



# A Topological Explication of Empirical Simplicity

{Kevin T. Kelly, Konstantin Genin}

# Ockham's Razor

- “Presume no more complexity than necessary.”



# Three Fundamental Questions

1. What is **simplicity**?
2. What is **Ockham's razor**?
3. What is its **epistemic justification**?

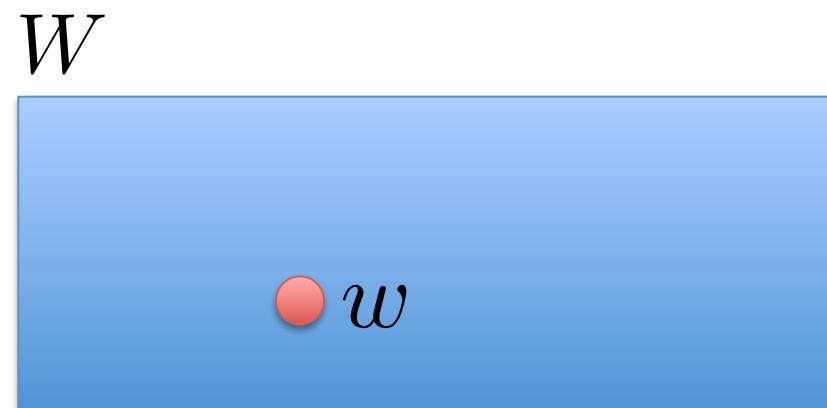




# 1. INFORMATION TOPOLOGY

# Worlds

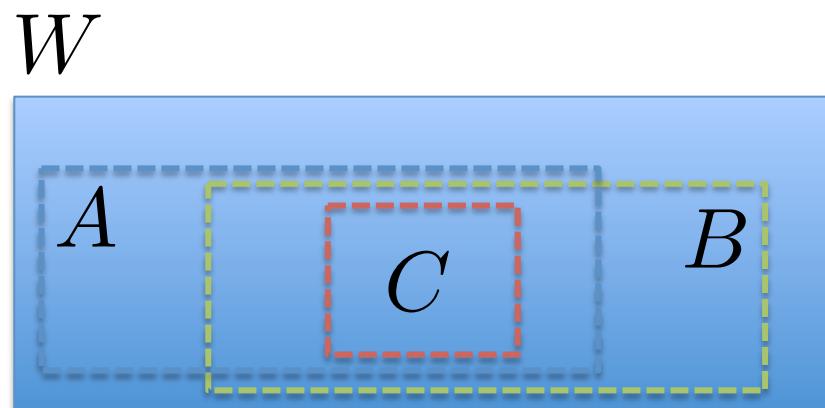
- The points in  $W$  are **possible worlds**.



# The Structure of Information

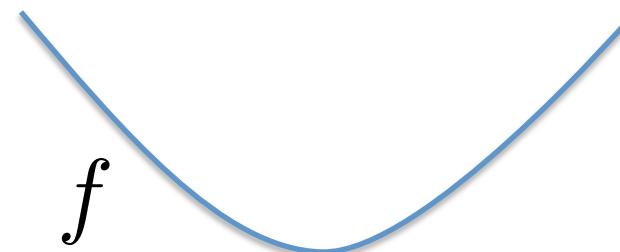
An **information basis**  $\mathcal{I}$  is a **countable** set of information states such that :

1. each world makes **some** information state true;
2. each **consistent pair** of information states is entailed by a stronger information state.



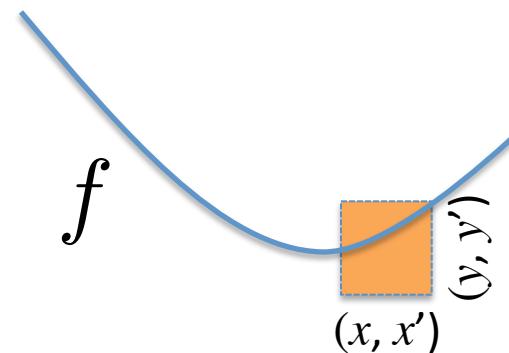
# Example: Equations

- **Worlds** = functions  $f : \mathbb{R} \rightarrow \mathbb{R}$ .



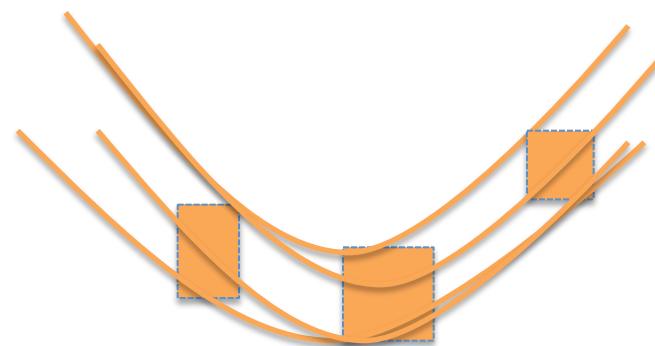
# Example: Possible Laws

- An **observation** is a joint measurement.



# Example: Possible Laws

- The **information state** is the set of all functions that touch each observation.

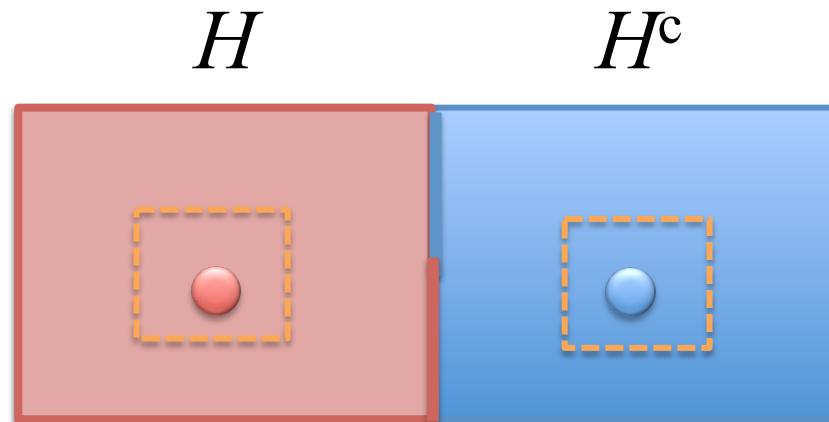


# $H$ will be Verified in $w$

$w$  is an **interior [exterior] point** of  $H$

iff  $H$  **will be** verified [refuted] in  $w$

iff there is  $E \in \mathcal{I}(w)$  s.t.  $H$  is **verified [refuted]** by  $E$ .

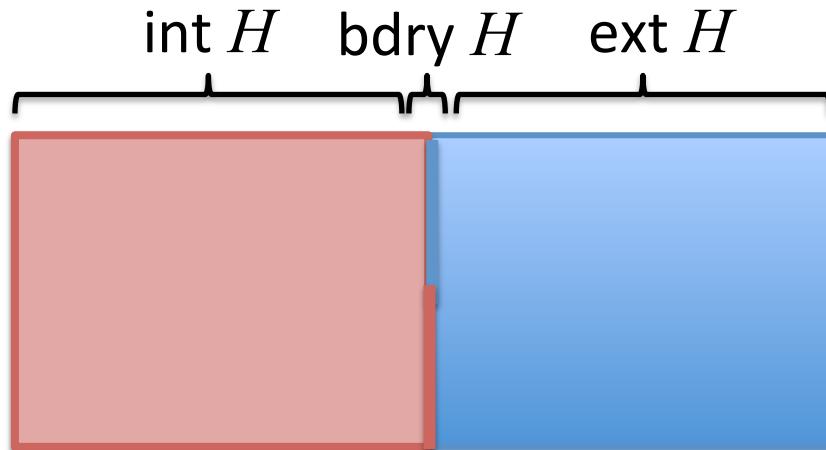


# $H$ will be Verified/Refuted

**int**  $H$  := the proposition that  $H$  **will be verified**.

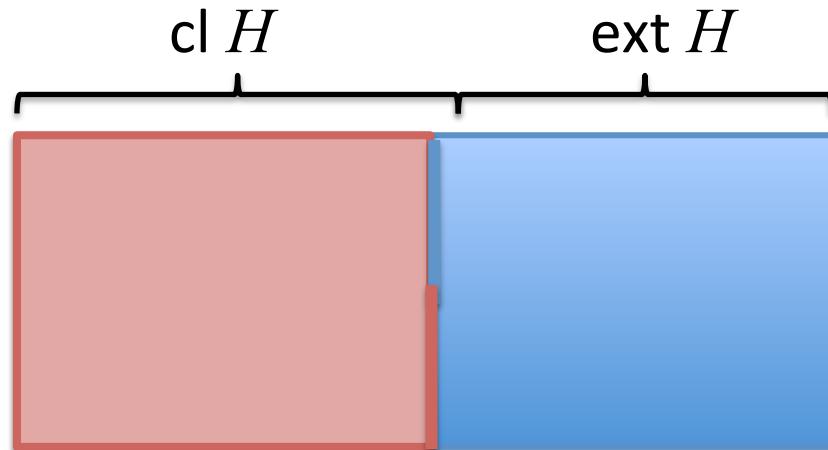
**ext**  $H$  := the proposition that  $H$  **will be refuted**.

**bdry**  $H$  := the proposition that  $H$  **will never be decided**.



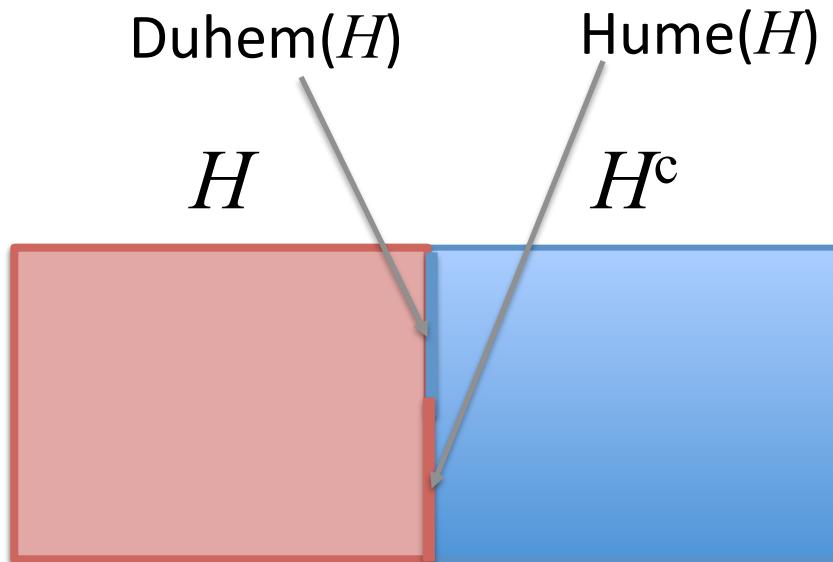
# $H$ will Never be Decided

$\text{cl } H$  := the proposition that  $H$  **will never be refuted**.



# Hume and Duhem

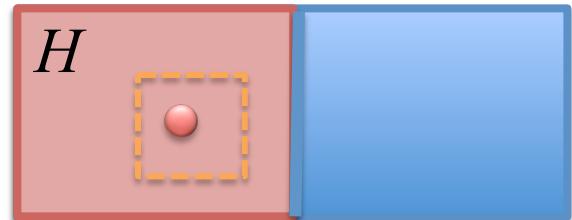
- $\text{bdry}(H) \cap H =$  “you face **Hume’s problem** w.r.t.  $H$ ”;
- $\text{bdry}(H) \cap H^c =$  “you face **Duhem’s problem** w.r.t.  $H$ ”



# Verifiability, Refutability, Decidability

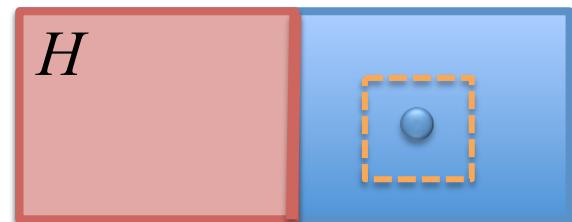
$H$  is **verifiable (open)** iff  $H \subseteq \text{int}(H)$ .

i.e., iff  $H$  will be **verified** however  $H$  is **true**.

