Clean Application Compartmentalization with SOAAP

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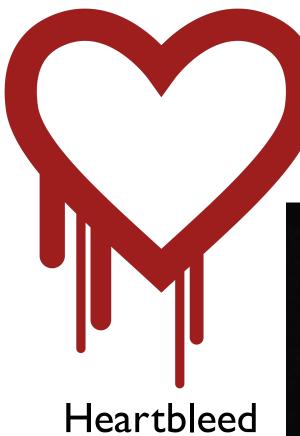
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Vulnerabilities galore...



[Insert next big vulnerability here]





Mitigate both known and unknown vulnerabilities





Principle of least privilege

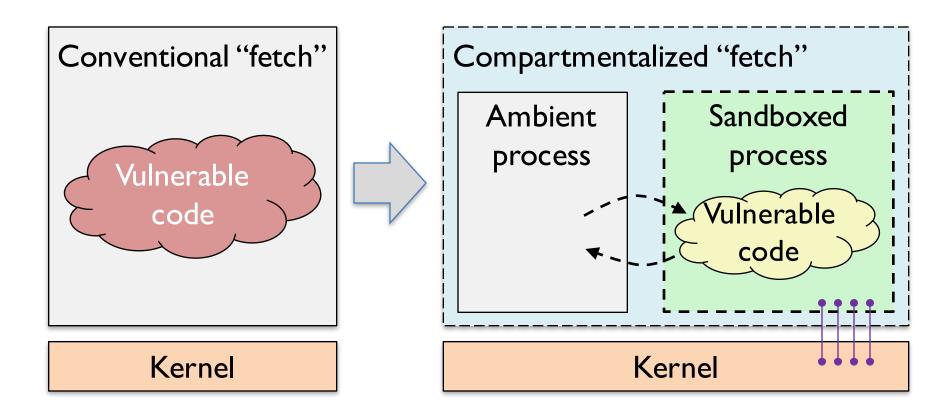
Every program and every privileged user of the system should operate using the least amount of privilege necessary to complete the job.

Saltzer 1974 - CACM 17(7)
Saltzer and Schroeder 1975 - Proc. IEEE 63(9)
Needham 1972 - AFIPS 41(1)





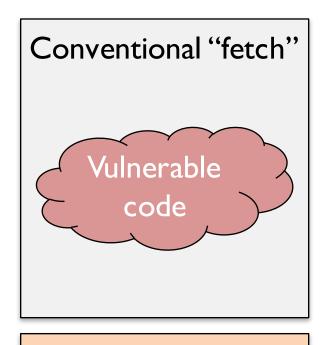
Application Compartmentalization







Application Compartmentalization



When a conventional application is compromised, its ambient rights are leaked to the attacker, e.g., full network and file system access.

Kernel

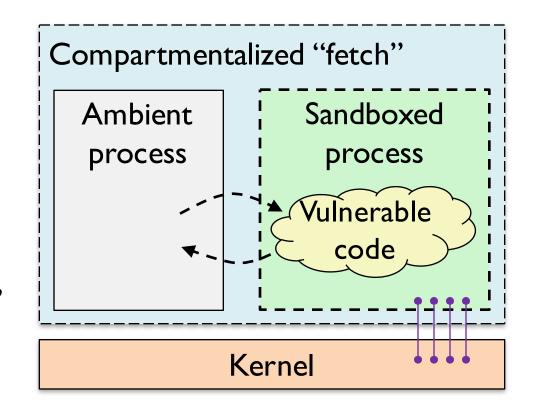




Application Compartmentalization

When a compartmentalized application is compromised, only rights held by the exploited component leak to the attacker.

