

$$= 2 \cdot 2^m (2^m \cdot 2 - 1) + (2^{m+1})^2$$

$$= 2 \cdot [2^{2m+1} - 2 \cdot 2^m + 2^{2m+2}]$$

$$= 2^{m+1} (2^{m+1} + 2^{m+1} - 1)$$

$$= 2^{m+1} (2 \cdot 2^{m+1} - 1)$$

$$P_2 + P_1 =$$

$$2 =$$

$$P_2 = 2 \cdot \text{truborg} = \text{truborg} \leftarrow b = 1 \text{ (good)}$$

$$\text{truborg} \leftarrow \text{truborg} + \text{mid} = \text{mid}$$

$$P_1 \leftarrow P_1 + 2 =$$

$$P_1 = 2 =$$

$$201 =$$

$$P_2 = \text{truborg} = \text{truborg} \leftarrow b = 1 \text{ (good)}$$

$$\text{truborg} \leftarrow \text{truborg} + \text{mid} = \text{mid}$$

$$P_2 \leftarrow P_2 + 201 =$$

$$201 =$$

$$201 \text{ return } 92$$

In the loop:

The loop will repeat a times

For times (t operators) & multiplication

addition & P

$$n \cdot (1 + 2) = \text{multiplication}$$





Show:

$$\begin{aligned} S(k+1) &= S(k) + 2k+1 \quad (\text{rr}) \\ &= k^2 + 2k+1 \quad (\text{ih}) \\ &= (k+1)^2 \end{aligned}$$

(41)

$$\begin{aligned} S(1) &= 1 & S(2) &= 3 \\ S(n) &= S(n-1) + S(n-2) \quad n \geq 3 \end{aligned}$$

Explain: A string of length  $n$  which has no two consecutive zeros can be generated in 2 ways:

+ Append 1 at the end of a string length  $n-1$  which has no two consecutive zeros  
 $\Rightarrow$  there are  $S(n-1)$  strings

+ Append 10 at the end of a string length  $n-2$  which has no two consecutive zeros  
 $\Rightarrow$  there are  $S(n-2)$  strings

$$\Rightarrow S(n) = S(n-1) + S(n-2)$$

(48)

Base case:  $S(1) = 1$  ( $2 \cdot 1 - 1 = 1$  (true))

Assume:  $S(k) = k$  ( $2k - 1$ ) with  $k \geq 2^m$

Show:

$$S(2^{m+1}) = 2 \cdot S(2^m) + (2^{m+1})^2 \quad (\text{rr})$$



$$\Leftrightarrow \log_{\frac{6}{5}} 4 + 1 \leq n$$

$$\Leftrightarrow 8.6 \leq n$$

$$\Rightarrow \boxed{n=9}$$

Conclusion

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$$S(1) = 1$$

$$S(n) = S(n-1) + n \quad n \geq 2$$

Base case:  $n=1$ :  ~~$S(1)$~~   $S(1) = \frac{1 \cdot (1+1)}{2} = 1$

Assume:  $S(k) = \frac{k(k+1)}{2}$

Show:

$$\begin{aligned} S(k+1) &= S(k) + k+1 \text{ (rr)} \\ &= \frac{k(k+1)}{2} + k+1 \end{aligned}$$

$$\begin{aligned} &= \cancel{k} \left( \frac{\cancel{k}+1}{2} + 1 \right) (k+1) \left[ \frac{\cancel{k}}{2} + 1 \right] \\ &= \frac{(k+1)(k+2)}{2} \end{aligned}$$

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$$S(1) = 1$$

$$S(n) = S(n-1) + 2n-1$$

Base case:  $n=1$ :  $S(1) = 1^2 = 1$

Assume:  $S(k) = k^2$





b)

Base case:  $n=1$ :  $SC(1) = 10^{1+2} = 1000$  (true)

Assume:  $SC(k) = 10^{k+2}$

Show:

$$SC(k+1) = 10 \cdot SC(k) \text{ (recurrence relation)}$$

$$= 10 \cdot 10^{k+2} \text{ (inductive hyp)}$$

$$= 10^{k+3}$$

c)  ~~$S(21) = 10^{21+2} = 10^{23}$~~

(21)

a)  $S(1) = 3$

$$S(n) = 5 \cdot S(n-1) - 6^{n-2} \quad n \geq 2$$

b)

Base case:  $n=1$ :  $S(1) = 4 \cdot 5^{1-1} - 6^{1-1} = 3$  (true)

Assume:  $S(k) = 4 \cdot 5^{k-1} - 6^{k-1}$

Show:

$$S(k+1) = ~~4 \cdot 5~~ 5 \cdot S(k) - 6^{k-1} \text{ (rr)}$$

$$= 5 \cdot (4 \cdot 5^{k-1} - 6^{k-1}) - 6^{k-1} \text{ (ih)}$$

$$= 4 \cdot 5^k - 5 \cdot 6^{k-1} - 6^{k-1}$$

$$= 4 \cdot 5^k - 6^k$$

c)  $S(n) \leq 0$

$$\Leftrightarrow ~~4 \cdot 5~~ 4 \cdot 5^{n-1} - 6^{n-1} \leq 0$$

$$\Leftrightarrow 4 \leq \left(\frac{6}{5}\right)^{n-1}$$

$$\Leftrightarrow \log \left(\frac{6}{5}\right)^4 \leq n-1$$



①

Base case:  $n = 1 : S(n) = 5 \cdot 1 = 5$  (true)

Assume:  $S(k) = 5 \cdot k$

Show:  $S(k+1) = S(k) + 5$  (recurrence relation)  
 $= 5 \cdot k + 5$  (inductive hypothesis)  
 $= 5(k+1)$

⑥

Base case:  $n = 1 : S(1) = 1^2 = 1$  (true)

Assume:  $S(k) = k^2$

Show:  $S(k+1) = S(k) + (2k+1)$   
 $= k^2 + 2k + 1$   
 $= (k+1)^2$

⑫

Base case:  $n = 1 : S(1) = 3^{1-1} \cdot (1+1)! = 2$  (true)

Assume:  $S(k) = 3^{k-1} \cdot (k+1)!$

Show:  $S(k+1) = 3(k+2) \cdot S(k)$  (recurrence relation)  
 $= 3(k+2) \cdot 3^{k-1} \cdot (k+1)!$  (inductive hyp)  
 $= 3^k \cdot (k+2)!$

⑮

a)  $S(1) = 1000$   
 $S(n) = 10 \cdot S(n-1) \quad n \geq 2$