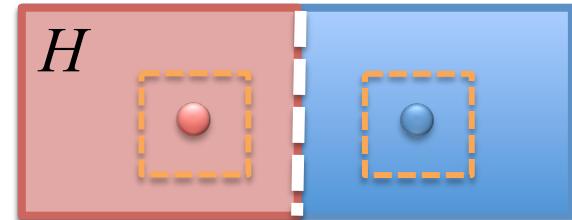


$H$  is **refutable (closed)** iff  $\text{cl}(H) \subseteq H$ .

i.e., iff  $H$  will be **refuted** however  $H$  is **false**.

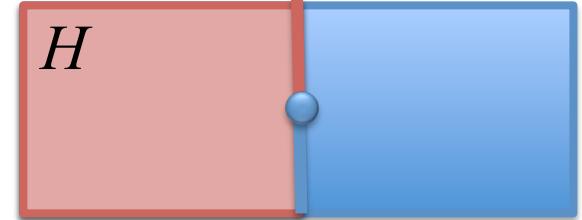


$H$  is **decidable (clopoen)** iff  $H$  is both  
verifiable and refutable.



# Veri-futability

$H$  is **veri-futable (locally closed)** iff  $H$  will be verified to be refutable, however  $H$  is true.



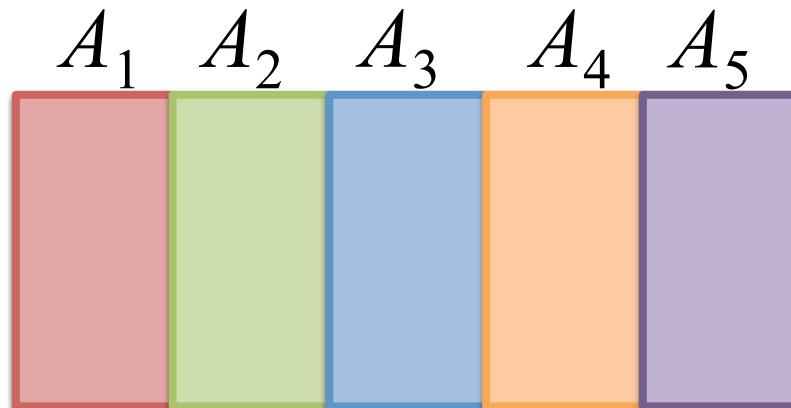
## Scientific models.

- E.g., “linear”, “quadratic”.

## **2. INDUCTIVE METHODS**

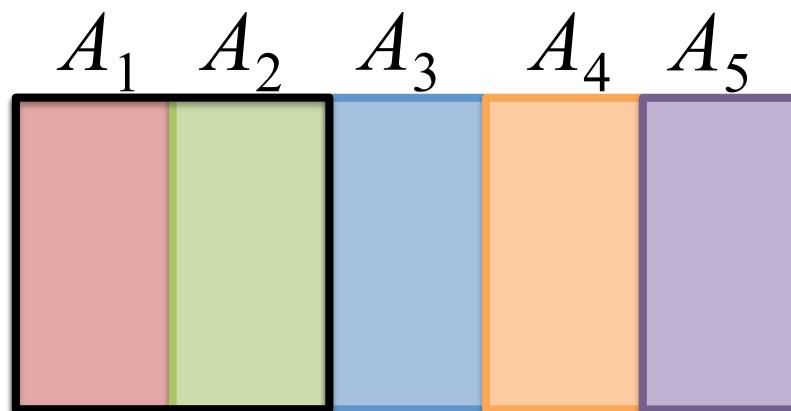
# Questions

- A question  $Q$  partitions  $W$  into a countable set of possible answers.
- Inquiry seeks the true answer.



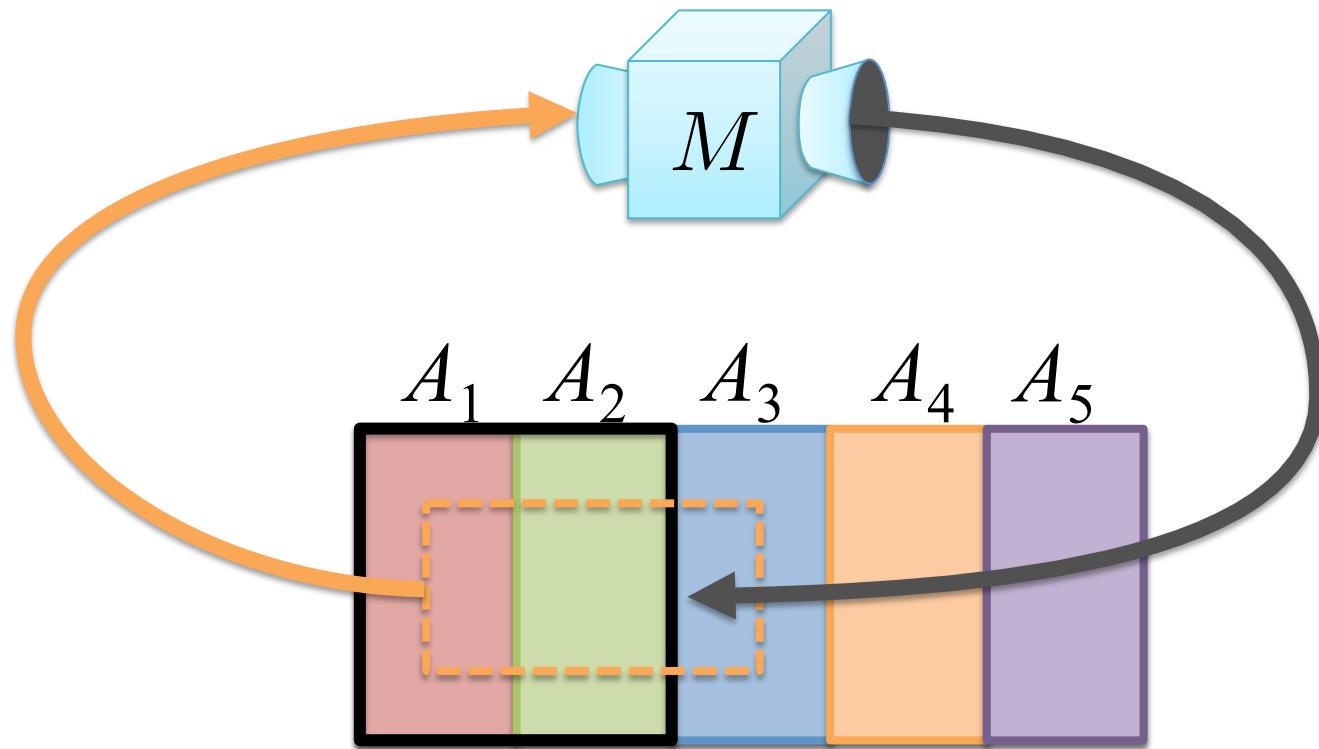
# Relevant Response

- A **disjunction** of answers.



# Inductive Methods

Information **in**, relevant response **out**.



# Verification Methods

**Verification method.** In **every** world  $w$ :

1.  $w \in H$ :  $M$  converges to  $H$  without error.
2.  $w \in H^c$ :  $M$  always concludes  $W$ .

**Refutation method.** In **every** world  $w$ :

1.  $w \in H$ :  $M$  always concludes  $W$ .
2.  $w \in H^c$ :  $M$  converges to  $H^c$  without error.

**Decision method.** does both.

# Fundamental Correspondence

## Proposition.

open	=	verifiable	=	meth. verifiable;
closed	=	refutable	=	meth. refutable;
clopen	=	decidable	=	meth. decidable.

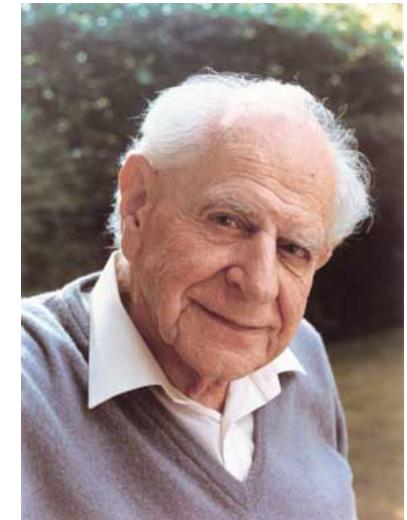
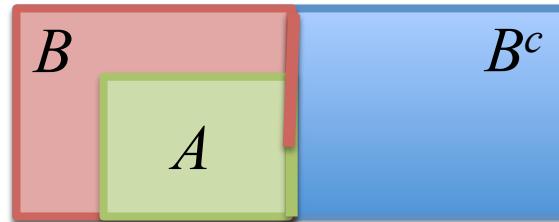




### **3. EMPIRICAL SIMPLICITY**

# Popper's Simplicity Order

- Every information state that **refutes**  $B$  refutes  $A$ .
- **Equivalently:** every information state compatible with  $A$  is compatible with  $B$ .

$$A \preceq B \text{ iff } A \subseteq \text{cl}B.$$


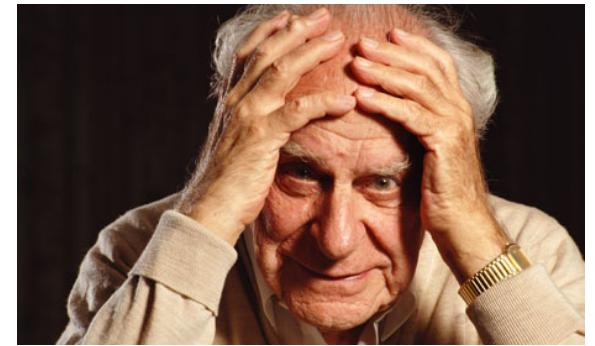
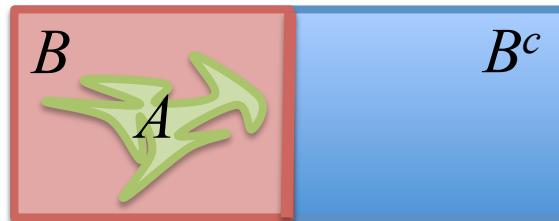
Karl Popper

# The “Tack-on” Objection

- It's wrong that **stronger** theories are **simpler**.



Clark Glymour



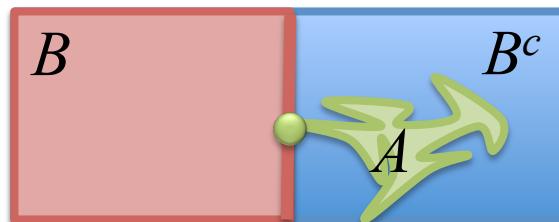
Karl Popper

# Our Slight Revision

- It's possible that you face the problem of induction from  $A$  to  $B$ .

$$A \triangleleft B \text{ iff } A \cap \text{cl } B \setminus B \neq \emptyset.$$

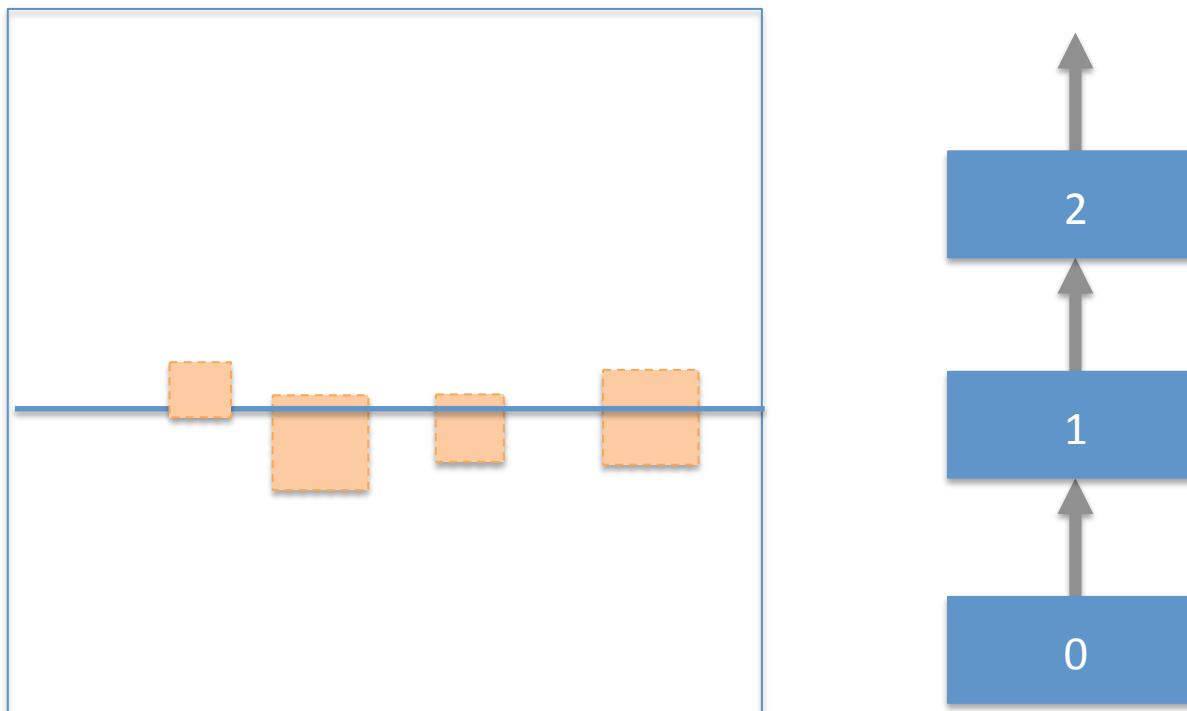
- Strict order if every answer is verifiable.



