 for(int i = 0; i \le n; ++i) {
   int mul = i == 0 || i == n ? 1 : 2 + i % 2 * 2; // 1 4 2 4 2 ... 4 1
   nowe[i] = !old || i % 2 ? f(low + i * jump) : old[i/2];
   total += nowe[i] * mul: // uses a global function ld f(ld x)
 return total * (high - low) / n / 3;
ld rec(ld low, ld high, ld prv, const vector<ld> & old) {
 ld mid = (low + high) / 2;
 vector<ld> left, right;
 ld L = simp(low, mid, old.data(), left);
 ld R = simp(mid. high. old.data() + old.size() / 2. right);
 ld del = L + R - prv;
                                                   // try without del/15
 if(abs(del) < 15 * eps) return L + R + del / 15; // eps ~ required abs precision</pre>
 return rec(low, mid, L, left) + rec(mid, high, R, right);
ld integral(ld low, ld high) {
 vector<ld> old;
 ld prv = simp(low, high, 0, old);
 return rec(low, high, prv, old);
                                      Math/Wzory.tex
```

```
 \int \sqrt{x^2+1} \, dx = \frac{1}{2} \left( x \sqrt{x^2+1} + \operatorname{arcsinh} x \right) + c \quad \int \sqrt{1-x^2} \, dx = \frac{1}{2} \left( x \sqrt{1-x^2} + \operatorname{arcsin} x \right) + c   \int \frac{1}{ax^2+bx+c} \, dx = \frac{2}{\sqrt{4ac-b^2}} \operatorname{arctan} \frac{2ax+b}{\sqrt{4ac-b^2}} \quad (\Delta < 0)   \int \frac{x}{ax^2+bx+c} \, dx = \frac{1}{2a} \ln |ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}   \int \tan x \, dx = -\ln |\cos x| + c   (\operatorname{arcsin} x)' = \frac{1}{\sqrt{1-x^2}}, \qquad (\operatorname{arccos} x)' = -\frac{1}{\sqrt{1-x^2}}   \operatorname{tgamma}(\mathbf{t}) = \Gamma(t) = \int_0^\infty e^{t-1} e^{-x} \, dx   \frac{1}{\pi} = 0.31831, \quad \pi^2 = 9.86960, \quad \frac{1}{\pi^2} = 0.10132, \quad \frac{1}{e} = 0.36788, \quad \gamma = 0.577215664901532   H_n = \ln n + \gamma + \frac{1}{2n} - \frac{1}{12n^2} + O(n^{-4})   \ln n! = n \ln n - n + \frac{1}{2} \ln(2\pi n) + \frac{1}{12n} - \frac{1}{360n^3} + \frac{1}{1260n^5} - O(n^{-7})
```

Jeśli  $n = 2^{a_0} \cdot p_1^{2 \cdot a_1} \cdot \ldots \cdot p_r^{2 \cdot a_r} \cdot q_1^{b_1} \cdot \ldots \cdot q_s^{b_s}$  i  $B = \prod (b_i + 1)$ , gdzie  $p_i$  to liczby postaci  $4 \cdot k + 3$ , a  $q_i$  to liczby postaci  $4 \cdot k + 1$ , to liczba sposobów na zapisanie n jako sumy dwóch kwadratów liczb naturalnych wynosi:

- 0, jeśli któraś z liczb  $a_i$  nie jest całkowita
- $\frac{B}{2}$ , jeśli B jest parzyste
- $\frac{B-(-1)^{a_0}}{2}$ , jeśli B jest nieparzyste

$$S_j = \sum_{1 \le i_1 < i_2 < \dots < i_j \le n} |A_{i_1} \cap A_{i_2} \cap \dots A_{i_j}|$$

 $\sum_{j\geqslant k} \binom{j}{k} (-1)^{j+k} S_j$ - liczba elementów należących do dokładnie k zbiorów

 $\sum_{i \geq k} {j-1 \choose k-1} (-1)^{j+k} S_j$  - liczba elementów należących do co najmniej k zbiorów.

