

Representation of Guarded Command Language

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1 Introduction

This documentation was created for Assignment 3 of 3MI3, it's purpose is to help understand the fundamentals behind how (and why) my program runs. Guarded Command Language is a language defined for predicate transformer semantics. We will attempt to implement this language using Ruby and Clojure.

2 Representation

Our representation builds on top of standard expressions like integers, variables, and operations of expressions. We also introduce testing expressions like True/False, and boolean operations on them. Lastly we add our statements, which include a skip, variable assignment, composition, an if loop and a do loop (both loops taking a list of test and statement tuples).

2.1 Ruby

For Ruby, we will classify our 3 main terms of expressions, tests and statements as classes:

```
class GCEExpr; end
class GCTest; end
class GCStmt; end
```

Then, we create classes for each subterm as subclasses that inherit their respective classes, note our skip, true and false classes take no arguments as they are base classes.

```
class GCConst < GCEExpr
  attr_reader :val

  def initialize(x)
    unless x.is_a?(Integer)
      throw "Constructing a constant out of a non integer value"
    end
    @val = x
  end
end

class GCVar < GCEExpr
  attr_reader :val

  def initialize(x)
    unless x.is_a?(Symbol)
      throw "Constructing a variable out of a non symbol value"
    end
    @val = x
  end
end

class GCOp < GCEExpr
  attr_reader :val1
  attr_reader :val2
  attr_reader :op

  def initialize(x, y, z)
    unless x.is_a?(GCEExpr) and y.is_a?(GCEExpr) and [:plus, :minus, :times, :div].include? z
      throw "Constructing an operation out of incorrect values"
    end
    @val1 = x
    @val2 = y
    @op = z
  end
end
```

```

class GCComp < GCTest
  attr_reader :val1
  attr_reader :val2
  attr_reader :op

  def initialize(x, y, z)
    unless x.is_a?(GCEExpr) and y.is_a?(GCEExpr) and [:eq, :less, :greater].include?(z)
      throw "Constructing a comparison out of incorrect values"
    end
    @val1 = x
    @val2 = y
    @op = z
  end
end

class GCAnd < GCTest
  attr_reader :val1
  attr_reader :val2

  def initialize(x, y)
    unless x.is_a?(GCTest) and y.is_a?(GCTest)
      throw "Constructing AND with non GCEExpr values"
    end
    @val1 = x
    @val2 = y
  end
end

class GCOrr < GCTest
  attr_reader :val1
  attr_reader :val2

  def initialize(x, y)
    unless x.is_a?(GCTest) and y.is_a?(GCTest)
      throw "Constructing OR with non GCEExpr values"
    end
    @val1 = x
    @val2 = y
  end
end

class GCTrue < GCTest
end

class GCFalse < GCTest
end

class GCSkip < GCStmt; end

class GCAssign < GCStmt
  attr_reader :var
  attr_reader :expr

  def initialize(x, y)

```

```

        unless x.is_a?(Symbol) and y.is_a?(GCEExpr)
          throw "Constructing an assignment out of incorrect values"
        end
        @var = x
        @expr = y
      end
    end

class GCCompose < GCStmt
  attr_reader :val1
  attr_reader :val2

  def initialize(x, y)
    unless x.is_a?(GCStmt) and y.is_a?(GCStmt)
      throw "Constructing a compose out of non GCStmt values"
    end
    @val1 = x
    @val2 = y
  end
end

class GCIf < GCStmt
  attr_reader :guards

  def initialize(x)
    unless checker?(x)
      throw "Constructing IF out of incorrect values"
    end
    @guards = x
  end

  def checker?(val)
    val.each do |v|
      unless v[0].is_a?(GCTest) and v[1].is_a?(GCStmt)
        return false
      end
    end
    return true
  end
end

class GCDo < GCStmt
  attr_reader :guards

  def initialize(x)
    unless checker?(x)
      throw "Constructing DO out of incorrect values"
    end
    @guards = x
  end

  def checker?(val)
    val.each do |v|
      unless v[0].is_a?(GCTest) and v[1].is_a?(GCStmt)

```

```

        return false
      end
    end
    return true
  end
end

```

2.2 Clojure

For Clojure, we define our subterms as records. We do not need to account for the main classes as we assume all inputs are well typed.

```

(defrecord GCConst [c])
(defrecord GCVar [v])
(defrecord GCOp [e1 e2 op])

(defrecord GCComp [e1 e2 op])
(defrecord GCAnd [t1 t2])
(defrecord GCOr [t1 t2])
(defrecord GCTrue [])
(defrecord GCFalse [])

(defrecord GCSkip [])
(defrecord GCAssign [v e])
(defrecord GCCompose [s1 s2])
(defrecord GCIf [guards])
(defrecord GCDo [guards])

