

No-repeat Substring (hard)

We'll cover the following ^

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Problem Statement

Given a string, find the **length of the longest substring**, which has **no repeating characters**.

Example 1:

Input: String="aabccbb"

Output: 3

Explanation: The longest substring without any repeating characters is "abc".

Example 2:



Input: String="abbbb"

Output: 2

Explanation: The longest substring without any repeating characters is "ab".

Example 3:

Input: String="abccde"

Output: 3

Explanation: Longest substrings without any repeating characters are "abc" & "cde".

Try it yourself

Try solving this question here:

Java	Python3	JS	C++
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```
1 import java.util.*;
2
3 class NoRepeatSubstring {
4     public static int findLength(String str) {
5         // TODO: Write your code here
6         return -1;
7     }
8 }
9
```



Solution

This problem follows the **Sliding Window** pattern, and we can use a similar dynamic sliding window strategy as discussed in Longest Substring with K Distinct Characters

(<https://www.educative.io/collection/page/5668639101419520/5671464854355968/5698217712812032/>).

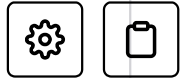
We can use a **HashMap** to remember the last index of each character we have processed. Whenever we get a repeating character, we will shrink our sliding window to ensure that we always have distinct characters in the sliding window.

Code

Here is what our algorithm will look like:

Java	Python3	C++	JS
<pre>1 import java.util.*; 2 3 class NoRepeatSubstring { 4 public static int findLength(String str) { 5 int windowStart = 0, maxLength = 0; 6 Map<Character, Integer> charIndexMap = new HashMap<>(); 7 // try to extend the range [windowStart, windowEnd] 8 for (int windowEnd = 0; windowEnd < str.length(); windowEnd++) { 9 char rightChar = str.charAt(windowEnd); 10 // if the map already contains the 'rightChar', shrink the window from the start 11 // we have only one occurrence of 'rightChar' 12 if (charIndexMap.containsKey(rightChar)) { 13 // this is tricky; in the current window, we will see the duplicate character 'rightChar' 14 // and if 'windowStart' is already ahead of the index of 'rightChar', 15 // we should move 'windowStart' to the index of 'rightChar' + 1 16 windowStart = Math.max(windowStart, charIndexMap.get(rightChar) + 1); 17 } 18 charIndexMap.put(rightChar, windowEnd); // insert the 'rightChar' at 'windowEnd' 19 // update the maximum length so far 20 maxLength = Math.max(maxLength, windowEnd - windowStart + 1); 21 } 22 return maxLength; 23 } 24 }</pre>			

```
18     maxLength = Math.max(maxLength, windowEnd - windowStart);
19 }
20
21 return maxLength;
22 }
23
24 public static void main(String[] args) {
25     System.out.println("Length of the longest substring:");
26     System.out.println("Length of the longest substring:");
27     System.out.println("Length of the longest substring:");
28 }
29 }
30
```



Time Complexity

The above algorithm's time complexity will be $O(N)$, where 'N' is the number of characters in the input string.

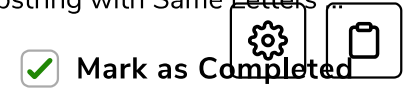
Space Complexity

The algorithm's space complexity will be $O(K)$, where K is the number of distinct characters in the input string. This also means $K \leq N$, because in the worst case, the whole string might not have any repeating character, so the entire string will be added to the **HashMap**. Having said that, since we can expect a fixed set of characters in the input string (e.g., 26 for English letters), we can say that the algorithm runs in fixed space $O(1)$; in this case, we can use a fixed-size array instead of the **HashMap**.

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