

COSC6364

Lecture 5: QR factorization

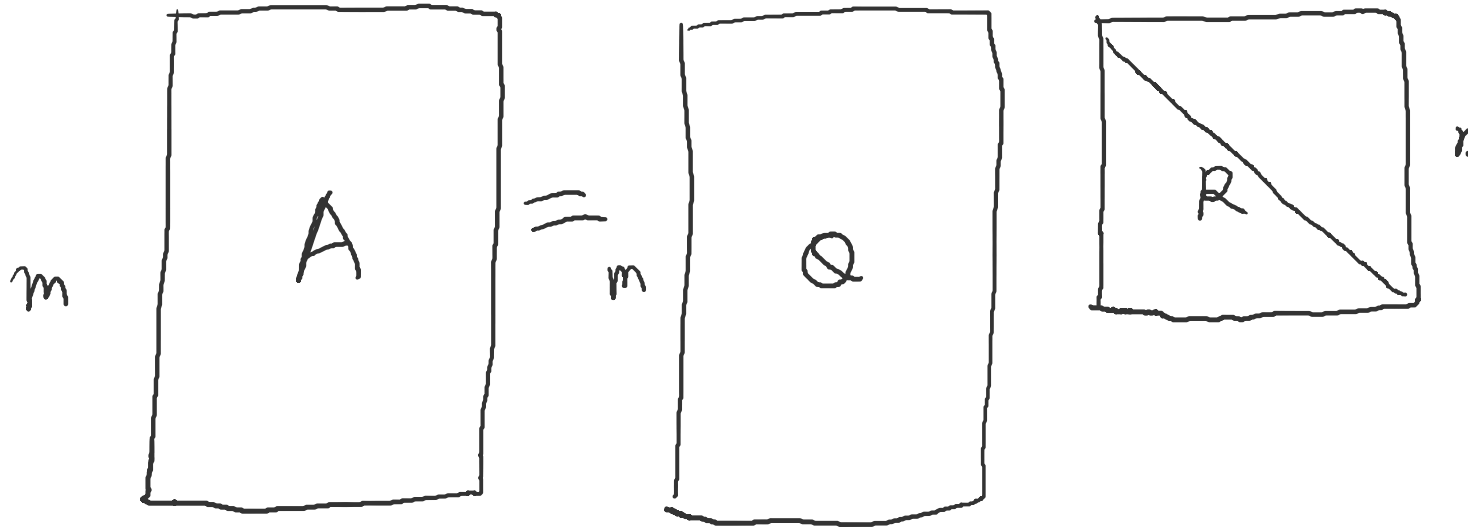
NLA Chapter 7, 8, 10

Probably the most important matrix factorization in NLA

QR factorization: (the reduce QR form)

$$A = \hat{Q}\hat{R}$$

Where \hat{Q} columns are orthonormal and \hat{R} is upper triangular



Reduced QR Factorization

- We are interested in the **column spaces** of a matrix A .
- In fact, we are interested in the successive spaces spanned by the columns $\langle a_1, a_2, \dots \rangle$ of A
$$\langle a_1 \rangle \subset \langle a_1, a_2 \rangle \subset \langle a_1, a_2, a_3 \rangle \subset \dots$$
- Notation: $\langle x, y, \dots \rangle$ denotes the subspace spanned by vectors inside the brackets.
- The idea of QR: to construct a sequence of orthonormal vectors q_1, q_2, \dots that span these successive spaces:
- $\langle a_1 \rangle = \langle q_1 \rangle$
 $\langle a_1, a_2 \rangle = \langle q_1, q_2 \rangle$
 $\langle a_1, a_2, a_3 \rangle = \langle q_1, q_2, q_3 \rangle$
$$\dots$$
- How to express these in terms of matrix?