 $\sum\limits_{j\geqslant 1} (-2)^{j-1} S_j$  - liczba elementów należących do nieparzyście wiele zbiorów

Zaczynając w punkcie 0, idąc w każdym ruchu o -1 z prawdopodobieństwem p i o +1 z 1-p, kończąc w -A lub B prawdopodobieństwo dojścia do B wynosi

 $\frac{1+r+r^2+\ldots+r^{A-1}}{1+r+r^2+\ldots+r^{A+B-1}} = \frac{1-r^A}{1-r^{A+B}}^* = 1 - \frac{1+r'+r'^2+\ldots+r'^{B-1}}{1+r'+r'^2+\ldots+r'^{A+B-1}} = 1 - \frac{1-r'^B}{1-r'^{A+B}}^*,$ (\*działa o ile  $p \neq \frac{1}{2}$ ) gdzie  $r = \frac{p}{1-p}, r' = \frac{1}{r},$  dla  $p = \frac{1}{2}$  równe  $\frac{A}{A+B}$ 

Dla  $x_1, x_2, \ldots, x_n$  macierz  $a_{i,j} = x_i^{j-1}$  ma wyznacznik  $\prod_{i < j} (x_j - x_i)$  i odwrotność

$$b_{i,j} = (-1)^{i \frac{\sum_{A \subset [n] \setminus \{j\}, |A| = n - i} \prod_{k \in A} x_k}{\prod_{k \neq j} (x_k - x_j)}}$$

## Miscellaneous/De Bruijn.cpp

//Generuje słowo cykliczne, Działa dla alph >= 2
vi de\_bruijn(int len, int alph){
 vi res, lyn{0};

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while (lyn[0] != alph - 1){
   int r = siz(lvn):
   if (len % r == 0)
     for (int c : lyn)
        res.pb(c);
    for (int i = r: i <= len - 1: ++i)
     lyn.pb(lyn[i - r]);
   while (lyn.back() == alph - 1)
     lyn.pop back();
    lvn.back()++:
 res.pb(alph - 1);
 return res;
                                 Miscellaneous/Loops.cpp
sim> bool next inc(c beg, c end, int bou) \{ //[0, bou) \}
 for (c i = end; i != beq;) {
   int v = *--i:
   if (v + 1 < bou) {
     iota(i, end, v + 1);
     return true;
   bou = *i:
 }
 return false;
sim> bool next nondec(c beg, c end, int bou) {//[0, bou)
 for (c i = end; i != beq;) {
   int v = *--i;
   if (v + 1 < bou) {
     fill(i. end. v + 1):
     return true;
 return false;
sim> bool next any(c beg, c end, int bou) {
 if (!bou) return false:
 for (c i = end: i != bea:) {
   if (++*--i == bou) *i = 0;
   else return true;
 return false:
sim>bool next diff(c beg,c end,int bou){//starts with 0,1,2...
 return next permutation(beg. end) || next inc(beg. end. bou);
sim>bool next split(c beg, c end) {//starts with all 0s, ends with 0,1,2...
 vector <int> seen(distance(beg, end));
 for (c i = beg; i != end; ++i) seen[*i]++;
 for (c i = end; i != beg;) {
   if (seen[*--i] > 1) {
     (*i)++:
     fill(++i. end. 0):
     return true:
 }
 return false:
//for (int m = (1 << k) - 1; m < (1 << n); m = nux(m)) - k-podzbiory [0,n)
sim>inline c nux(c m){
 if (!m) return numeric limits<c>::max();
 c A = m \& -m:
 c B = \sim ((A - 1) ^ m);
 c C = B \& -B:
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c D = (C >> (1 + ctz(A))) - 1;
  return C | (m \& \sim (C - 1)) | D:
V/for (unt i = 0; i < (1u << n); i = rrh(i)) maska bez sasiednich zapalonych, rosnaco
sim> inline c rrh(c m) {
 c a = \sim (m \mid m >> 1):
 cb=a\&-a:
 return (m | b) &~ (b - 1);
V/for (unt i = (1u << (n + 1)) / 3; ~i; i = lyl(i)) maska bez sąsiednich, malejąco
sim> inline c lvl(c m) {
 if (!m) return -1;
 c x = m \& -m;
  return (m ^ x) | ((x << 1) / 3):
                               Number Theory/Cornacchia.cpp
                       // Typ, w którym mieści się modulo.
using T = int;
using T2 = long long; // Typ, w którym mieści się kwadrat modulo.
struct Pierwiastek \{//\ Dla\ pierwszego\ p > 2,\ szuka\ pierwiastków\ modulo\ p.
