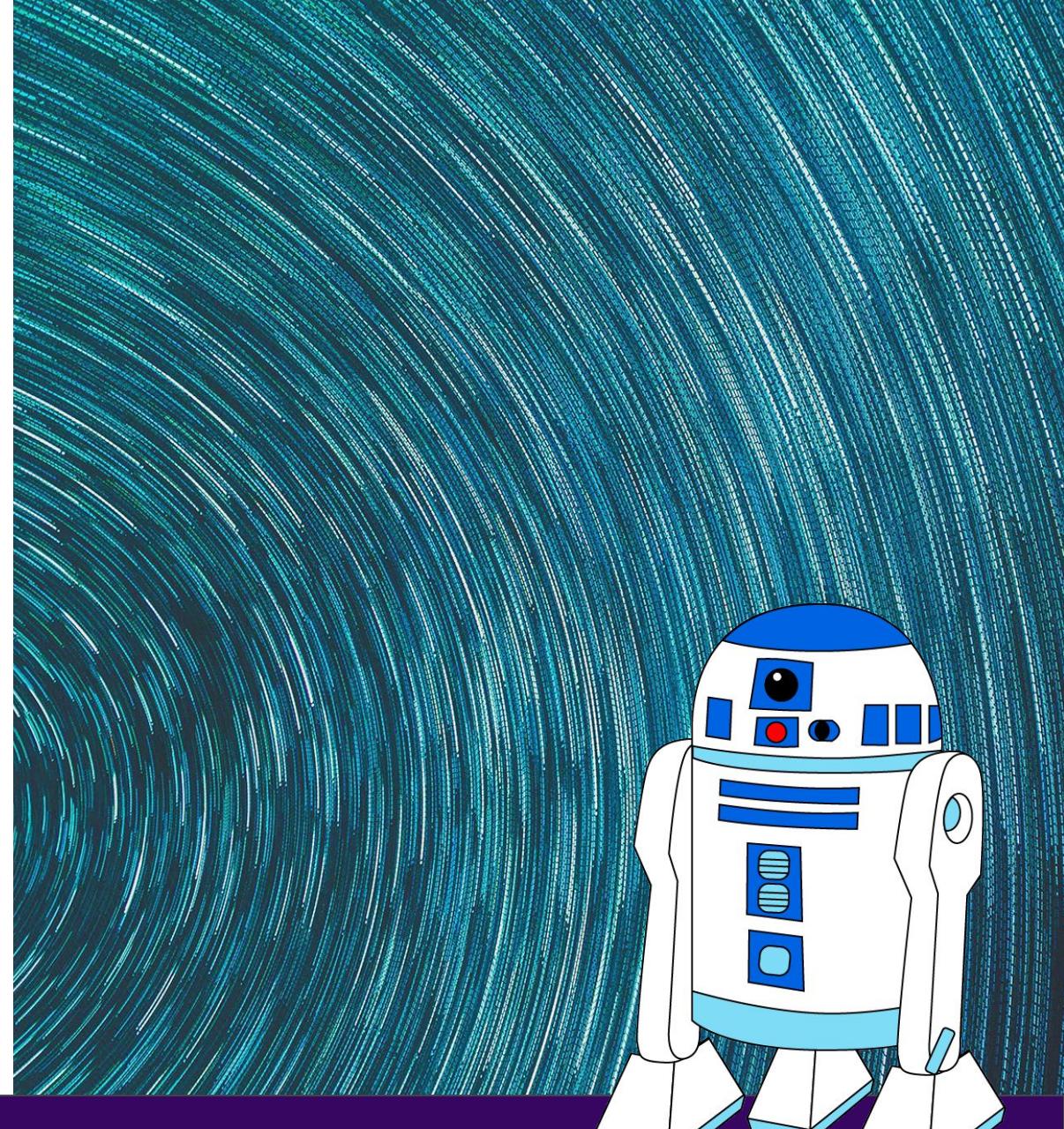


CIS 421/521:  
ARTIFICIAL INTELLIGENCE

# Search Problems

Professor Chris Callison-Burch



# Problem Solving Agents & Problem Formulation

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AIMA 3.1-3.3

# Reflex Agents

A simple reflex agent is one that selects an action based only on the **current percept**.

It **ignores** the rest of the **percept history**.



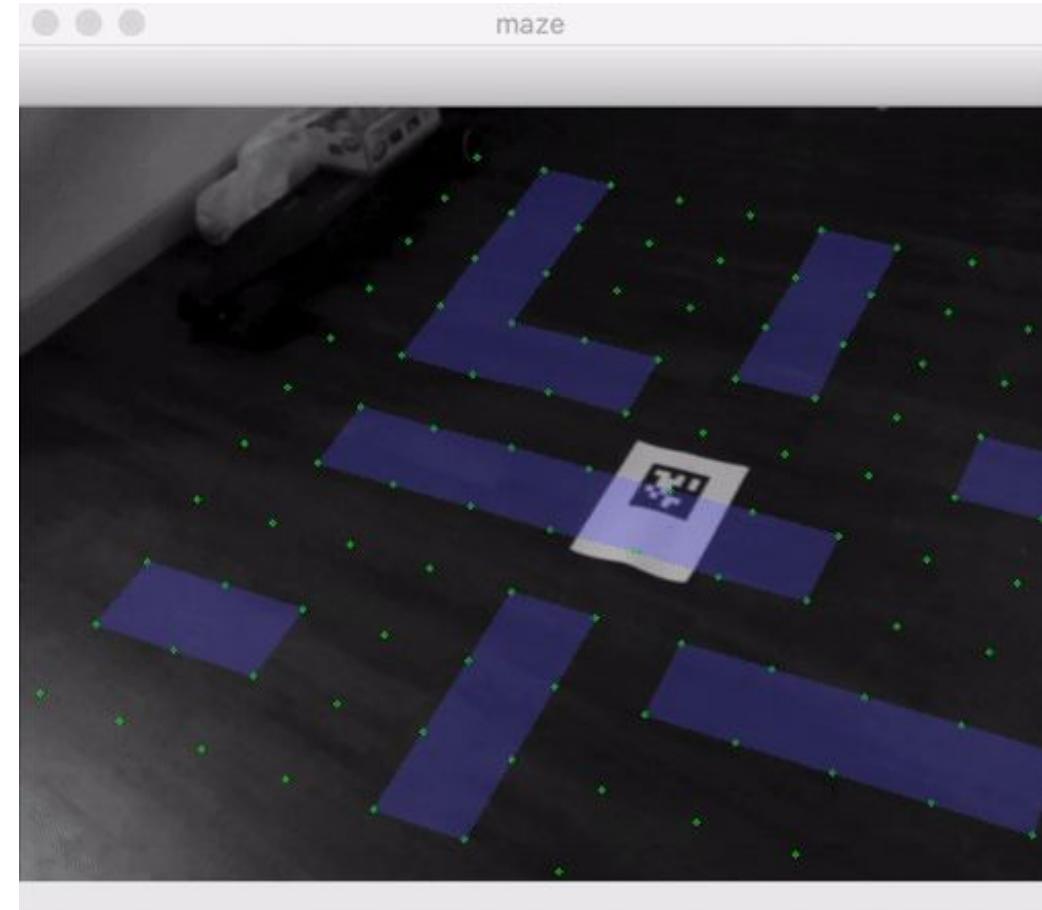
# Problem-Solving Agent

A problem-solving agent must **plan ahead**.

