

# Induction and Deduction in Statistics

{Kevin T. Kelly, Konstantin Genin}

Carnegie Mellon University

Belgrade 2016

# The First Cut in the Philosophical Pie

- All the **objects of human...enquiry** may naturally be divided into **two kinds**, to wit,
- ***Relations of Ideas***, and
- ***Matters of Fact*.**

David Hume, *Enquiry*, Section IV, Part 1.



# The First Cut in the Philosophical Pie

- Any ... **inference** in science belongs to one of **two kinds**:
  1. either it yields **certainty** in the sense that the **conclusion is necessarily true**, provided that the premises are true,
  2. or it does not.
- The first kind is that of **deductive inference**...
- The second kind will here be called '**inductive inference**'.
- R. Carnap, *The Continuum of Inductive Methods*, 1952, p. 3 .



# The First Cut in the Philosophical Pie

- **Deductive inference:**

- Truth preserving.
- Stable (monotonic).
- Non-ampliative.

- **Inductive inference:**

- Everything else.

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

## Induction

- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

## Induction

- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$

Assuming **deterministic** data!



# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

## Induction

- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$

Assuming **stochastic** data...

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

## Induction

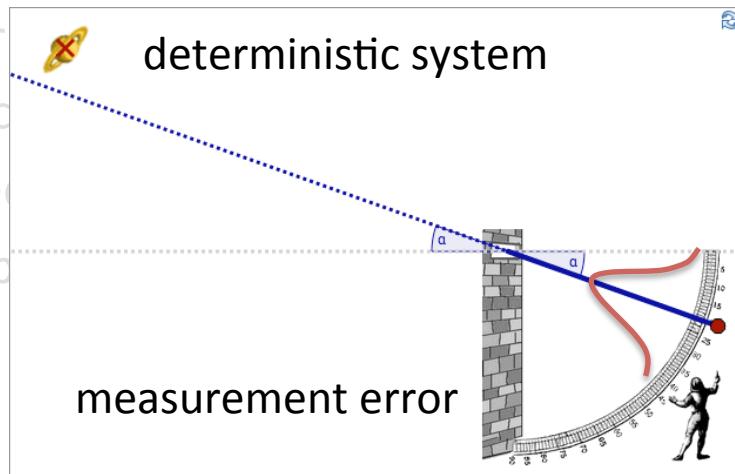
- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

Assuming **stochastic** data...

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$



## Induction

- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

and **all continuous measurement is stochastic!**

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$



$$\begin{aligned} \frac{dx}{\sqrt{x^3 + 1}} &= \frac{dx}{\sqrt{x^3 + 1}} = \left[ \frac{dx}{x + \epsilon} \right] \cdot \frac{\sqrt{x^3 + 1}}{dx - \epsilon dx} dt = \\ &\frac{6t}{\epsilon} \left( \frac{t^3 + 1}{t - 1} - \frac{1}{t + 1} \right) dt = 6 \left( t^2 + 1 + \frac{1}{t - 1} \right) dt \\ &= 6 \left[ \frac{(t^2 + 1)^2}{2} + \sqrt{t^2 + 1} \cdot (n \sqrt{t^2 + 1} + 1) \right] + C \end{aligned}$$

## Induction

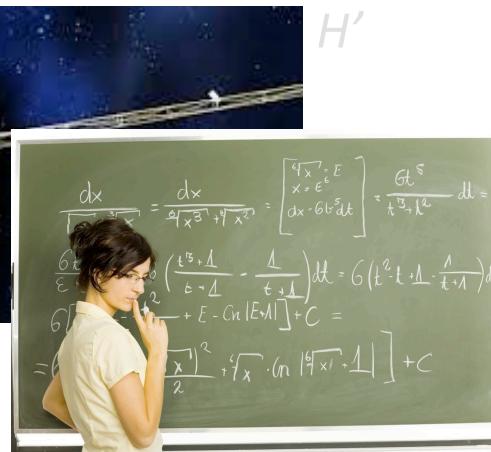
- Inferring universal  $H$
- Choosing between universal  $H_0$ ,  $H_1$ ,  $H_2$ , ...
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H$ ,  $H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$

Also, **real calculation** occasions **probable error** from numerical approximation, human error, and cosmic rays.

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$



## Induction

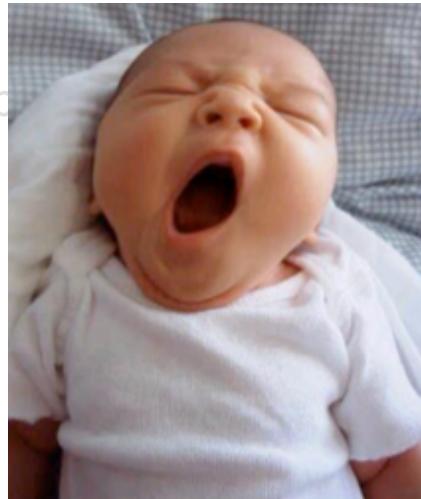
- Inferring universal  $H$
- Choosing between universal  $H_0$ ,  $H_1$ ,  $H_2$ , ...
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H$ ,  $H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$
- Calculation

Also, **real calculation** occasions **probable error** from numerical approximation, human error, and cosmic rays.

# Inference in Science

## Deduction

- Calculation
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$



## Induction

- Inferring universal  $H$
- Choosing between universal  $H_0, H_1, H_2, \dots$
- Refuting universal  $H$
- Verifying existential  $H$
- Deciding between universal  $H, H'$
- Predicting  $E$  from  $H$
- Hypotheses compatible with  $E$
- Calculation

Boooring!

# A More Revealing First Cut

Induction Deduction

Ideal

Statistical

Refuting universal  $H_0$   
Verifying existential  $H_1$   
Deciding between universal  $H_0, H_1$   
Predicting  $E$  from  $H$   
Hypotheses compatible with  $E$   
Ideal calculation

Rejecting simple  $H_0$   
Accepting composite  $H_1$   
Deciding between simple  $H_0, H_1$   
Direct inference from simple  $H$   
Confidence interval  
Real calculation

Inferring universal  $H_0$   
Choosing between universal  $H_0, H_1, H_2, \dots$

Inferring simple  $H_0$   
Model selection

# Ideal Methods

## Deductive

- Stable
- Guaranteed to avoid error



## Inductive

- Unstable
- Not guaranteed to avoid error



# Statistical Methods

## Deductive

- Stable **in chance**
- Guaranteed **low chance of error**



## Inductive

- Unstable **in chance**
- No guarantee of **low chance of error.**



# Deeper Question

Can one **represent** deductive statistical methods as literally **deducing** their conclusions from **statistical information?**

# Deeper Question

Can one **represent** deductive statistical methods as literally **deducing** their conclusions from **statistical information?**

**Yes.**

# The Structure of Ideal Information

✗ Logic

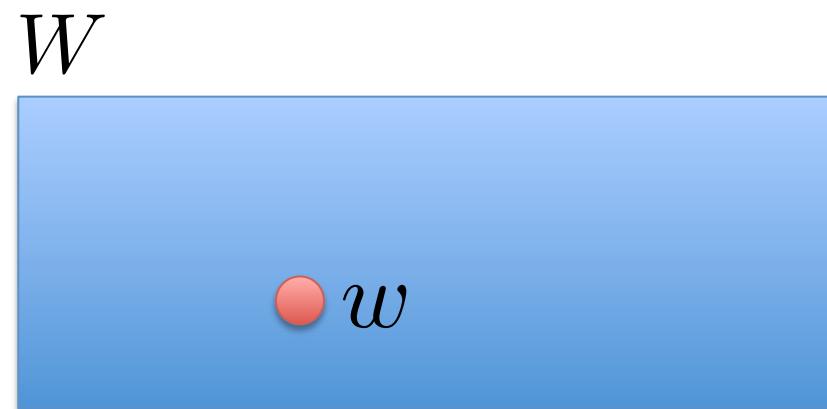
✗ Probability

✓ Topology



# Worlds

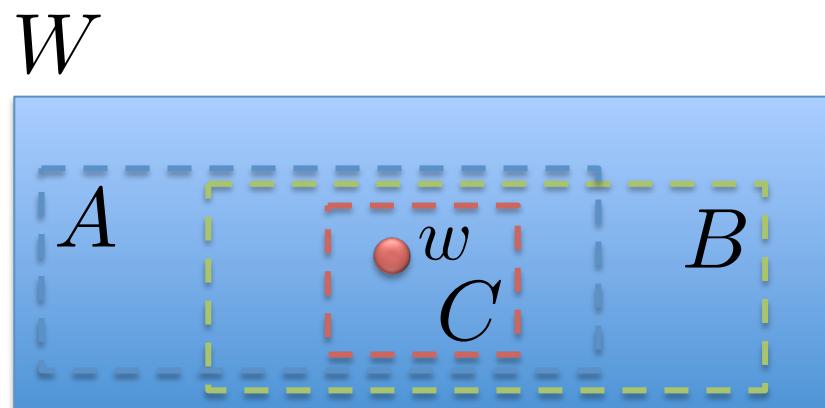
- The points in  $W$  are **possible worlds**.



