

High-Precision Computation: Mathematical Physics and Dynamics

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Abstract

At the present time, IEEE 64-bit floating-point arithmetic is sufficiently accurate for most scientific applications. However, for a rapidly growing body of important scientific computing applications, a higher level of numeric precision is required. Such calculations are facilitated by high-precision software packages that include high-level language translation modules to minimize the conversion effort. This paper presents an overview of recent applications of these techniques and provides some analysis of their numerical requirements. We conclude that high-precision arithmetic facilities are now an indispensable component of a modern large-scale scientific computing environment.

1 Introduction

Virtually all present-day computer systems, from personal computers to the largest supercomputers, implement the IEEE 64-bit floating-point arithmetic standard, which provides 53 mantissa bits, or approximately 16 decimal digit accuracy. For most scientific applications, 64-bit arithmetic is more than sufficient, but for a rapidly expanding body of applications, it is not. In this paper we will examine a variety of situations where high-precision arithmetic is useful:

1. *Ill-conditioned linear systems.* Many innocent-looking problems involve ill-conditioned linear systems that give rise to numerical errors with 64-bit arithmetic.
2. *Large summations.* Anomalous results often stem from the loss of associativity in summations, particularly when executed on a parallel computer environment where the order of summation cannot be controlled [71].

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3. *Long-time simulations.* Almost any kind of physical simulation, if performed over many time intervals, will eventually depart from reality, due to cumulative round-off error, in addition to errors arising from discretization of time and space.
4. *Large-scale simulations.* Computations that are well-behaved on modest-sized problems, such as those run on a single-CPU system, may exhibit significant numerical errors when scaled up massively parallel systems.
5. *Resolving small-scale phenomena.* It is often necessary to employ a very fine-scale resolution to “zoom” in on the phenomena in question.
6. “Experimental mathematics” computations. Numerous recent results in experimental mathematics could not be obtained except by very high precision computations.

With regards to item 1, it should be kept in mind that the vast majority of persons currently performing numerical computations are not experts in numerical analysis, and this fact is not likely to change anytime soon. For example, in 2010 at the University of California, Berkeley, a total of 219 students enrolled in the two sections of Math 128A, a one-semester introductory numerical analysis course required of applied math majors, but only 24 enrolled in Math 128B, a more advanced course. By contrast, in the same year a total of 870 seniors graduated in the Division of Mathematical and Physical Sciences (including Mathematics, Physics and Statistics), the College of Chemistry and the College of Engineering (including Computer Science). If we add to this list graduates in other fields with computational components, such as biology, geology, medicine and social sciences, we conclude that fewer than 2% of the Berkeley graduates each year who likely will be using numerical computations in their career work have advanced training in numerical analysis.

Thus, for the foreseeable future, almost all technical computing will be performed by persons who have had only basic training in numerical analysis, or none at all. High-precision arithmetic is an attractive option for such users, because even in situations where numerically better behaved algorithms are known in the literature that may resolve a numerical problem, it is often both easier and more reliable to simply increase the precision used for the existing algorithm, using tools such as those described in Section 2. And, as we will see below, there are problems for which no known algorithmic change can rectify the numerical difficulties encountered.

1.1 Extra precision versus algorithm changes

The following example illustrates some of the issues involved. Suppose one wishes to recover the integer polynomial that produces the result sequence $(1, 32771, 262217, 885493, 2101313, 4111751, 7124761)$ for integer arguments $(0, 1, \dots, 6)$. While there are several ways to approach this problem, many scientists and engineers will employ a least-squares scheme, since this is a very familiar tool in scientific data analysis, and efficient library software is readily available. Indeed, this approach is suggested in a widely used reference [70, pg. 44]. In this approach, one

constructs the $(n + 1) \times (n + 1)$ linear system

$$\begin{bmatrix} n+1 & \sum_{k=1}^n x_k & \cdots & \sum_{k=1}^n x_k^n \\ \sum_{k=1}^n x_k & \sum_{k=1}^n x_k^2 & \cdots & \sum_{k=1}^n x_k^{n+1} \\ \vdots & \vdots & \ddots & \vdots \\ \sum_{k=1}^n x_k^n & \sum_{k=1}^n x_k^{n+1} & \cdots & \sum_{k=1}^n x_k^{2n} \end{bmatrix} \begin{bmatrix} a_0 \\ a_1 \\ \vdots \\ a_n \end{bmatrix} = \begin{bmatrix} \sum_{k=1}^n y_k \\ \sum_{k=1}^n x_k y_k \\ \vdots \\ \sum_{k=1}^n x_k^n y_k \end{bmatrix}, \quad (1)$$

where (x_k) are the integer arguments and (y_k) are the sequence values. Then one solves for (a_1, a_2, \dots, a_n) using, for example, *LINPACK* [45] or *LAPACK* [44] software.

In the specific problem mentioned above, a double-precision (64-bit) floating-point implementation of the least-squares scheme succeeds in finding the correct polynomial coefficients, which, after rounding to the nearest integer, are $(1, 0, 0, 32769, 0, 0, 1)$, or, in other words, $f(x) = 1 + (2^{15} + 1)x^3 + x^6$. Unfortunately, this scheme fails to find the correct polynomial for a somewhat more difficult problem, namely to find the degree-8 polynomial that generates the 9-long sequence $(1, 1048579, 16777489, 84941299, 268501249, 655751251, 1360635409, 2523398179, 4311748609)$, for integer arguments $(0, 1, \dots, 8)$. The program finds approximate degree-8 polynomial coefficients, but they are not correct, even after rounding to the nearest integer — too much floating-point round-off error has occurred.

Numerical analysts may point out here that this approach is not the best scheme for this type of problem, in part because the Vandermonde matrix system (1) is known to be ill-conditioned. A more effective approach in the cases such as this is to employ the Lagrange interpolating polynomial, which, given a set of $n + 1$ data points $(x_0, y_0), (x_1, y_1), \dots, (x_n, y_n)$, is defined as $L(x) = \sum_{j=0}^n y_j p_j(x)$, where

$$p_j(x) = \prod_{0 \leq i \leq n, i \neq j} \frac{x - x_i}{x_j - x_i}. \quad (2)$$

In the problem at hand, $x_j = j$ for $0 \leq j \leq n$. The chief sources of numerical error here are the summations inherent in the formula $L(x) = \sum_{j=0}^n y_j p_j(x)$ (see item 2 above).

This scheme, implemented with 64-bit IEEE arithmetic, correctly deduces that the 9-long data sequence above is produced by the polynomial $1 + (2^{20} + 1)x^4 + x^8$. However, this scheme fails when given the more challenging 13-long input data vector $(1, 134217731, 8589938753, 97845255883, 549772595201, 2097396156251, 6264239146561, 15804422886323, 35253091827713, 71611233653971, 135217729000001, 240913322581691, 409688091758593)$, which is generated by $1 + (2^{27} + 1)x^6 + x^{12}$.

The state-of-the-art algorithm in this area, as far as the present authors are aware, is a technique due to James Demmel and Plamen Koev [42], which accurately solves “totally positive” systems such as (1), where the determinant of any square submatrix is positive. A *Matlab* implementation of this scheme is available at [62]. We found that this program solves the degree-6 and degree-8 problems mentioned above, but, like the Lagrange polynomial scheme, fails for the degree-12 problem.

However, there is another approach to these problems: simply modify the source code of any reasonably effective solution scheme to invoke higher-precision arithmetic. For example, when

Algorithm	Precision (digits)	Problem degree		
		6	8	12
Least-squares	16	succeeded	failed	failed
	31	succeeded	succeeded	succeeded
Lagrange	16	succeeded	succeeded	failed
	31	succeeded	succeeded	succeeded
Demmel-Koev	16	succeeded	succeeded	failed

Table 1: Success and failure of various polynomial data fit schemes

we modified our Fortran-90 least-squares scheme to employ double-double precision (approximately 31-digit accuracy), using the QD software [59] mentioned in Section 2, we were able to correctly solve all three problems (degrees 6, 8 and 12). Converting the Lagrange polynomial scheme to use double-double arithmetic was even easier, and the resulting program also solved all three problems without incident. These results are summarized in Table 1. No entry is listed for the Demmel-Koev scheme with 31-digit arithmetic, because we relied on a 16-digit *Matlab* implementation, although we have no reason to doubt that it would also succeed.

2 High-precision software

Efficient algorithms are known for performing, to any desired precision, the basic arithmetic operations, square and n -th roots, and most transcendental functions [30, pp. 215–245], [31, pp. 299–318], [32, 33, 34, 37]. Until recently, utilizing high-precision arithmetic required one to rewrite a scientific application with individual subroutine calls for each arithmetic operation. The difficulty of writing and debugging such code has deterred all but a few computational scientists and mathematicians from using such tools.

