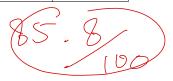
A T		
Name:		

- This quiz is **closed book**, but you may have one $8.5 \times 11''$ sheet with notes in your own handwriting on both sides.
- Calculators are not allowed.
- You may assume all of the results presented in class.
- Please show your work. Partial credit cannot be given for a wrong answer if your work isn't shown.
- Write your solutions in the space provided. If you need more space, write on the back of the sheet containing the problem. Please keep your entire answer to a problem on that problem's page.
- Be neat and write legibly. You will be graded not only on the correctness of your answers, but also on the clarity with which you express them.
- If you get stuck on a problem, move on to others. The problems are not arranged in order of difficulty.
- The exam ends at 9:30 PM.

Problem	Points	Grade	Grader
1	10	0	
2	10	10	
3	20	20	
4	15	20	
5	20	20	
6	25	25	
7	10	5	
8	10	8	
Total	120	102	



Problem 1. [10 points]

Consider these two propositions:

$$P: (A \lor B) \Rightarrow C$$

$$Q: (\neg C \Rightarrow \neg A) \lor (\neg C \Rightarrow \neg B)$$

Which of the following best describes the relationship between P and Q? Please circle exactly one answer. 7 CATAVTCATB

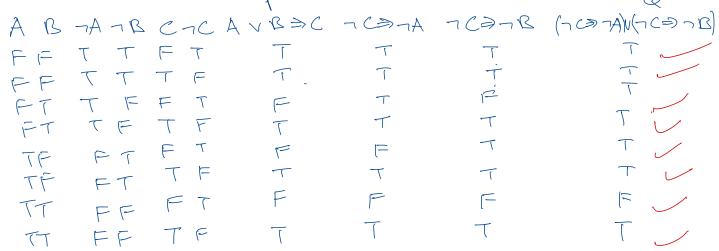
- 1. *P* and *Q* are equivalent
- 2. $P \Rightarrow Q$
- 3. $Q \Rightarrow P$
- 4. All of the above

5) None of the above

FT T

T

Draw a truth table to illustrate your reasoning. You can use as many columns as you need.







Problem 2. [10 points]

Let $G_0 = 1$, $G_1 = 3$, $G_2 = 9$, and define

$$G_n = G_{n-1} + 3G_{n-2} + 3G_{n-3} \tag{1}$$

for $n \ge 3$. Show by induction that $G_n \le 3^n$ for all $n \ge 0$.

By stong induction,

Let PLM):= Gn=Gn-1+3Gn-2+36n-2, Gn < 3° for Yn E72. Go=1, G1=3, G2=9.

Bax case, n=3: 62=62+36,+36,

Indichustep: Aggume pen) is the for KEZt, OSKEn,

For PCN+1), Gn+1=Gn+3Gn-1+3Gn-2

$$= 2^{n}(2+\frac{1}{3})$$

> pcn) stock Une72t.

Problem 3. [20 points]

In the game of Squares and Circles, the players (you and your computer) start with a shared finite collection of shapes: some circles and some squares. Players take turns making moves. On each move, a player chooses any two shapes from the collection. These two are replaced with a single one according to the following rule:

A pair of identical shapes is replaced with a square. A pair of different shapes is replaced with a circle.

At the end of the game, when only one shape remains, you are a winner if the remaining shape is a circle. Otherwise, your computer wins.

(a) [5 pts] Prove that the game will end.

Let SCn) be the number of sharpes on turn n. Lemma: For a DEIN, if a < b, then SCa) > SCb). Of: If a < b, then there are b-a turns made outler turn a: By induction on the number of turns made after a, Bove case: 6-9=1. For one turn, the player picks too shapes. Any two shapes are removed and I shape is added. the total number of shapes decrease by 1 > SCO) 7 Sb. Industice step: Suppose lemma 15 true for some b-a= k, then for b-a+1, One turn before, there are SCD) shapes. By the IH, SCD) < SCO). For ten b+1, shapes decree by I as shown similarly in termina 1. → 2(b+1) < 3(b) < S(a) > SCD+1) < SCO ⇒ scan7SCb) for Ya, b∈IN, acb Sinc SCO (SCO) for ValbEN, aCb, SCO) & shelly decreasy reach turn at some point, SCD) reaches a minimum where either SCD) comput decrean further 3) game ends.

(b) [15 pts] Prove that you will win if and only if the number of circles initially is odd.

Cicle Square Odd loesn't north

Cemma: It no of circles is odd initially, there will always be an odd amount of circle Br the entre park.

Pt: By cases on rethals of making a more.

Let n'be the no of andes. n'is add. Let m' be the no. of squaes.

Case 1: Pick both works: It 2 circles pickel, a 2 quare is odded - thre will be

n-2 circles > odd, m+1 square >> case 1 15 odd.

core 2: pich cuche & squam: a arroleis added.

Thre are n-1+1 crows 2 n crows 2 odd

m-1 2 quares.

Case 3: proh both squares. 2 1 square removed.

n choles, M-1 2 quars.

> policieles.

-. Lemma is the.

At each tern, there will be pieces removed. Ance the game must end. The game must end the game must end and another cicles. Thus there will always be I arrole left of the game start with each arroles.

Problem 4. [15 points]

(1x=1 (mod 113)

6

(a) [8 pts] Find a number $x \in \{0, 1, ..., 112\}$ such that $11x \equiv 1 \pmod{113}$.

