SER 334 A Session

Exam 2 Review Session

February 8th 2024

7:00 pm - 9:00 pm MST

Agenda

Practice Exam

Sample Problems

Q&A

SI Session Expectations

Thanks for coming to the **SER 334** SI session. We have a packed agenda and we are going to try to get through as many of our planned example problems as possible. This session will be recorded and shared with others.

- If after this you want to see additional examples, please visit the drop-in tutoring center.
- We will post the link in the chat now and at the end of the session.
 - tutoring.asu.edu
- Please keep in mind we are recording this session and it will be made available for you to review 24-48 hours after this session concludes.
- Finally, please be respectful to each other during the session.

Interact with us:

Zoom Features



Zoom Chat

- Use the chat feature to interact with the presenter and respond to presenter's questions.
- Annotations are encouraged



Opens: Saturday
February 10th
@ 2 AM

Exam Info Page

Closes: Sunday February 11th @ 11:59 PM

Practice Exam

75 minutes

~26 questions



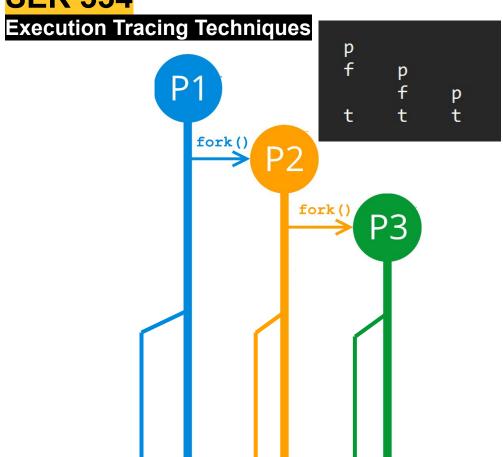
Pro Tip: Take a peek at the programming questions before getting started



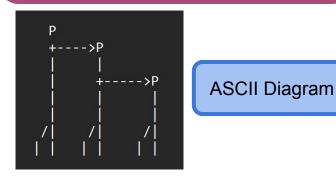
Processes

- 1. Imagine that you are trying to write a program that performs a data-intensive computation, and which needs to support some level of recoverability. Would you be better off using shared data or message passing? Explain. [5 points] [Acuña, Pilcher]
 - (a) Shared data because every thread has individual access to the memory
 - (b) Message passing because there will be a record of each message sent which can be saved
 - (c) Shared data because it would require less memory to be used than message passing
 - (d) Message passing because if a message gets corrupted, it only affects one other process

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Any other suggestions or techniques for tracing during the exam?



р		
fork	р	
	fork	р
thread	thread	thread

2. Trace the following program. As pids come into existence, assume they come from 2600, 2601, 2602, 2603, and so on. [Acuña, Silberschatz]

```
int main() {
   pid t pid, pid1;
   pid = fork();
   if (pid = 0) {
       pid1 = getpid();
       printf("child: pid = %d", pid); /* A */
       printf("child: pid1 = %d", pid1); /* B */
   else
       pid1 = getpid();
       printf("parent: pid = %d", pid); /* C */
       printf("parent: pid1 = %d", pid1); /* D */
       wait (NULL);
   return 0;
```

temp1 150 temp2 temp3 temp1 150 temp2 temp3

From the program above, identify the values of the process ids at lines A, B, C, and D. [10 points]

3. Trace the previous program and then list all of the possible outputs (just write the A/B/C/D text; no need for pids) that could be generated. (Assume that fork will always succeed.) [10 points] [Acuña]

```
int main() {
  pid t pid, pid1;
  pid = fork();
   if (pid == 0) {
       pid1 = getpid();
       printf("child: pid = %d", pid); /* A */
       printf("child: pid1 = %d", pid1); /* B */
   else
       pid1 = getpid();
       printf("parent: pid = %d", pid); /* C */
       printf("parent: pid1 = %d", pid1); /* D */
       wait (NULL);
   return 0;
```

4. Consider the following program.

```
int main() {
    pid_t pid;

pid = fork();
    if (pid == 0) {
        fork();
        thread_create(...);
    }
    thread_create(...);
    return 0;
}
```

How many unique threads are created in the program above? [5 points] [Acuña, Silberschatz]

- (a) 2
- (b) 3
- (c) 4
- (d) 5

- 5. Is it faster to switch between processes or threads within those processes? Explain. [5 points] [Acuña]
 - (a) Processes they are permanently assigned to a specific CPU so switching never occurs.
 - (b) Processes they have kernel level privileges to access the PCB handling functions.
 - (c) Threads they don't require the call stack to be saved/restored since the memory is shared among all threads for a given process.
 - (d) Threads they are light weight and don't require saving/restoring many resources.

- 6. Consider the task of implementing a mutex, specifically the unlock functionality (like is supported by "int pthread_mutex_unlock(pthread_mutex_t *mutex)"). Which data structure would you prefer to use to determine the next waiting process to wake up? [5 points] [Acuña]
 - (a) Hash tables it will mean it's O(1) to find the next process.
 - (b) Binary Tree it mimics the parent-child relationships that are created by calls to fork().
 - (c) Stack it dynamically resizes unlike the other data structures.
 - (d) Queue it would also enable the mutex to support bounded waiting time.

