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## Compression

reduce amount of data to represent same info

data: conveys information

information: ?

redundant data: irrelevant or repeated

$b + b' \equiv \cancel{20}$  bits in 2 rep's  
for same info

relative data redundancy:

$$R = 1 - \frac{1}{C}$$

where  $C = \frac{b}{b'}$  compression ratio

e.g.,  $b = 100$  &  $b' = 10$

then  $C = 10 \Rightarrow$  10 bits in  $b$  & 1 bit in  $b'$

90% redundant

$b_2 \setminus 4$

$b = \#$  bits in image

principal types of redundancy :

Coding code is system of symbols used to represent information  
each piece of info is assigned a sequence of code symbols (code word)  
 $\#$  symbols in code word  $\Rightarrow$  its length  
usually 8-bits in image

spatial spatially correlated  
(temporal) video

irrelevant : not used by vision system

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Coding redundancy

$r_k$  : random variable in  $[0, L-1]$  used to represent intensities in  $M \times N$  image

$p(r_k)$  : probability of  $r_k$

$$= \frac{n_k}{MN}$$

$$k = 0, \dots, L-1$$

$n_k \equiv$  count of  $r_k$

average # of bits / pixel

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$$L_{avg} = \sum_{k=0}^{L-1} l(r_k) p_r(r_k)$$

$l(r_k) \equiv$  # bits to represent  $r_k$

total # bits:  $M N L_{avg}$

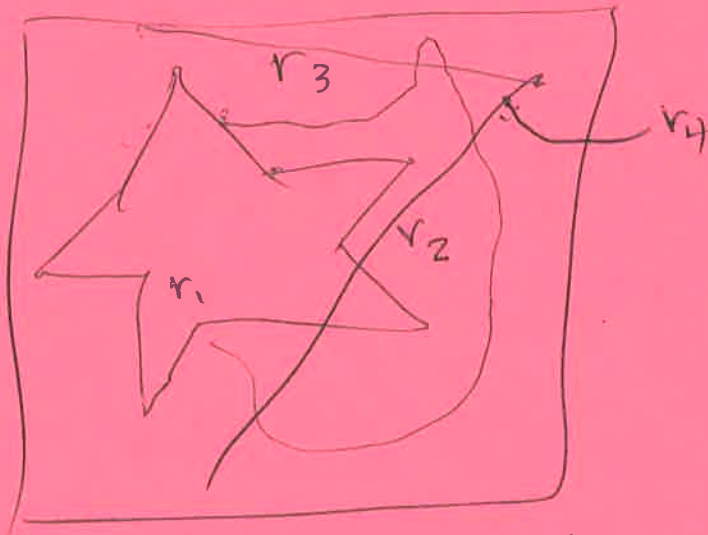
e.g., if  $8 = m$  bits, then

$$\begin{aligned} L_{avg} &= \sum_{k=0}^{L-1} m p_r(r_k) \\ &= m \sum_{k=0}^{L-1} p_r(r_k) \\ &= m \end{aligned}$$

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Consider image in book 8.1



	$p_r(r_k)$	Code1	$l_1(r_k)$	Code2	$l_2(r_k)$
$r_1 = 87$	.25	01010111	8	01	2
$r_2 = 128$	.47	11101000	8	.1	1
$r_3 = 186$	.25	10111010	8	000	3
$r_4 = 255$	.03	11111111	8	001	3

$L_{avg} =$

$$C = \frac{8}{1.81} = 4.42$$

$$R = 1 - \frac{1}{4.42} = 0.774$$

$$0.25(2) + 0.47(1) + .28(3) = 1.81 \text{ bits}$$

what about using 2 bits?

$$C = \frac{8}{2} = 4 \quad \sim 10\% \text{ worse}$$

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Most images are of some thing, and the intensities are in a narrow range. thus, better to use short codes for these

information theory modeled as prob. process  
 $E$ : random event  $I(E)$ : info contained by  $E$

$$I(E) = \log \frac{1}{P(E)} = -\log P(E)$$

$$> p = [0 : 0.01 : 1];$$

$$> b = -\log_2(p); \rightarrow \text{measured in bits}$$

$$> b(56)$$

$$> b(26)$$

$$> b(\text{end})$$

average info per source output (Events)  $a_j$  ← source symbols

entropy:

$$H = - \sum_{j=1}^J P(a_j) \log P(a_j)$$

for image

$$\tilde{H} = - \sum_{k=0}^{L-1} p_r(r_k) \log_2 p_r(r_k)$$

how can this measure be used? To pick compression method

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## Complexity measures

\* How difficult to describe or build system  
description length of string  
shortest computer program that generates string