# The Structure of Information

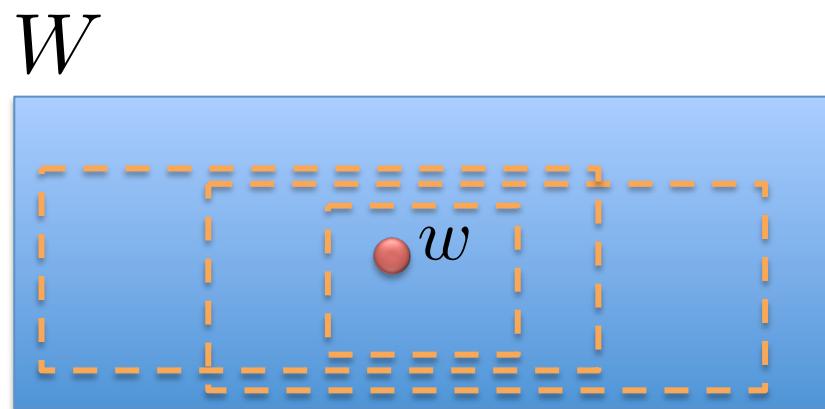
An **information basis**  $\mathcal{I}$  is a **countable** set of propositions called **information states** such that :

1. each world makes **some** information state true;
2. each pair of **true** information states is **entailed** by a true information **state**.



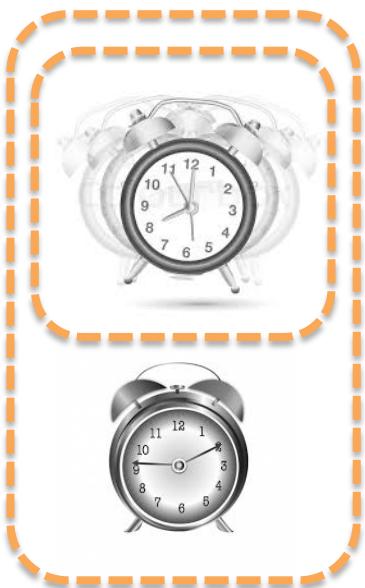
# The Structure of Information

$$\mathcal{I}(w) := \{E \in \mathcal{I} : w \in E\}.$$



# Simplest Example: Alarm Clock

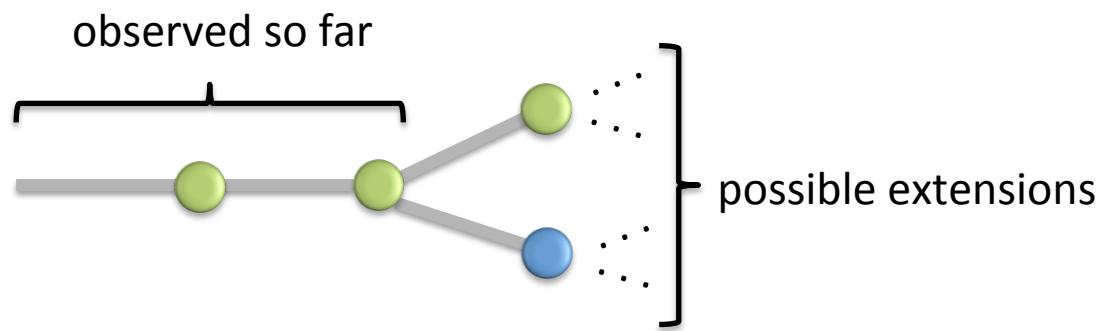
- The theorist is **awakened** by her graduate students only when her theory is refuted.



# Example: Sequential Binary Experiment

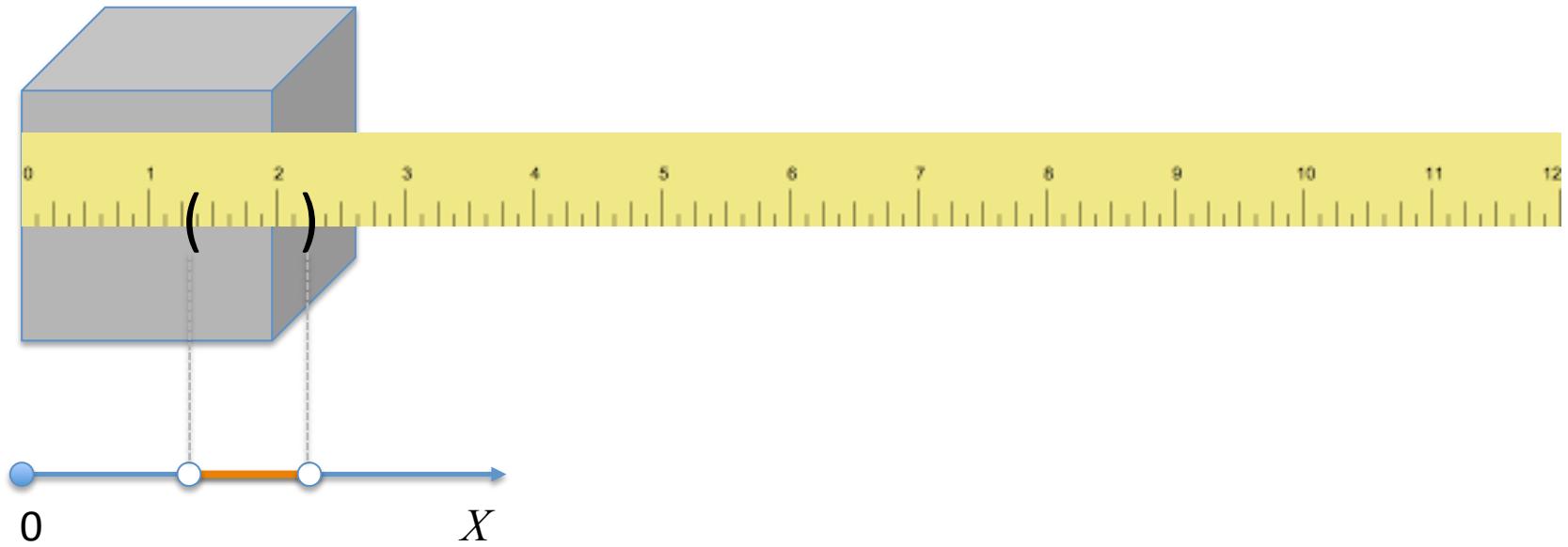
**Worlds** = infinite discrete sequences of outcomes.

**Information states** = cones of possible extensions:



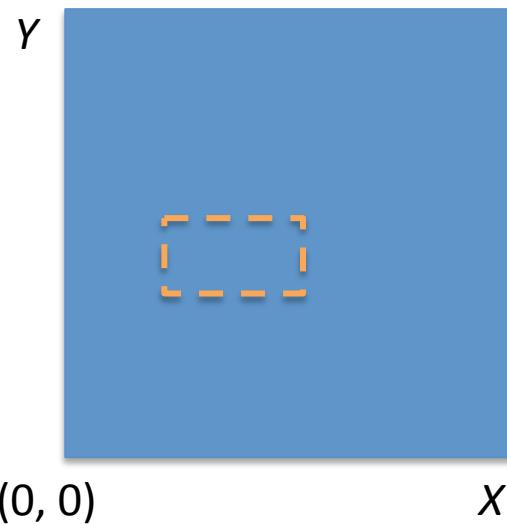
# Example: Measurement of $X$

- **Worlds** = real numbers.
- **Information states** = open intervals.



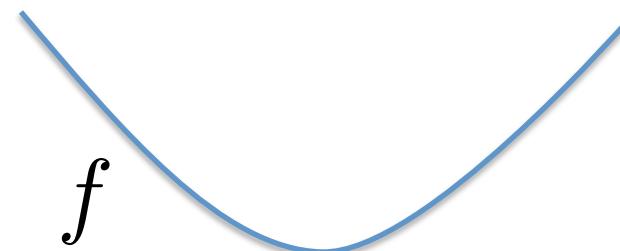
# Example: Joint Measurement

- **Worlds** = points in real plane.
- **Information states** = open rectangles.



