Social Network Analysis - Optimization, Scaling, and Final Evaluation

Algorithms and Data Structures

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1 Overview

1.1 Problem Statements

A social network is incredibly useful for maintaining friendships, reaching out, and following your favorite celebrity. In particular, Twitter is one social media that allows people to share and exchange ideas freely and easily. Twitter offers many features that allow interactions between end-users and followers, and one of them would be Retweet. Retweet is an important concept as it determines how trends form and how tweets go viral. Distributing content is as simple as a click of a button, and social network analysis can help us understand the feature's behavior behind the scenes at a larger scale.

1.2 Solution

Social network analysis is the process of investigating social structures using networks and graph theories. It combines various techniques for analyzing the structure of social networks and theories that aim to explain the underlying dynamics and patterns observed in these structures. It is an inherently interdisciplinary field originally from social psychology, statistics, and graph theory. Today, we will investigate, analyze, and learn more about the Social Network's algorithm to find influential users in a social network and how influential users grow daily.

2 Optimization of Data Structures

The shortest path[1] is the problem of finding a path between two articles in a graph such that the sum of the weights of its constituent edges is minimized. For social network analysis, we solved the shortest path problem, which was used to find the recommended friends and the shortest path between two cliques via Bread First Search (BFS).

2.1 Recommended friends

Before optimization, we must understand the implementation of the recommended friends feature. The feature has 3 steps:

- Step 1: Find all the combinations of all the node's neighbors
- Step 2: If there is no connection between two neighbor nodes, we would recommend them and increase further if there are more intermediate node
- Step 3: Prioritize the top 10 pairs of nodes that have the highest intermediate nodes (or bypass connections instead of direct connections)

Preprint. Under review.

```
# recommendedFriends will recommend friends who shared

# recommendedFriends will recommend friends who shared

# If there is no connection, we will recommend them

def recommendedFriends(self) -> list[tuple]:

recommendedFriends = defaultdict(int)

for node, _ in self.fbGraph.nodes(data=True):

# Create a combination of 2 with all the neighbor nodes that the node has

# and determine if there is a path between them. If not,

# add them

for firstNeighborFriend, secondNeighborFriend in combinations(self.fbGraph.neighbors(node), 2):

| if not self.pathExistBFS(firstNeighborFriend, secondNeighborFriend):

# Identify the top 10 pairs of users

sortedRecommendedFriends = sorted(recommendedFriends.values())

top10Pairs = [pair for pair, count in recommendedFriends.items() if count > sortedRecommendedFriends[-10 return top10Pairs

# pathExistBFS will use breath first search

# https://www.geeksforgeeks.org/breadth-first-search-or-bfs-for-a-graph/#
```

Figure 2: Recommend Friends in Python

The time complexity of the algorithm is expected to be $O(V \times C(V, 2) \times (V + E))$ in total and O(V + E) for BFS, where V is the number of vertices and E is the number of edges. For unsorted nodes with random degrees, the edges are small enough to be negligible. However, with sorting nodes based on their degree (the number of edges connecting to the node), the edge has increased from 20 to 3824 (a total of 190.2% increasing.

```
Number of nodes 100
Number of edges 3824
Number of nodes ['u698', 'u7034', 'u7473', 'u3233', 'u3468', 'u6664', 'u3683', 'u5395', 'u607', 'u3248', 'u1262', 'u
1011', 'u2170', 'u2925', 'u3283', 'u376', 'u1687', 'u3219', 'u186', 'u1270', 'u7302', 'u1695', 'u1481', 'u255', 'u26
4', 'u115', 'u157', 'u644', 'u1617', 'u527', 'u724', 'u1327', 'u1469', 'u1339', 'u8655', 'u8819', 'u218', 'u4819', 'u741', 'u1678', 'u1588', 'u9745', 'u1975', 'u855', 'u817', 'u1414', 'u639', 'u6441', 'u1080', 'u741', 'u6701', 'u167', 'u37
```

Figure 3: Sorted nodes with the highest degree (highest edge)

