

Representation of Simply-Typed λ -Calculus

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November 29, 2020

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1 Introduction

This documentation was created for Assignment 2 of 3MI3, it's purpose is to help understand the fundamentals behind how (and why) my program runs. Lambda calculus is a formal system for expressing computation based on function abstraction and application. We will attempt to create a typed interpretation of the lambda calculus (also known as Simply-Typed λ -Calculus) using Scala and Ruby.

2 Representation

Our representation adds on to the standard untyped λ calculus expressions (**variable**, **abstraction**, **application**) with some additional terms (**zero**, **successor**, **isZero**, **true**, **false**, **test**) as well as a way to represent these terms with a type.

2.1 Scala

For Scala, we will classify our main term `STTerm` and type `STType` as a trait and our typed terms and custom types as classes that extend this trait. Our variable term will only take an integer input since a variable will be represented by an integer, and `True/False/Zero` are all objects that need no parameters but the rest of the terms will take inputs that match the type term as well as the type of the custom types that we create.

```
sealed trait STTerm
case class STVar(index: Int) extends STTerm {
  override def toString() = index.toString()
}
case class STApp(t1: STTerm, t2: STTerm) extends STTerm {
  override def toString() = "(" + t1.toString() + ") (" + t2.toString() + ")"
}
case class STAbs(t: STType, t1: STTerm) extends STTerm {
  override def toString() = "lambda " + t.toString() + " . " + t1.toString()
}
case object STZero extends STTerm {
  override def toString() = "0"
}
case class STSuc(t: STTerm) extends STTerm {
  override def toString() = "S " + t.toString()
}
case class STIsZero(t: STTerm) extends STTerm {
  override def toString() = t.toString() + " ?"
}
case object STTrue extends STTerm {
  override def toString() = "True"
}
case object STFalse extends STTerm {
  override def toString() = "False"
}
case class STTest(t1: STTerm, t2: STTerm, t3: STTerm) extends STTerm {
  override def toString() = "test " + t1.toString() + " " + t2.toString() + " " + t3.toString()
}

sealed trait STType
case object STNat extends STType {
  override def toString() = "nat"
}
case object STBool extends STType {
  override def toString() = "bool"
}
case class STFun(dom: STType, codom: STType) extends STType {
  override def toString() = "(" + dom.toString() + ") -> (" + codom.toString() + ")"
}
```

2.2 Ruby

For Ruby, we define our main term `STTerm` and type `STType` as classes and our typed terms and custom types as classes that extend this main class. Our term parameters are initialized with "attr_reader" to allow accessing. Moreover, for individual terms and types, we define "==" functions for comparisons between two terms/types.

```
class STTerm
end

class STVar < STTerm
  attr_reader :index

  def initialize(i)
    unless i.is_a?(Integer)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @index = i
  end

  def ==(r); r.is_a?(STVar) && r.index == @index end
end

class STApp < STTerm
  attr_reader :t1
  attr_reader :t2

  def initialize(t1, t2)
    unless t1.is_a?(STTerm) && t2.is_a?(STTerm)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @t1 = t1
    @t2 = t2
  end

  def ==(r); r.is_a?(STApp) && r.t1 == @t1 && r.t2 == @t2 end
end

class STAbs < STTerm
  attr_reader :t
  attr_reader :t1

  def initialize(t, t1)
    unless t.is_a?(STType) && t1.is_a?(STTerm)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @t = t
    @t1 = t1
  end

  def ==(r); r.is_a?(STAbs) && r.t == @t && r.t1 == @t1 end
end

class STZero < STTerm
  def ==(r); r.is_a?(STZero) end
end
```

```

end

class STSuc < STTerm
  attr_reader :t

  def initialize(t)
    unless t.is_a?(STTerm)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @t = t
  end

  def ==(r); r.is_a?(STSuc) && r.t == @t end
end

class STIsZero < STTerm
  attr_reader :t

  def initialize(t)
    unless t.is_a?(STTerm)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @t = t
  end

  def ==(r); r.is_a?(STIsZero) && r.t == @t end
end

class STTrue < STTerm
  def ==(r); r.is_a?(STTrue) end
end

class STFalse < STTerm
  def ==(r); r.is_a?(STFalse) end
end

class STTest < STTerm
  attr_reader :t1
  attr_reader :t2
  attr_reader :t3

  def initialize(t1, t2, t3)
    unless t1.is_a?(STTerm) && t2.is_a?(STTerm) && t3.is_a?(STTerm)
      throw "Constructing a lambda term out of non-lambda terms"
    end
    @t1 = t1
    @t2 = t2
    @t3 = t3
  end

  def ==(r); r.is_a?(STTest) && r.t1 == @t1 && r.t2 == @t2 && r.t3 == @t3 end
end

class STType end

