

# Space Optimal Vertex Cover in Dynamic Streams

(with an overview of graph streaming)

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Joint work with Vihan Shah (Rutgers University)

- 1 Introduction
- 2 Streaming Models
- 3 Matchings in Graph Streams
- 4 Space Optimal Vertex Cover

# 1 Introduction

## 2 Streaming Models

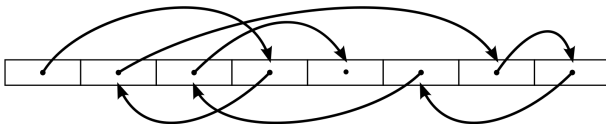
## 3 Matchings in Graph Streams

## 4 Space Optimal Vertex Cover

# Classic Setting

## Assumption

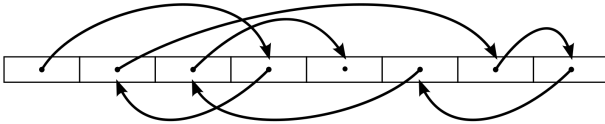
Classical algorithms rely on the assumption that they have a random access to the input of the algorithm





## Assumption (Infeasible)

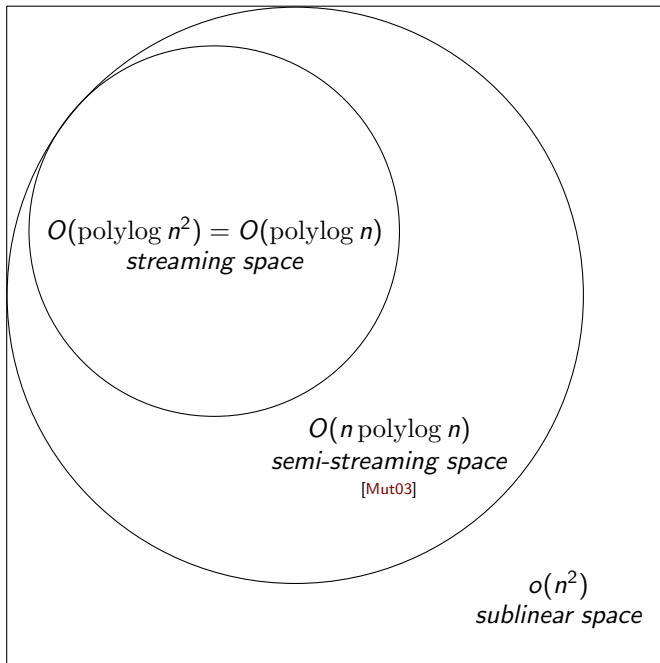
Classical algorithms rely on the assumption that they have a random access to the input of the algorithm



## Definition (Graph Streaming)

A  $n$ -vertex graph is presented as a sequence of edges to an algorithm uses space **sublinear** in the size of the input ( $o(n^2)$ ).







1 Introduction

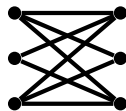
2 Streaming Models

3 Matchings in Graph Streams

4 Space Optimal Vertex Cover

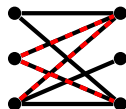
## Insertion-Only [FKM<sup>+</sup>04] (finite stream)

$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$	$e_8$
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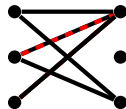
## Sliding Window [CMS13] (infinite stream)