In the past 10 years or so, several high-precision software packages have been produced that include high-level language interfaces that make such code conversions relatively painless. These packages typically utilize custom datatypes and operator overloading features, which are available in languages such as C++ and Fortran-90, to facilitate conversion. Here are some high-precision arithmetic software packages that are freely available on the Internet, listed in alphabetical order. The ARPREC [19], QD [59] and MPFUN90 packages are available from the first author’s website: <http://crd-legacy.lbl.gov/~dhbailey/mpdist>.

- ARPREC. This package includes routines to perform arithmetic with an arbitrarily high level of precision, including many algebraic and transcendental functions. High-level language interfaces are available for C++ and Fortran-90, supporting real, integer and complex datatypes.
- GMP. This package includes an extensive library of routines to support high-precision integer, rational and floating-point calculations. GMP has been produced by a volunteer effort and is distributed under the GNU license by the Free Software Foundation. It is available at <http://gmplib.org>.

- MPFR. The MPFR library is a C library for multiple-precision floating-point computations with exact rounding, and is based on the GMP multiple-precision library. Additional information is available at <http://www.mpfr.org>.
- MPFR++. This is a high-level C++ interface to MPFR. Additional information is available at <http://perso.ens-lyon.fr/nathalie.revol/software.html>. A similar package is GMPFRXX, available at <http://math.berkeley.edu/~wilken/code/gmpfrxx>.
- MPFUN90. This is similar to ARPREC in user-level functionality, but is written entirely in Fortran-90 and provides a Fortran-90 language interface.
- QD. This package includes routines to perform “double-double” (approx. 31 digits) and “quad-double” (approx. 62 digits) arithmetic. High-level language interfaces are available for C++ and Fortran-90, supporting real, integer and complex datatypes. This software is much faster than using arbitrary precision software when only 31 or 62 digits are required.

Just as an example of the simple case, the QD package, which provides double-double and quad-double arithmetic, is based on the following algorithms for the accurate addition and multiplication of two IEEE 64-bit operands using rounded arithmetic, due to Knuth [61] and Dekker [41]:

```

function [x, y] = TwoSum(a; b)
x = f1(a + b)
z = f1(x - a)
y = f1((a - (x - z)) + (b - z))

function [x, y] = Split(a)
c = f1(factor · a)   (in double precision factor = 227 + 1)
x = f1(c - (c - a))
y = f1(a - x)

function [x, y] = TwoProd(a; b)
x = f1(a · b)
[a1, a2] = Split(a)
[b1, b2] = Split(b)
y = f1(a2 · b2 - (((x - a1 · b1) - a2 · b1) - a1 · b2))

```

In the above, `f1` stands for the floating-point evaluation using rounded arithmetic. These algorithms satisfy the following error bounds [68] (where \mathbb{F} denotes the set of floating-point numbers and u the rounding unit of the computer):

Theorem 1 *For $a, b \in \mathbb{F}$ and $x, y \in \mathbb{F}$, `TwoSum` and `TwoProd` verify*

$$\begin{aligned}
[x, y] &= \text{TwoSum}(a, b), x = \text{fl}(a + b), x + y = a + b, |y| \leq u|x|, |y| \leq u|a + b|, \\
[x, y] &= \text{TwoProd}(a, b), x = \text{fl}(a \times b), x + y = a \times b, |y| \leq u|x|, |y| \leq u|a \times b|.
\end{aligned}$$

One downside of using high-precision software is that such facilities greatly increase computer run times, compared with using conventional 64-bit arithmetic. For example, computations using double-double precision arithmetic typically run five to ten times slower than with 64-bit arithmetic. This figure rises to at least 25 times for the quad-double arithmetic, to more than 100 times for 100-digit arithmetic, and to well over 1000 times for 1000-digit arithmetic. However, in some cases high-precision arithmetic is only needed in one or two places in the code (such as in a summation loop), so that the total run time is not greatly increased.

3 Applications of high-precision arithmetic

3.1 High-precision solutions of ordinary differential equations

One central question of planetary theory is whether the solar system is stable over cosmological time frames (many millions or billions of years). Planetary orbits are well known to exhibit chaotic behavior. Indeed, as Isaac Newton once noted, “The orbit of any one planet depends on the combined motions of all the planets, not to mention the actions of all these on each other. To consider simultaneously all these causes of motion and to define these motions by exact laws allowing of convenient calculation exceeds, unless I am mistaken, the forces of the entire human intellect.” [47, p. 121].

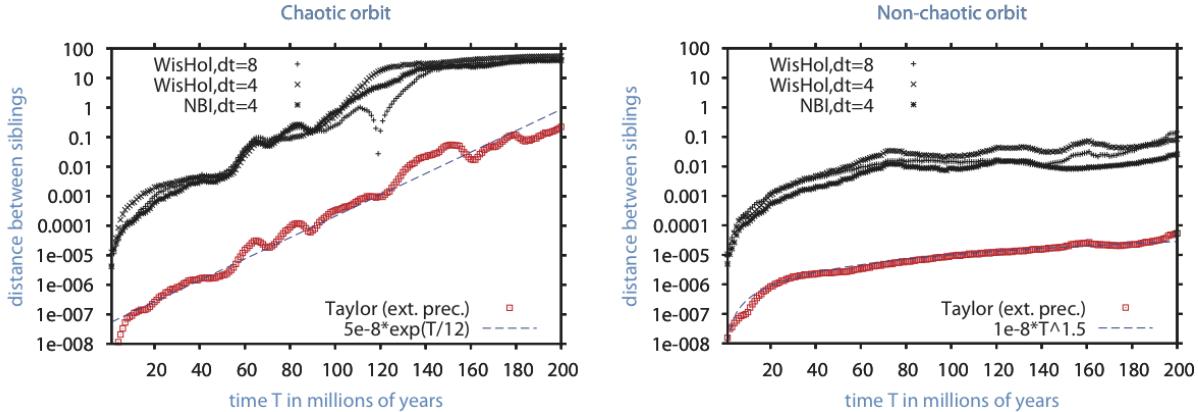


Figure 1: Divergence between nearby trajectories, integrated with four different numerical integrators (the Wisdom-Holman symplectic integrator with two stepsizes, the NBI’s 14th order Cowell-Sörmer integrator and the Taylor method to check the results). Left: a chaotic trajectory with a Lyapunov time of about 12 million years. Right: a trajectory showing no evidence of chaos over 200My. Both trajectories are within observational uncertainty of the outer planetary positions. (Reproduced with permission from [57])

Scientists have studied this question by performing very long-term simulations of planetary motions, often using special-purpose computer systems [4]. These simulations typically do fairly well for long periods, but then fail at certain key junctures, such as when two planets pass fairly close to each other. Researchers have found that double-double or quad-double arithmetic is required to avoid severe numerical inaccuracies, even if other techniques are employed to

reduce numerical error [63]. A team led by W. Hayes studied solar system orbits using various numerical ordinary differential equation (ODE) integrators, checked to higher precision using a Taylor series integrator, performed using 19-digit Intel extended precision [57] (see Figure 1).

High-precision arithmetic has also arisen in the study of dynamical systems, such as in the study of the bifurcations and stability of periodic orbits. Runge-Kutta schemes have been widely used for such calculations, but during the last few years the Taylor method, augmented with high-precision arithmetic, has emerged as a preferred method [76].

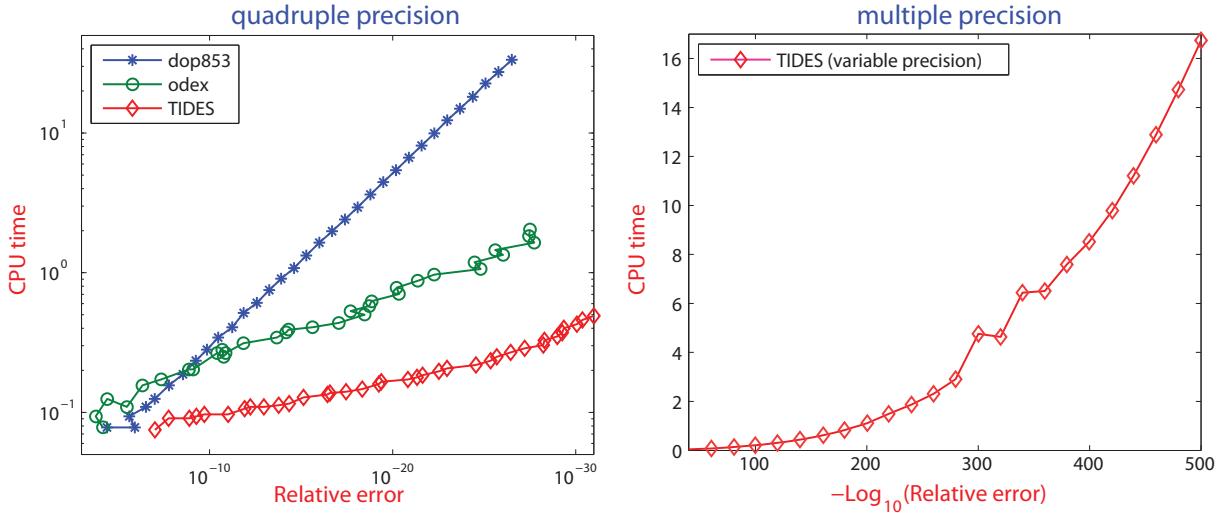


Figure 2: Left: Precision vs. CPU time diagram in quadruple precision for the numerical integration of the unstable periodic orbit LR for the Lorenz model using a Runge-Kutta code (`dop853`), an extrapolation code (`odex`) and a Taylor series method (`TIDES`). Right: Precision vs. CPU time diagram for the multiple-precision numerical integration of an unstable periodic orbit for the Lorenz model using the `TIDES` code.