# Example: Quantitative Laws

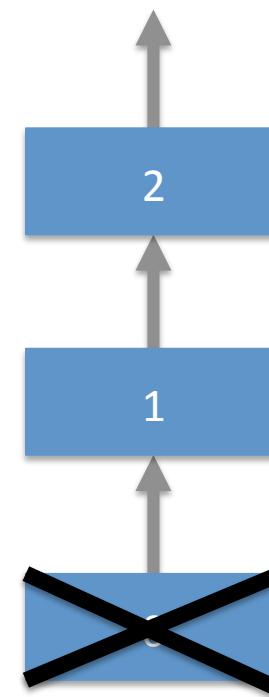
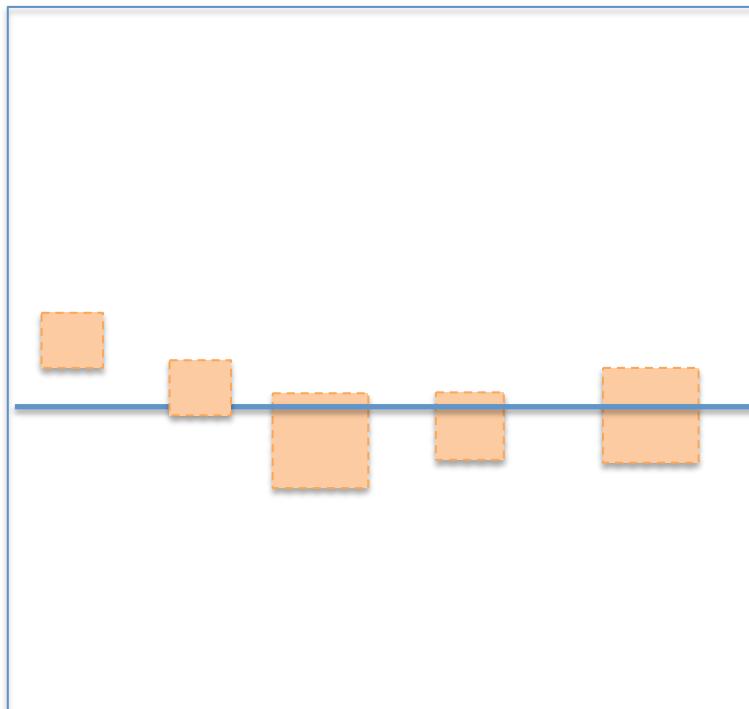
$\mathcal{Q}$  = What is the true polynomial degree?

$\mathcal{I}$  = finitely many inexact measurements.



# Example: Quantitative Laws

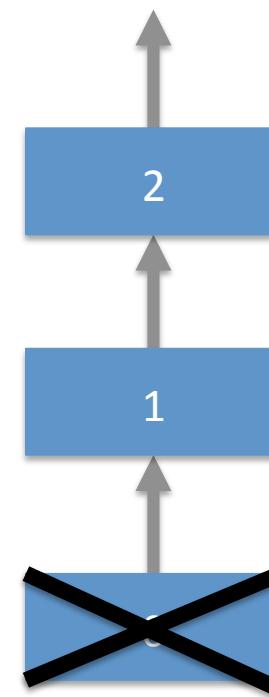
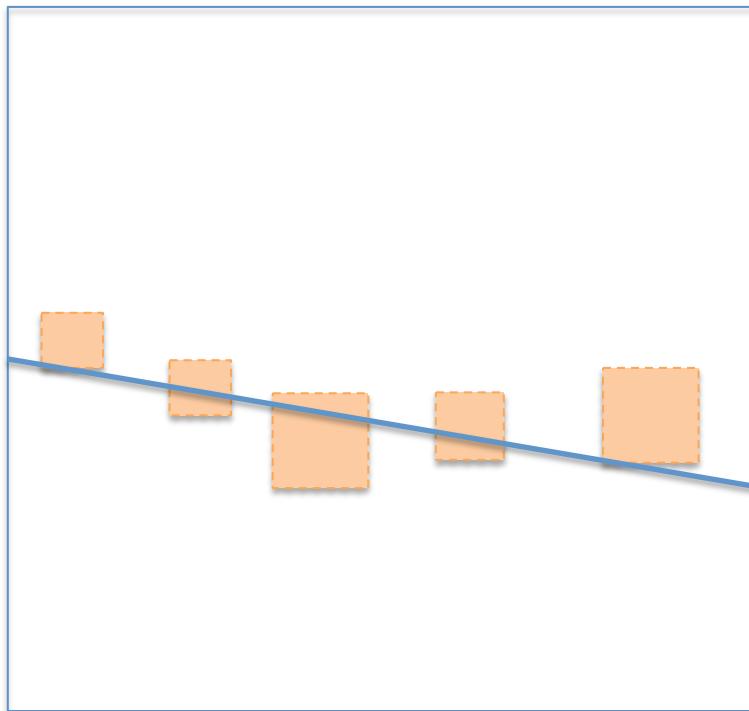
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# Example: Quantitative Laws

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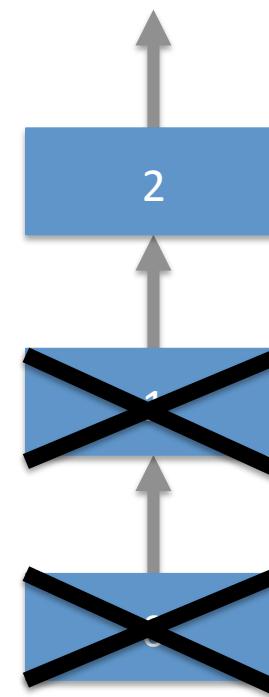
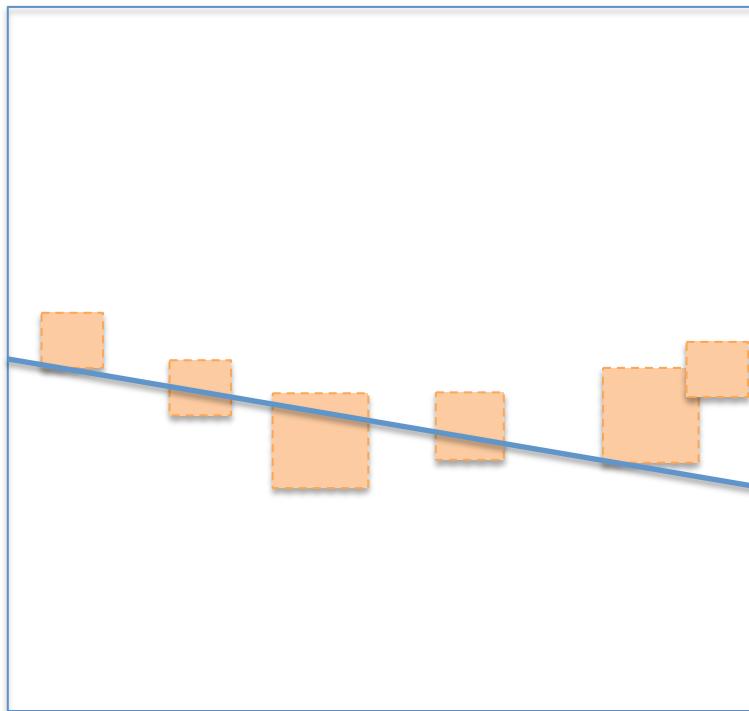
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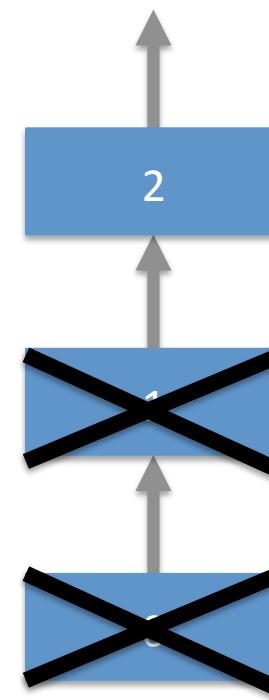
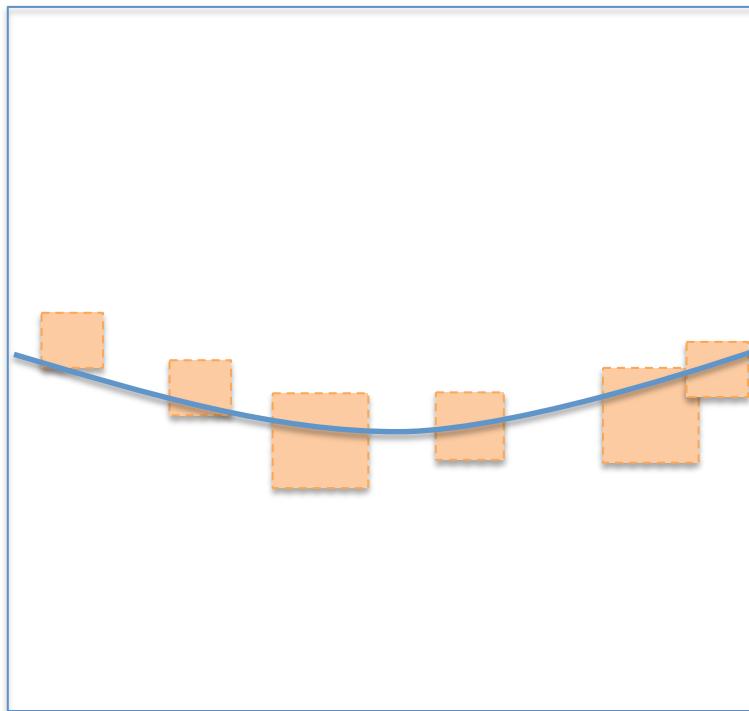
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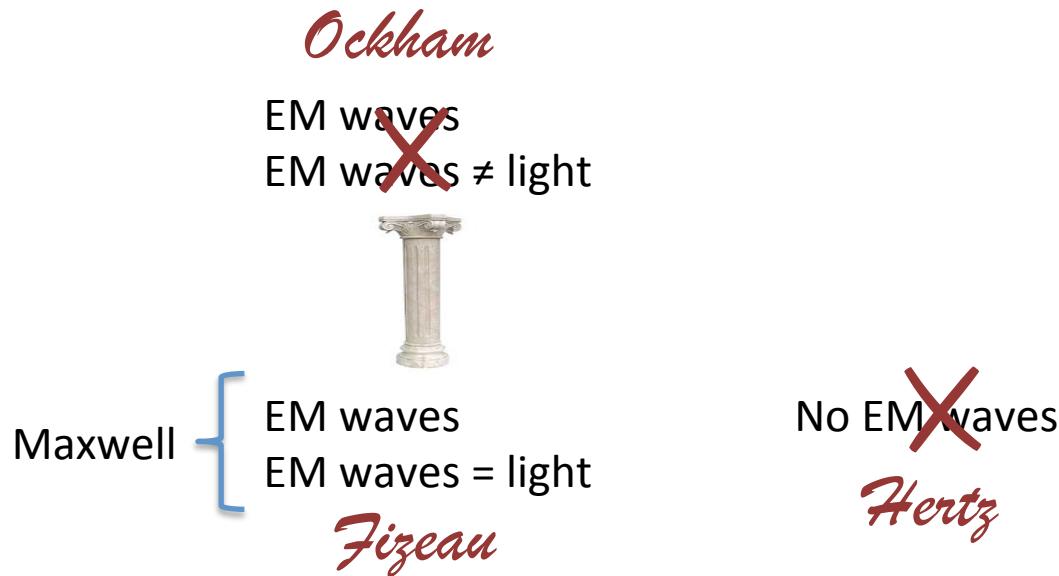


# Topological Simplicity

1. Motivated by the **problem of induction**.
2. Depends only on the **structure** of possible **information**.
3. Independent of **notation**.
4. Independent of **parameterization**.
5. Independent of **prior probabilities**.
6. Non-trivial in **0-dimensional** spaces.