The computational process that it undertakes is called **search**.

It will consider a **sequence of actions** that form a **path** to a **goal state**.

Such a sequence is called a **solution**.



# Impact of Task Environments

The properties of the task environments change the types of solutions that we need.

If an environment is:

- **Fully observable**
- **Deterministic**
- **Known environment**

*The solution to any problem in such an environment is a fixed sequence of actions.*

In environments that are

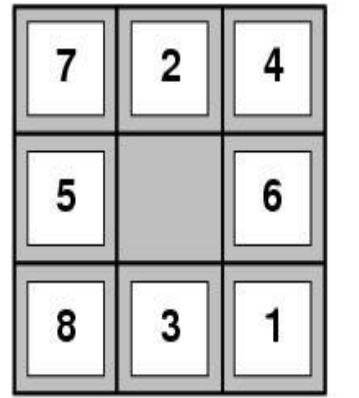
- **Partially observable** or
- **Nondeterministic**

The solution must recommend different future actions depending on the what percepts it receives. This could be in the form of a *branching strategy*.

# Example search problem: 8-puzzle

Formulate **goal**

- Pieces to end up in order as shown...



Start State



Goal State



Formulate **search problem**

- States:** configurations of the puzzle ( $9!$  configurations)
- Actions:** Move one of the movable pieces ( $\leq 4$  possible)
- Performance measure:** minimize total moves

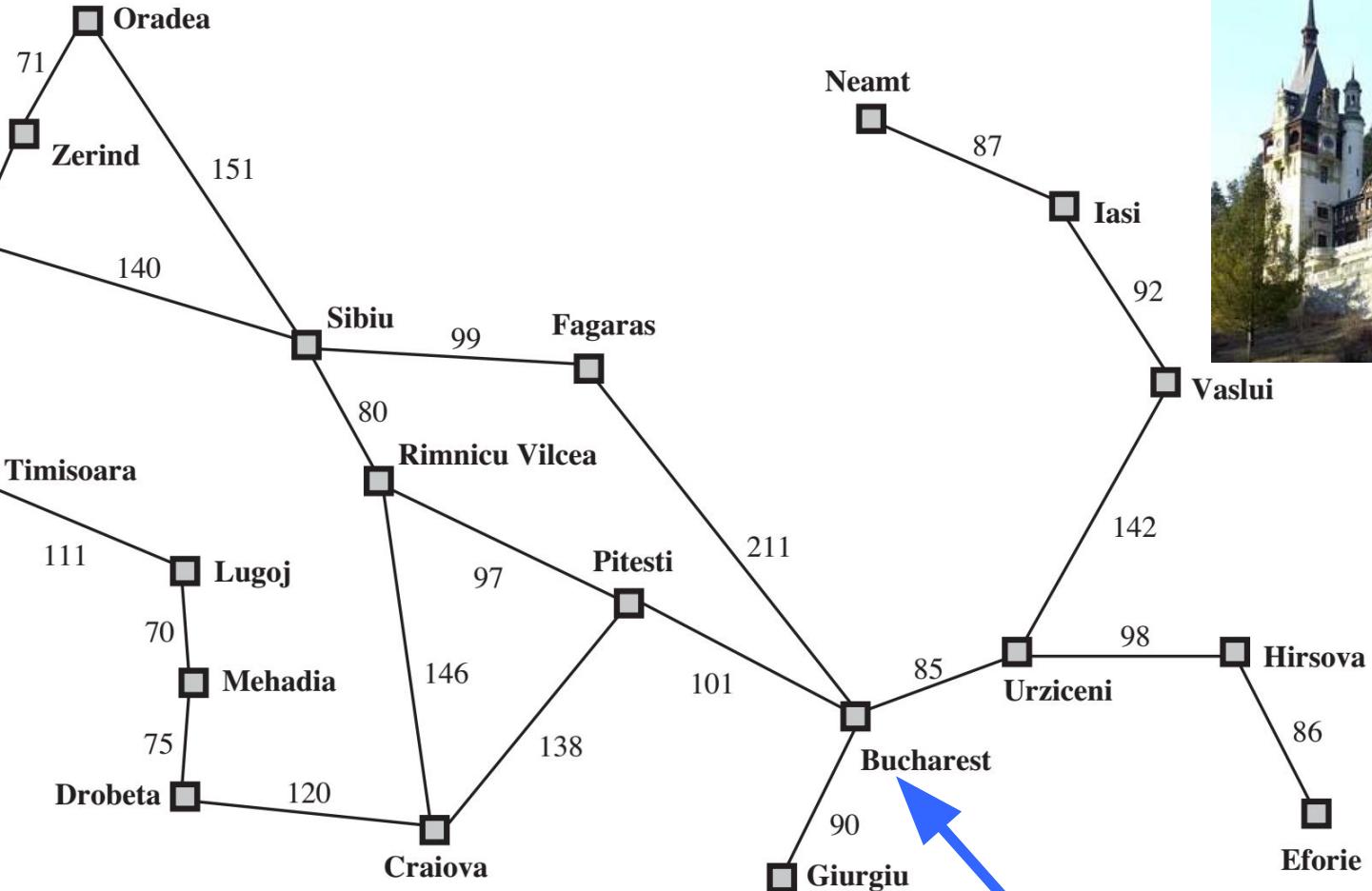
Find **solution**

- Sequence of pieces moved: 3,1,6,3,1,...

# Example search problem: Holiday in Romania



You are here



You need  
to be here

# Holiday in Romania

On holiday in Romania; currently in Arad

- Flight leaves tomorrow from Bucharest

Formulate **goal**

- Be in Bucharest

Formulate **search problem**

- States: various cities
- Actions: drive between cities
- Performance measure: minimize travel time / distance

Find **solution**

- Sequence of cities; e.g. Arad, Sibiu, Fagaras, Bucharest, ...

# More formally, a problem is defined by:

1. **States**: a set  $S$

2. An **initial state**  $s_i \in S$

3. **Actions**: a set  $A$

$\forall s \text{ } Actions(s) = \text{the set of actions that can be executed in } s, \text{ that are applicable in } s.$

4. **Transition Model**:  $\forall s \forall a \in Actions(s) \text{ } Result(s, a) \rightarrow s_r$

$s_r$  is called a **successor** of  $s$

$\{s_i\} \cup Successors(s_i)^* = \text{state space}$

5. **Path cost** (*Performance Measure*): Must be additive, e.g. sum of distances, number of actions executed, ...

