Machine-Level Programming III: Switch Statements and IA32 Procedures

15-213: Introduction to Computer Systems
6th Lecture, Sep. 9, 2010

Instructors:

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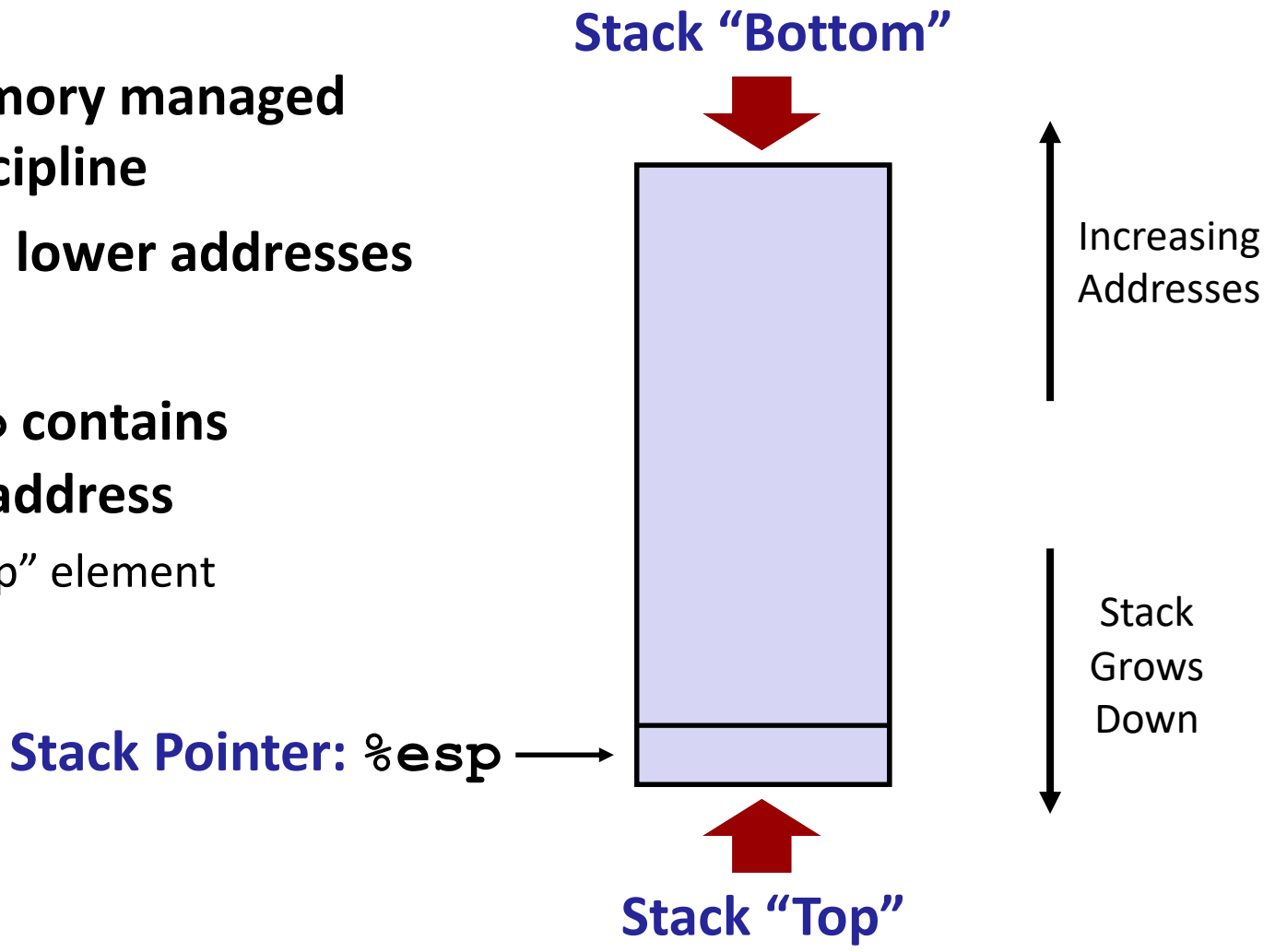
Today

- **IA 32 Procedures**
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers

IA32 Stack

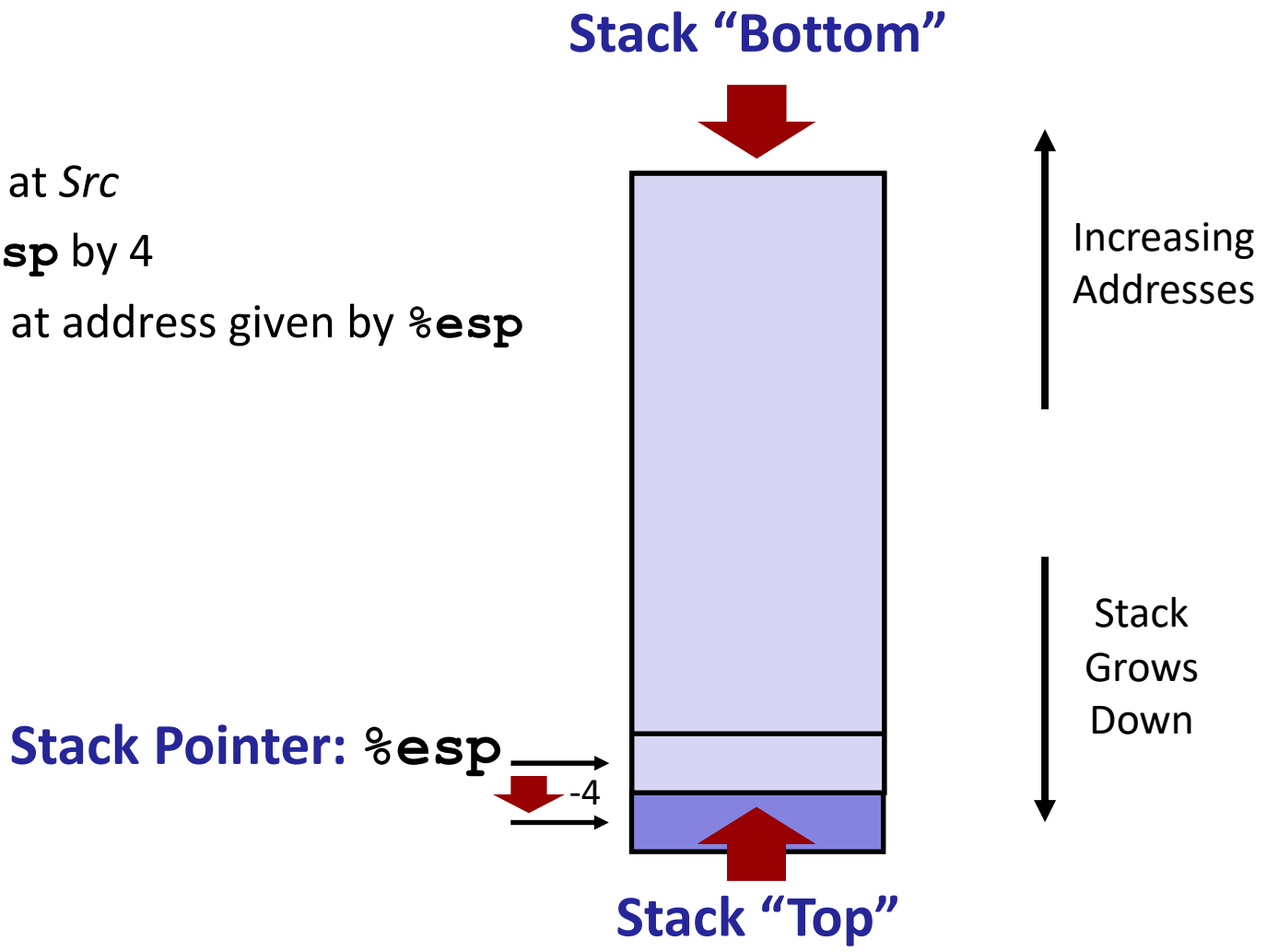
- Region of memory managed with stack discipline
- Grows toward lower addresses

- Register `%esp` contains lowest stack address
  - address of "top" element

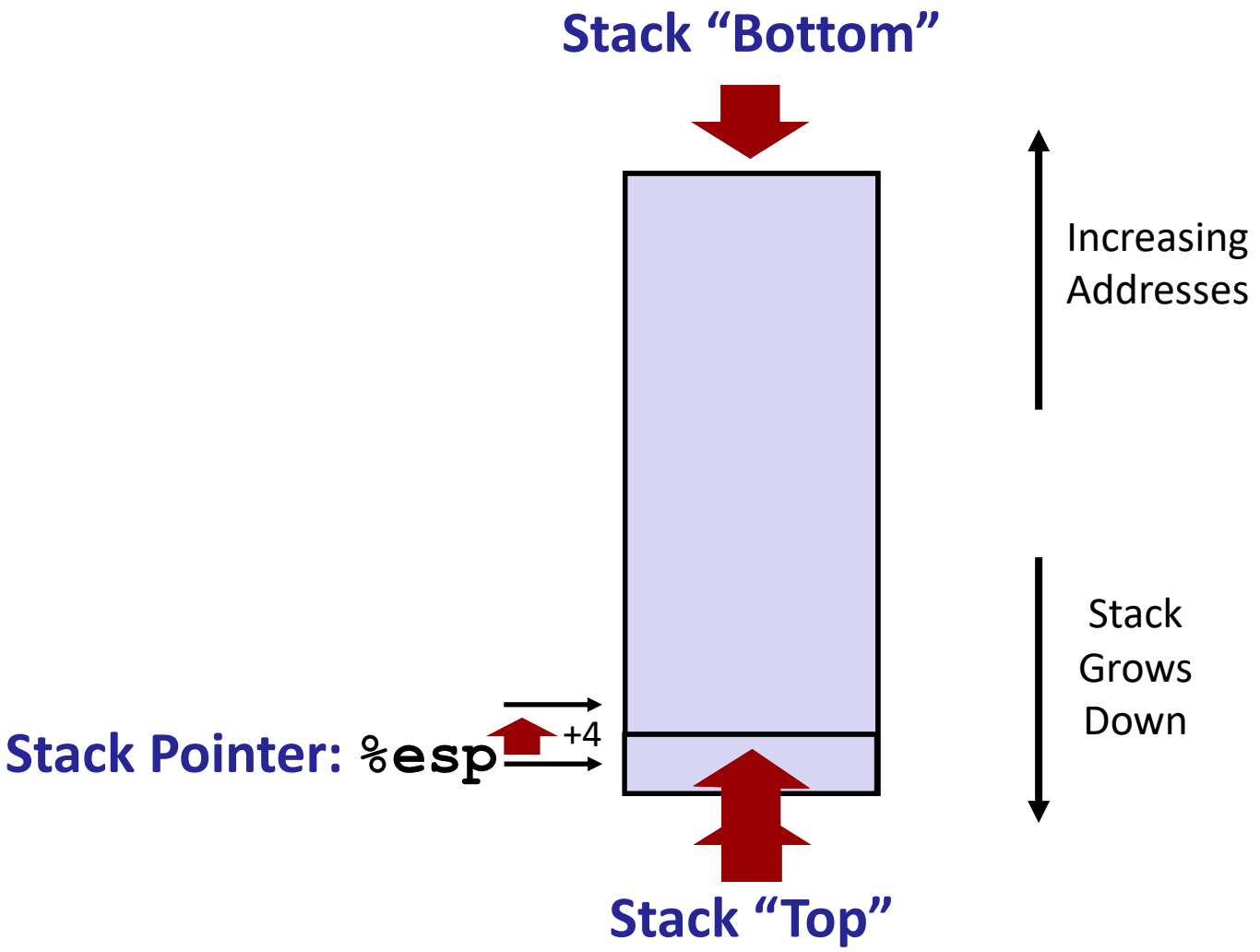


IA32 Stack: Push

- **`pushl` *Src***
  - Fetch operand at *Src*
  - Decrement **`%esp`** by 4
  - Write operand at address given by **`%esp`**