Reduced QR Factorization cont'd

- To be precise, assume $A \in \mathbb{R}^{m \times n}$ ($m \geq n$) has full rank n . We want the sequence q_1, q_2, \dots, q_n to have the property:
 $\langle q_1, q_2, \dots, q_j \rangle = \langle a_1, a_2, \dots, a_j \rangle, \quad j = 1, \dots, n$
- This amounts to

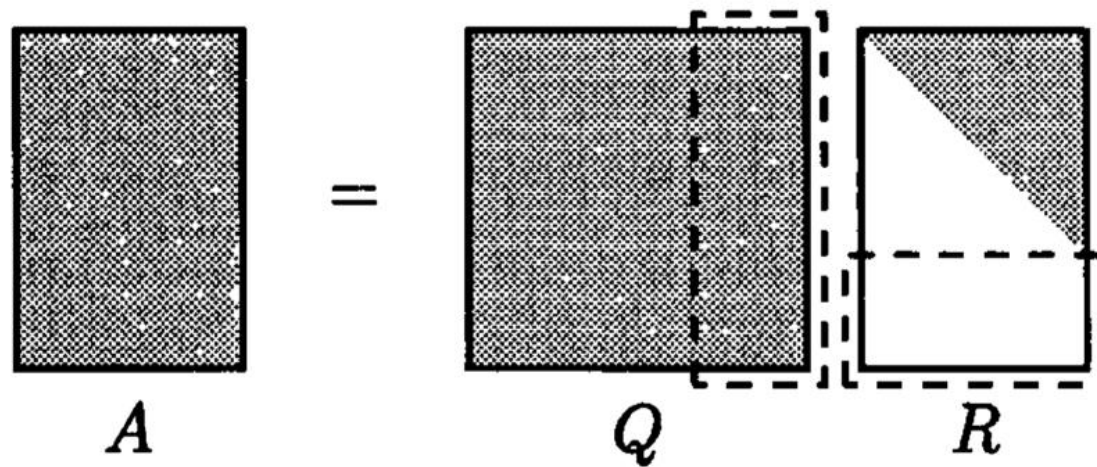
$$\left[\begin{array}{c|c|c|c} a_1 & a_2 & \cdots & a_n \end{array} \right] = \left[\begin{array}{c|c|c|c} q_1 & q_2 & \cdots & q_n \end{array} \right] \left[\begin{array}{cccc} r_{11} & r_{12} & \cdots & r_{1n} \\ & r_{22} & & \vdots \\ & & \ddots & \\ & & & r_{nn} \end{array} \right],$$

- Or, $A = \hat{Q}\hat{R}$, where \hat{Q} is a $m \times n$ orthonormal columns, and \hat{R} is $n \times n$ upper triangular matrix. This is the reduced QR factorization.

Full QR factorization

- The full QR factorization extends $m - n$ columns to \hat{Q} to make it square $m \times m$. (Analogous to Full SVD vs. Reduced SVD)

Full QR Factorization ($m \geq n$)



- The silent columns of Q forms $\text{range}(A)^\perp$

Gram-Schmidt Orthogonalization

- Look at the QR factorization in the form:

$$a_1 = r_{11}q_1,$$

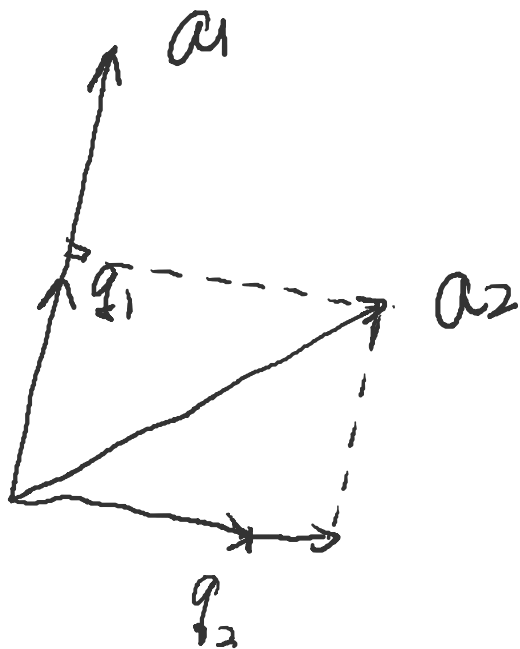
$$a_2 = r_{12}q_1 + r_{22}q_2,$$

$$\vdots$$

$$a_n = r_{1n}q_1 + r_{2n}q_2 + \cdots + r_{nn}q_n$$

- Gram-Schmidt orthogonalization computes the r_{ij} from these equations, from top to bottom.
- Can you figure out how to compute $r_{11}, r_{12}, r_{22}, \dots$? (Hint: note that the $\{q_i\}$ are orthonormal!)