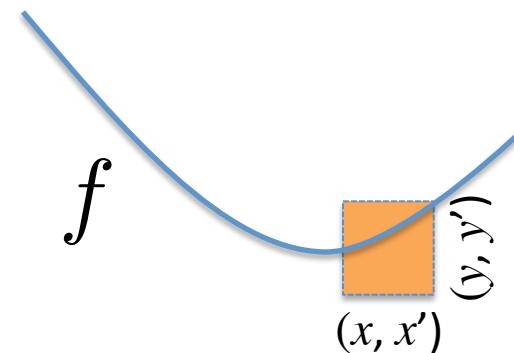
# Example: Equations

- **Worlds** = functions  $f : \mathbb{R} \rightarrow \mathbb{R}$ .



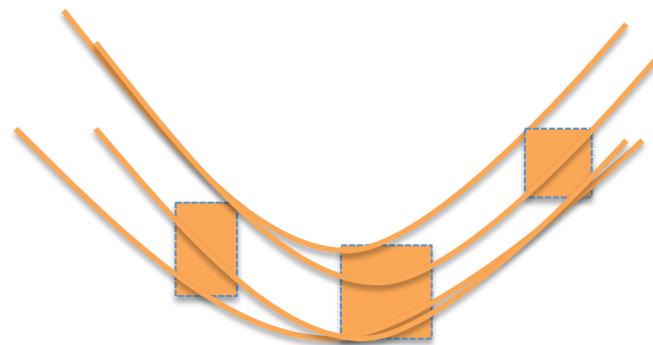
# Example: Laws

- An **observation** is a joint measurement.



# Example: Laws

- The **information state** is the set of all worlds that touch each observation.

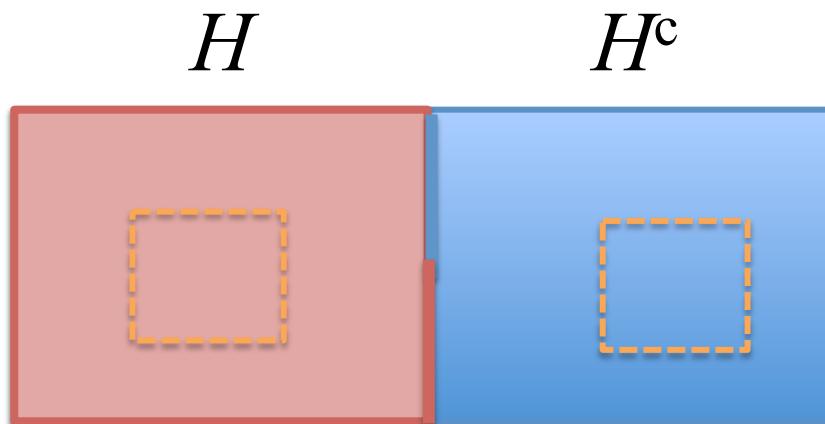


# Deductive Verification and Refutation

$H$  is **verified** by  $E$  iff  $E \subseteq H$ .

$H$  is **refuted** by  $E$  iff  $E \subseteq H^c$ .

$H$  is **decided** by  $E$  iff  $H$  is either verified or refuted by  $E$ .

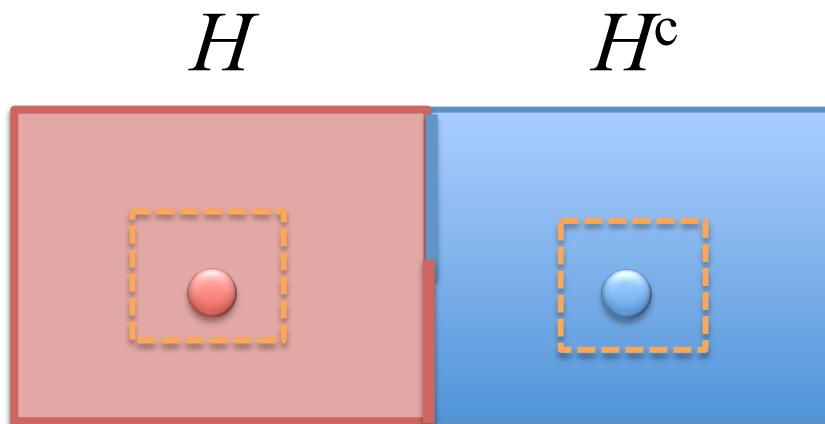


# Will be Verified

$w$  is an **interior [exterior] point** of  $H$  iff

iff  $H$  **will be** verified [refuted] in  $w$

iff there is  $E \in \mathcal{I}(w)$  s.t.  $H$  is **verified [refuted]** by  $E$ .

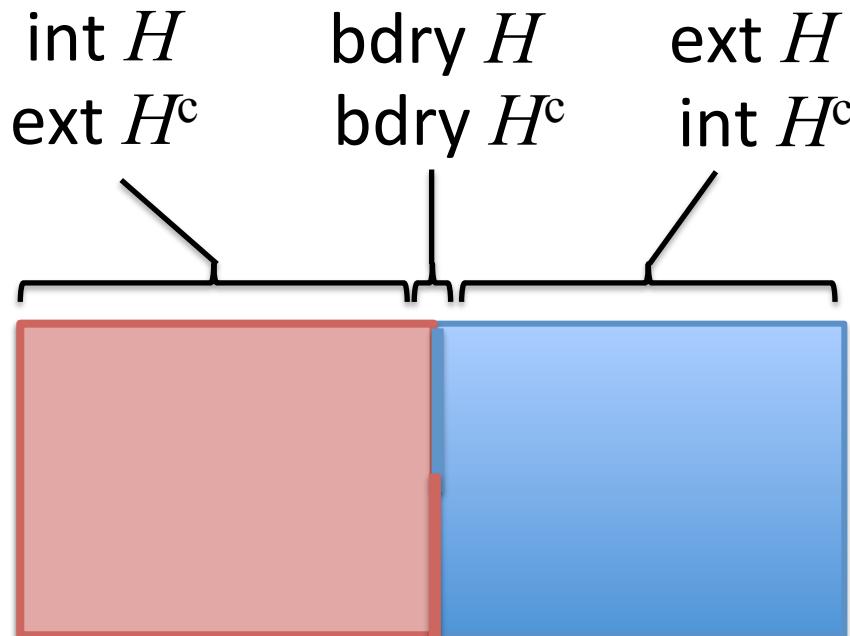


# Will be Verified

**int**  $H$  := the proposition that  $H$  **will be verified**.

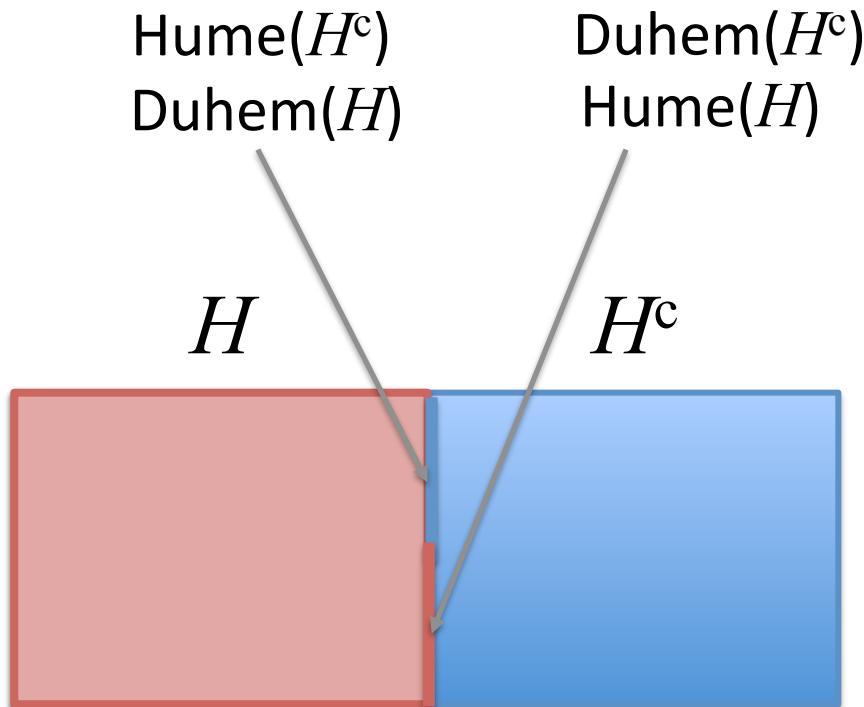
**ext**  $H$  := the proposition that  $H$  **will be refuted**.

