CSE 6643 Homework 5

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1 Power Method [30 pts]

We consider a matrix \boldsymbol{A} such that

$$\mathbf{Q}^T \mathbf{A} \mathbf{Q} = \operatorname{diag}(\lambda_1, \dots, \lambda_m), \tag{1}$$

where Q is orthogonal. We denote by q_i the ith column of Q.

We now consider the power method and the sequence $oldsymbol{
u}^{(k)}$ defined as

$$\boldsymbol{z}^{(k)} = \boldsymbol{A} \boldsymbol{\nu}^{(k-1)} \tag{2}$$

$$\boldsymbol{\nu}^{(k)} = \boldsymbol{z}^{(k)} / \|\boldsymbol{z}^{(k)}\|_2,\tag{3}$$

where we assume that $\|\boldsymbol{\nu}^{(0)}\|_2=1$. Assume that we have θ_k such that

$$\cos(\theta_k) = \mathbf{q}_1^T \mathbf{\nu}^{(k)} \tag{4}$$

with $\cos(\theta_0) \neq 0$. Prove that

$$1 - \cos(\theta_k)^2 \le \frac{1}{a_1^2} \sum_{i=2}^m a_i^2 \left(\frac{\lambda_i}{\lambda_1}\right)^{2k}, \quad a_i = \boldsymbol{q}_i^T \boldsymbol{\nu}^{(0)}$$
 (5)

$$\cos(\theta_k) = q_1^T \frac{A^k \nu^{(0)}}{\|A^k \nu^{(0)}\|_2} \qquad (\text{def. of power method})$$

$$= q_1^T \frac{A^k \nu^{(0)}}{\|A^k \nu^{(0)}\|_2} \qquad (\text{def. of power method})$$

$$= \frac{q_1^T Q \Lambda^k Q^T \nu^{(0)}}{\|Q \Lambda^k Q^T \nu^{(0)}\|_2} \qquad (\text{def. of } A)$$

$$= \frac{q_1^T Q \Lambda^k a}{\|Q \Lambda^k a\|_2} \qquad (a \equiv Q^T \nu^{(0)} \rightarrow a_i = q_i^T \nu^{(0)})$$

$$\cos(\theta_k)^2 = \frac{(q_1^T Q \Lambda^k a)^2}{\|Q \Lambda^k a\|_2^2} \qquad (\text{square both sides})$$

$$= \frac{(q_1^T Q \Lambda^k a)^2}{\|Q \Lambda^k a\|_2^2} \qquad (\text{apply transpose in den.})$$

$$= \frac{(q_1^T Q \Lambda^k a)^2}{\Lambda^{2k} a^T a} \qquad (\text{simplify den.})$$

$$= \frac{(q_1^T Q \Lambda^k a)^2}{\sum_{i=1}^m \lambda_i^{2k} a_i^2} \qquad (\text{expand den.})$$

$$= \frac{(q_1^T Q \Lambda^k a)^2}{\sum_{i=1}^m \lambda_i^{2k} a_i^2} \qquad (\text{expand den.})$$

$$= \frac{\lambda_1^{2k} a_1^2}{\sum_{i=1}^m \lambda_i^{2k} a_i^2} \qquad (\text{evaluate num.})$$

$$1 - \cos(\theta_k)^2 = 1 - \frac{\lambda_1^{2k} a_1^2}{\sum_{i=1}^m \lambda_i^{2k} a_i^2} \qquad (\text{negate \& add 1 to both sides})$$

$$= \frac{\sum_{i=1}^m \lambda_i^{2k} a_i^2}{\sum_{i=1}^m \lambda_i^{2k} a_i^2} \qquad (\text{common den.})$$

$$= \frac{\sum_{i=1}^m \lambda_i^{2k} a_i^2}{\lambda_1^{2k} a_1^2 + \sum_{i=2}^m \lambda_i^{2k} a_i^2} \qquad (\text{rearrange summation terms})$$

$$\leq \frac{\sum_{i=2}^m \lambda_i^{2k} a_i^2}{\lambda_i^{2k} a_i^2} = \frac{1}{a_1^2} \sum_{i=1}^m a_i^2 \left(\frac{\lambda_i}{\lambda_i}\right)^{2k} \qquad (\text{QED})$$

2 The LU iteration algorithm [30 pts]

We consider the following iteration, starting with some full rank $G_0 \in \mathbb{C}^{m \times m}$:

$$Z_k = AG_{k-1}, (6)$$

$$G_k R_k = Z_k$$
, LU factorization with no pivoting, (7)

where G_k is lower-triangular and R_k is upper-triangular. We define

$$T_k = G_k^{-1} A G_k. (8)$$

Find an algorithm that computes the sequence T_k in a manner similar to the QR iteration. This algorithm is a variant of the QR iteration. The eigenvalues of A appear on the diagonal of T_k .