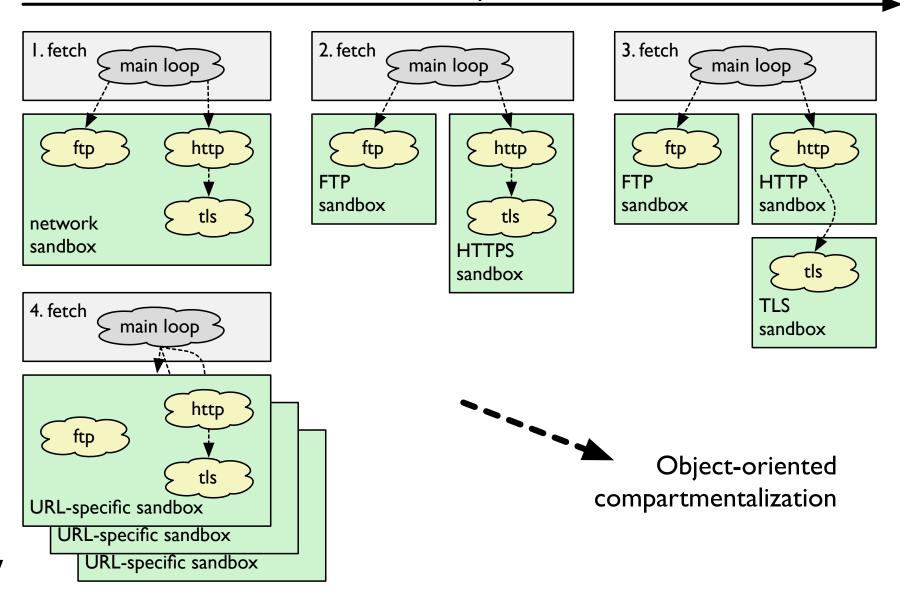
Most vulnerabilities will no longer yield significant rights, and attackers must exploit many vulnerabilities to meet their goals.





