



Simplicity and Scientific Questions

{Kevin T. Kelly, Konstantin Genin}

Ockham's Razor

- **Ockham:** “Pluralitas non est ponenda sine neccesitate.”
- **Science:** “Presume no more complexity than necessary.”



Ockham's Razor

- **Ockham:** “Pluralitas non est ponenda sine neccesitate.”
- **Science:** “Presume no more complexity than necessary.”
- But what **is** simplicity?
- And why **rely** on it?



Indispensable for:

- Theory choice
- Inductive inference, e.g. language learning.
- Statistical model selection
- Causal discovery from non-experimental statistical data.

Epistemic Justification

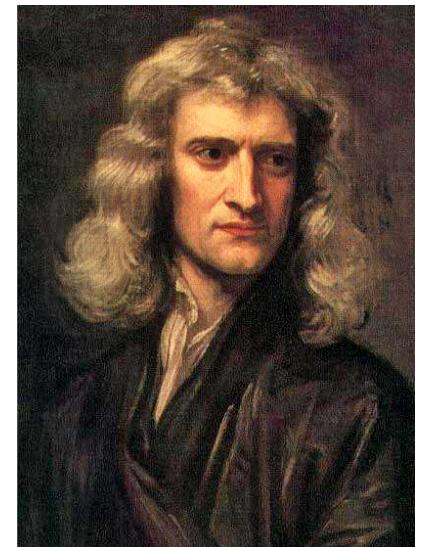
Traditionally, epistemic justification has been
too strictly conceived:

. . . justifying an epistemic principle requires
answering an epistemic question: why are
parsimonious theories more likely to be true?
(Baker, 2013)

Epistemic Justification

When your standards are too high, you are led
either to **metaphysics**,

Nature is pleased with simplicity and affects not the
pomp of superfluous causes (Newton, 1833).



Epistemic Justification

... or **despair**.

[N]o one has shown that any of these rules is more likely to pick out true theories than false ones. It follows that none of these rules is epistemic in character (Laudan, 2004).

Epistemic Justification

Theoretical virtues do not indicate the truth the way litmus paper indicates pH.

Epistemic Justification

Theoretical virtues do not indicate the truth the way litmus paper indicates pH. We can make progress if we don't demand the impossible:

The fact that the truth of the predictions reached by induction cannot be guaranteed does not preclude a justification in a weaker sense (Carnap, 1945).



Epistemic Justification

Truth-indicativeness is too strong a standard. But mere convergence to the truth in the limit is too weak to mandate any behavior in the short run.

Reichenbach is right ... that any procedure, which does not [converge in the limit] is inferior to his rule of induction. However, his rule ... is far from being the only one possessing that characteristic. The same holds for an infinite number of other rules of induction. ... Therefore we need a more general and stronger method for examining and comparing any two given rules of induction ... (Carnap, 1945)

Epistemic Justification

Truth-
Indicative

?

Converges
In the limit

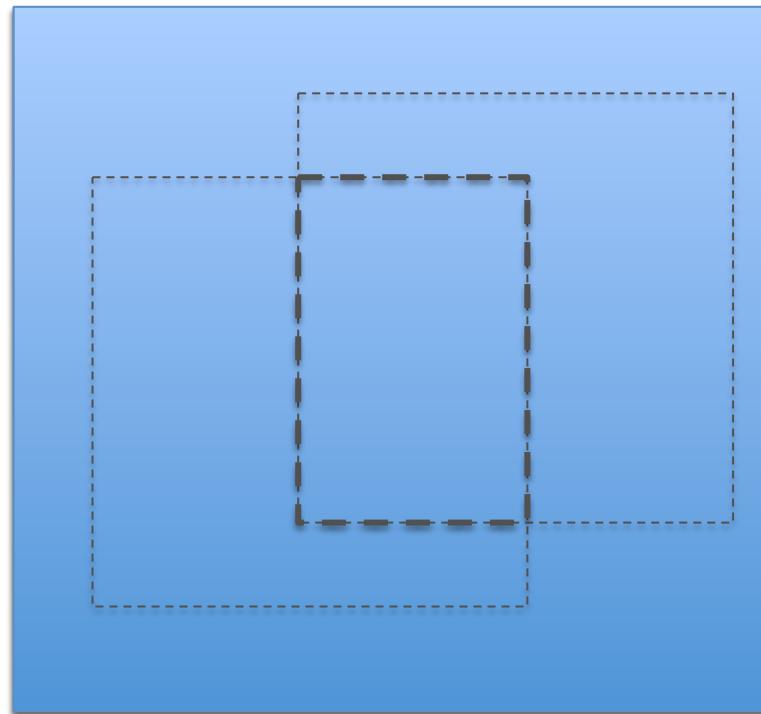
Is there something in between?



1. EMPIRICAL PROBLEMS

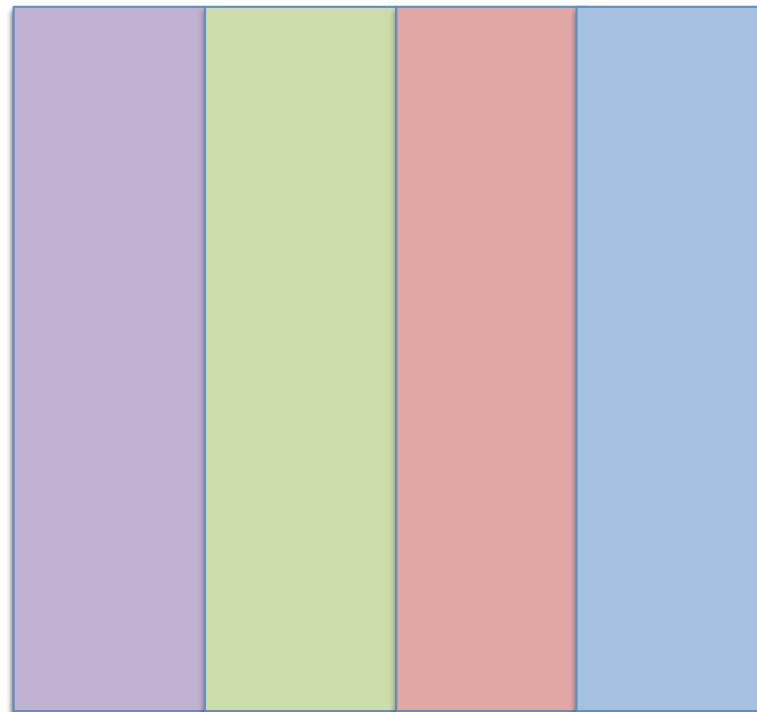
Information Spaces

- W is a set of possible worlds.
- \mathcal{I} is a set of propositional information states.
 1. Covers W .
 2. Closed under finite conjunction.



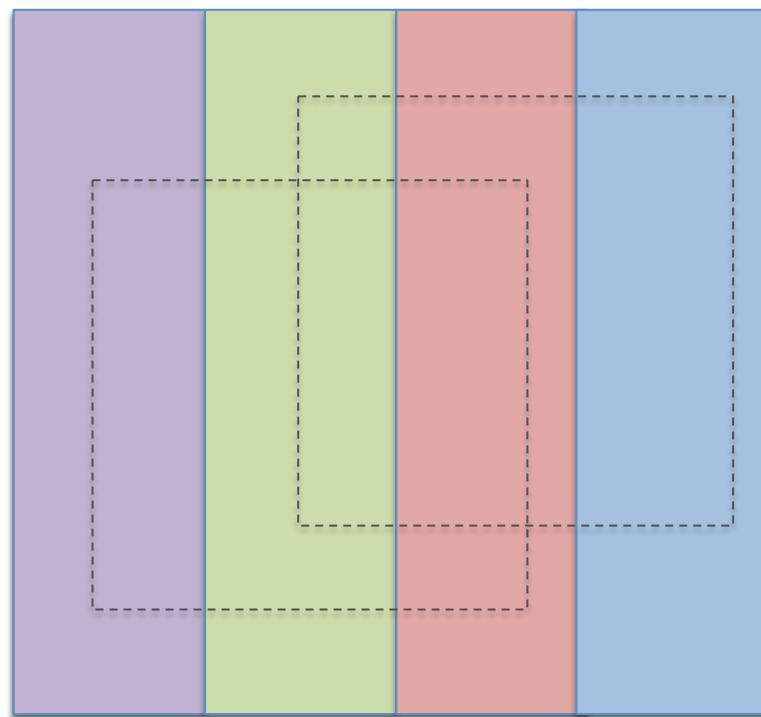
Questions

- A **question** partitions W into countably many possible answers (Hamblin 1958)
- **Relevant responses** are disjunctions of answers.



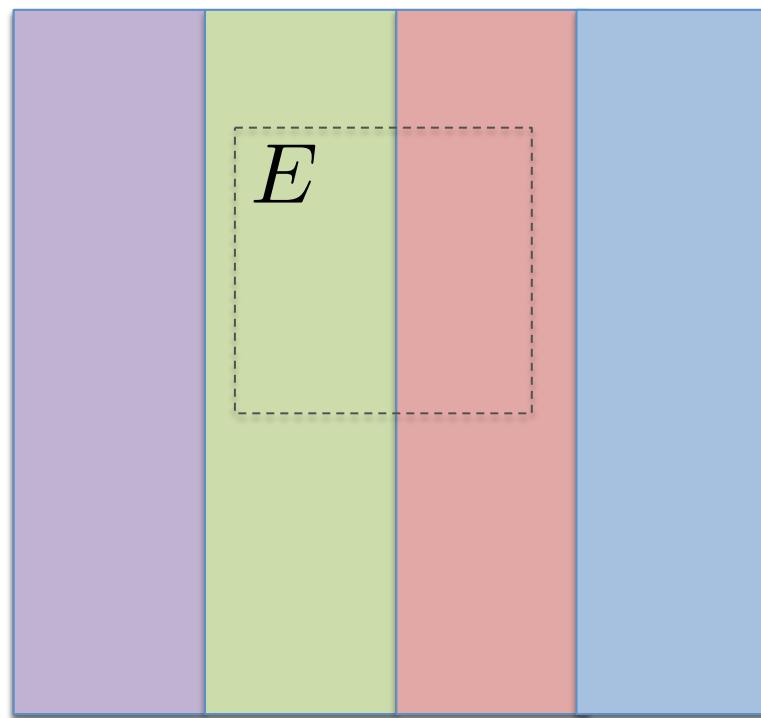
Empirical Problem

$$\mathfrak{P} = (W, \mathcal{I}, \mathcal{Q}).$$



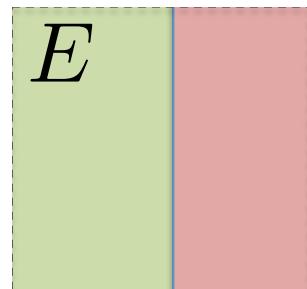
Problem Restriction

\mathfrak{P}



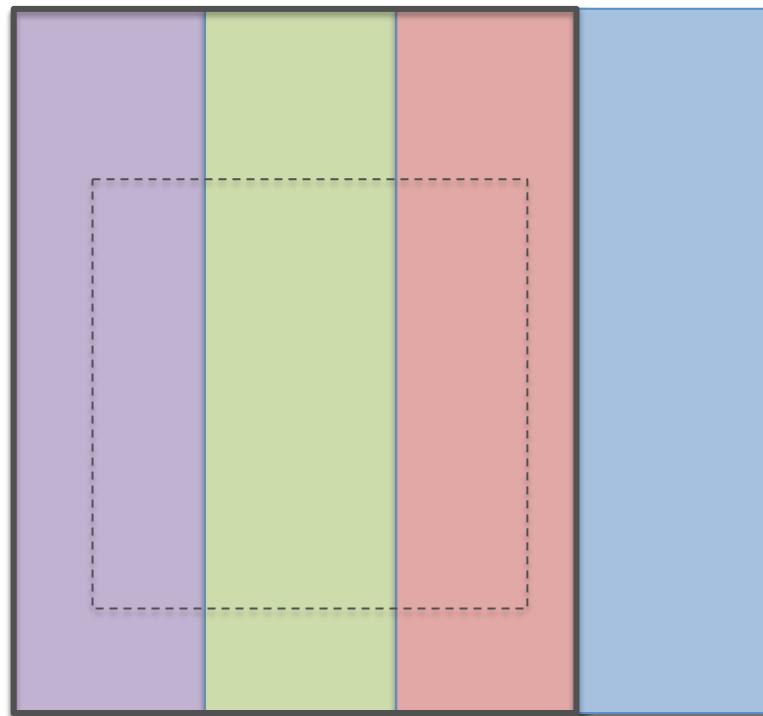
Problem Restriction

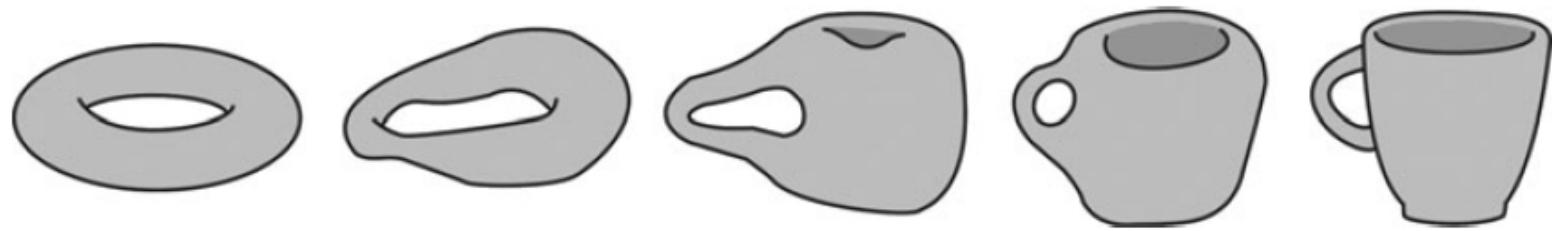
$\mathfrak{P}|_E$



Relevant Response Given E

Disjunction of answers **compatible** with E .

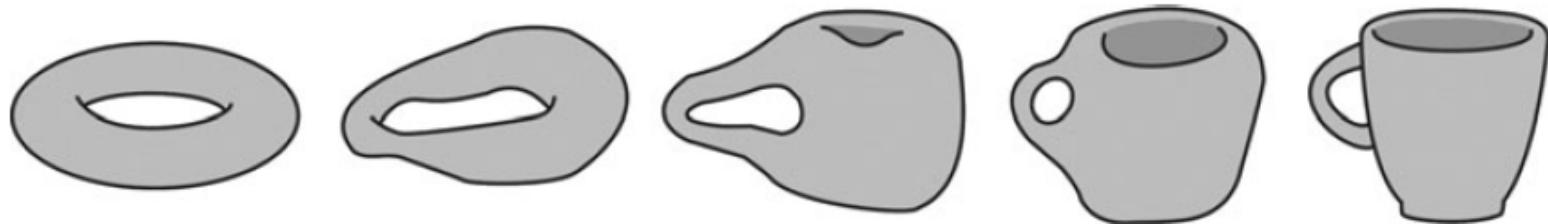




2. INFORMATION TOPOLOGY

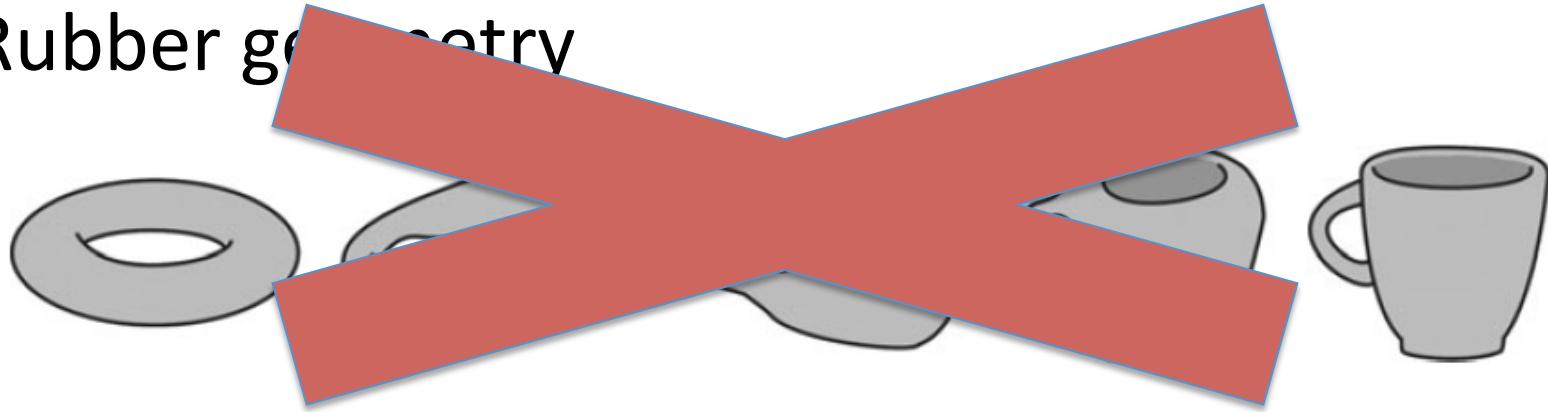
Topology

Rubber geometry



Topology

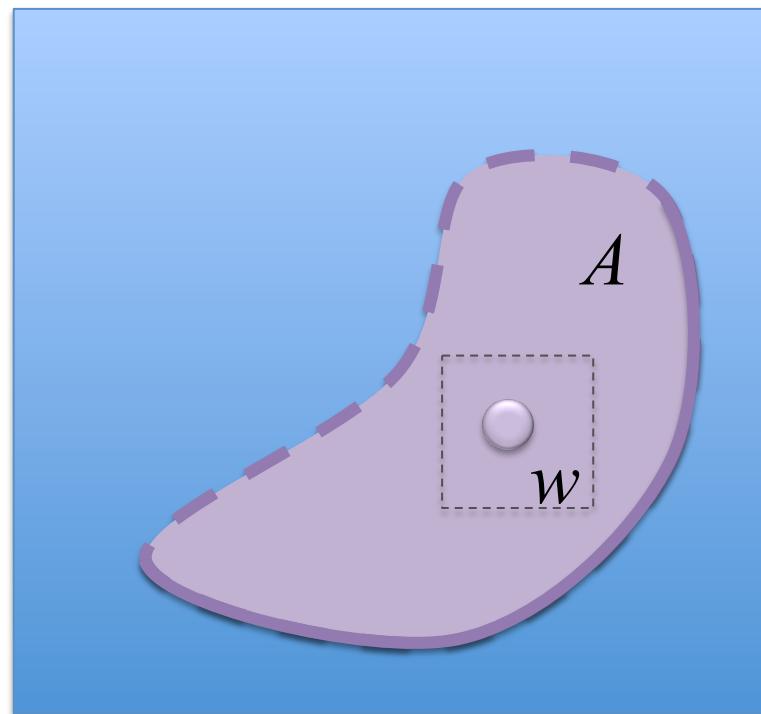
Rubber geometry



The logic of **verification**.