\* Amount system is organized or  
amount of interesting structure



order - disorder related to  
information + entropy

Shannon: entropy is high if many states  
have equal probability

next symbol has most info when  
all are equally likely

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fidelity criteria

quantify loss : objective vs. subjective

$$f(x, y) \xrightarrow{\text{compress}} f_c(x, y) \xrightarrow{\text{decompress}} \hat{f}(x, y)$$

error  $e(x, y) = \hat{f}(x, y) - f(x, y)$

$e_{\text{total}}$   $= \sum_{x=1}^M \sum_{y=1}^N [\hat{f}(x, y) - f(x, y)]$

root-mean-square error

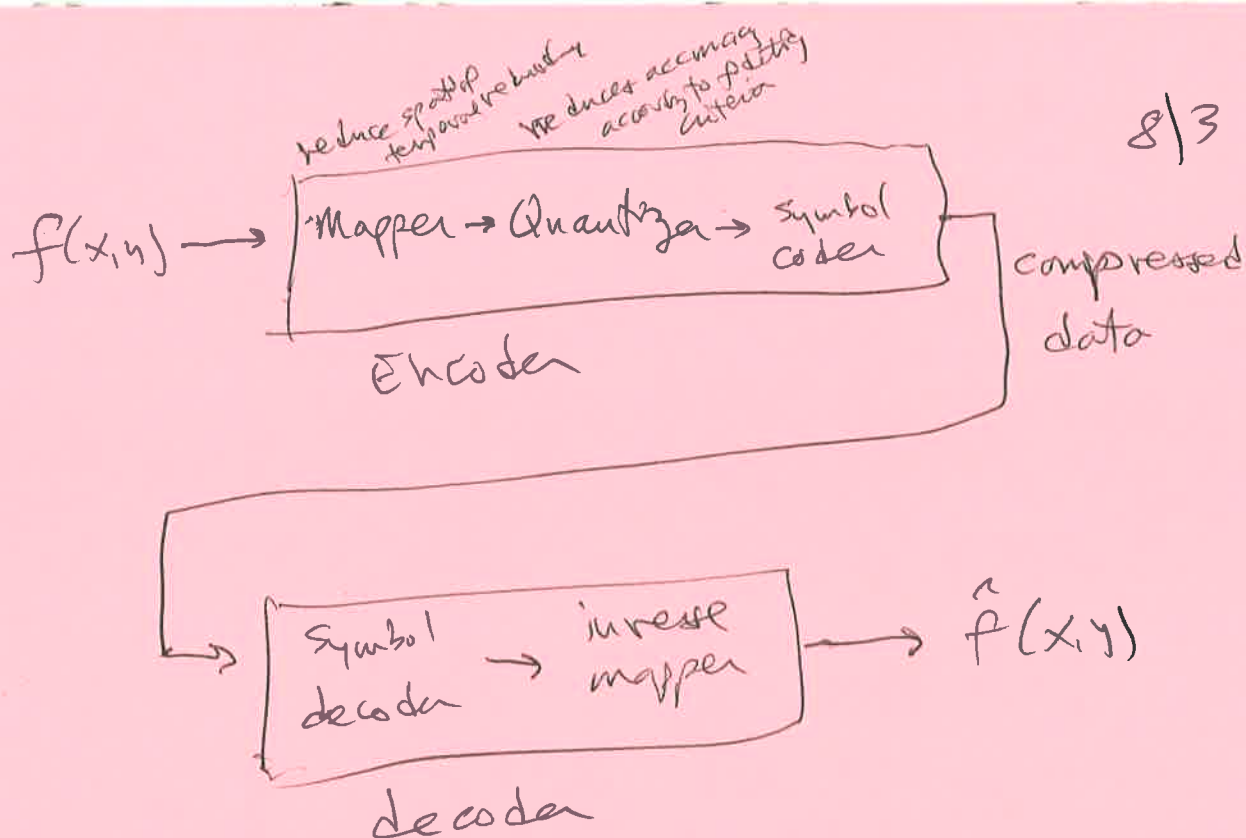
$$e_{\text{rms}} = \left[ \frac{1}{MN} \sum_{x=1}^M \sum_{y=1}^N [\hat{f}(x, y) - f(x, y)]^2 \right]^{1/2}$$

mean squared ~~error~~ <sup>signal-to-noise ratio</sup>

$$\text{SNR}_{\text{ms}} = \frac{\sum_{x=1}^M \sum_{y=1}^N \hat{f}(x, y)^2}{\sum_{x=1}^M \sum_{y=1}^N [\hat{f}(x, y) - f(x, y)]^2}$$

