

FAULT TOLERANT SYSTEMS

<http://www.ecs.umass.edu/ece/koren/FaultTolerantSystems>

Part 3 - Resilient Structures Chapter 2 - HW Fault Tolerance

Part.3 .1

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M-of-N Systems

- ◆ An **M-of-N** system consists of **N** identical modules
- ◆ Fails when fewer than **M** modules are functional
- ◆ Best-known example - **The Triplex (TMR)**
 - * Three identical modules whose outputs are voted on
- ◆ This is a **2-of-3** system: as long as a majority of the processors produce correct results, the system will be functional

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Reliability of M-of-N Systems

- ◆ **N** identical modules
- ◆ **R(t)** - reliability of an individual module
- ◆ The reliability of the system is the probability that **N-M** or fewer modules have failed by time **t** (or - at least **M** are functional)

$$R_{m-of-n}(t) = \sum_{i=0}^{N-M} \binom{N}{i} (1-R(t))^i R(t)^{N-i}$$

$$= \sum_{i=M}^N \binom{N}{i} R(t)^i (1-R(t))^{N-i}$$

where

$$\binom{N}{i} = \frac{N!}{i!(N-i)!}$$

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Correlated Failures in M-of-N Systems

- ◆ Key to the high reliability - statistical independence of failures in modules
- ◆ Correlated failure can greatly diminish reliability
- ◆ **Example:** q_{cor} - probability that the entire system suffers a global failure

$$R_{m-of-n-cor}(t) = (1 - q_{cor}) \sum_{i=M}^N \binom{N}{i} R(t)^i (1-R(t))^{N-i}$$

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M-of-N Systems - Modes of Correlation

- ◆ If system is not designed carefully, the correlated failure factor can dominate the overall failure probability
- ◆ Different modes of correlation among modules exist - not necessarily a global failure
- ◆ Correlated failure rates are extremely difficult to estimate
- ◆ From now on we will assume statistically independent failures in modules

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Reliability of TMR - Triple Modular Redundant Cluster

- ◆ M-of-N system with M=2, N=3 - system good if at least two modules are operational
- ◆ A voter picks the majority output
- ◆ Voter can fail - reliability of voter $R_{vot}(t)$

$$\begin{aligned}
 R_{tmr}(t) &= R_{vot}(t) \sum_{i=0}^1 \binom{3}{i} (1-R(t))^i R(t)^{3-i} \\
 &= R_{vot}(t) \sum_{i=2}^3 \binom{3}{i} R(t)^i (1-R(t))^{3-i} \\
 &= R_{vot}(t) (3R^2(t) - 2R^3(t))
 \end{aligned}$$

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Reliability of TMR - Constant Failure Rates

- ◆ $R(t) = e^{-It}$
- ◆ Assuming no voter failures - $R_{\text{vot}}(t) = 1$

$$R_{\text{tmr}}(t) = 3e^{-2It} - 2e^{-3It}$$

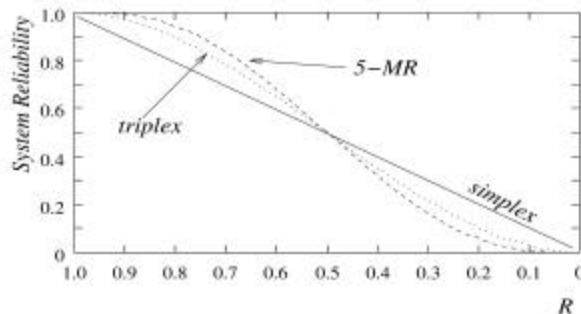
$$MTTF_{\text{tmr}} = \int_0^{\infty} R_{\text{tmr}}(t) \cdot dt = \frac{5}{6I} < \frac{1}{I} = MTTF_{\text{simplex}}$$

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NMR - N-Modular Redundant Cluster

- ◆ M-of-N cluster with N odd and $M = (N+1)/2$
- ◆ Assume voter failure rate negligible - $R_{\text{vot}}(t) = 1$



- ◆ Below $R=0.5$ - redundancy becomes a disadvantage
- ◆ Usually $R \gg 0.5$ - triplex offers significant reliability gains

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Compensating & Non-overlapping Faults

- ◆ **Conservative assumption** - every failure of voter leads to an erroneous output and any failure of two modules is fatal
- ◆ **Counter Example** - one module produces a permanent logical 1 and a second module has a permanent logical 0 - **TMR** will function properly
 - * These are compensating faults
- ◆ A similar situation may arise regarding certain faults within the voter circuit
- ◆ **Another example** - **non-overlapping faults** - one module has a faulty adder and another module has a faulty multiplier
- ◆ If the circuits are disjoint, they are unlikely to generate wrong outputs simultaneously

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Voters

- ◆ A voter receives inputs X_1, X_2, \dots, X_N from an **M-of-N** cluster and generates a representative output
- ◆ **Simplest voter** - bit-by-bit comparison of the outputs producing the majority vote
- ◆ This only works when all functional processors generate outputs that match bit by bit
 - * Processors must be identical, be synchronized and use the same software
- ◆ **Otherwise** - two correct outputs can diverge slightly, in the lower significant bits