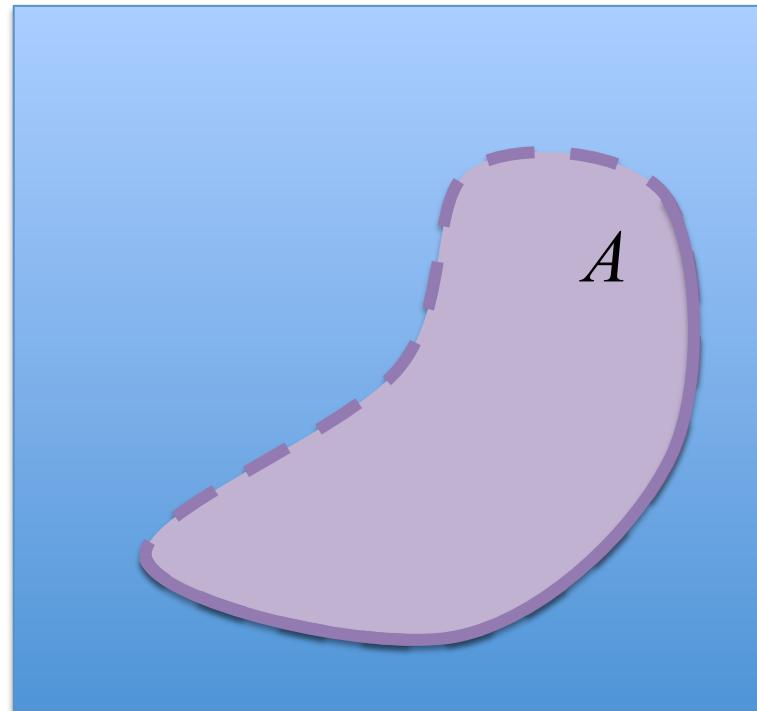
w is an Interior Point of A

W presents information that verifies A .



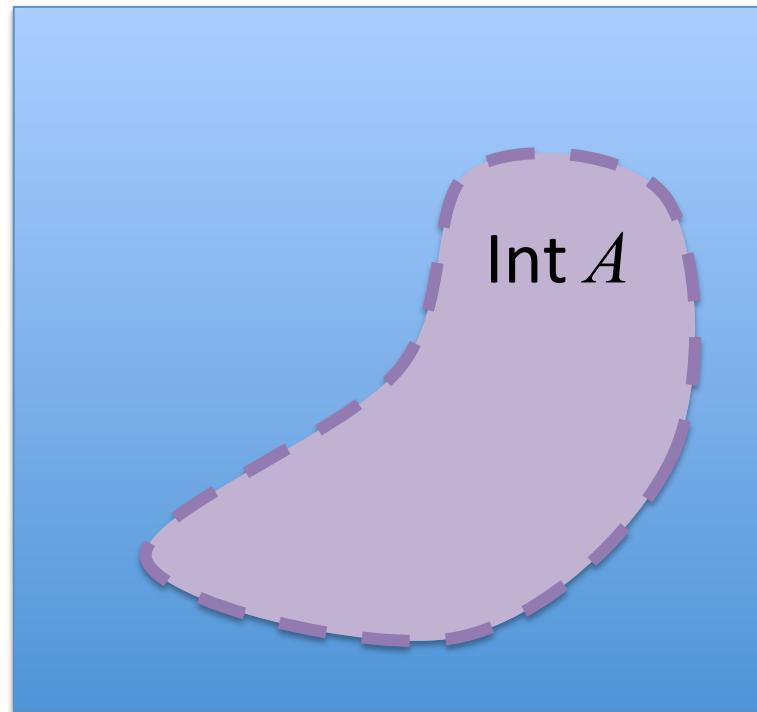
Interior of A

$\text{Int } A = \text{it will be verified that } A.$



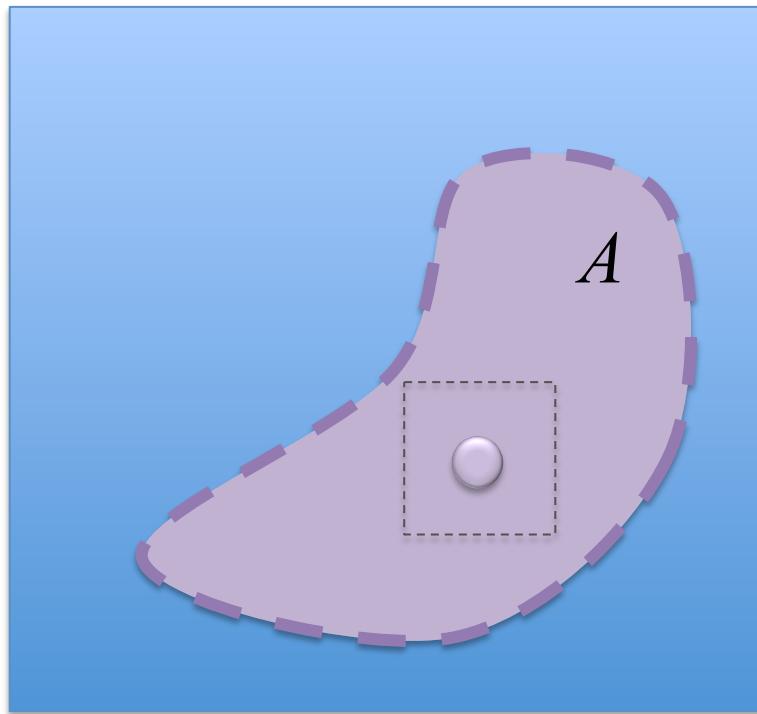
Interior of A

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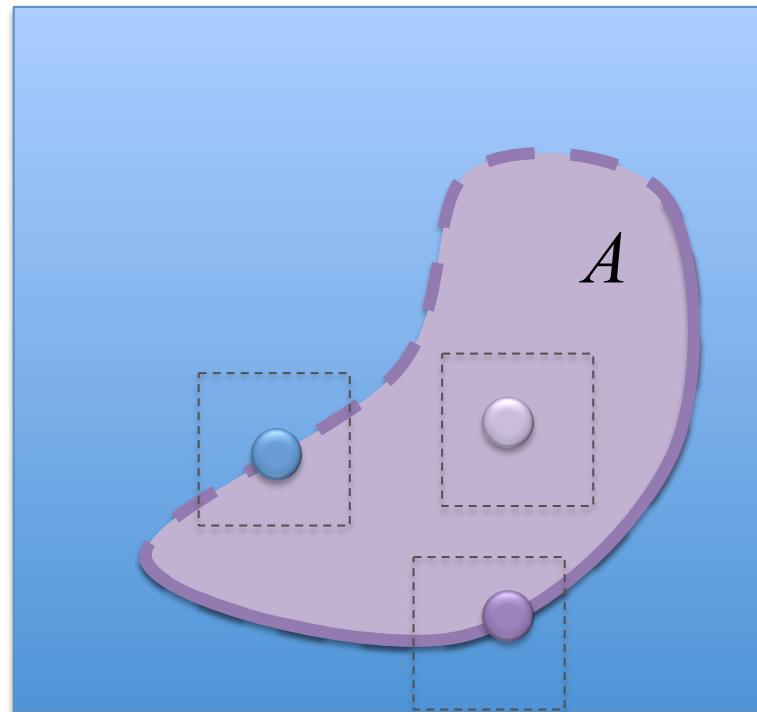
Open = Verifiable

A is open iff A entails that A will be verified.



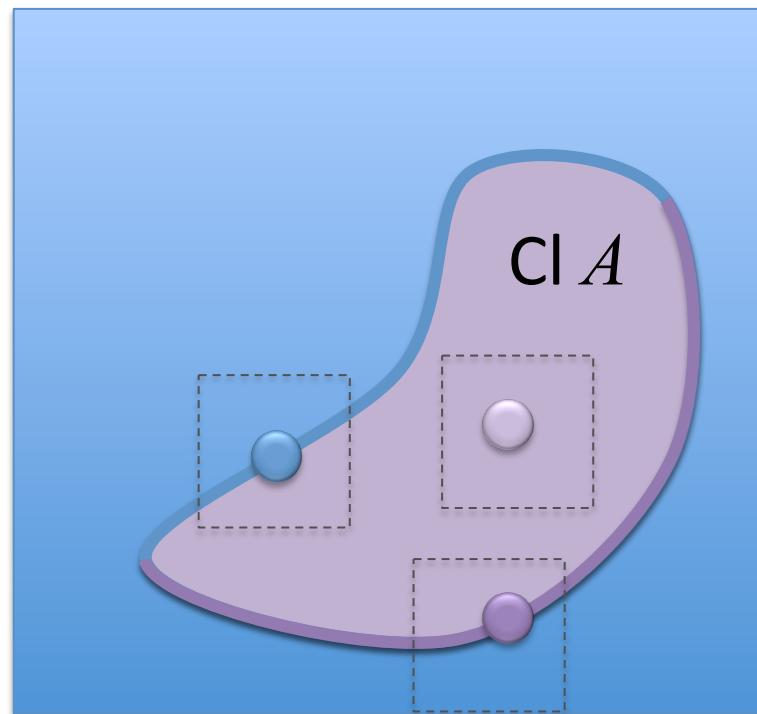
Closure of A

$\text{Cl } A = A$ will never be refuted.



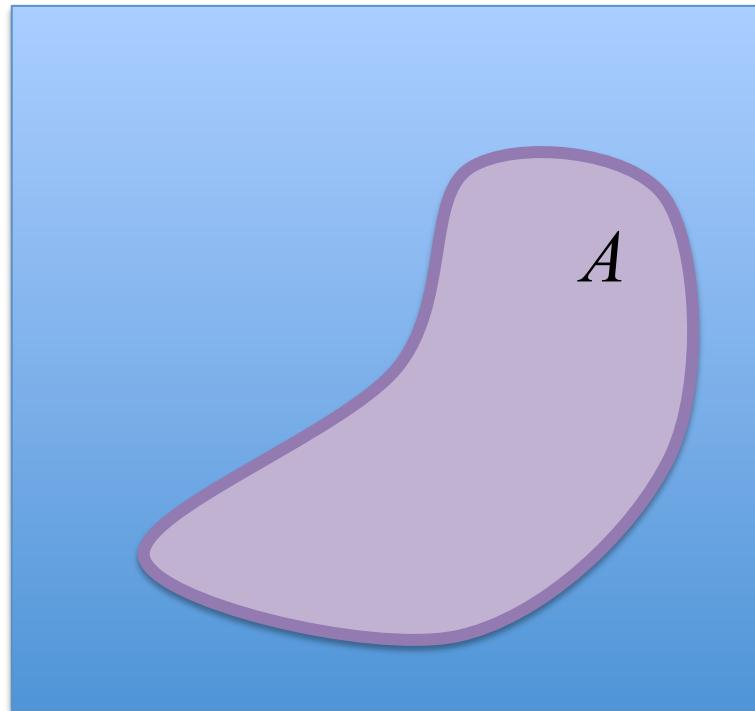
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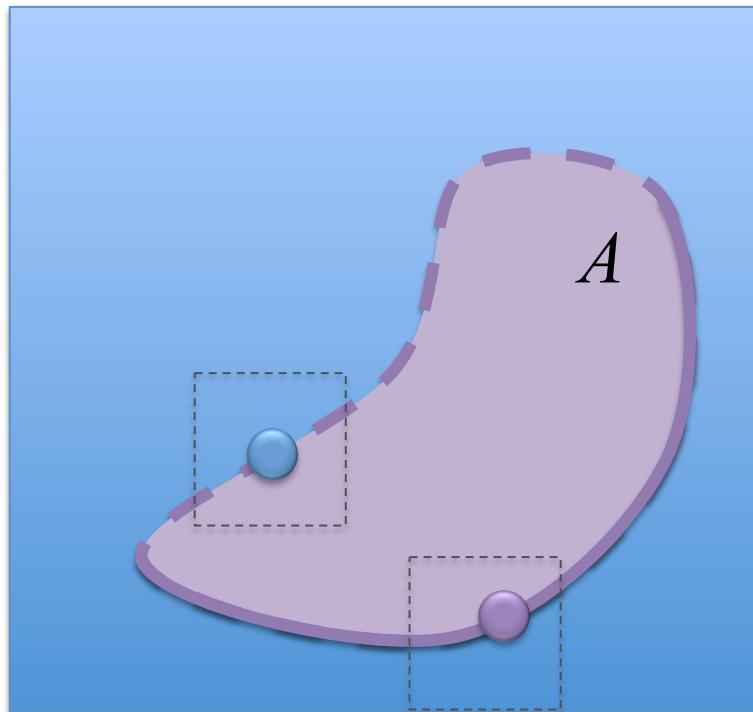
Closed = Refutable

A is closed iff $\neg A$ entails that A will be refuted.



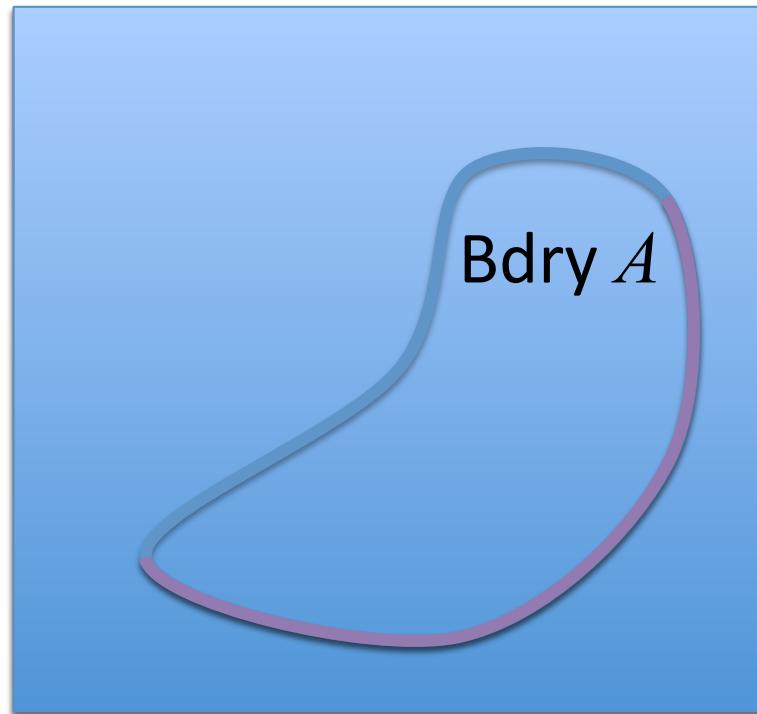
Boundary of A

$\text{Bdry } A = \text{ "}A\text{ will never be decided"}$.



