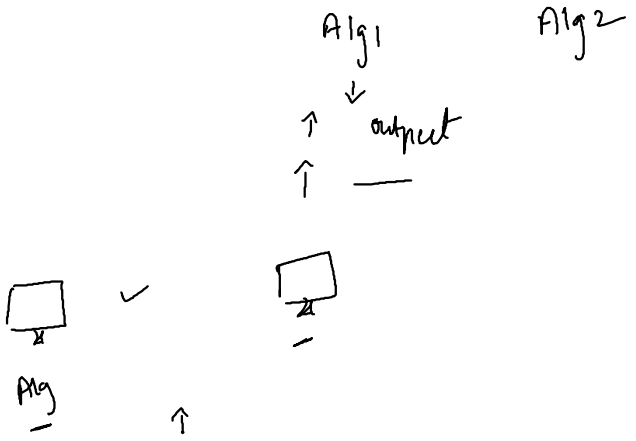
Asymptotic Notations

TC => O(n), Ω(n), θ(n) } ×

1) Big O = Worst case
2) Big Ω = Best case ✓
3) Big Theta(θ) = Tight estimate (worst + best)



why AN?

5 { Alg1 → 1 sec ✓
5 { Alg2 → 10 sec

1 ... worse — ?
×

tool

1) Upper Bound
maximum
metric 50
≤ 50 51 ×
Pizza → 30 min
O

2) Lower Bound
Overestimate
Pizza => 30 min
to
(=> 2 hrs)
✓ Ω

3) Tight Bound
Exact steps
Bus 10 min
Θ

n=5
0,1,2,3,4 Code Example

```
for(int i=0 ; i<n ; i++) {
    cout<<i;
}
```

O(n)
Ω O(n²) ×
O(n)

TB => 20 min ✓
UB => 30 min ✓
LB => 2 ghante ✓

1) Big O Notation (WorstCase) - Upper Bound (O)

Maximum time
-ve ← → +ve

TC = O(n)

My 1 ghante,

Big O
Maths → Big O
Industry → order of

17 MD → Big O

Algi → + additional tareeke se chai banana

1 → t_1 O(n)

2 → $2t_1$

3 → $3t_1$

:

n → nt_1

$T(n) = nt_1 + b$

Big O → mujhe worst case ka andaza de do → maximum

$= n$

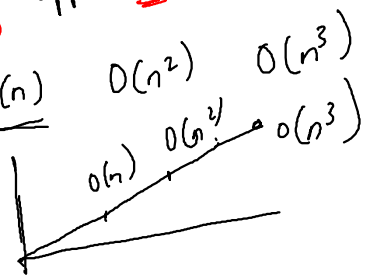
$\Rightarrow O(n)$ → ic (2 min) X (min X)

Big O

Upper Bound

Upper Limit

Note: $O(n) < O(n^2) < O(n^3)$



$O(n) < O(n^2) < O(n^3)$

O(n) O(n²)

min

$O(n) \rightarrow O(n^2)$

Example

1) $f(n) = 3n + 5$

$\Rightarrow n$

$\Rightarrow O(n)$

Bigo → maximum time → worst case → upper bound

guarantee → code isse zyaada slow nahi chalega.

2) ID → Big O → order of → worst case

3 log

1st Ananya → n (fastest)

2nd Bhavna → n^2 (thoda slow)

3rd Suresh → n^3 (bahut slow)

1 Race

3 Race

Ananya → Price

O(n) $O(n^2)$ $O(n^3)$

O(1)

O(1), O(n)

$T(n) = 2n + 5$

$\Rightarrow n$

$\Rightarrow O(n)$

O(n)

$O(n^2)$ Medium

$O(n^3)$ Bahut slow

Big O ✓

Code kitna slow ho sakta hai worse-case me, uska upper limit batata hai

→ Code isse zyaada slow nahi chalega

→ Shayad fast chale → but not slower.

Example: Sharmaji ka beta (Rohan)