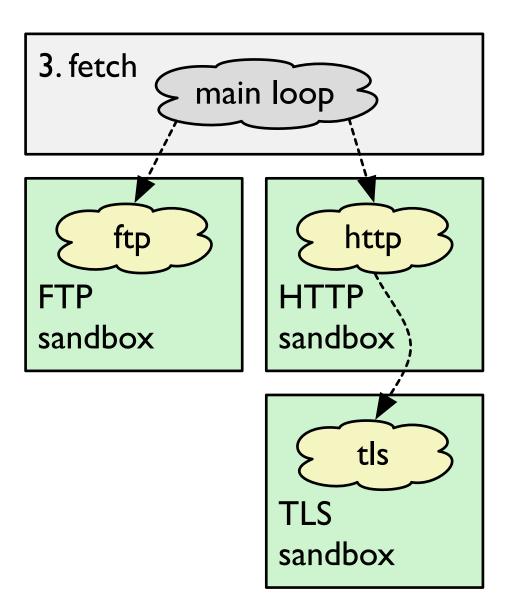


Code-centered compartmentalization



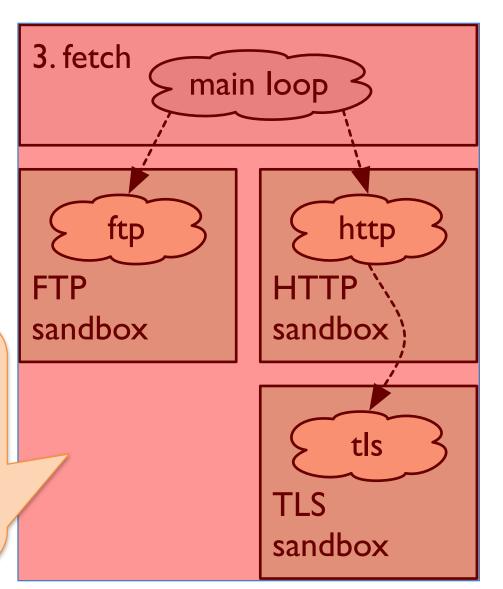










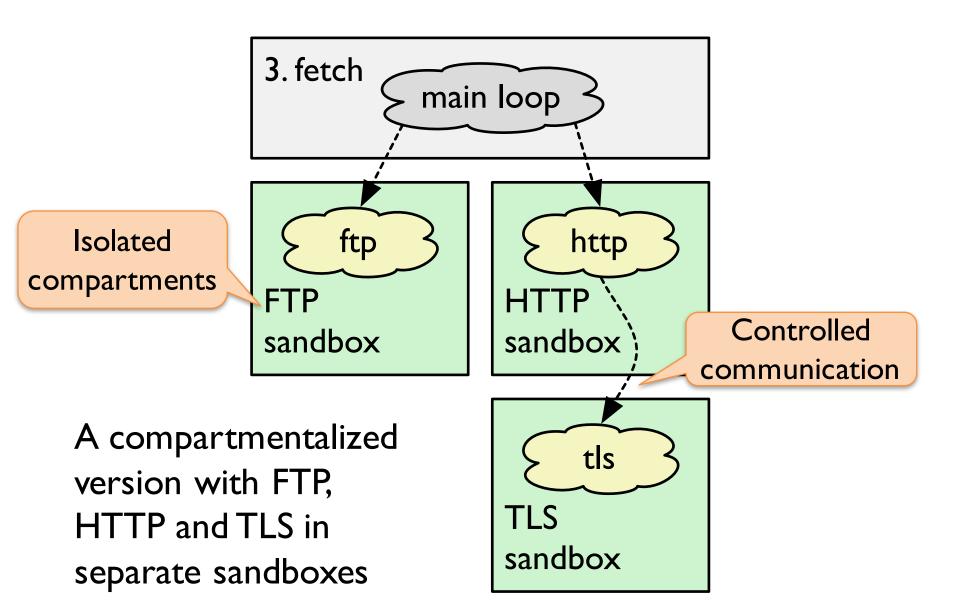




Before, a single exploit anywhere in the program would have leaked ambient rights



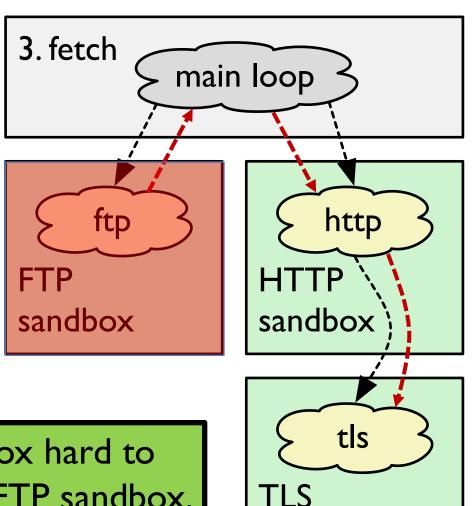






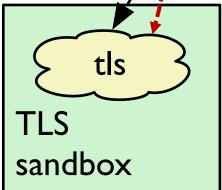


"exploit FTP to access TLS state"





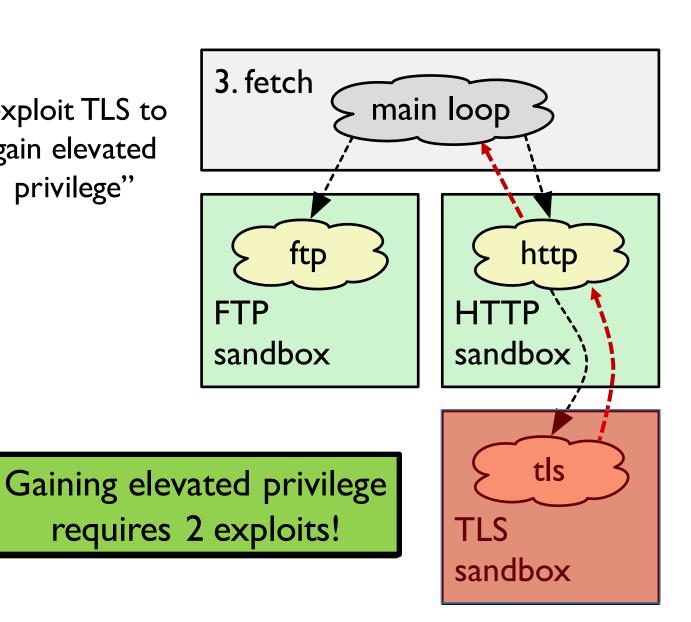
TLS sandbox hard to access from FTP sandbox. Requires 3 exploits!







"exploit TLS to gain elevated privilege"







Compartmentalization is hard!

• "local" program turned into a distributed one

- Preserving functional correctness
- Mapping security model to sandboxing substrate
- Different compartmentalization tradeoffs

Hard to change and maintain





Onus is on the developer

"It's up to you to understand these elements [of App Sandbox] and then to use your understanding to create a plan for adopting App Sandbox."