# Example: EM Unification

Actual history (M. Morrisson).



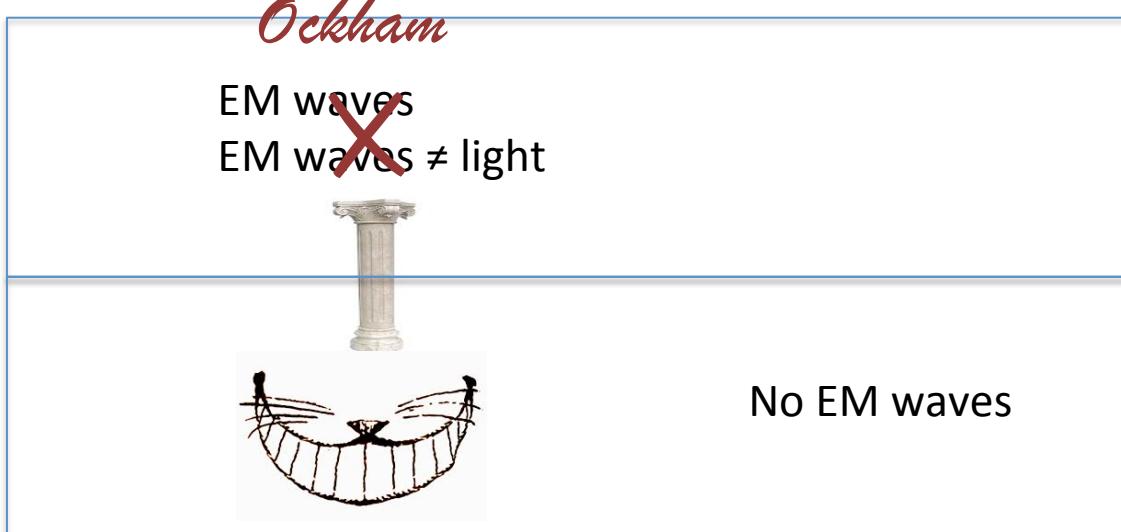
# Example: EM Unification

Hypothetical history (H. Lin).



# Example: EM Unification

- If simplicity is a **ranking**, then Hertz is pre-empted by Ockham (AGM, Spohn).



# Example: EM Unification

- But Hertz can **settle** the question, so **wait**.

*Hertz?*

EM waves

EM waves  $\neq$  light

No EM waves



## 4. OCKHAM'S RAZOR

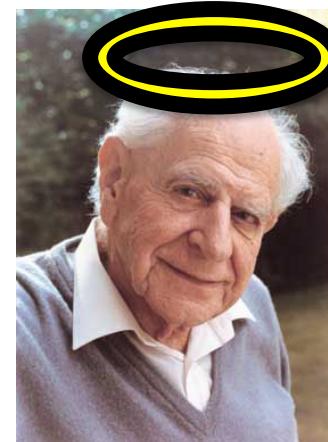
# Ockham's Razor

- Output a **simplest relevant response** given  $E$ .
  - Allows for **suspension of judgment**.
  - Works for **infinite descending simplicity chains**.



# *Popper's Razor*

- Output a relevant response that is **refutable** (**closed**) given  $E$ .



# Error Razor

- “Err on the side of simplicity”.
- In arbitrary world  $w$ , never produce a relevant response  $B$  such that the true answer  $A_w$  is strictly simpler than  $B$ .

# All the Same Razor!

**Proposition.**

Ockham's razor = Popper's razor = error razor.





## 5. OCKHAM'S RAZOR JUSTIFIED

# Inductive Justification

Infer straight to the truth



Too **strong!**

Convergence in the limit



Too **weak!**

# Inductive Justification

Infer straight to the truth



Too **strong!**

**Straightest possible** convergence



**Just right!**

- Feasible;
- Mandates short-run norms.

Convergence in the limit



Too **weak!**

# Departures from Straightness



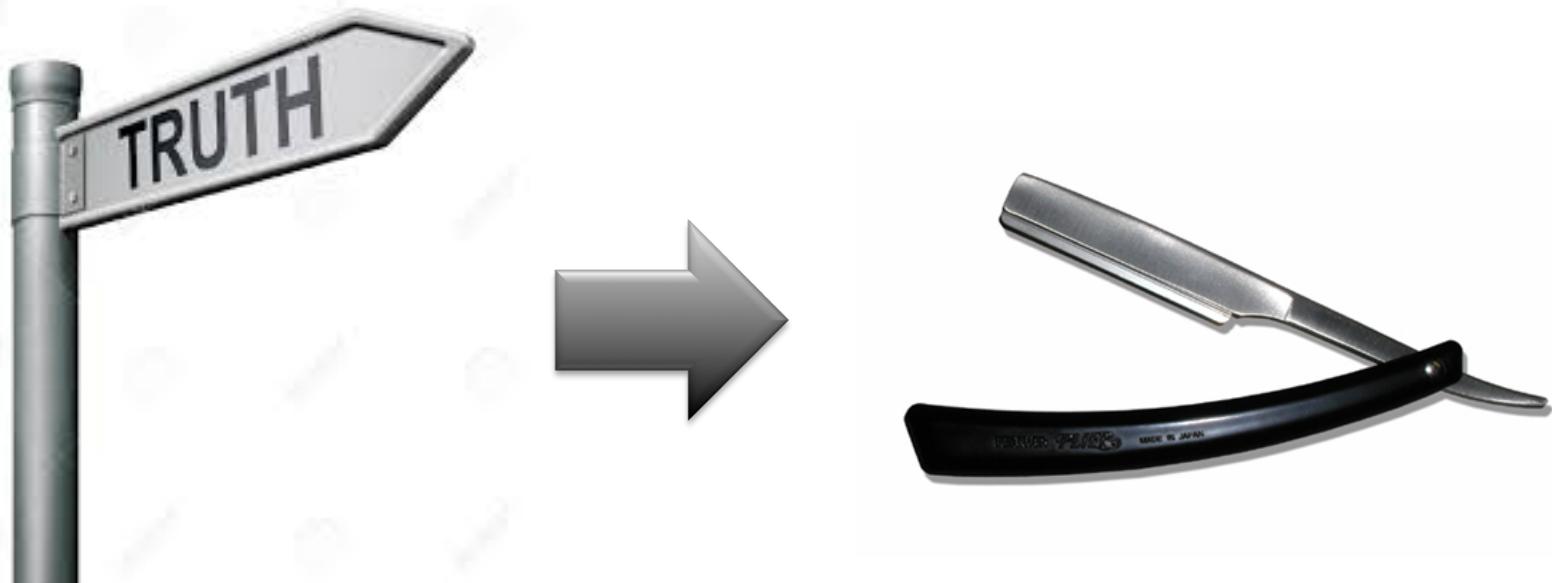
Bad



Worse!

# Thesis

- Ockham's razor is **necessary** for avoiding doxastic cycles.



# Doxastic Cycles

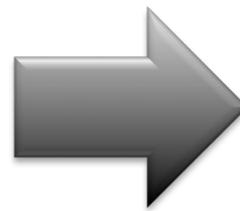
- Each relevant response **contradicts** the preceding.
- The last response **entails** the **first**.



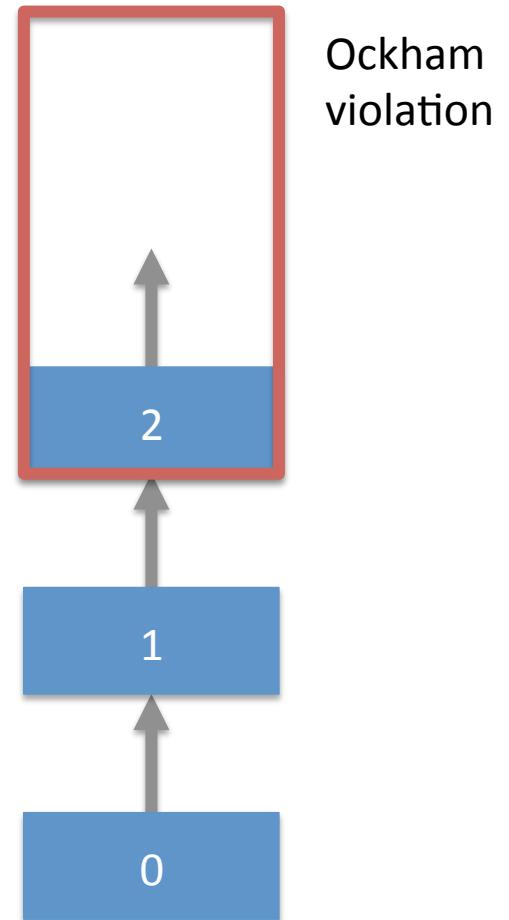
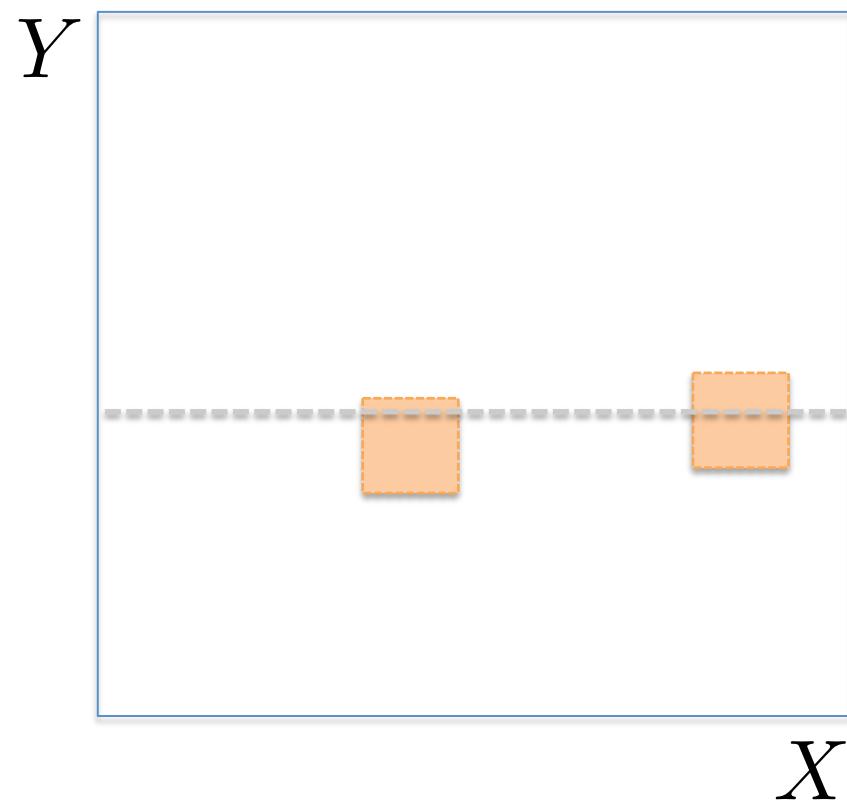
# Main Result 1



- **Proposition:** Every *cycle-free* solution satisfies Ockham's razor.