IA32 Stack: Pop



Procedure Control Flow

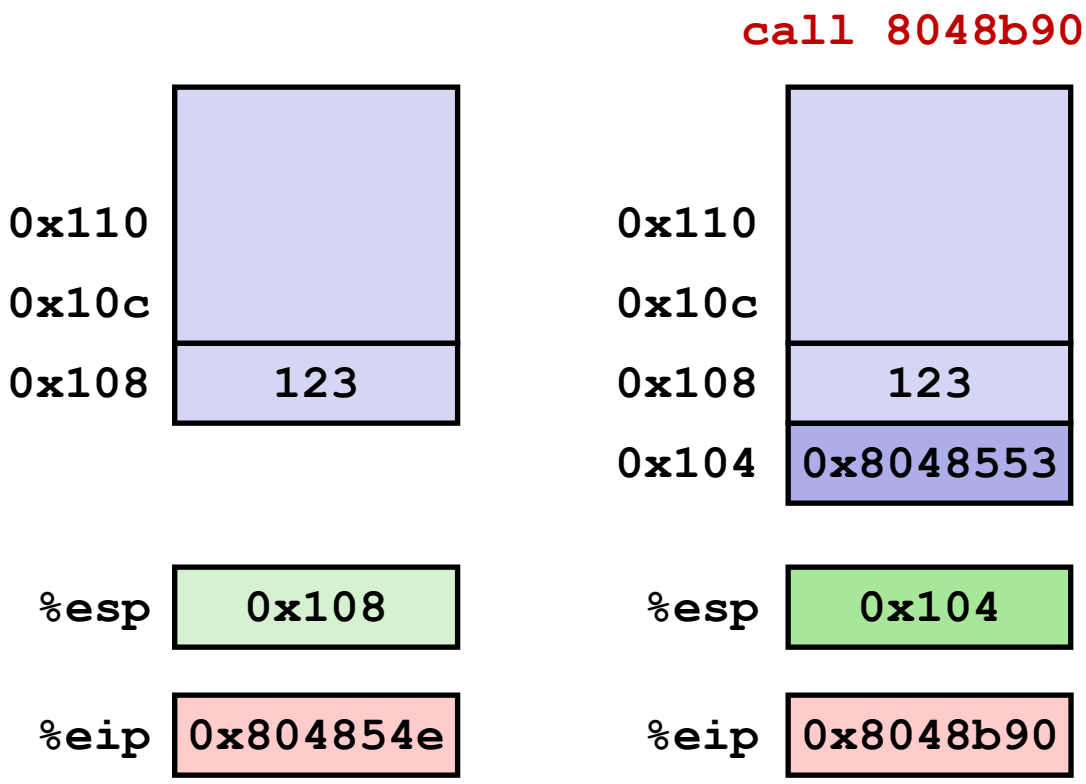
- **Use stack to support procedure call and return**
- **Procedure call: `call label`**
  - Push return address on stack
  - Jump to ***label***
- **Return address:**
  - Address of the next instruction right after call
  - Example from disassembly

```
804854e:  e8 3d 06 00 00     call   8048b90 <main>
8048553:  50                 pushl  %eax
```

  - Return address = `0x8048553`
- **Procedure return: `ret`**
  - Pop address from stack
  - Jump to address

Procedure Call Example

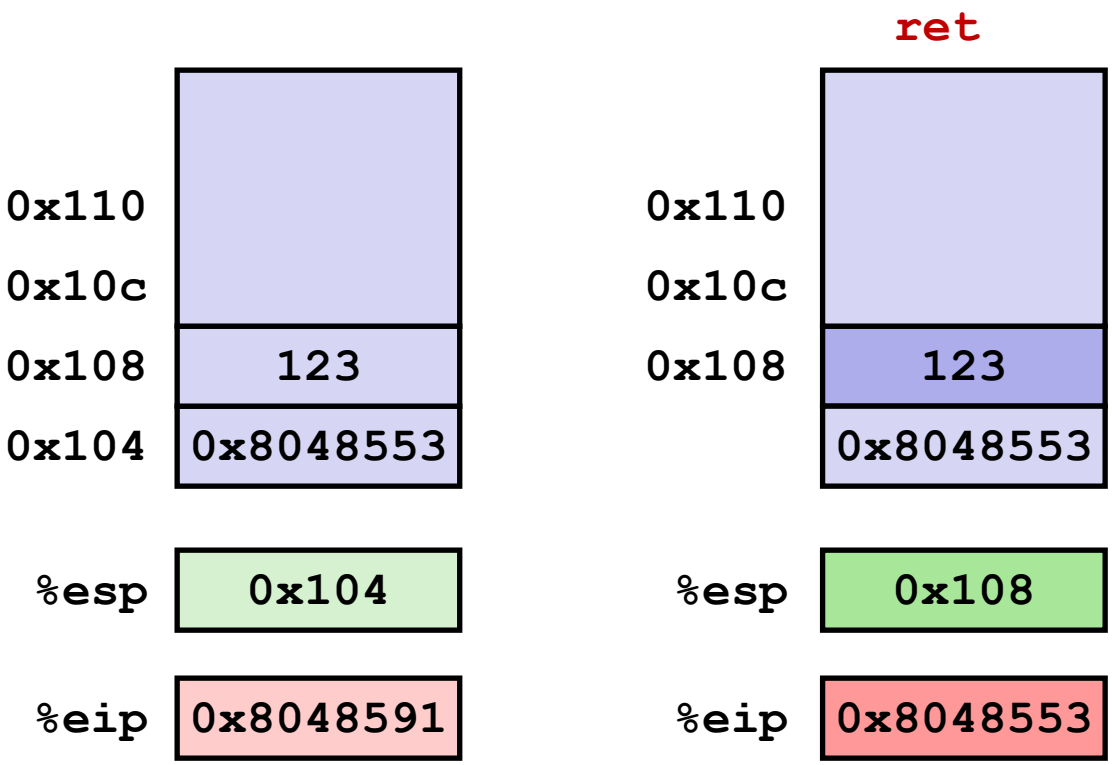
```
804854e:    e8 3d 06 00 00        call   8048b90 <main>
8048553:    50                    pushl  %eax
```



%eip: program counter

Procedure Return Example

```
8048591:      c3                          ret
```



%eip: program counter

Stack-Based Languages

- **Languages that support recursion**
  - e.g., C, Pascal, Java
  - Code must be "*Reentrant*"
    - Multiple simultaneous instantiations of single procedure
  - Need some place to store state of each instantiation
    - Arguments
    - Local variables
    - Return pointer
- **Stack discipline**
  - State for given procedure needed for limited time
    - From when called to when return
  - Callee returns before caller does
- **Stack allocated in *Frames***
  - state for single procedure instantiation

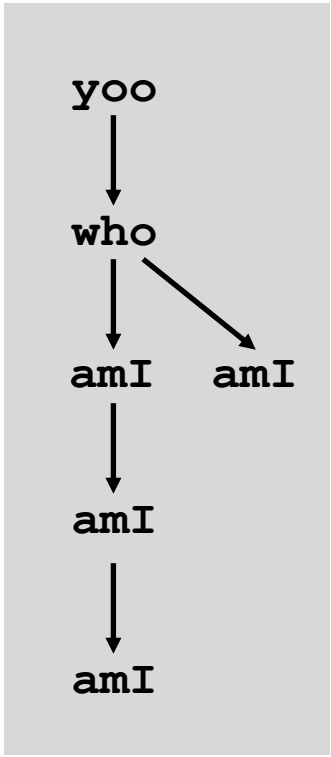
Call Chain Example

```
yoo(…)
{
   •
   •
   who();
   •
   •
}
```

```
who(…)
{
   • • •
   amI();
   • • •
   amI();
   • • •
}
```

```
amI(…)
{
   •
   •
   amI();
   •
   •
}
```

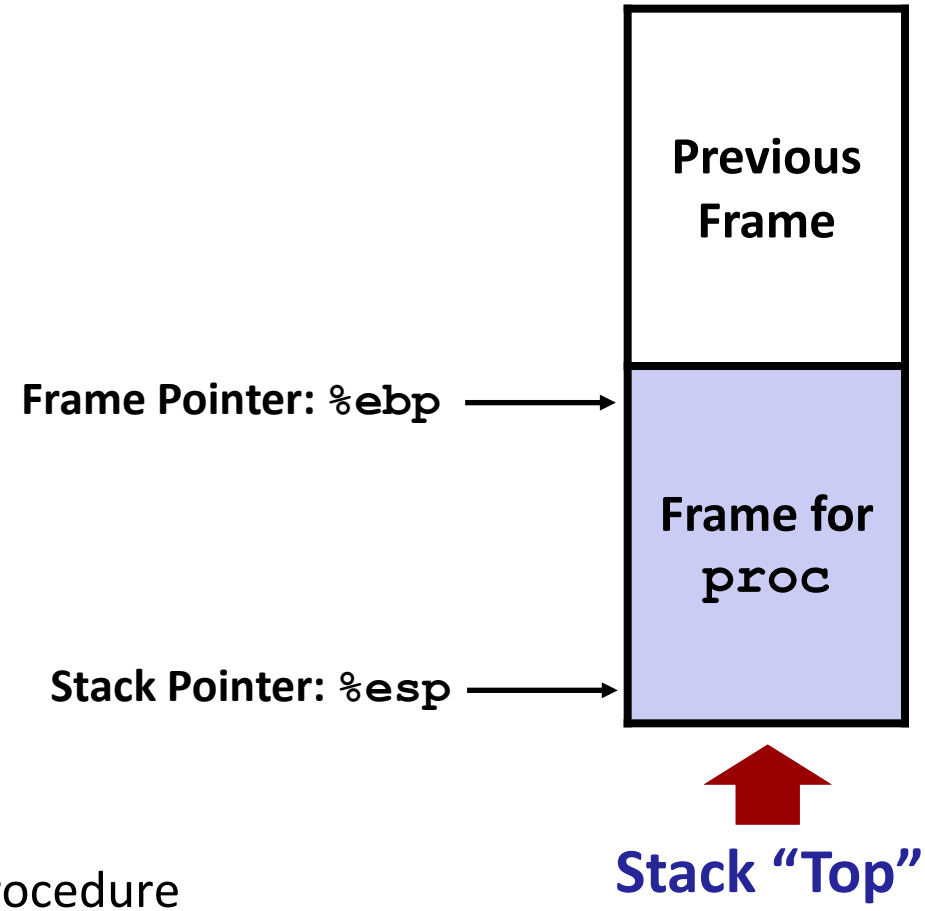
Example Call Chain



Procedure `amI()` is recursive

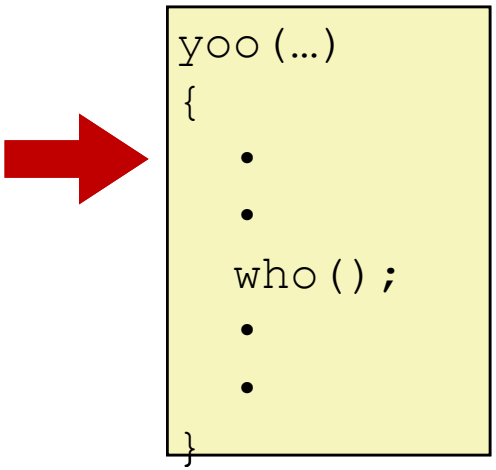
Stack Frames

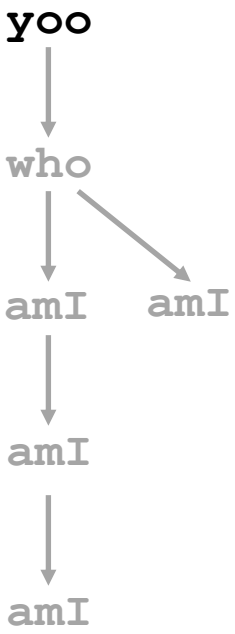
- **Contents**
  - Local variables
  - Return information
  - Temporary space

