Chapter 2

Instructions: Language of the Computer



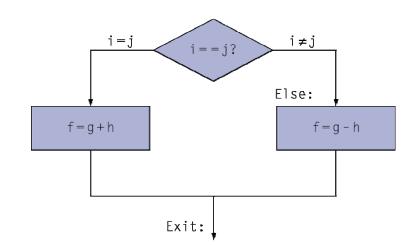
Conditional Operations

- Branch to a labeled instruction if a condition is true
 - Otherwise, continue sequentially
- beq rs, rt, L1
 - if (rs == rt) branch to instruction labeled L1;
- bne rs, rt, L1
 - if (rs != rt) branch to instruction labeled L1;
- j L1
 - unconditional jump to instruction labeled L1

Compiling If Statements

C code:

- f, g, ... in \$s0, \$s1, ...
- Compiled MIPS code:



```
bne $s3, $s4, Else
add $s0, $s1, $s2
j Exit
```

Else: sub \$s0, \$s1, \$s2

Exi t: *...

Assembler calculates addresses

Compiling Loop Statements

C code:

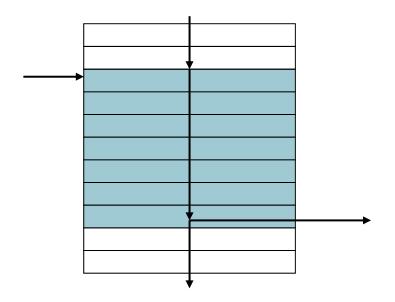
```
while (save[i] == k) i += 1;
```

- i in \$s3, k in \$s5, address of save in \$s6
- Compiled MIPS code:

```
Loop: sll $t1, $s3, 2
add $t1, $t1, $s6
lw $t0, 0($t1)
bne $t0, $s5, Exit
addi $s3, $s3, 1
j Loop
Exit: ...
```

Basic Blocks

- A basic block is a sequence of instructions with
 - No embedded branches (except at end)
 - No branch targets (except at beginning)



- A compiler identifies basic blocks for optimization
- An advanced processor can accelerate execution of basic blocks

More Conditional Operations

- Set result to 1 if a condition is true
 - Otherwise, set to 0
- slt rd, rs, rt
 - if (rs < rt) rd = 1; else rd = 0;
- slti rt, rs, constant
 - if (rs < constant) rt = 1; else rt = 0;</p>
- Use in combination with beq, bne

```
slt $t0, $s1, $s2 # if ($s1 < $s2)
bne $t0, $zero, L # branch to L
```

Branch Instruction Design

- Why not bl t, bge, etc?
- Hardware for <, ≥, ... slower than =, ≠</p>
 - Combining with branch involves more work per instruction, requiring a slower clock
 - All instructions penalized!
- beq and bne are the common case
- This is a good design compromise

Signed vs. Unsigned

- Signed comparison: sl t, sl ti
- Unsigned comparison: sl tu, sl tui
- Example

 - \$1 = 0000 0000 0000 0000 0000 0000 0001
 - slt \$t0, \$s0, \$s1 # signed
 -1 < +1 ⇒ \$t0 = 1</pre>
 - sltu \$t0, \$s0, \$s1 # unsigned +4,294,967,295 > +1 ⇒ \$t0 = 0