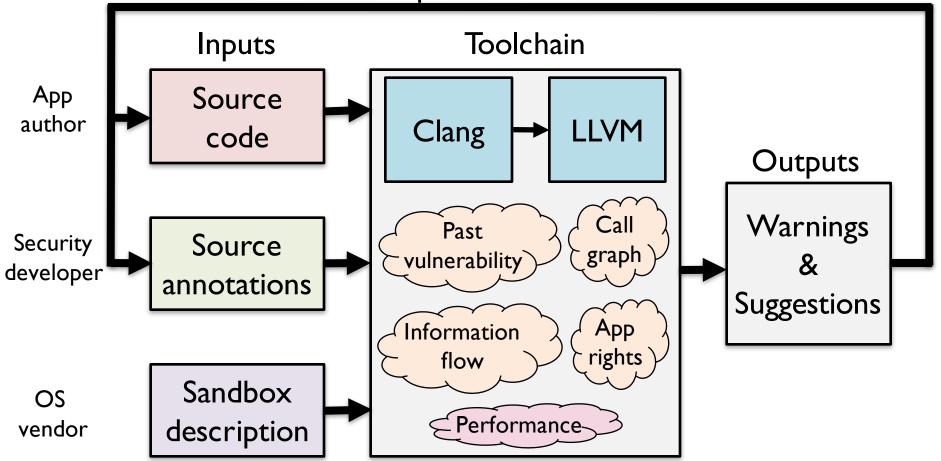
- Apple App Sandbox Design Guide





Security-Oriented Analysis of Application Programs

Repeated refinement







Example: Confidentiality/IFC

```
char* server_key __soaap_classify("secret");
extern void compute session key(char*, char*);
void main() {
 while (...) {
  accept_connection();
  soaap_sandbox_persistent("session")
void accept connection() {
 char session_key[256] __soaap_private;
 compute session key(session key, server key);
```





Example: Confidentiality/IFC

```
Classify tag
char* server_key __soaap_classify("secret");
extern void compute_session_key(char*,char*);
void main() {
 while (...) {
  accept_connection();
                               Sandbox entrypoint
  soaap_sandbox_persistent("session")
                                                Private state
void accept connection() {
 char session_key[256] __soaap_private;
 compute session key(session key, server key);
```





Example: Past vulnerabilities/ supply-chain trojans

```
soaap_provenance("some vendor")
  soaap_sandbox_ephemeral("parser")
void parse(__soaap_fd_permit(read) int ifd, DOMTree* out) {
  if (...) {
     soaap_vuln_pt("CVE-2005-ABC");
  soaap_vuln_fn("CVE-2005-DEF")
void not sandboxed() {
```





Example: Past vulnerabilities/ supply-chain trojans

```
Provenance
                                                Sandbox entrypoint
        soaap_provenance("some vendor")
                                                         Delegated rights
        soaap_sandbox_ephemeral("parser")
     void parse(__soaap_fd_permit(read) int ifd, DOMTree* out) {
        if (...) {
           soaap_vuln_pt("CVE-2005-ABC");
                                                     Past vulnerable point
                   Past vulnerable
                      function
        soaap_vuln_fn("CVE-2005-DEF")
     void not sandboxed() {
```



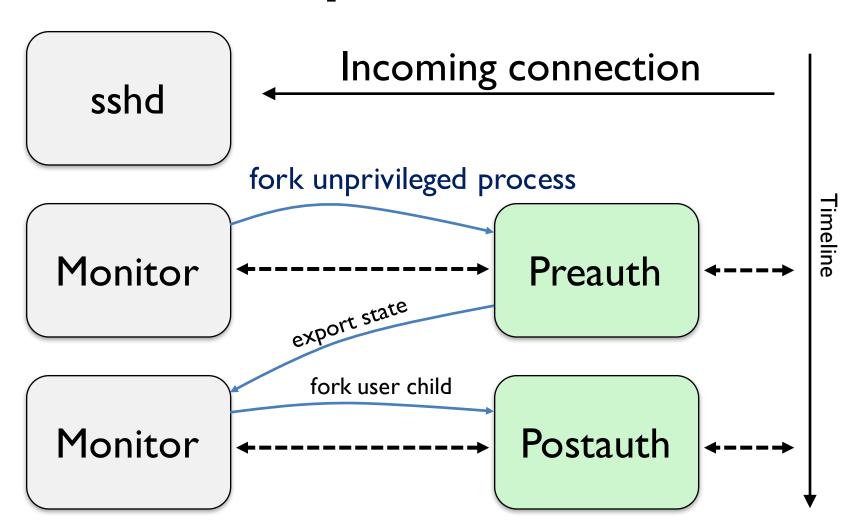


Case studies

- Fetch design-space exploration
- Okular large-scale new compartmentalization
- OpenSSH long-term maintenance
- Chromium sandboxing effectiveness



