

Lab_Offsets

CMPUT 229

University of Alberta

Outline

1 Lab_Offsets: New Concepts and Assignment

- Branch Offset
- ASCII and Hexadecimal
- ASCII, Hex and CString
- Subroutines
- Masking
- Shifting
- Lab 2: Hexadecimal and Branches
- Assignment Tips
- Questions?

Lab_Offsets: Decoding Branch Instructions

- In this lab you will learn to decode a MIPS Branch instruction.

Branches

Branches in MIPS are control structures (i.e., if-then structures in higher level programming)

Instruction	What does it do?
bgez \$s, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s \geq 0$.
bgezal \$s, offset	Sets $\$ra = PC$, jumps to $PC + \text{offset} \times 4$ if $\$s \geq 0$.
bltz \$s, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s < 0$.
bltzal \$s, offset	Sets $\$ra = PC$, jumps to $PC + \text{offset} \times 4$ if $\$s < 0$.
beq \$s, \$t, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s = \t .
bne \$s, \$t, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s \neq \t .
blez \$s, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s \leq 0$.
bgtz \$s, offset	Jumps to $PC + \text{offset} \times 4$ if $\$s > 0$.

Branch Offset

- Then branch **offset** indicates where to jump.
- If a branch is taken, the PC is set to $PC + offset \times 4$.
- Offset: how many words up or down we need to move.
 - $Offset \in [-2^{15}, 2^{15} - 1]$
- Encoding:

31 30 29 28 27 26	25 24 23 22 21	20 19 18 17 16	15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0
OpCode	Register s	Register t	Branch Offset (16-bit Immediate)

Branch Offset: An Example

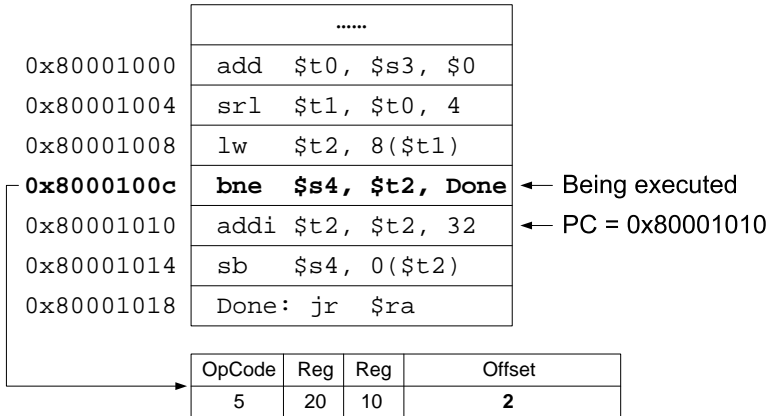
- Suppose there is a program containing a branch-not-equals, bne.
- Suppose $\$s4 = \$t2$.

0x80001000	add \$t0, \$s3, \$0
0x80001004	srl \$t1, \$t0, 4
0x80001008	lw \$t2, 8(\$t1)
0x8000100c	bne \$s4, \$t2, Done
0x80001010	addi \$t2, \$t2, 32
0x80001014	sb \$s4, 0(\$t2)
0x80001018	Done: jr \$ra

Branch Offset: An Example

	
0x80001000	add \$t0, \$s3, \$0	
0x80001004	srl \$t1, \$t0, 4	
0x80001008	lw \$t2, 8(\$t1)	
0x8000100c	bne \$s4, \$t2, Done	← Being executed
0x80001010	addi \$t2, \$t2, 32	← PC = 0x80001010
0x80001014	sb \$s4, 0(\$t2)	
0x80001018	Done: jr \$ra	

Branch Offset: An Example



Branch Offset: An Example

	
0x80001000	add \$t0, \$s3, \$0	
0x80001004	srl \$t1, \$t0, 4	
0x80001008	lw \$t2, 8(\$t1)	
0x8000100c	bne \$s4, \$t2, Done	← Being executed
0x80001010	addi \$t2, \$t2, 32	
0x80001014	sb \$s4, 0(\$t2)	
0x80001018	Done: jr \$ra	← PC (to be executed)

$$PC = 0x80001010 + 2 * 4 = 0x80001018$$

Hexadecimal: Why?

- Hexadecimal is used for easy conversion to and from binary.
- It is easier to read and represent HEX than binary.
- Also: conversion between hexadecimal and ASCII is easy.

Simple Concepts of Hex and ASCII

- All data in a computer is represented in binary (i.e., only 0s and 1s).
- Hexadecimal: one way to read/print groups of bits rather than individual bits. Each hex value represents 4 binary digits.
- American Standard Code for Information Interchange (ASCII): a representation standard for characters and symbols.
- `.ascii` predirectives in the data segment are null-terminated strings (similar to C strings). A “string” in MIPS must have a null character at the end.