As a result, we encounter the scaling issue: running the recommended friends feature takes from 0.076 to 18.63s (a total of 24457% increasing) with sorting nodes based on their degree. Fortunately, Python offers a Profiler [2] (including CPU Profile and Memory Profiler) that helps us identify the feature's bottleneck in the execution time. A total of 4400152 calls, and each takes 20.857s.

```
Ordered by: cumulative time

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```

Figure 4: Profiler for recommended friends

To optimize the algorithm, instead of finding whether there aren't any connections between the node's neighbors, we will traverse the path during the initialization, recording the path between two nodes via cache if there is a connection. Afterward, we will only use the recorded path as a baseline for the recommended friends feature by referring the friends via the current node if there are no connections between them.

```
# recommendedFriends will recommend friends who shared

# but hasn't been friends yet by creating a combination between the current user's neighbor

# If there is no connection, we will recommend them

def recommendedFriends(self) -> list[tuple]:

recommendedFriends = defaultdict(int)

for node, _in self.fbGraph.nodes(data=True):

# Create a combination of 2 with all the neighbor nodes that the node has

# and determine if there is a path between them. If not,

# add them

for firstNeighborFriend, secondNeighborFriend in combinations(self.fbGraph.neighbors(node), 2):

if (firstNeighborFriend, secondNeighborFriend) not in self.traversePath:

# if not self.pathExistBFS(firstNeighborFriend, secondNeighborFriend):

| recommendedFriends[(firstNeighborFriend, secondNeighborFriend)] += 1

# Identify the top 10 pairs of users

sortedRecommendedFriends = sorted(recommendedFriends.values())

top10Pairs = [pair for pair, count in recommendedFriends.items() if count > sortedRecommendedFriends[-16]

return top10Pairs
```

Figure 4: Optimization for recommended friends

As a result of the optimization, the time complexity has been improved to $O(V^2x(V+E))$; however, in exchange for the space complexity of O(V * P) where V is the number of vertices and P is the maximum path a vertice can have, which leads to the following graph:

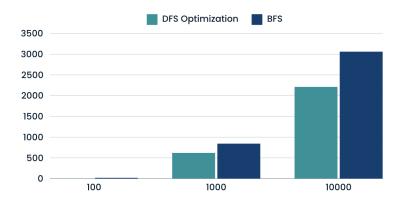


Figure 5: BFS Optimization for sorted degree Execution Time

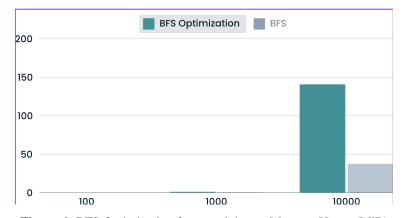


Figure 6: BFS Optimization for sorted degree Memory Usage (MiB)

However, for unsorted nodes with random degrees, the edges are small enough to be negligible, which results in similar execution time for both optimized BFS with cache and original BFS.



Figure 7: BFS Optimization for unsorted degree Execution Time (second)

2.2 Connect communities

Before optimization, we must understand how to connect community features. The feature has 3 steps:

- Step 7: Find the cliques[3] of the social network
- Step 8: Find the bottleneck user which serves as a "bridge-spanning role" in the network using betweenness centrality [8]
- Step 9: Calculate the minimum distance between the two bottleneck users to see if there are any connections.