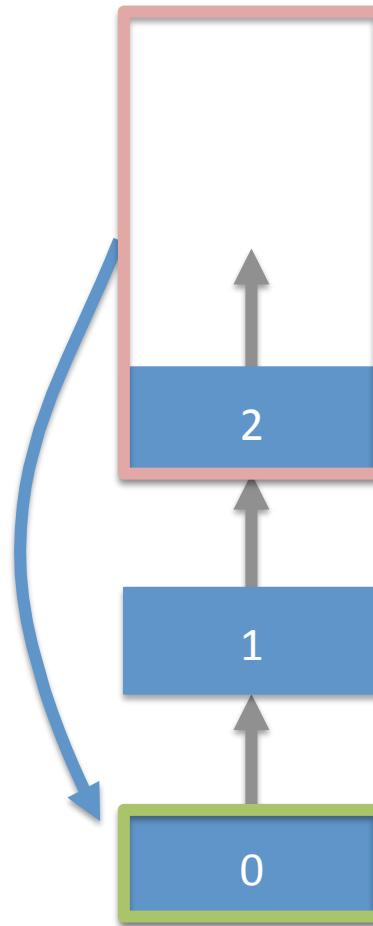
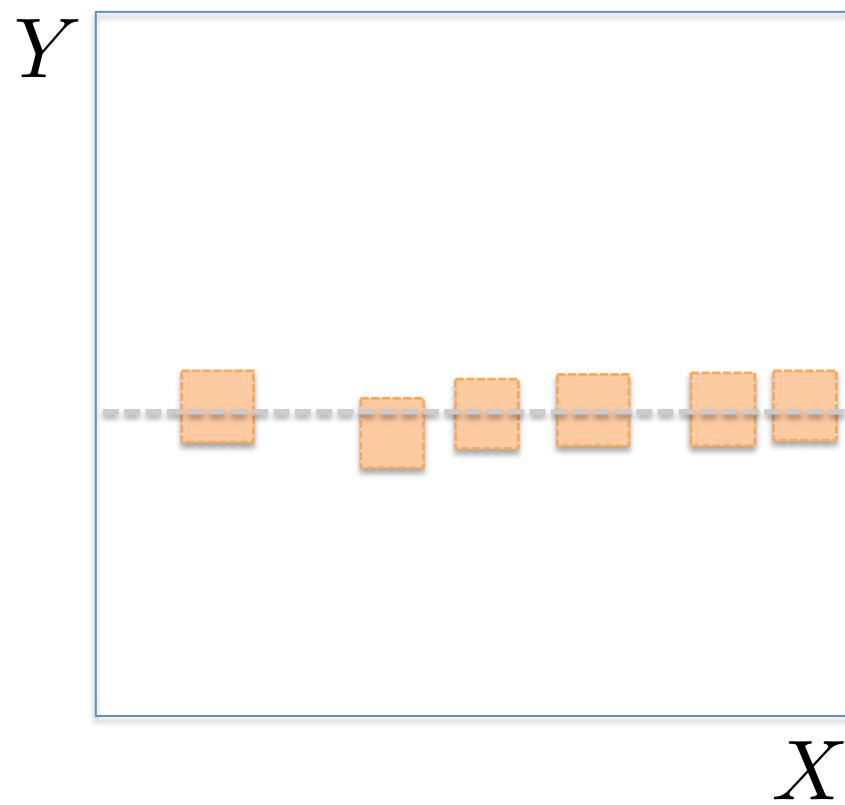


# The Idea



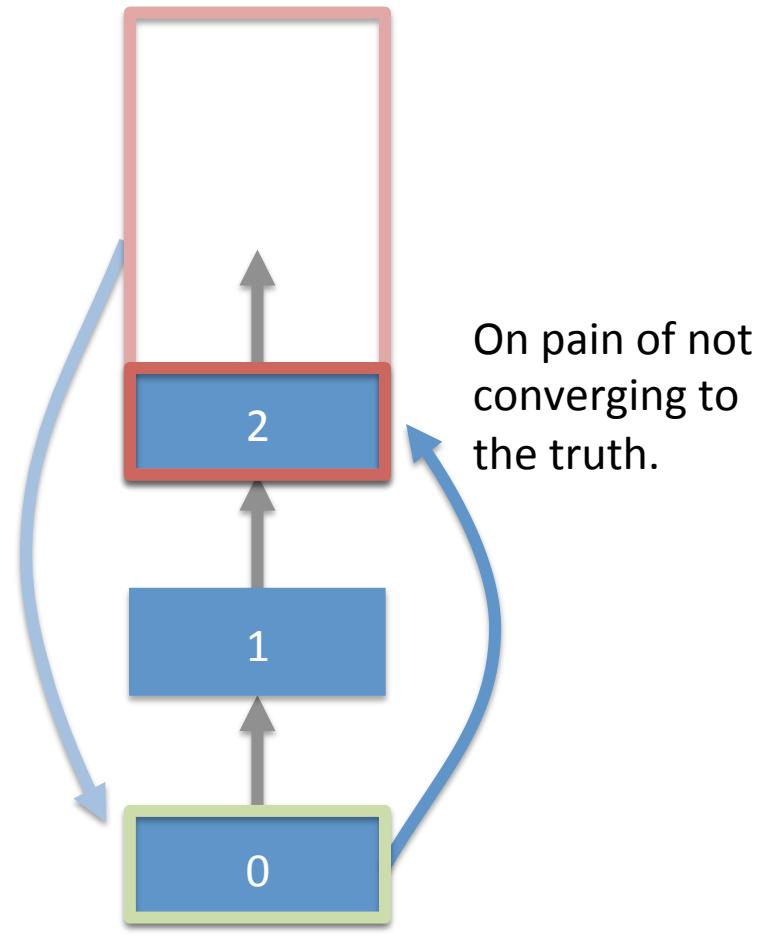
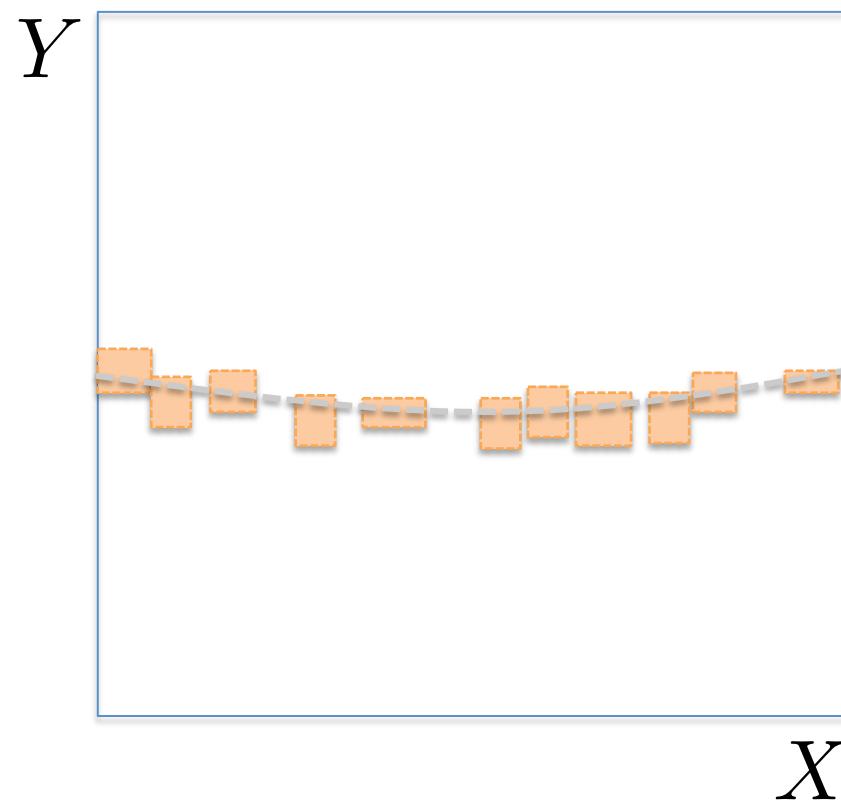
Ockham  
violation

# The Idea



On pain of not  
converging to  
the truth.

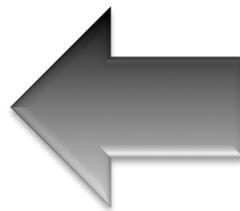
# The Idea



# Main Result 2



- **Proposition** (Baltag, Gerasimczuk, and Smets): Every solvable question is **refinable** to a verifiable question with a **cycle-free** solution.





## 6. OCKHAM'S STATISTICAL RAZOR

# Skepticism

The above account...

“may be okay if the candidate theories are **deductively** related to observations, but when the relationship is **probabilistic**, I am **skeptical** ...”.

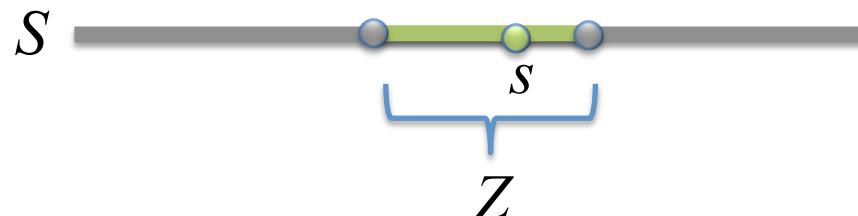


Elliott Sober, *Ockham's Razors*, 2015

# Statistics

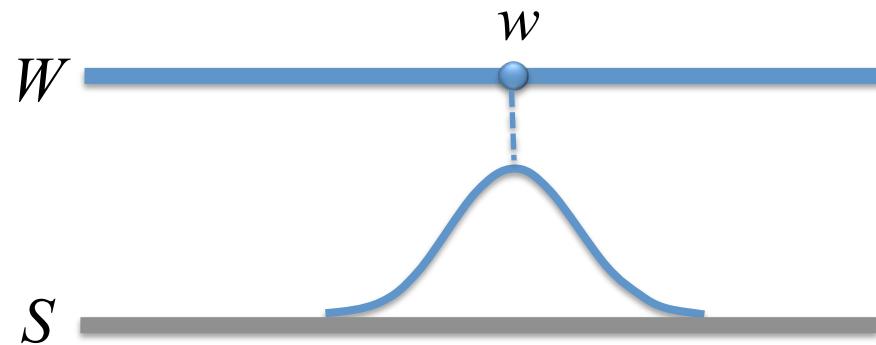
- The **sample space**  $S$  always comes with its **own topology**  $\mathcal{T}$ .
- $\mathcal{T}$  reflects what is **verifiable** about the **sample itself**.

$s$  definitely falls within **open interval**  $Z$ .



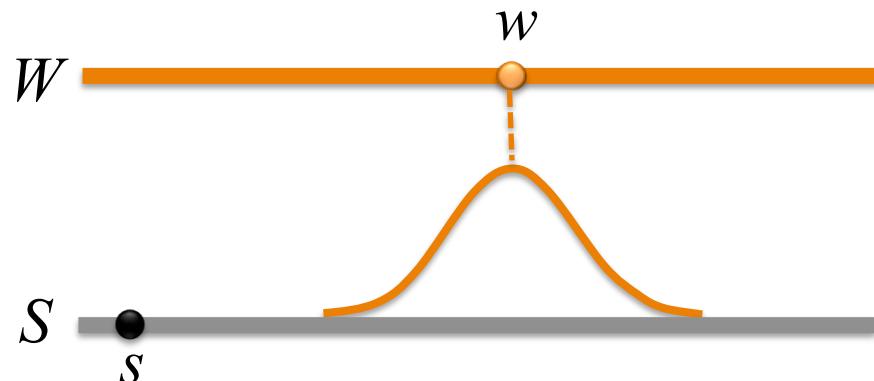
# Statistics

- Worlds are probability measures over  $\mathcal{T}$ .



# The Difficulty

- Every sample is logically consistent with all worlds!
- So it seems that statistical information states are all trivial!



# Response

- Solve for the **unique** topology whose **open sets** are **exactly** the **statistically verifiable** propositions.