 Pierwiastek(T p) : s(0), p(p), m(p - 1) {//template <int p> jeśli stałe p
    assert(p > 2);
    while (m \% 2 == 0) \{ m /= 2; s++; \}
    c = 2:
    while (Pot(c, p / 2) == 1) c = rand() % (p - 1) + 1; // rand() musi być duży
 T Pot(T a. T pot) const {
   T r = 1:
    while (pot) {
     if (pot \& 1) r = r * (T2) a % p;
     a = a * (T2) a % p;
     pot >>= 1:
    return r;
    // Zwraca a**pot % p.
 T Licz(T a) const { // Znajduje pierwiastek z a modulo p.
   if (a == 0) return 0;
                                       // Sprawdza, czy a jest reszta
    if (Pot(a, p / 2) != 1) return -1; // kwadratowa. Jeśli nie, zwraca -1.
   Tz = Pot(c. m):
    T v = Pot(a. m / 2):
    T u = (T2) v * v % p;
    v = (T2) v * a % p;
    u = (T2) u * a % p:
    for (int i = s - 1; i >= 1; i --) {
      if (Pot(u, T(1) << (i - 1)) != 1) {
        u = (T2) u * z % p;
        u = (T2) u * z % p:
        v = (T2) v * z % p:
      z = (T2) z * z % p;
    return v; // Pierwiastkami liczby a są: {v, (p - v) % p}.
 private: int s;
   T p. m. c:
template<class num t>
inline num t isgrt(num t k) {
 num t r = sart(k) + 1;
 while (r * r > k) r--:
 return r;
pii cornacchia(int p. int d) \{//dla\ p\ pierwszego,\ znaiduie\ a,\ b:\ p=a^2+db^2\}
 int x0 = Pierwiastek(p) Licz(p - d);
 if (x0 == -1) return \{-1, -1\};
  x0 = max(x0, p - x0);
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614.889.782.588.491.410: 15 prime divisors, largest=47
  int a = p, b = x0;
  int l = isart(p):
  while (b > l) {
   int r = a \% b;
   a = b;
   b = r:
 int c = p - b * b;
 if (c % d) return {-1, -1},
 c /= d:
 int cc = isqrt(c);
 if (cc * cc != c) return {-1, -1};
 return {b, cc};
                                  Number Theory/CRT.cpp
void eukl(ll &x, ll &y, ll a, ll b) {
 if (!a) { x = 0; y = 1; return; }
 eukl(y, x, b % a, a);
 x -= y * (b / a);
ll mno(ll a, ll b, ll mod) {/*...*/}
pll crt2(ll p. ll a. ll g. ll b) \{ //x == p \pmod{a}, x == q \pmod{b} \}
 if (a==-1) return {-1, -1},// Zarówno wynik, jak i argumenty są
 ll x, y;
                          // postaci x = first (mod second).
                           // Jeśli kongruencja jest niespełnialna
 eukl(x, y, a, b);
 ll nwd=x*a+v*b:
                           // to zwraca (-1, -1).
 if ((p%nwd)!=(q%nwd))
   return {-1, -1};
 a/=nwd; b/=nwd;
 ll nww=a*b:
 ll ret=mno(x*a, q/nwd, nww)+mno(y*b, p/nwd, nww);
 if ((ret%=nww)<0) ret+=nww;</pre>
 return {ret*nwd+(p%nwd), nww*nwd};
                                Number Theory/Euklides.cpp
ll trzesienie(ll a, ll b, ll c) {
 if (c<0) return 0;</pre>
 if (a>b) swap(a, b):
 ll p=c/b;
 ll k=b/a;
 ll d=(c-p*b)/a;
  return trzesienie(b-k*a, a, c-a*(k*p+d+1))+(p+1)*(d+1)+k*p*(p+1)/2;
 //counts pairs of nonnegative integers (x, y) such that ax + by <= c
                            Number Theory/Fibonacci Cycle.txt
Znaleziony cykl niekoniecznie jest najmniejszym cyklem, wiadomo jednak, że cykl(m) <= 6m
Dla m. które nie sa postaci 2 * 5^r. zachodzi ograniczenie: cvkl(m) <= 4m.