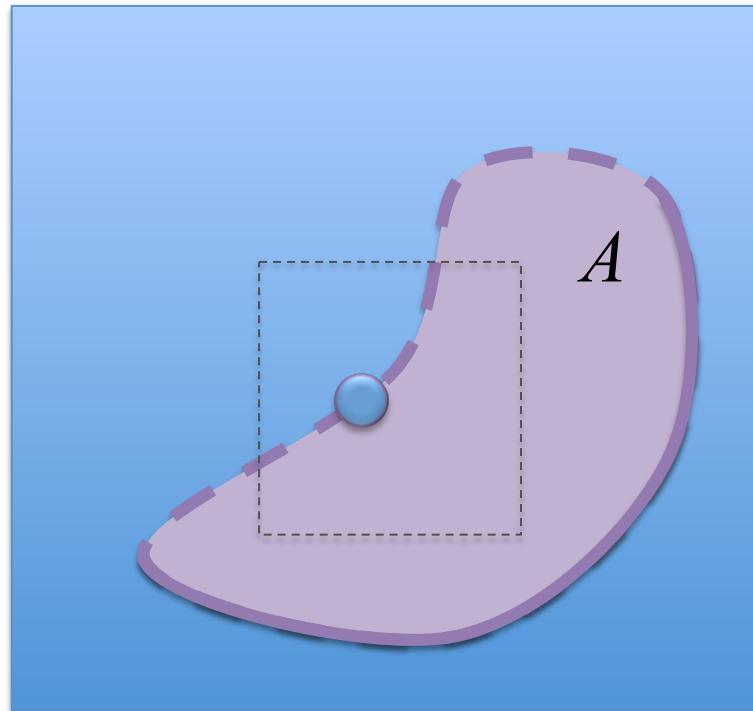
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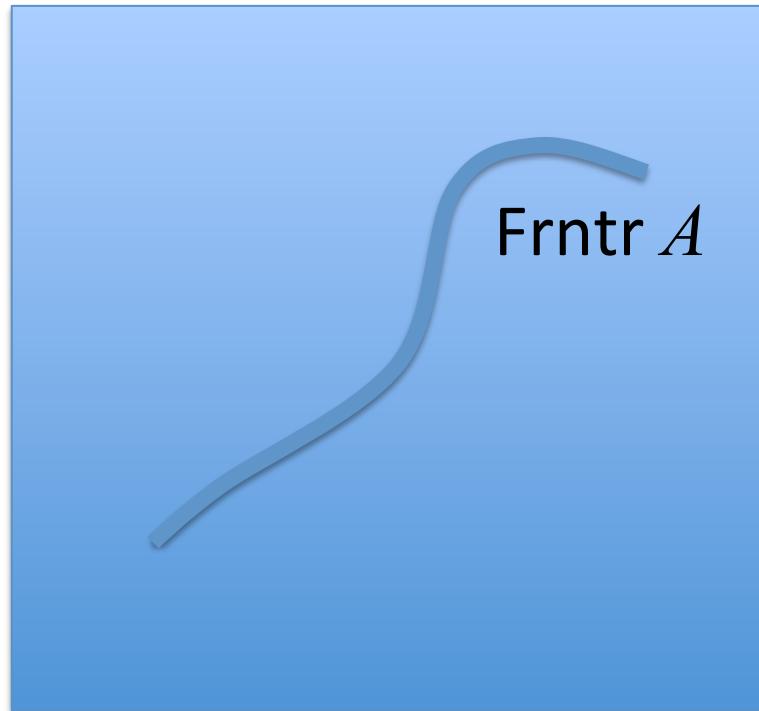
Frontier of A

$\text{Frntr } A = A$ is false, but will never be refuted.



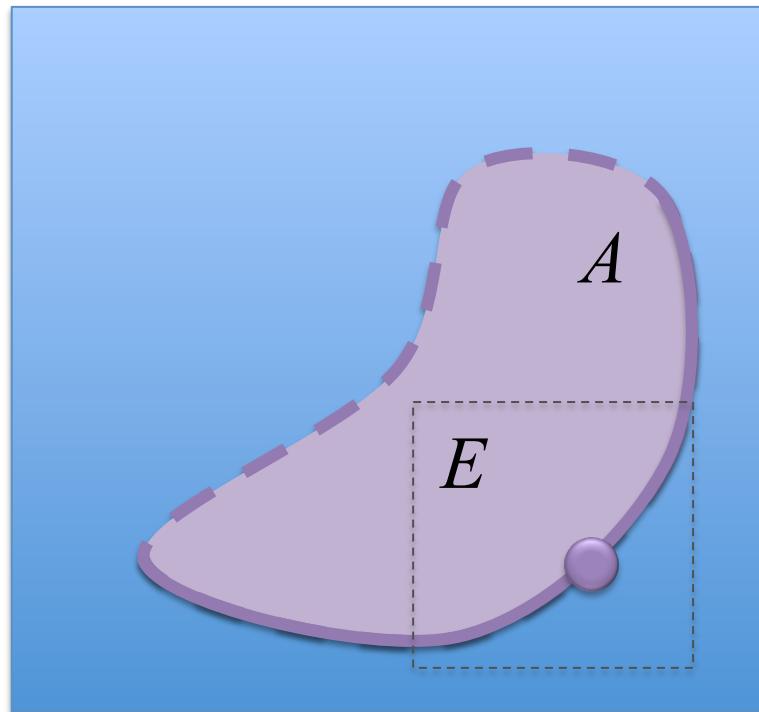
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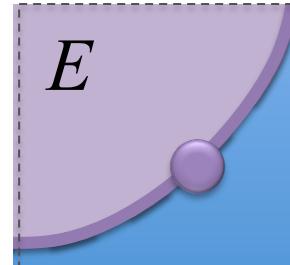
Locally Closed

A is locally closed iff A entails that A will become refutable.



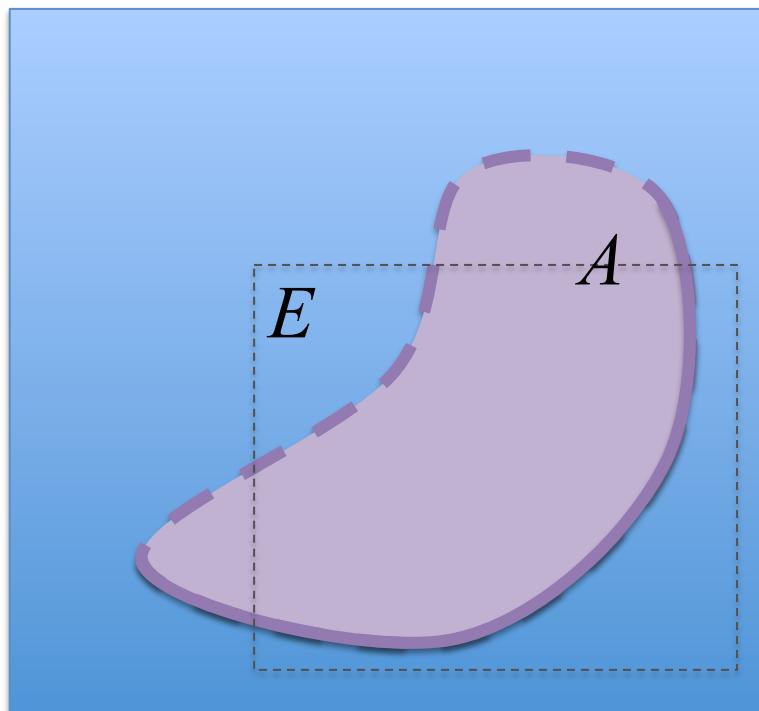
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Local Closure and Default Reasoning

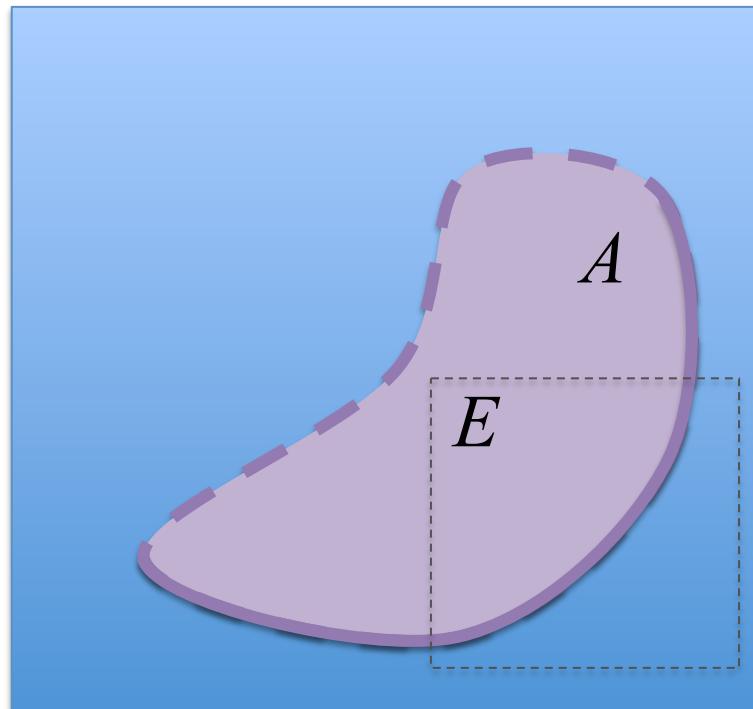
E is a default reason for A iff
 A is refutable but not refuted given E .



E is not a
default reason
for A .

Locally Closed

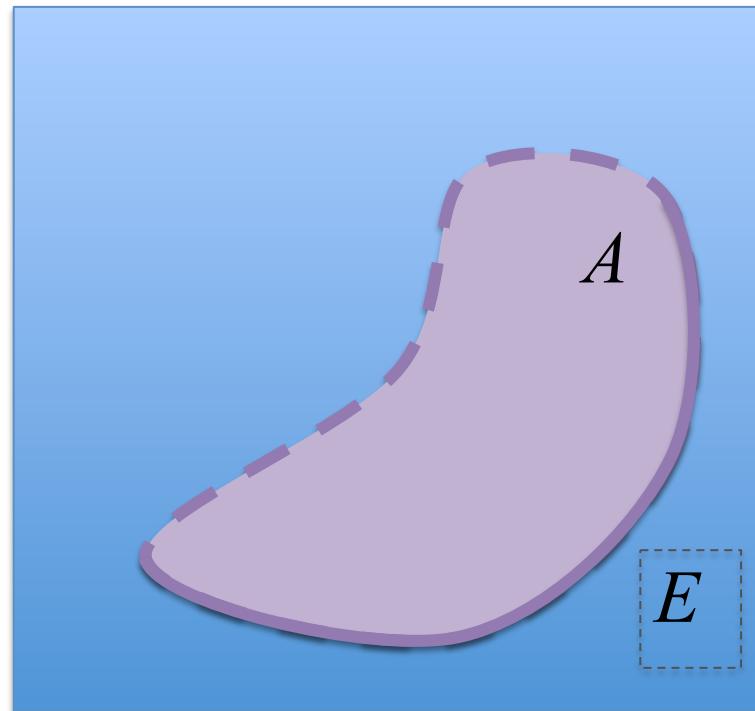
E is a default reason for A iff
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E is a default
reason for A .

Locally Closed

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E refutes A .

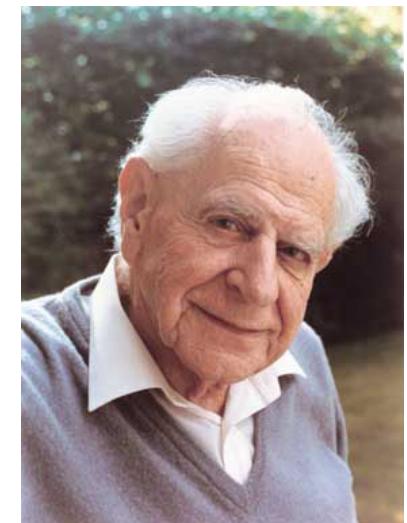
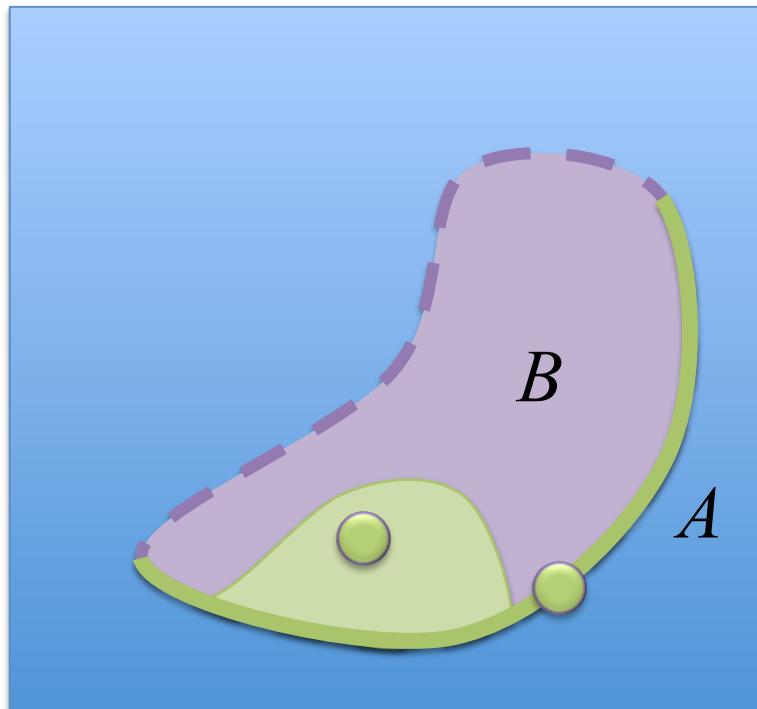


3. EMPIRICAL SIMPLICITY

Simplicity = Specialization Preorder

$A \preceq B$ iff $A \subseteq \text{cl}B$

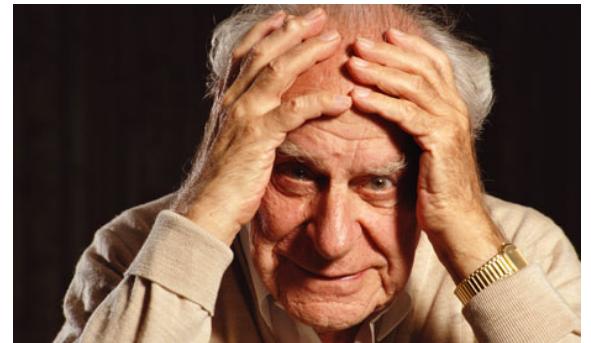
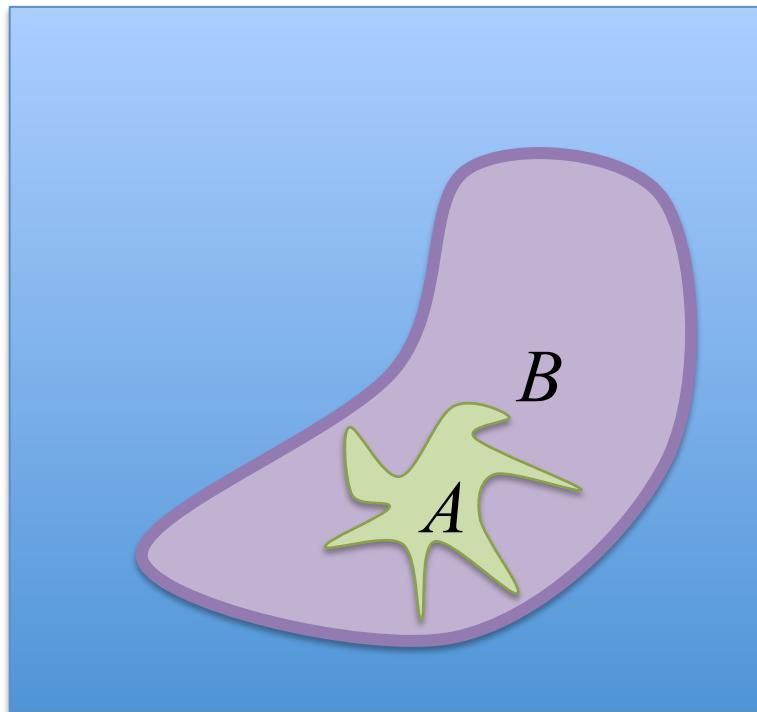
iff all information compatible with A is compatible with B
iff all information refuting B also refutes A .