```
# connectCommunities will find the two largest communities based on the cliques and
# determine the bottleneck user based on the number of friends or followers
# they have. Afterwards, if there is a path or similar interest between them, we can
# recommend the bottle neck's user to other followers/friends
# to increase the reach
def connectCommunities(self):
    cliques = sorted(nx.find_cliques(self.fbGraph), key=lambda x: len(x))
    largestClique = self.fbGraph.subgraph(set(cliques[-1])).copy()
    secondLargestClique = self.fbGraph.subgraph(set(cliques[-2])).copy()
    largestCliqueBetweenessCentrality = nx.betweenness_centrality(largestClique)
    secondLargestCliqueBetweenessCentrality = nx.betweenness_centrality(secondLargestClique)
    largestCliqueBottleneckUser = max(largestCliqueBetweenessCentrality, key=largestCliqueBetweenessCentrality
    secondLargestCliqueBottleneckUser = max(secondLargestCliqueBetweenessCentrality, key=secondLargestCliqueBottleneckUser, second fr self.pathExistBFS(largestCliqueBottleneckUser, second fr self.pathExistBFS(largestC
```

Figure 8: Connect communities in Python

Since Facebook social network is an indirect graph, therefore, it would work best with BFS instead of other algorithms such as Floyd–Warshall algorithm, Bellman-Ford algorithm, or even Dijkstra's algorithm:

	BFS	Dijkstra's	Bellman-Ford	Floyd Warshall
Time Complexity	O(V + E)	$O(V + E) \log V$	O(V + E)	$O(V^3)$
Recommended Graph Size	Large	Large/Medium	Medium/ Small	Small
All pairs shortest path	Only unweighted graphs	Ok	Bad	Yes
Negative cycles	No	No	Yes	Yes
Shortest path with weighted edges	Bad	Best algorithm	Works	Bad in general
Shortest path with unweighted edges	Best algorithm	Ok	Bad	Bad in general

Figure 9: The shortest path algorithm's analysis

The time complexity of the algorithm is expected to be $O(3^{\frac{v}{3}} + v + e)$ [4] in total and O(v + e) for BFS.

qWe can use the same aforementioned algorithm to cache all the traverse paths of all the nodes inside the social network and determine if there is a connection between two nodes, and we can also know the minimum distance between the two.

```
def connectCommunities(self)-> int:
    cliques = sorted(nx.find_cliques(self.fbGraph), key=lambda x: len(x))
    largestClique = self.fbGraph.subgraph(set(cliques[-1])).copy()
    secondLargestClique = self.fbGraph.subgraph(set(cliques[-2])).copy()
    largestCliqueBetweenessCentrality = nx.betweenness_centrality(largestClique)
    secondLargestCliqueBetweenessCentrality = nx.betweenness_centrality(secondLargestClique)
    secondLargestCliqueBetweenessCentrality = nx.betweenness_centrality(secondLargestClique)
    largestCliqueBottleneckUser = max(largestCliqueBetweenessCentrality, key=largestCliqueBetweenessCentrality
    secondLargestCliqueBottleneckUser)
    #IstanceBetweenivoUsers = self.getDistanceWithCurrentNode(largestCliqueBottleneckUser)
    #IstanceBetweenivoUsers = self.getDistanceWithCurrentNode(largestCliqueBottleneckUser, secondLargestCliqueBottleneckUser)
    ##IstanceBetweenivoUsers = self.getDistanceWithCurrentNode(largestCliqueBottleneckUser, secondLargestCliqueBottleneckUser, secondLargestCliqueB
```

Figure 10: Optimization for connecting communities

As the cache is using BFS to traverse the path behind the scenes, we will increase the execution time when compared to the original BFS from 0.845s to 4.6s(a total of 444% increasing) However, with Dijkstra O(V + E) log V, there will be slower execution time but negligible for the addition log V from 0.845s to 0.877s (a total of 3.6% increasing. The result has been shown below:

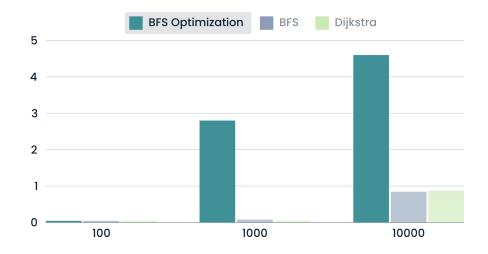


Figure 11: Execution time for Optimized BFS, BFS and Dijkstra

Therefore, we can say BFS with the unweighted graph works best for our current use case when determining the connection between the two largest communities.

3 Advanced Testing and Validation

To maintain backward compatibility and documentation for future use, we will develop a set of test cases. There are 5 features complex enough to develop a test case for it:

• Shortest path with BFS:

Figure 12: Test finding path between two nodes with BFS

• Shortest path with BFS Optimization:

Figure 13: Test recording all traversal paths with BFS Optimization

• Find the largest community:

```
def test_FindLargestCommunity(self):
     dataset = downloadDataset()
      with open(dataset[1], 'rb') as f:
           fbGraph = pickle.load(f)
      nodes = [node for node in list(fbGraph.nodes())[:1000]]
      fbSubGraph = fbGraph.subgraph(nodes)
      sn = SocialNetwork(fbSubGraph)
      largestClique = set(sorted(nx.find\_cliques(fbSubGraph), \; key=lambda \; x: \; len(x))[-1])
      expectedFacebookLargestClique = fbSubGraph.subgraph(largestClique).copy()
       for node in list(expectedFacebookLargestClique.nodes()):
                   {\tt expectedFacebookLargestClique.add\_nodes\_from(fbSubGraph.neighbors(node))}
                  \label{lem:expected} \textbf{expectedFacebookLargestClique.add\_edges\_from(zip([\underline{node]}*len(list(fbSubGraph.neighbors(node))), fbSuparties from the properties of the properties o
      self.assertIsNotNone(expectedFacebookLargestClique)
     self.assertEqual(expectedFacebookLargestClique.nodes(), sn.findLargestCommunities().nodes())
def test_FindLargestCommunityWithEmptyGraph(self):
      fbGraph = nx.Graph()
      sn = SocialNetwork(fbGraph)
       self.assertIsNone(sn.findLargestCommunities())
```

Figure 14: Test for finding largest community

• Find important people:

```
def test_FindImportantPeople(self):
 dataset = downloadDataset()
 with open(dataset[1], 'rb') as f:
  fbGraph = pickle.load(f)
 nodes = [node for node in list(fbGraph.nodes())[:1000]]
 fbSubGraph = fbGraph.subgraph(nodes)
 sn = SocialNetwork(fbSubGraph)
 importantPeople = sn.findImportantPeople()
 degCent = nx.degree_centrality(fbSubGraph)
 degCent = {k: v for k, v in sorted(degCent.items(), key=lambda item: item[1], reverse=True)}
 maxDegCent = max(list(degCent.values()))
 expectedImportantPeople = [(n, dc) for n, dc in degCent.items() if dc == maxDegCent]
 self.assertEqual([expectedImportantPeople], importantPeople)
def test_FindImportantPeopleWithEmptyGraph(self):
 fbGraph = nx.Graph()
 sn = SocialNetwork(fbGraph)
 importantPeople = sn.findImportantPeople()
 self.assertEqual([], importantPeople)
```

Figure 15: Test for finding important users

• Connect communities:

```
dataset = downloadDataset()
      with open(dataset[1], 'rb') as f:
          fbGraph = pickle.load(f)
      nodes = [node for node in list(fbGraph.nodes())[:1000]]
      fbSubGraph = fbGraph.subgraph(nodes)
      sn = SocialNetwork(fbSubGraph)
     cliques = sorted(nx.find_cliques(fbSubGraph), key=lambda x: len(x))
     largestClique = fbSubGraph.subgraph(set(cliques[-1])).copy()
     secondLargestClique = fbSubGraph.subgraph(set(cliques[-2])).copy()
     largest Clique Between ess Centrality = nx.between ness\_centrality (largest Clique)
     second Largest Clique Between ess Centrality = nx.between ness\_centrality (second Largest Clique) \\
      largest Clique Bottleneck User = \max(largest Clique Betweeness Centrality, \ key = largest Clique Betweeness Centrality) + largest Clique Betweeness Centrality + largest Clique Betweeness + largest + larg
     secondLargestCliqueBottleneckUser = max(secondLargestCliqueBetweenessCentrality, key=secondLargestClique
     expectedDistanceBetweenTwoUsers = sn.getDistanceWithCurrentNode(largestCliqueBottleneckUser, secondLarge
      self.assertEqual(expectedDistanceBetweenTwoUsers, sn.connectCommunities())
def test_ConnectCommunitiesWithEmptyGraph(self):
      fbGraph = nx.Graph()
      sn = SocialNetwork(fbGraph)
      self.assertEqual(-1, sn.connectCommunities())
```

Figure 16: Test for connecting communities and the shortest distance

• Find recommended friends:

```
def test_FindRecommendedFriends(self):
                                                                                                                                                                          fbSubGraph
                                                                                                                                                                                                                                                     AB 世 台
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
        fbGraph = pickle.load(f)
    nodes = [node for node in list(fbGraph.nodes())[:1000]]
     fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    recommendedFriends = sn.recommendedFriends()
    expectedRecommendedFriends = set()
    expectedRecommendedFriends = defaultdict(int)
     for node, _ in fbSubGraph.nodes(data=True):
                 for firstNeighborFriend, secondNeighborFriend in combinations(fbSubGraph.neighbors(node), 2):
                            if not fbSubGraph.has_edge(firstNeighborFriend, secondNeighborFriend):
                                         expectedRecommendedFriends[(firstNeighborFriend, secondNeighborFriend)] += 1
    sortedExpectedRecommendedFriends = sorted(expectedRecommendedFriends.values())
    {\tt expectedRecommendedFriends} = {\tt [pair \it for \it pair, \it count \it in \it expectedRecommendedFriends.items() \it if \it count \it > \it sorted \it expected \it Pair \it expected \it Pair \it expected \it Pair \it expected \it expected
    self.assertGreater(len(recommendedFriends), 0)
     self.assertEqual(expectedRecommendedFriends, recommendedFriends)
def test_FindRecommendedFriendsWithEmptyGraph(self):
    fbGraph = nx.Graph()
     sn = SocialNetwork(fbGraph)
     self.assertEqual([], sn.recommendedFriends())
```

Figure 17: Test for finding recommend friends

When running the test in parallel, we will have the following result:

```
unitTests = [
    'test_FindRecommendedFriends',
    'test_ShortestPathWithEmptyGraph',
    'test_FindImportantPeople',
    'test_GonnectCommunities',
    'test_FindLargestCommunity',
    'test_ShortestPathBFS',
    'test_TraversePathWillAllNodes'
]

with ThreadPoolExecutor() as executor:
    # Execute only the specified tests in parallel
    futures = [executor.submit(runTests, TestSocialNetworkAnalysis, test) for test in unitTests]
    for future in futures:
    | future.result()
```

Figure 18: Run unit test in parallel

Figure 19: Unit test result

As a result of these tests, if there are any features being optimized even more or any set of bug fixes/improvements, the test will provide a set of bulletproof when releasing a new version. Moreover, it also unit tests edge cases (e.g., paths with empty nodes or empty edges) to provide more robust functionalities.

However, we also need to perform a corresponding stress test [5] (a performance test in the future to give a user transparency of how the algorithm uses their allocation resource) in order to identify the breaking point of a system and improve it under extreme load conditions.

• Find the largest community:

```
def test_StressTestingFrom1000To20000NodesForFindingLargestCommunities(self):
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
    fbGraph = pickle.load(f)
    for numberOfNodes in range(1000, 20000, 1000):
    nodes = [node for node in list(fbGraph.nodes())[:numberOfNodes]]
    fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    sn.findLargestCommunities()
```

Figure 20: Stress Test for finding largest community

• Find important people:

```
def test_StressTestingFrom1000To20000NodesForFindingImportantPeople(self):
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
    fbGraph = pickle.load(f)
    for numberOfNodes in range(1000, 20000, 1000):
    nodes = [node for node in list(fbGraph.nodes())[:numberOfNodes]]
    fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    sn.findImportantPeople()
```

Figure 21: Stress test for finding important users

• Connect communities:

```
def test_StressTestingFrom1000To20000NodesForConnectingComunities(self):
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
    fbGraph = pickle.load(f)
    for numberOfNodes in range(1000, 20000, 1000):
    nodes = [node for node in list(fbGraph.nodes())[:numberOfNodes]]
    fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    sn.connectCommunities()
```

Figure 22: Stress test for connecting communities

• Find recommended friends:

```
def test_StressTestingFrom1000To20000NodesForRecommendedFriendsWithUnsortedData(self):
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
    fbGraph = pickle.load(f)
    for numberOfNodes in range(1000, 20000, 1000):
    nodes = [node for node in list(fbGraph.nodes())[:numberOfNodes]]
    fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    sn.recommendedFriends()

def test StressTestingFrom1000To20000NodesForRecommendedFriendsWithSortedDegreeData(self):
    dataset = downloadDataset()
    with open(dataset[1], 'rb') as f:
    fbGraph = pickle.load(f)
    for numberOfNodes in range(1000, 20000, 1000):
    nodes = [node for node in list(fbGraph.nodes())[:numberOfNodes]]
    fbSubGraph = fbGraph.subgraph(nodes)
    sn = SocialNetwork(fbSubGraph)
    sn.recommendedFriends()
```

Figure 23: Stress test for finding recommended friends

When running the test in parallel, we will have the following result:

```
VERMINAL

187.15. 'u4289', 'u1387', 'u8625', 'u3468', 'u19261', 'u1719', 'u1481' 'u6552', 'u18813', 'u5122', 'u7459'

98', 'u1244', 'u255', 'u2886', 'u4288', 'u7377', 'u8692', 'u3484', 'u741', 'u1687', 'u1899', 'u6441', 'u5382', 'u5395', 'u5377', 'u775', 'u187', 'u1899', 'u6441', 'u5382', 'u5395', 'u5377', 'u3479', 'u4879', 'u4877', 'u5787', 'u7538', 'u5188', 'u5488', 'u4847', 'u5489', 'u4871', 'u4879', 'u4879', 'u4879', 'u5479', 'u5479',
```

Figure 24: Stress testing result

4 Final Evaluation and Performance Analysis

For the optimization of both shortest path problems with recommended friends and connect communities, we have noticed a increase in execution time (from 0.076 to 18.63s - a total of 24457% increasing). However, increasing memory for storing all the traversal paths beforehand (from 37.5 MiB to 140 MiB - a total of 3.75% increasing)



Figure 25: Reduction in execution time for Optimized BFS for sorted data (second)

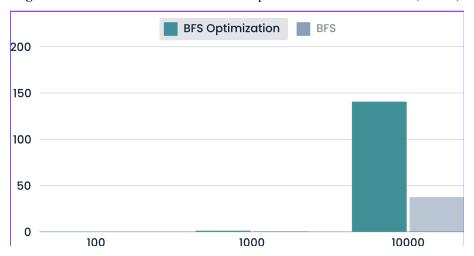


Figure 26: Increasing in memory usage for Optimized BFS for sorted data (MiB)

However, even though we have optimized the data structures, there are some areas we can improve on:

- The dataset from social networks is static; therefore, we would need to find a way to always append the new traversal path when there is a new user joining it.
- The solution only optimizes the features that include traversing to all the paths or share similar functionalities. However, for features that only need to determine a connection between two users, such as connect communities, it will have the drawback of over-engineering the solutions.
- The optimizing solutions only work towards an unweighted graph. However, if there are any
 weights or different data structures, the optimization and the current implementation won't
 work.

• In a constrained memory environment, the space complexity of O(V * P), where P is the number of paths a node can have, will become a problem compared to the O(V) of the original solution.

Therefore, to avoid the aforementioned problems, we can optimize it by:

- We will store all the traversal paths of all the nodes during the initial utilization. When there is a user being added, or there are new connections between two users as a result of recommended friends features via a Pub/Sub system [6] (e.g. Kafka, RabbitMQ, etc.), we will add the node and all the traversal path of that node to the cache.
- Instead of adding every user and its traversal path to the cache, we can prioritize different characteristics of a node to determine if the current node is needed or node (e.g if the user is not active or has less than 5 friends or the user does not active in connecting friends[7]). The worst case will be the same. However, we would reduce a small amount of memory.
- Instead of storing everything in a cache, we can store a subset of similar traits user and their traversal path in a separate node. Whenever there is a new connection or a new user comes in, we will put the connections/users in the corresponding node based on certain factors (e.g., number of shared friends).

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