Jeśli nie zrobi się Nww, to wiadomo, że cykl(m) <= m * 2**(1 + #dz.pierw.m).
cycle(2) = 3, cycle(5) = 20, cycle(p == 1 \text{ or } 9 \text{ mod } 10) = p - 1
cycle(p == 3 \text{ or } 7 \text{ mod } 10) = 2(p+1) cycle(p^k) = p^{k} - 1 * cycle(k)
cvcle(product p i^a i) = lcm(cvcle(p i^a^i))
                           Number Theory/number theory data.txt
2: 1 prime divisors, largest=2
6: 2 prime divisors, largest=3
30: 3 prime divisors. largest=5
210: 4 prime divisors, largest=7
2.310: 5 prime divisors, largest=11
30.030: 6 prime divisors, largest=13
510.510: 7 prime divisors. largest=17
9.699.690: 8 prime divisors, largest=19
223.092.870: 9 prime divisors, largest=23
6.469.693.230: 10 prime divisors, largest=29
200.560.490.130: 11 prime divisors. largest=31
7.420.738.134.810: 12 prime divisors, largest=37
304.250.263.527.210: 13 prime divisors, largest=41
13.082.761.331.670.030: 14 prime divisors, largest=43
```

```
pi(10)=4
                     pi(20)=8
                                             pi(40)=12
                                                                    pi(75)=21
pi(100)=25
                     pi(200)=46
                                             pi(400) = 78
                                                                    pi(750)=132
pi(1000)=168
                     pi(2000)=303
                                             pi(4000)=550
                                                                    pi(7.500)=950
pi(10000)=1.229
                     pi(20000)=2.262
                                             pi(40000)=4.203
                                                                    pi(75000)=7.393
pi(100000)=9.592
                     pi(200000)=17.984
                                             pi(400000)=33.860
                                                                    pi(750000)=60.238
pi(1000000)=78.498 pi(2000000)=148.933
                                             pi(4000000)=283.146
                                                                    pi(7.500000)=508.261
pi(10000000)=664.579 pi(20000000)=1.270.607 pi(40000000)=2.433.654 pi(75000000)=4.394.304
pi(100000000)=5.761.455
                                             pi(200000000)=11.078.937
pi(400000000)=21.336.326
                                             pi(750000000)=38.703.181
pi(1000000000)=50.847.534
                                             pi(2000000000)=98.222.287
pi(4000000000)=189.961.812
                                             pi(7.500000000)=345.826.612
pi(10000000000)=455.052.511
                                             pi(20000000000)=882,206,716
pi(40000000000)=1.711.955.433
                                             pi(75000000000)=3.125.641.477
pi(100000000000)=4.118.054.813
                                             pi(200000000000)=8.007.105.059
pi(400000000000)=15.581.005.657
                                             pi(750000000000)=28.514.968.374
840=2^3*3*5*7
                                          32 divisors
7.560=2^3*3^3*5*7
                                          64 divisors
                                         240 divisors
720.720=2^4*3^2*5*7*11*13
                                          448 divisors
8.648.640=2^6*3^3*5*7*11*13
73.513.440=2^5*3^3*5*7*11*13*17
                                         768 divisors
735.134.400=2^6*3^3*5^2*7*11*13*17
                                         1.344 divisors
6.983.776.800=2^5*3^3*5^2*7*11*13*17*19
                                                              2.304 divisors
97.772.875.200=2^6*3^3*5^2*7^2*11*13*17*19
                                                              4.032 divisors
963.761.198.400=2^6*3^4*5^2*7*11*13*17*19*23
                                                              6.720 divisors
9.316.358.251.200=2^6*3^3*5^2*7*11*13*17*19*23*29
                                                              10.752 divisors
