

Week 1 May 6

Wednesday, May 10, 2023 5:09 PM

1. Consider the irrational number Pi the ratio of a circle's circumference to its diameter?

1. How can you compute Pi via a digital computer? *Hint: what is the series expansion for Pi?*

Series Expansion for π :

$$\frac{1}{1+w} = 1 - w + w^2 - w^3 + w^4 - \dots = \sum_{n=0}^k (-1)^n (w)^n = \sum_{n=0}^k (-w)^n$$

for $w = x^2$

$$\frac{1}{1+x^2} = 1 - x^2 + x^4 - x^6 + x^8 - \dots = \sum_{n=0}^k (-w^2)^n$$

Integrate both sides

$$\arctan x = x - \frac{x^3}{3} + \frac{x^5}{5} - \frac{x^7}{7} + \dots$$

for $x=1$

$$\frac{\pi}{4} = 1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \dots = \sum_{n=0}^k (-1)^n \frac{1}{2n+1}$$

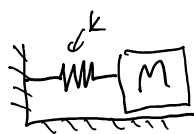
$$\pi = 4 \sum_{n=0}^k (-1)^n \frac{1}{2n+1}$$

Steps: - raise to the power - multiply by (-1)
 - Add - Sum
 - divide

2. How can you compute Pi via an analog computer? *Hint: what are the differential equations for sinusoidal motion? How can this be used to compute Pi?*

analog computers use physical phenomena (e.g. electrical, mechanical)

One such mechanical system is a mass on a spring



assuming frictionless

Hooke's Law

$$F = -kx$$

$$a = -\frac{k}{m}x$$

$$\Rightarrow x'' = -\frac{k}{m}x \quad \leftarrow \text{second order homogeneous eq}$$

$$x'' + \frac{k}{m}x = 0$$

$$\Rightarrow c^2 + \frac{k}{m} = 0$$

$\frac{k}{m} \geq 0 \therefore$ complex roots

$$c = \pm \sqrt{-\frac{k}{m}}$$

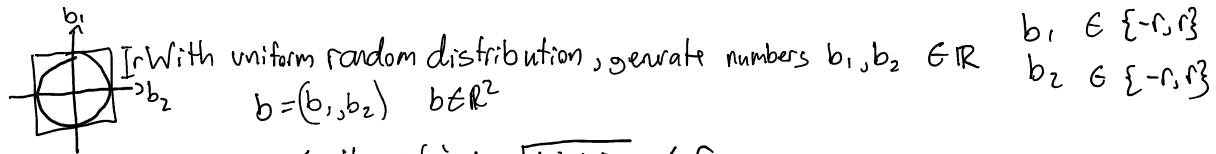
$$c = \pm \sqrt{\frac{k}{m}} i$$

$$x = e^{0t} [C_1 \cos(\sqrt{\frac{k}{m}} t) + C_2 \sin(\sqrt{\frac{k}{m}} t)]$$

$$x = C \sin(\omega t + \phi)$$

3. How can you compute Pi with the use of a random number generator? *Hint: what is the ratio the area of a square to that of an*

3. How can you compute Pi with the use of a random number generator? ^{steps: - measure, - multiply, - add} Hint: what is the ratio the area of a square to that of an inscribed circle? If you draw two random numbers, how can you tell if they define a coordinate within a square or an inscribed circle?



b is 'in the circle' if $\sqrt{b_1^2 + b_2^2} \leq r$
 b is 'out the circle' if $\sqrt{b_1^2 + b_2^2} > r$

Take a random b n times

$$\frac{\text{Area of circle}}{\text{Area of square}} = \frac{\pi r^2}{(2r)^2} = \frac{\pi}{4}$$

$$\frac{\pi}{4} \approx \frac{\frac{\# \text{ Points in circle}}{\# \text{ Points out circle}}}{\frac{\# \text{ Points in circle}}{\# \text{ Points in circle}}} \rightarrow \pi \approx 4 \times \frac{\# \text{ Points in circle}}{\# \text{ Points out circle}}$$