GS-Ortho



$$q_1 = \frac{a_1}{\|a_1\|}, \quad r_{11} = \|a_1\|$$

$$v_2 = a_2 - (q_1^T a_2) q_1, \quad q_2 = \frac{v_2}{\|v_2\|}$$

$$r_{12} = q_1^T a_2, \quad r_{22} = \|v_2\|$$

$$v_3 = a_3 - (q_1^T a_3) q_1 - (q_2^T a_3) q_2$$

$$r_{13} = q_1^T a_3, \quad r_{23} = q_2^T a_3$$

$$r_{33} = \|v_3\|$$

$$v_j = a_j - (q_1^T a_j) q_1 - (q_2^T a_j) q_2 - \dots - (q_{j-1}^T a_j) q_{j-1}$$

$$r_{1j} = q_1^T a_j, \quad r_{2j} = q_2^T a_j, \quad \dots, \quad r_{j-1,j} = q_{j-1}^T a_j$$

$$r_{jj} = \|v_j\|$$

Classical Gram-Schmidt (CGS): Careful, may not be stable!

Conceptually simple to understand but numerical unstable

QR factorization by CGS:

Algorithm 7.1. Classical Gram–Schmidt (unstable)

for $j = 1$ **to** n

$$v_j = a_j$$

for $i = 1$ **to** $j - 1$

$$r_{ij} = q_i^* a_j$$

$$v_j = v_j - r_{ij} q_i$$

$$r_{jj} = \|v_j\|_2$$

$$q_j = v_j / r_{jj}$$

Existence and Uniqueness of (reduced) QR

- Does QR factorization always exist?
 - Yes! You just saw how to create one...
- But is it unique?
 - Pretty much yes, but under the following conditions:
 - $A \in \mathbb{R}^{m \times n}$ ($m \geq n$) is full rank
 - If you insist $r_{jj} > 0$

Solution of $Ax = b$ by QR Factorization

How to solve x from $Ax = b$?

If $A \in \mathbb{R}^{m \times m}$ is non-singular, then we can rewrite the equation as

$$\begin{aligned} QRx &= b \\ Rx &= Q^T b \end{aligned}$$

And the last equation is easy to solve because R is upper triangular. So here's the algorithm to solve $Ax = b$ by QR:

1. Compute QR factorization of $A = QR$
2. Compute $y = Q^T b$
3. Solve $Rx = y$

This is actually an excellent way to solve a linear equation but not the standard one; the standard one is LU with partial pivoting.

Gram-Schmidt Orthogonalization

- GS is a “triangular orthogonalization”, which makes the columns of a matrix orthonormal via a sequence of right triangular matrix multiplications.
- Soon we’ll learn another approach of doing QR factorization called “orthogonal triangularization” (Householder QR).

GS as a sequence of orthogonal projections

Let $A \in \mathbb{R}^{m \times n}$, $m \geq n$ be a full rank matrix with columns $\{a_i\}$.

We can express the GS as a sequence of orthogonal projections:

$$q_1 = \frac{P_1 a_1}{\|P_1 a_1\|}, q_2 = \frac{P_2 a_2}{\|P_2 a_2\|}, \dots$$

Each P_j denotes an orthogonal projector that projects a_j orthogonally onto the space orthogonal to $\langle q_1, q_2, \dots, q_{j-1} \rangle$.

What is the projector in terms of q_j ?

$$P_j = I - Q_{j-1} Q_{j-1}^T$$

Why? (Hint: what does the projector $q q^T$ do? What does $I - q q^T$ do?)

Modified Gram-Schmidt (MGS)

- CGS is easy to understand but unfortunately not very stable.
- A mathematically equivalent version (MGS) however is stable.
- The CGS computes

$$v_j = P_j a_j$$

- Key idea (why?):

$$v_j = P_{\perp q_{j-1}} \cdots P_{\perp q_2} P_{\perp q_1} a_j$$

- Initiate $v_j = a_j, j = 1, \dots, n$.
- In the i iteration, apply $P_{\perp q_i}$ to all $v_j, j > i$.
- And that's it! Essentially, different order of updating v_j is the difference between CGS and MGS.