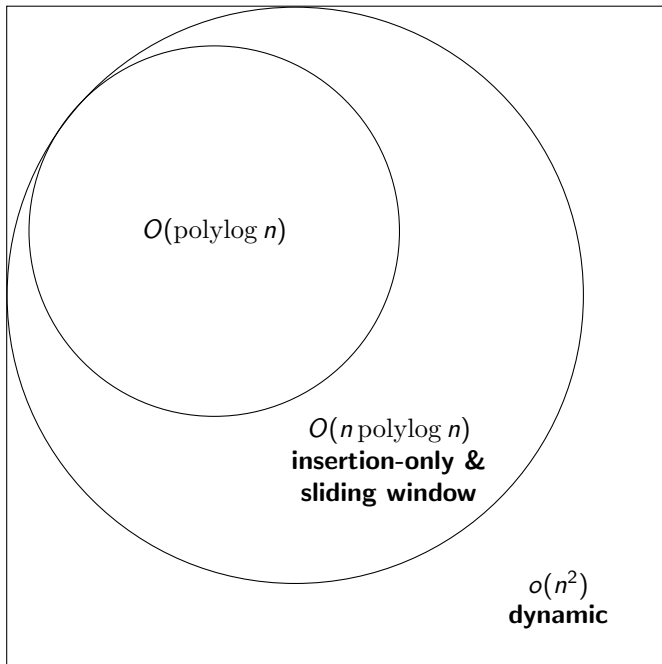
$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$	...
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## Dynamic [AGM12] (finite stream)

$e_1$	$e_2$	$e_3$	$\overline{e_1}$	$\overline{e_3}$	$e_4$	$e_5$	$e_1$
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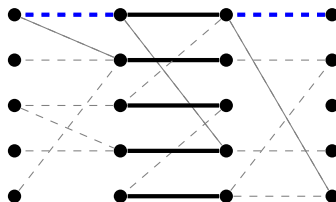


- 1 Introduction
- 2 Streaming Models
- 3 Matchings in Graph Streams**
- 4 Space Optimal Vertex Cover

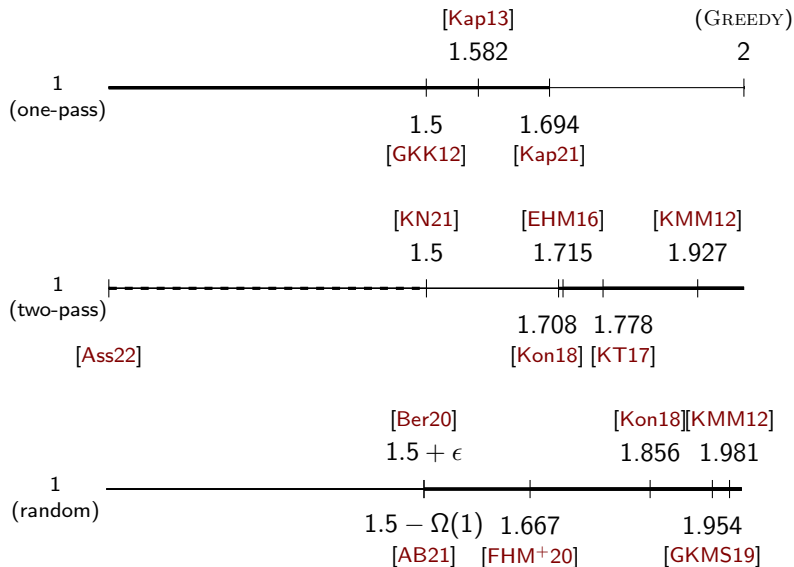
# Simple and Powerful

## GREEDY Matching:

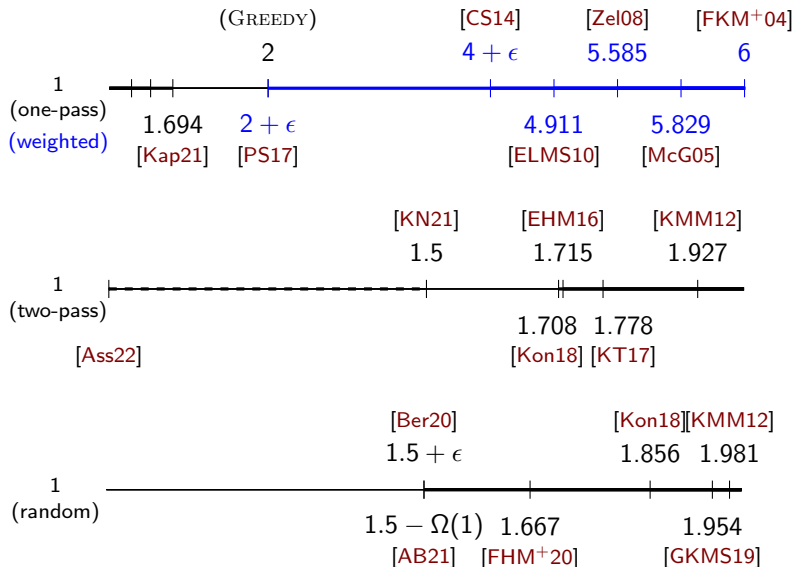
- ① Add edge if neither endpoint is matched
- Maximal
- 2-approximation