```

3 Stack Machine

In Ruby, we define a stack machine implementation to carry out evaluation of statements. This stack machine comprises of a command stack where we store instructions to perform, a result stack to store temporary results and the memory which maps variables to integers.

```
def stackEval(commands, result, memory)

  if commands.length() == 0
    return memory
  end

  command = commands.shift
  case command
  when GCConst
    result.unshift(command.val)
    stackEval(commands, result, memory)
  when GCVar
    stackEval(commands, result.unshift(memory.call(command.val)), memory)
  when GCOp
    commands.unshift(command.op).unshift(command.val1).unshift(command.val2)
    stackEval(commands, result, memory)
  when GCComp
    commands.unshift(command.op).unshift(command.val1).unshift(command.val2)
    stackEval(commands, result, memory)
  when GCAnd
    commands.unshift(:and).unshift(command.val1).unshift(command.val2)
    stackEval(commands, result, memory)
  when GCOr
    commands.unshift(:or).unshift(command.val1).unshift(command.val2)
    stackEval(commands, result, memory)
  when GCTrue
    result.unshift(true)
    stackEval(commands, result, memory)
  when GCFalse
    result.unshift(false)
    stackEval(commands, result, memory)
  when GCSkip
    stackEval(commands, result, memory)
  when GCAssign
    commands.unshift(:assign).unshift(command.expr)
    result.unshift(command.var)
    stackEval(commands, result, memory)
  when GCCompose
    commands.unshift(command.val2).unshift(command.val1)
    stackEval(commands, result, memory)
  when GCIf
    allguards = command.guards
    result.unshift(:ifstop)
    commands.unshift(:if)
    allguards.each do |guard|
      commands.unshift(guard[1])
      commands.unshift(:ifguard)
      commands.unshift(guard[0])
    end
  end
end
```

```

    stackEval(commands, result, memory)
when GCD0
    allguards = command.guards
    result.unshift(:dostop)
    commands.unshift(command)
    commands.unshift(:do)
    allguards.each do |guard|
        commands.unshift(guard[1])
        commands.unshift(:doguard)
        commands.unshift(guard[0])
    end
    stackEval(commands, result, memory)
when :doguard
    bval = result.shift
    stmt = commands.shift
    if bval == true
        result.unshift(stmt)
    end
    stackEval(commands, result, memory)
when :do
    trueguards = []
    result.each do |guard|
        trueguards.push(guard)
        break if guard == :dostop
    end
    removestop = trueguards.pop
    if trueguards.length() > 0
        r = rand(0..trueguards.length()-1)
        commands.unshift(trueguards[r])
    else
        removedo = commands.shift
    end
    stackEval(commands, result, memory)
when :ifguard
    bval = result.shift
    stmt = commands.shift
    if bval == true
        result.unshift(stmt)
    end
    stackEval(commands, result, memory)
when :if
    trueguards = []
    result.each do |guard|
        trueguards.push(guard)
        break if guard == :ifstop
    end
    removestop = trueguards.pop
    if trueguards.length() > 0
        r = rand(0..trueguards.length()-1)
        commands.unshift(trueguards[r])
        stackEval(commands, result, memory)
    else
        stackEval(commands, result, memory)
    end
end

```

```

when :plus
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1+val2)
  stackEval(commands, result, memory)
when :times
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1*val2)
  stackEval(commands, result, memory)
when :minus
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1-val2)
  stackEval(commands, result, memory)
when :div
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1/val2)
  stackEval(commands, result, memory)
when :eq
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1==val2)
  stackEval(commands, result, memory)
when :less
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1<val2)
  stackEval(commands, result, memory)
when :greater
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1>val2)
  stackEval(commands, result, memory)
when :and
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1&&val2)
  stackEval(commands, result, memory)
when :or
  val1 = result.shift
  val2 = result.shift
  result.unshift(val1||val2)
  stackEval(commands, result, memory)
when :assign
  solution = result.shift
  variable = result.shift
  stackEval(commands, result, updateState(memory, variable, solution))
end
end

def emptyState
  lambda { |c|

```



```

        0
    }
end

def updateState(sigma, x, n)
    lambda { |c|
        if (c == x)
            n
        else
            sigma.call(c)
        end
    }
end

```

The explanation is very straightforward for most of the terms. Constants/True/False move directly to the result stack, variables are searched inside the memory for a mapping. Operators (integer and boolean) are a little different where they evaluate each individual argument and using an operational symbol, pop the results from the result stack to evaluate further and push back into the result stack. Assignments simply add onto the memory with the specified variable and expression. If and Do, however, can get a little complicated. I implemented If by first pushing into the result stack an if symbol and the test statement of each guard, if it evaluates to True we then pop that from the result stack and push the statement of this true guard. Once the list is traversed, we pop out all the true statements at which point we randomly select one and add it to the command stack to be evaluated. Do is similar to If, except it has the entire Do loop on the stack right after the loops returns a True (to be performed again) If none are true, the Do loop is removed from the stack.