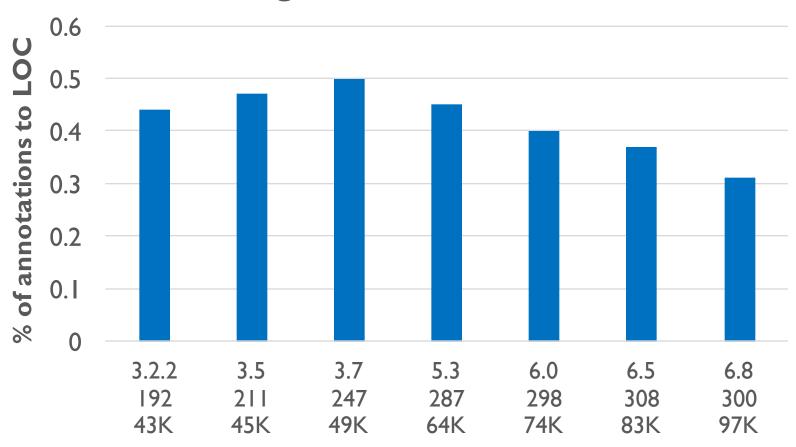








Percentage of annotations to LOC

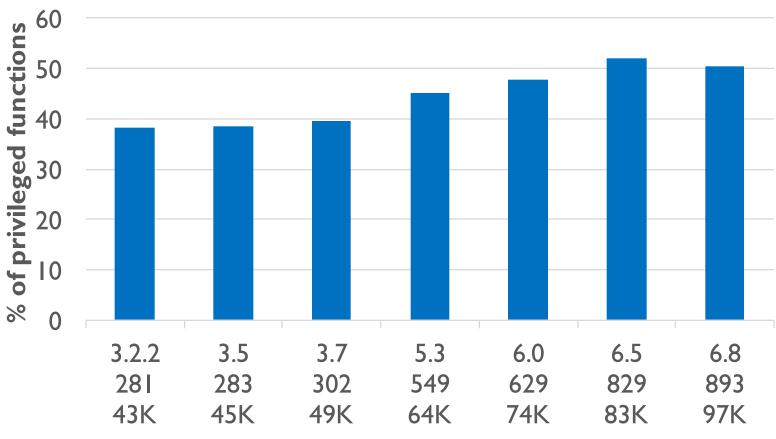


OpenSSH Version, No. of annotations and LOC





Percentage of privileged functions

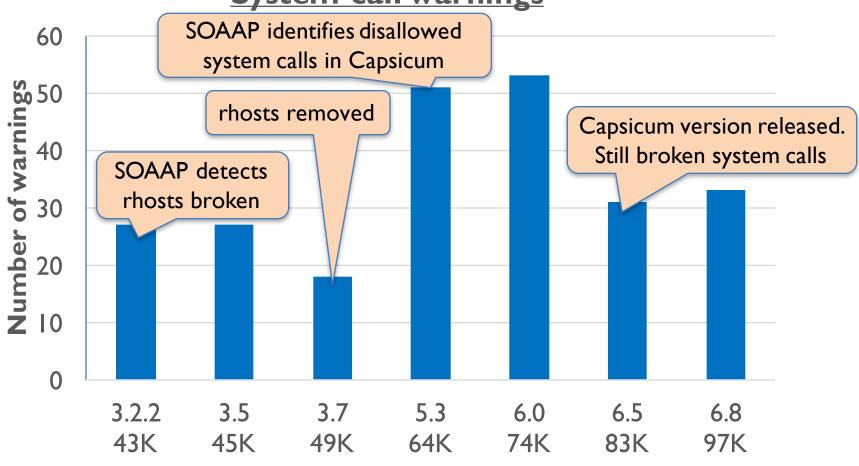


OpenSSH Version, No. of privileged funcs and LOC





System call warnings



OpenSSH Version and LOC





Chromium

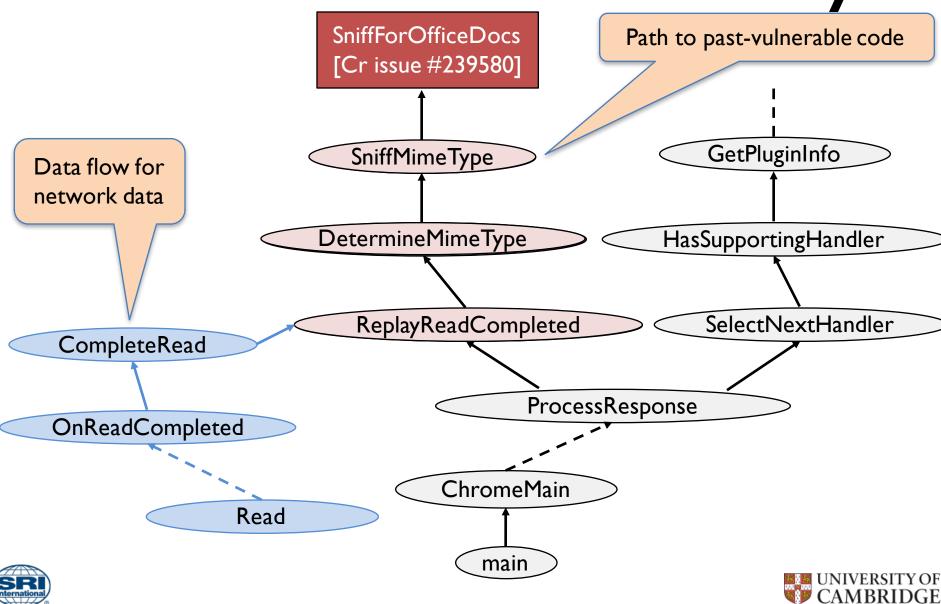
Browser process

User Interface Network Storage Renderer process Renderer process HTML Renderer **HTML** Renderer **V8** Engine **V8** Engine **DOM** Bindings **DOM Bindings**



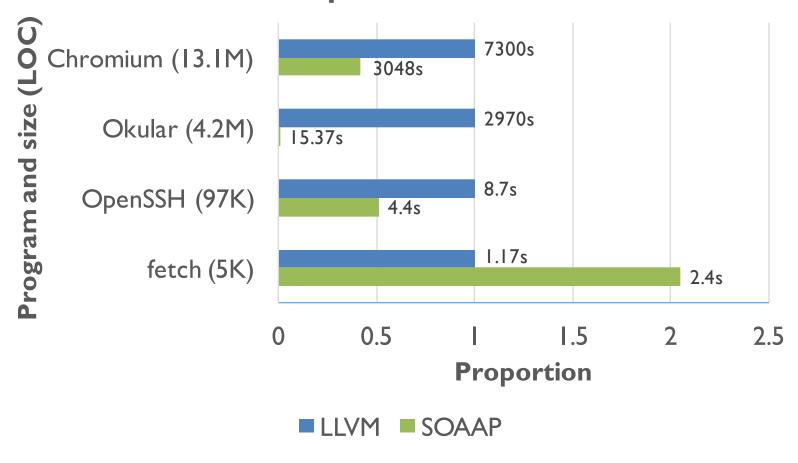


New Chromium vulnerability?



SOAAP tool performance

SOAAP running times as proportion of compile+link time







Conclusion

- Application compartmentalization is important but hard to get right
- Conceptual framework gives structure
- SOAAP enables reasoning about tradeoffs
- SOAAP can scale to multi-million LoC programs

http://bit.ly/soaap







Example SOAAP output (I)

```
$ make soaap
*** Sandboxed method "accept connection" read data
*** value of class: [secret] but only has clearances for: []
+++ Line 28 of file session.c