Sir Karl Popper

The “Tack-on” Objection

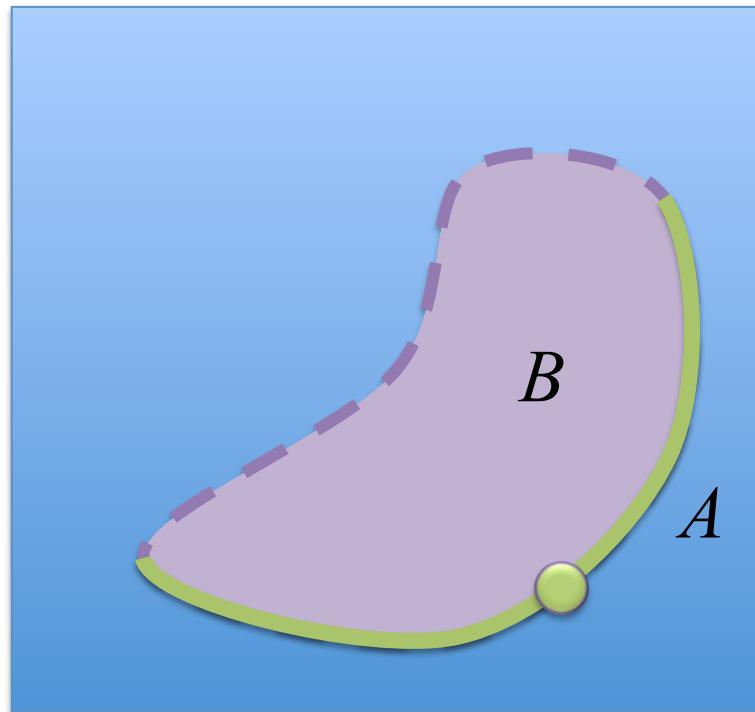
- Adding complex principles to a simple theory doesn't make it simpler (Glymour 1980).



Empirical Simplicity

$A \triangleleft B$ iff $A \cap \text{Frntr}B \neq \emptyset$.

iff A is consistent with: B is false,
but will never be refuted.

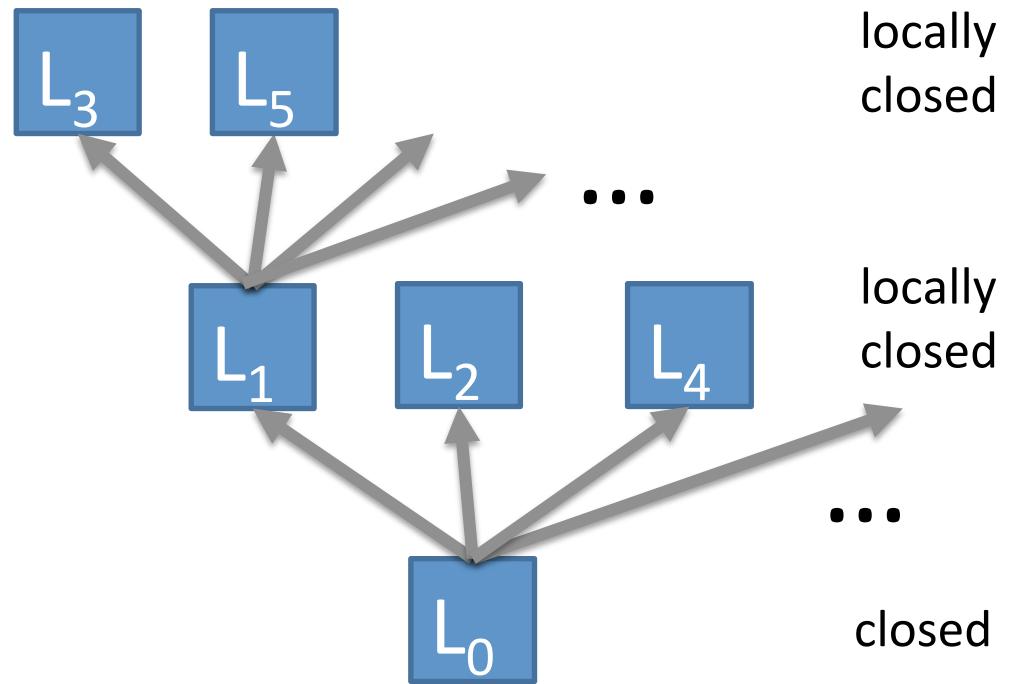


Example: Language Learning (Gold)

\mathcal{Q} = Which language?

\mathcal{I} = finitely many grammatical sentences.

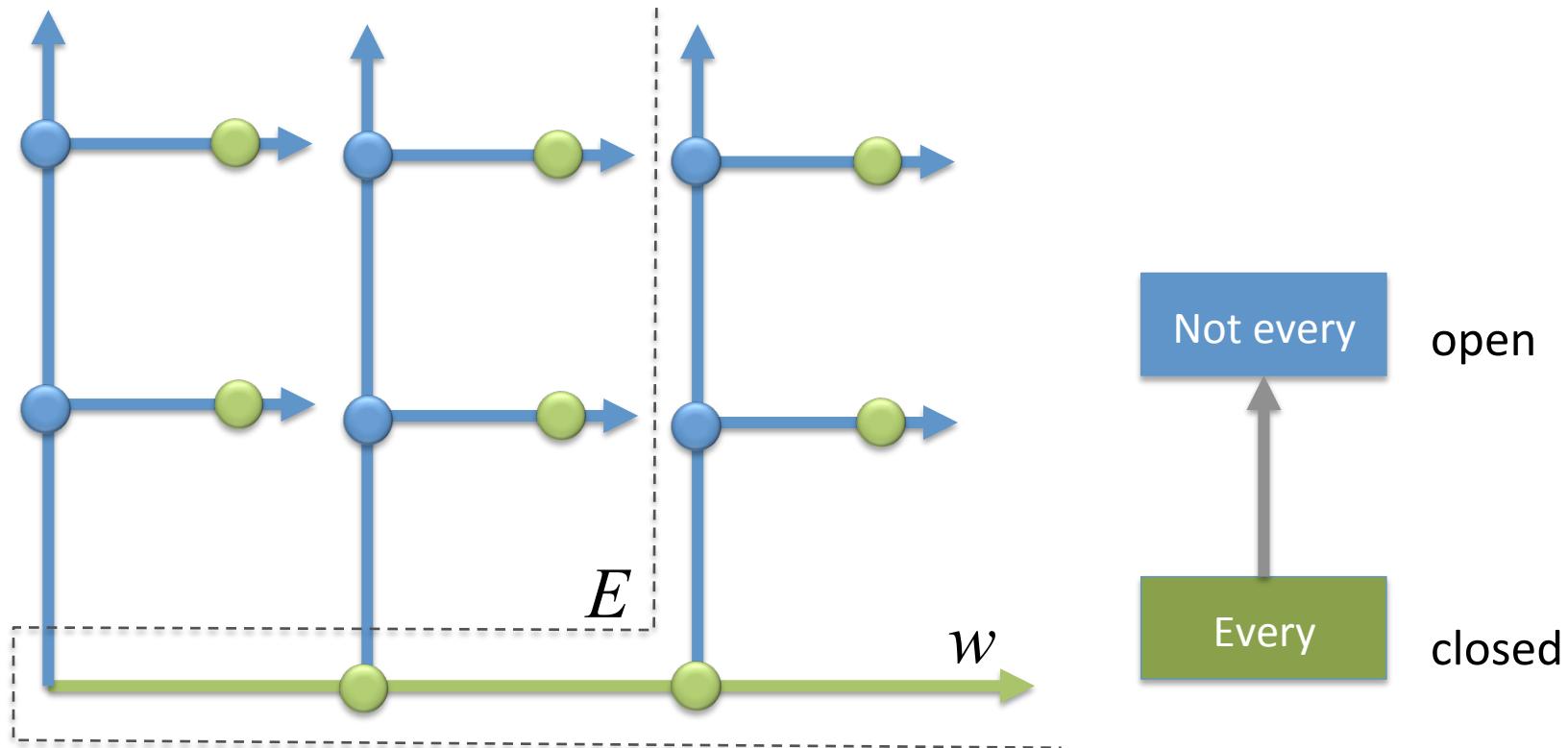
s_1, s_2, s_3, \dots



Example: The Baire Space

\mathcal{Q} = Will every outcome be green?

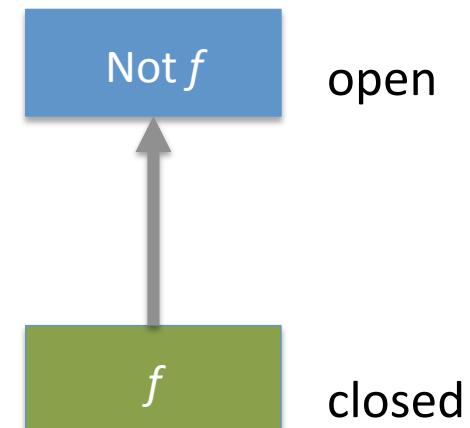
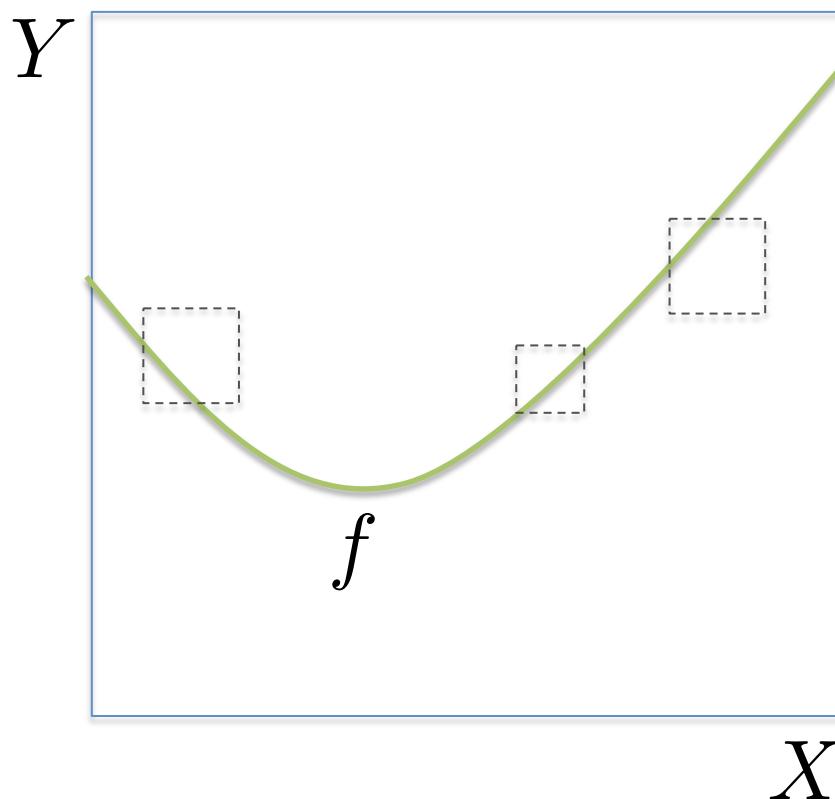
\mathcal{I} = observation histories.



Example: Continuous Laws

\mathcal{Q} = Does $Y = f(X)$?.

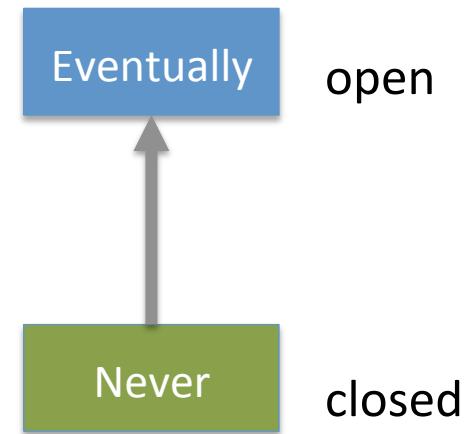
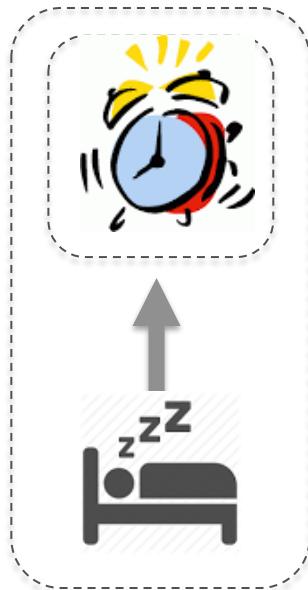
\mathcal{I} = finitely many inexact measurements.



Example: Sierpinski Space

Q = Will it ever ring?

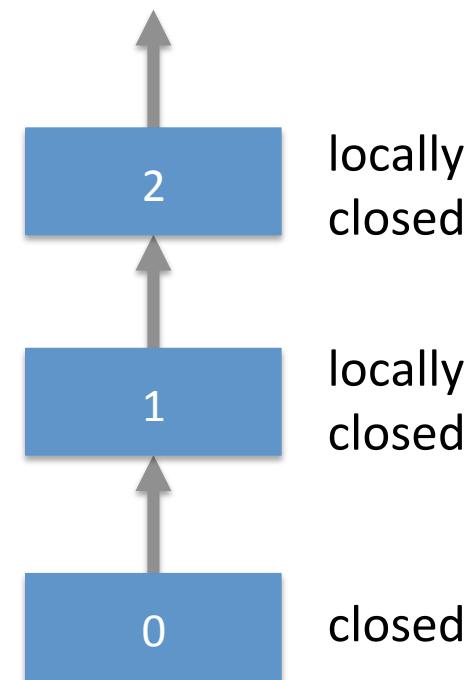
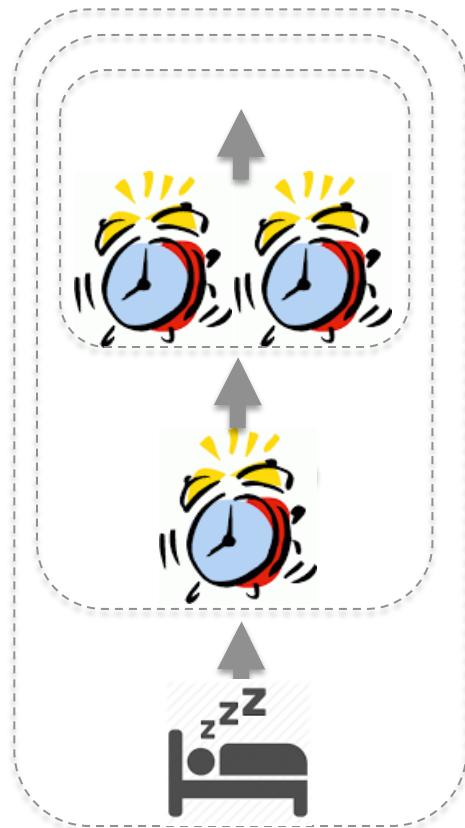
I = alarm or no alarm yet.



Example: Upward Topology

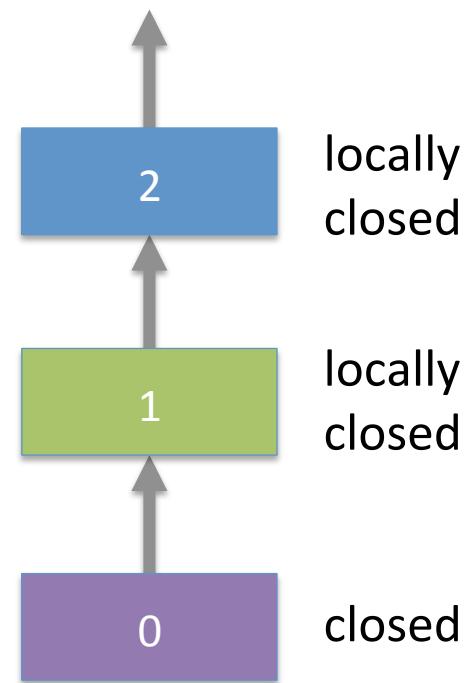
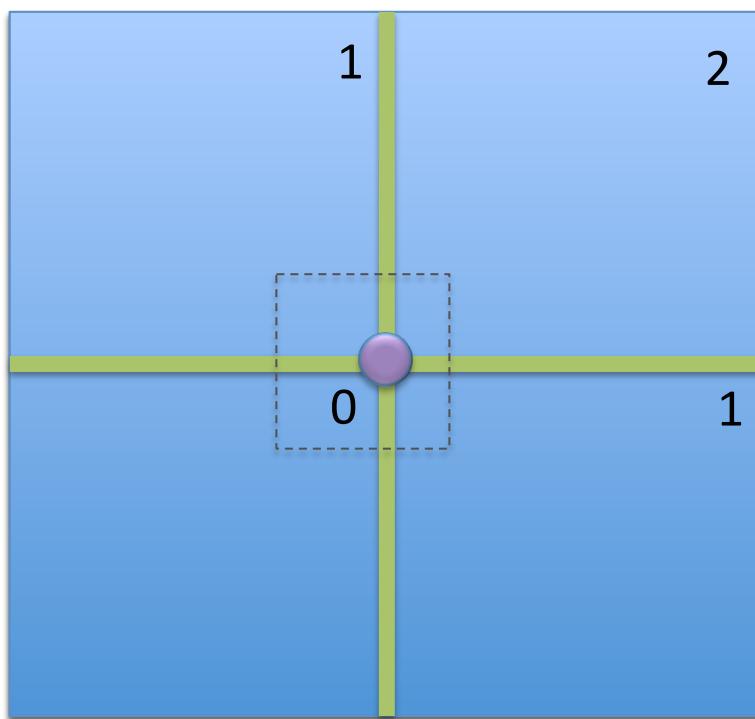
Q = How many alarms?