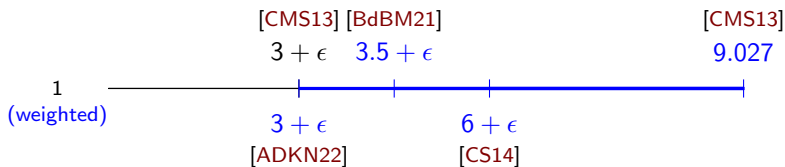
# Insertion-Only



# Insertion-Only



# Sliding Window

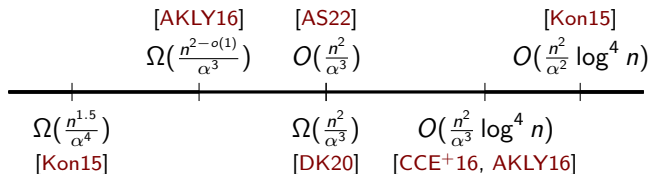




# Dynamic (Insertion-Deletion)

Goal:

- Find an  $\alpha$ -approximation

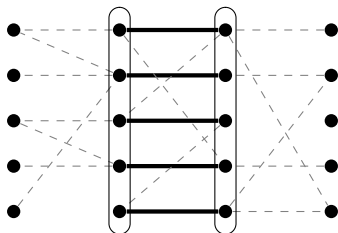


# Relation to Vertex Cover

GREEDY

MatchingVertex Cover:

- Maximal
- 2-approximation



# Relation to Vertex Cover (Results)

	matching	vertex cover
insertion-only	$[1.694, 2]$	$[1, 2]$ GREEDY
sliding window	$[1, 3 + \epsilon]$	$[1, 3 + \epsilon]$ [Sub21]
dynamic	$\Theta(\frac{n^2}{\alpha^3})$	$\Theta(\frac{n^2}{\alpha^2})$ [DK20] & this work

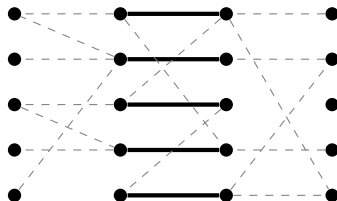
All UBs use GREEDY in some way!

- 1 Introduction
- 2 Streaming Models
- 3 Matchings in Graph Streams
- 4 Space Optimal Vertex Cover

# What makes dynamic graph streams hard?

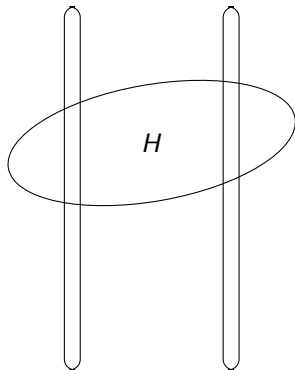
GREEDY Matching:

- **Not** Maximal
- **0**-approximation
- 2-approximation



**Note:** Deterministically returning a single edge requires  $\Omega(n^2)$  bits of space for dense graphs.

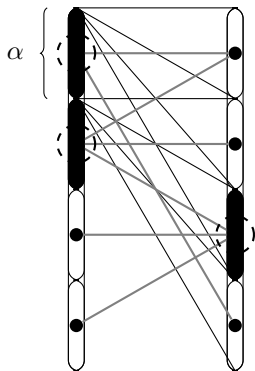
# What techniques are used?