MGS:

Mathematically equivalent to CGS but numerically superior

Algorithm 8.1. Modified Gram–Schmidt

for $i = 1$ **to** n

$$v_i = a_i$$

for $i = 1$ **to** n

$$r_{ii} = \|v_i\|$$

$$q_i = v_i / r_{ii}$$

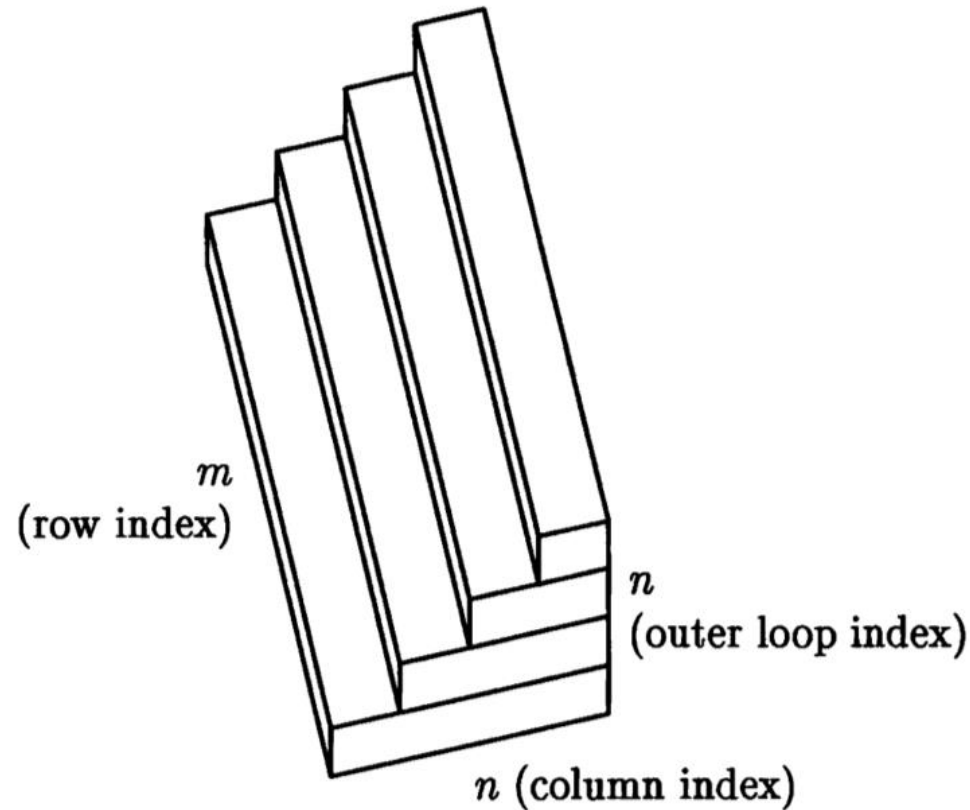
for $j = i + 1$ **to** n

$$r_{ij} = q_i^* v_j$$

$$v_j = v_j - r_{ij} q_i$$

How much does Gram-Schmidt orthogonalization cost?

- Unit of cost: floating point operations (FLOP)
- Both CGS and MGS cost about the same: $2mn^2$ FLOPs.



Gram-Schmidt as Triangular Orthogonalization

The way to look at MGS in matrix form:

The first outer iteration:

$$\left[\begin{array}{c|c|c|c} v_1 & v_2 & \cdots & v_n \end{array} \right] \left[\begin{array}{cccc} \frac{1}{r_{11}} & \frac{-r_{12}}{r_{11}} & \frac{-r_{13}}{r_{11}} & \cdots \\ & 1 & & \\ & & 1 & \\ & & & \ddots \end{array} \right] = \left[\begin{array}{c|c|c|c} q_1 & v_2^{(2)} & \cdots & v_n^{(2)} \end{array} \right].$$

In general, each iteration is a right upper triangular matrix multiplication: $AR_1R_2\cdots R_n = \hat{Q}$.

$$R_2 = \left[\begin{array}{cccc} 1 & & & \\ & \frac{1}{r_{22}} & \frac{-r_{23}}{r_{22}} & \cdots \\ & & 1 & \\ & & & \ddots \end{array} \right], \quad R_3 = \left[\begin{array}{cccc} 1 & & & \\ & 1 & & \\ & & \frac{1}{r_{33}} & \cdots \\ & & & \ddots \end{array} \right], \dots$$

Householder Triangularization

Householder triangularization (HT) is another method to compute QR factorization. Compared to GS orthogonalization,

- HT is more numerically stable.
- Loss of ability as basis for iterative methods
- HT is cheaper than GS (less FLOPs)
- It's orthogonal triangularization, rather than triangular orthogonalization:

Householder vs. Gram-Schmidt

The GS iteration applies a succession of elementary triangular matrices R_k on the right of A :

$$A \underbrace{R_1 R_2 \dots R_n}_{\hat{R}^{-1}} = \hat{Q}$$

so that the resulting matrix has orthonormal columns.