$c(x, a, y)$  is the step cost, assumed  $\geq 0$

- (where action  $a$  goes from state  $x$  to state  $y$ )

6. **Goal test**:  $Goal(s)$

Can be implicit, e.g.  $checkmate(s)$

$s$  is a *goal state* if  $Goal(s)$  is true

# Vacuum World

**States:** A state of the world says which objects are in which cells.

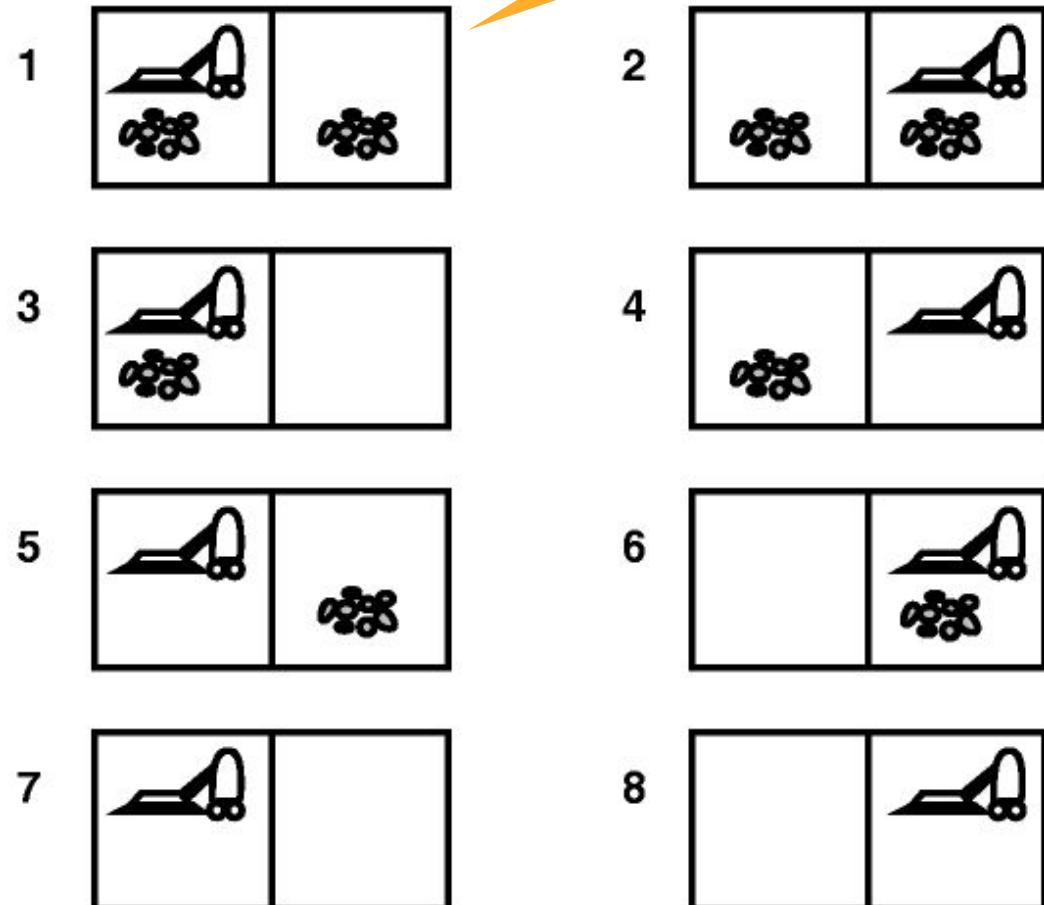
In a simple two cell version,

- the agent can be in either cell
- each cell can have dirt or not

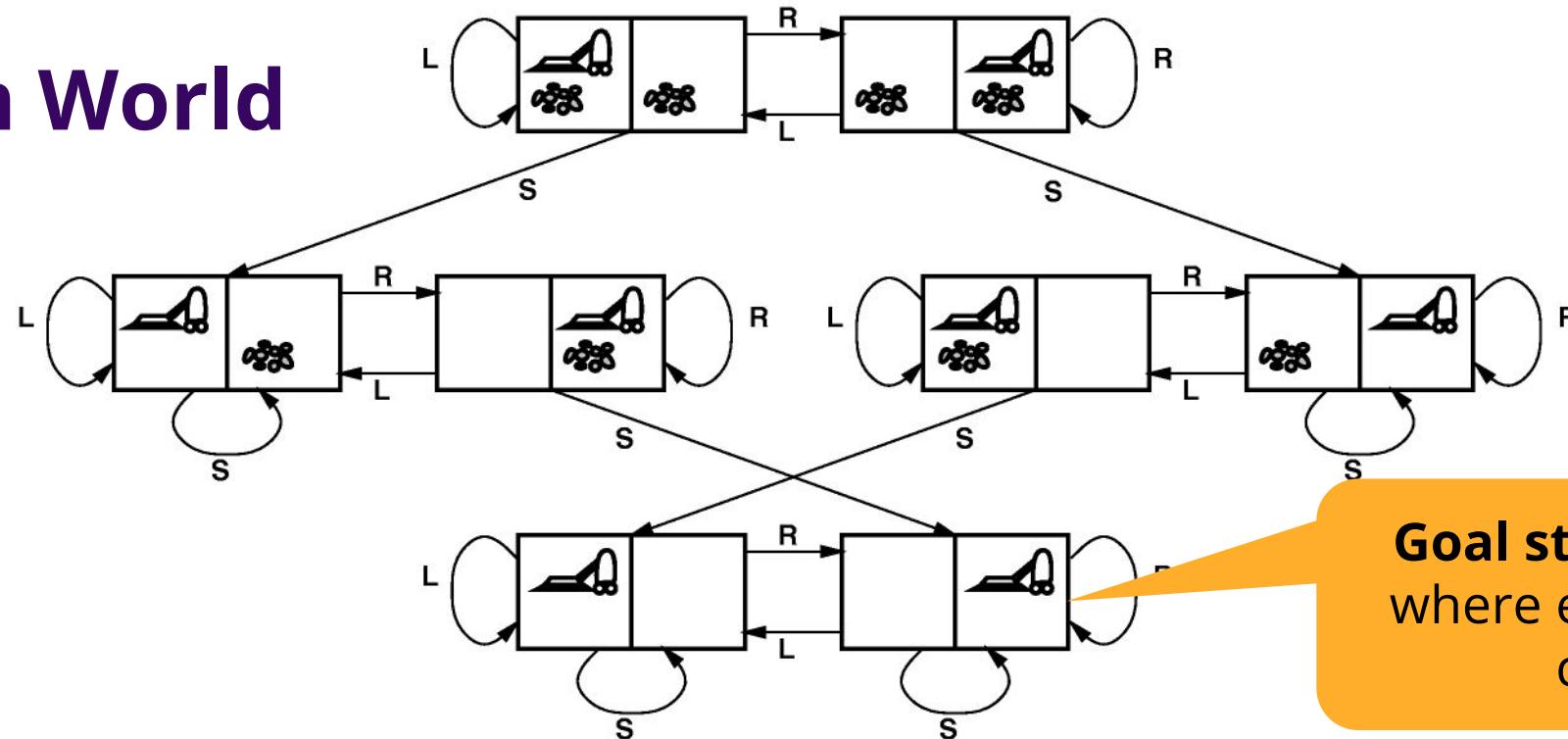
2 cells \* 2 positions for agent \* 2 possibilities for dirt = 8 states.