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Plurality Voting

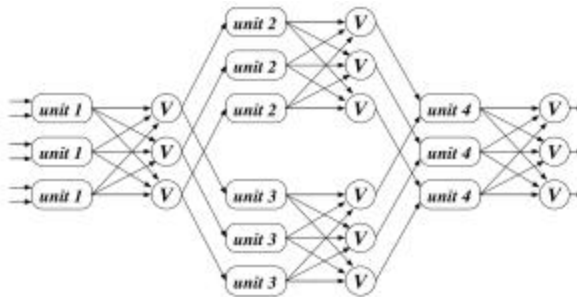
- ◆ We declare two outputs X and Y as practically identical if $|x-y| < d$ for some specified d
- ◆ A **k-plurality voter** looks for a set of at least k practically identical outputs, and picks any of them (or their median) as the representative
- ◆ **Example** - $d = 0.1$, five outputs
- ◆ 1.10, 1.11, 1.32, 1.49, 3.00
- ◆ The subset {1.10, 1.11} would be selected by a **2-plurality voter**

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Variations on NMR

- ◆ Unit-level
Modular
Redundancy

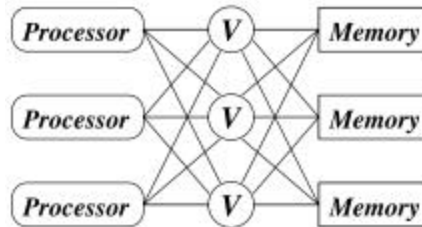


- ◆ Voters no longer critical - a single faulty voter is no worse than a single faulty unit
- ◆ The level of replication and voting can be lowered using additional voters - increasing the size and delay of the system

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Triplicated Processor/Memory System



- ◆ All communications (in either direction) between triplicated processors and triplicated memories go through majority voting
- ◆ Higher reliability than a single majority voting of triplicated processor/memory structure

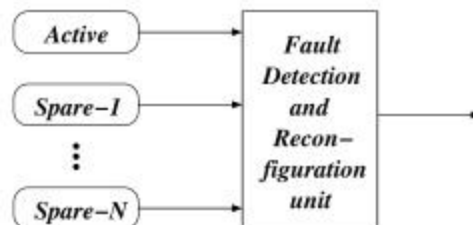
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Active/Dynamic Redundancy

- ◆ In previous examples - considerable extra hardware used to instantaneously mask errors
- ◆ In many cases, temporary erroneous results may be acceptable if
 - * system can detect an error
 - * replace the faulty module by a fault-free spare
 - * reconfigure itself
- ◆ This is called dynamic (or active) redundancy

Example:



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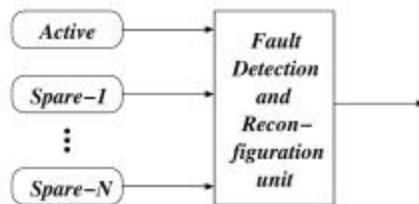
Reliability of Dynamic Redundancy - Powered Spares

- ◆ All N spare modules are active (powered) and have the same failure rate - resulting in a basic parallel system with $N+1$ modules
- ◆ System reliability is

$$R_{dynamic}(t) = R_{dru}(t)[1 - (1 - R(t))^{N+1}]$$

$R(t)$ - reliability of module

$R_{dru}(t)$ - reliability of Detection & Reconfiguration unit



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Dynamic Redundancy with Unpowered (Standby) Spares

- ◆ Spare modules are not powered (e.g., to conserve energy) and cannot fail until they become active
- ◆ C - coverage factor - probability that faulty active module is correctly diagnosed and disconnected, and good spare successfully connected
- ◆ Calculating exact reliability for the general case is complicated
- ◆ Reliability for a special case:

- * Very large N ; constant failure rate λ per active module
- * Rate of nonrecoverable faults is $(1-C)\lambda$
- * Reliability at time t - probability of no nonrecoverable faults up to time t

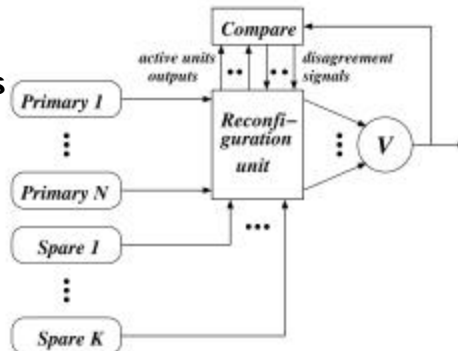
$$R_{dynamic}(t) = R_{dru}(t)e^{-(1-c)\lambda t}$$

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Hybrid Redundancy

- ◆ **NMR** masks permanent and intermittent failures but its reliability drops below that of a single module for very long mission times
- ◆ Hybrid redundancy overcomes this by adding spare modules to replace active modules once they become faulty
- ◆ A hybrid system consists of a core of **N** processors (**NMR**), and **K** spares



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Hybrid Redundancy - Reliability

- ◆ Reliability of a hybrid system with a **TMR** core and **K** spares is

$$R_{\text{hybrid}}(t) = R_{\text{vot}}(t)R_{\text{rec}}(t)[1 - mR(t)(1 - R(t))^{m-1} - (1 - R(t))^m]$$

- * **m=K+3** - total number of modules
- * **R_{vot}(t)** and **R_{rec}(t)** - reliability of voter and comparison & reconfiguration circuitry
- * Assuming: any fault in voter or comparison & reconfiguration circuit will cause a system fault
- ◆ In practice, not all faults in these circuits will be fatal: the reliability will be higher

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Sift-Out Modular Redundancy

- ◆ Like **NMR** all **N** modules are active but simpler than hybrid redundancy
- ◆ Comparing output of each module to outputs of other still operational modules
 - * A module whose output disagrees with other is switched out
- ◆ Sift-out should not be too aggressive - most failures are transient
- ◆ Purge a module only if it produces incorrect outputs over a sustained period of time