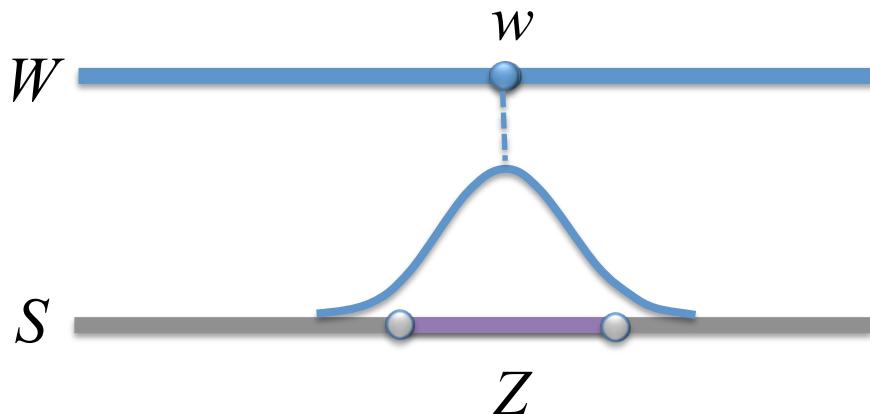
Topology



Statistics

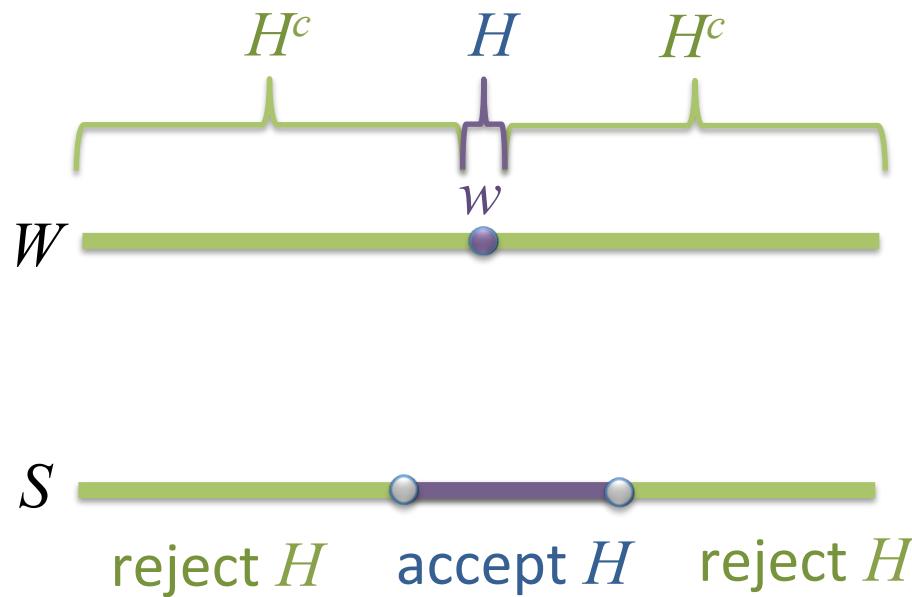
# Feasible Sample Events

- It's impossible to tell whether a point right on the **boundary** of  $Z$  is in or out of  $Z$ .
- $Z$  is **feasible** iff the chance of the boundary is zero in every world.



# Feasible Tests

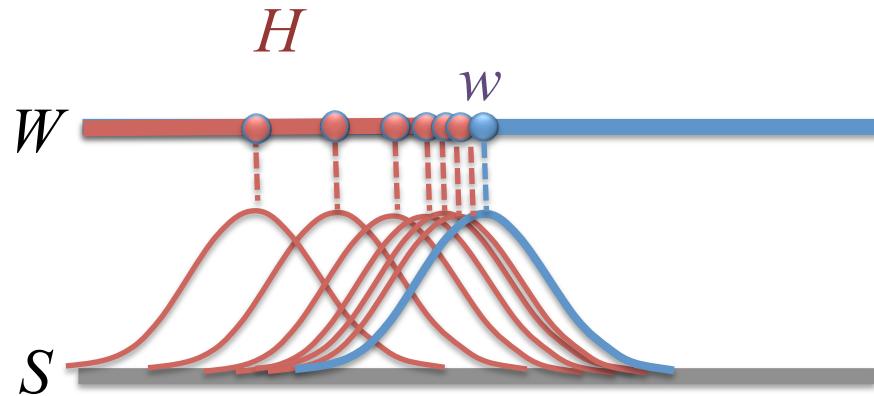
A **feasible test**  $M$  of a statistical hypothesis  $H$  is a measurable function from samples to  $\{\text{not-}H, W\}$  with a **feasible rejection region**.



# Statistical Information Topology

$w \in \text{cl } H$  iff there exists sequence  $(w_n)$  in  $H$ , such that for all feasible tests  $M$ :

$$\lim_{n \rightarrow \infty} p_{w_n}(M \text{ rejects}) \rightarrow p_w(M \text{ rejects}).$$



# Weak Topology

**Proposition:** If  $\mathcal{T}$  has a basis of feasible zones, then  
statistical information topology = the standard, **weak**  
topology.

# Statistical Verification Methods

- A **statistical verification method** for  $H$  at **level**  $\alpha > 0$  is a sequence  $(M_n)$  of **feasible** tests of  $\text{not-}H$  such that for all  $n$ :
  1. if  $w \in H$ :  $M_n$  converges in probability to  $H$ ;
  2. If  $w \in H^c$ :  $M_n$  concludes  $W$  with probability at least  $1-\alpha$ .
- $H$  is **statistically verifiable** iff  $H$  has a statistical **verification** method at **each**  $\alpha > 0$ .

# The Topology of Statistical Methodology



**Proposition.** If  $\mathcal{T}$  has a basis of feasible regions,

1. **open** = statistically verifiable.
2. **closed** = statistically refutable.
3. **clopen** = statistically decidable.

Topology



Statistics

# Simplicity

**Same** as before!

I.e., it is possible that  $A$  is true but  $B$  is never statistically refuted.

# Ockham's Razor

- “Simplest **compatible** with the data” is trivial since **every** answer is logically **compatible** with **every** sample.



# Ockham's Razor

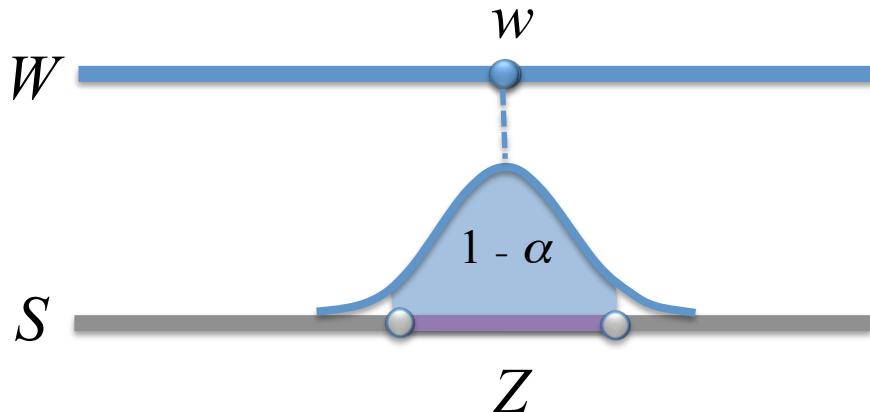
- “Simplest **compatible** with the data” is trivial since **every** answer is logically **compatible** with **every** sample.
- **Solution.** The **error razor** is defined in terms of **truth** rather than **compatibility** with current information, so it still makes sense!

# Ockham's $\alpha$ -Razor

Probabilistic version of the error-razor:

A statistical method is  $\alpha$ -Ockham iff the chance that it outputs an answer more complex than the true answer is bounded by  $\alpha$ .

Agrees with significance for simple vs. complex binary questions!



# Ockham's $\alpha$ -Razor

Statistical method  $(M_n)$  is  $\alpha$ -Ockham iff for all worlds  $w$ , sample sizes  $n$  and relevant responses  $A$ :

if  $Q_w \triangleleft A$ , then  $p_w^n[M_n = A] \leq \alpha$ .

# Reversals in Chance

Method  $(M_n)$  performs the sequence  $(A, B)$  at  $\alpha$  iff there is a world  $w$  and two sample sizes such that:

the **gain** in chance of outputting  $B$

pro-rated by

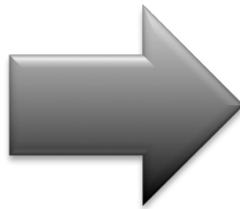
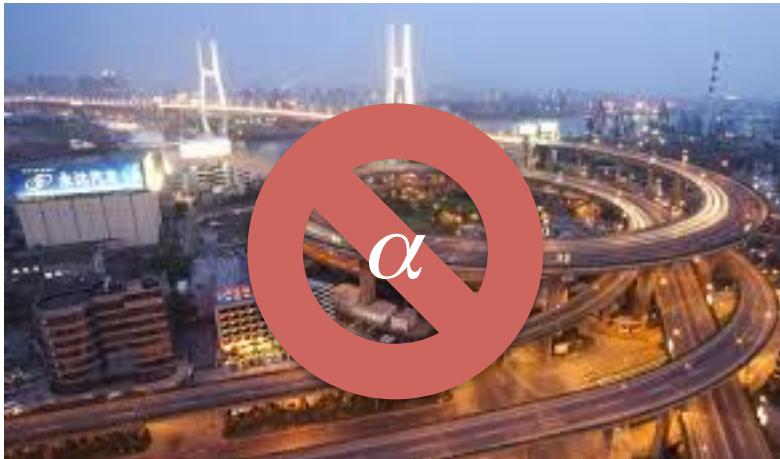
the **loss** in chance of outputting  $A$ ,

is at least  $\alpha$ .

# Main Result



- **Proposition:** Ockham's  $\alpha$  razor is necessary for avoiding  $\alpha$  cycles in chance.
- Valid for **all** solvable problems.



# Conjecture (with Simulations)



- Every solvable question is **refinable** to a verifiable question such that has an  $\alpha$ -cycle-free solution, for all  $\alpha > 0$ .



# Summary and Discussion

1. Simplicity is a **topological** feature of problems.
2. Topological system is **notation-independent**.
3. Ockham's razor is **necessary** for **optimally straight** convergence to the truth.
4. The same holds for **statistical** inductive inference.
5. Optimally straight convergence is **weak**, but its implications for scientific method are **strong**.

# **THE BAYESIAN MIRACLE**

# It Would be a Miracle if...

...the **parameters** of the complex theory were **tuned** to  
**mimic** the predictions of the simple theory.



# The Miracle is in You

On simple data  $E$  there is parameter setting  $\theta$  such that:

$$p(E \mid Comp(\theta)) \approx p(E \mid Simp).$$

So the **miracle** is your own **prior prejudice**.

$$p(Simp) \gg p(Comp(\theta)).$$



But that **is** Ockham's razor, not an **epistemic justification** of it.

# **THE “OVER-FITTING ARGUMENT”**

# Accuracy

- Our national pastime.



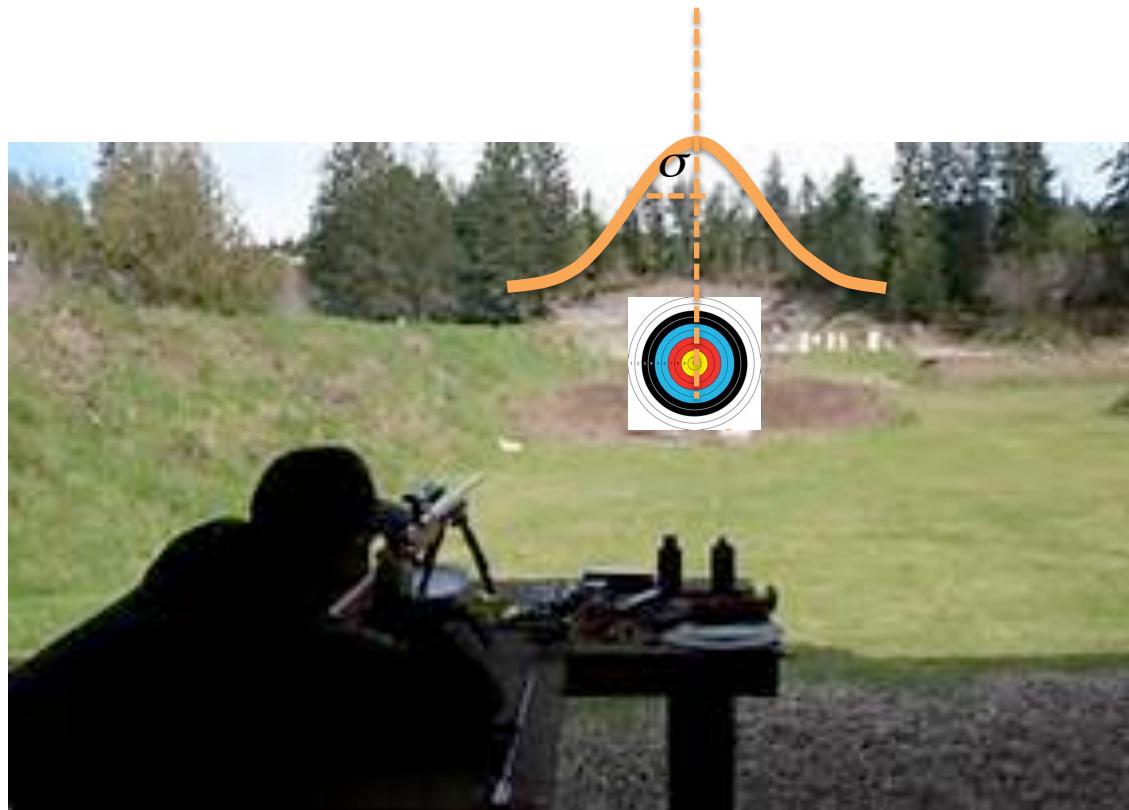
# Analysis of Inaccuracy

- $\text{MSE} = \text{bias}^2 + \text{variance}$ .



