

CSE310 Written Homework #1

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1a) Time spent to get to first router: $t = \frac{8 \text{ Mb}}{2 \text{ Mb/s}} = \boxed{4s}$

Time spent in total: $4s \cdot 3 \text{ links} = \boxed{12s}$

1b) $\frac{1 \cdot 10^4 \text{ bits}}{2 \cdot 10^6 \text{ bits/s}} = \frac{1}{2 \cdot 10^2} s = \boxed{0.005s}$

The second packet will be fully received at the first switch at time $t = \boxed{0.01s}$

1c)

$$\begin{aligned} t &= \frac{(n + p + 1)L}{R} \\ t &= \frac{(3 + 800 + 1) \cdot 10^4 \text{ bits}}{2 \cdot 10^6 \text{ bits/s}} \\ t &= \frac{8.04 \cdot 10^6 \text{ bits}}{2 \cdot 10^6 \text{ bits/s}} \\ t &= \boxed{4.02s} \end{aligned}$$

By utilizing message segmentation, the time taken to transmit the message is lowered by a factor of 3. This is because we are fully utilizing all links in the network as much as possible, instead of using one link at a time.

$$\begin{aligned} 2a) \quad R &= 2 \cdot 10^6 \text{ bits/s} \\ d_{prop} &= \frac{\text{link length}}{\text{propagation speed}} = \frac{2 \cdot 10^7 \text{ m}}{2.5 \cdot 10^8 \text{ m/s}} = \frac{2}{25} s = \boxed{0.08 s} \end{aligned}$$

$$R \cdot d_{prop} = 0.08 \cdot 2 \cdot 10^6 = 0.16 \cdot 10^6 = \boxed{1.6 \cdot 10^5}$$

2b) When the first packet is sent through the link, it spends 0.08 seconds in the link. During this 0.08 seconds, the link continues to accept bits at a rate equal to the bandwidth (2Mbps). So the total number of bits in the link is $0.08s \cdot 2 \cdot 10^6 \text{ b/s} = \boxed{1.6 \cdot 10^5 \text{ bits}}$.

2c) The bandwidth-delay product is the maximum number of bits that can be in the link at any given time.

$$2d) w_{bit} = \frac{m}{R \cdot d_{prop}} = \frac{2 \cdot 10^7 m}{1.6 \cdot 10^5 m/b} = \frac{200}{1.6} = \boxed{125 m}$$

If a bit is 125m wide, then a bit is a little more than twice as wide as a football field.

$$2e) w_{bit} = \frac{m}{R \cdot d_{prop}} = \boxed{\frac{m}{R \cdot \frac{m}{s}}}$$

3)

$$F = 1.5 \cdot 10^{10} \text{ bits}$$

$$u_s = 3.0 \cdot 10^7 \text{ bits/s}$$

$$d_i = 2.0 \cdot 10^6 \text{ bits/s}$$

$$\text{For client-server scenario: } D_{cs} = \max \left(\frac{N \cdot F}{u_s}, \frac{F}{d_i} \right)$$

N , the number of users, is represented on the left hand side, has values of 10, 100, and 1000. u , the upload speed, has values of $3 \cdot 10^5$, $7 \cdot 10^5$, and $2 \cdot 10^6$ and is represented on the bottom. All units in the table are seconds.

10	7500	7500	7500
100	50000	50000	50000
1000	500000	500000	500000
	$3 \cdot 10^5$	$7 \cdot 10^5$	$2 \cdot 10^6$

$$\text{For peer-to-peer scenario: } D_{p2p} = \max \left(\frac{F}{u_s}, \frac{F}{d_i}, \frac{N \cdot F}{\sum u_i} \right)$$

Again, N is represented on the left and u at the bottom.

10	500000	214285.71	75000
100	5000000	2142857.14	750000
1000	50000000	21428571.43	7500000
	$3 \cdot 10^5$	$7 \cdot 10^5$	$2 \cdot 10^6$