In contrast, the Householder method applies a succession of elementary unitary matrices Q_k on the left of A :

$$\underbrace{Q_n \dots Q_2 Q_1}_{Q^*} A = R$$

so that the resulting matrix is upper triangular.

Triangularization by Introducing Zeros

At the heart of Householder Triangularization is the Householder reflector (orthogonal matrices) that introduces zeros below the diagonal in the k th column while preserving all the zeros previously introduced.

$$\begin{bmatrix} \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \\ \times & \times & \times \end{bmatrix} \xrightarrow{Q_1} \begin{bmatrix} \times & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \\ 0 & \times & \times \end{bmatrix} \xrightarrow{Q_2} \begin{bmatrix} \times & \times & \times \\ & \times & \times \\ & 0 & \times \\ & 0 & \times \\ & 0 & \times \end{bmatrix} \xrightarrow{Q_3} \begin{bmatrix} \times & \times & \times \\ & \times & \times \\ & & \times \\ & & 0 \\ & & 0 \end{bmatrix} \quad (10.1)$$

$A \qquad Q_1 A \qquad Q_2 Q_1 A \qquad Q_3 Q_2 Q_1 A$

Householder Reflector (1/3)

How do we construct such orthogonal matrices Q_1, Q_2, \dots ?

Here's a standard approach. Each Q_k is taken to be:

$$Q_k = \begin{bmatrix} I & 0 \\ 0 & F \end{bmatrix}$$

where $I \in \mathbb{R}^{(k-1) \times (k-1)}$, and $F \in \mathbb{R}^{(m-k+1) \times (m-k+1)}$ is orthogonal.

Q_k introduces zeros into the k th column. We choose F to be a particular matrix called Householder Reflector

Householder Reflector (2/3)

Suppose, at the beginning of step k , the entries $k, k + 1, \dots, m$ of the k th column are given by vector $x \in \mathbb{R}^{m-k+1}$. The Householder reflector F should effect the following map:

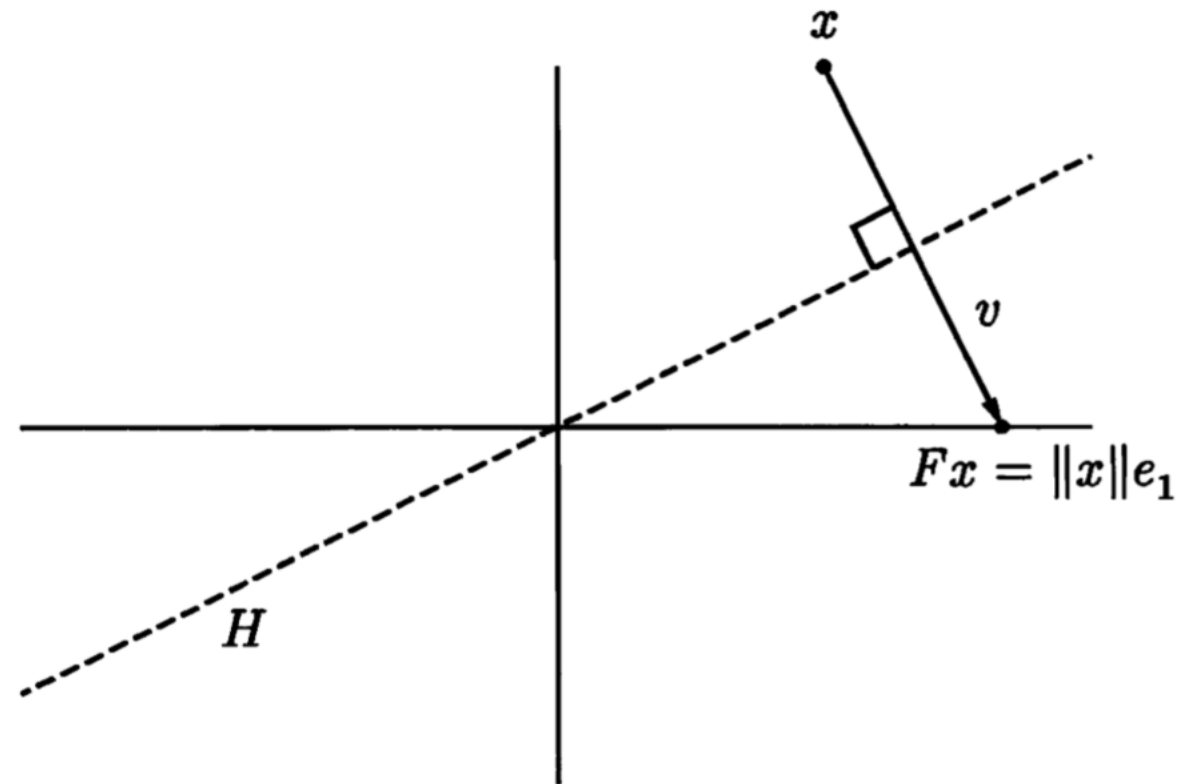
$$x = \begin{bmatrix} \times \\ \times \\ \times \\ \vdots \\ \times \end{bmatrix} \xrightarrow{F} Fx = \begin{bmatrix} \|x\| \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \|x\|e_1. \quad (10.3)$$