The Taylor method is as follows [21, 24, 39]. Consider the initial value problem $\dot{\mathbf{y}} = \mathbf{f}(t, \mathbf{y})$. The value of the solution at t_i (that is, $\mathbf{y}(t_i)$) is approximated by \mathbf{y}_i from the n -th degree Taylor series of $\mathbf{y}(t)$ at $t = t_i$ (the function \mathbf{f} must be a smooth function). So, denoting $h_i = t_i - t_{i-1}$,

$$\begin{aligned} \mathbf{y}(t_0) &=: \mathbf{y}_0, \\ \mathbf{y}(t_i) &\simeq \mathbf{y}_{i-1} + \mathbf{f}(t_{i-1}, \mathbf{y}_{i-1}) h_i + \dots + \frac{1}{n!} \frac{d^{n-1} \mathbf{f}(t_{i-1}, \mathbf{y}_{i-1})}{dt^{n-1}} h_i^n =: \mathbf{y}_i. \end{aligned}$$

The problem is thus reduced to the determination of the Taylor coefficients $\{1/(j+1)! d^j \mathbf{f}/dt^j\}$. This may be done quite efficiently by means of the automatic differentiation (AD) techniques. Note that the Taylor method has several good features (for details see [21, 22, 24]).

In the Figure 2 we present some comparisons on the Lorenz model [67] for the classical Saltzman's parameter values using the Taylor method (`TIDES` code) and the well established codes `dop853` (a Runge-Kutta code) and `odex` (an extrapolation code) developed by Hairer and Wanner [53]. We observe that in quadruple precision the Taylor method is the fastest and,

as expected, the `odex` code is more efficient than the Runge-Kutta code (note that `odex` is a variable order code, as `TIDES`, and so it is more adaptable than the fixed order method). In double precision the most efficient code is the Runge-Kutta code, but for high precision the Taylor series method is the only reliable method among the standard methods. Note that the computer time for a high-precision numerical integration of one period ($T = 1.55865$) of the LR unstable periodic orbit (in symbolic dynamics notation one loop around the left equilibrium point, and one around the right one [80]) maintaining 500 digits is just around 16 seconds using a normal desktop computer, a quite reasonable time.

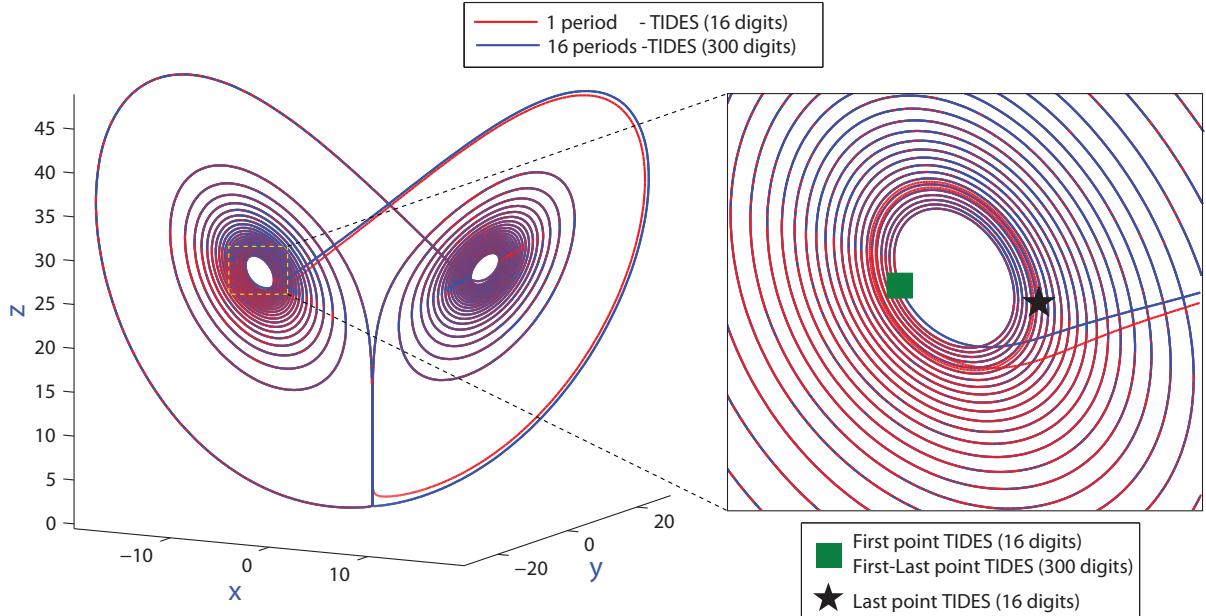


Figure 3: Numerical integration of the $L^{25}R^{25}$ unstable periodic orbit for the Lorenz model during 16 time periods using the `TIDES` code with 300 digits and 1 time period using double precision.

One may wonder whether such high accuracy is required in these kind of systems. To illustrate this need, we show in Figure 3 the results of 16 time periods computing using the the `TIDES` code with 300 digits and 1 time period using double precision for the numerical simulation of the $L^{25}R^{25}$ unstable periodic orbit for the Lorenz model. Note that we lose more than 16 digits on each period (the period of the orbit is $T = 33.890206423038$ and the largest Lyapunov exponent $\lambda = 0.958$, so $\exp(\lambda T) \approx 1.5324 \cdot 10^{16}$), and therefore it is not possible to simulate any period of this orbit in double precision. The double precision orbit is not periodic (see the zoom) and it also loses the symmetry of the correct orbit. The `TIDES` (Taylor series Integrator for Differential EquationS) used here [1, 27] is available from <http://gme.unizar.es/software/tides> (or send email to `tides@unizar.es` or `rbarrio@unizar.es`).

3.2 High precision arithmetic in recurrences

Computations involving recurrences are often highly unstable [49]. A classical example of unstable recurrence is the evaluation of the Bessel function of first kind $J_i(x)$ [3] by means of the three-term recurrence

$$J_{n+1}(x) = \frac{2n}{x} J_n(x) - J_{n-1}(x). \quad (3)$$

In this case, virtually all numerical significance is lost after just a few iterations, and, what's more, there is no way of using extended precision to improve the results. The reason of this disastrous build-up of errors [49] is due to the fact that the Bessel function of the first kind, but also the Bessel function of the second kind $Y_i(x)$ are solutions of the recurrence relation (3) and $J_i(x)/Y_i(x) \sim (x/2)^{2i}/(2(i!)^2)$ as $i \rightarrow \infty$, and so $J_i(x)$ is a (highly) minimal solution at infinity. This implies that any error is extremely amplified and the numerical solution goes quite fast to the dominant one. In this case we have to look for another completely different algorithm.

In other circumstances the recurrence is more stable, and high-precision arithmetic can produce useful results. One example here is the evaluation of coefficients of orthogonal polynomials. Such expansions have application in almost all mathematical and physical disciplines, including approximation theory, spectral methods, representation of potentials and others. In the last few years, researchers have studied different extensions, like orthogonal polynomials in Sobolev spaces [43]. One particular case of interest is when measures related to derivatives are purely atomic, with a finite number of mass points. That is, given a set of K evaluation points $\{c_1, \dots, c_K\}$ (the support of the discrete measure), a set of indexes that indicate the maximum order of derivatives in each evaluation point $\{r_1, \dots, r_K\}$, and a set of non-negative coefficients $\{\lambda_{ji} \mid j = 1, \dots, K; i = 0, \dots, r_j\}$, we define the Sobolev inner product

$$\langle p, q \rangle_W = \int_{\mathbb{R}} p(x) q(x) d\mu_0(x) + \sum_{j=1}^K \sum_{i=0}^{r_j} \lambda_{ji} p^{(i)}(c_j) q^{(i)}(c_j), \quad \lambda_{ji} \geq 0. \quad (4)$$

We are interested in evaluating a finite series of orthogonal polynomials with respect to a discrete Sobolev inner product. Some algorithms for such calculations were proposed in [26, 28], but the resulting algorithms are slightly unstable, and so, a combination of double and multiple precision is required. This may be done in such a way that the theoretical error bounds permit us to use high-precision just on the unstable cases, and so the computational complexity does not grows significantly.