4 Small step semantics

We will use Clojure to now reduce our terms by one step (small step) and return the statement as well as its resulting state memory.

```
(defn drop-nth [n coll]
  (concat
    (take n coll)
    (drop (inc n) coll)))

(defn emptyState []
  (fn [x] 0))

(defn updateState [sig x n]
  (fn [c] (if (= c x) n ((sig) c))))

(defn reduce [some]
  (let [command (.stmt some)]
    (cond
      (instance? GCVar command)
        (Config. ((.sig some) (.v command)) (.sig some))
      (instance? GCOp command)
        (if (instance? GCCConst (.e1 command))
          (if (instance? GCCConst (.e2 command))
            (cond
              (= :plus (.op command)) (Config. (+
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
              (= :minus (.op command)) (Config. (-
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
              (= :times (.op command)) (Config. (*
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
              (= :div (.op command)) (Config. (/
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
            )
          (let [newcommand (reduce (Config. (.e2 command) (.sig some)))]
            (Config. (GCOp. (.e1 command) (.stmt newcommand) (.op command))
              (.sig newcommand))
          ))
        (let [newercommand (reduce (Config. (.e1 command) (.sig some)))]
          (Config. (GCOp. (.stmt newercommand) (.e2 command) (.op command))
            (.sig newercommand))
          ))
      (instance? GCComp command)
        (if (instance? GCCConst (.e1 command))
          (if (instance? GCCConst (.e2 command))
            (cond
              (= :eq (.op command)) (Config. (=
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
              (= :less (.op command)) (Config. (<
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
              (= :greater (.op command)) (Config. (>
                (.c (.e1 command)) (.c (.e2 command))) (.sig some))
            )
          (let [newcommand (reduce (Config. (.e2 command) (.sig some)))]
            (Config. (GCOp. (.e1 command) (.stmt newcommand) (.op command))
              (.sig newcommand))
          ))
    ))
```

```

        (.sig newcommand))
    ))
    (let [newercommand (reduce (Config. (.e1 command) (.sig some)))]
      (Config. (GCOp. (.stmt newercommand) (.e2 command) (.op command))
        (.sig newercommand))
    ))
(instance? GCAnd command)
  (if (or (instance? GCTrue (.t1 command)) (instance? GCFalse (.t1 command)))
    (if (or (instance? GCTrue (.t2 command)) (instance? GCFalse (.t2 command)))
      (if (and (instance? GCTrue (.t1 command)) (instance? GCTrue (.t2 command)))
        (Config. (GCTrue.) (.sig some))
        (Config. (GCFalse.) (.sig some)))
      (let [newcommand (reduce (Config. (.t2 command) (.sig some)))]
        (Config. (GCAnd. (.t1 command) (.stmt newcommand)) (.sig newcommand))
      ))
    (let [newcommand (reduce (Config. (.t1 command) (.sig some)))]
      (Config. (GCAnd. (.stmt newcommand) (.t2 command))
        (.sig newcommand))
    ))
(instance? GCOr command)
  (if (or (instance? GCTrue (.t1 command)) (instance? GCFalse (.t1 command)))
    (if (or (instance? GCTrue (.t2 command)) (instance? GCFalse (.t2 command)))
      (if (or (instance? GCTrue (.t1 command)) (instance? GCTrue (.t2 command)))
        (Config. (GCTrue.) (.sig some))
        (Config. (GCFalse.) (.sig some)))
      (let [newcommand (reduce (Config. (.t2 command) (.sig some)))]
        (Config. (GCAnd. (.t1 command) (.stmt newcommand)) (.sig newcommand))
      ))
    (let [newcommand (reduce (Config. (.t1 command) (.sig some)))]
      (Config. (GCAnd. (.stmt newcommand) (.t2 command))
        (.sig newcommand))
    ))
(instance? GCAssign command)
  (if (instance? GCConst (.e command))
    (Config. (GCSkip.) (updateState (.sig some) (.v command) (.c (.e command))))
    (let [newcommand (reduce (Config. (.e command) (.sig some)))]
      (Config. (GCAssign. (.v command) (.stmt newcommand)) (.sig newcommand))
    ))
(instance? GCCompose command)
  (if (instance? GCSkip (.s1 command))
    (Config. (.s2 command) (.sig some))
    (let [newcommand (reduce (Config. (.s1 command) (.sig some)))]
      (Config. (GCCompose. (.stmt newcommand) (.s2 command)) (.sig newcommand))
    ))
(instance? GCIf command)
  (if (= 0 (count (.guards command)))
    (Config. (GCSkip.) (.sig some))
    (let [random (rand-int (count (.guards command)))]
      (if (instance? GCTrue (get (get (.guards command) random) 0))
        (Config. (get (get (.guards command) random) 1) (.sig some))
        (if (instance? GCFalse (get (get (.guards command) random) 0))
          (Config. (GCIf. (drop-nth random (.guards command))) (.sig some))
          (let [newcommand (reduce (Config. (get (get (.guards command)
            random) 0) (.sig some)))]

