

UNIVERSITY OF HOUSTON

PROGRAMMING ASSIGNMENT 2

COSC 3320
Algorithms and Data Structures

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Solution

DRAFT

1 Pseudocode and Explanation

Notice that this problem is effectively just the [Select](#) algorithm on the *distances from the origin*. Letting $\text{dist}_0^2(x, y) = x^2 + y^2$, our algorithm is simply

Algorithm 1 [Select](#)

```

1: def SELECT(points, k):
    Input ▷ An array points of  $n$  ordered pairs and the desired order statistic,  $k$ 
    Output ▷ The value of the  $k^{\text{th}}$  order statistic
2: if |points| ≤ 1:
3:     return points[0]
4: else:
5:     Partition points into  $\lfloor n/5 \rfloor$  groups of 5 elements each and a leftover group of up to 4 elements.
6:     Find the median of each group using MergeSort with comparison function  $\text{dist}_0^2$ 
7:     medians ← the set of these  $\lfloor n/5 \rfloor$  medians
8:     pivot ← SELECT(medians,  $\lfloor n/10 \rfloor$ )
9:     left ← { $p \in \text{points} \mid \text{dist}_0^2(p) < \text{dist}_0^2(\text{pivot})$ }
10:    right ← { $p \in \text{points} \mid \text{dist}_0^2(p) > \text{dist}_0^2(\text{pivot})$ }
11:    if |left| =  $k - 1$ :
12:        return p
13:    else if |left| >  $k - 1$ :
14:        return SELECT(left, k)
15:    else:
16:        return SELECT(right,  $k - |\text{left}| - 1$ )

```

Then we can simply return the k closest points to the origin by comparing with [SELECT](#)(points, k). However, we must be careful about duplicates: say there are ℓ points less than the k^{th} closest point and t points equidistant. Then, by definition, we must have $\ell + t \geq k$. If $t = 1$ then $\ell = k - 1$ (of which all distances being distinct is a special case). Either way, we must choose all ℓ points and then any $k - \ell$ points equidistant to our k^{th} closest point.

$$\underbrace{\{p_1, p_2, \dots, p_\ell\}}_{\ell \text{ points}} \underbrace{\{p_{\ell+1}, p_{\ell+2}, \dots, p_k, \dots, p_{\ell+t}\}}_{k-\ell \text{ points}}$$

$\ell + (k - \ell) = k \text{ points}$

Algorithm 2 [k-Closests](#)

```

1: def k-CLOSESTS(points, k):
    Input ▷ An array points of  $n$  ordered pairs and an integer  $k$ 
    Output ▷ The  $k$  closest points to the origin
2: if |points| ≤ k:
3:     return points
4: else:
5:     cutoff ← SELECT(points, k)
6:     left ← { $p \in \text{points} \mid \text{dist}_0^2(p) < \text{dist}_0^2(\text{cutoff})$ }
7:     equal ← { $p \in \text{points} \mid \text{dist}_0^2(p) = \text{dist}_0^2(\text{cutoff})$ }
8:      $\ell \leftarrow |\text{left}|$ 
9:     append  $k - \ell$  points of equal to left
10:    return left

```

2 Analysis

Correctness follows from correctness of the [Select](#) from the textbook. The runtime of [Select](#) follows from the same analysis. Specifically, our runtime is $\Theta(n)$.