Translation between ASCII and Hex

■ Hex to ASCII

Input: 0x4

Output: '4'

INPUT	ASCII	Binary
-------	-------	--------

$(00000100)_2$	0x34	$(00110100)_2$
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0000 0100 OR 0011 0000 = 0011 0100

Translation between ASCII and Hex

■ ASCII to Hex

Input: '4'

Output: 0x4

ASCII	Binary	OUTPUT
0x34	$(00110100)_2$	$(00000100)_2$

0011 0100 AND 0000 1111 = 0000 0100

Subroutines

- *Subroutines* in MIPS are calls to another portion of the code.
- **Why?** ← Just like “procedures” or “functions” in higher level languages.
- `jal` is an instruction to make the jump to a symbol (marker for a piece of code).
- `jr` returns control to the caller.
- Conventions: `$a` registers for arguments, `$v` for return values.
- Conventions will be examined in more in detail in upcoming labs.

Masking Operations

- The smallest segment of memory that can be loaded/stored in main memory is *one byte*.
- There is no instruction in MIPS to set/get individual bits inside a word.
- Masks allow us to read and set individual bits and perform some arithmetic functions.

Masking Operations - AND

AND

- 1 When using ones in the mask, the corresponding bits are preserved or queried.
- 2 When using zeros in the mask, the corresponding bits are turned off.

Example 1: Querying a bit.

Mask:	0000 1000
Number:	1111 0101
Result:	0000 0000

Example 2: Extracting information.

Mask:	0000 1111
Number:	1010 1000
Result:	0000 1000

Masking Operations - OR

OR

- 1 When using ones in the mask, the corresponding bits are set to one.
- 2 When using zeros in the mask, the corresponding bits are preserved.

Example 1: Set to one.

Mask:	0000 1111
Number:	1101 0101
Result:	1101 1111

Example 2: Preserving bits.

Mask:	0000 1111
Number:	1010 0000
Result:	1010 1111

Masking Operations - XOR

OR

- 1 When using ones in the mask, the corresponding bits are inverted (toggled).
- 2 When using zeros in the mask, the corresponding bits are preserved.

Mask:	0000 1111
Number:	1101 0101
Result:	1101 1010

Shifting and Rotating Bits

- Shifting is basically moving bits from one location to another.
- Shifting can also be viewed as an arithmetic operation
 - Each time you shift to the **left** a bit you multiply **times 2**.
 - Each time you shift to the **right** a bit you divide **by 2**.
- *Rotating* is the same as shifting, but bits get carried back to the opposite side.

Shift:	4 ₁₀ left logical
Number:	1101 0101
Result:	0101 0000

Shift:	4 ₁₀ right logical
Number:	1101 0101
Result:	0000 1101

Shift:	4 ₁₀ right arithmetic
Number:	1101 0101
Result:	1111 1101

Shift:	4 ₁₀ rotate left
Number:	1101 0101
Result:	0101 1101

Shifting as an Arithmetic Operation

- In any positional base- b numeral system, adding zeros to the right implies multiplying times the base (power of the radix).
 - $1040_{10} * 10_{10} = 10400_{10}$ (sll 1).
 - $0011_2 * 10_2 = 0110_2$, (sll 1) ($3_{10} * 2_{10} = 6_{10}$).
 - $001F_{16} * 10_{16} = 01F0_{16}$, (sll 1) ($31_{10} * 16_{10} = 496_{10}$).
- Analogously, shifting to the right implies a division by the base.
 - $0110_2 / 10_2 = 0011_2$, (srl 1) ($6_{10} / 2_{10} = 3_{10}$).
 - Caveat: for signed numbers in two's complement you need an **arithmetic shift** instead of a **logical shift**.
 - $1011_2 / 10_2 = 1101_{16}$, (sra 1) ($-5_{10} / 2_{10} \neq -3_{10}$).

Lab 2 Assignment

You need to create a **subroutine** called `disassembleBranch` that:

- Takes an argument `$a0` → the address of a branch instruction.
- Translates from a binary representation to text.
- Output:
 - If instruction is not a branch → no output is generated.
 - If instruction is a MIPS branch, → the instruction must be printed in the screen.
- Please refer to the lab specification for all details on the output.

Assignment Tips

- Read specifications very carefully. Pay special attention to what you have to include - we don't want a `main` method.
- Test your assignments on the lab machines before you submit. That's where we'll be marking them.
- Use the test cases provided to debug your code.
- Look at the marksheet to get an idea of how the grading will be done.
- Style marks are easy marks. Format your code like the `example.s` file we provided, and write good comments.
- Be sure to submit code that runs and loads; otherwise, you will lose marks.

Questions?