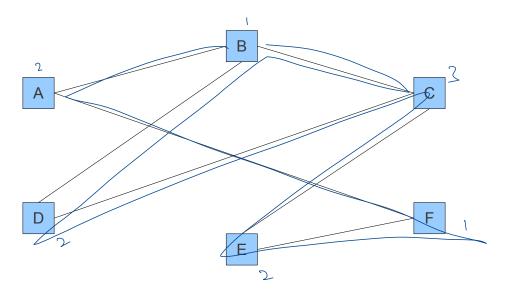
(b) [7 pts] Find a number $y \in \{0,1,\ldots,112\}$ such that $11^{112111} \equiv y \pmod{113}$ (Hint: Note that 113 is a prime.) *Note that 113 is a prime.)*

$$|| ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2} ||^{1/2$$

Problem 5. [20 points]

Consider the simple graph G given in figure 1.

Figure 1: Simple graph G



(a) [4 pts] Give the diameter of *G*.



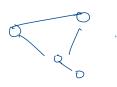
(b) [4 pts] Give a Hamiltonian Cycle on *G*.



(c) [4 pts] Give a coloring on G and show that it uses the smallest possible number of colors. $A = 2 \cdot B = 1 \cdot C = 3 \cdot E = 2 \cdot F = 1 \cdot D = 2$

By contadiction, suppore a 2 colony of G exists.

For C, since C & adjacent to B, D, E, C must be colored X and BDE colored y.
Have, B is adjacent to D, have either B or D must be colored or I forc is colored or I tower this is a contradiction on ber C to be colored X. Bond D must both be y.
Thus 3 colors is smallest.



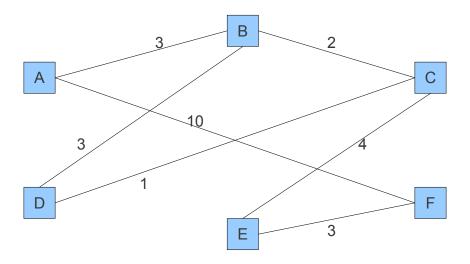
(d) [4 pts] Does *G* have an Eulerian cycle? Justify your answer.

No. Eclesion implies all vertex mot have even degrees. However, the are cortism the odd degree as not Ecleman.



Now consider graph *H*, which is like *G* but with weighted edges, in figure 2:

Figure 2: Weighted graph H



(e) [4 pts] Give a list of edges reflecting the order in which one of the greedy algorithms presented in class (i.e. in lecture, recitation, or the course text) would choose edges when finding an MST on *H*.

Problem 6. [25 points] Let G be a graph with m edges, n vertices, and k components. Prove that G contains at least m - n + k cycles. (Hint: Prove this by induction on the number of edges, m)

By induction on the number of edges, Base are: m=1.

1-n+k=k+1-n. Since ledge only, there is only, I component $\Rightarrow k+1=2 \Rightarrow the are 2-n$ eycles 2-n<0 $\Rightarrow posposition is true.$

Jerlider step. mt1,

It m+1 edges, support we remove an edge, then there are m-n+4 cucles by the induction hypothesis. Let this growth have m edges, nvertices, ke components.

Now we add back the edge. This edge can either be a unique path between 2 comporants, or it can be an edge between 2 connected components.

If ingle path, then k decreases by 1. There are no changes in the no. of sides.

Let k, be the amount of components for the initial graph.

K = ko-1. [Le number of Godes is M+1-n+k-1 = m-n+ko.]

The number of Godes is M+1-n+k-1 = m-n+ko.

If edge between 2 comparedo, then no-of cycles by papasition is m+1-n+k=m-n+k+1

) I more than before

> for since #readdition of this edge mist make at (east 1 rear cycle

i. propositor is true her all and of edges.

Ď.

Problem 7. [10 points] For the following sum, find an upper and a lower bound that differ by at most 1.

$$\sum_{i=1}^{\infty} \frac{1}{\sqrt{i^3}}$$

$$\int_{1}^{\infty} \frac{1}{i^2 z} dz = \int_{1}^{\infty} i^{-\frac{z}{2}} dz = -\frac{1}{2} \int_{1}^{\infty} i^{-\frac{z}{2}} dz = -$$

Problem 8. [10 points] State whether each of the following claims is True or False and prove your answer.

(a) [2 pts] $x \ln x$ is O(x)

$$\chi(\eta\chi=\mathcal{O}(\chi)) \Rightarrow \lim_{z\to\infty} \frac{\chi(\eta\chi)}{\chi} < \infty$$

$$\lim_{x\to 0} x \ln x = \lim_{x\to \infty} (ux = \infty)$$

(b) [2 pts] x/100 is o(x)

$$\frac{\chi}{\cos^2 = \cos^2 x} = \cos^2 x$$

(c) [2 pts] x^{n+1} is $\Omega(x^n)$

$$Z^{nel} = \mathcal{Q}(\chi^n) \Rightarrow \lim_{\chi \to \infty} \frac{\chi^{nel}}{\chi^n} > 0.$$

$$\lim_{\chi \to \infty} \frac{\chi^{nel}}{\chi^n} = \lim_{\chi \to \infty} \chi = \infty$$

$$\Rightarrow \lim_{\chi \to \infty} \chi = 0.$$

(d) [4 pts] n! is $\Theta(n^n)$.

$$\begin{aligned}
& \prod_{n \geq 0} & \bigoplus_{n \geq 0} & \prod_{n \geq 0} &$$

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