$\text{SNR}_{\text{rms}}$   $= \sqrt{\text{SNR}_{\text{ms}}}$





codec

if  $f = \hat{f}$  error free, lossless, info preserving

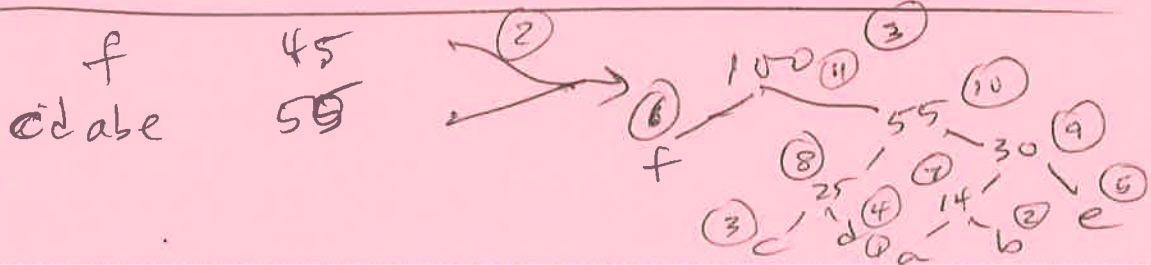
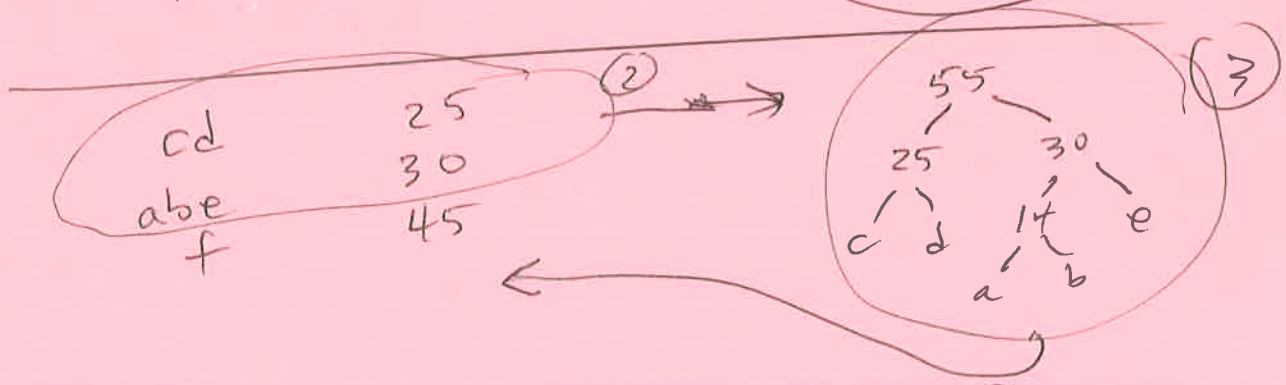
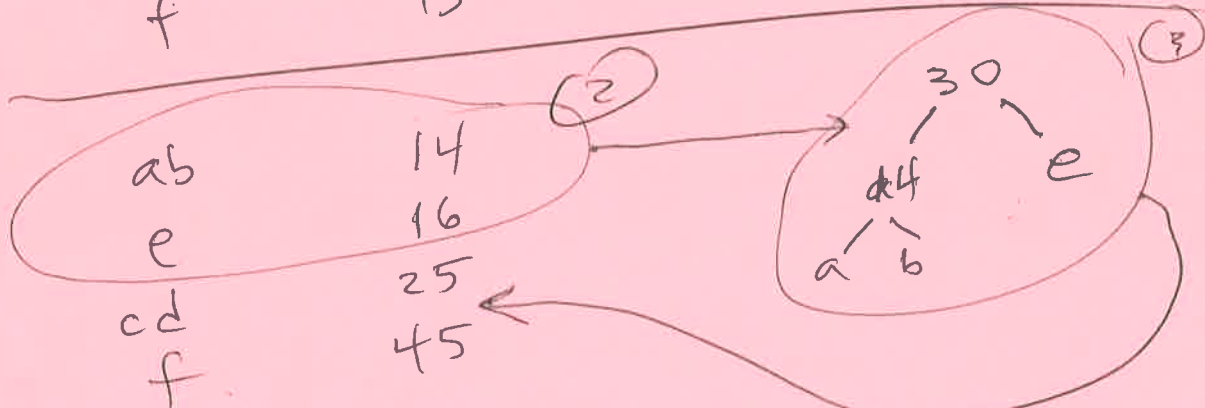
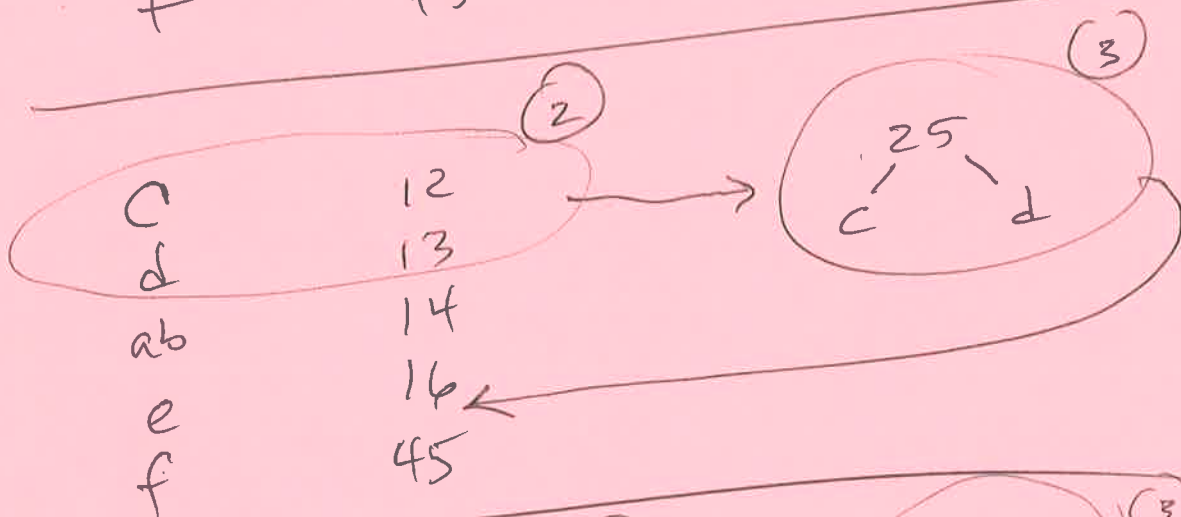
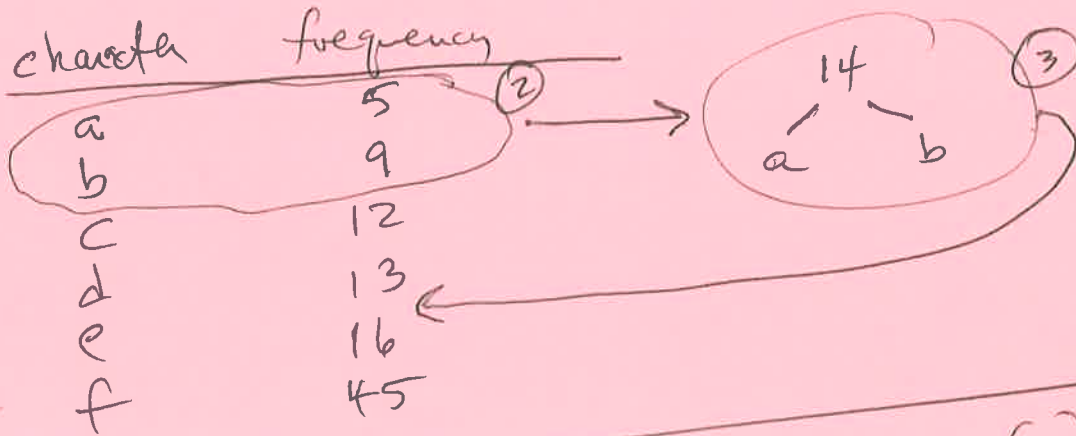
else lossy