7. In pthreads, mutex support involves three functions and a new type:

```
typedef pthread_mutex_t
int pthread_mutex_init(pthread_mutex_t *mutex, const pthread_mutexattr_t *attr);
int pthread_mutex_lock(pthread_mutex_t *mutex);
int pthread mutex unlock(pthread mutex t *mutex);
```

How does this functionality address (or does not address) the issues of Mutual Exclusion, Progress, and Bounded Waiting Time? [10 points] [Acuña]

8. The following code snippet models a user checking out a book from a library:

```
struct Book{
    char* title;
    int user;
    int is_checked_out;
}
void check_out(int userId, int book_num, Book* books){
    Book book = books[book_num];
    while(book.is_checked_out);
    book.is_checked_out = 1;
    book.user = userId;
    books[book_num] = book;
}
```

If this code were to be used in a multithreaded program (where multiple users/threads may try to check out books at the same time) would there be a race condition? If so, describe this race condition. Assume that the book being processed is initally not checked out. (The state is the contents of the books array after use of this function.) [10 points] [Edgar]

- 9. How does a monitor function differently from a semaphore? [5 points] [Pilcher]
 - (a) a monitor only allows one process to be in its queue at a time while a semaphore allows multiple processes
 - (b) a monitor and a semaphore function the same except a monitor is represented as a class
 - (c) a monitor allows a set number of processes to be active at a time while a semaphore only allows one process to be active at a time
 - (d) a monitor allows one process to be active at a time while a semaphore can allow multiple processes to be active at the same time

10. Consider the semaphore based solution we had for the reader-writer problem:

What functionality does the logic involving read_count and rc_mutex provide? [5 points] [Acuña,

semaphore rw_mutex = 1;
semaphore rc_mutex = 1;
int read_count = 0;

//writer process
do {
 wait(rw_mutex);
 /* writing is performed */
 signal(rw_mutex);
} while (true);

//reader process

do {

Pilcher]

- (a) It guarantees there is never a negative amount of readers
- (b) It ensures there is only one writer process at one time
- (c) It allows multiple processes to read at the same time
- (d) It ensures only one process can read at a time

if (read_count == 1)
 wait(rw_mutex);
signal(rc_mutex);
/* reading is performed */

/* reading is p wait (rc_mutex);

wait(rc_mutex);
read count++;

wait(rc_mutex);
read_count--;

read_count == ;
if (read_count == 0)
signal(rw_mutex)

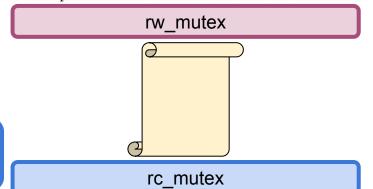
signal(rw_mutex);
signal(rc_mutex);

signal(rc_m
} while (true);

read count

Pilcher

10. Consider the semaphore based solution we had for the reader-writer problem:



- (a) It guarantees there is never a negative amount of readers
- (b) It ensures there is only one writer process at one time
- (c) It allows multiple processes to read at the same time
- (d) It ensures only one process can read at a time

```
semaphore rw mutex = 1;
semaphore rc mutex = 1;
int read count = 0;
```

```
//writer process
                                                                         do {
                                                                              wait (rw mutex);
                                                                              /* writing is performed */
                                                                              signal(rw mutex);
                                                                            while (true);
                                                                         //reader process
                                                                         do {
                                                                              wait (rc mutex);
                                                                              read count++;
                                                                              if (read count == 1)
What functionality does the logic involving read_count and rc_mutex provide? [5 points] [Acuña,
                                                                                   wait (rw mutex);
                                                                              signal(rc mutex);
                                                                              /* reading is performed */
```

wait (rc mutex); read count --;

if (read count = 0)

signal(rc mutex);

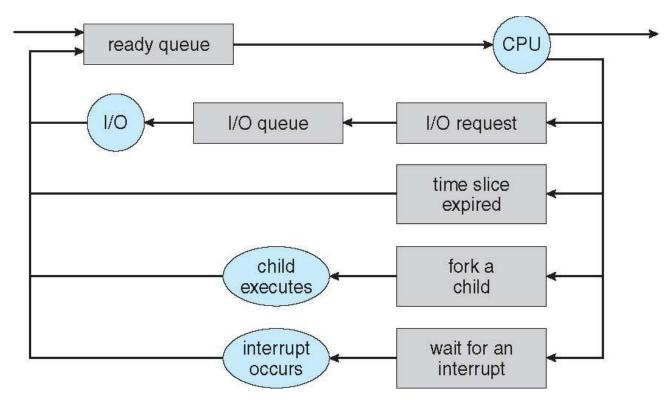
while (true);

signal (rw mutex);

Module 5 Sample



3. [Bahremand] Explain why the queue diagram in slide 9 has a continuous cycle that flows between the listed queues and resources? Is a cycle necessary? [2 points]