**bdry**  $H$  := the proposition that  $H$  **will never be decided**.



# Will be Verified

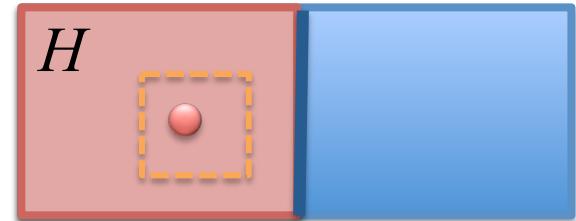
- $\text{bdry}(H) \cap H$  = “you face **Hume’s problem** w.r.t.  $H$ ”;
- $\text{bdry}(H) \cap H^c$  = “you face **Duhem’s problem** w.r.t.  $H$ ”



# Verifiability, Refutability, Decidability

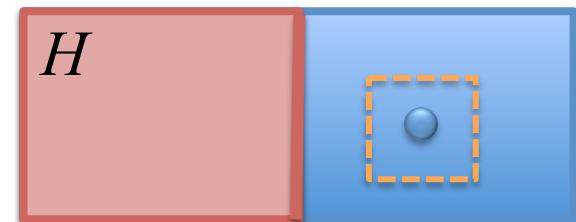
$H$  is **verifiable** iff  $H \subseteq \text{int}(H)$ .

i.e., iff  $H$  will be **verified** however  $H$  is **true**.

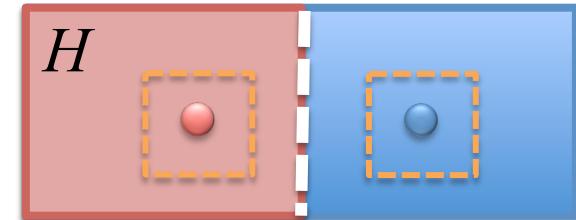


$H$  is **refutable** iff  $\text{cl}(H) \subseteq H$ .

i.e., iff  $H$  will be **refuted** however  $H$  is **false**.



$H$  is **decidable** iff  $H$  is both **verifiable** and **refutable**.



# Methods

- A **verification method** for  $H$  is an **inference rule**  $V(E) = A$  such that in **every** world  $w$ :
  1.  $w \in H$ :  $V$  converges to  $H$  without error.
  2.  $w \in H^c$ :  $V$  always concludes  $W$ .

# Methods

- A **verification method** for  $H$  is an inference rule  $V(E) = A$  such that in every world  $w$ :
  1.  $w \in H$ :  $V$  converges to  $H$  without error.
  2.  $w \in H^c$ :  $V$  always concludes  $W$ .
- A **refutation method** for  $H$  is just a verification method for  $H^c$ .
- A **decision method** for  $H$  converges to  $H$  or to  $H^c$  without error.

# Methods

- A **verification method** for  $H$  is an inference rule  $V(E) = A$  such that in every world  $w$ :
  1.  $w \in H$ :  $V$  converges to  $H$  without error.
  2.  $w \in H^c$ :  $V$  always concludes  $W$ .
- A **refutation method** for  $H$  is just a verification method for  $H^c$ .
- A **decision method** for  $H$  converges to  $H$  or to  $H^c$  without error.
- $H$  is **methodologically verifiable [refutable, decidable]** iff  $H$  has a method of the corresponding kind.

# Verification is Deductive

**Proposition** (truth preservation and non-ampliativity).

If  $V$  is a verifier, refuter or decider for  $H$  and  $V(E) = A$ ,  
then  $E \subseteq A$ .

**Proposition** (monotonicity).

If there is a verifier, refuter or decider for  $H$ , then there is a monotonic one that never drops  $H$  or  $H^c$  after having concluded it.

# Topology

Let  $\mathcal{I}^*$  denote the closure of  $\mathcal{I}$  under union.

**Proposition:**

If  $\mathfrak{I} = (W, \mathcal{I})$  is an information basis  
then  $\mathfrak{I} = (W, \mathcal{I}^*)$  is a topological space.

- $H$  is open iff  $H \in \mathcal{I}^*$ .
- $H$  is closed iff  $H^c$  is open.
- $H$  is clopen iff  $H$  is both closed and open.

# Methodology = Topology

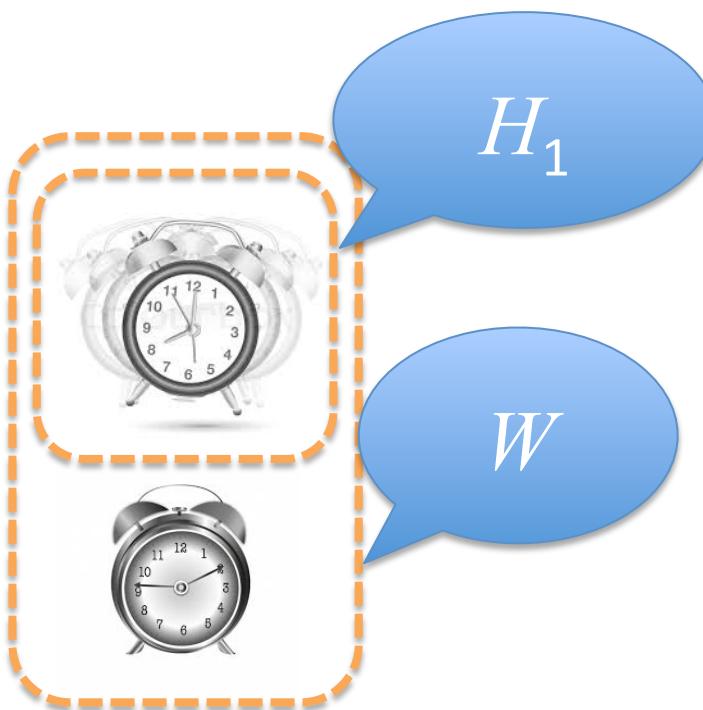
## Proposition.

1. **open** = verifiable = methodologically verifiable.
2. **closed** = refutable = methodologically refutable.
3. **clopen** = decidable = methodologically decidable.

# Simplest Example

$H_0$  = “I will **never** be awakened” is **closed**.

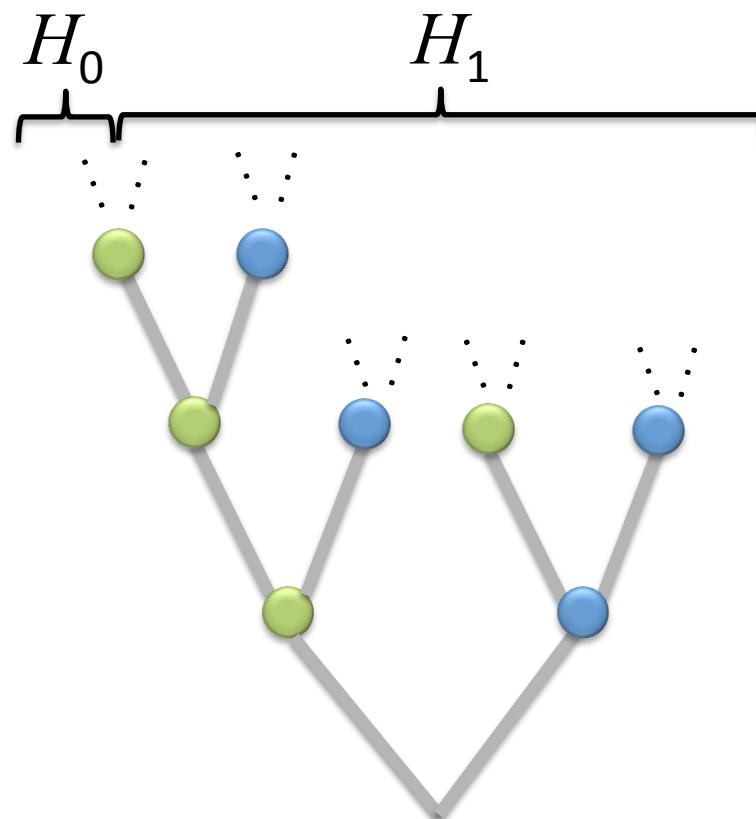
$H_1$  = “I will **eventually** be awakened” is **open**.



# Sequential Examples

$H_0$  = “every outcome is green” is **closed**.