97.821.761.637.600=2^5*3^4*5^2*7^2*11*13*17*19*23*29
                                                              17.280 divisors
                                                              26.880 divisors
866.421.317.361.600=2^6*3^4*5^2*7*11*13*17*19*23*29*31
8.086.598.962.041.600=2^8*3^3*5^2*7^2*11*13*17*19*23*29*31
                                                                        41.472 divisors
74.801.040.398.884.800=2^6*3^3*5^2*7^2*11*13*17*19*23*29*31*37
                                                                        64.512 divisors
897.612.484.786.617.600=2^8*3^4*5^2*7^2*11*13*17*19*23*29*31*37
                                                                        103.680 divisors
                                   Number Theory/Phi.cpp
struct Coprimes {
  vector <ll> w, dp;
  int qdz(ll v) {
    if (v \le w.back()/v)
      return v-1:
    return w.size()-w.back()/v;
  il phi(ll n) {
    for (ll i=1; i*i<=n; i++) {
      w.push back(i);
      if ((n/i)!=i)
        w.push back(n/i):
    sort(w.begin(), w.end());
    for (ll i : w) {
      dp.push back(i*(i+1)/2):
      for (ll j=1; j*j<=i; j++) {
        if (i>1) dp.back()-=dp[gdz(j)]*(i/j-i/(j+1));
        if (j>1 \&\& j<=i/(j+1)) dp.back()-=dp[gdz(i/j)];
    return dp.back();
  ll ask(ll v) { return dp[qdz(v)]: \frac{1}{v}=\frac{n}{u} for some u
                                    Number Theory/Pi.cpp
struct Primes {
  vector <ll> w. dp:
  int gdz(ll v) {
    if (v<=w.back()/v)</pre>
      return v-1:
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if ((n/i)!=i)
       w.push back(n/i);
    sort(w.begin(), w.end());
    for (ll i : w)
     dp.push back(i-1):
    for (ll i=1; (i+1)*(i+1) <=n; i++) {
     if (dp[i]==dp[i-1]) continue;
     for (int i=(int)w.size()-1: w[i]>=(i+1)*(i+1): i--)
       dp[j] -= dp[gdz(w[j]/(i+1))] - dp[i-1];
    return dp.back();
 il ask(ll v) { return dp[gdz(v)]; }//v==n/u for some u
                                  Number Theory/Rho.cpp
// szvbkie. faktorvzuje liczby z przedzialu [1e18. 1e18 + 1e5] w 9s na cfie
namespace rho {
const int maxy = 50:
const int maxp = 1e6 + 7; // preprocesujemy rozklad do maxp - 1
int ptot:
int d[maxp]:
int pr[maxp];
vector<ll> ans; // rozklad, czyscimy zazwyczaj po kazdym wywolaniu decompose()
inline ll mod(ll a. ll n) {
 if (a >= n) a -= n:
 return a;
inline ll add(ll a, ll b, ll n) {
 a += b:
 mod(a, n);
 return a;
inline ll mul(ll x, ll y, ll p) { // uwaga, ta funkcja rzuca overflow, tak ma byc
 ll ret = x * y - (ll)((ld)x * y / p + 0.5) * p;
 return ret < 0 ? ret + p : ret;
Il fast(ll x, ll k, ll p) {
 ll ret = 1 \% p;
 for (; k > 0; k >>= 1, x = mul(x, x, p)) (k \& 1) \&\& (ret = mul(ret, x, p));
 return ret:
bool rabin(ll n) {
 if (n == 2) return 1;
 if (n < 2 | | !(n \& 1)) return 0:
 ll s = 0. r = n - 1:
 for (; !(r & 1); r >>= 1, ++s)
  for (int i = 0: pr[i] < n && pr[i] < maxv: ++i) {
   ll cur = fast(pr[i], r, n), nxt;
   for (int j = 0; j < s; ++j) {
     nxt = mul(cur, cur, n);
     if (nxt == 1 \&\& cur != 1 \&\& cur != n - 1) return false:
     cur = nxt:
   if (cur != 1) return false;
 return true:
ll factor(ll n) {
```

return w.size()-w.back()/v;

for (ll i=1; i\*i<=n; i++) {</pre>

w.push back(i);

il pi(ll n) {

```
static ll seg[maxp];
  while (true) {
   ll x = rand() % n, y = x, c = rand() % n;
    ll *px = seq, *py = seq, tim = 0, prd = 1;
    while (true) {
     *py++ = y = add(mul(y, y, n), c, n);
     *pv++ = y = add(mul(y, y, n), c, n);
     if ((x = *px++) == y) break;
     ll tmp = prd;
     prd = mul(prd, abs(y - x), n);
     if (!prd) return gcd(tmp, n);
     if ((++tim) == maxv) {
       if ((prd = qcd(prd, n)) > 1 \&\& prd < n) return prd;
   if (tim \&\& (prd = qcd(prd, n)) > 1 \&\& prd < n) return prd;
void decompose(ll n) { // glowna funkcja
 if (n < maxp) {
   while (n > 1) ans.pb(d[n]), n \neq d[n];
 } else if (rabin(n))
   ans.pb(n);
 else {
   ll fact = factor(n):
   decompose(fact), decompose(n / fact);