\mathcal{I} = cumulative alarms.



Example: Euclidean Metric Topology

$Q =$ How many parameters are free?



Example: Quantitative Laws

\mathcal{Q} = What is the true polynomial degree?

\mathcal{I} = finitely many inexact measurements.

$$Y = \sum_{i=0}^N a_i X^i.$$

degree



$$Y = a_0.$$

$$Y = a_0 + a_1 X.$$

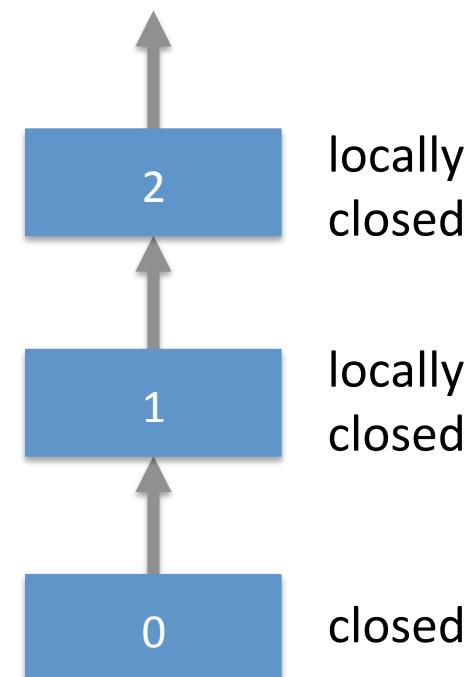
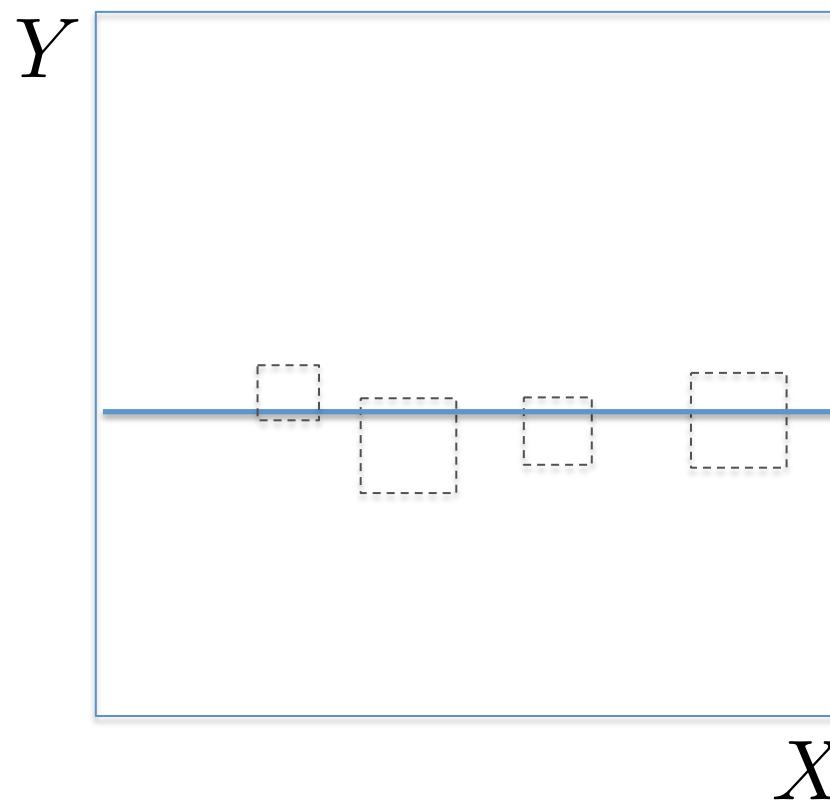
$$Y = a_0 + a_1 X + a_2 X^2.$$

⋮

Example: Quantitative Laws

Q = What is the true polynomial degree?

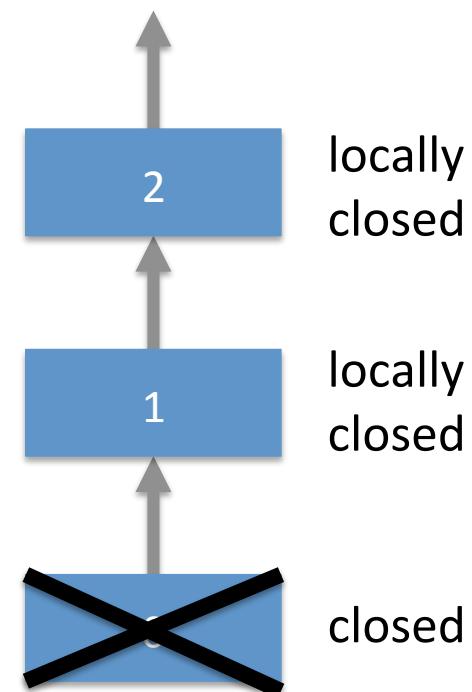
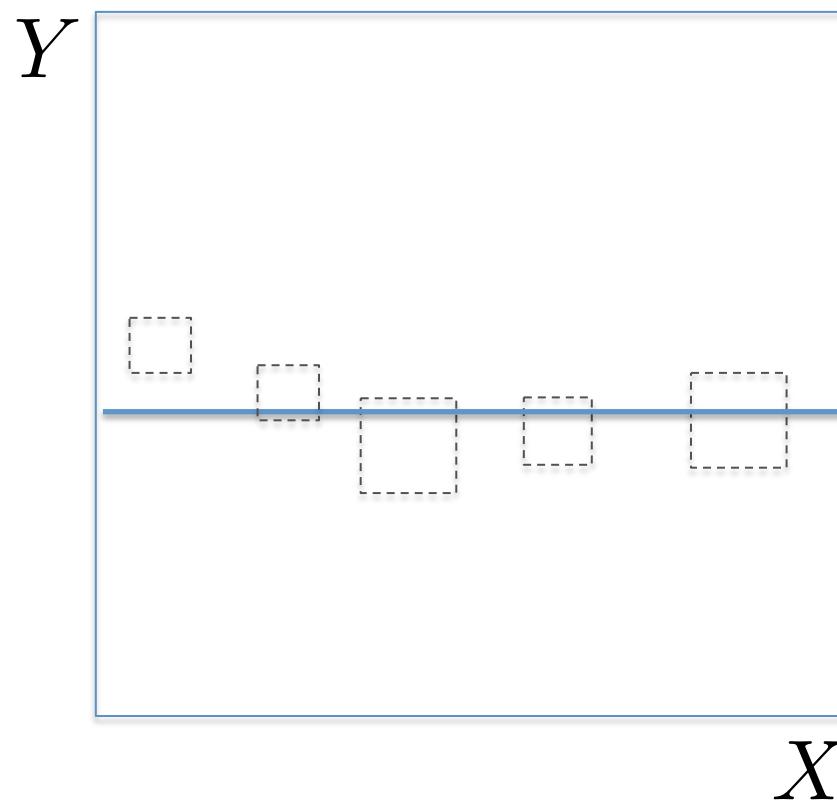
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Example: Quantitative Laws

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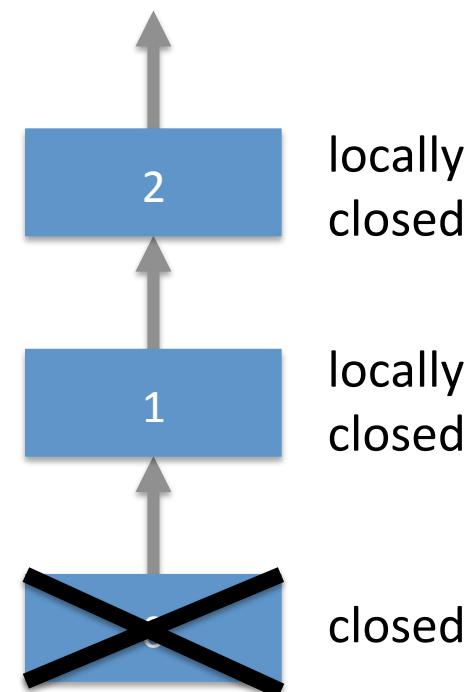
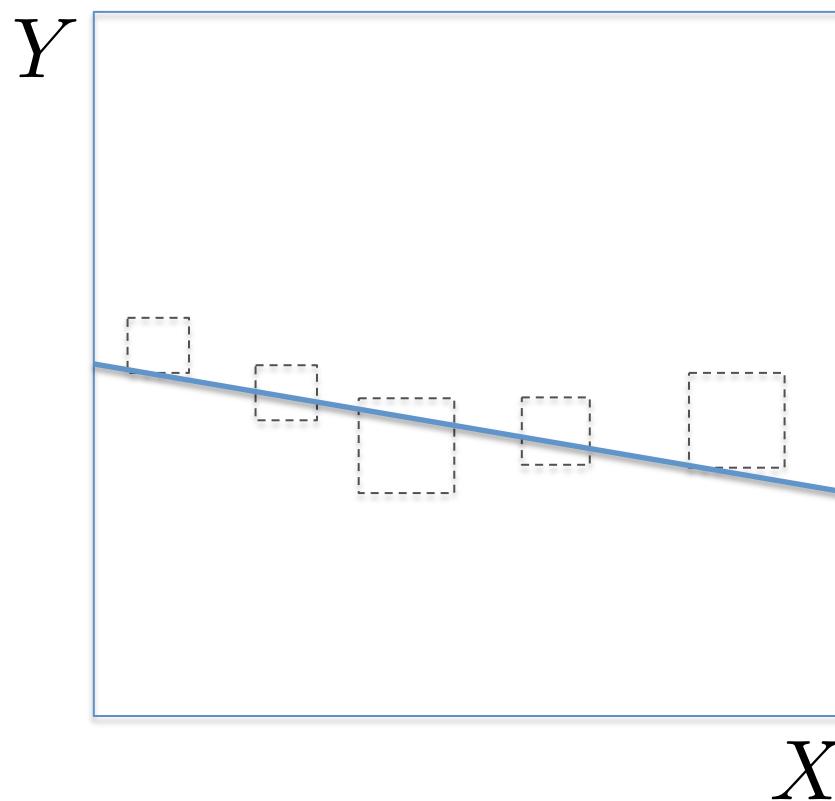
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Example: Quantitative Laws

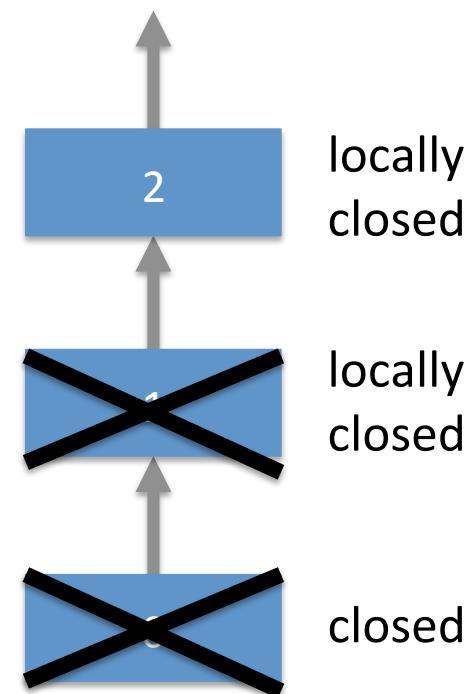
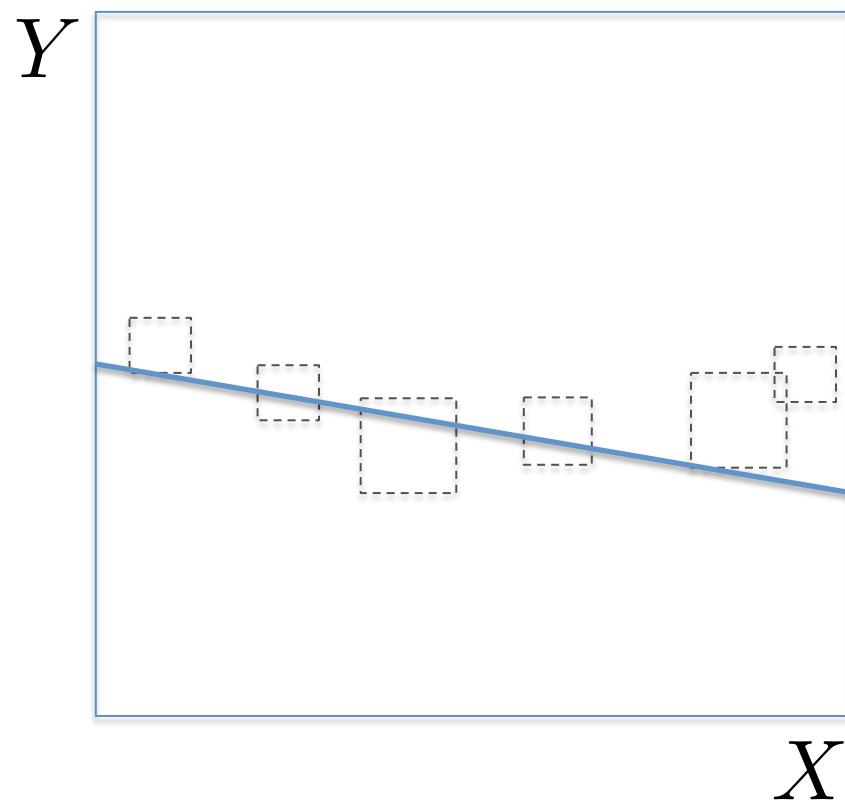
\mathcal{Q} = What is the true polynomial degree?