image file format: organization standard

image container: multiple types of image data

compression standard: de (compression) procedures





## Huffman Coding

algorithm ([www.geeksforgeeks.org](http://www.geeksforgeeks.org))

### Huffman Tree

input: array of unique characters & frequencies

output: Huffman tree

1. create sorted<sup>by frequencies</sup> list of characters (ascending)
2. Extract<sup>(remove)</sup> two lowest frequency nodes,  $n_1$  &  $n_2$
3. Create new node,  $n$ , whose frequency is sum from  $n_1$  &  $n_2$  with  $n_1$  as left child &  $n_2$  as right child

Insert ~~new~~ into sorted list

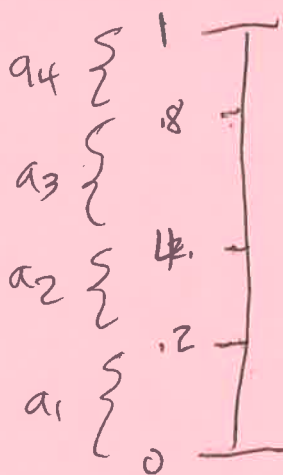
4. Repeat 2 & 3 until only 1 node

### Huffman Codes

Traverse tree to every leaf; for left move assign 0, for right assign 1. Code is sequence generated by path.