In Figure 4 we show the behavior of some theoretical error bounds [28]: a backward error bound, the running error bound and the relative error in a multiple-precision evaluation of a Sobolev series. Note that we present relative error bounds and relative rounding errors, that is, for $q(x) \not\approx 0$ we divide by $|q(x)|$. We have up to degree 50 of the function $f(x) = (x+1)^2 \sin(4x)$ in Chebyshev-Sobolev orthogonal polynomials, considering one mass point $c = 1$ up to first derivative in the discrete part of the inner product. In the figures on the left we use double precision (53 bits on the mantissa) and on the right we use multiple precision (96 bits on the mantissa for $x < -0.5$ (on the left of the vertical line) and 64 for $x > -0.5$). The turning point $x = -0.5$ is the point where the relative running error in double precision is greater than 10^{-10} . These results make it clear that the combination of rounding error bounds (in this case

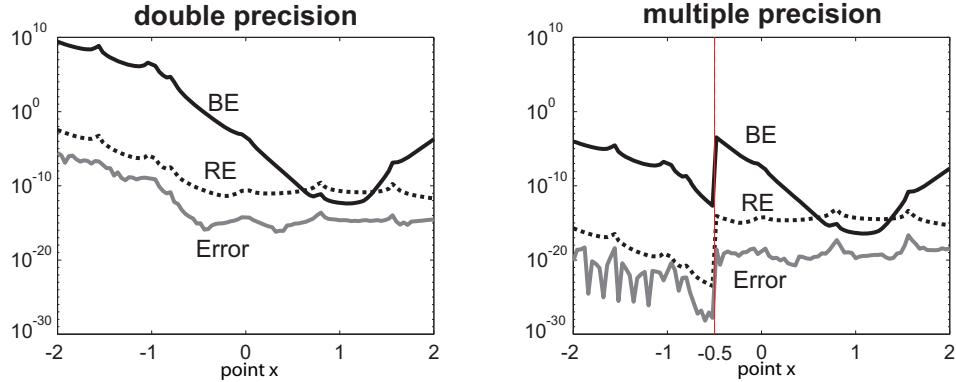


Figure 4: Behavior of the theoretical error bounds (BE a backward error bound and RE for the running error bound) and the relative error in the double- and multiple-precision evaluation of the Chebyshev-Sobolev approximation of degree 50 of the function $f(x) = (x + 1)^2 \sin(4x)$, where the discrete Sobolev measure have one mass point $c = 1$ up to 1st derivative in the discrete part of the inner product. In the figure on the left we use double precision and on the right multiple-precision (on the left of the vertical line we use 96 bits on the mantissa and 64 on the right part). (Reproduced with permission from [28].)

the running error bound) and multiple-precision libraries permits us to evaluate Sobolev series accurately.

Another situation where high precision is useful is in evaluating ill-conditioned polynomials. For instance, numerical errors are encountered when evaluating the polynomial $p(x) = (x - 0.75)^7(x - 1)^{10}$ close to one of its multiple roots. One solution is to find an optimal polynomial basis, although this may not be practical in many real-world situations. Another option is to use a good algorithm (e.g., Horner's algorithm for power series, the de-Cateljau's algorithm for the Bernstein basis and Clenshaw's algorithm for classical orthogonal polynomial basis), implemented with high-precision arithmetic. A third option, which is quite attractive when one does not want to deal with high-precision software, is to employ “compensated” algorithms that recently emerged in stability analysis [68, 72]. This approach permits one to use double precision arithmetic, yet still maintain the quality of the numerical evaluations with a relative error on the order of the rounding unit u , plus the conditioning of the problem times the square of the rounding unit. For instance, recently Graillat *et al.* [51] developed a “compensated” version of the Horner's algorithm. Also, H. Jiang *et al.* [60] developed a “compensated” version of Clenshaw's algorithm [38] to evaluate a finite series of Chebyshev orthogonal polynomials $p(x) = \sum_{j=0}^n a_j T_j(x)$. For this compensated algorithm (and all the other ones) it is possible to prove the following relative error bounds:

Theorem 2 [60] *Let $p(x) = \sum_{i=0}^n a_i T_i(x)$ be a polynomial in Chebyshev form. If the condition number for polynomial evaluation of $p(x)$ at entry x is defined by*

$$\text{cond}(p, x) = \frac{\tilde{p}(|x|)}{|p(x)|} = \frac{\sum_{j=0}^n |a_j| \tilde{T}_j(|x|)}{|\sum_{j=0}^n a_j T_j(x)|}, \quad (5)$$

with $\tilde{T}_j(|x|)$ the absolute polynomials associated with $T_j(x)$ [60], then the relative forward error bounds of the Clenshaw algorithm and compensated Clenshaw algorithm are such that

$$\frac{|\text{Clenshaw}(p, x) - p(x)|}{|p(x)|} \leq \mathcal{O}(u) \cdot \text{cond}(p, x), \quad (6)$$

$$\frac{|\text{CompClenshaw}(p, x) - p(x)|}{|p(x)|} \leq u + \mathcal{O}(u^2) \cdot \text{cond}(p, x). \quad (7)$$

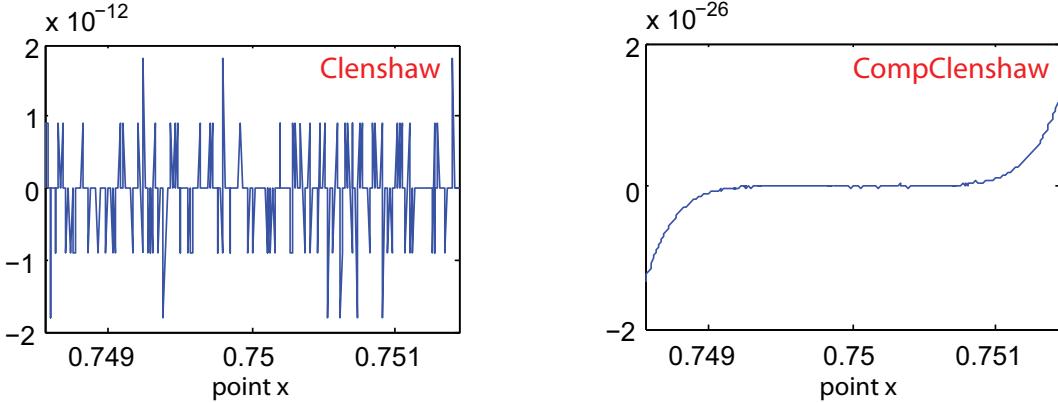


Figure 5: Evaluation of $p(x) = (x - 0.75)^7(x - 1)^{10}$ in the neighborhood of the multiple root $x = 0.75$, using Clenshaw (left) and Compensated Clenshaw (right). (Reproduced with permission from [60]).

This theorem shows one particularly nice feature of compensated algorithms, namely that the effect of the conditioning of the problem is delayed up to second order in the rounding unit u , yielding highly accurate (in relative error) computations.

Figure 5 presents the evaluation of the polynomial $p(x) = (x - 0.75)^7(x - 1)^{10}$ for 400 equally spaced points in the interval $[0.74855, 0.75145]$. It is clear that the compensated Clenshaw's algorithm gives a much smoother solution than the original Clenshaw's algorithm. Moreover, the relative error is always (except when $p(x)$ is very close to zero) of the order of the rounding unit u . This is often a crucial consideration in algorithms for locating zeros of polynomials in floating point arithmetic, because oscillations like the ones presented on the left figure can make impossible to obtain accurate results.

While compensated algorithms are often quite effective, they are not suitable for all situations, and so the use of high-precision software such as the QD library [59] is sometimes required.

3.3 High precision arithmetic in dynamical systems

In the words of Henri Poincaré, periodic orbits form the “skeleton” of a dynamical system and provide much useful information. Therefore, the search for periodic orbits is a quite old problem and numerous numerical and analytical methods have been designed for them. Here we mention

just two methods that have been used with high-precision in the literature: the Lindstedt-Poincaré technique [79] and one of the most simple and powerful method to find periodic orbits, namely the systematic search method [23], where one takes advantage of symmetries of the system to find symmetric periodic orbits [64].