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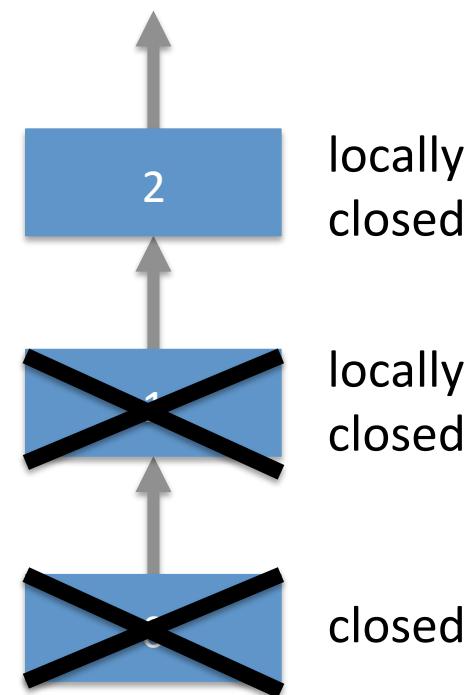
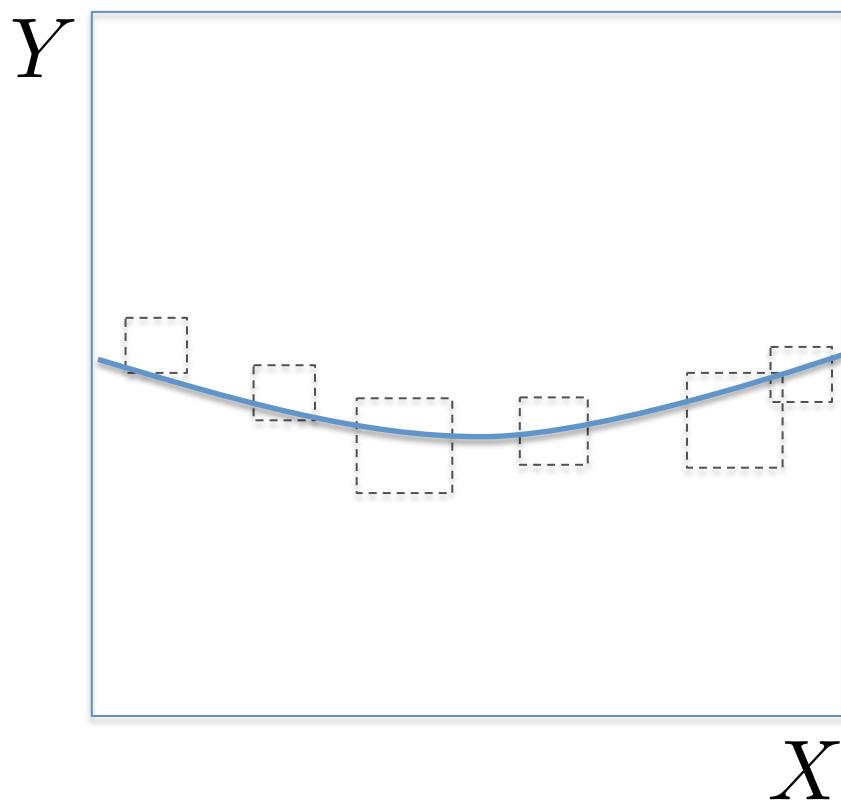
Example: Quantitative Laws

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Example: Competing Paradigms

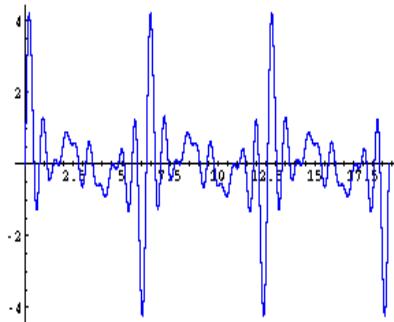
Polynomial paradigm

$$Y = \sum_{i=0}^N a_i X^i.$$



Trigonometric polynomial paradigm

$$Y = \sum_{i=0}^N a_i \sin(iX) + b_i \cos(iX).$$



Example: Competing Paradigms

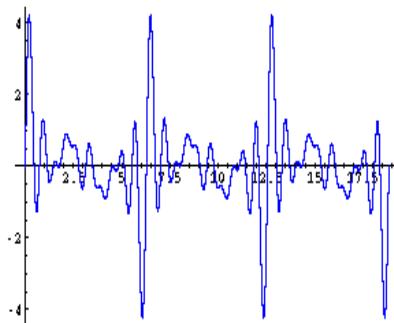
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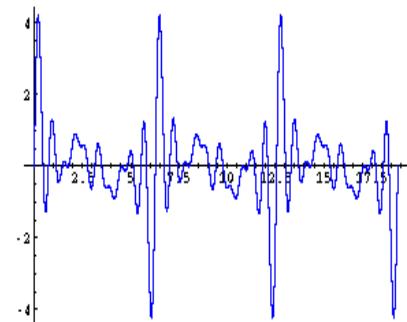
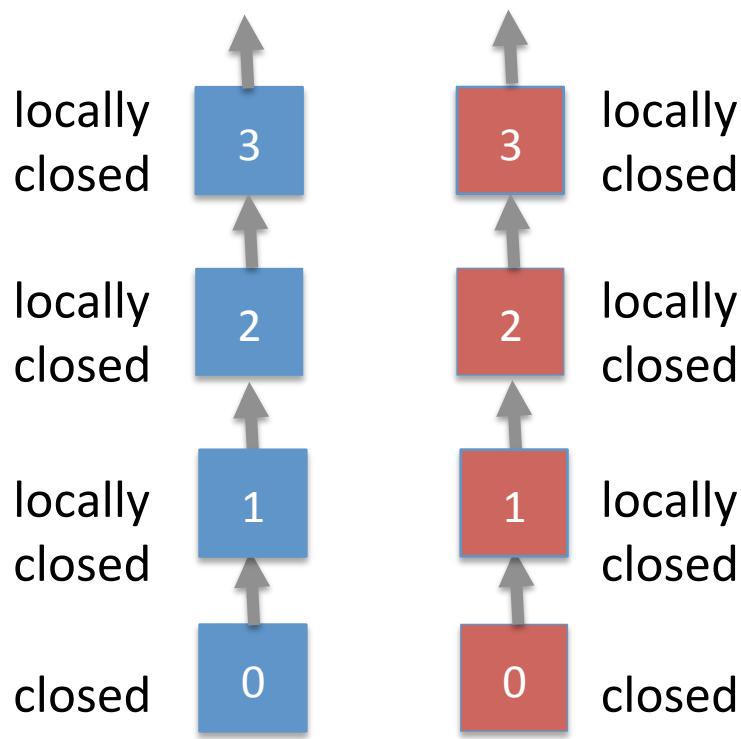
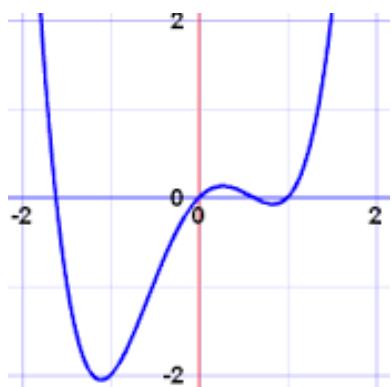
$$Y = \sum_{i=0}^N a_i \sin(iX) + b_i \cos(iX).$$



Example: Competing Paradigms

\mathcal{Q} = which degree and which paradigm is true?

\mathcal{I} = finitely many inexact measurements.



4. INDUCTIVE METHODS

Reasoning

Deductive

- Monotonic
- Non-ampliative



Reasoning

Deductive

- Monotonic
- Non-ampliative

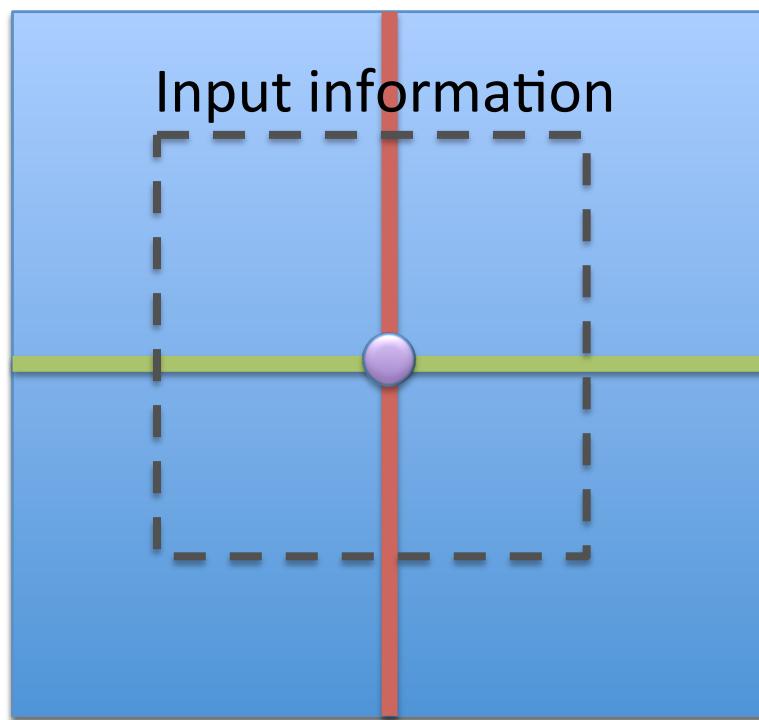


Inductive

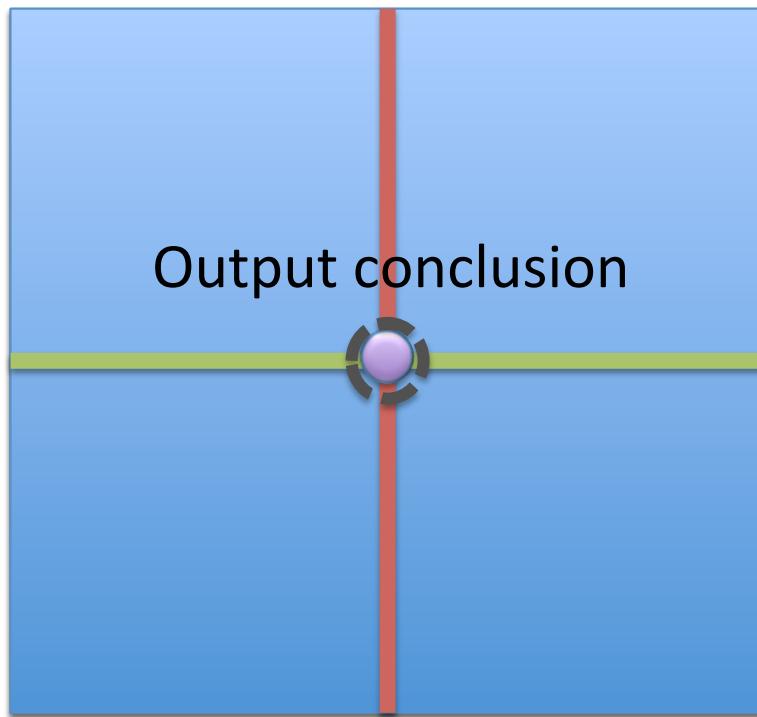
- Non-monotonic
- Ampliative



Inductive Inference

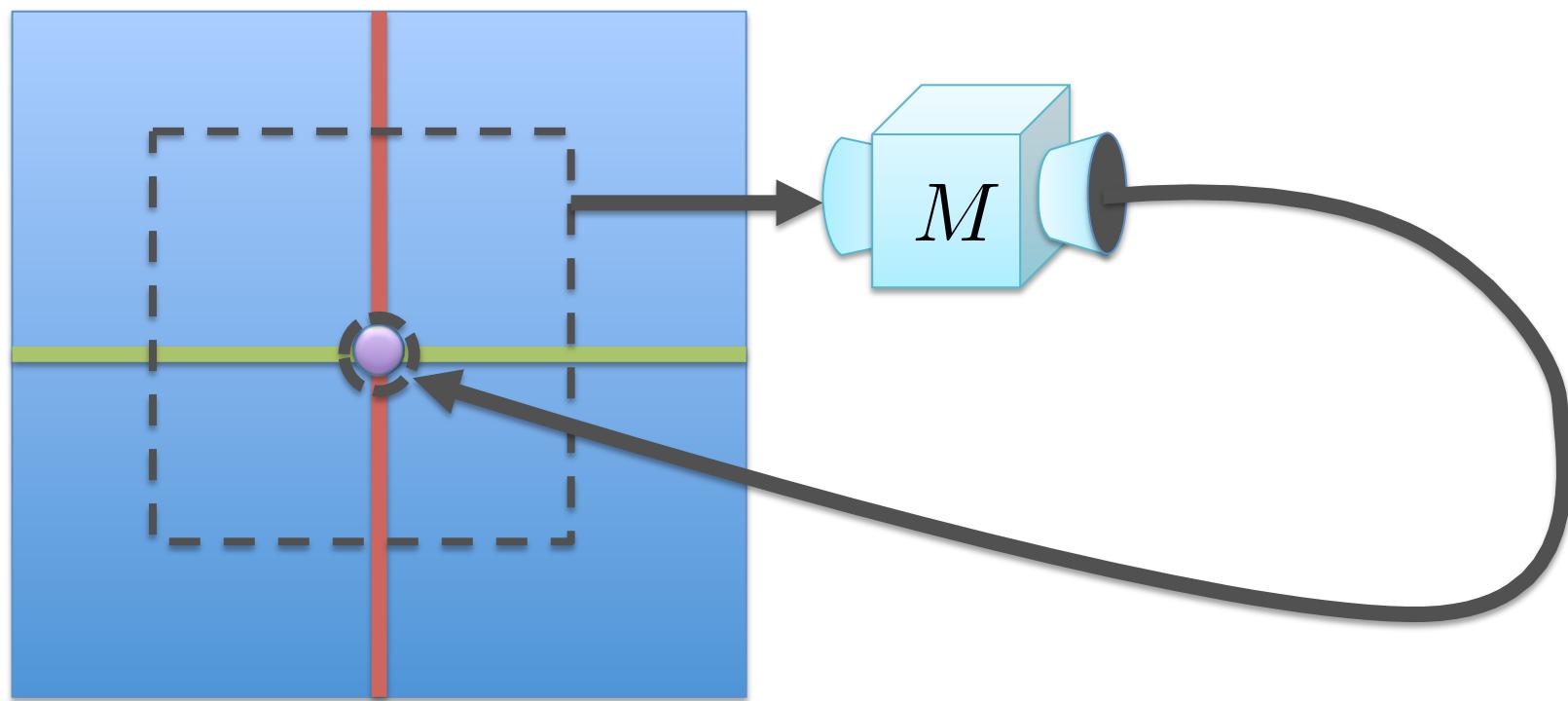


Inductive Inference



Inductive Methods

Information in, relevant response out.

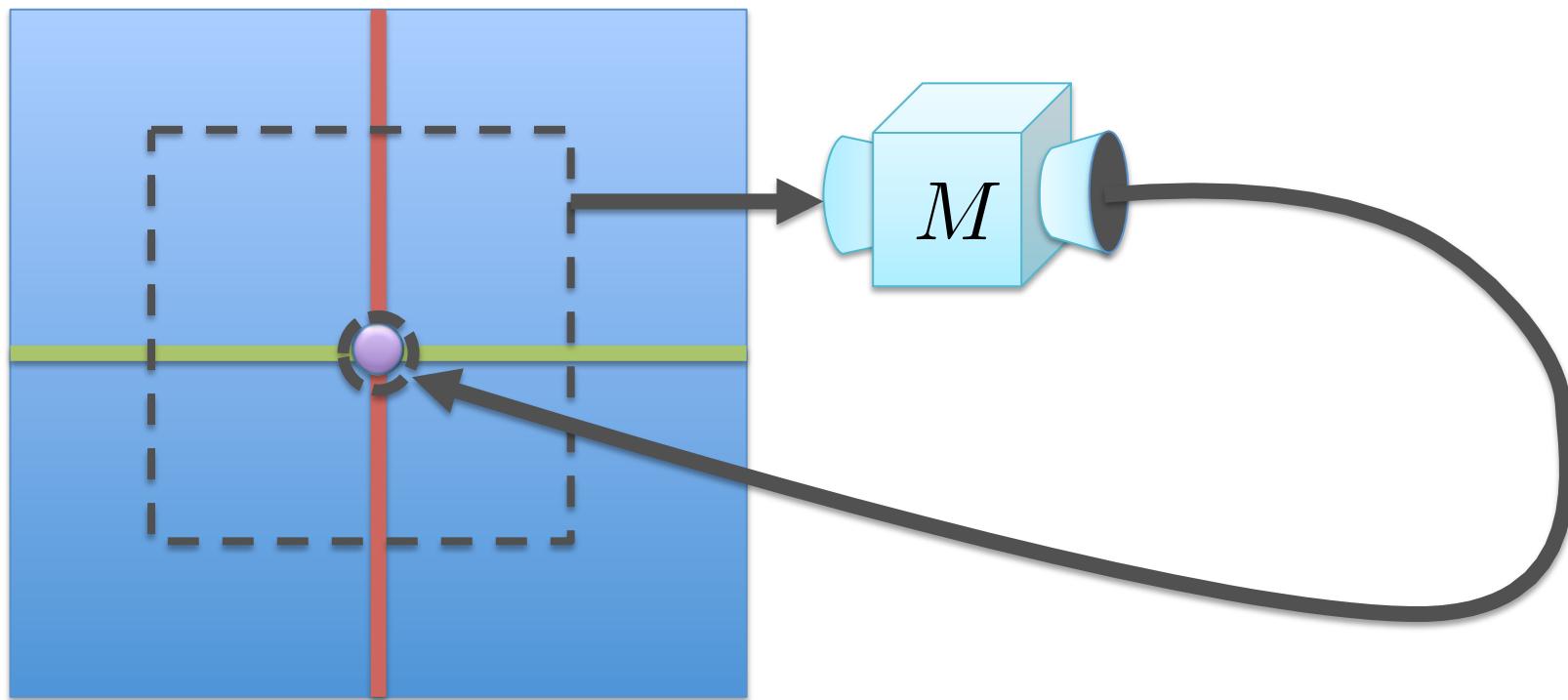


Solution in the Limit

M solves a problem in the limit iff

each world w presents information such that

M produces the true answer in w on any further information in w .