From Piazza:

You just need to show your algorithm and are not required to prove the algorithm reveals the eigenvalues on the diagonal of T_k . It is sufficient to use the given equations to derive a sequence of T_k . You need to find a manner similar to the QR iteration to compute T_k . That is, the computation of T_k should be similar to R_kG_k . Note that here $T_k \neq R_kG_k$, but you will find that when changing the formula a little bit, e.g, inserting a matrix between R_k and G_k , it will make the equality hold. This modification should not bring much additional computational cost.

This algorithm consists of iteratively doing matrices multiplication, applying LU factorization, and then doing two different matrices multiplications to compute T_k . Its relationship to QR factorization only concerns the manner of computation.

Goal: Find a matrix \boldsymbol{B} s.t. $\boldsymbol{T}_k = \boldsymbol{G}_k^{-1} \boldsymbol{A} \boldsymbol{G}_k = \boldsymbol{R}_k \boldsymbol{B} \boldsymbol{G}_k$.

The QR Iteration algorithm:

- $T_{k-1} = U_k R_k$
- $T_k = R_k U_k$

3 Convergence of Orthogonal Iteration [40 pts]

In this problem, you will explore the convergence behavior of the orthogonal iteration algorithm. In the problems below, only the asymptotic convergence will match the theoretical estimates. For small k, you may see deviations. To simplify grading, please use the provided (empty) file HW4_your_code.jl for your code. Please also attach all results to your report, both plots and print statements. The TAs should be able to grade your homework without running your code.

(a) [10 pts]

Write Julia code to create a matrix in $\mathbb{R}^{m \times m}$ of size m = 8 with eigenvalues $1, 3, 9, \dots, 3^{m-1}$. Explain your code (as comments and in your report) and describe what it does.

(b) [10 pts]

Implement the orthogonal iteration algorithm. Print the values along the diagonal of \mathbf{R}_k at each iteration k for $k = 1, \dots, 5$. Print each number using at most four significant digits.

(c) [10 pts]

Considering entry p along the diagonal, plot the convergence of the pth eigenvalue. Choose p = 1, 2, and 3. Compare with the theoretical rate of convergence at step k for entry p, which

is given by

$$\max(|\lambda_{p+1}/\lambda_p|^k, |\lambda_p/\lambda_{p-1}|^k), \quad 1$$

$$|\lambda_2/\lambda_1|^k, \quad p = 1 \tag{10}$$

$$\left|\lambda_m/\lambda_{m-1}\right|^k, \quad p=m. \tag{11}$$

Use a semi-log plot (meaning that the y-axis of your plot should be logarithmic).

(d) [10 pts]

Consider the block (p+1:m,1:p) in the matrix

$$\boldsymbol{A}_k = \boldsymbol{Q}_k^T \boldsymbol{A} \boldsymbol{Q}_k. \tag{12}$$

Plot the 2-norm of this block as a function of k for p=4. Compare with the analytical estimate, which states that it should decay like $|\lambda_{p+1}/\lambda_p|^k$.

4 LU and QR iteration [20 bonus pts]

We consider the LU iteration applied to a symmetric matrix. Assume that $A_0 \in \mathbb{R}^{m \times m}$ is symmetric and positive-definite. We produce a sequence of matrices A_i using the following algorithm:

$$\mathbf{A}_i = \mathbf{G}_i^T \mathbf{G}_i, \quad \mathbf{A}_{i+1} = \mathbf{G}_i \mathbf{G}_i^T, \tag{13}$$

where G_i are upper-triangular matrices.

Consider now A', the matrix obtained after one step of the QR iteration, that is,

$$A_0 = QR, \quad A' = RQ. \tag{14}$$

We assume that the diagonal of R is positive.

(a) [10 bonus pts]

Use A_0^2 to show that

$$R = G_1 G_0 \tag{15}$$

(b) [10 bonus pts]

Show that

$$\mathbf{A}' = \mathbf{A}_2 \tag{16}$$

5 Sensitivity of Eigenvalues [20 bonus pts]

We are interested in determining the sensitivity of eigenvalues with respect to perturbations in the matrix.

Prove that if $A \in \mathbb{C}^{m \times m}$ is diagonalizable with $A = V \Lambda V$, and if μ is an eigenvalue of A + E,

$$\min_{\lambda \in \lambda(\mathbf{A})} |\mu - \lambda| \le \kappa_p(\mathbf{V}) \|\mathbf{E}\|_p. \tag{17}$$

Here $\|\cdot\|_p$ denotes the *p*-norm and $\kappa_p(\mathbf{V}) \coloneqq \|\mathbf{V}\|_p \|\mathbf{V}^{-1}\|_p$ is the condition number associated with the *p*-norm.

Hint: assume that $\mathbf{I} + \mathbf{F}$ is singular. Then, there is an $\mathbf{x} \neq \mathbf{0}$ such that $\mathbf{x} + \mathbf{F}\mathbf{x} = \mathbf{0}$. Therefore, $\mathbf{F}\mathbf{x} = -\mathbf{x}$ and $\|\mathbf{F}\|_p \geq 1$. For the proof, identify the proper singular matrix and use this result.