Theorem 3 *Let $o(\mathbf{x})$ be an orbit of a flow of an autonomous vector field $d\mathbf{x}/dt = f(\mathbf{x})$ with a reversal symmetry S (thus $dS(\mathbf{x})/dt = -f(S(\mathbf{x}))$). Then, an orbit $o(\mathbf{x})$ intersects $\text{Fix}(S) := \{\mathbf{x} \mid S(\mathbf{x}) = \mathbf{x}\}$ in precisely two points if and only if the orbit is periodic (and not a fixed point) and symmetric with respect to S .*

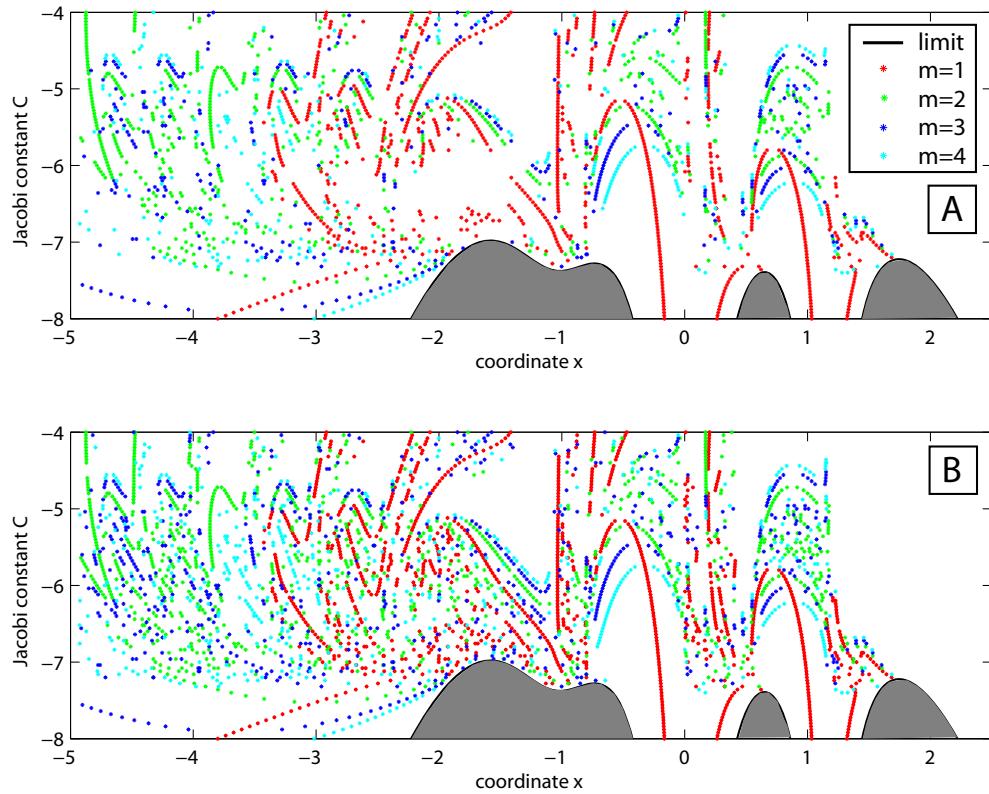


Figure 6: Symmetric periodic orbits (m denotes the multiplicity of the periodic orbit) in the most chaotic zone of the $7 + 2$ Ring problem using double (A) and quadruple (B) precision. (Reproduced with permission from [23]).

The above results were already known by Birkhoff, DeVogelaere and Strömgren (among others) and were used to find symmetric periodic orbits.

The usage of high-precision numerical integrators in the determination of periodic orbits is required in the search of highly unstable periodic orbits. For instance, in Figure 6 we show the computed symmetric periodic orbit for the $7 + 2$ Ring problem using double and quadruple precision [25]. The $(n + 2)$ -body Ring problem [25] describes the motion of an infinitesimal

particle attracted by the gravitational field of $n+1$ primary bodies, n in the vertices of a regular polygon that is rotating on its own plane about the center with a constant angular velocity. Each point on the figures corresponds to the initial conditions of one symmetric periodic orbit, and the grey area corresponds to regions of forbidden motion (delimited by the limit curve). Note that in order to avoid “false” initial conditions it is useful to check if the initial conditions generate a periodic orbit up to a given tolerance level. But in the case of highly unstable periodic orbits we may lose several digits in each period, so that double precision is not enough in many unstable cases, resulting in gaps in the figure.

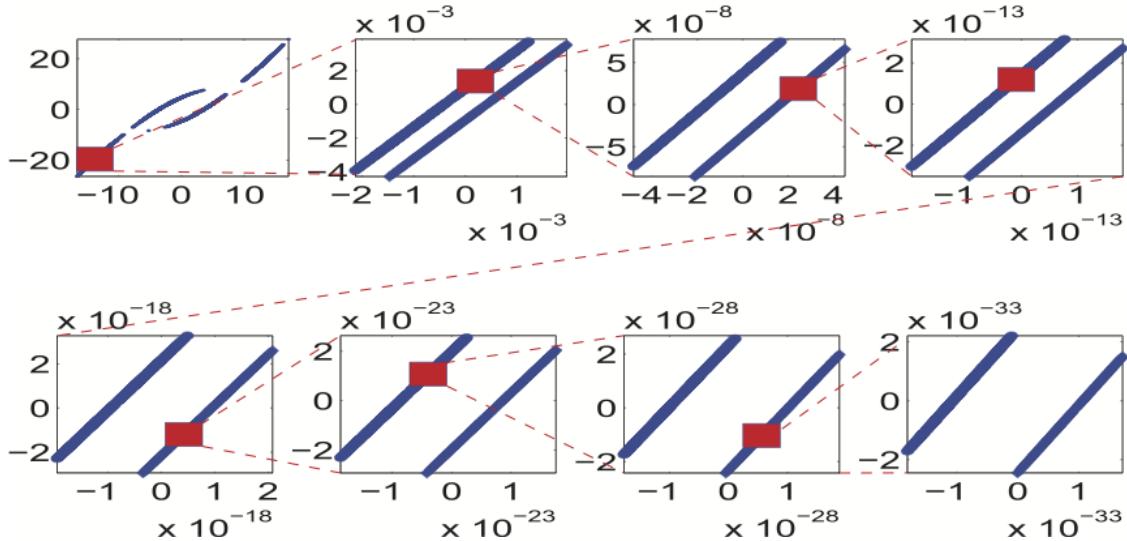


Figure 7: Fractal property of the Lorenz attractor. On the first plot, the intersection of an arbitrary trajectory on the Lorenz attractor with the section $z = 27$. The plot shows a rectangle in the $x - y$ plane. All later plots zoom in on a tiny region (too small to be seen by the unaided eye) at the center of the red rectangle of the preceding plot to show that what appears to be a line is in fact not a line. (Reproduced with permission from [81]).

The Lindstedt-Poincaré method [79] for computing periodic orbits is based on the Lindstedt-Poincaré technique of perturbation theory, Newton’s method for solving nonlinear systems and Fourier interpolation. D. Viswanath [80] uses this algorithm in combination with high-precision libraries to obtain periodic orbits for the Lorenz model at the classical Saltzman’s parameter values. This procedure permits one to compute, to high accuracy (more than 100 digits of precision), highly unstable periodic orbits (for instance the orbit with symbolic dynamics $LRL^2R^2 \dots L^{15}R^{15}$ has a leading characteristic multiplier 3.06×10^{59} , which means that we can expect that at each period we lose around 59 digits of precision). For these reasons, high-precision arithmetic plays a fundamental role in the study of the fractal properties of the Lorenz attractor (see Figure 7) and in a consistent formal development of complex singularities of the Lorenz system using psi series [80, 81].