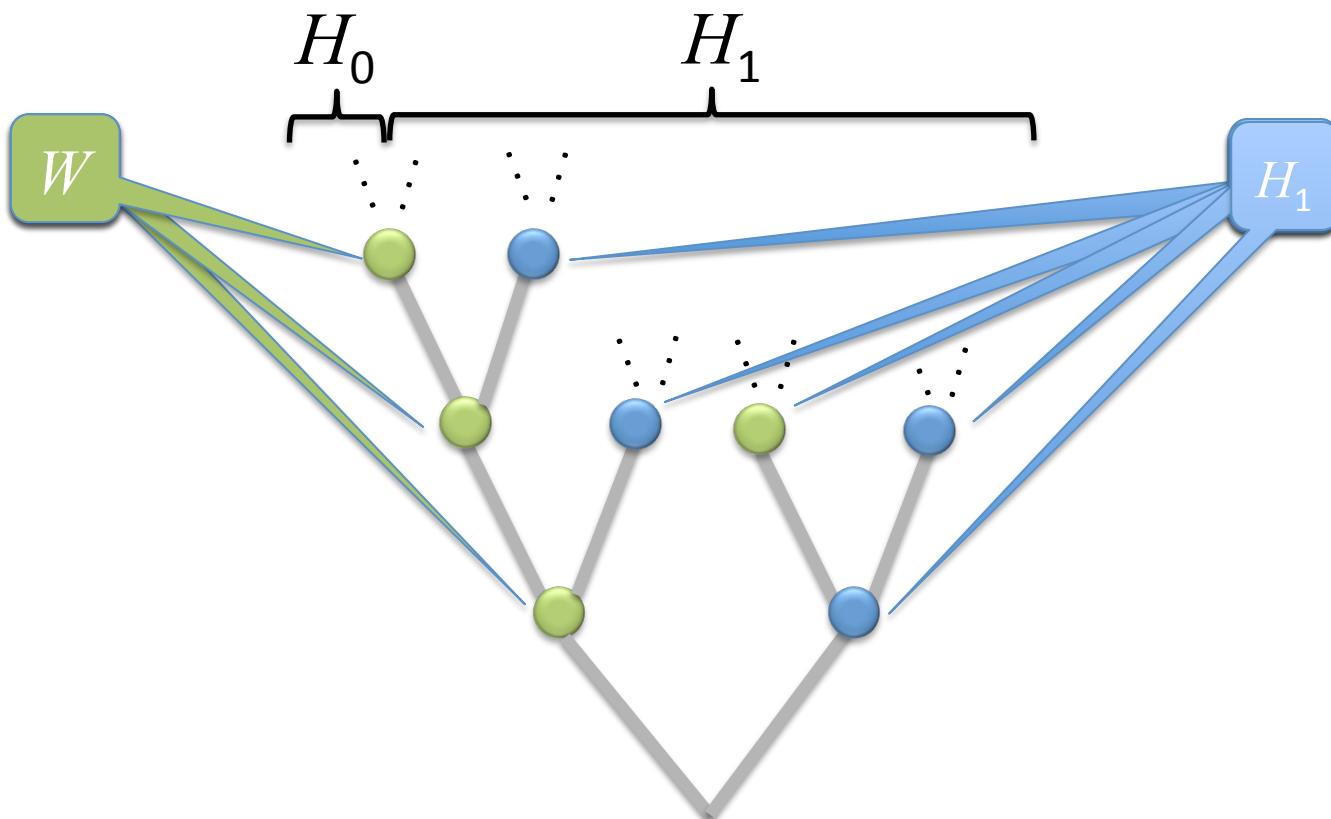
$H_1$  = “some outcome is blue” is **open**.



# Sequential Examples

$H_0$  = “every outcome is green” is **closed**.

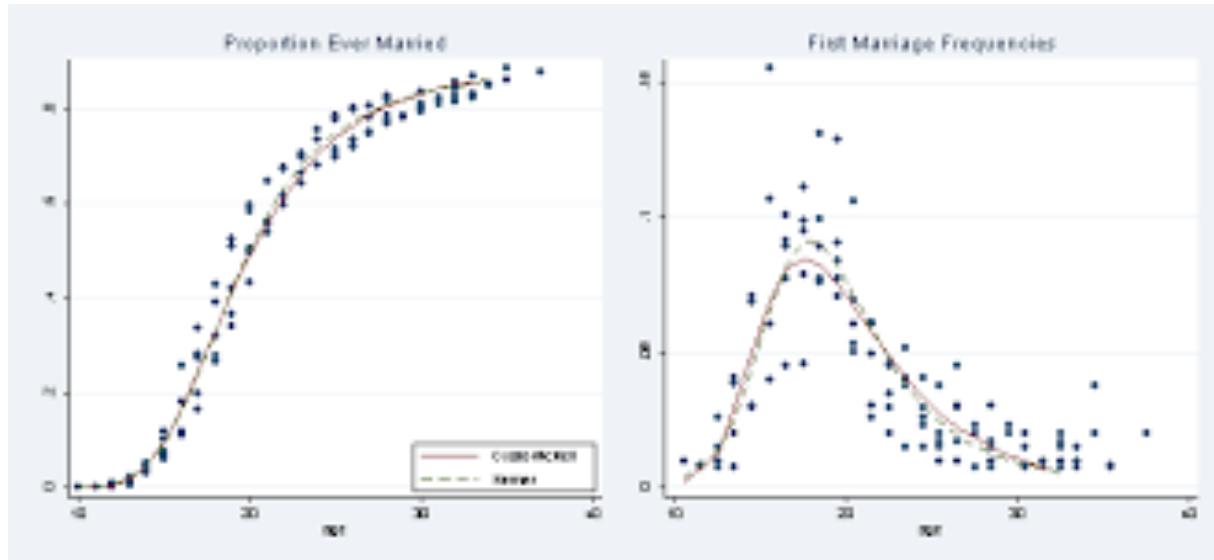
$H_1$  = “some outcome is blue” is **open**.



# **STATISTICAL DEDUCTION AND ITS TOPOLOGY**

# Statistical Methodology

- Does information topology also shed light on statistically deductive methods and problems?



# Skepticism

The approach...

“may be okay if the candidate theories are **deductively** related to observations, but when the relationship is **probabilistic**, I am **skeptical** ...”.



Elliott Sober, *Ockham's Razors*, 2015

# Skepticism

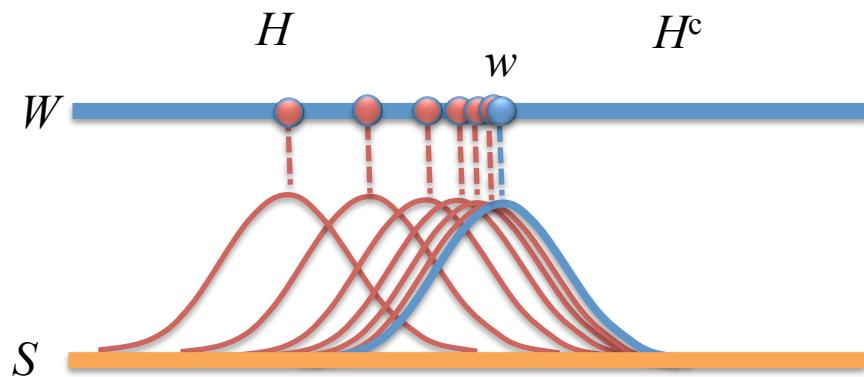
- If it's interesting to **guess** that something is **impossible**...
- then it's surely interesting to **demonstrate** that it is **necessary**.



Elliott Sober, *Ockham's Razors*, 2015

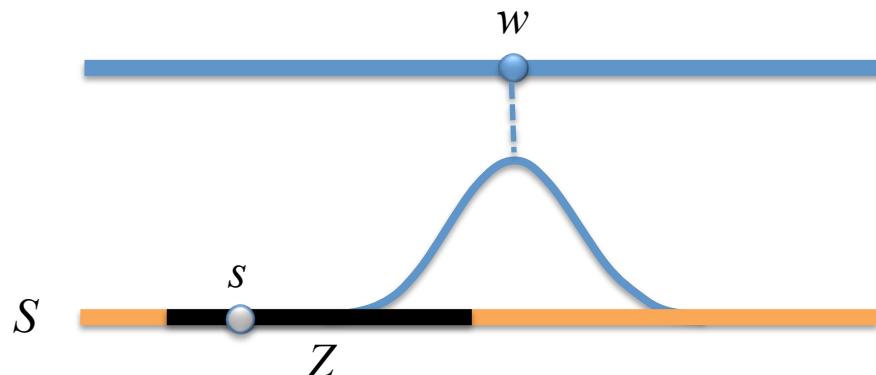
# Statistical Information Topology

Possibilities **nearer** to the truth should be **harder** to rule out by statistical information.



# Gathering Statistical Information

1. The sample space  $S$  has its own (metrizable) topology.
2. Choose a sample event  $Z$  over  $S$ .
3. Obtain sample  $s$ .
4. Observe whether  $Z$  occurs.



