

## Gains from Trade in Two-Sided Markets

Today we study a new setting, the two-sided market setting, and a new objective, gains from trade.

**Definition 1.** In *two-sided markets*, we have  $n$  buyers and  $m$  sellers, and a platform facilitating trade. Each seller owns some item(s) and has private values for them  $\mathbf{s}_j$  and will not sell below these values. Each buyer has private values for item(s)  $\mathbf{b}_i$  and will not buy above these values.

Let's simplify significantly for now and focus on the simplest possible setting: one buyer and one seller for one item. This setting is called *bilateral trade*.

**Definition 2.** In the *bilateral trade* setting, there is one seller with one item for sale, and item  $s \sim F_S$  for their own item. There is also one buyer with item  $b \sim F_B$  for the item. The platform's job is to determine a price for the buyer to pay,  $p^B$ , and a payment to the seller  $p^S$ .

We need to review our standard concepts in this new setting and make sure that we understand them, and see if anything additional is necessary.

**Utility.**

**Budget Balance.**

**Many Single-Dimensional Buyers and Sellers.** Now, we consider the setting with  $m$  identical sellers, each seller  $j$  with one item and one value  $s_j \sim F_S$  for their item. There are  $n$  buyers, each with a value  $b_i$  for any item, where  $b_i \sim F_B$ .

**Welfare.**

**Gains from Trade.**

**OPT vs. Constrained-OPT.** Our goal is to maximize GFT, and we would like the mechanism that does so to be

1. Dominant-Strategy Incentive-Compatible
2. Ex-Post Individually Rational
3. Weakly Budget-Balanced

In economics, they call the allocation that is the solution to the unconstrained optimization problem of maximizing GFT “first-best.” They call the mechanism that is the solution to the constrained optimization problem of maximizing GFT *subject to* (1-3) “second-best.”

**Theorem 1** (Myerson Satterthwaite [3]). *Even for 1 buyer, 1 seller, and 1 item, the allocation that maximizes GFT (and thus welfare) may not be implementable by any mechanism satisfying (1-3). That is, first-best is not always attainable.*

**The Optimal Allocation.**

decreasing		increasing
$b^{(1)}$	$\geq$	$s^{(1)}$
$b^{(2)}$	$\geq$	$s^{(2)}$
$\vdots$		$\vdots$
$b^{(q)}$	$\geq$	$s^{(q)}$
$b^{(q+1)}$	$\leq$	$s^{(q+1)}$
$\vdots$		$\vdots$
$b^{(n)}$	$\leq$	$s^{(m)}$

Figure 1: The optimal allocation.

**Claim 1.** The (post-trade) welfare is equal to the sum of the highest  $m$  values in the population.

**The Buyer Trade Reduction(BTR) Mechanism [1].** The simple prior-free mechanism we will use is as follows, inspired by McAfee’s Trade Reduction mechanism [2]:

1. Solicit all buyer and seller values.
2. Compute the optimal allocation on the reported values.
3. Buy items at some  $p^S$ ; sell items to buyers at some  $p^B$ .

(a)

(b)

Let  $BTR(n, m)$  denote the GFT from this mechanism in a market with  $n$  buyers and  $m$  sellers.

**Observation 2** (DSIC+IR). *This mechanism is DSIC and ex-post IR because we set prices only using the values of non-winning agents, so winning agents pay prices lower than their values that they cannot impact.*

**Observation 3** (Budget Balance). *Setting prices according to (3a) or (3b) satisfies weak budget balance.*

**Claim 2.** BTR reduces if and only if the  $m + 1^{\text{st}}$  highest-valued agent is a seller.

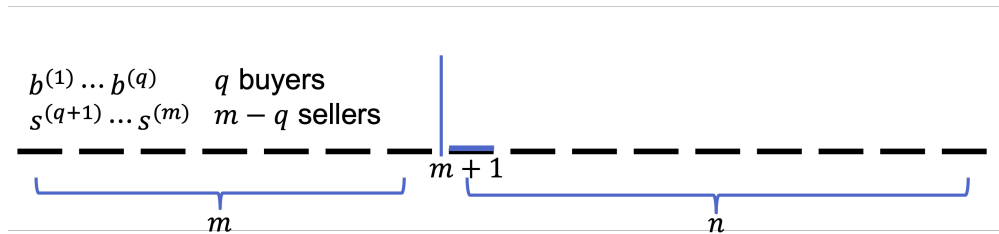


Figure 2: When BTR reduces.

**Theorem 4** (Babaioff G. Gonczarowski [1]). *When buyers and sellers are drawn i.i.d. from some distribution  $F$ , given an initial market with  $n$  buyers and  $m$  sellers, running Buyer Trade Reduction on a market with 1 additional buyer yields at least as much GFT as the optimal GFT in the initial market.*

$$BTR(n+1, m) \geq \text{OPT}(n, m).$$

*Proof.* Approach: Aim to show that

$$\text{OPT}(n+1, m) - \text{OPT}(n, m) \geq \text{OPT}(n+1, m) - BTR(n+1, m).$$

## References

- [1] Moshe Babaioff, Kira Goldner, and Yannai A. Gonczarowski. Bulow-klemperer-style results for welfare maximization in two-sided markets. In Shuchi Chawla, editor, *Proceedings of the 2020 ACM-SIAM Symposium on Discrete Algorithms, SODA 2020, Salt Lake City, UT, USA, January 5-8, 2020*, pages 2452–2471. SIAM, 2020.
- [2] R Preston McAfee. A dominant strategy double auction. *Journal of economic Theory*, 56(2):434–450, 1992.
- [3] Roger B Myerson and Mark A Satterthwaite. Efficient mechanisms for bilateral trading. *Journal of economic theory*, 29(2):265–281, 1983.