With  $n$  cells, there are  $n*2^n$  states.

One state is designated as the initial state



# Vacuum World



**Goal states:** States where everything is clean.

## **Actions:**

- *Suck*
  - *Move Left*
  - *Move Right*
  - *(Move Up)*
  - *(Move Down)*

**Action Cost:** Each action costs 1

## Transition:

Suck – removes dirt

Move – moves in that direction, unless agent hits a wall, in which case it stays put.

# Solutions & Optimal Solutions

A **solution** is a sequence of **actions** from the **initial state** to a **goal state**.

**Optimal Solution:** A solution is **optimal** if no solution has a lower **path cost**.

# Art: Formulating a Search Problem

Decide:

Which properties matter & how to represent

- *Initial State, Goal State, Possible Intermediate States*

Which actions are possible & how to represent

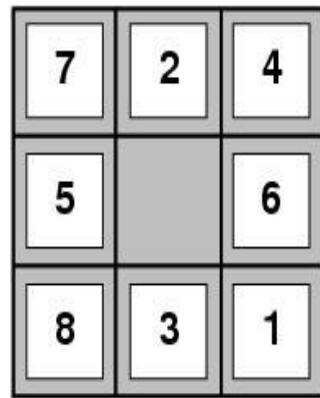
- *Operator Set: Actions and Transition Model*

Which action is next

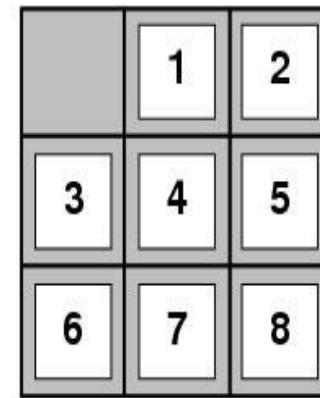
- *Path Cost Function*

**Formulation greatly affects combinatorics of search space and therefore speed of search**

# Example: 8-puzzle



Start State



Goal State

States? List

Initial state?

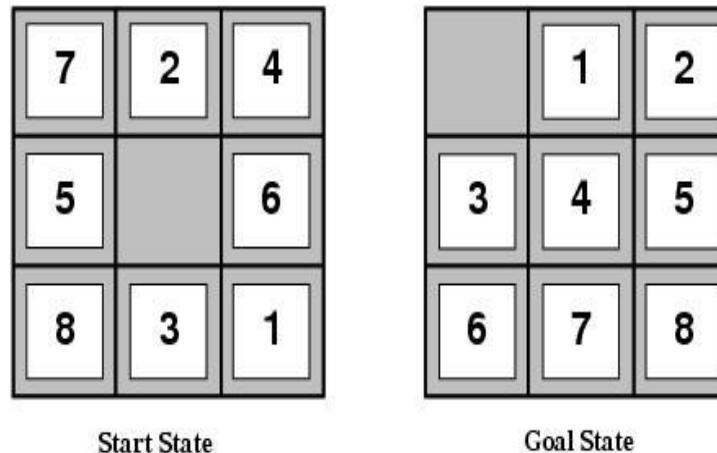
Actions? {Left}

Transition Model?

Goal test? Che

Path cost? Nur

# Example: 8-puzzle



States? List of 9 locations- e.g., [7,2,4,5,-,6,8,3,1]

Initial state? [7,2,4,5,-,6,8,3,1]

Actions? {*Left, Right, Up, Down*}

Transition Model? ...

Goal test? Check if goal configuration is reached

Path cost? Number of actions to reach goal

# Hard subtask: Selecting a state space

Real world is absurdly complex

State space must be **abstracted** for problem solving

(abstract) **State** = set (equivalence class) of real-world states

(abstract) **Action** = equivalence class of combinations of real-world actions

- e.g. *Arad → Zerind* represents a complex set of possible routes, detours, rest stops, etc
- The abstraction is valid if the path between two states is reflected in the real world

Each abstract action should be “easier” than the real problem

# Useful Concepts

**State space:** the set of all states reachable from the initial state by *any* sequence of actions

- *When several operators can apply to each state, this gets large very quickly*
- *Might be a proper subset of the set of configurations*

**Path:** a sequence of actions leading from one state  $s_j$  to another state  $s_k$

**Frontier:** those states that are available for *expanding* (for applying legal actions to)

**Solution:** a path from the initial state  $s_i$  to a state  $s_f$  that satisfies the goal test

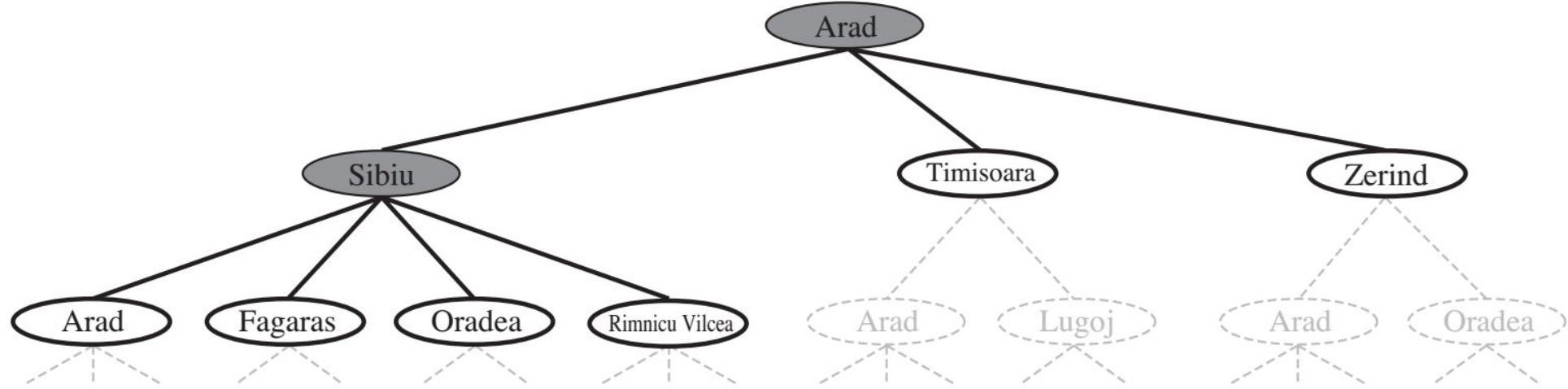
# Basic search algorithms: *Tree Search*

Generalized algorithm to solve search problems

**Enumerate in some order all possible paths from the initial state**

- Here: search through ***explicit tree generation***
  - ROOT= initial state.
  - Nodes in search tree generated through ***transition model***
  - Tree search treats different paths to the same node as distinct

# Generalized tree search



```
function TREE-SEARCH(problem, strategy) return a solution or failure
```

  Initialize frontier to the *initial state of the problem*

  do

    if the frontier is empty then return *failure*

      choose leaf node for expansion according to strategy & remove from frontier

      if node contains goal state then return *solution*

      else expand the node and add resulting nodes to the frontier

The strategy determines  
search process!