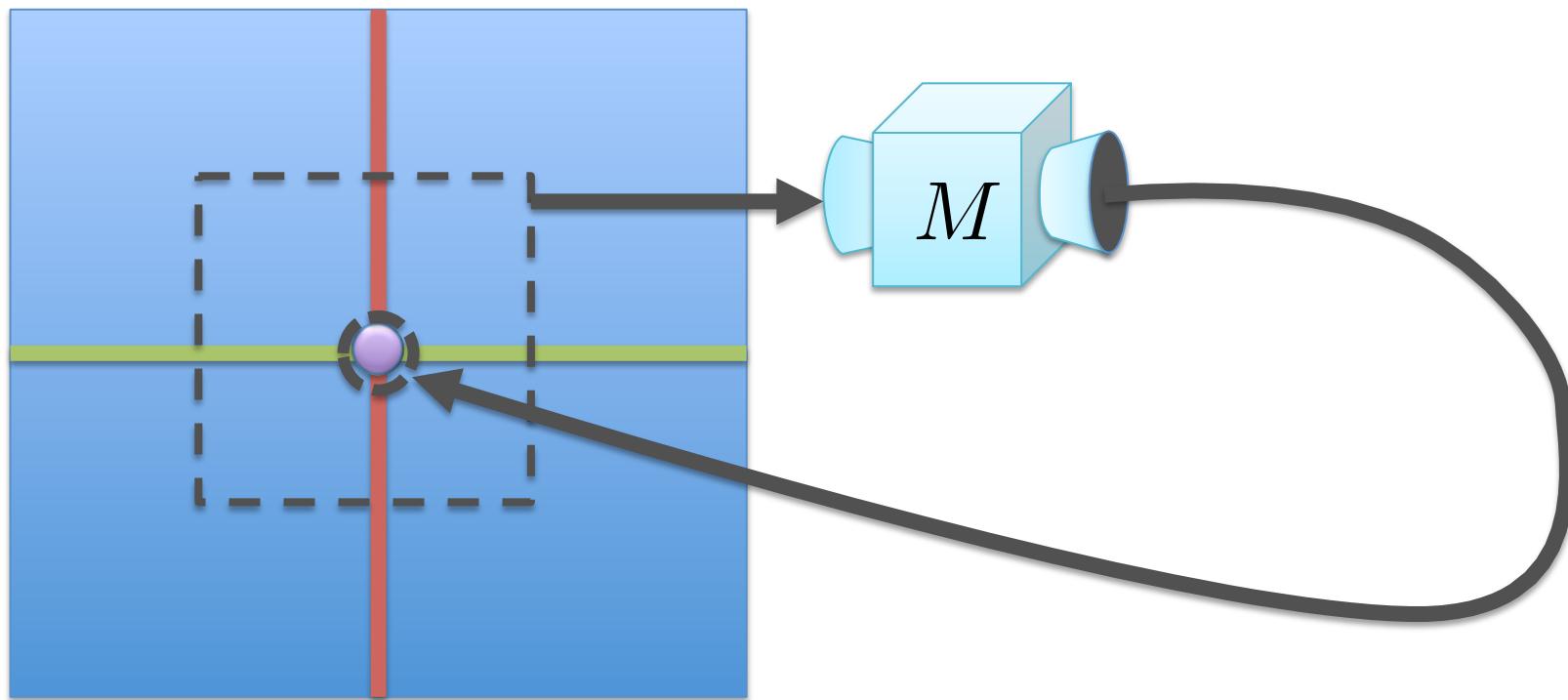


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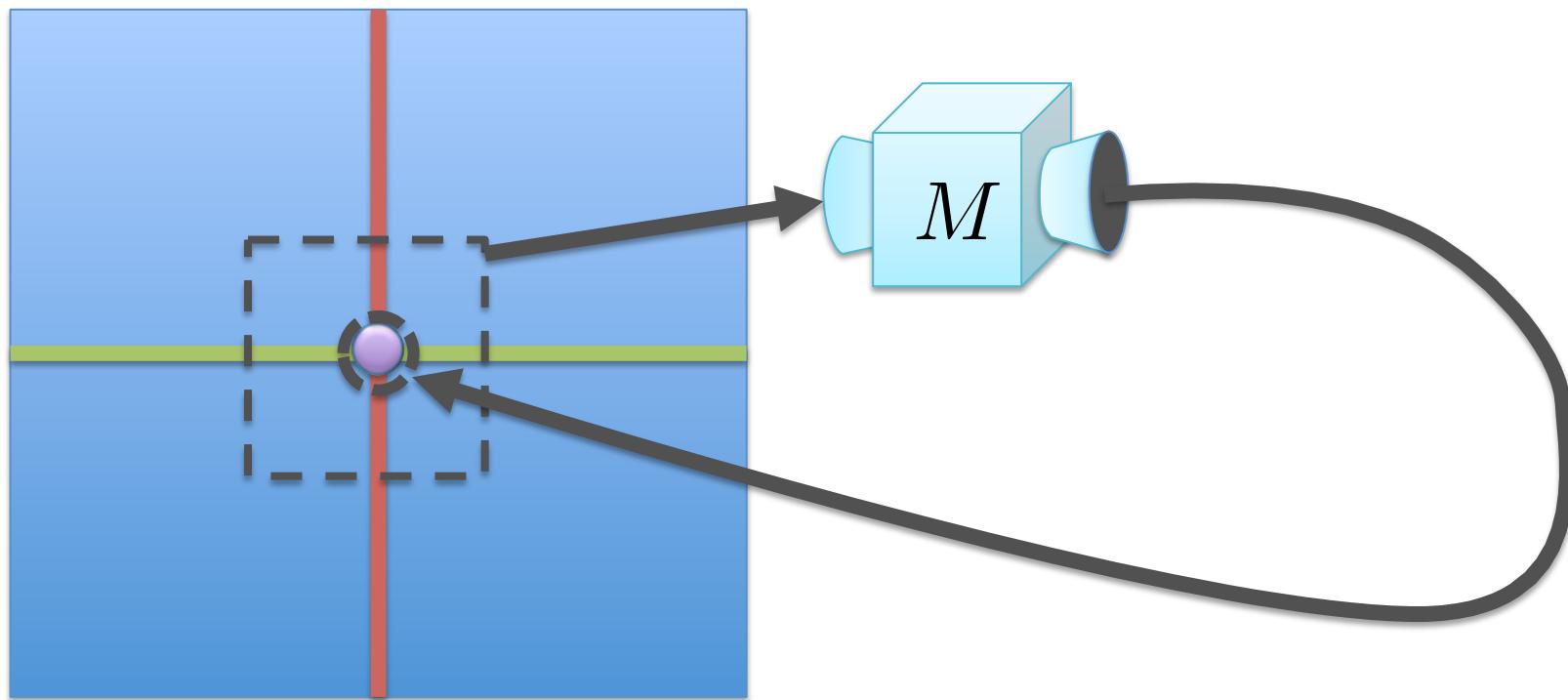


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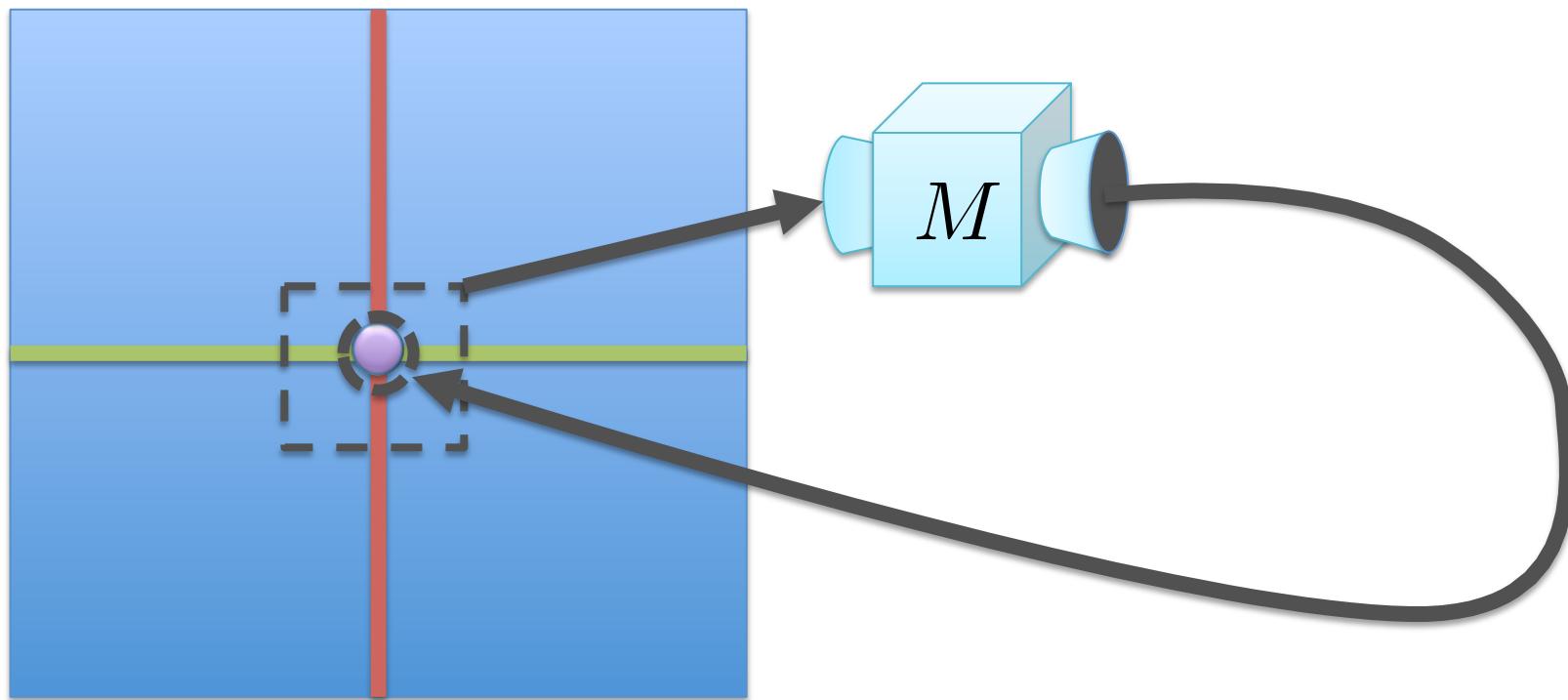


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Reliability and Knowledge

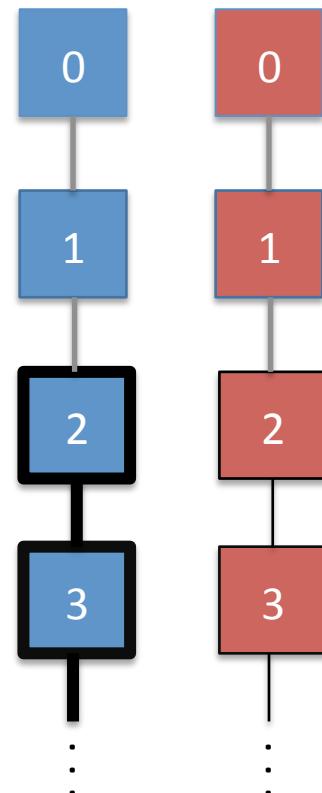
If **knowing** the answer to a scientific question entails that your method is a **solution**, then whether you know the answer depends intrinsically on the **question**.



5. OCKHAM'S RAZOR

Ockham's Razor

- Output a **simplest relevant response** given E .
 - Allows for **suspension of judgment**.
 - Makes sense for **infinite descending chains**.



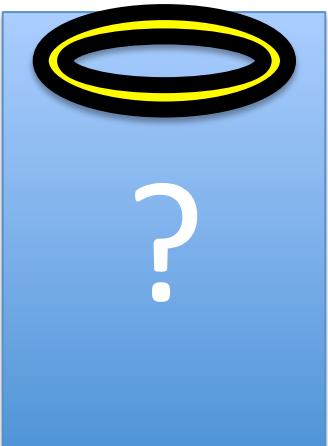
Popper's Razor

- Output a relevant response that is **refutable** (**closed**) given E .



Error Razor

- “Err on the side of simplicity”.
- In arbitrary world w , never produce a relevant response B such that the true answer A_w is strictly simpler than B .



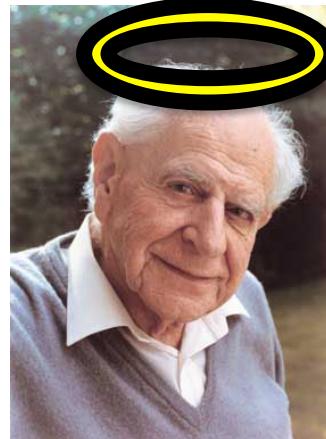
Equivalence

Proposition.

Ockham's razor = Popper's razor = error razor.



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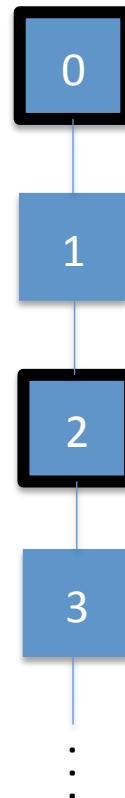


Patience

- Never rule out a simplest relevant response given E .
 - Says that **simplicity** is the **only** reason for inductive leaps beyond experience.
 - Logically independent of Ockham's razor.



Patient but
not Ockham



Patience

- Never rule out a simplest relevant response given E .
 - Says that **simplicity** is the **only** reason for inductive leaps beyond experience.
 - Logically independent of Ockham's razor.



Ockham but
not patient



Error patience

- In arbitrary world w , never output a relevant response that rules out all answers as simple as A_w .



Equivalence

- **Proposition:** Error patience is equivalent to patience.



6. OCKHAM'S RAZOR JUSTIFIED

Reliability

Deductive

- Converge to the truth
directly



Reliability

Deductive

- Converge to the truth
directly



Inductive

- Converge to the truth
indirectly



Goldilocks Philosophy



Straight



Too strong!



Just right?

Arbitrarily crooked



Too weak!

Straightest Possible Convergence

Straight



Too strong!

Couldn't be straighter



Just right!

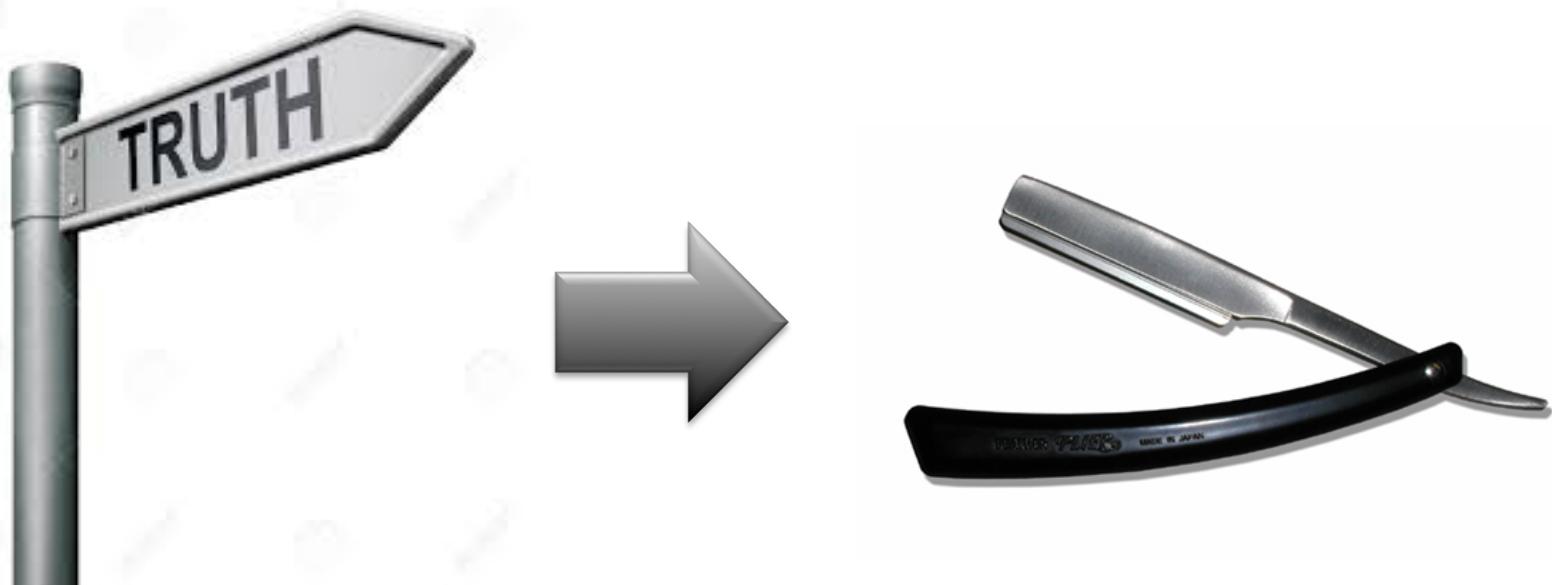
Arbitrarily crooked



Too weak!