*** Sandboxed method "accept connection" executing in
*** sandboxes: [session] may leak private data through
*** the extern function "compute session key"
+++ Line 28 of file session.c
*** Persistent sandbox "session" contains private data that
*** may leak when the sandbox is reused. Consider using an
*** ephemeral sandbox instead or scrub the memory region
*** before control returns
```







Example SOAAP output (2)

```
$ make soaap
*** Method "not sandboxed" had past vulnerability
*** "CVE-2005-DEF" but is not sandboxed. Another
*** vulnerability here could leak ambient authority
*** to an attacker including full network and file
*** system access
*** Sandboxed method "parse" has a past-vulnerability
*** annotation for "CVE-2005-ABC". Another vulnerability
*** here would only leak the following:
+++ Read access to file descriptor "ifd"
```





Chromium security disparity

	OS	Sandbox	LoC	FS	IPC	NET	S≠S′	Priv
DAC	Windows	DAC ACLs	22,350	\triangle	\triangle	X	×	\
	Linux	chroot()	600	>	X	×	/	X
MAC	Mac OS X	Sandbox	560	/	<u>^</u>	/	/	/
	Linux	SELinux	200	>	\triangle	/	×	X
Сар	Linux	seccomp	11,300	\triangle	/	/	/	/
	FreeBSD	Capsicum	100	>	/	/	\	>

Watson et al., "Capsicum: practical capabilities for UNIX"