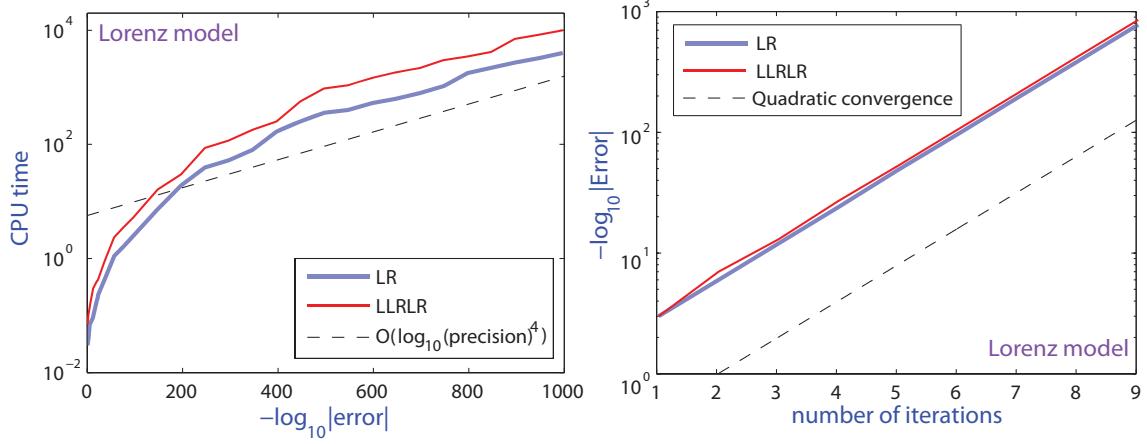


Figure 8: Computational relative error vs. CPU time and number of iterations in a 1000-digit computation of the periodic orbits LR and LLRLR of the Lorenz model. (Reproduced with permission from [2]).

A simpler option to compute high-precision periodic orbits has been proposed recently in [2], where the use of the Taylor series method permits to apply modified versions of the Newton method to obtain periodic orbits with more than 1000 precision digits. Figure 8 presents, as an example, the computational relative error vs. CPU time and number of iterations in a 1000-digit computation of the periodic orbits LR and LLRLR in the Lorenz model.

Another area of dynamical systems that often requires high precision is the study of splitting of separatrices in area preserving maps [50, 54]. Numerical difficulties arise because this phenomena can exhibit exponentially small splitting. One of the most common examples is the standard map defined by $(x, y) \mapsto (\hat{x}, \hat{y})$, where

$$\hat{y} = y + \varepsilon \sin x, \quad \hat{x} = x + \hat{y},$$

and ε is a small positive constant. This map can be obtained, for example, by a simple time discretization (a symplectic Euler of discretization step $\sqrt{\varepsilon}$) of the pendulum equation $\dot{x} = y, \dot{y} = \sin x$ [54]. The phase space structure of both systems, the continuous case and the map, are very different (except for small values of ε). In fact, the pendulum problem is an integrable system and its phase space is very regular (see Figure 9). There is a unique separatrix that connect the hyperbolic fixed point at 0 and at 2π , that is, the unstable manifold at 0 coincide with the stable manifold at 2π . When we see the map, the two manifolds do not coincide and so the separatrix splits (splitting of separatrices). Now we have transverse intersection points that gives homoclinic points and that imply the existence of complex dynamics or chaotic motion. Therefore the study of this phenomena of splitting of separatrices gives a deep information about the system, and so related with this, it is important to study the angle between the stable and the unstable separatrices at the intersection points. If the angle does not vanish we may affirm that this phenomena occurs. In Figure 9 we illustrate also the phenomena with two other maps (the quadratic map and the asymmetric cubic map [50]).

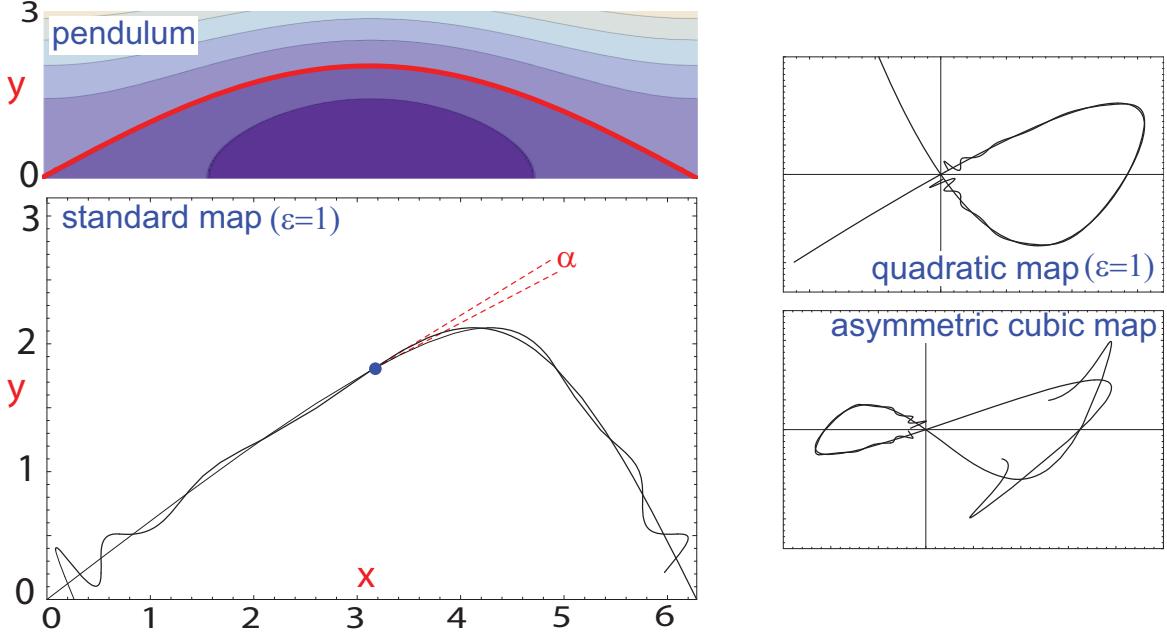


Figure 9: Left: Phase-space for the pendulum equations with the separatrix in red and the discrete version (standard map) for $\varepsilon = 1$ with the stable and the unstable separatrices. Right: stable and the unstable separatrices for the quadratic map and the asymmetric cubic map. (Partially reproduced with permission from [50].)

An asymptotic formula for the angle between the stable and the unstable separatrices for the standard map at the primary homoclinic point was given by Lazutkin [66]:

$$\alpha = \frac{\pi}{\varepsilon} e^{-\frac{\pi^2}{\sqrt{\varepsilon}}} (1118.8277059409\dots + \mathcal{O}(\sqrt{\varepsilon})).$$

As a result, the separatrices are transversal, but the angle between them is exponentially small compared to ε . This leads to severe problems in numerical simulations. Gelfreich and Simó [50] use a homoclinic invariant ω that gives the area of a parallelogram defined by two vectors tangent to the stable and the unstable manifolds at the homoclinic point. While ω in the standard map can be represented by an asymptotic series, one question is what happens when we use several generalizations of the standard map. In [50], the authors employed high-precision computation of the homoclinic invariant and consecutive extraction of coefficients of an asymptotic expansion, in order to obtain a numerical evidence that various different types of asymptotic expansions arise in this class of problems. These results are unachievable using standard double precision; in some numerical simulations 1000-digit precision was required. In the literature there are other numerous examples of high-precision computation of this phenomena of exponentially small splitting of separatrices.

3.4 High precision arithmetic in experimental mathematics

In this section we give a selection of five less directly applied applications:

Very high-precision computations (typically 100 to several thousand digits) have proven to be an essential tool in “experimental mathematics” [30, 7]. One of the key techniques used here is the PSLQ integer relation detection algorithm [14], which, given an n -long vector (x_i) of real numbers (presented as a vector of high-precision values), attempts to recover the integer coefficients (a_i) , not all zero, such that

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n = 0 \quad (8)$$

(to available precision), or else determines that there are no such integers (a_i) of a given size. Perhaps the best-known application of PSLQ in experimental mathematics is the 1996 computer-based discovery of what is now known as the “BBP” formula for π :

$$\pi = \sum_{k=0}^{\infty} \frac{1}{16^k} \left(\frac{4}{8k+1} - \frac{2}{8k+4} - \frac{1}{8k+5} - \frac{1}{8k+6} \right). \quad (9)$$

This formula has the remarkable property that it permits one to calculate binary (or hexadecimal) digits beginning at the n -th digit, without needing to calculate any of the first $n-1$ digits, using a simple scheme that requires very little memory and no multiple-precision arithmetic software [6], [30, pp. 135–143]. Since 1996, numerous other formulas of this type have been found using PSLQ and then subsequently proven [5]. For example, a binary formula is known for Catalan’s constant $G = \sum_{n \geq 0} (-1)^n / (2n+1)^2$, and both binary and ternary (base-3) formulas are known for π^2 [13].

In 2010, Tse Wo Zse, a researcher with Yahoo! Cloud Computing, used a variant of this formula to compute a string of hexadecimal digits of π beginning at the 500 trillionth digit (corresponding to the two quadrillionth binary digit) [78]. In 2011, IBM researchers used BBP-type formulas to calculate base-64 digits of π^2 ; equation, base-729 digits of π^2 , and base-4096 digits of G , in each case beginning at the ten trillionth position and validated by a second independent computation [13]. These computations were performed on IBM’s benchmark machine in Rochester Minnesota.