 }
void init() { // wywolujemy przed pierwsza faktoryzacja, raz na calv program
 for (int i = 2; i * i < maxp; ++i)
   if (!d[i])
     for (int j = i * i; j < maxp; j += i) d[j] = i;
 for (int i = 2; i < maxp; ++i)
   if (!d[i]) {
     d[i] = i;
     pr[ptot++] = i:
   }
  // namespace rho
                                     Strings/Aho.cpp
const int MAXN = 404, sigma = 26;
int term[MAXN], len[MAXN], to[MAXN][sigma], link[MAXN], sz = 1;
void add str(string s) {
 int cur = 0:
 for(auto c: s)
   if(!to[cur][c - 'a'])
     to[cur][c - 'a'] = sz++;
     len[to[cur][c - 'a']] = len[cur] + 1;
   cur = to[curl[c - 'a']:
 term[cur] = cur;
void push links() {
 int que[szl:
 int st = 0, fi = 1;
 que[0] = 0;
 while(st < fi)</pre>
    int V = que[st++];
    int U = link[V]:
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if(!term[V]) term[V] = term[U];
   for(int c = 0: c < sigma: c++)
     if(to[V][c]) {
       link[to[V][c]] = V ? to[U][c] : 0;
       que[fi++] = to[V][c];
     } else {
       to[V][c] = to[U][c];
 }
                                   Strings/Eertree.cpp
const int maxn=1000*1000+7, alfa=26;
int len[maxn], link[maxn], to[maxn][alfa];
int sz. last. n:
char s[maxn]:
void init() {
 s[n++]=len[1]=-1;
 link[0]=1:
 sz=2:
int get link(int v) {
 while (s[n-len[v]-2]!=s[n-1]) v=link[v]:
 return v:
void add letter(char c) {
 s[n++]=c-='a':
 last=get link(last):
 if(!to[last][c]) {
   len[sz]=len[last]+2;
   link[sz]=to[get link(link[last])][c]:
   to[last][c]=sz++:
 last=to[last][c];
int main() {
 init();
 //~ for(int i=1; i<=nn; i++)
   //~ add letter(tekst[i]):
                                 Strings/Graf Podslow.cpp
struct suffix automaton {
 vector<map<char.int>> edges:
 vi link. length:
 int last;
 suffix automaton(string s) {
   edges.pb(map<char.int>()):
   link.pb(-1):
   length.pb(0);
   last = 0:
   for (int i=0: i<siz(s): i++) {</pre>
     edges.pb(map<char,int>());
     length.pb(i+1);
     link.pb(0);
     int r = siz(edges) - 1:
     int p = last;
     while (p \ge 0 \&\& !edges[p].count(s[i])) {
       edges[p][s[i]] = r;
       p = link[p]:
     if (p != -1) {
       int q = edges[p][s[i]];
       if (length[p] + 1 == length[a]) {
         link[r] = q;
       else {
```

```
edges.pb(edges[q]);
         length.pb(length[p] + 1);
         link.pb(link[a]):
         int qq = siz(edges)-1;
         link[q] = link[r] = qq;
          while (p \ge 0 \&\& edges[p][s[i]] == g) {
           edges[p][s[i]] = qq;
           p = link[p];
     last = r;
 }
                                  Strings/Lyndon.cpp
V/1) Przyjmuje słowo s (wypełnione na pozycjach 0, 1, ..., n-1).
     Dzieli słowo s na pewną liczbę słów Lyndona p 1, ... p k tak, że:
     p \ 1 \ge p \ 2 \ge \dots \ge p \ k \ (leksykograficznie)
     Podział jest zapisywany w tablicy b - na i-tej pozycji jest true,
     jeśli nastąpiło cięcie przed i-ta literką.
V/ 2) Znaiduie minimalne leksykograficznie przesuniecie cykliczne słowa.
// 3) Znajduje minimalny leksykograficznie sufiks słowa.
void lyndon(char * s, // Słowo zaczynające się na pozycji 0:
                       // 2) s powinno być sklejone: xx.
                       // Długość słowa s (liczac ew. podwojenie).
           int n.
           int& suf, // 3) pozycja minimalnego leksykograficznie sufiksu.
           int& cyk, // 2) pozycja minimalnego leksykograficznie przes. cykl.
           bool* b) { // Tablica ciecia b.
  for (int i = 0: i < n: i++) b[i] = false: // wykomentui, jeśli nie 1)
  int p = 0, k = 0, m = 1:
  while (p < n) {
   if (m == n \text{ or } s[m] < s[k]) {
     if (p < n / 2) cyk = p; // wykomentuj, jeśli nie 2)
     while (p \le k) {
       p += m - k;
       if (p < n) {
         suf = p: // wvkomentui, ieśli nie 3)
         b[p] = true; // wykomentuj, jeśli nie 1)
     m = (k = p) + 1:
   } else if (s[m++] != s[k++]) k = p;
                                 Strings/Manacher.cpp
// @s[0..n-1] - napis długości @n.