# Non-Ockham Empirical Estimates

**Variance** but no **bias**.



# Ockham Estimates

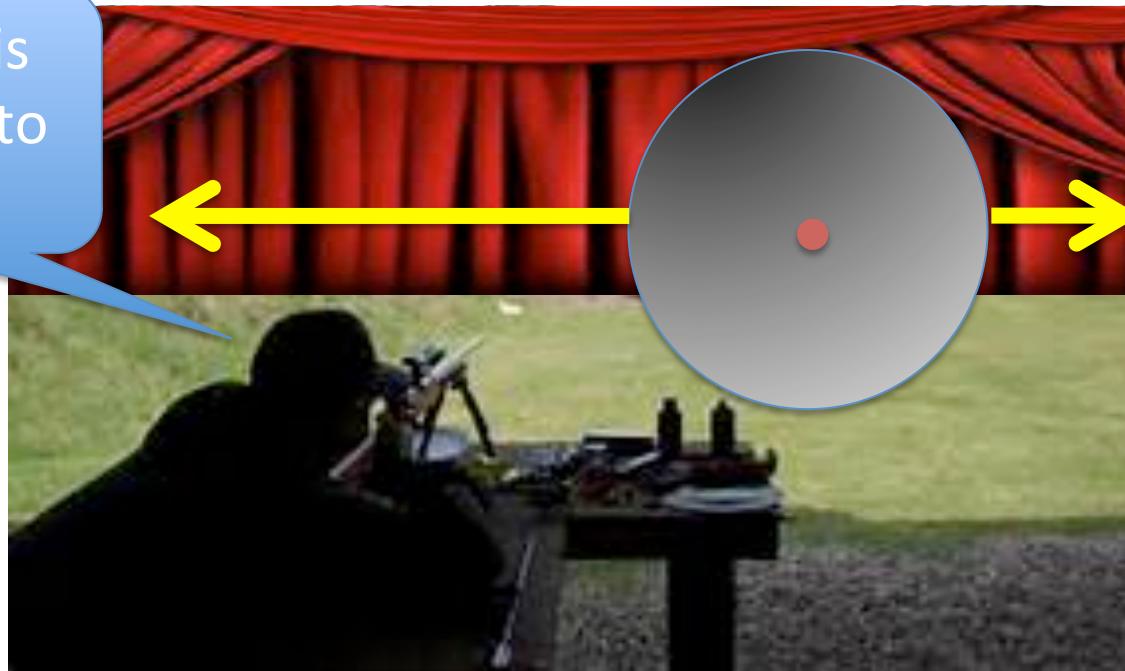
- Like shooting through a **funnel**.
- Small **variance** and **bias** **if** the simple theory is **approximately true**.



# But in Science...

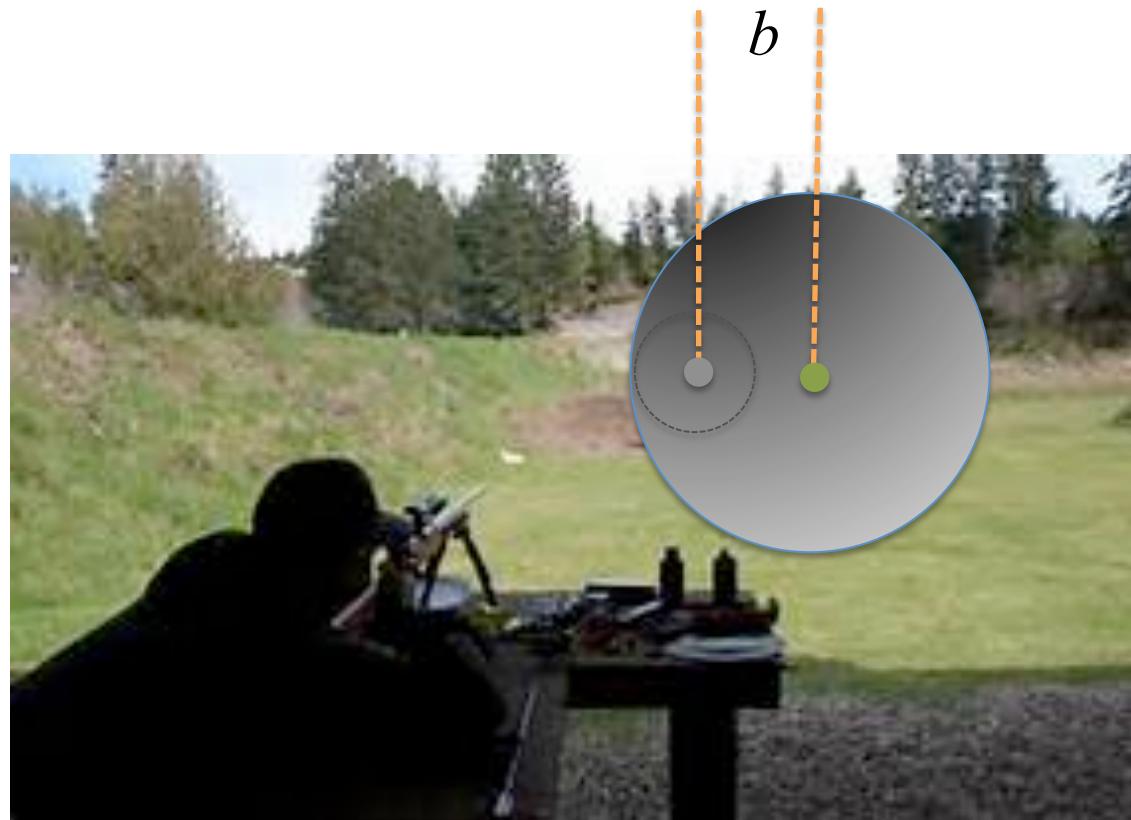
- The funnel must be installed before you see the target!

How is this  
supposed to  
work?



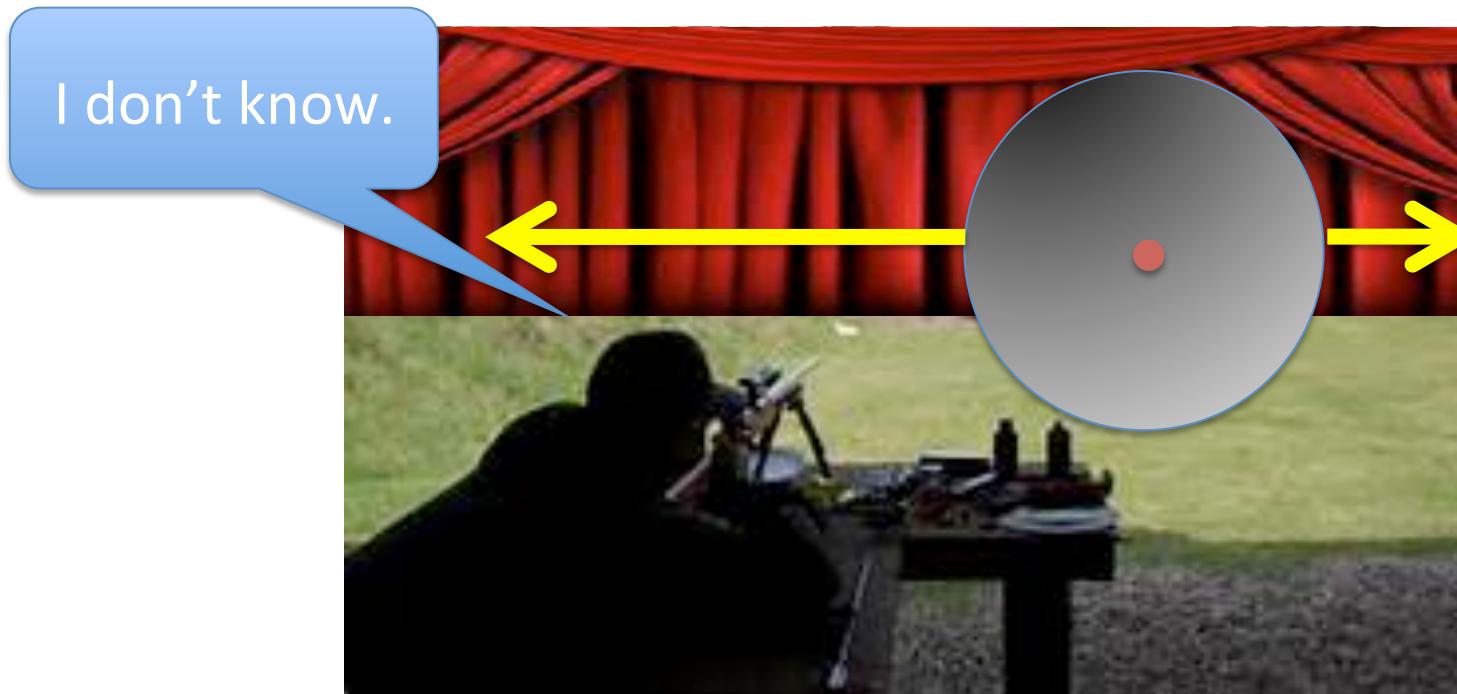
# The Curtain Rises

- If the funnel is not **nearly centered**, it makes **good shots **worse****.



# The Elusive Overfitting Argument

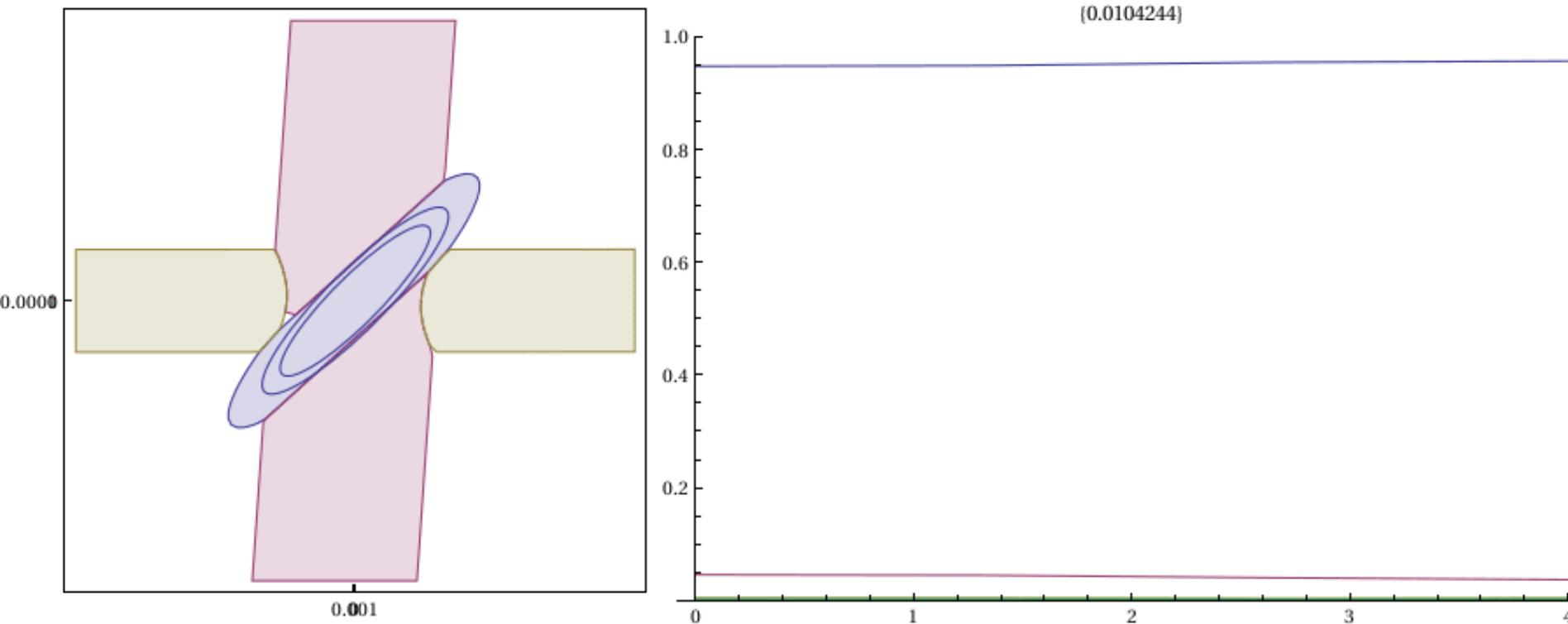
- So how does **blindly** installing the funnel make you **more** accurate?