```

```

class STNat < STType
  def ==(type); type.is_a?(STNat) end
  def to_s; "nat" end
end

class STBool < STType
  def ==(type); type.is_a?(STBool) end
  def to_s; "bool" end
end

class STFun < STType
  attr_reader :dom
  attr_reader :codom

  def initialize(dom, codom)
    unless dom.is_a?(STType) && dom.is_a?(STType)
      throw "Constructing a type out of non-types"
    end
    @dom = dom; @codom = codom
  end

  def ==(type); type.is_a?(STFun) && type.dom == @dom && type.codom == @codom end
  def to_s; "(" + dom.to_s + ") -> (" + codom.to_s + ")" end
end

```

3 Typechecking

Now that we have defined our terms, we can create a function `typecheck` that checks if our term obeys the type rules. This is done by checking if the parameters match the typerules and the parameters of those and so on (using recursion). If at any point, typerules do not match, we note this and come back to our main function to return false for the term being properly typed (both languages use different ways to check this). While typechecking, we keep track of variable types through the use of an environment (I have gone with a List implementation).

3.1 Scala

For Scala, we define a new function `typecheck` and a function `typeOf` within it. `typecheck` calls the `typeOf` with the empty environment and `typeOf` pattern matches the parameter with different `STTerms` and recurses until either failure of typerules or reaching the end. I also use the implementation of `Either[]` in this function to keep track of failed typerules (represented by `None`). If we encounter `None` at some point, the function simply returns `False` otherwise `True`

```
def typecheck(term: STTerm): Boolean = {
  def typeOf(env: List[STType], exp: STTerm): Option[STType] = exp match {
    case STVar(index) => (env.length > index) match {
      case true => Some(env(index))
      case false => None
    }
    case STAbs(t, t1) => {
      val newenv = List(t) ++ env
      val newtype = typeOf(newenv, t1)
      newtype match {
        case Some(value) => Some(STFun(t, value))
        case None => None
      }
    }
    case STApp(t1, t2) => typeOf(env, t1) match {
      case Some(value) => value match {
        case STBool => None
        case STNat => None
        case STFun(dom, codom) => typeOf(env, t2) match {
          case Some(value) if (value == dom) => Some(codom)
          case Some(value) => None
          case None => None
        }
      }
      case None => None
    }
    case STZero => Some(STNat)
    case STSuc(t) => typeOf(env, t) match {
      case Some(value) if (value == STNat) => Some(STNat)
      case Some(value) => None
      case None => None
    }
    case STIsZero(t) => typeOf(env, t) match {
      case Some(value) if (value == STNat) => Some(STBool)
      case Some(value) => None
      case None => None
    }
    case STTrue => Some(STBool)
  }
}
```

```

    case STFalse => Some(STBool)
    case STTest(t1, t2, t3) => typeOf(env, t1) match {
      case Some(value1) if (value1 == STBool) => typeOf(env, t2) match {
        case Some(value2) => typeOf(env, t3) match {
          case Some(value3) if (value2 == value3) => Some(value2)
          case Some(value) => None
          case None => None
        }
        case Some(value) => None
        case None => None
      }
      case Some(value) => None
      case None => None
    }
  }

typeOf(List[STType](), term) match {
  case Some(value) => true
  case None => false
}
}

```

3.2 Ruby

In Ruby, instead of the Either[] implementation, we use Nil to represent our failed typerules. Each term contains the typeOf function that has been passed the environment. The main typecheck function resides in the STTerm class since every term has the same code for the initial "passing an empty environment" part and we can be less redundant this way:

```
class STTerm
  def typecheck
    (typeOf(Array.new)) ? true : false
  end
end

class STVar < STTerm
  def typeOf(env)
    (env.length > @index) ? env[@index] : nil
  end
end

class STApp < STTerm
  def typeOf(env)
    val1 = @t1.typeOf(env)
    val2 = @t2.typeOf(env)
    if (val1.is_a?(STFun))
      if (val2 == val1.dom)
        val1.codom
      else
        nil
      end
    else
      nil
    end
  end
end

class STAbs < STTerm
  def typeOf(env)
    newenv = env.unshift(@t)
    (@t1.typeOf(newenv)) ? STFun.new(@t, @t1.typeOf(newenv)) : nil
  end
end

class STZero < STTerm
  def typeOf(env); STNat.new end
end

class STSuc < STTerm
  def typeOf(env)
    (@t.typeOf(env) == STNat.new) ? STNat.new : nil
  end
end

class STIsZero < STTerm
  def typeOf(env)
    (@t.typeOf(env) == STNat.new) ? STBool.new : nil
  end
end
```



```

    end
end

class STTrue < STTerm
  def typeOf(env); STBool.new end
end

class STFalse < STTerm
  def typeOf(env); STBool.new end
end

class STTest < STTerm
  def typeOf(env)
    type1 = @t1.typeOf(env)
    type2 = @t2.typeOf(env)
    type3 = @t3.typeOf(env)
    if (type1 == STBool.new) && (type2 == type3)
      type2
    else
      nil
    end
  end
end
end