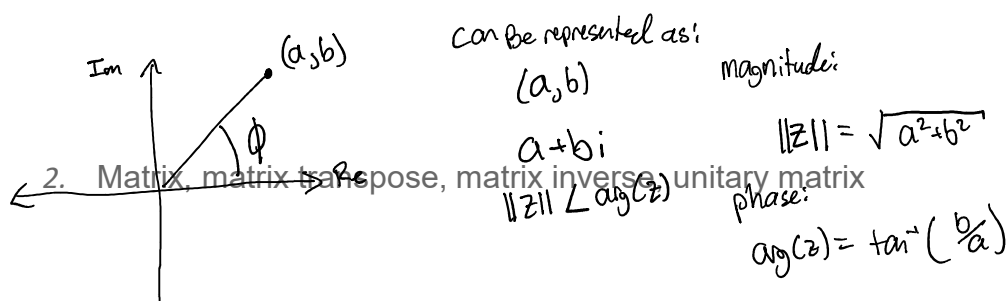
Steps: generate b_1, b_2

4. For each of the above algorithms, analyze them in terms of ^{square, subtract, sqrt, compare} representation of information, time (i.e., number of steps) to compute the value (i.e., execute the algorithm), the quality of the result, how to ^{add} improve quality of result, and the primitive operations required.

continued above

2. Review the following elementary concepts:

- Complex numbers, their representation, visualization, magnitude and phase





$$u \cdot v = 0$$

$$\arg(z) = \tan^{-1}\left(\frac{b}{a}\right)$$

$$A \in \mathbb{R}^{3 \times 3}$$

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$

$$A^{-1} \cdot A = I$$

$$A^{-1} \cdot \begin{cases} \text{must be square} \\ \det(A) \neq 0 \end{cases} \left. \begin{array}{l} \text{non-singular} \\ \text{i.e.} \\ \text{invertable} \end{array} \right\}$$

$$\hookrightarrow \det(B \in \mathbb{R}^{2 \times 2}) = \det\left(\begin{bmatrix} a & b \\ c & d \end{bmatrix}\right) = ad - bc$$

$$A^T = \begin{bmatrix} a_{11} & a_{21} & a_{31} \\ a_{12} & a_{22} & a_{32} \\ a_{13} & a_{23} & a_{33} \end{bmatrix}$$

$$\det(A) = a_{11} \begin{vmatrix} a_{22} & a_{23} \\ a_{32} & a_{33} \end{vmatrix} - a_{12} \begin{vmatrix} a_{21} & a_{23} \\ a_{31} & a_{33} \end{vmatrix} + a_{13} \begin{vmatrix} a_{21} & a_{22} \\ a_{31} & a_{32} \end{vmatrix}$$

Unitary matrix

- Square matrix
- Maintains vector length/norm, orthogonality

$$\text{satisfies } A^H = A^{-1} \Rightarrow A^H A = I$$

Conjugate Transpose Inverse

$$\text{eg. } U = \frac{1}{2} \begin{bmatrix} 1-i & 1+i \\ 1+i & 1-i \end{bmatrix}$$