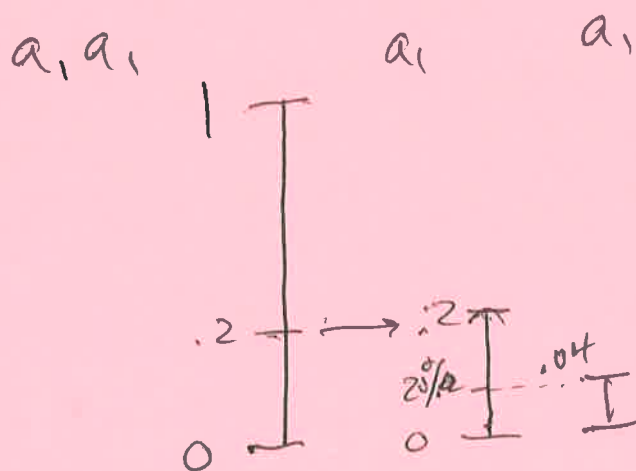
# Arithmetic coding

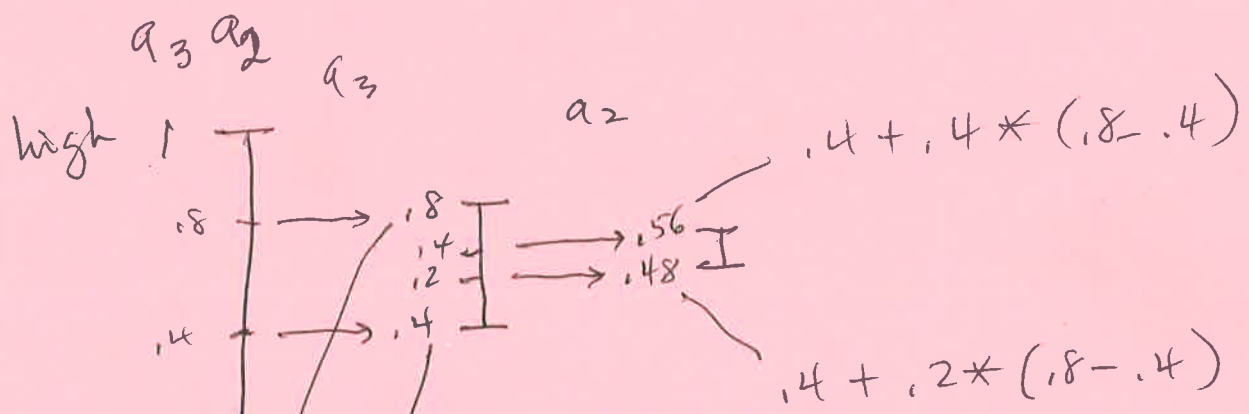
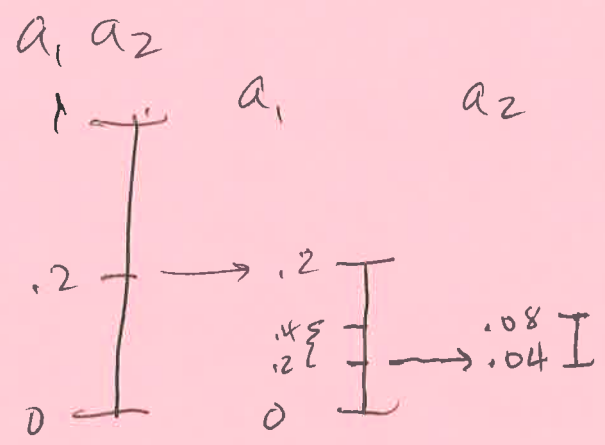
Given symbols, probabilities and a sequence of symbols, convert to an interval in  $[0, 1]$ .



if string has 1 symbol, then know which ~~one~~ by interval  
 $a_1 \rightarrow [0, .2)$      $a_2 \rightarrow [.2, .4)$     ...

if string has 2 symbols then find interval as:  
 1<sup>st</sup> symbol restricts to its interval,  
 2<sup>nd</sup> symbol restricts to its interval on 1<sup>st</sup> interval





$$low = low + interval(s, 1) * (high - low)$$

$$high = low + interval(s, 2) * (high - low)$$

see week 7 code