

# Codeforces 837E (2100)

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<https://codeforces.com/problemset/problem/837/E>

Accepted: <https://codeforces.com/contest/837/submission/360157274>

## §1 Solution

### §1.1 Explanation

#### Lemma 1.1

For any two natural numbers  $a$  and  $b$ ,

$$\gcd(a, b) \mid (b - k \gcd(a, b))$$

for any  $k \in \mathbb{Z}^+$

However, it is not true that for all  $k$

$$\gcd(b - k \gcd(a, b)) = \gcd(a, b)$$

So we need to compute this extra factor in the gcd of  $(b - k \gcd(a, b))$  and  $a$ . Or in other words, we need to find the minimum  $k$  such that,

$$\gcd\left(\frac{a}{\gcd(a, b)}, \frac{b}{\gcd(a, b)} - k\right) > 1$$

However we know that,

$$\gcd\left(\frac{a}{\gcd(a, b)}, \frac{b}{\gcd(a, b)} - k\right) \mid \frac{a}{\gcd(a, b)}$$

Therefore we could iterate over the divisors of

$$\frac{a}{\gcd(a, b)}$$

and take the minimum  $k$ . Iterating over divisors should take  $\mathcal{O}(\sqrt{a})$  time and we will do this atmost  $\mathcal{O}(\log a)$  times. Therefore, the total time complexity of this algorithm is  $\mathcal{O}(\sqrt{a} \log a)$

### §1.2 Code

```
1 void solve() {  
2     ll a, b;  
3     std::cin >> a >> b;  
4 }
```

```
5  ll ans = 0;
6  while (b > 0) {
7      ll d = std::gcd(a, b);
8
9      a /= d;
10     b /= d;
11
12     if (a == 1) {
13         ans += b;
14         break;
15     }
16
17     ll best = LLONG_MAX;
18     for (ll i = 1; i * i <= a; i++) {
19         if (a % i == 0) {
20             if (i > 1) {
21                 best = std::min(best, b % i);
22             }
23
24             if (i * i != a) {
25                 best = std::min(best, b % (a / i));
26             }
27         }
28     }
29
30     ans += best;
31     b -= best;
32 }
33
34 std::cout << ans << '\n';
35 }
```