// @r[0..2n-2] - tablica promieni palindromów.
V/s:abaabaacaabbba
void Manacher(const char* s, int n, int* r) {
 for (int i = 0, m = 0, k = 0, p = 0; i < 2 * n - 1; m = i + + - 1) {
   while (p < k \text{ and } i / 2 + r[m] != k)
     r[i++] = min(r[m--], (k + 1 - p++) / 2);
   while (k + 1 < n \text{ and } p > 0 \text{ and } s[k + 1] == s[p - 1])
     k++, p--;
    r[i] = (k + 1 - p++) / 2;
bool is pal(int p, int k) { //Przedział [p, k] obustronnie domknięty
 return r[p + k] >= (k - p + 1) / 2;
                                   Strings/Pref.cpp
void Pref(const char* s, int n, int* p) {
 p[0] = n;
```

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int i = 1, m = 0;
 while (i < n) {
   while (m + i < n \text{ and } s[m + i] == s[m]) m++:
   p[i++] = m:
   \mathsf{m} = \mathsf{max}(\mathsf{m} - 1, 0);
   for (int k = 1: p[k] < m: m--) p[i++] = p[k++]:
                                      Strings/Suf.cpp
void Sufar(const int* s. int n. int alpha. int* sa. int* lcp = nullptr) {
 if (n > 0) sa[0] = 0:
 if (n <= 1) return;</pre>
 vi roz(alpha + 1), wsk(alpha), typ(n + 1), ids(n, -1), news, pos;
 auto star = [&](int i) -> bool { return tvp[i] == 3; }:
 auto Indukui = [\&]() \rightarrow void {
   copy(roz.begin(), roz.end() - 1, wsk.begin());
   sa[wsk[s[n - 1]]++] = n - 1;
   for(int i = 0: i < n: ++i)
     if (sa[i] > 0 \&\& !typ[sa[i] - 1])
       sa[wsk[s[sa[i] - 1]]++] = sa[i] - 1;
   copy(roz.begin() + 1, roz.end(), wsk.begin());
   for(int i = n-1: i >= 0: --i)
     if (sa[i] > 0 \&\& typ[sa[i] - 1])
       sa[--wsk[s[sa[i] - 1]]] = sa[i] - 1;
 };
 tvp[n] = 3:
 for(int i = n-1; i >= 0; --i) {
   sa[i] = -1;
   roz[s[i] + 1]++;
   if (i != n - 1 \&\& s[i] < s[i + 1] + !!tvp[i + 1]) {
     typ[i] = 1;
   else if (typ[i + 1]) {
     tvp[i + 1] = 3:
 partial sum(all(roz), roz.begin());
 copy(roz.begin() + 1, roz.end(), wsk.begin());
 for(int i = 0; i < n; ++i) if (star(i)) sa[--wsk[s[i]]] = i:
 Indukui():
 int nast id = 0, b = -1;
 for(int \bar{i} = 0; i < n; ++i) {
   int a = sa[i]:
   if (!star(a)) continue:
   if (b >= 0) while (a == sa[i] || !star(a) || !star(b)) {
     if (star(a) != star(b) || s[a++] != s[b++]) {
       nast id++:
       break:
   } }
   ids[b = sa[i]] = nast id;
 for (int i = 0: i < n: i++) {
   if (ids[i] == -1) continue;
   news.pb(ids[i]);
   pos.pb(i):
 vi new sa(siz(news));
 Sufar(news.data(), siz(news), nast id + 1, new sa.data());
 fill(sa, sa + n, -1);
 copy(roz.begin() + 1, roz.end(), wsk.begin());
 reverse(all(new sa));
 for (int j : new sa) sa[--wsk[s[pos[i]]]] = pos[i];
 Indukui():
 if (lcp) {
   for (int i = 0; i < n; ++i) ids[sa[i]] = i;</pre>
   for (int i = 0, k = 0; i < n; i++, k = max(0, k - 1)) {
```

```
if (ids[i] == n - 1) { k = 0; continue; }
     const int i = sa[ids[i] + 1]:
     while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]) k++;
     lcp[ids[i]] = k;
} } }
                                  Strings/Ukkonen.cpp
template <typename Char>
struct Ukkonen {
 // Musi być ściśle większe niż jakakolwiek długość słowa.