Thesis

- Ockham's razor is **necessary** for **straightest** convergence to the truth.



The Straightest Path

Do not fear for me.
Make straight your
own path to
destiny.

Sophocles

meetville.com

أَهْدِنَا الصِّرَاطَ الْمُسْتَقِيمَ ①

GUIDE US TO THE STRAIGHT PATH

AL-FATIHAH : 6



sperience

Trust in the LORD with all your heart
and lean not on your own understanding;
in all your ways acknowledge him,
and he will make your paths straight.

Proverbs 3:5-6

Fools dwelling in darkness, but
thinking themselves wise, go round
and round, by tortuous paths, like the
blind led by the blind.

Katha Upanishad

Two Departures from Straightness



Course-reversals



Cycles

Doxastic Reversal Sequence

- A finite sequence of relevant responses in which each entry **contradicts** its predecessor.



Doxastic Cycle Sequences

- A reversal sequence whose terminal entry entails its first entry.



Straightest Convergence

- M produces reversal [cycle] sequence

$$\mathbf{s} = (R_1, \dots, R_k)$$

iff there exist information states

$$\mathbf{e} = (E_1 \supset \dots \supset E_k)$$

such that

$$M(\mathbf{e}) = (M(E_1), \dots, M(E_k))$$

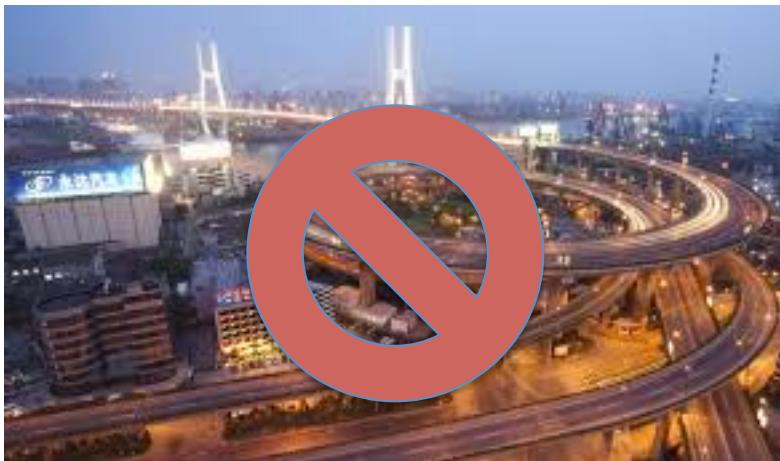
is a reversal [cycle] sequence such that $M(E_i) \subseteq R_i$,
for i from 1 to k .

Straightest Convergence

- Solution M is reversal [cycle] optimal iff:
every solution produces each reversal [cycle]
sequence by M .

Main Result 1

- **Proposition (Baltag, Gerasimczuk, and Smets):** Every solvable problem is refinable to a problem with a **cycle-free** solution.

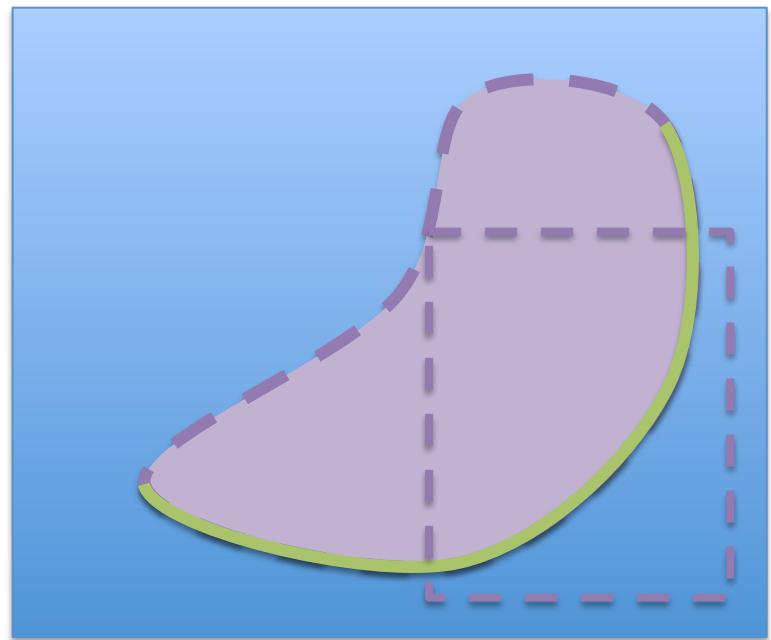


Main Result 2

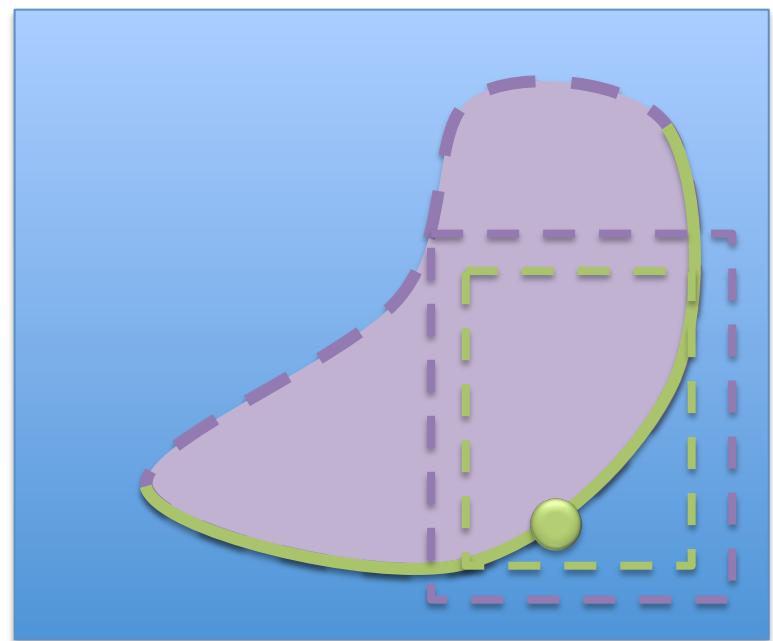
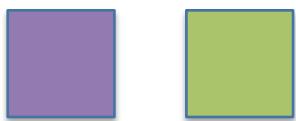
- **Proposition:** Every *cycle-free* solution satisfies Ockham's razor.



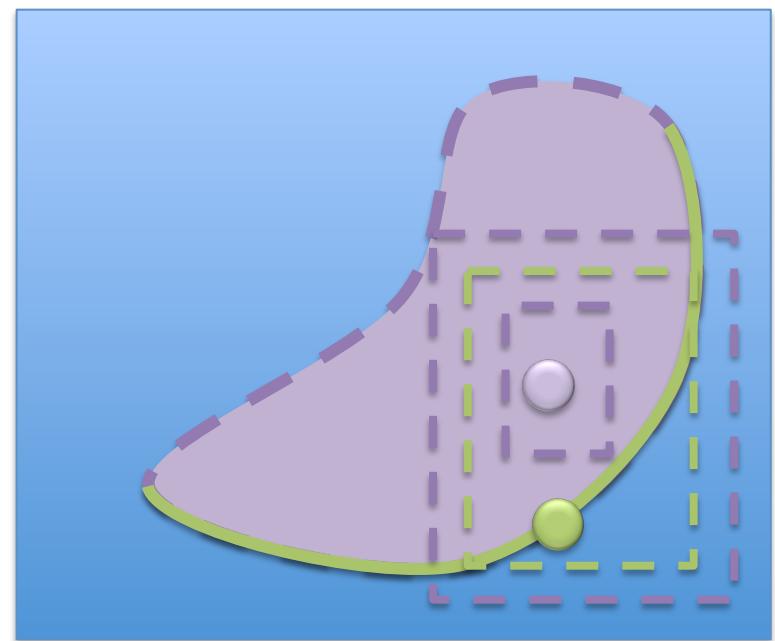
The Idea



The Idea



The Idea



Main Result 3

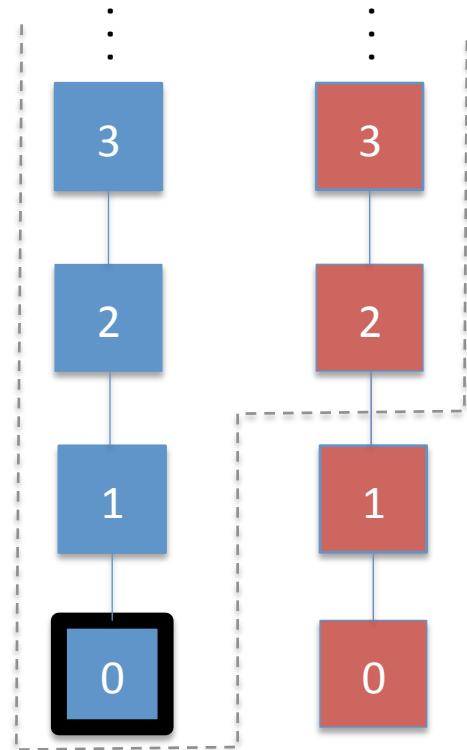
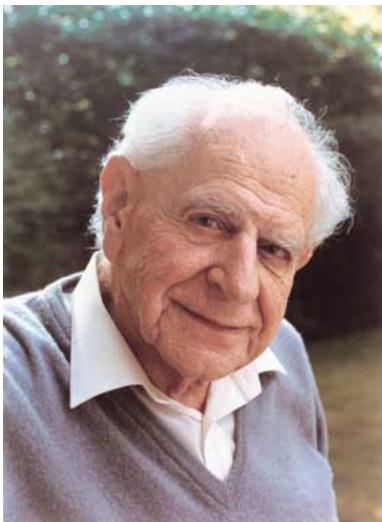
- We can **characterize** the solvable problems that have reversal optimal solutions.
- The characterization depends on the basis, so it is **not topological**.

Main Result 3

- We can **characterize** the solvable problems that have reversal optimal solutions.

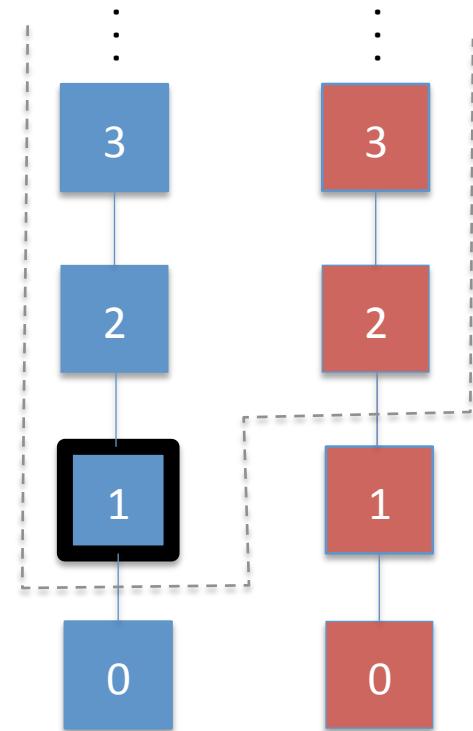
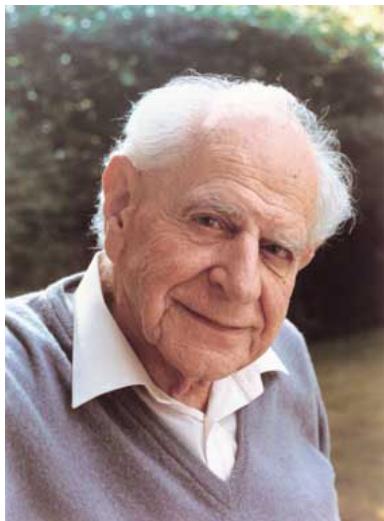
No Reversal Optimal Solution

“Popper”:
choose the
paradigm
with **fewer**
free
parameters.



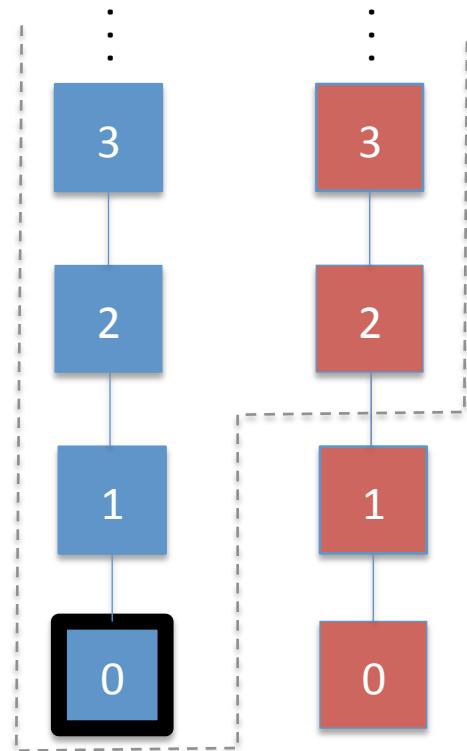
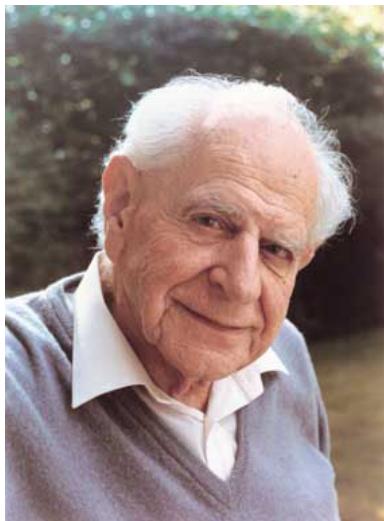
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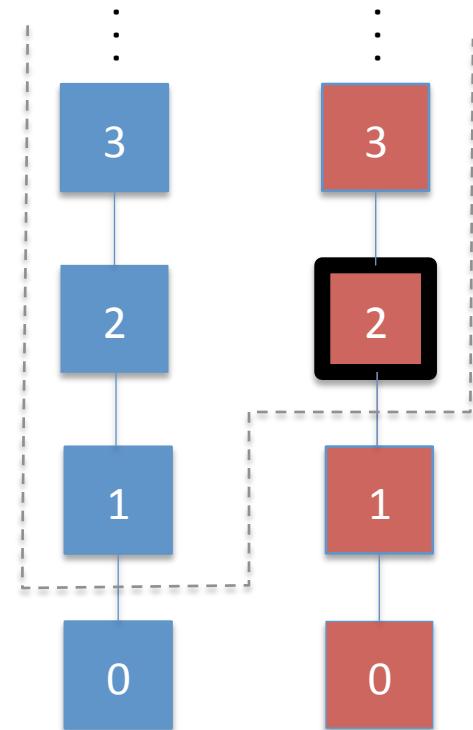
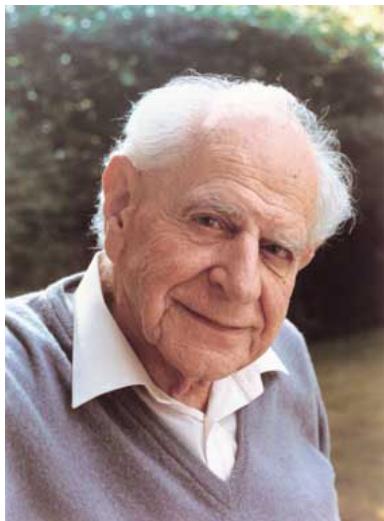


Lakatos:
choose the
paradigm that
was **adjusted**
least recently.



No Reversal Optimal Solution

“Popper”:
choose the
paradigm
with **fewer**
free
parameters.



Lakatos:
choose the
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Main Result 4

- **Proposition:** a solution is reversal optimal only if it is patient.



Contextual Justification

- If patience is truth-conducive in **your** problem, its feasibility in some **other** problem is irrelevant.

Summary and Discussion

1. Simplicity is a topological feature of problems.
2. Ockham's razor is necessary for cycle-optimal convergence to the true answer.
3. Patience is necessary for reversal-optimal convergence to the true answer.
4. Optimally straight convergence is weak, but its implications for scientific method are strong.
5. The same holds for statistical inductive inference.
 1. Significance → a small tolerance for reversals and cycles.
 2. Power → drop theories you are destined to drop sooner.