# 8-Puzzle: States and Nodes

A **state** is a (representation of a) **physical configuration**

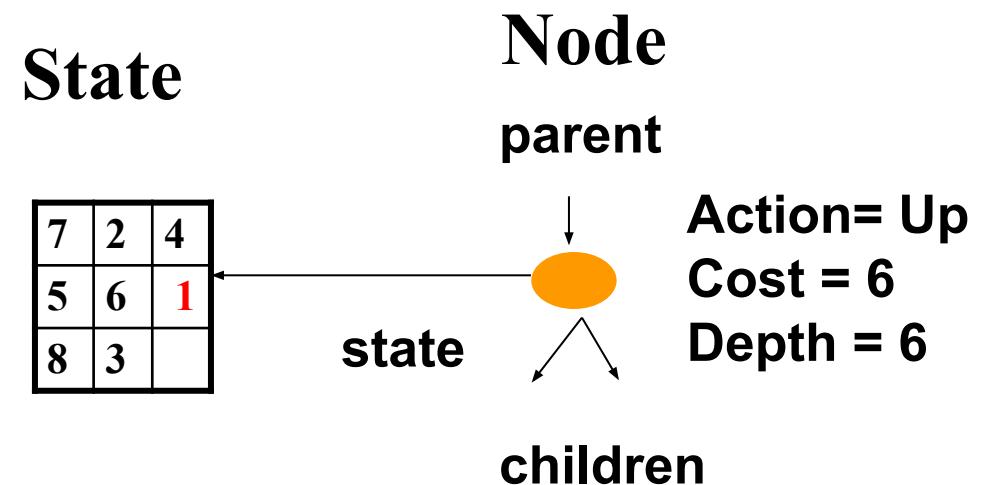
A **node** is a data structure constituting **part of a search tree**

- Also includes *parent*, *children*, *depth*, *path cost*  $g(x)$
- Here  $\text{node} = \langle \text{state}, \text{parent-node}, \text{children}, \text{action}, \text{path-cost}, \text{depth} \rangle$

States do not have parents, children, depth or path cost!

The EXPAND function

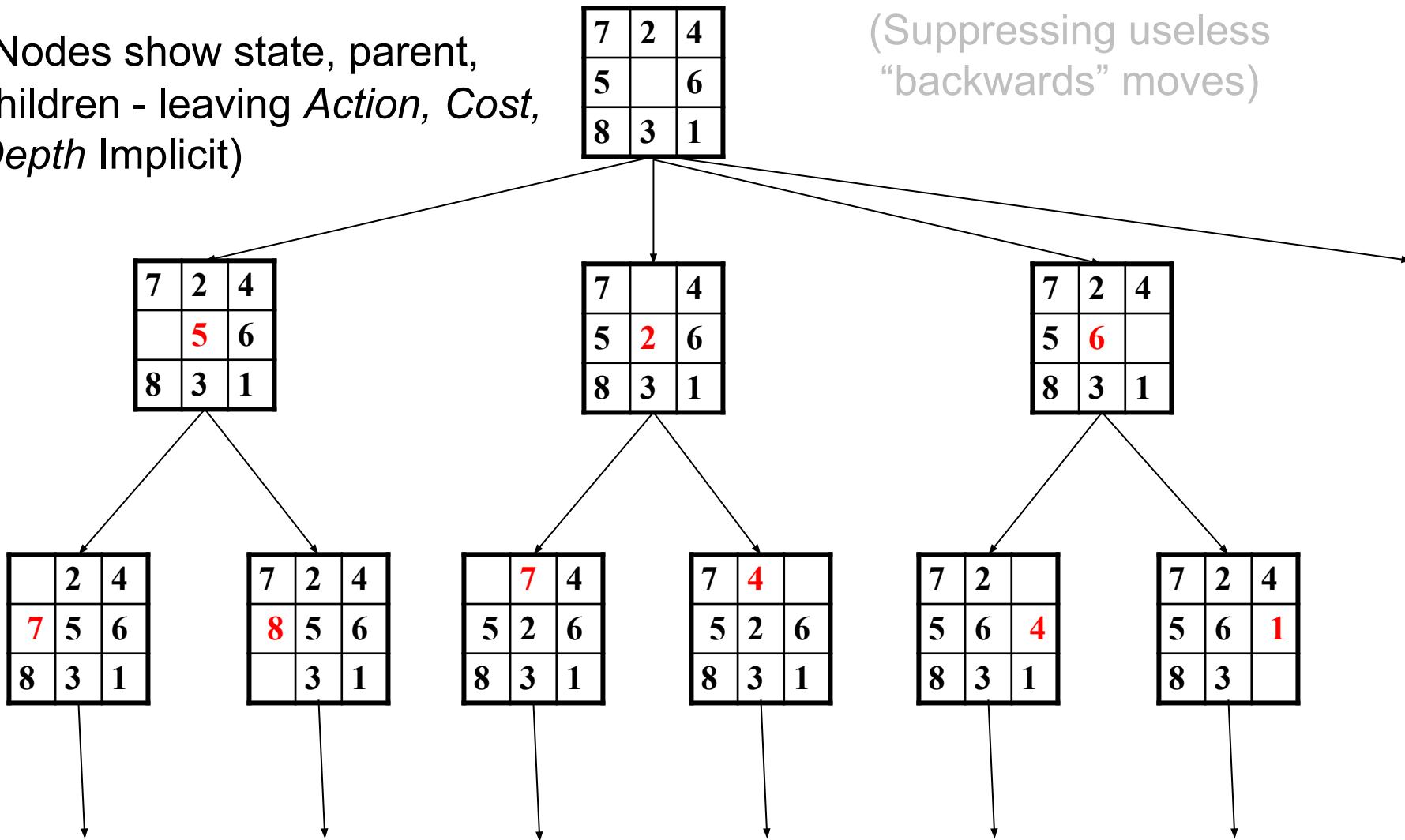
- uses the Actions and Transition Model to create the corresponding states
  - creates new nodes,
  - fills in the various fields