(Why should the first component of Fx be $\|x\|$?)

Householder Reflector (3/3)

The reflector F will reflect across the hyperplane H orthogonal to $v = \|x\|e_1 - x$.

A hyperplane is a generalization of plane in 2D (e.g. 3-dimensional subspace in 4D space). A hyperplane can be characterized as all the vectors orthogonal to a fixed vector (normal vector).



Projector and Reflector

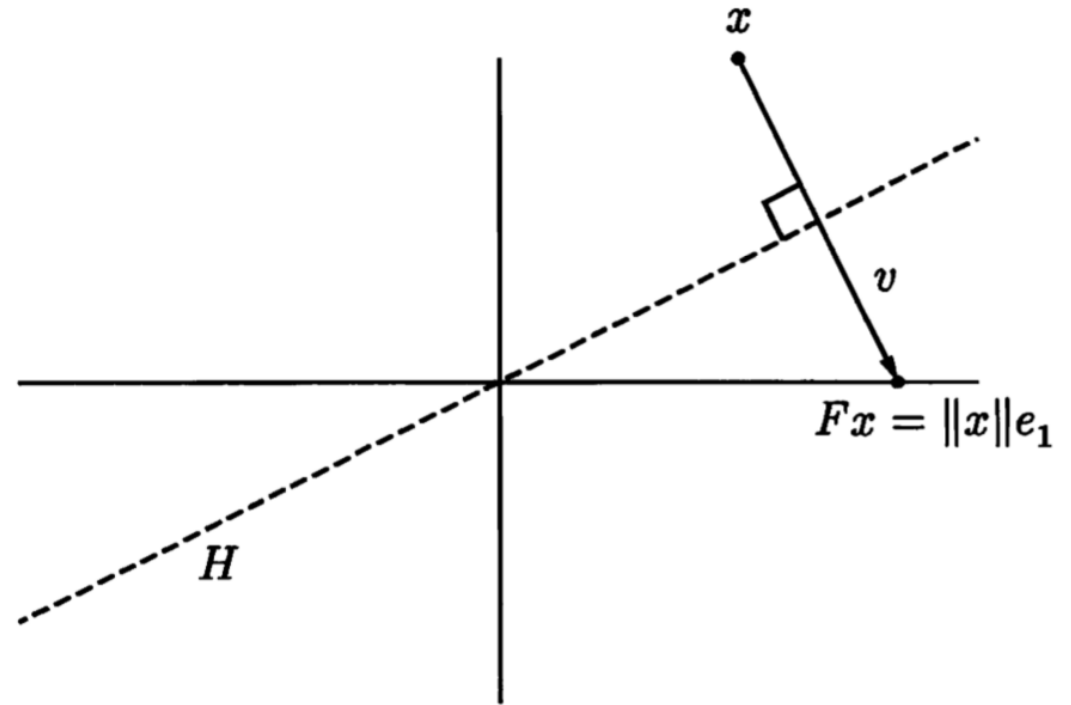
Remember the projectors? What is the projector that maps a vector y to $\langle v \rangle$? What projector maps y to H ?

$$P = I - \frac{vv^T}{v^T v}$$

To reflect y across H , we must not stop at projection onto H ; we must go exactly twice as far in the same direction:

$$F = I - 2 \frac{vv^T}{v^T v}$$

(Which one of P, F is orthogonal?)