```

For the sake of clear documentation, I have shown the classes with their respective typeOf functions within. In reality, you would simply add the function into the already created classes from earlier.

4 Translation to untyped λ calculus

We will now translate our Simply typed terms into Untyped terms. This is done by removing the type for Abstraction and using the λ definitions for our other terms (True/False/Zero etc.) as well as combining them with other terms occasionally using STApp.

4.1 Scala

In Scala, this is simply implemented by a single function that maps each STTerm with it's corresponding ULTerm or representation using ULTerms:

```
def eraseTypes(term: STTerm): ULTerm = term match {  
  case STVar(index) => ULVar(index)  
  case STAbs(t, t1) => ULAbs(eraseTypes(t1))  
  case STApp(t1, t2) => ULApp(eraseTypes(t1), eraseTypes(t2))  
  case STZero => ULAbs(ULAbs(ULVar(0)))  
  case STSuc(t) => ULApp(ULAbs(ULAbs(ULAbs(ULApp(ULVar(1), ULApp(ULApp(ULVar(2), ULVar(1)),  
    ULVar(0)))))), eraseTypes(t))  
  case STIsZero(t) => ULApp(ULAbs(ULApp(ULApp(ULVar(0), ULAbs(eraseTypes(STFalse))),  
    eraseTypes(STTrue))), eraseTypes(t))  
  case STTrue => ULAbs(ULAbs(ULVar(1)))  
  case STFalse => ULAbs(ULAbs(ULVar(0)))  
  case STTest(t1, t2, t3) => ULApp(ULApp(ULApp(ULAbs(ULAbs(ULAbs(ULApp(ULApp(ULVar(2),  
    ULVar(1)), ULVar(0))))), eraseTypes(t1)), eraseTypes(t2)),  
    eraseTypes(t3))  
}
```

4.2 Ruby

In Ruby, we can simply define our function within each term class that returns the respective ULTerm translation:

```
class STVar < STTerm
  def eraseTypes
    ULVar.new(@index)
  end
end

class STApp < STTerm
  def eraseTypes
    ULApp.new(@t1.eraseTypes, @t2.eraseTypes)
  end
end

class STAbs < STTerm
  def eraseTypes
    ULAbs.new(@t1.eraseTypes)
  end
end

class STZero < STTerm
  def eraseTypes
    ULAbs.new(ULAbs.new(ULVar.new(0)))
  end
end

class STSuc < STTerm
  def eraseTypes
    ULApp.new(ULAbs.new(ULAbs.new(ULAbs.new(ULApp.new(ULVar.new(1), ULApp.new(
      ULApp.new(ULVar.new(2), ULVar.new(1), ULVar.new(0))))), @t.eraseTypes)
    end
  end
end

class STIsZero < STTerm
  def eraseTypes
    ULApp.new(ULAbs.new(ULApp.new(ULApp.new(ULVar.new(0), ULAbs.new(STFalse.eraseTypes)),
      STTrue.eraseTypes)), @t.eraseTypes)
  end
end

class STTrue < STTerm
  def eraseTypes
    ULAbs.new(ULAbs.new(ULVar.new(1)))
  end
end

class STFalse < STTerm
  def eraseTypes
    ULAbs.new(ULAbs.new(ULVar.new(0)))
  end
end
```

```

class STTest < STTerm
  def eraseTypes
    ULApp.new(ULApp.new(ULApp.new(ULAbs.new(ULAbs.new(ULAbs.new(ULApp.new(ULApp.new(
      ULVar.new(2), ULVar.new(1)), ULVar.new(0))))) , @t1.eraseTypes), @t2.eraseTypes),
      @t3.eraseTypes)
  end
end

```

For the sake of clear documentation, I have shown the classes with their respective `typeOf` functions within. In reality, you would simply add the function into the already created classes from earlier.