For a subgraph  $H$  of  $G$  with  $m$  edges:

- ① Check if  $H$  is empty
  - A counter using  $\Theta(\log m)$  bits
- ② Retrieve an edge if one is present
  - An  $L_0$ -sampler using  $\Theta(\log^3 n)$  bits
  - Neighbourhood sampler

# $\alpha$ -Approx Det. Dynamic Vertex Cover [DK20]



- ① Vertex groups of size  $\alpha$ 
  - about  $\frac{n}{\alpha}$  groups
- ② Check if there is an edge between each pair of groups
  - about  $\frac{n^2}{\alpha^2}$  pairs
- ③ Return vertices of the group-level vertex cover

**Space:**  $O(\frac{n^2}{\alpha^2})$  counters, each using  $O(\log \alpha)$  bits. Hence,  $O(\frac{n^2}{\alpha^2} \log \alpha)$  bits.

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$\Omega(\frac{n^2}{\alpha^2})$	$O(\frac{n^2}{\alpha^2} \log \alpha)$
[DK20]	[DK20]

# What's the issue?

Problem:

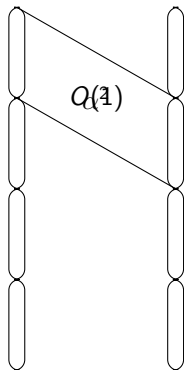
- Counters use  $O(\log \alpha)$  bits.
- Each counter counts upto  $\alpha^2$  edges.

Goal:

- Counters to use  $O(1)$  bits.
- Counters to count upto constant many edges

How?

- Sparse graph (GREEDY!)
- Randomly partition (easy)



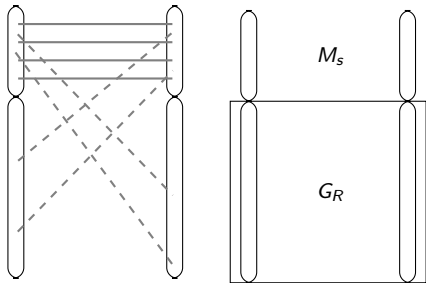


# Sparsness properties of GREEDY

## Lemma

Let  $G$  be a  $n$ -vertex graph with  $m$  edges and let  $M_s$  be a GREEDY matching on  $s \leq m$  uniform randomly sampled edges. Then, for  $G_R = G[V \setminus V(M_s)]$ ,

$$\Delta(G_R) \leq \frac{C \cdot m \cdot \log n}{s}.$$



Proof (sketch).

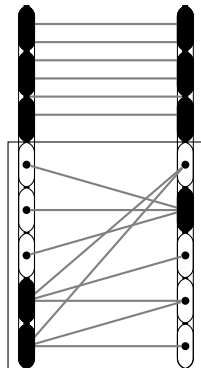
For any vertex  $v \in G$ ,

- $v \in G_R$
- $\deg_{G_R}(v) > \frac{C \cdot m \cdot \log n}{s}$

both do not occur w.h.p. □

# Algorithm

- ① Randomly partition vertices ( $\frac{n}{\alpha}$  groups)
- ② Sample  $O(\frac{n^2}{\alpha^2 \log^3 n})$  edges randomly and run GREEDY on them to get  $M$
- ③ Check if an edge is present between pairs and compute group-level vertex cover
- ④ Return vertices of the covering groups including those with matched vertices

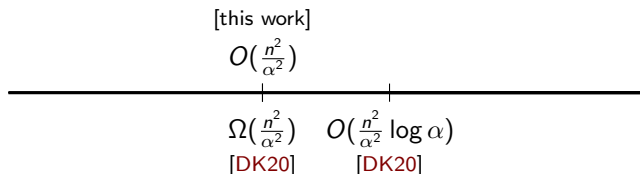


## Limitation:

- Sparseness properties are only sufficient for  $\alpha \ll n^{\frac{1}{3.5}}$

## Extension:

- Non-uniform sampling instead of uniform sampling, i.e., using neighbourhood edge sampling [AS22]
- Results are an average degree bound instead of max degree bound
- Full range, i.e.,  $\alpha \ll n$





## Open Questions:

- Deterministic techniques or LB instead
- Other problems like dominating set and spectral sparsification

# Thank You

# Questions?

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