# 8-Puzzle Search Tree

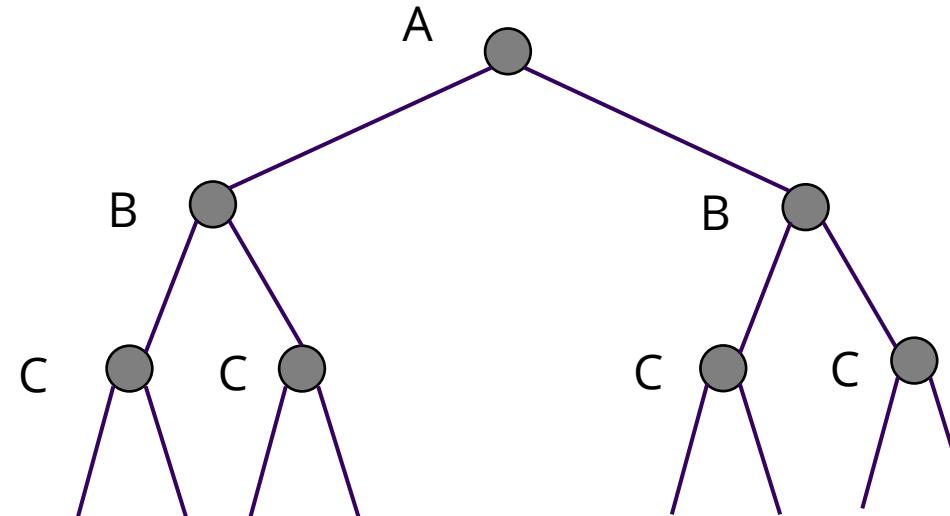
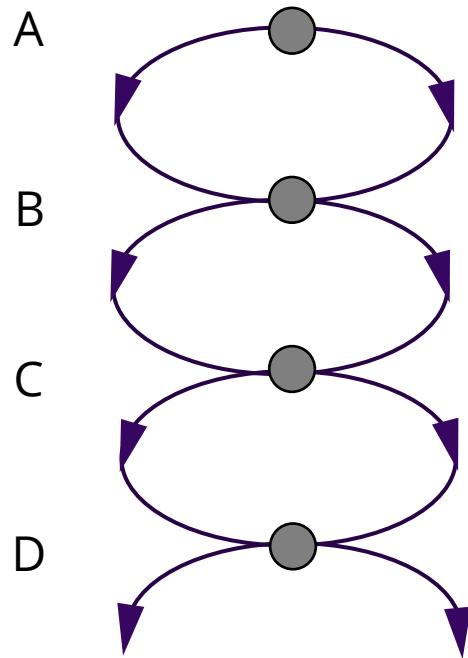
(Nodes show state, parent,  
children - leaving *Action, Cost,  
Depth* Implicit)

(Suppressing useless  
“backwards” moves)

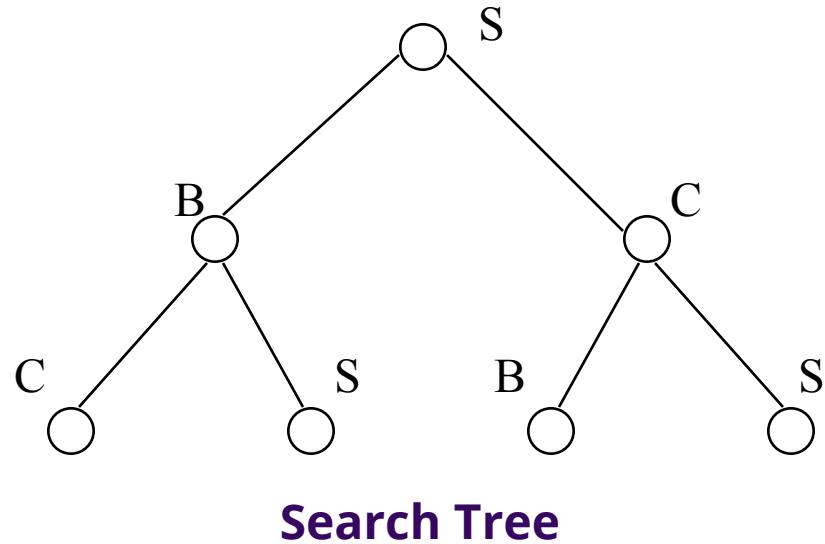
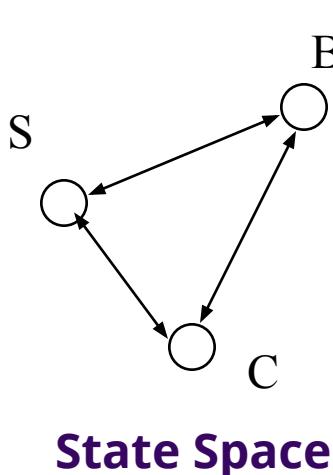


# Problem: Repeated states

Failure to detect ***repeated states*** can turn a linear problem into an ***exponential*** one!



# Solution: Graph Search!



## Graph search

- Simple Mod from tree search: *Check to see if a node has been visited before adding to search queue*
  - must keep track of all possible states (can use a lot of memory)
  - e.g., 8-puzzle problem, we have  $9!/2 \approx 182K$  states

# Graph Search vs Tree Search

**function** TREE-SEARCH(*problem*) **returns** a solution, or failure

  initialize the frontier using the initial state of *problem*

**loop do**

**if** the frontier is empty **then return** failure

    choose a leaf node and remove it from the frontier

**if** the node contains a goal state **then return** the corresponding solution  
    expand the chosen node, adding the resulting nodes to the frontier

**function** GRAPH-SEARCH(*problem*) returns a solution, or failure

  initialize the frontier using the initial state of *problem*

**initialize the explored set to be empty**

**loop do**

**if** the frontier is empty **then return** failure

    choose a leaf node and remove it from the frontier

**if** the node contains a goal state **then return** the corresponding solution

**add node to the explored set**

    expand the chosen node, adding the resulting nodes to the frontier

**only if not in the frontier of explored set**

An informal description of the general tree-search and graph-search algorithms. The parts of GRAPH-SEARCH marked in bold italics are the additions needed to handle repeat states.