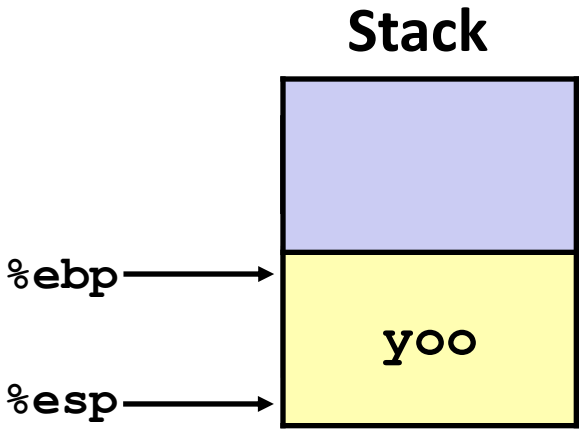


- **Management**
  - Space allocated when enter procedure
    - "Set-up" code
  - Deallocated when return
    - "Finish" code

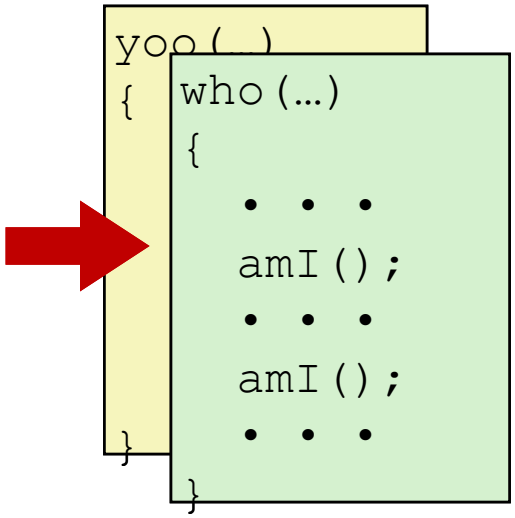
Example

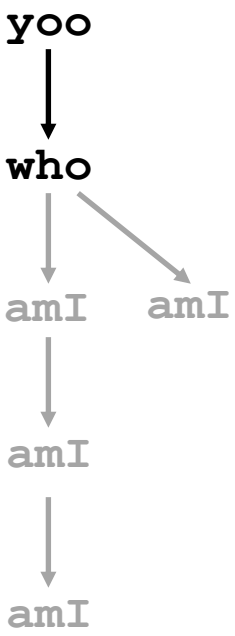


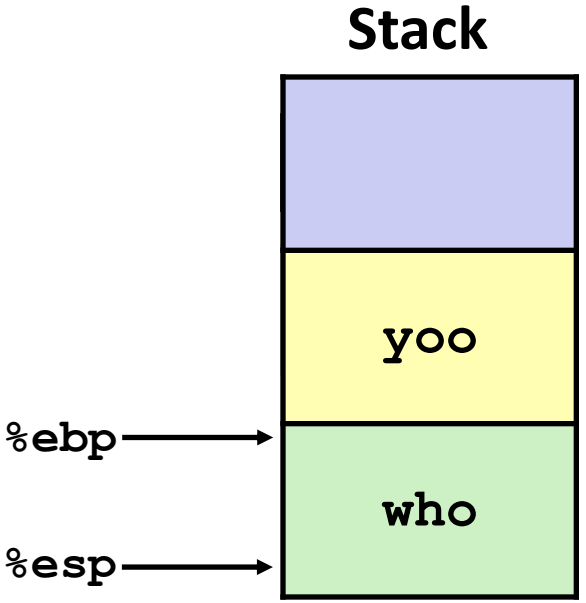




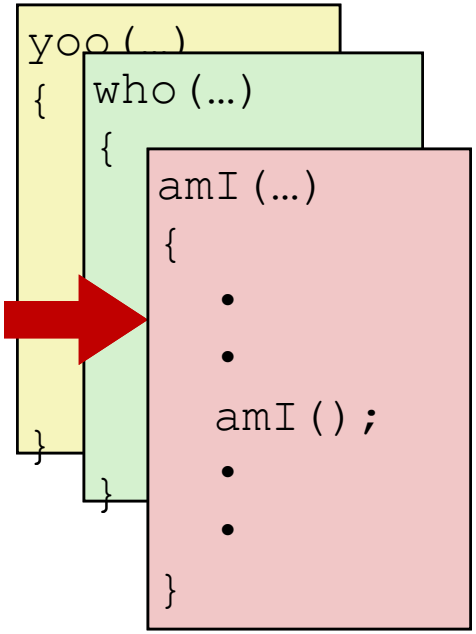
Example

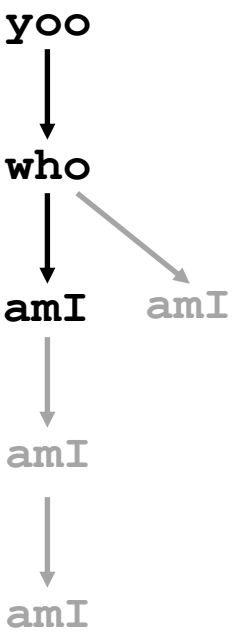


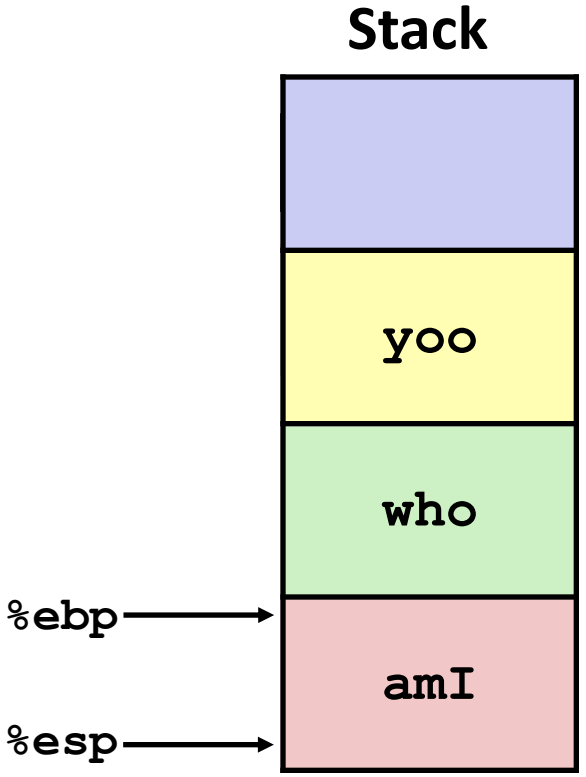




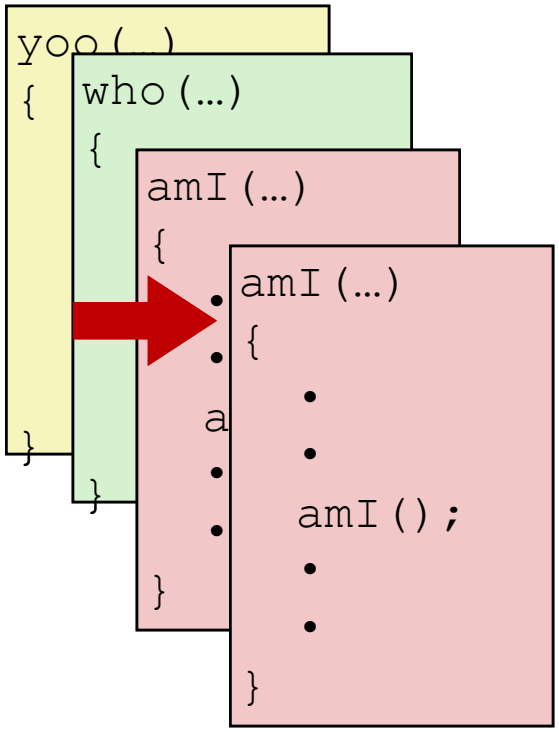
Example


yoo( )
{
who(…)
{
amI(…)
{
•
•
amI();
•
•
}
}
}


yoo
who
amI
amI
amI
amI


Stack
yoo
who
amI
%ebp
%esp

Example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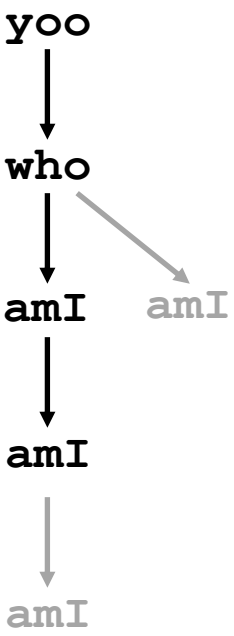
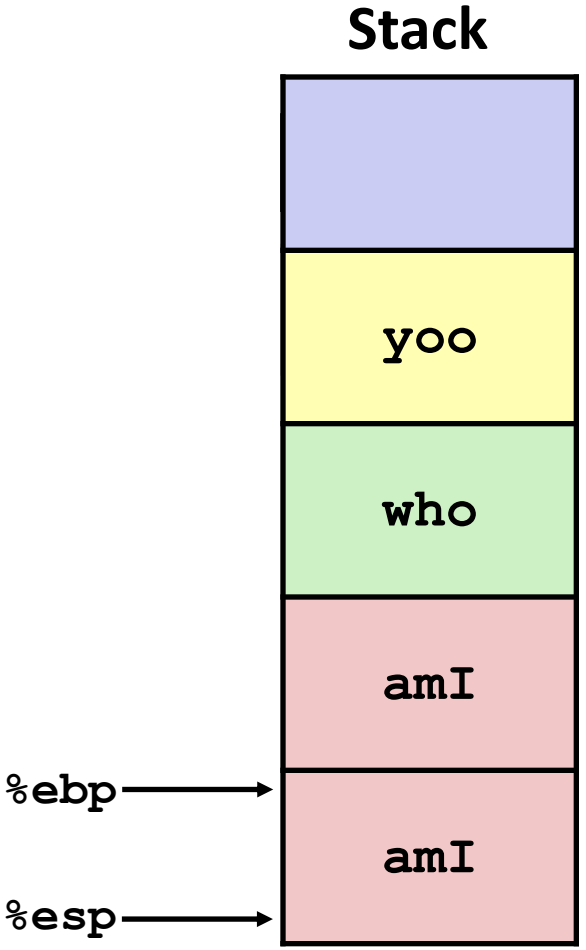
who(…)
{
amI(…)
{
amI(…)
{
amI();
}
}