Bus

Big O ⇒ 30 min

✓10 min ✓

✓30 min ✓

$f(n) = n^2 + n + 3$

$= n^2 + n$

$\Rightarrow n^2$

Max $\Rightarrow O(n^2)$ ✓

$n = 2$ ✓

$n^2 = 4$

$n = 2$

✓ **f(n) is O(g(n)) if there exist constants c and n_0 such that:**

$f(n) \leq c \times g(n)$ for all $n \geq n_0$

$f(n) \Rightarrow$ zyada se zyaada

↙ ↘ input

time

$g(n) \rightarrow$ Big O

$f(n) \leq c\,g(n)$

Example

$f(n) = n^2 + n + 3$

$= O(n^2)$

$g(n) = n^2$

$n = 1$, $c = 5$

$5 \leq 1$ ✗

$5 \leq 5$ ✓

$c = 5$ ∀ $n \geq n_0$

$f(n) = 1^2 + 1 + 3 \Rightarrow 5$

$g(n)\ n^2 = 1^2 = 1$

$= 5(1) \Rightarrow 5$

27 $f(n) = n^3 + 1$

$O(n^3)$ ✓

✓

$g(n) = n^3$

$f(n) \leq c\,g(n)$

$n \geq n_0$

$c = 1$, $c = 2$

$1^3 + 1 = 2$

$1^3 = 1$ ✗

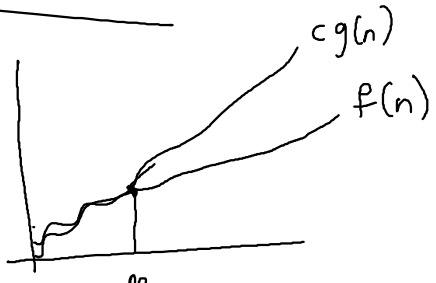
$1^3 + 1 = 2$

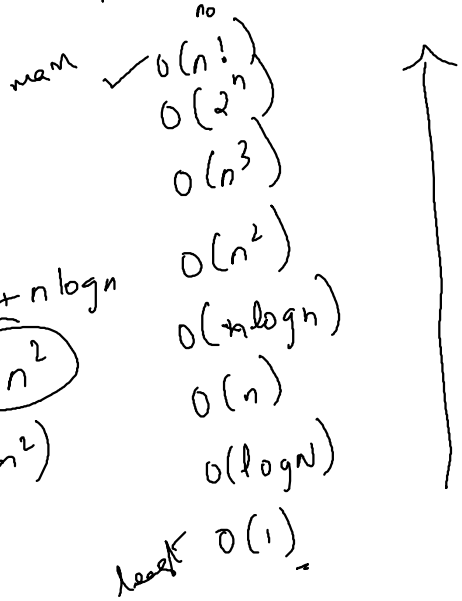
$2^3 \Rightarrow 8$

$2 \leq 8$ ✓

$O(n^3)$ → $O(n^4)$ ⇒ $O(n^5)$ — Big

→ $O(n^3)$





⇒ Different Big O

1) Constant time $O(1)$
2) Linear time $O(n)$
3) Quadratic $O(n^2)$
4) Cubic $O(n^3)$
5) logarithmic $O(\log n)$

$n + \log N + n\log n + n^2$ ⇒ $O(n^2)$

2) Big Omega Notation Ω (Best case / Lower bound)

↳ sabse kam time mai kitna jaldi kaam hosakta hai

lucky = favorite

key = 3

⇒

-1

$\Omega g(n)$

f(n) is Ω(g(n)) if there exist constants c and n_0 such that:
$c \times g(n) \leq f(n)$ for all $n \geq n_0$

1) $f(n) = 5n^2 + 4n + 2$
$= \Omega(n^2)$ ✓

$5n^2$ → quadratic
$4n$ → Linear time
$+2$ constant

$f(n) \geq n^2$ ∀ $n \geq$ some value

$n^2 \leq f(n)$ ∀ $n \geq$ some n_0

$n = 1000$
$5(1000)^2 + 4(1000) + 2$
$= 5000000 + 4000 + 2$
$= 5004002$

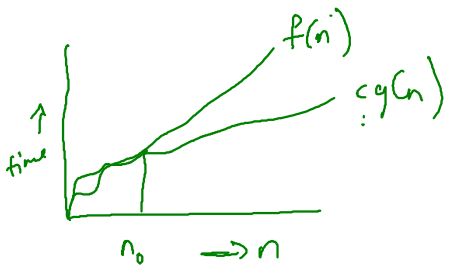
$f(n) = 5n^2 + 4n + 1$

$= \Omega(n^2)$

$g(n) = n^2$

$c = 1$ =)

$5n^2 \geq n^2$, $\boxed{c = 1}$



$f(n) \in \Omega(n^2)$

3) Big Theta Θ → Tight bound

↳ Exact steps.

25-30

$f(n) = n^2 + n + 3$

Big O = n^2 ✓

Big Ω = n^2 ✓

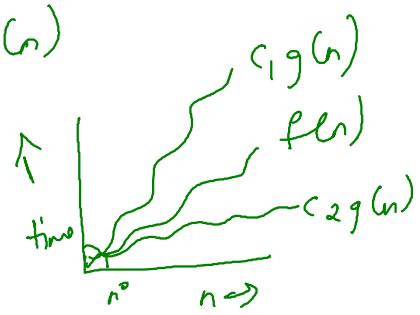
Big theta Θ = n^2

$f(n)$ is $\Theta g(n)$ if it is both $Og(n)$ & $\Omega g(n)$

$f(n) \leq c_1 g(n)$

$c_2 g(n) \leq f(n)$

$c_2 g(n) \leq f(n) \leq c_1 g(n)$



Big O → worst case (zyada se zyada time)

Big Ω → Best case (kam se kam time)

Big Θ → Tight bound (fix time range)