In an unexpected turn of events, it has been found that these computer-discovered formulas have implications for the age-old question of whether (and why) the digits of certain well-known math constants are statistically random. In particular, one of the present researchers and Richard Crandall found that the question of whether constants such as π and $\log 2$ are 2-normal (i.e., every string of m binary digits appears, in the limit, with frequency 2^{-m}) reduces to a conjecture about the behavior of a certain explicit pseudorandom number generator that is related to the respective BBP-type formula for that constant [15], [30, pp. 163–178]. This same line of investigation has led to a formal proof of normality for an uncountably infinite class of explicit real numbers [16], the simplest instance of which is

$$\alpha_{2,3} = \sum_{n=1}^{\infty} \frac{1}{3^n 2^{3^n}},$$

which is provably 2-normal (and provably *not* 6-normal) [9]. The normality of π and some related constants is analyzed statistically and graphically in [10]. In particular, very strong evidence for the normality of π based on analysis of nearly 16 trillion bits is presented. An illustration is provided in Figure 10, while ten billion digits may be visually explored at <http://gigapan.org/gigapans/99214>.

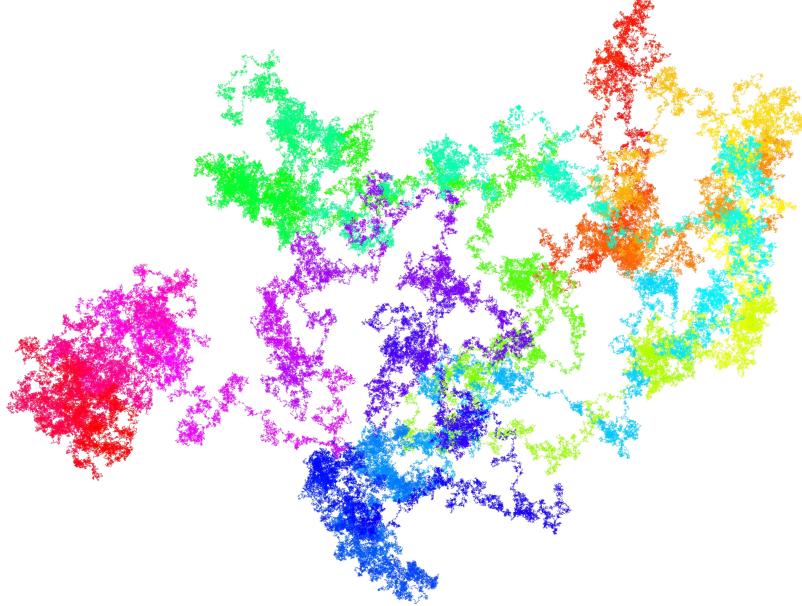


Figure 10: A walk on the first billion base 4 digits of Pi. The digits 0,1,2,3 are coded to up, down, left or right respectively.

Very high-precision computations, combined with the PSLQ algorithm, have been remarkably effective in recognizing (in terms of analytic formulas) certain classes of definite integrals that arise in mathematical physics settings. In one study, the tanh-sinh quadrature scheme [18, 77], implemented using the ARPREC software [19], was employed to study the following classes of integrals [11, 12]. Here, the D_n integrals arise in the Ising theory of mathematical physics, and the C_n have tight connections to quantum field theory:

$$\begin{aligned} C_n &= \frac{4}{n!} \int_0^\infty \cdots \int_0^\infty \frac{1}{\left(\sum_{j=1}^n (u_j + 1/u_j)\right)^2} \frac{du_1}{u_1} \cdots \frac{du_n}{u_n} \\ D_n &= \frac{4}{n!} \int_0^\infty \cdots \int_0^\infty \frac{\prod_{i < j} \left(\frac{u_i - u_j}{u_i + u_j}\right)^2}{\left(\sum_{j=1}^n (u_j + 1/u_j)\right)^2} \frac{du_1}{u_1} \cdots \frac{du_n}{u_n} \\ E_n &= 2 \int_0^1 \cdots \int_0^1 \left(\prod_{1 \leq j < k \leq n} \frac{u_k - u_j}{u_k + u_j} \right)^2 dt_2 dt_3 \cdots dt_n, \end{aligned}$$

where (in the last line) $u_k = \prod_{i=1}^k t_i$.

Evaluating these n -dimensional integrals to high precision presents a daunting computational challenge, but in the first case we were able to show that the C_n integrals can be written as one-dimensional integrals:

$$C_n = \frac{2^n}{n!} \int_0^\infty p K_0^n(p) dp,$$

where K_0 is the *modified Bessel function* [3]. After computing C_n to 1000-digit accuracy for various n , we were able to identify the first few instances of C_n in terms of well-known constants, e.g., $C_4 = 7\zeta(3)/12$, where ζ denotes the Riemann zeta function. When we computed C_n for fairly large n , for instance

$$C_{1024} = 0.63047350337438679612204019271087890435458707871273234\dots,$$

we found that these values rather quickly approached a limit. By using the new edition of the *Inverse Symbolic Calculator*, available at <http://carma-lx1.newcastle.edu.au:8087>, this numerical value can be identified as

$$\lim_{n \rightarrow \infty} C_n = 2e^{-2\gamma},$$

where γ is Euler's constant, which we were subsequently able to prove [11]. Some results were also obtained for the D_n and E_n , although these computations were considerably more difficult.

A more recent study considered, for complex s , the n -dimensional *ramble integrals* [8]

$$W_n(s) = \int_{[0,1]^n} \left| \sum_{k=1}^n e^{2\pi x_k i} \right|^s dx, \quad (10)$$

which occur in the theory of uniform random walk integrals in the plane, where at each step a unit-step is taken in a random direction. Integrals such as (10) are the s -th moment of the distance to the origin after n steps. It is shown in [35] that when $s = 0$ the first derivatives of these integrals can be written as

$$W'_n(0) = \log(2) - \gamma - \int_0^1 (J_0^n(x) - 1) \frac{dx}{x} - \int_1^\infty J_0^n(x) \frac{dx}{x} \quad (11)$$

$$= \log(2) - \gamma - n \int_0^\infty \log(x) J_0^{n-1}(x) J_1(x) dx, \quad (12)$$

where $J_n(x)$ denotes the Bessel function of the first kind.

Due to the oscillatory nature of these integrals, they present substantial challenges for high-precision numerical integration. One approach that we have found effective for these integrals is known as the Sidi mW extrapolation algorithm, as described in a 1994 paper by Lucas and Stone [65] (which in turn is based on two earlier papers by Sidi [73, 74]), combined with tanh-sinh quadrature and Gaussian quadrature [8]. Using this scheme, we were able to evaluate these integrals to 1000-digit accuracy, at least when n is odd, using the ARPREC software [19]. This scheme is not very effective when n is even, but in this case we were able to compute

modestly high precision results (50–100 digits) by employing asymptotic formulas for the Bessel function. In response to this ineffectiveness, Sidi [75] has made an analysis and proposed a more sophisticated scheme which should redress the situation.

These results were used to verify several other studies. For instance, our result when $n = 6$ matched to 80-digit precision a computation based on a conjecture due to Villegas [36]. Similarly, for $n = 4$ our 80-digit result agrees to full precision with the closed form given in [35].

Our calculations also confirmed, to 600-digit precision, the following amazing conjecture based on one of Villegas, [36]:

$$W_5'(0) \stackrel{?}{=} \left(\frac{15}{4\pi^2} \right)^{5/2} \int_0^\infty \{ \eta^3(e^{-3t})\eta^3(e^{-5t}) + \eta^3(e^{-t})\eta^3(e^{-15t}) \} t^3 dt, \quad (13)$$

where

$$\eta(q) = q^{1/24} \prod_{n \geq 1} (1 - q^n) = q^{1/24} \sum_{n=-\infty}^{\infty} (-1)^n q^{n(3n+1)/2}. \quad (14)$$

While the intuitive genesis of equation (13) lies in algebraic K-theory, it is fair to say that there is no inkling of how to prove it.