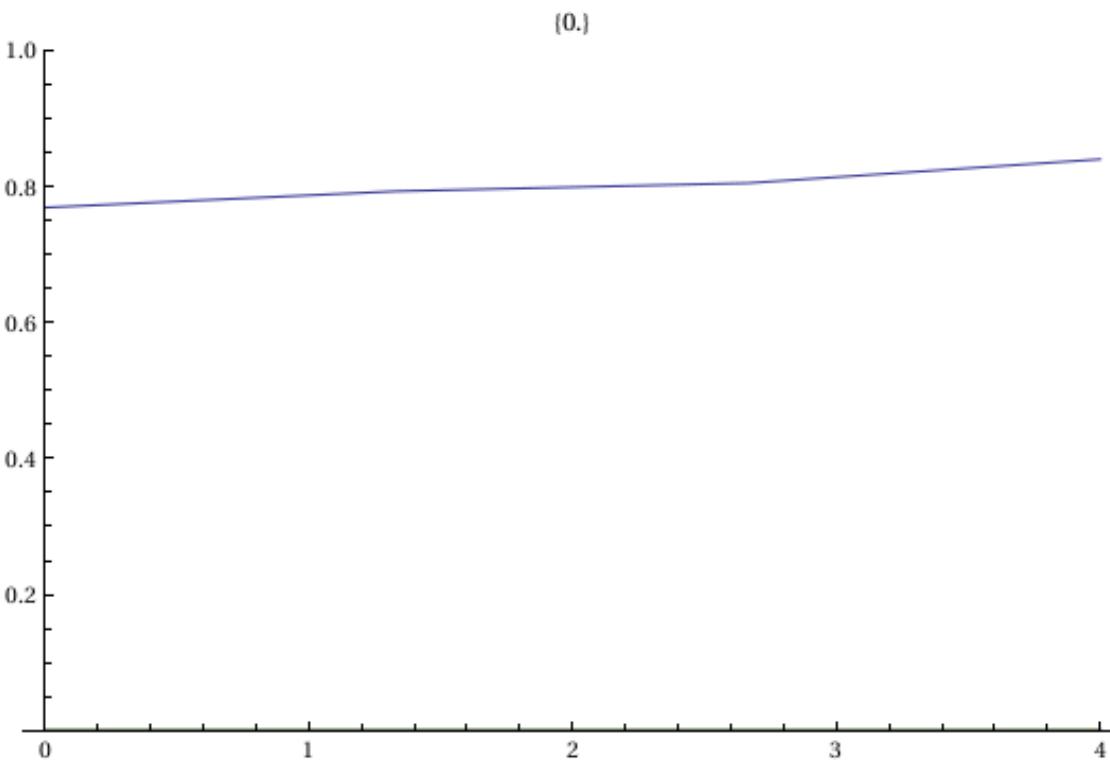
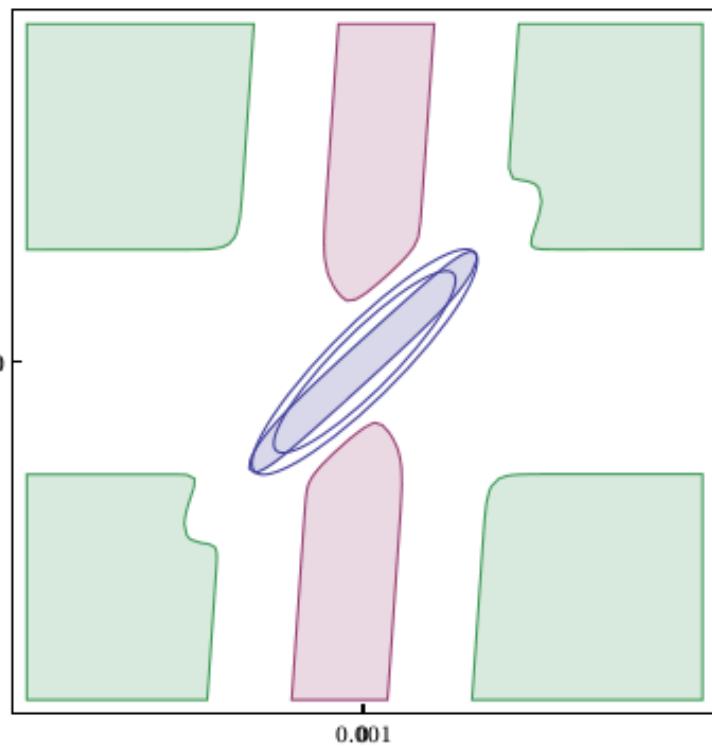
# **SIMULATIONS**

# Bayesian Mode

- Method: Bayesian mode.
- Prior bias toward simplicity, Gaussian priors on parameters.

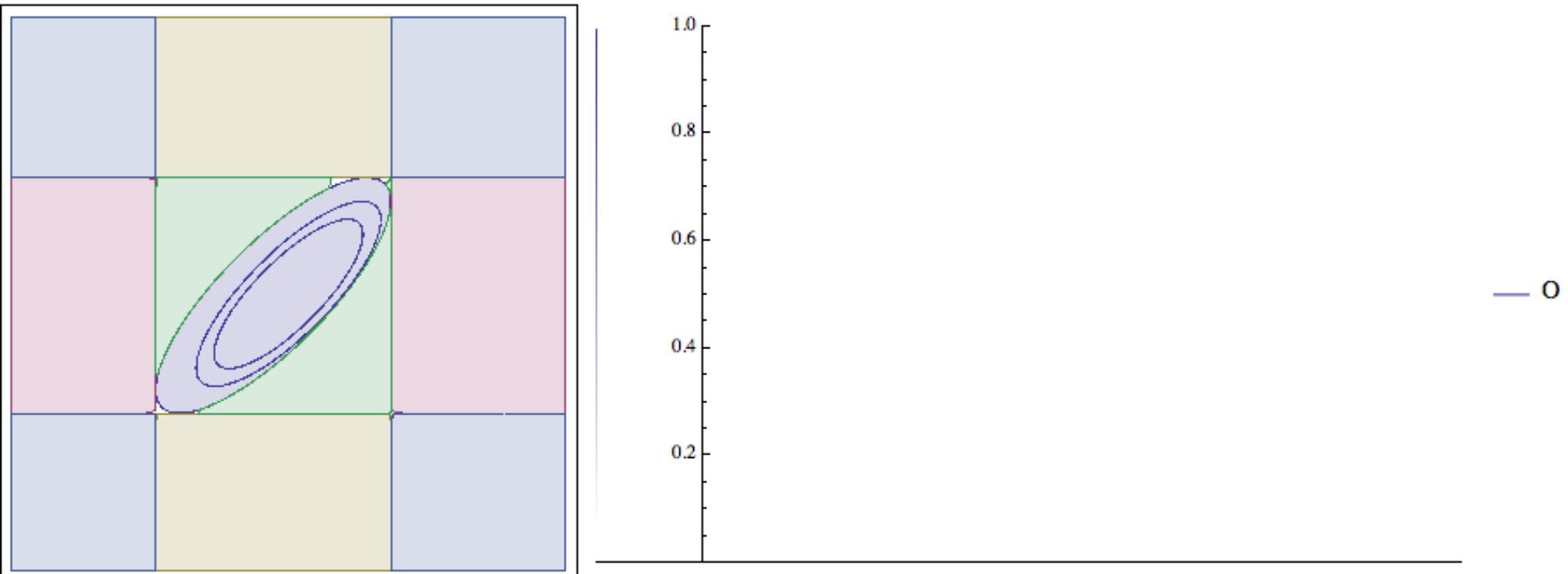


# Bayesian, 95% Threshold



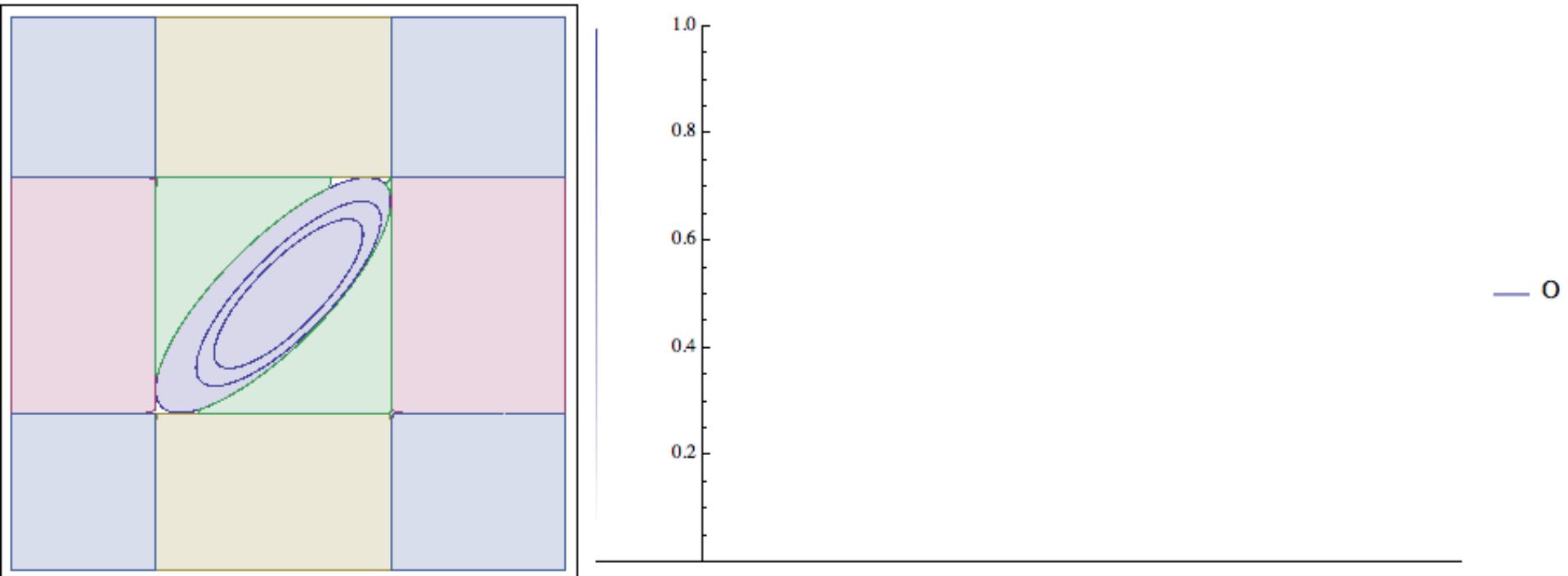
# Frequentist Ockham

- Nested tests



# Error Statistics Reinterpreted

- “Significance” = tolerance on cycles and reversals in chance.
- “Power” = if you are destined to drop a model, get it over with a.s.a.p.



# **SIMPLICITY AND PARADIGMS**

# Example: Competing Paradigms

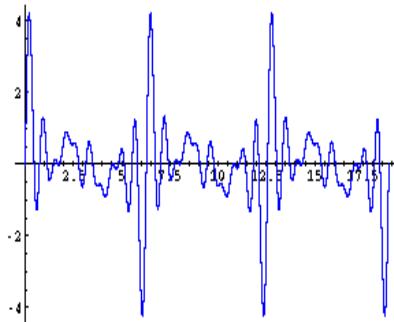
Polynomial paradigm

$$Y = \sum_{i=0}^N a_i X^i.$$



Trigonometric polynomial paradigm

$$Y = \sum_{i=0}^N a_i \sin(iX) + b_i \cos(iX).$$



# Example: Competing Paradigms

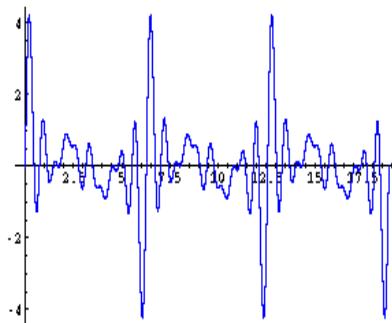
Polynomial paradigm

$$Y = \sum_{i=0}^N a_i X^i.$$

degree

Trigonometric polynomial paradigm

$$Y = \sum_{i=0}^N a_i \sin(iX) + b_i \cos(iX).$$



# Example: Competing Paradigms

$\mathcal{Q}$  = which degree and which paradigm is true?

$\mathcal{I}$  = finitely many inexact measurements.