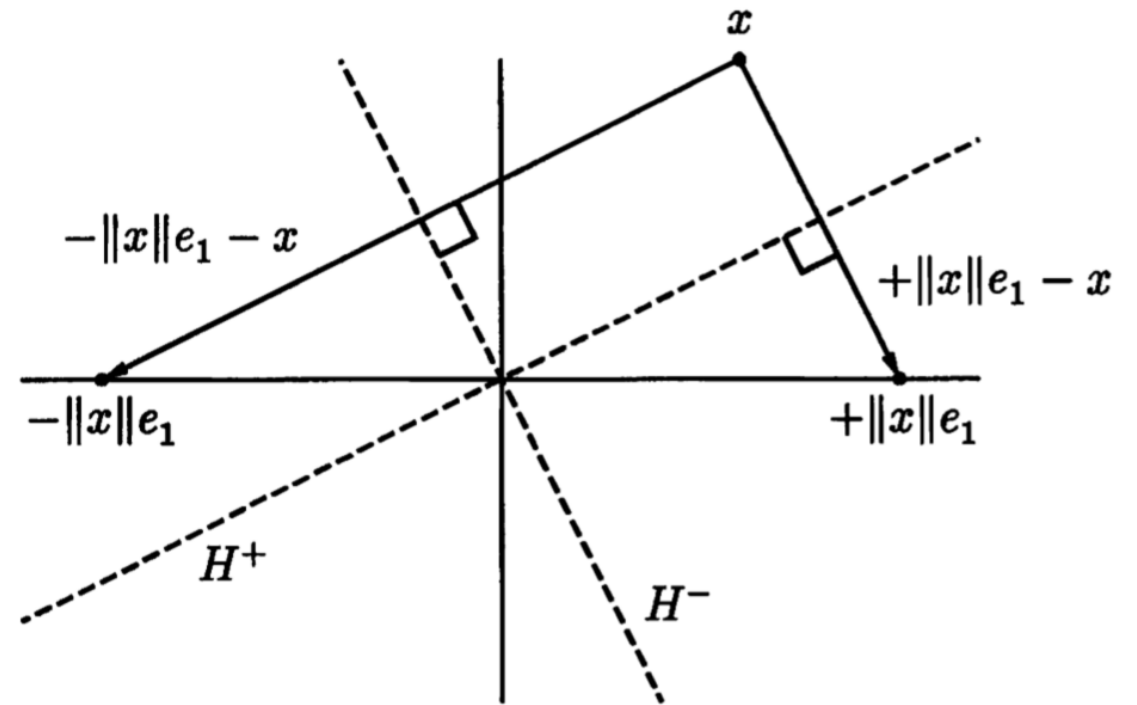


The Better of Two Reflectors

In fact we can have two reflectors... is there a difference?

Two possible images on the axis, $\pm \|x\|e_1$. We'll always want the furthest one from x . Why?

It turns out that we will calculate $v = \pm \|x\|e_1 - x$. If they are close the result is very sensitive to rounding errors.



Householder QR

Given a matrix $A \in \mathbb{R}^{m \times n}$, $m \geq n$, the following algorithm computes a QR factorization of A . The upper triangular part of R will overwrite A . The Q factor is stored as Householder vectors v_1, \dots, v_n .

We can further store the Householder vectors in lower triangular part of A so that we completely overwrite A and need no extra memory.

Algorithm 10.1. Householder QR Factorization

for $k = 1$ **to** n

$$x = A_{k:m,k}$$

$$v_k = \text{sign}(x_1)\|x\|_2 e_1 + x$$

$$v_k = v_k / \|v_k\|_2$$

$$A_{k:m,k:n} = A_{k:m,k:n} - 2v_k(v_k^* A_{k:m,k:n})$$

Applying or Forming Q

OK, but how do I get my good Q from the Householder vectors, you might ask?

It turns out that, you probably can get away without forming Q , if all you need to multiply Q or Q^T by something (applying).