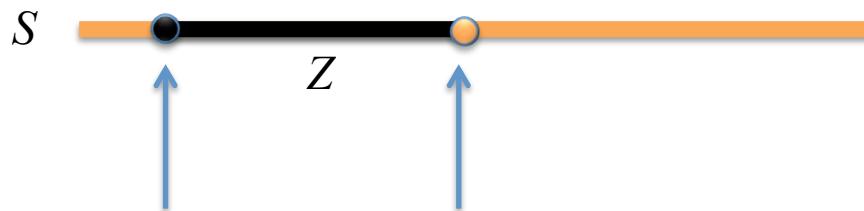
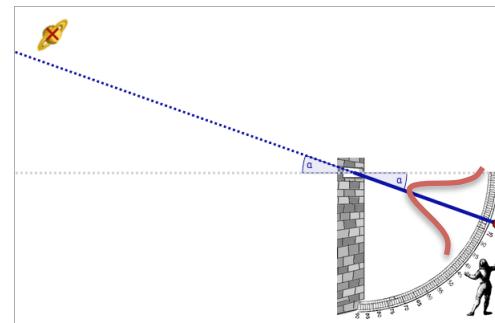
# Feasible Sample Events

- But **every** non-trivial  $Z$  on the real line has **boundary points**.



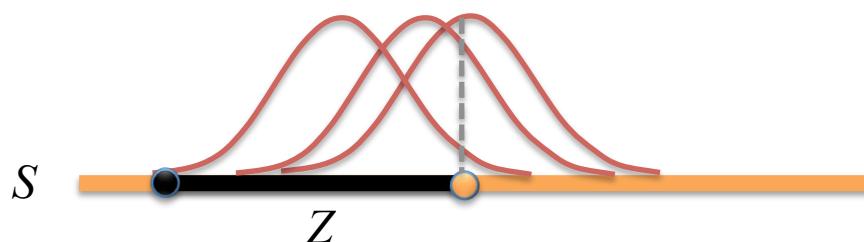
# Feasible Sample Events

- You can't really determine whether a sample hits exactly on the boundary.



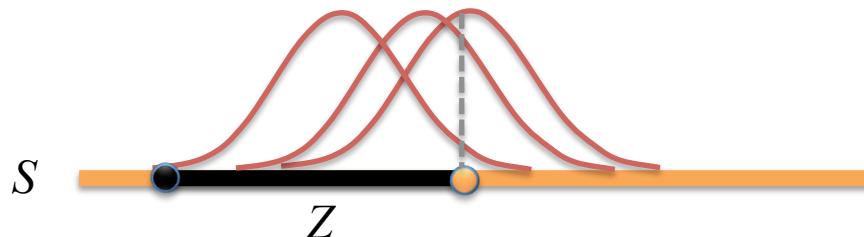
# Feasible Sample Events

- That doesn't matter **statistically** as long as the boundary carries 0 probability.
- So  $Z$  is an observationally **feasible** sample event iff  $p(\text{bdry } Z) = 0$ , for each  $p$  in  $W$ .
- I.e, **feasible**  $Z$  is **almost surely clopen** (decidable) in  $S$ .



# Feasible Statistical Models

- $S$  is **feasible** for  $W$  iff  
 $S$  has a **countable topological basis** of **feasible zones**.

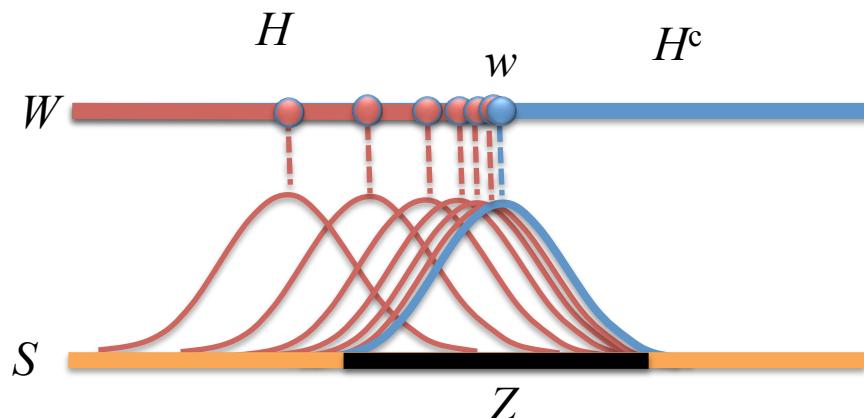


# Statistical Information Topology

$w \in \text{cl}(H)$  iff

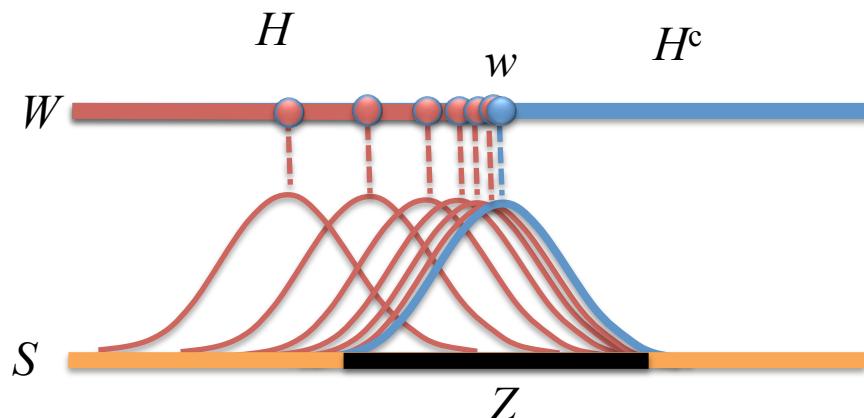
$H$  contains a sequence of worlds  $w_0, \dots, w_n, \dots$  such that  
for every feasible sample event  $Z \subseteq S$ :

$$\lim_{i \rightarrow \infty} p_{w_i}(Z) = p_w(Z).$$



# For Those Who Care

**Proposition:** Assuming that  $S$  is **feasible** for  $W$ ,  
statistical information topology = weak topology.



# IID Sampling

## Proposition.

If  $S$  is **feasible** for  $W$ , then:

1.  $S^N$  is **feasible** for the IID product measures  $w^N$  such that  $w \in W$ ;
2. The information topology on  $W^N$  is **homeomorphic** to the information topology on  $W$ .

# Feasible Statistical Methods

A **feasible statistical method** at sample size  $N$  is a function  $M^N$  from **sample events** in  $S^N$  to **propositions** over  $W$  such that:

$(M^N)^{-1}(H)$  is **feasible**.

A **feasible statistical method** is a collection

$$\{M^N : N \in \mathbf{N}\}$$

of feasible statistical methods at each sample size.

# Statistical Verification Methods

- A **statistical verification method** at **level**  $\alpha > 0$  for  $H$  is a **feasible** statistical method  $\{V^N : N \in \mathbb{N}\}$  with range  $\{W, H\}$  such that:
  1. for  $w \in H$ :  $\lim_{N \rightarrow \infty} p_w^N(V^N = H) = 1$ ;
  2. for  $w \notin H$ :  $p_w^N(V^N = H) \leq \alpha$ , for all  $N$ .
- A statistical **refutation** method at level  $\alpha > 0$  for  $H$  is a statistical verification method for  $H^c$ .
- A statistical **decision** method at level  $\alpha > 0$  for  $H$  is **both**.
- $H$  is **statistically verifiable [refutable, decidable]** iff  $H$  has a statistical verification [refutation, decision] method at **each** level  $\alpha > 0$ .

# The Topology of Statistical Methodology

**Proposition.** Suppose that  $S$  is **feasible** for  $W$ . Then:

1. **open** = statistically verifiable.
2. **closed** = statistically refutable.
3. **clopen** = statistically decidable.

# Stability

**Conjecture:** The methods can be constructed to be **monotonic** in chance of producing  $H$ .

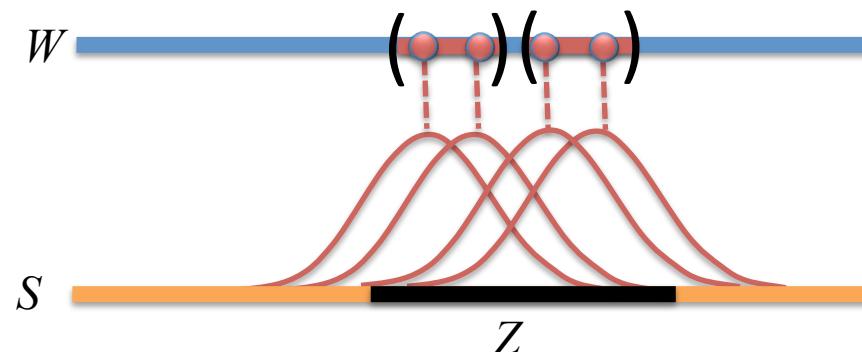
**Conjecture:** the  $\alpha$  level can be made to converge **monotonically** to 0.