yoo
who
amI
amI
amI
amI


Stack
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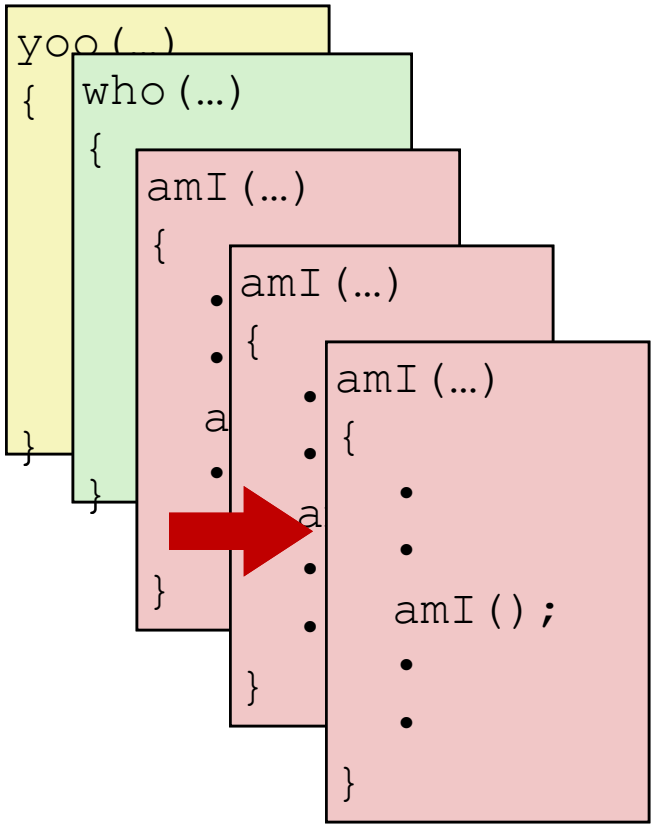
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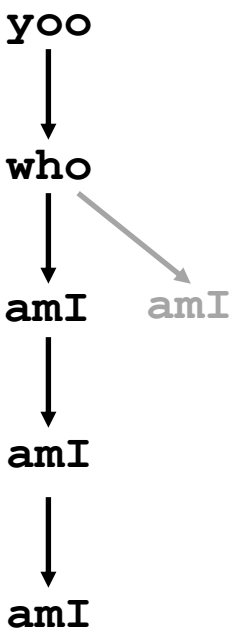
yoo
who
amI
amI
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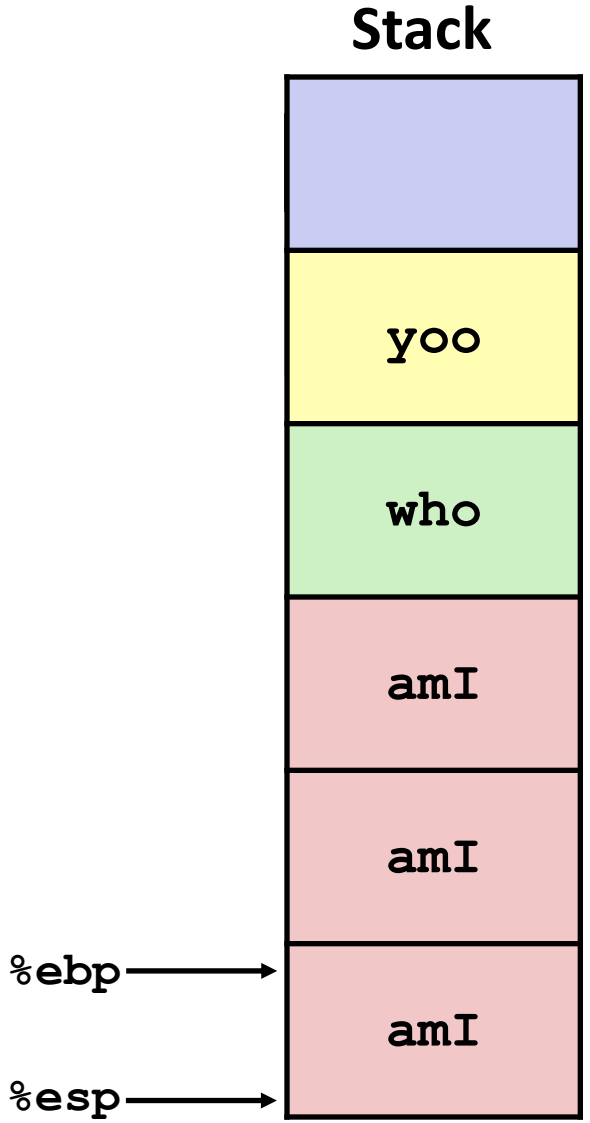
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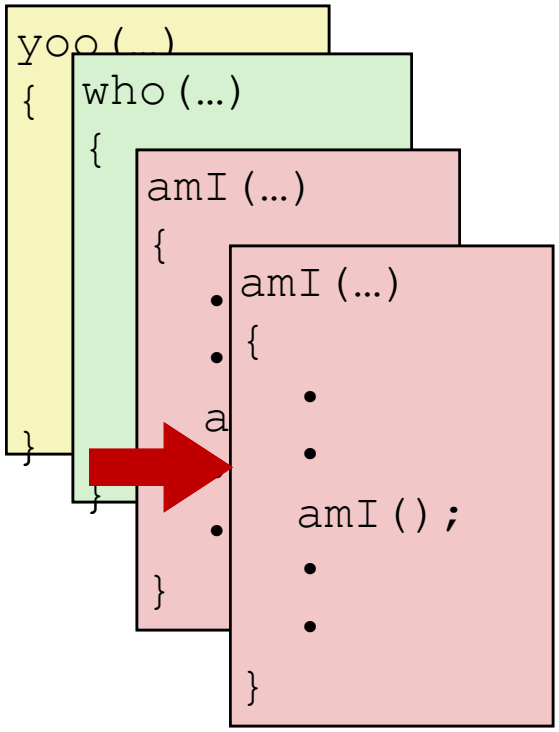
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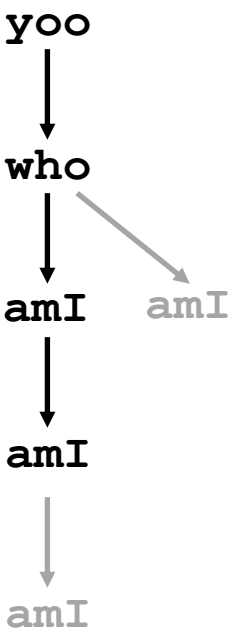
yoo
who
amI
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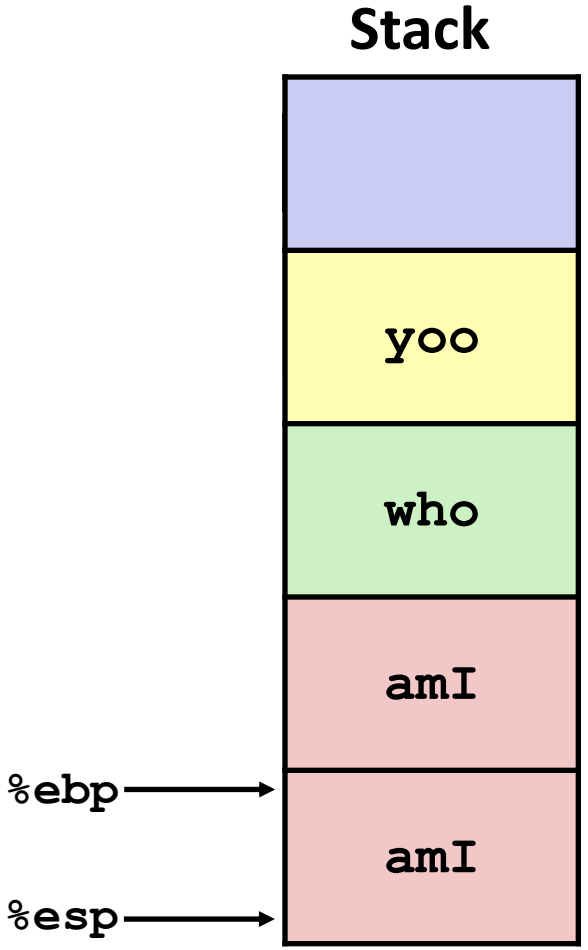
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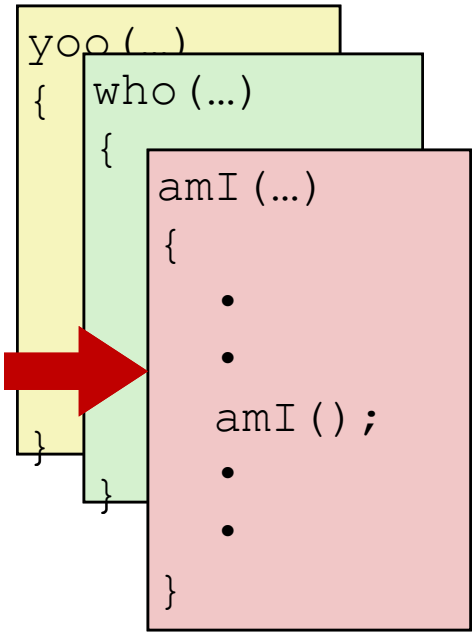
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amI();
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•
}
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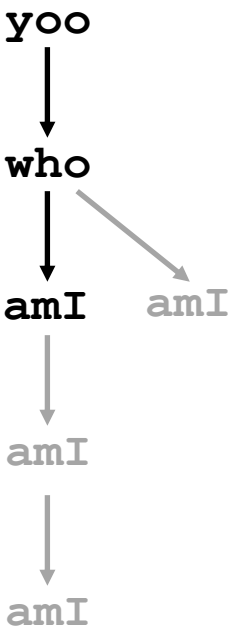
}