  static constexpr int kInfinity = numeric limits<int>::max();
  struct Node {
   map<Char, pair<Node*, pii>> transition;
   Node* suflink;
 }:
 // Ta metoda jest wywoływana zawsze gdy tworzona jest krawędź {node}[a, +oo).
 void CreateLeafCallback(Node* node, int a) {
   // printf("CreateLeafCallback({%p}[%d, +oo))\n", node, a);
 // Ta metoda jest wywoływana zawsze gdy krawędź {node}[a, b] zamienia się
 // w dwie krawędzie: {node}[a, c-1], {middle}[c, b].
 void SplitEdgeCallback(Node* node, int a, int b, Node* middle, int c) {
   // printf("SplitEdgeCallback({%p}|%d, %d] -> {%p}|%d, %d] + {%p}|%d, %d]\n".
   //
            node, a, b, node, a, c - 1, middle, c, b):
 }
  // vector<unique ptr<Node>> nodes to delete; // Odkomentować w celu usuwania.
 Node* NewNode() {
   Node* node = new Node():
   // nodes to delete.insert(node);
                                              // Odkomentować w celu usuwania.
   return node;
 vector<Char> text: // Słowo powinno zaimować indeksv [0..n-1].
 Node* root:
 Node* pin;
 Node* last explicit node: // Ostatni wierzchołek ..explicit''.
  // "reserve" warto ustawić na maksymalna dlugość słowa, ale wcale nie trzeba.
  Ukkonen(const int reserve = 0) : root(nullptr), pin(nullptr) {
   text.reserve(reserve):
   root = NewNode():
   pin = NewNode();
    root->suflink = pin;
    last explicit node = root:
   last length = 0:
  void Canonize(Node** s, int* a, int b) {
   if (b < *a) return:</pre>
   pair<Node*, pii> t = (*s)->transition[text[*a]];
   Node* sp = t.st;
    int ap = t.nd.st;
    int bp = t.nd.nd:
    while (bp - ap <= b - *a) {
     *a = *a + bp - ap + 1;
     *s = sp:
     if (*a <= b) {
       t = (*s)->transition[text[*a]];
       sp = t.st;
       ap = t.nd.st;
       bp = t.nd.nd:
   }
 bool TestAndSplit(Node* s. int a. int b. Char c. Node** ret) {
   if (a <= b) {
     pair<Node*, pii>& t = s->transition[text[a]];
     Node* sp = t.st;
```

```
int ap = t.nd.st;
   int bp = t.nd.nd:
   if (c == text[ap + b - a + 1]) {
     *ret = s;
     return true;
   *ret = NewNode():
   t.nd.nd = ap + b - a;
   t st = *ret;
   (*ret)->transition[text[ap + b - a + 1]] =
        mp(sp, mp(ap + b - a + 1, bp));
   SplitEdgeCallback(s, ap, bp, *ret, ap + b - a + 1);
   return false;
  return s->transition.find(c) != s->transition.end();
void Update(Node** s. int* a. int i) {
 Node* oldr = root:
 Node* r;
 bool end = TestAndSplit(*s, *a, i - 1, text[i], &r);
  while (!end) {
   CreateLeafCallback(r, i);
   r->transition[text[i]] = mp(nullptr, mp(i, kInfinity));
   if (oldr != root) oldr->suflink = r;
   oldr = r:
   *s = (*s) -> suflink;
   Canonize(s, a, i - 1);
   end = TestAndSplit(*s, *a, i - 1, text[i], &r);
 if (oldr != root) oldr->suflink = *s:
// Dodaje kolejna litere do drzewa.
void AddLetter(Char z) {
 const int i = siz(text);
 text.pb(z);
 auto it = pin->transition.find(z);
 if (it == pin->transition.end())
   pin->transition[z] = mp(root, mp(i, i));
 Update(&last explicit node, &last length, i);
 Canonize(&last explicit node, &last length, i);
// Zamienia wszystkie krawędzie: [x, +oo) -> [x, text.size()-1].
void ClearInfinities(Node* node = nullptr) {
```

if (node == nullptr) node = root; for (auto& it : node->transition) { if (it.nd.nd.nd == kInfinity) it.nd.nd.nd = siz(text) - 1; else ClearInfinities(it.nd.st);

constexpr int Ukkonen<Char>::kInfinity;
int main() { // Przykład użycia.
 string s = "abcdefgh#";

Ukkonen<char> u(s.size() /\* reserve \*/);
for (char c : s) u.AddLetter(c):

}

template <typename Char>

u.ClearInfinities():

return 0;

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