3. Vector space vector norms

$$U = \frac{1}{2} \begin{bmatrix} 1-i & 1+i \\ 1+i & 1-i \end{bmatrix}$$

$$U^H = \frac{1}{2} \begin{bmatrix} 1+i & 1-i \\ 1-i & 1+i \end{bmatrix}$$

multiply & show $U^H U = I$

Vector Space properties

$$\text{norm: } V \in \mathbb{R}^{n \times 1}$$

$$\|V\| = \sqrt{V_1^2 + V_2^2 + \dots + V_n^2}$$

1. Closure under vector space

2. Closure under scalar multiplication

3. Associativity of addition

4. Boolean algebra

4. Commutativity of addition

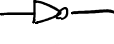



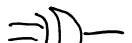


5. Identity element of addition

6. Existence of additive Inverse

7. Logic operations (AND, OR, NOT, NAND, NOR, XOR, etc.);

8. Compatibility of scalar multiplications

1 - On/active
 0 - off/inactive

\bar{x} not	Not 	$\begin{array}{c c} 1 & 0 \\ \hline 0 & 1 \end{array}$
$x+y$ OR	OR 	$\begin{array}{c c} 0 & 0 & 1 \\ \hline 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \end{array}$
truth tables: xy AND	And 	$\begin{array}{c c} 0 & 0 & 0 \\ \hline 0 & 1 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 1 \end{array}$
$\bar{x} \rightarrow y$ implication	Nand 	$\begin{array}{c c} 0 & 0 & 1 \\ \hline 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{array}$
$x \oplus y$ exclusive OR	XOR 	$\begin{array}{c c} 0 & 0 & 0 \\ \hline 0 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{array}$
$x \equiv y$ equivalence	NOR 	$\begin{array}{c c} 0 & 0 & 1 \\ \hline 0 & 1 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{array}$
	XNOR 	$\begin{array}{c c} 0 & 0 & 1 \\ \hline 0 & 1 & 0 \\ 1 & 0 & 0 \\ 1 & 1 & 1 \end{array}$

DeMorgan's rule;

x	y	\bar{x}	$x+y$	xy	$x \rightarrow y$	$x \oplus y$	$x \equiv y$
0	0	1	0	0	1	0	1
0	1	1	1	0	1	1	0
1	0	0	1	0	0	1	0
1	1	0	1	1	1	0	1

parity and XORs

$$(A \cup B)' = A' \cap B'$$



Parity - Check for errors during data transmission/storage
 - Odd parity - make total # of ones odd (inc. parity)
 - Even parity - make total # of ones even

XORs - output the inverse of input

- functional completeness

- functional completeness
 sum-of-products any logical or arithmetic operation can be expressed as
 a combination of only And & not operators

- functional completeness
 product-of-sums
 ↳ same as above

- equivalent to logical AND function

3. Let x be a vector in \mathbb{R}^N , and let U be a unitary matrix that can operate on x . What is the 2-norm of x ? What is the 2-norm of Ux ?

$$F(A,B) = AB + A'B = AB + AB$$

 output true for a specific input

- equivalent to logical OR function

$$F(A,B) = (A+B)(A+B')(A'+B)(A'+B')$$

output false for a specific input

4. Code search() that takes an input list and a keying function f as input, and behaves as described in the lecture.

$$x \in \mathbb{R}^{n \times 1}$$

$$\|x\|^2 = \sqrt{x_1^2 + x_2^2 + \dots + x_N^2}$$

$$\|Ux\|^2 = \|x\|^2$$

5. Code sort() that takes an input list and a two-input compare function f as input, and behaves as described in any computer algorithms resource.

```

def search(input_list, f):
    for x in input_list:
        if f(x):
            return x

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6. Let u be a signed binary word. Design the function $\text{sqrt}(u)$, that returns the square root of u to two bits of precision (and full dynamic range). Show the truth table and the logic equations.

Truth Table:

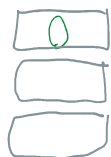
U	Binary	Sqrt(u)	Sqrt(u) Binary
0	0000	0	00
1	0001	1	01
2	0010	1	01
3	0011	1	01
4	0100	2	10
5	0101	2	10
6	0110	2	10
7	0111	3	11
8	01000	3	11
9	01001	3	11
10	01010	3	11
11	01011	3	11
12	01100	3	11
13	01101	4	11
14	01110	4	11
15	01111	4	11

Logic eqn

$Y_1 = 1$ if U_1 or U_2 or U_3 or U_4 are true
 \Rightarrow only consider U_2, U_3

$$Y_1 = U_2 + U_3$$

Logic Eqns:



0
0
0
0
0
0
0
0
0
0
0

Y_0 all except
 • all zero \Rightarrow NAND
 • $\bar{U}_0 \bar{U}_1 U_2 \bar{U}_3$
 • $U_0 \bar{U}_1 U_2 \bar{U}_3$
 • $U_0 U_1 U_2 U_3$

$$Y_0 = (\bar{U}_0 \bar{U}_1 U_2 \bar{U}_3) \cdot (U_0 \bar{U}_1 U_2 \bar{U}_3) + (\bar{U}_0 U_1 U_2 \bar{U}_3) \cdot (U_0 U_1 U_2 U_3)$$

$$u_4 u_3 u_2 u_1 u_0$$

$$\chi_1 \chi_0$$

$$\gamma_0 =$$