yoo
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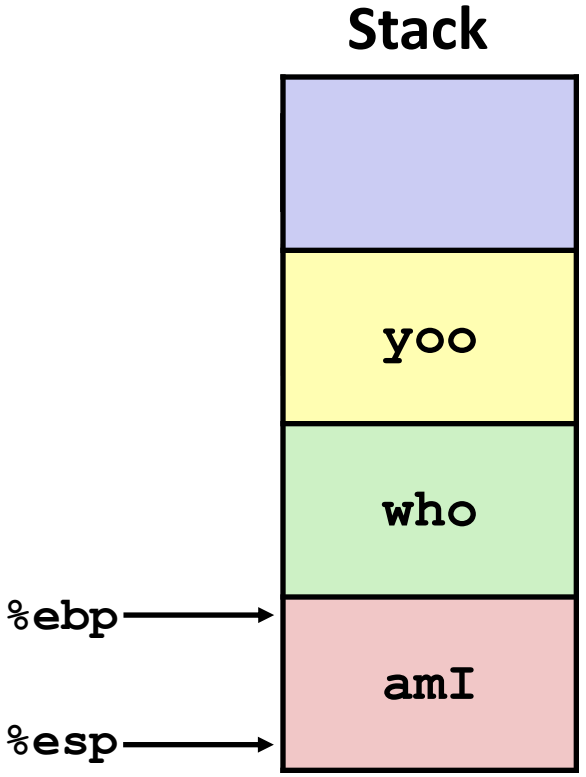
Stack
yoo
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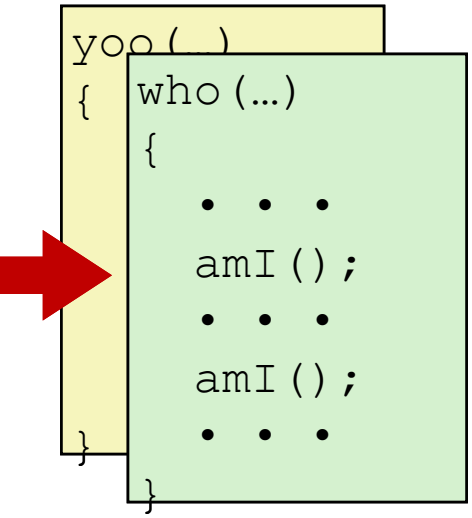
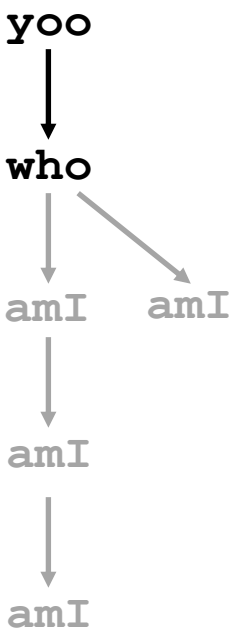
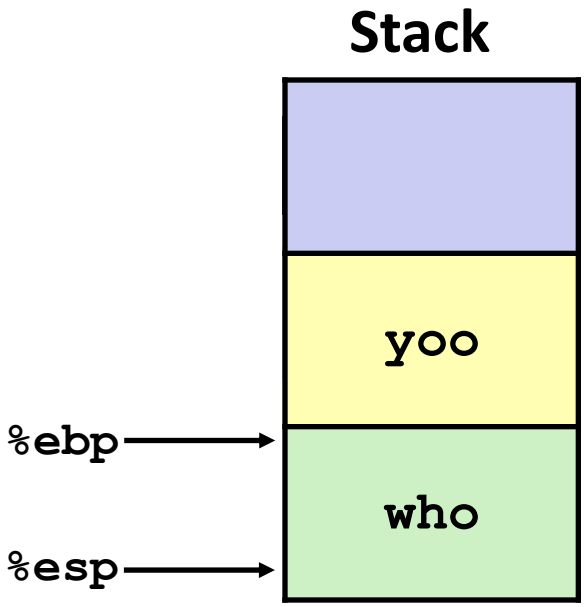
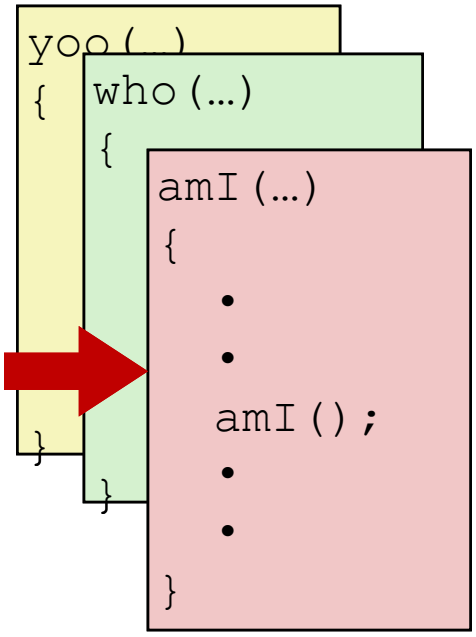
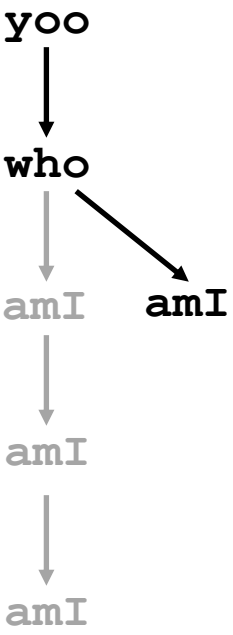


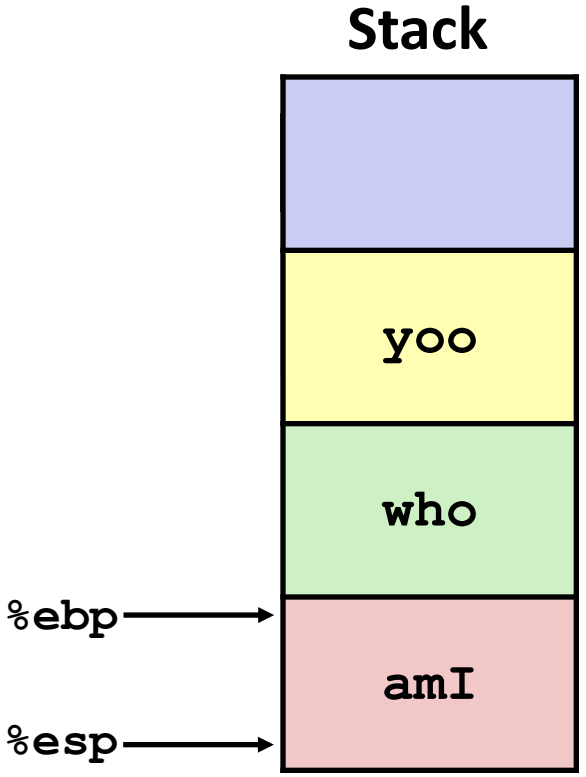




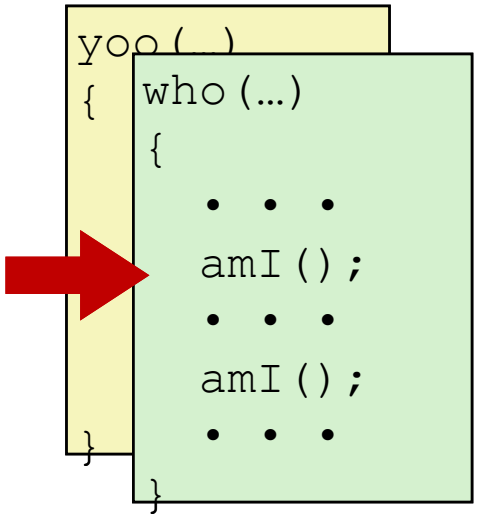
Example

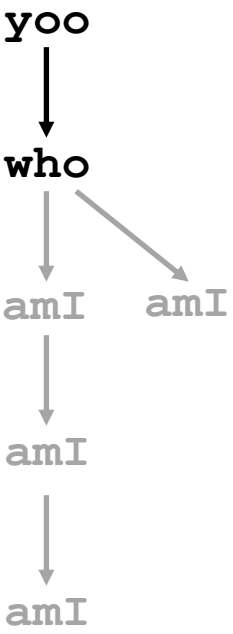

yoo(…)
{
who(…)
{
amI(…)
{
•
•
amI();
•
•
}
}
}

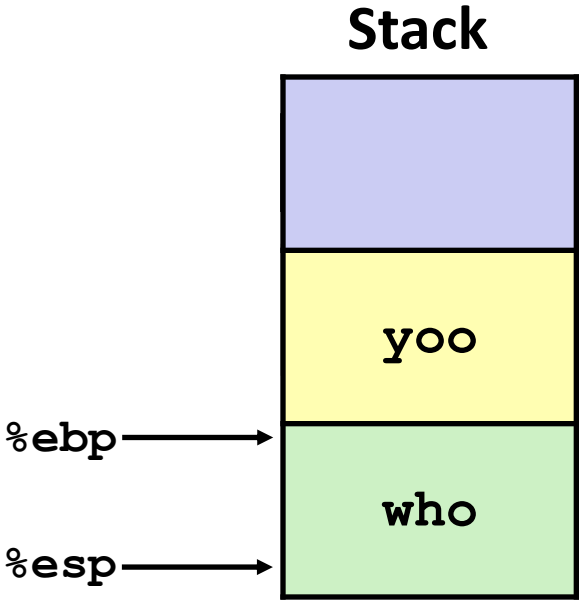

yoo
who
amI
amI
amI
amI


Stack
yoo
who
amI
%ebp
%esp

Example







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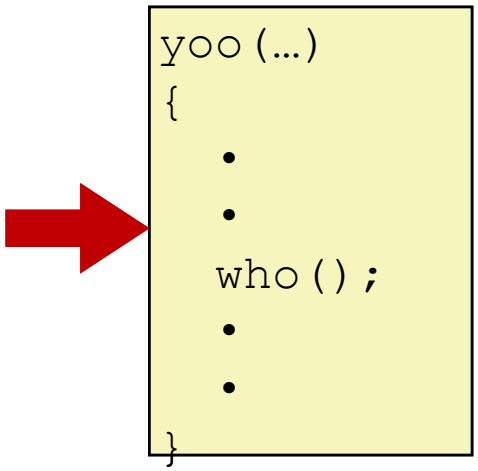
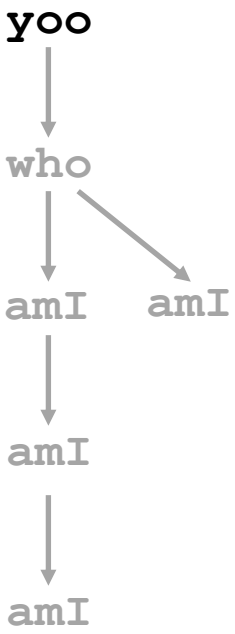
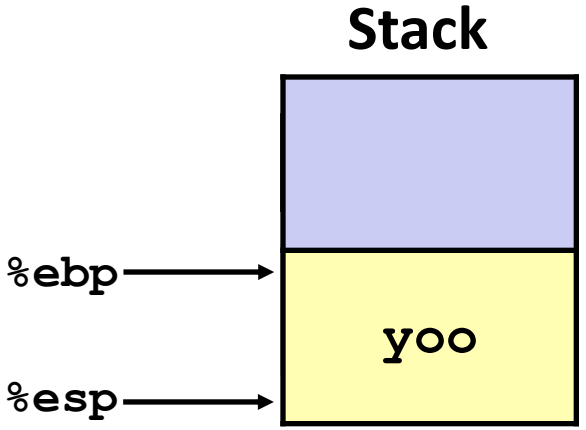


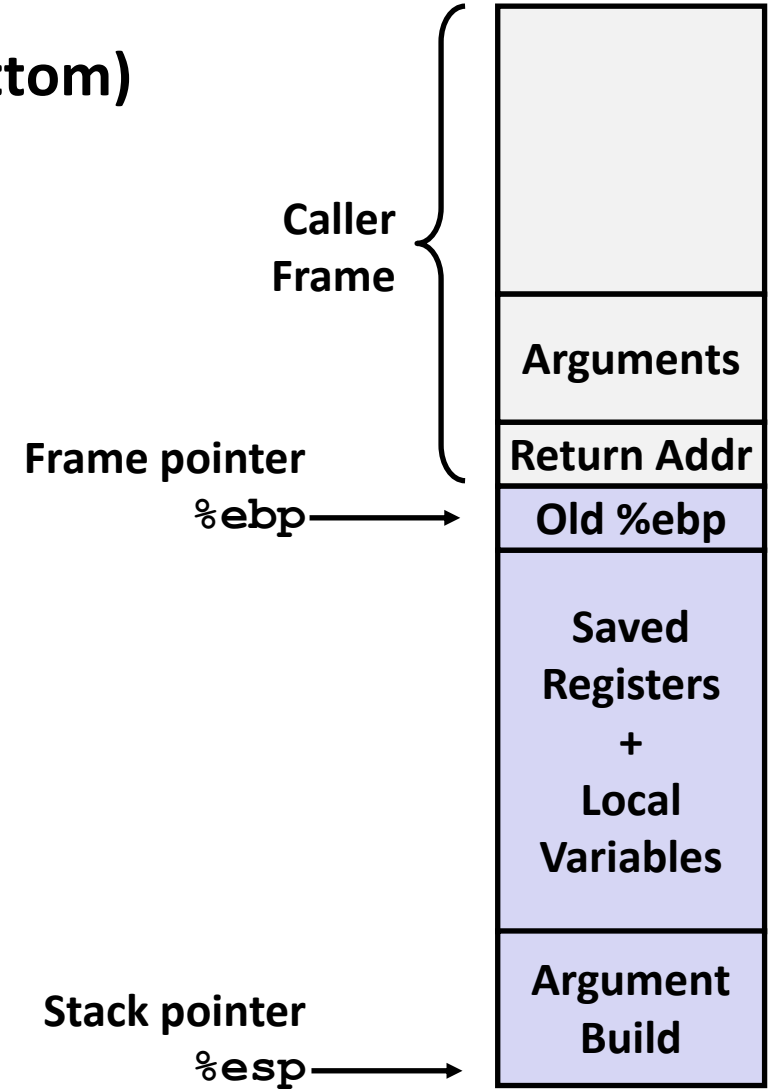




IA32/Linux Stack Frame

- **Current Stack Frame ("Top" to Bottom)**
  - "Argument build:"
    Parameters for function about to call
  - Local variables
    If can't keep in registers
  - Saved register context
  - Old frame pointer