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5. [Acuña] Consider the following program written with the pthreads library: _What is this programming trying to compute and how does it uses threads?

```
Exam 2 Review
```

```
#include <pthread.h>
#include <stdio.h>
int t1, t2, input;
void* runner(void *param) {
    int result, upper = (int)param;
    int a = 1, b = 1, c = 1;
    for (int i = 2; i < input-upper; i++) {</pre>
        c = a + b;
        a = b:
        b = c;
    if(upper == 1) {
        t1 = c:
    } else {
        t2 = c:
    pthread_exit(0);
```

```
int main(int argc, char *argv[]) {
    pthread t tid1, tid2;
    pthread attr t attr;
    input = 10;
    pthread attr init(&attr);
    pthread create(&tid1, &attr, runner, 1);
    pthread create(&tid2, &attr, runner, 2);
    pthread join(tid1, NULL);
    pthread join(tid2, NULL);
    printf("result = %d\n", t1+t2);
```



5. [Acuña] Consider the code for Peterson's Solution. Notice that part of the algorithm has been commented out. Explain how this changes it's functionality. Will it still solve the critical section problem? Explain.

```
//shared memory
int turn = 0;
bool flag[2] = { false, false };
//for some process i
do {
    flag[i] = true;
    turn = j;
    while (flag[j] /* && turn == j */);
    //critical section
    flag[i] = false;
    //remainder section
} while (true);
```

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6. [Lisonbee] The bounded buffer problem is a result of producing and consuming work asynchronously at different and/or variable rates. Provided below is are the producer and consumer functions used by a program. Assume that the producer and consumer are running in parallel. In order to solve the bounded buffer problem demonstrated here, the appropriate calls to wait and signal need to be added. Rewrite the above code Exam 2 Review using wait and signal and initialize the 3 semaphores to the appropriate values (Note: you must use all three semaphores at least once in your calls to wait and signal).

```
int data[15], i0 = 0, i1 = 0;
semaphore mutex = ...;
semaphore empty = ...;
semaphore full = ...;
void producer() {
   while (1) {
        data[i0] = i0 * i0;
        i0 = ++i0 \% 15;
```

```
void consumer() {
    while (1) {
        printf("%d\n", data[i1]);
        i1 = ++i1 \% 15;
```

```
7. [Acuña] Consider the monitor based solution we had for the dining philosophers problem:
                                                                                              Module 8 Sample
                   monitor DiningPhilosophers {
                           enum { THINKING; HUNGRY, EATING) state [5];
                           cond t self [5];
                            void pickup(int i) {
Exam 2 Review
                                state[i] = HUNGRY;
                                test(i);
                                if (state[i] != EATING)
                                    self[i].wait();
                           void putdown(int i) {
                                state[i] = THINKING;
                                // test left and right neighbors
                                test((i + 4) \% 5);
                                test((i + 1) \% 5):
                            void test(int i) {
                                if ((state[(i + 4) \% 5] != EATING) \&\& //AND \longrightarrow OR
                                                                            //AND \longrightarrow OR
                                     (state[i] == HUNGRY) &&
                                     (state[(i + 1) \% 5] != EATING)) 
                                         state [i] = EATING;
                                         self[i].signal();
                            void initialization code() {
                                for (int i = 0; i < 5; i++)
                                     state[i] = THINKING;
```

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How would the functionality of this change if the conditional in test() used | instead of &&?

2. [Acuña] **Implement** the wait() operation for a semaphore using the test_and_set() instruction. Assume that you have two methods available: add_this_process(list) which adds the current process to a list, and block() which pauses the current process until it is signaled.

```
struct semaphore* S {
    static boolean lock;
    struct process_node* list;
    int value;
}
void wait(semaphore* S) {
```

SER 334 Scratch Space

Upcoming Events

SI Sessions:

- Sunday, February 11th at 7:00 pm MST Cancelled Good luck on Exam 2!
- Monday, February 12th at 7:00 pm MST

Review Sessions:

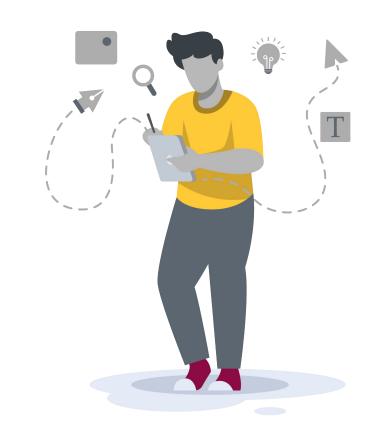
Exam 3 Review: TBD

Good luck on the Exam You got this!

Questions?

Survey:

http://bit.ly/ASN2324



More Questions? Check out our other resources!

tutoring.asu.edu



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^{*}Available slots for this pilot are limited

Additional Resources

- Course Repo
- Course Discord
- BMP File Format (Wiki)
- Linux Kernel API
- Bootlin Linux Cross Referencer
- Dining Philosophers Interactive
- Producer/Consumer Visual