The research on ramble integrals also led us to examine moments of elliptic integral functions of the form [8]:

$$I(n_0, n_1, n_2, n_3, n_4) = \int_0^1 x^{n_0} K^{n_1}(x) K'^{n_2}(x) E^{n_3}(x) E'^{n_4}(x) dx, \quad (15)$$

where the elliptic functions K, E and their complementary versions are given by:

$$\begin{aligned} K(x) &= \int_0^1 \frac{dt}{\sqrt{(1-t^2)(1-x^2t^2)}} & K'(x) &= K(\sqrt{1-x^2}) \\ E(x) &= \int_0^1 \frac{\sqrt{1-x^2t^2}}{\sqrt{1-t^2}} dt & E'(x) &= E(\sqrt{1-x^2}). \end{aligned} \quad (16)$$

To better understand these product integrals, we computed a large number of them (4389 individual integrals in total) to extreme precision — 1500 to 3000-digit precision — using the ARPREC software. We then discovered, using PSLQ, thousands of intriguing relations between

these numerical values, including the following limited selection [8]:

$$\begin{aligned}
& 81 \int_0^1 x^3 K^2(x) E(x) dx \stackrel{?}{=} -6 \int_0^1 K^3(x) dx - 24 \int_0^1 x^2 K^3(x) dx \\
& + 51 \int_0^1 x^3 K^3(x) dx + 32 \int_0^1 x^4 K^3(x) dx \\
& - 243 \int_0^1 x^3 K(x) E(x) K'(x) dx \stackrel{?}{=} -59 \int_0^1 K^3(x) dx + 468 \int_0^1 x^2 K^3(x) dx \\
& + 156 \int_0^1 x^3 K^3(x) dx - 624 \int_0^1 x^4 K^3(x) dx - 135 \int_0^1 x K(x) E(x) K'(x) dx \quad (17) \\
& - 20736 \int_0^1 x^4 E^2(x) K'(x) dx \stackrel{?}{=} 3901 \int_0^1 K^3(x) dx - 3852 \int_0^1 x^2 K^3(x) dx \\
& - 1284 \int_0^1 x^3 K^3(x) dx + 5136 \int_0^1 x^4 K^3(x) dx - 2592 \int_0^1 x^2 K^2(x) K'(x) dx \\
& - 972 \int_0^1 K(x) E(x) K'(x) dx - 8316 \int_0^1 x K(x) E(x) K'(x) dx. \quad (19)
\end{aligned}$$

These identities led to a detailed study by James Wan [82], who has been able to prove many but by no means all of them.

4 Other brief examples

We briefly summarize here a number of other applications of high-precision arithmetic that have been reported to us. For additional details, please see the listed references.

4.1 Supernova simulations

Recently Edward Baron, Peter Hauschildt, and Peter Nugent used the QD package [59] to solve for the non-local thermodynamic equilibrium populations of iron and other atoms in the atmospheres of supernovae and other astrophysical objects [20, 55]. Iron, for example, may exist as Fe II in the outer parts of the atmosphere, but in the inner parts Fe IV or Fe V could be dominant. Introducing artificial cutoffs leads to numerical glitches, so it is necessary to solve for all of these populations simultaneously. Since the relative population of any state from the dominant stage is proportional to the exponential of the ionization energy, the dynamic range of these numerical values can be large. Among various potential solutions, these authors found that using double-double (or, in some cases, quad-double) arithmetic to be the most straightforward and effective.

4.2 Climate modeling

It is well-known that climate simulations are fundamentally chaotic — if microscopic changes are made to the present state, within a certain period of simulated time the future state is completely

different. Indeed, ensembles of these calculations are required to obtain statistical confidence in global climate trends produced from such calculations. As a result, climate modeling codes quickly diverge from any “baseline” calculation, even if only the number of processors used to run the code is changed. For this reason, it is often difficult for researchers to compare results, or even to determine whether they have correctly deployed their code on a given system. Recently Helen He and Chris Ding found that almost all of the numerical variation in an atmospheric code occurred in a long inner product loop in the data assimilation step and in a similar operation in a large conjugate gradient calculation. He and Ding found that employing double-double arithmetic for these loops dramatically reduced the numerical variability of the entire application, permitting computer runs to be compared for much longer run times than before [58].

4.3 Coulomb n -body atomic system simulations

Numerous computations have been performed using high-precision arithmetic to study atomic-level Coulomb systems. For example, Alexei Frolov of Queen’s University in Ontario, Canada has used high-precision software to solve the generalized eigenvalue problem $(\hat{H} - E\hat{S})C = 0$, where the matrices \hat{H} and \hat{S} are large (typically $5,000 \times 5,000$ in size) and very nearly degenerate. Until recently, progress in this area was severely hampered by the numerical difficulties induced by these nearly degenerate matrices. Frolov found that by employing 120-digit arithmetic, “we can consider and solve the bound state few-body problems which have been beyond our imagination even four years ago” [17, 48].

4.4 Studies of the fine structure constant of physics

In the past few years, significant progress has been achieved in using high-precision arithmetic to obtain highly accurate solutions to the Schrodinger equation for the lithium atom. In particular, the non-relativistic ground state energy has been calculated to an accuracy of a few parts in a trillion, a factor of 1500 improvement over the best previous results. With these highly accurate wave functions, researchers have been able to test the relativistic and QED effects at the 50 parts per million (ppm) level and also at the one ppm level [83]. Along this line, a number of properties of lithium and lithium-like ions have also been calculated, including the oscillator strengths for certain resonant transitions, isotope shifts in some states, dispersion coefficients and Casimir-Polder effects between two lithium atoms. When some additional computations are completed, the fine structure constant may be obtained to an accuracy of 16 parts per billion [84].

4.5 Scattering amplitudes of quarks, gluons and bosons

An international team of physicists working on the Large Hadron Collider (LHC) is computing scattering amplitudes involving quarks, gluons and gauge vector bosons, in order to predict what results could be expected on the LHC. By default, these computations are performed using conventional double precision (64-bit IEEE) arithmetic. Then if a particular phase space point is deemed numerically unstable, it is recomputed with double-double precision. These researchers

expect that further optimization of the procedure for identifying unstable points may be required to arrive at an optimal compromise between numerical accuracy and performance. Their objective is to design a procedure where the number of digits in the higher precision calculation is dynamically set according to the instability of the point [46]. Three related applications of high-precision arithmetic are given in [29, 69, 40].

4.6 Detecting Strange Nonchaotic Attractors

In the study of dynamics of dissipative systems the detection of the attractors is quite important, because they are the visible invariant sets of the dynamics of the problem. An attractor is defined as *strange* if it is not a piecewise smooth manifold and *chaotic* if any orbit on it exhibits sensitive dependence on initial conditions. All the first examples of strange attractors in the literature were strange chaotic attractors, but soon some strange nonchaotic attractors (SNAs) were identified [52]. Several authors suggested that in the transition to chaos in quasiperiodically forced dissipative systems, in particular in the so called fractalization route in which a smooth torus seems to fractalize, strange nonchaotic attractors appear. In [56], Haro and Simó showed that in truth some of these attractors are nonstrange. These authors found that multiprecision arithmetic with more than 30 digits was needed to reliably study this behavior at very small scales. Therefore, in this case (and in many cases) the SNAs is not produced via the fractalization route, but what is evident is that this phenomena requires a very high-precision numerical simulation to give a correct information of what really happens on the systems.

5 Conclusion

For most scientific and engineering computations, either IEEE 32-bit or (more often) 64-bit floating-point arithmetic provides sufficient accuracy. But for a rapidly expanding body of applications, even 64-bit floating-point arithmetic is not sufficient. Typical situations that may require higher-precision arithmetic include:

1. Ill-conditioned linear systems.
2. Large summations.
3. Long-time simulations.
4. Large-scale simulations.
5. Resolving small-scale phenomena.
6. “Experimental mathematics” computations.

Performing such calculations with high-precision arithmetic once was a major challenge, but such tasks have been greatly facilitated by recently developed software packages that include high-level language translation modules to minimize the conversion effort. Run times often increase substantially when using high-precision arithmetic, but in many cases it suffices to

convert only a handful of key routines, and other portions of the computation can be done with conventional arithmetic. Moreover, thanks to the sophistication of modern computer algebra packages, it is often possible to do a portion of the high-precision component symbolically — thereby improving both accuracy and run times.

In this paper, we have described a number of specific applications where these situations arise, and where high-precision arithmetic is required. These include: (a) solution of certain types of ordinary differential equations, (b) evaluation of recurrences, (d) detection of exponentially small phenomena in dynamical systems, (d) computer-based discovery of new mathematical relations (such as the “BBP” formula for π), (e) supernova simulations, (f) climate modeling, (g) Coulomb n -body atomic system simulations, and others.

It is worth noting that all of these examples have arisen just in the past 10–15 years. Thus, we may be witnessing the birth of a new era of scientific computing, in which the numerical precision required for a computation is as important to the program design as are the algorithms and data structures.

We conclude that high-precision arithmetic facilities are now an indispensable component of a modern large-scale scientific computing environment. We hope that our survey and analysis of these computations will be useful to help further develop these facilities into truly usable and easy-to-use computational tools, and to identify additional classes of scientific computations where these tools are useful.

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