- **Caller Stack Frame**
  - Return address
    - Pushed by `call` instruction
  - Arguments for this call



Revisiting swap

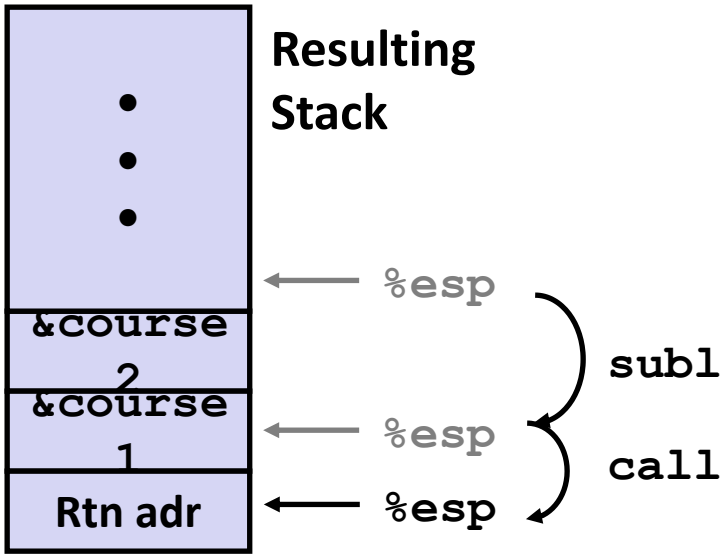
```
int course1 = 15213;
int course2 = 18243;

void call_swap() {
  swap(&course1, &course2);
}
```

Calling swap from call_swap

```
call_swap:
    • • •
    subl    $8, %esp
    movl    $course2, 4(%esp)
    movl    $course1, (%esp)
    call    swap
    • • •
```

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```



Revisiting swap

```
void swap(int *xp, int *yp)
{
  int t0 = *xp;
  int t1 = *yp;
  *xp = t1;
  *yp = t0;
}
```

```
swap:
    pushl %ebp
    movl  %esp, %ebp
    pushl %ebx

    movl  8(%ebp), %edx
    movl  12(%ebp), %ecx
    movl  (%edx), %ebx
    movl  (%ecx), %eax
    movl  %eax, (%edx)
    movl  %ebx, (%ecx)

    popl  %ebx
    popl  %ebp
    ret
```

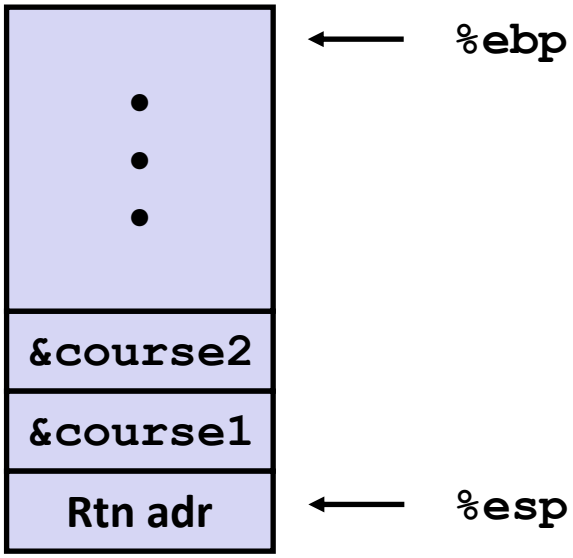
Set Up

Body

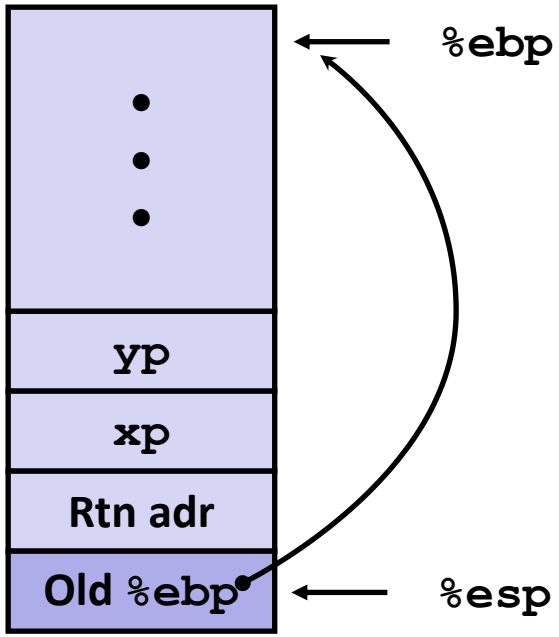
Finish

swap Setup #1

Entering Stack



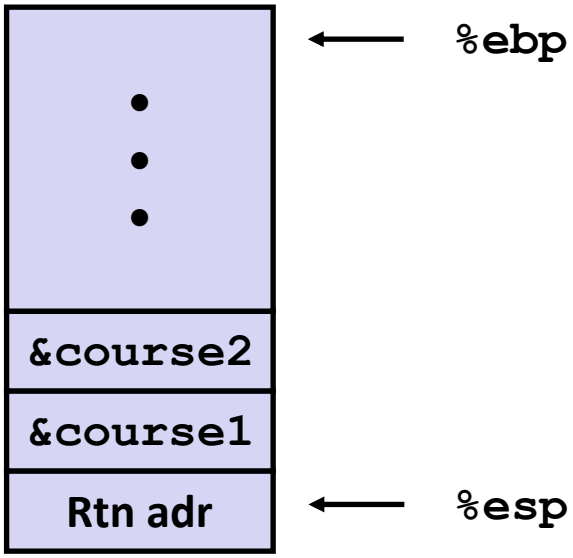
Resulting Stack



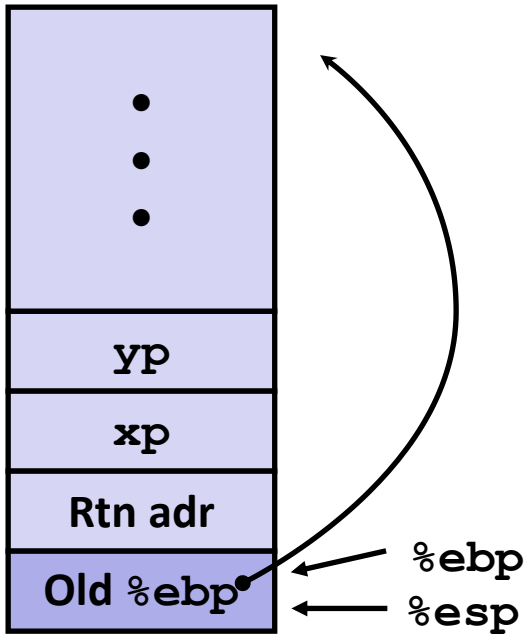

```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

swap Setup #2

Entering Stack



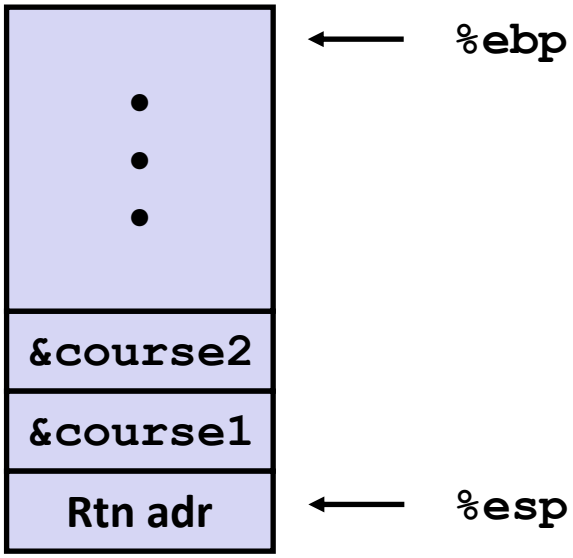
Resulting Stack



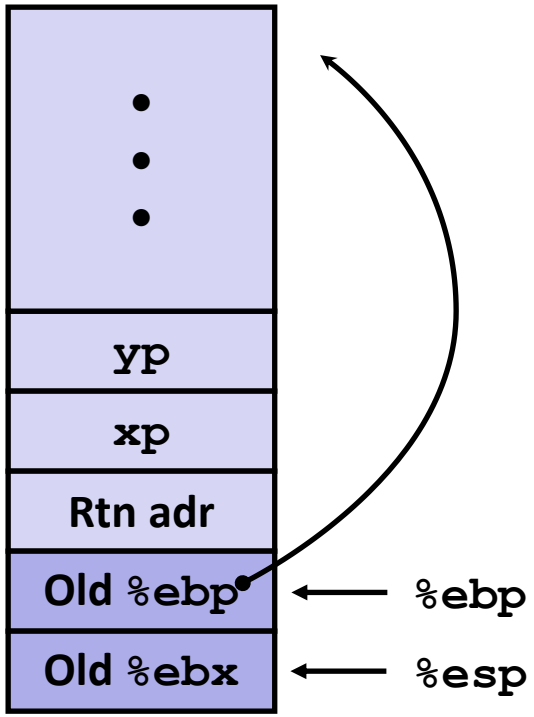