But if you insist, you can form the explicit Q by applying Q to the identity matrix I ...

Let's see.

Applying Q and Q^T

We know that:

$$Q^T = Q_n \cdots Q_2 Q_1$$

and its conjugate:

$$Q = Q_1 Q_2 \cdots Q_n$$

(why there's no transpose?)

Algorithm 10.2. Implicit Calculation of a Product Q^*b

for $k = 1$ **to** n

$$b_{k:m} = b_{k:m} - 2v_k(v_k^* b_{k:m})$$

Algorithm 10.3. Implicit Calculation of a Product Qx

for $k = n$ **downto** 1

$$x_{k:m} = x_{k:m} - 2v_k(v_k^* x_{k:m})$$

Forming Q

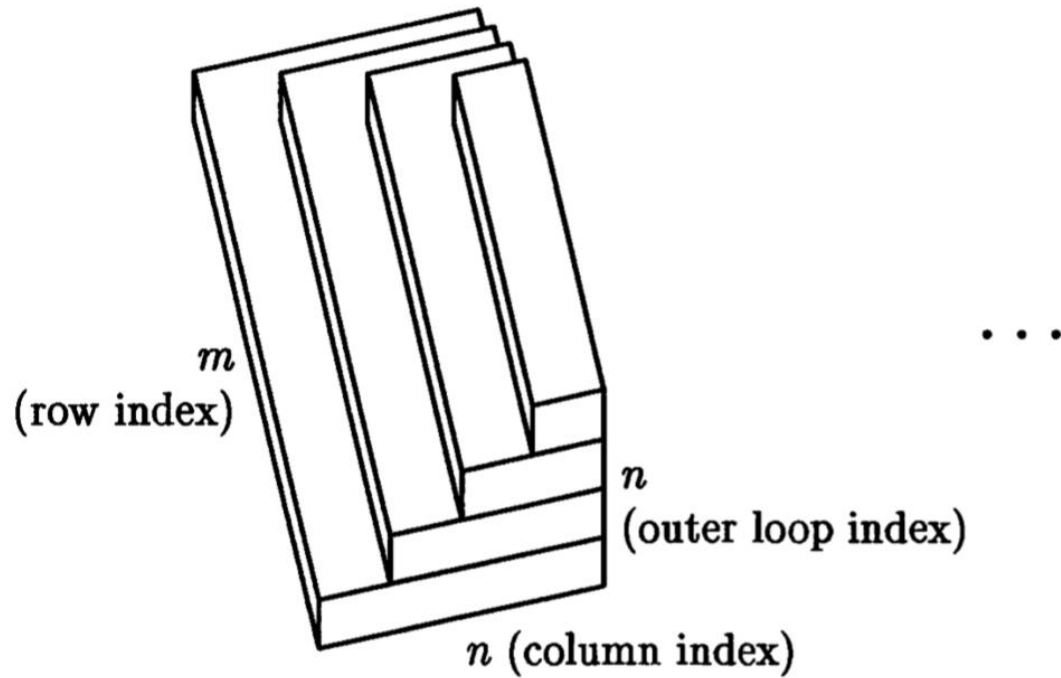
A bunch of ways:

1. QI
2. $Q^T I$ then transpose
3. IQ

The first method is the best. Why?

Operation Count

Householder QR:



$$\sim 2mn^2 - \frac{2}{3}n^3 \text{ flops}$$

Classical vs. Modified Gram-Schmidt

Numerical stability between CGS and MGS.

First we create a matrix with a widely varying singular values spaced by factor of 2 between 2^{-1} and 2^{-80} using Matlab:

```
[U,X] = qr(randn(80));
```

Set U to a random orthogonal matrix.

```
[V,X] = qr(randn(80));
```

Set V to a random orthogonal matrix.

```
S=diag(2.^(-1:-1:-80));
```

Set S to a diagonal matrix with exponentially graded entries.

```
A = U*S*V;
```

Set A to a matrix with these entries as singular values.

CGS and MGS to do QR

We plot the diagonal elements r_{ii} . Since $r_{ii} = \|P_i a_i\|$, we would estimate that r_{ii} closely matches the i -th singular value $\sigma_i = 2^{-i}$ within a constant factor.

Here's the result of the r_{ii} produced by CGS and MGS QR:

r_{ii} of CGS and MGS