# Uninformed Search Strategies

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AIMA 3.3-3.4



# ***Uninformed* search strategies:**

AKA “Blind search”

Uses only information available in problem definition

Informally:

***Uninformed search:*** All non-goal nodes in frontier look equally good

***Informed search:*** Some non-goal nodes can be ranked above others.

# Search Strategies

Review: **Strategy** = order of tree expansion

- Implemented by different queue structures (LIFO, FIFO, priority)

Dimensions for evaluation

- **Completeness**- always find the solution?
- **Optimality** - finds a least cost solution (lowest path cost) first?
- **Time complexity** - # of nodes generated (*worst case*)
- **Space complexity** - # of nodes simultaneously in memory (*worst case*)

Time/space complexity variables

- $b$ , *maximum branching factor* of search tree
- $d$ , *depth* of the shallowest goal node
- $m$ , maximum length of any path in the state space (potentially  $\infty$ )

# Introduction to *space* complexity

You know about:

- “Big O” notation
- *Time complexity*

*Space complexity* is analogous to time complexity

Units of space are arbitrary

- Doesn’t matter because Big O notation ignores constant multiplicative factors
- Plausible Space units:
  - One Memory word
  - Size of any fixed size data structure
    - For example, size of fixed size node in search tree

# Review: Breadth-first search

Idea:

- Expand *shallowest* unexpanded node

Implementation:

- *frontier* is FIFO (First-In-First-Out) Queue:
  - Put successors at the *end* of *frontier* successor list.



Image credit: Dan Klein and  
Pieter Abbeel  
<http://ai.berkeley.edu>

# Breadth-first search (simplified)

**function** BREADTH-FIRST-SEARCH(*problem*) **returns** a solution, or failure

*node*  $\leftarrow$  a node with STATE = *problem*.INITIAL-STATE, PATH-COST = 0

**if** *problem*.GOAL-TEST(*node*.STATE) **then return** SOLUTION(*node*)

*frontier*  $\leftarrow$  a FIFO queue with *node* as the only element

*explored*  $\leftarrow$  an empty set

**loop do**

**if** EMPTY?(*frontier*) **then return** failure

*node*  $\leftarrow$  POP(*frontier*) /\* chooses the shallowest node in *frontier* \*/

add *node*.STATE to *explored*

**for each** *action* **in** *problem*.ACTIONS(*node*.STATE) **do**

*child*  $\leftarrow$  CHILD-NODE(*problem*, *node*, *action*)

**if** *child*.STATE is not in *explored* or *frontier* **then**

**if** *problem*.GOAL-TEST(*child*.STATE) **then return** SOLUTION(*child*)

*frontier*  $\leftarrow$  INSERT(*child*, *frontier*)

Position within  
queue of new items  
determines search  
strategy

Subtle: Node inserted into  
queue only after testing to  
see if it is a goal state

# Breadth-first search (simplified)

```
function BREADTH-FIRST-SEARCH(problem) returns a solution, or failure
  node <- a node with STATE = problem.INITIAL-STATE, PATH-COST=0
  if problem.GOAL-TEST(node.STATE) then return SOLUTION(node)
  frontier <- a FIFO queue with node as the only element
  loop do
    if EMPTY?(frontier) then return failure
    node <- POP(frontier) // chooses the shallowest node in frontier
    add node.STATE to explored
    for each action in problem.ACTIONS(node.STATE) do
      child <- CHILD-NODE(problem, node, action)
      if problem.GOAL-TEST(child.STATE) then return SOLUTION(child)
      frontier <- INSERT(child, frontier)
```

Position within  
queue of new items  
determines search  
strategy

Subtle: *Node inserted into  
queue only after testing to  
see if it is a goal state*

From Figure 3.11 Breadth-first search (ignores loops, repeated nodes)

# Properties of breadth-first search

**Complete?** Yes (if  $b$  is finite)

**Time Complexity?**  $1+b+b^2+b^3+\dots+b^d = O(b^d)$

**Space Complexity?**  $O(b^d)$  (keeps every node in memory)

**Optimal?** Yes, if cost = 1 per step  
(not optimal in general)

$b$ : maximum branching factor of search tree

$d$ : depth of the least cost solution

$m$ : maximum depth of the state space ( $\infty$ )

# Exponential Space (and time) Not Good...

- Exponential complexity uninformed search problems *cannot* be solved for any but the smallest instances.
- (*Memory* requirements are a bigger problem than *execution* time.)

DEPTH	NODES	TIME	MEMORY
2	110	0.11 milliseconds	107 kilobytes
4	11110	11 milliseconds	10.6 megabytes
6	$10^6$	1.1 seconds	1 gigabytes
8	$10^8$	2 minutes	103 gigabytes
10	$10^{10}$	3 hours	10 terabytes
12	$10^{12}$	13 days	1 petabytes
14	$10^{14}$	3.5 years	99 petabytes

Assumes b=10, 1M nodes/sec, 1000 bytes/node

# Review: Depth-first search

Idea:

- Expand *deepest* unexpanded node

Implementation:

- *frontier* is LIFO (Last-In-First-Out) Queue:
  - Put successors at the *front* of *frontier* successor list.



Image credit: Dan Klein and  
Pieter Abbeel  
<http://ai.berkeley.edu>

# Properties of depth-first search

**Complete?** No: fails in infinite-depth spaces, spaces with loops

- Modify to avoid repeated states along path
  - complete in finite spaces

**Time?**  $O(b^m)$ : terrible if  $m$  is much larger than  $d$

- but if solutions are dense, may be much faster than breadth-first

**Space?**  $O(b*m)$ , i.e., linear space!

**Optimal?** No

*b: maximum branching factor of search tree*

*d: depth of the least cost solution*

*m: maximum depth of the state space ( $\infty$ )*

# Depth-first vs Breadth-first

Use depth-first if

- *Space is restricted*
- There are many possible solutions with long paths and wrong paths are usually terminated quickly
- Search can be fine-tuned quickly

Use breadth-first if

- *Possible infinite paths*
- Some solutions have short paths
- Can quickly discard unlikely paths

# Outline for today's lecture

Formulating Search Problems – An Example

Search Fundamentals

Introduction to Uninformed Search

- Review of Breadth first and Depth-first search

*Iterative deepening search (AIMA 3.4.4-5)*

- *Strange Subroutine: Depth-limited search*
- *Depth-limited search + iteration = WIN!!*

# Search Conundrum

## Breadth-first

- Complete,
- Optimal
- but uses  $O(b^d)$  space*

## Depth-first

- Not complete *unless m is bounded*
- Not optimal
- Uses  $O(b^m)$  time; terrible if  $m \gg d$
- but only uses  $O(b*m)$  space*

How can we get the best of both?

# Depth-limited search: A building block

Depth-First search *but with depth limit  $\ell$*

- i.e. nodes at depth  $\ell$  *have no successors*.
- No infinite-path problem!

If  $\ell = d$  (by luck!), then optimal

- But:
  - If  $\ell < d$  then incomplete 😞
  - If  $\ell > d$  then not optimal 😞

Time complexity:  $O(b^l)$

Space complexity:  $O(bl)$  😊

# Iterative deepening search

A general strategy to find best depth limit  $l$ .

- Key idea: use *Depth-limited search* as subroutine, with increasing  $l$ .

```
For  $l = 0$  to  $\infty$  do
    depth-limited-search to level l
    if it succeeds
        then return solution
```

- *Complete & optimal*: Goal is always found at depth  $d$ , the depth of the shallowest goal-node.