```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

swap Setup #3

Entering Stack

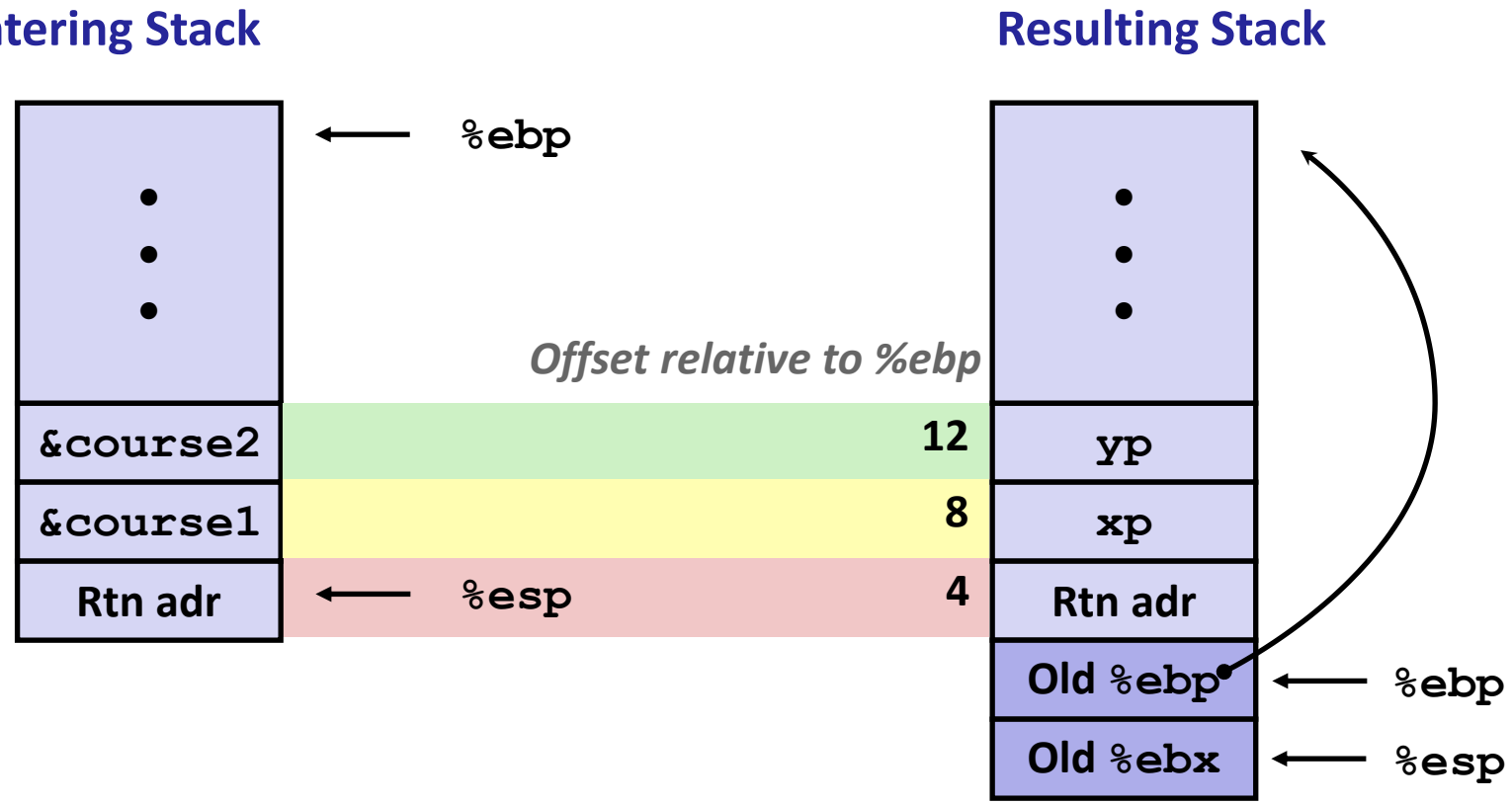


Resulting Stack




```
swap:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
```

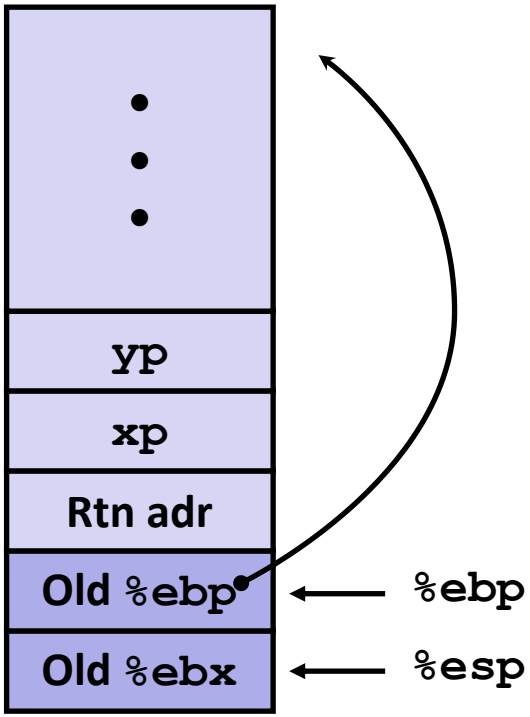
swap Body




```
movl 8(%ebp),%edx    # get xp
movl 12(%ebp),%ecx   # get yp
. . .
```

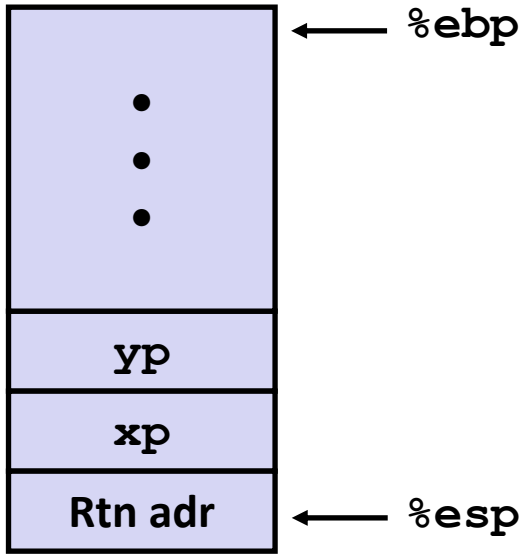
swap Finish

Stack Before Finish




```
popl    %ebx
popl    %ebp
```

Resulting Stack



- **Observation**
  - Saved and restored register **%ebx**
  - Not so for **%eax**, **%ecx**, **%edx**

Disassembled swap

```
08048384 <swap>:
 8048384:  55                   push   %ebp
 8048385:  89 e5                mov    %esp,%ebp
 8048387:  53                   push   %ebx
 8048388:  8b 55 08             mov    0x8(%ebp),%edx
 804838b:  8b 4d 0c             mov    0xc(%ebp),%ecx
 804838e:  8b 1a                mov    (%edx),%ebx
 8048390:  8b 01                mov    (%ecx),%eax
 8048392:  89 02                mov    %eax,(%edx)
 8048394:  89 19                mov    %ebx,(%ecx)
 8048396:  5b                   pop    %ebx
 8048397:  5d                   pop    %ebp
 8048398:  c3                   ret
```

Calling Code

```
 80483b4:  movl   $0x8049658,0x4(%esp) # Copy &course2
 80483bc:  movl   $0x8049654,(%esp)    # Copy &course1
 80483c3:  call   8048384 <swap>       # Call swap
 80483c8:  leave                       # Prepare to return
 80483c9:  ret                         # Return
```

Today

- Switch statements
- **IA 32 Procedures**
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers

Register Saving Conventions

- **When procedure `yoo` calls `who`:**
  - `yoo` is the ***caller***
  - `who` is the ***callee***
- **Can register be used for temporary storage?**

```
yoo:
    • • •
    movl $15213, %edx
    call who
    addl %edx, %eax
    • • •
    ret
```

```
who:
    • • •
    movl 8(%ebp), %edx
    addl $18243, %edx
    • • •
    ret
```

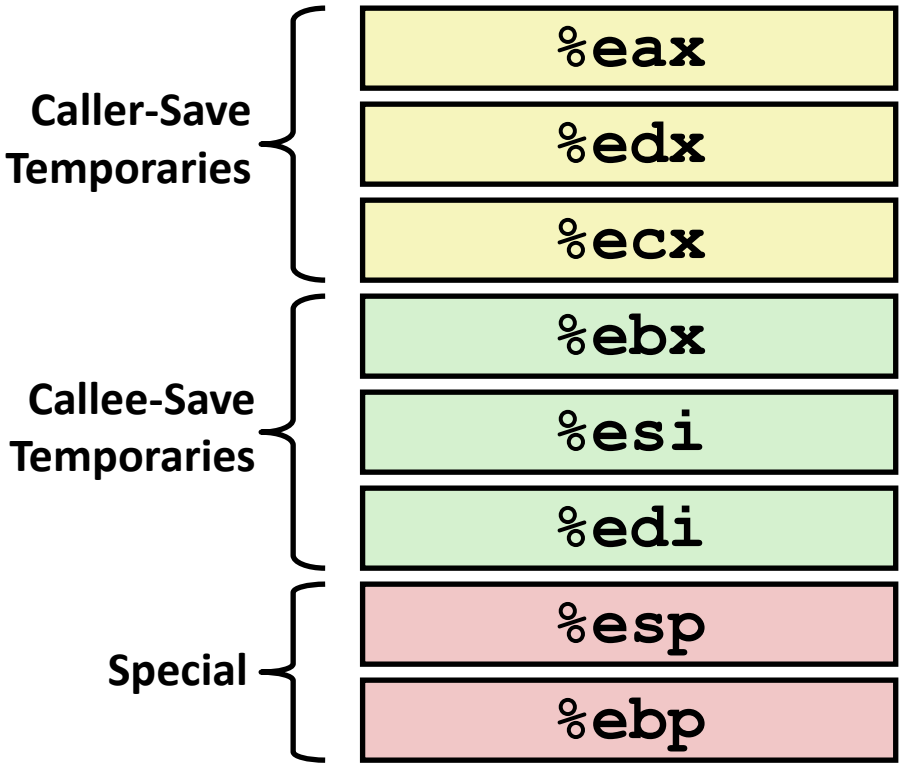
  - Contents of register `%edx` overwritten by `who`
  - This could be trouble → something should be done!
    - Need some coordination

Register Saving Conventions

- **When procedure `yoo` calls `who`:**
  - `yoo` is the ***caller***
  - `who` is the ***callee***
- **Can register be used for temporary storage?**
- **Conventions**
  - ***"Caller Save"***
    - Caller saves temporary values in its frame before the call
  - ***"Callee Save"***
    - Callee saves temporary values in its frame before using

IA32/Linux+Windows Register Usage

- **`%eax`, `%edx`, `%ecx`**
  - Caller saves prior to call if values are used later
- **`%eax`**
  - also used to return integer value
- **`%ebx`, `%esi`, `%edi`**
  - Callee saves if wants to use them
- **`%esp`, `%ebp`**
  - special form of callee save
  - Restored to original values upon exit from procedure



Today

- **IA 32 Procedures**
  - Stack Structure
  - Calling Conventions
  - Illustrations of Recursion & Pointers

Recursive Function

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

- **Registers**
  - `%eax`, `%edx` used without first saving
  - `%ebx` used, but saved at beginning & restored at end

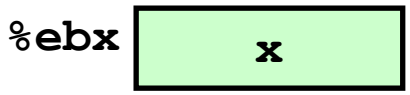
```
pcount_r:
    pushl %ebp
    movl%esp, %ebp
    pushl %ebx
    subl$4, %esp
    movl8(%ebp), %ebx
    movl$0, %eax
    testl %ebx, %ebx
    je  .L3
    movl%ebx, %eax
    shrl%eax
    movl%eax, (%esp)
    callpcount_r
    movl%ebx, %edx
    andl$1, %edx
    leal(%edx,%eax), %eax
.L3:
    addl$4, %esp
    popl%ebx
    popl%ebp
    ret
```