# Information Basis

Define **intervals** of worlds w.r.t.  $Z$ :

$$E_Z(a, b) = \{v \in W : a < p_v(Z) < b\}.$$



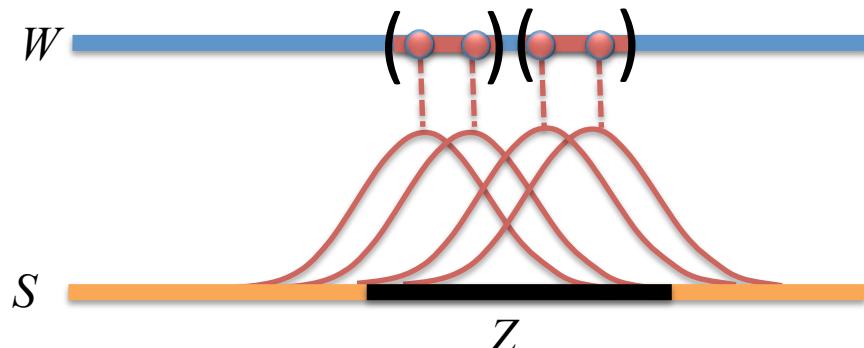
# Information Basis

**Proposition.** Let  $\mathcal{B}$  be a **feasible, countable basis** for  $S$  that (w.l.o.g) is closed under finite intersection.

Then:

$$\mathcal{I} = \{E_Z(a, b) : a, b \in \mathbb{Q} \wedge a < b \wedge Z \in \mathcal{B}\}$$

is a **countable basis** for the statistical information topology on  $W$ .



# Ideal Statistical Information

Think of statistically verified basis elements as ideal statistical information available for other methods.

Let  $\{E_0, \dots, E_i, \dots\}$  enumerate  $\mathcal{I}$ .

Let  $V_i$  statistically verify  $E_i$  at level  $\alpha/2^i$ .

# Statistical Verification as Deduction

Let  $H$  be **open** (statistically verifiable).

Let  $V_H(s) = H$  iff there exists  $E_i \subseteq H$  such that  $V_i(s) = E_i$ .

**Proposition.**  $V$  statistically verifies  $H$  at level  $\alpha$ .

**Moral.** You can safely **think of** statistical verification of  $H$  as literal **deduction** of  $H$  from **ideal statistical information**  $E_i$ .

Similarly for **refutation** and **decision**.

# Statistical Verification as Deduction

	Ideal	Statistical
<b>closed</b>	universal $H_0$	simple null $H_0$
<b>open</b>	existential $H_1$	composite alternative $H_1$
<b>clopen</b>	exhaustive universal $H_0, H_1$ .	exhaustive simple $H_0, H_1$ .

# Statistical Verification as Deduction



# Statistical Verification as Deduction





# EXTENSION TO STATISTICAL INDUCTIVE INFERENCE

# Inductive Inference

Ideal method  $D$  **converges** to  $H$  in  $w$  iff

there exists  $E \in \mathcal{I}(w)$  such that

for all  $F \in \mathcal{I}(w)$  for which  $F \subseteq E$ ,

$$D(F) \subseteq H.$$

Statistical method  $D$  **converges** to  $H$  in  $w$  iff

$$\lim_{N \rightarrow \infty} p_w^N(D = H) = 1.$$

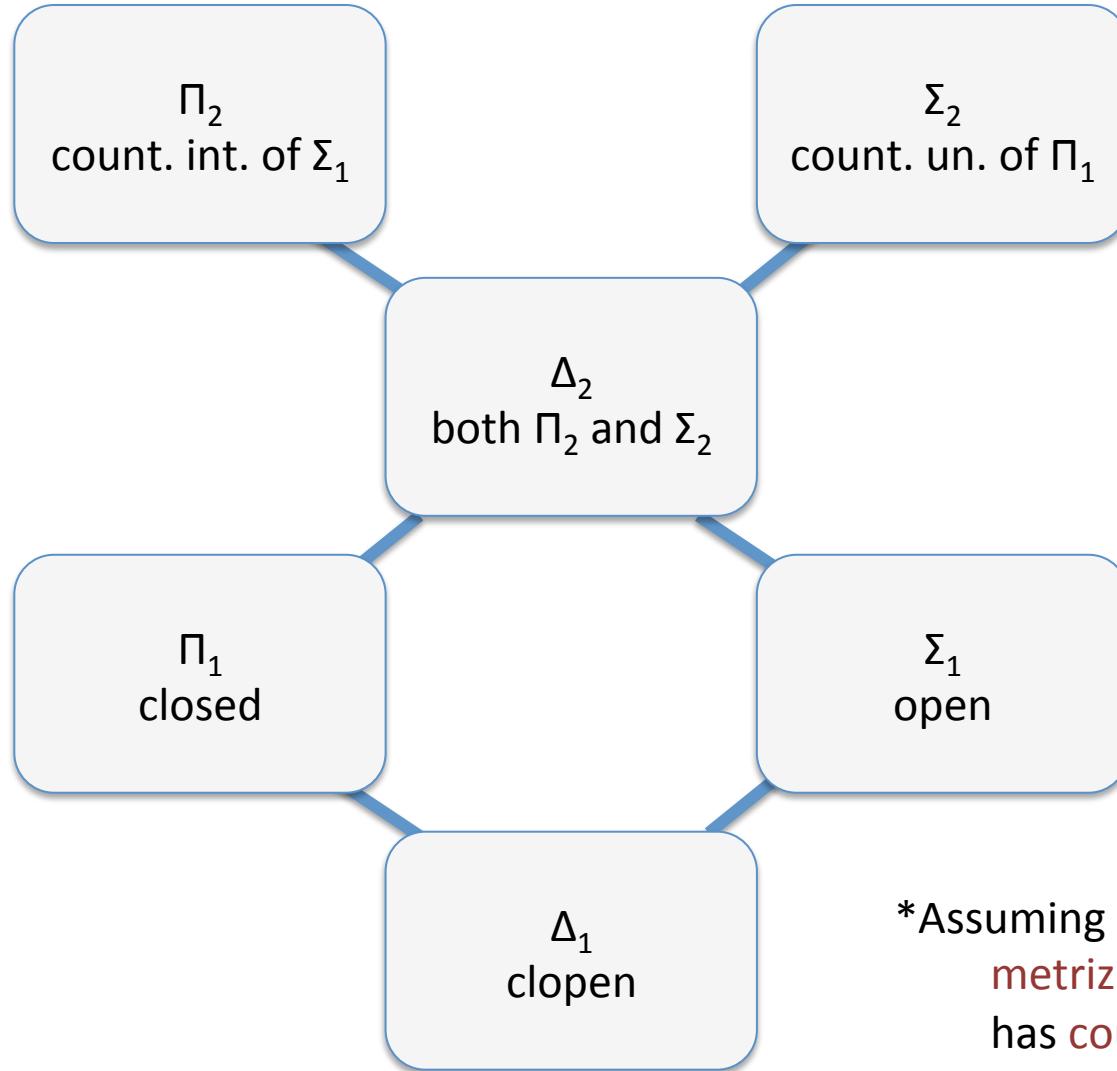
# Inductive Inference

- **$D$  verifies  $H$  in the limit** iff:

$$w \in H \Leftrightarrow D \text{ converges to } H.$$

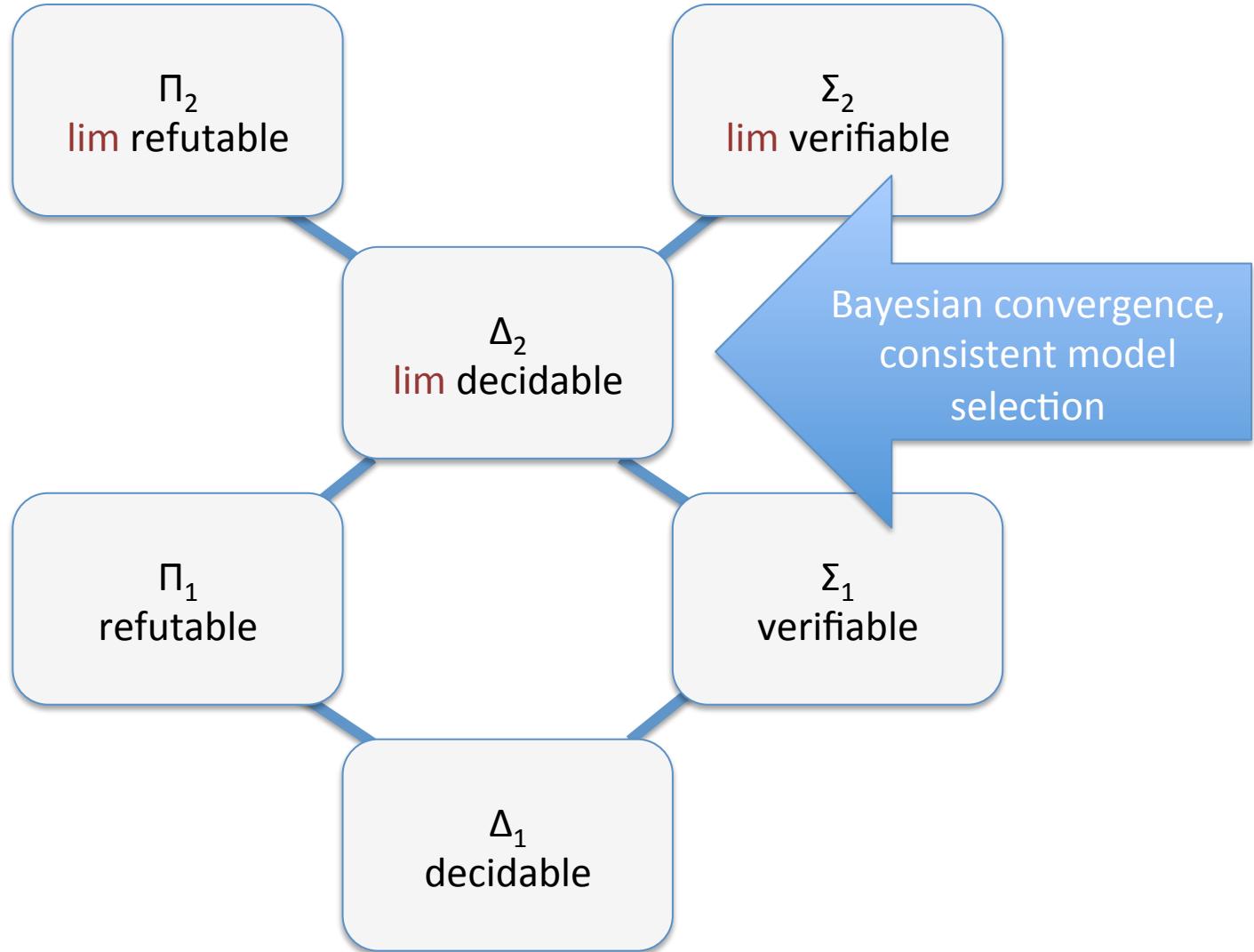
- **$D$  refutes  $H$  in the limit** iff  $D$  verifies  $H^c$  in the limit.
- **$D$  decides  $H$  in the limit** iff  $D$  both verifies and refutes  $H$  in the limit.

# Borel Hierarchy\*



\*Assuming the topology is:  
**metrizable;**  
has **countable basis**.

# Both Ideally and Statistically





# EXTENSION TO OCKHAM'S RAZOR

# Simplicity

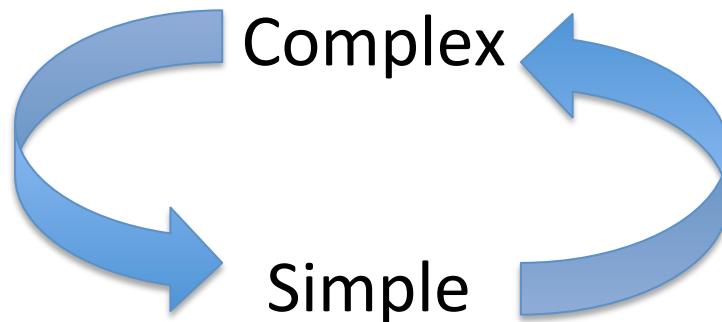
- Simplicity can be defined topologically:

$$A < B \Leftrightarrow A \cap \text{cl}(B) \cap B^c \neq \emptyset.$$

# Epistemic Argument for the Razor

Ideal case:

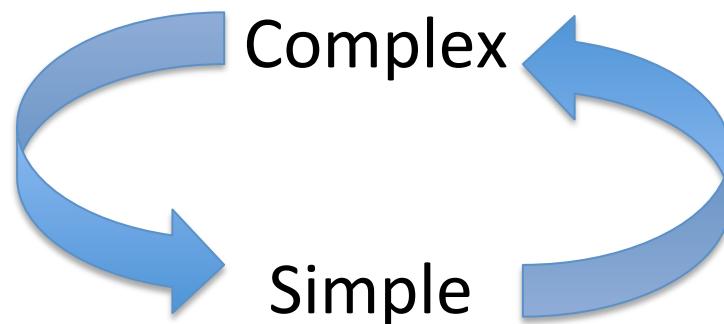
- If you **violate Ockham's razor** then
  1. either you **fail to converge** to the truth or
  2. nature can force you into a **cycle of opinions** (complex-simple-complex), even though such cycles are **avoidable**.



# Epistemic Argument for the Razor

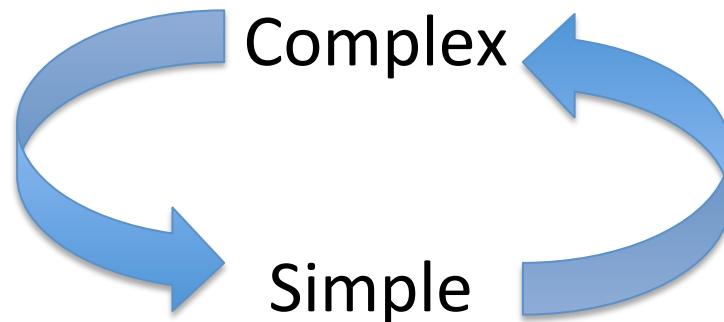
## Statistical case:

- If you **violate Ockham's razor** with chance  $\alpha$ , then
  1. either you **fail to converge** to the truth in chance or
  2. nature can force you into an  $\alpha$ -**cycle of opinions** (complex-simple-complex), even though such cycles are avoidable.



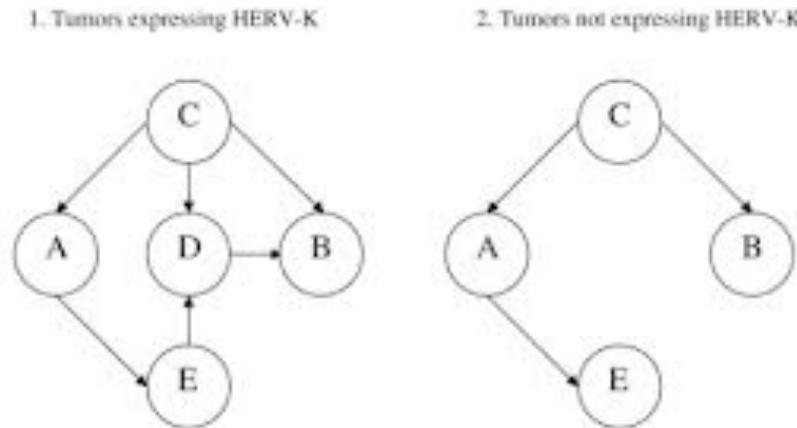
# No Assumption that Reality is Simple

Indeed, by **favoring** a **complex** hypothesis, you incur the cycle in a **complex** world!



# Application: Causal Inference from Non-experimental Data

- Causal network inference from retrospective data.
- That is an inductive problem.
- The search is strongly guided by Ockham's razor.
- We have the only non-Bayesian foundation for it.



Symbols: "A" = cause (YFV); "B" = outcome (cancer); "C" = confounders (recreational solar exposure and high social class); "D" and "E" = mediators (HERV-K antigen and immune response).

# Application: Science

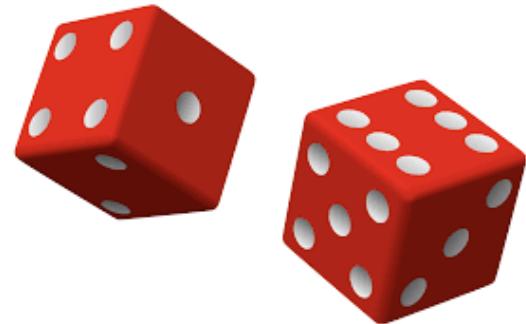
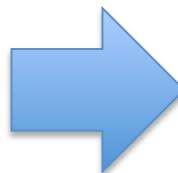
- All scientific conclusions are supposed to be counterfactual.
- Scientific inference is strongly simplicity biased.
- Standard ML accounts of Ockham's razor do not apply to such inferences (J. Pearl).
- Our account does.

# A New Objective Bayesianism?

How much **prior bias toward simple models** is necessary to avoid  $\alpha$ -cycles?

# A Method for Methodology

1. Develop basic methodological ideas in **topology**.
2. Port them to **statistics** via the **statistical information topology**.



# Some Concluding Remarks

1. Information topology is the **structure** of the scientist's **problem context**.
2. The apparent **analogy** between statistical and ideal verifiability reflects **shared topological structure**.
3. Indeed, one can **think of** basis elements as **propositional statistical information** from which statistical conclusions can be **literally deduced**.
4. Thereby, **ideal logical/topological ideas** can be **ported** in a direct and uniform fashion to statistics.