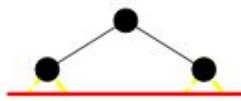
*Could this possibly be efficient?*

# Nodes constructed at each deepening

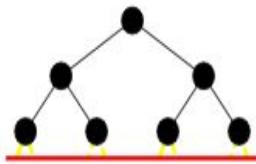
Depth 0: 0 (Given the node, doesn't *construct* it.)



Depth 1:  $b^1$  nodes



Depth 2:  $b$  nodes +  $b^2$  nodes



Depth 3:  $b$  nodes +  $b^2$  nodes +  $b^3$  nodes

...

# Total nodes constructed:

Depth 0: 0 (Given the node, doesn't *construct* it.)

Depth 1:  $b^1 = b$  nodes

Depth 2:  $b$  nodes +  $b^2$  nodes

- Depth 3:  $b$  nodes +  $b^2$  nodes +  $b^3$  nodes
- ...

Suppose the first solution is the last node at depth 3:

Total nodes constructed:

**3**\* $b$  nodes + **2**\* $b^2$  nodes + **1**\* $b^3$  nodes

# ID search, Evaluation: Time Complexity

- More generally, the time complexity is
  - $(d)b + (d-1)b^2 + \dots + (1)b^d = O(b^d)$

*As efficient in terms of  $O(\dots)$  as Breadth First Search:*

- $b + b^2 + \dots + b^d = O(b^d)$

# ID search, Evaluation

Complete: YES (no infinite paths) 😊

Time complexity:

$$O(b^d)$$

Space complexity:

$$O(bd^{\frac{d}{2}}) 😊$$

Optimal: YES if step cost is 1. 😊

# Summary of algorithms

Criterion	Breadth-First	Depth-First	Depth-limited	Iterative deepening
Complete?	<b>YES</b>	<b>NO</b>	<b>NO</b>	<b>YES</b>
Time	$b^d$	$b^m$	$b^l$	$b^d$
Space	$b^d$	$bm$	$bl$	$bd$
Optimal?	<b>YES</b>	<b>NO</b>	<b>NO</b>	<b>YES</b>

# Next Up: Informed Search

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AIMA 3.5-3.6

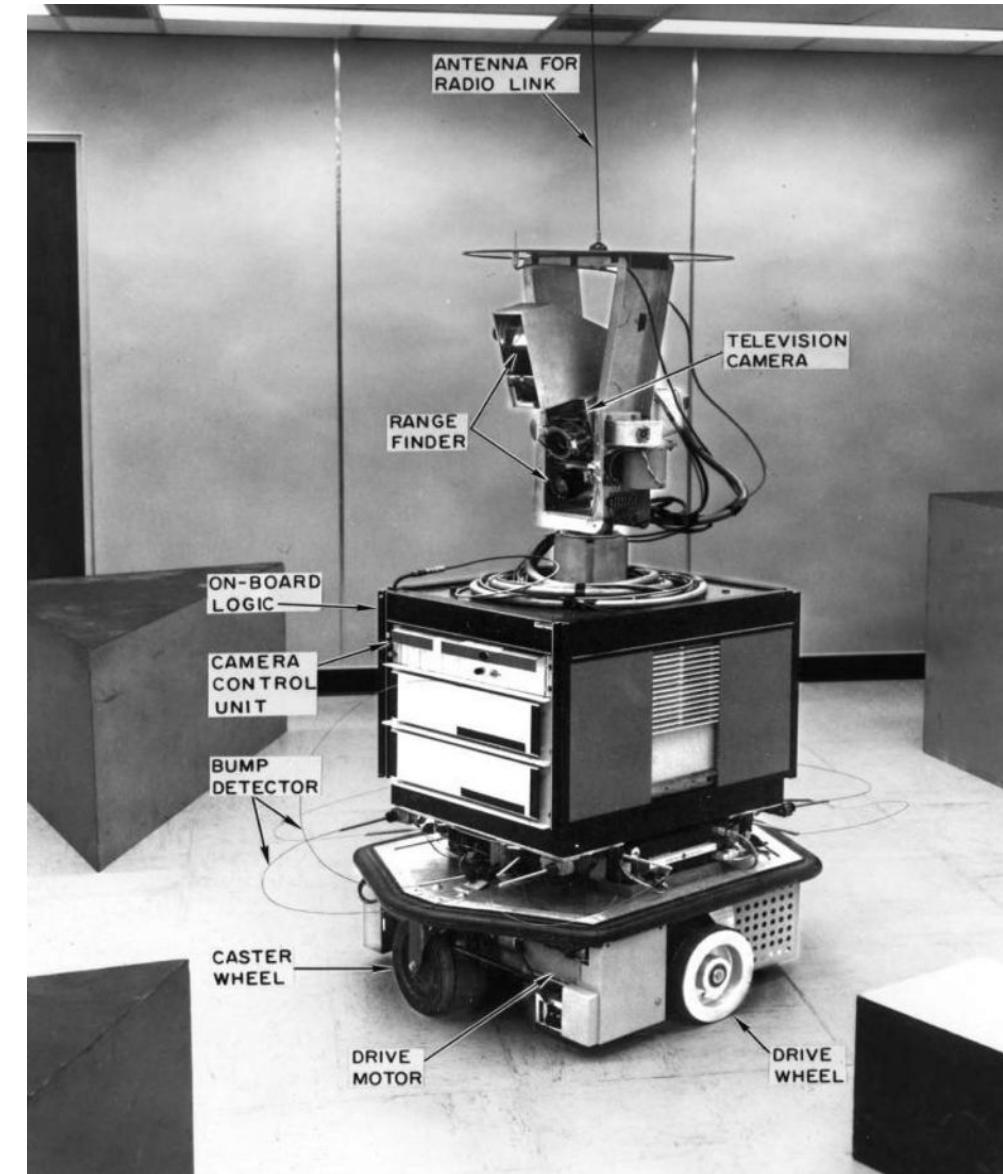
# Informed Search

An **informed search** strategy uses **domain-specific information** about the location of the goals in order to find a solution **more efficiently** than uninformed search.

Hints will come as part of a **heuristic function** denoted  $h(n)$ .

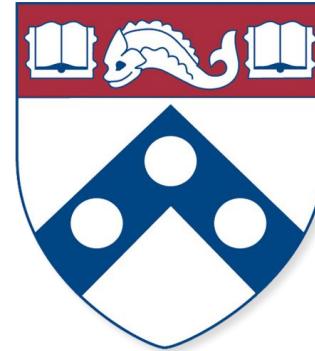
One of the most famous informed search algorithms is **A\*** which was developed for **robot navigation**.

Shakey the robot was developed at the Stanford Research Institute from 1966 to 1972.



<https://www.youtube.com/watch?v=7bsEN8mwUB8>





# Penn Engineering

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