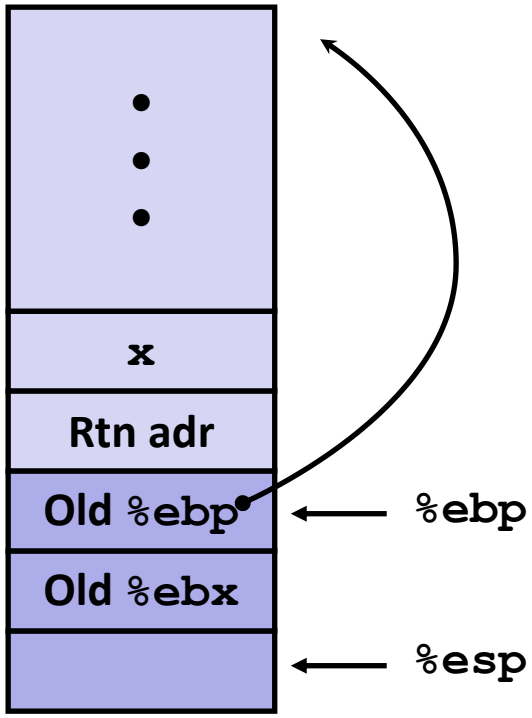
Recursive Call #1

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
pcount_r:
    pushl %ebp
    movl %esp, %ebp
    pushl %ebx
    subl $4, %esp
    movl 8(%ebp), %ebx
     • • •
```

- **Actions**
  - Save old value of **%ebx** on stack
  - Allocate space for argument to recursive call
  - Store x in **%ebx**



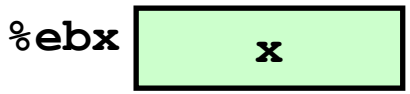


Recursive Call #2

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
    . . .
    movl   $0, %eax
    testl %ebx, %ebx
    je   .L3
    . . .
.L3:
    . . .
    ret
```

- **Actions**
  - If x == 0, return
    - with **%eax** set to 0

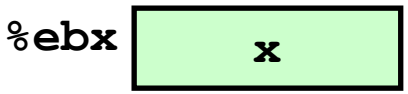


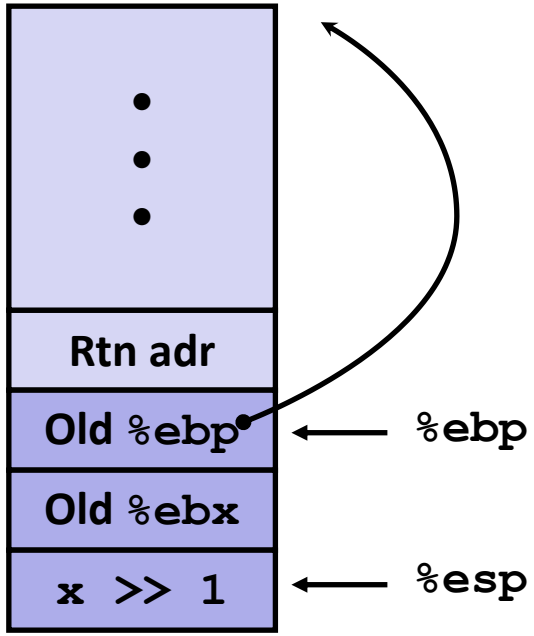
Recursive Call #3

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
  • • •
movl   %ebx, %eax
shrl   %eax
movl   %eax, (%esp)
call   pcount_r
  • • •
```

- **Actions**
  - Store x >> 1 on stack
  - Make recursive call
- **Effect**
  - %eax set to function result
  - %ebx still has value of x



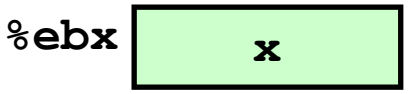


Recursive Call #4

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
  • • •
movl    %ebx, %edx
andl    $1, %edx
leal    (%edx,%eax), %eax
  • • •
```

- Assume
  - %eax holds value from recursive call
  - %ebx holds x
- Actions
  - Compute (x & 1) + computed value
- Effect
  - %eax set to function result

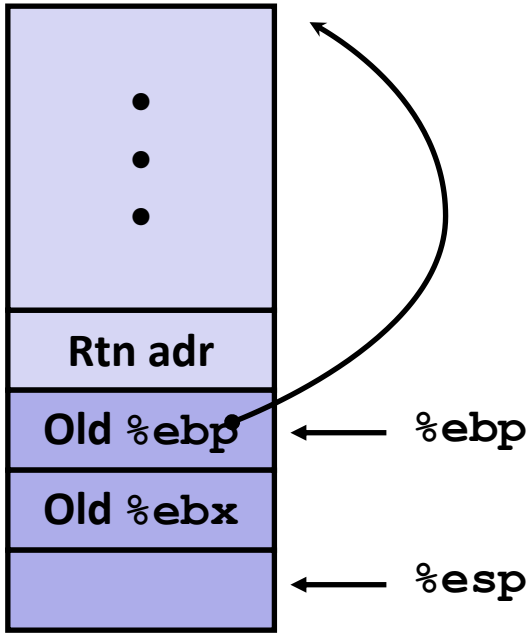


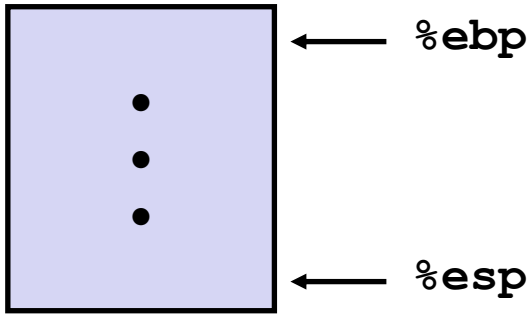
Recursive Call #5

```
/* Recursive popcount */
int pcount_r(unsigned x) {
  if (x == 0)
    return 0;
  else return
    (x & 1) + pcount_r(x >> 1);
}
```

```
    • • •
L3:
    addl$4, %esp
    popl%ebx
    popl%ebp
    ret
```

- **Actions**
  - Restore values of %ebx and %ebp
  - Restore %esp





%ebx

Old %ebx

Observations About Recursion

- **Handled Without Special Consideration**
  - Stack frames mean that each function call has private storage
    - Saved registers & local variables
    - Saved return pointer
  - Register saving conventions prevent one function call from corrupting another's data
  - Stack discipline follows call / return pattern
    - If P calls Q, then Q returns before P
    - Last-In, First-Out
- **Also works for mutual recursion**
  - P calls Q; Q calls P

Pointer Code

Generating Pointer

```
/* Compute x + 3 */
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

Referencing Pointer

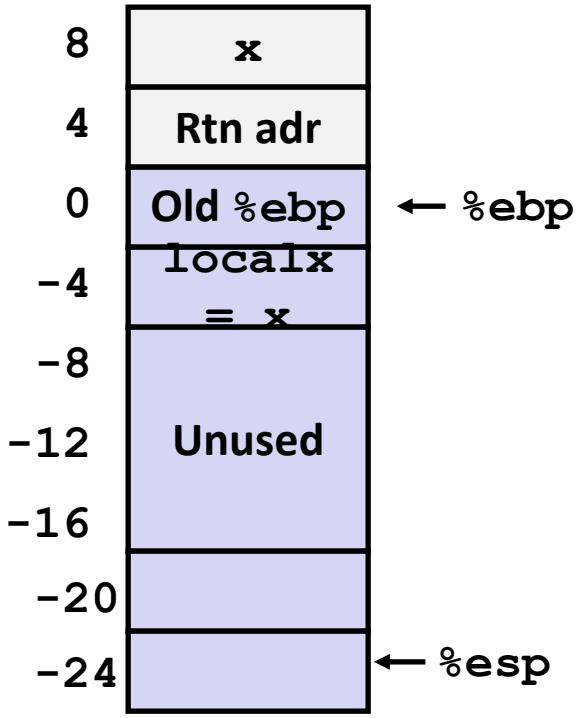
```
/* Increment value by k */
void incrk(int *ip, int k) {
  *ip += k;
}
```

- add3 creates pointer and passes it to `incrk`

Creating and Initializing Local Variable

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- Variable localx must be stored on stack
  - Because: Need to create pointer to it
- Compute pointer as -4(%ebp)



First part of add3

```
add3:
   pushl %ebp
   movl  %esp, %ebp
   subl  $24, %esp     # Alloc. 24 bytes
   movl  8(%ebp), %eax
   movl  %eax, -4(%ebp) # Set localx to x
```

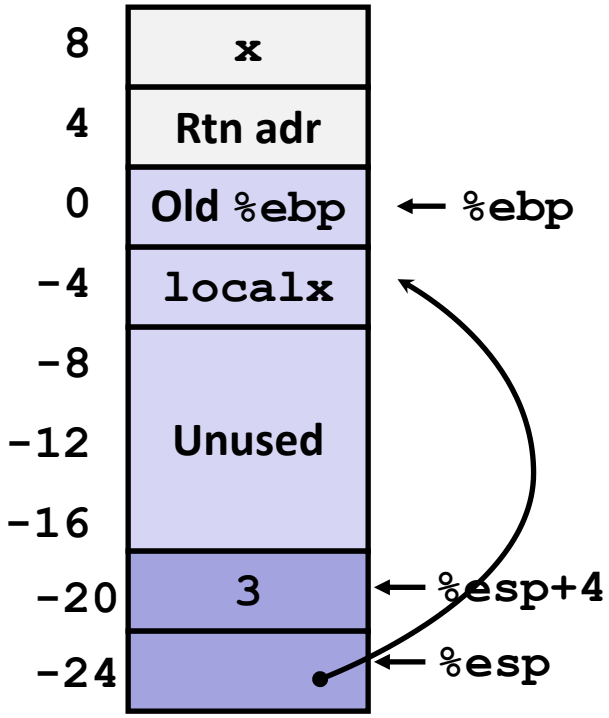
Creating Pointer as Argument

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- Use leal instruction to compute address of localx

Middle part of add3

```
  movl $3, 4(%esp)   # 2nd arg = 3
  leal -4(%ebp), %eax# &localx
  movl %eax, (%esp)  # 1st arg = &localx
  call incrk
```



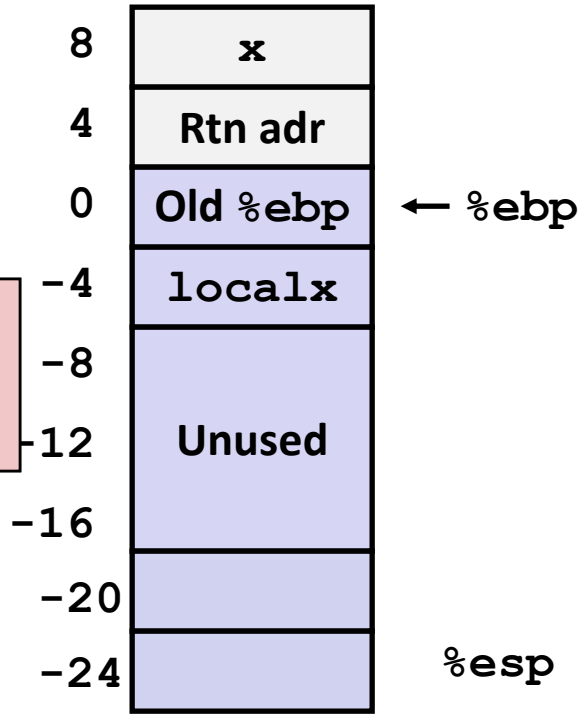
Retrieving local variable

```
int add3(int x) {
  int localx = x;
  incrk(&localx, 3);
  return localx;
}
```

- Retrieve localx from stack as return value

Final part of add3

```
movl -4(%ebp), %eax # Return val= localx
leave
ret
```



IA 32 Procedure Summary

- **Important Points**
  - Stack is the right data structure for procedure call / return
    - If P calls Q, then Q returns before P
- **Recursion (& mutual recursion) handled by normal calling conventions**
  - Can safely store values in local stack frame and in callee-saved registers
  - Put function arguments at top of stack
  - Result return in **`%eax`**
- **`Pointers are addresses of values`**
  - **`On stack or global`**