```

```

        (Config. (assoc-in (.guards command) [random 0]
        (.stmt newcommand)) (.sig newcommand))
    )))
  ))
  (instance? GCDo command)
    (if (= 0 (count (.guards command)))
      (Config. (GCSkip.) (.sig some))
      (let [allguards (.guards command)
            newguards (.guards command)
            length (count allguards)
            ifguards (loop [x 0]
                          (when (< x length)
                            (assoc-in (newguards) [x 1] (GCCompose.
                                                             (get (get allguards x) 1) command))
                            (recur (+ x 1))
                          ))])
        (Config. (GCIf. ifguards) (.sig some))
      ))
  )))

```

Most of our code simply checks for the existence of a constant integer or a skip. If it exists for both arguments, the whole expression is evaluated, otherwise the non constant is reduced. If and Do is a little tricky however, If chooses a random guard from the list of guards and checks if it's test is True/False. If true, reduces the statement and if false, removes the guard from the list of guards and repeats. If it is neither true or false, it reduces the test and redoes the If with the new reduced test. If we traverse the entire list without encountering any true guard, we simply reduce down to a skip. Do is similar (though it has been a trouble to implement in Clojure) and we can utilize the If command as well. We simply take the list of guards and create a new If loop of those guards. However, for each guard statement we compose the entire while loop at the end of it. This way it keeps running the true guard statement until no true guards are left at which point it reduces to skip.

5 Big step semantics

We start off by writing a wellscoped method that checks in the environment of local+global variables to see whether our input is well scoped or not.

```
def wellScoped(program)
  environment = program.globals
  helper(program.stmt, environment)
end

def helper(prog, env)
  case prog
  when GCSkip
    true
  when GCLocal
    helper(prog.stmt, env.push(prog.var))
  when GCAssign
    env.include?(prog.var) && helper(prog.expr, env)
  when GCCompose
    helper(prog.val1, env) && helper(prog.val2, env)
  when GCIf
    prog.each do |guard|
      helper(guard[0], env) && helper(guard[1], env)
    end
  when GCDo
    prog.each do |guard|
      helper(guard[0], env) && helper(guard[1], env)
    end
  when GCComp
    helper(prog.val1, env) && helper(prog.val2, env)
  when GCAnd
    helper(prog.val1, env) && helper(prog.val2, env)
  when GCOr
    helper(prog.val1, env) && helper(prog.val2, env)
  when GCTrue
    true
  when GCFalse
    true
  when GCConst
    true
  when GCVar
    env.include?(prog.val)
  when GCOp
    helper(prog.val1, env) && helper(prog.val2, env)
  end
end
```

We start with a list of global variables and work our way through each case by checking its arguments for well scoped. If we encounter a local variable, we reset this local variables for the purpose of checking the scope of its arguments, we then reset it to its original value after we are done so as to not clash with a possible global variable of the same name.

We then define an eval method to evaluate our input and return the memory state that maps variable to integers. We take the help of a reduce function that evaluates each term using the memory state (this is the big step of the semantics) to help us with simpler evaluation.

```

def eval(program)
  environment = Hash.new do |hash, key|
    raise("Accessing a variable that doesn't exist!")
  end
  program.globals.each do |guard|
    environment[guard] = "something"
  end
  end_env = helper2(program.stmt, environment)
end

def helper2(prog, env)
  case prog
  when GCSkip
    return env
  when GCAssign
    variable = prog.var
    value = reduce(prog.expr, env)
    env[variable] = value
    return env
  when GCCompose
    res1 = helper2(prog.val1, env)
    helper2(prog.val2, res1)
  when GCLocal
    if env.include?(prog.var)
      tmp = env[prog.var]
      env[prog.var] = "no_value"
      newenv = helper2(prog.stmt, env)
      newenv[prog.var] = tmp
      return newenv
    else
      env[prog.var] = "no_value"
      newenv = helper2(prog.stmt, env)
      newenv.delete(prog.var)
      return newenv
    end
  when GCIf
    trueguards = []
    allguards = prog.guards
    allguards.each do |guard|
      if reduce(guard[0], env) == true
        trueguards.push(guard[1])
      end
    end
    if trueguards.length() > 0
      helper2(trueguards.sample, env)
    else
      return env
    end
  when GCDo
    trueguards = []
    allguards = prog.guards
    allguards.each do |guard|
      if reduce(guard[0], env) == true
        trueguards.push(guard[1])
      end
    end
    if trueguards.length() > 0
      helper2(trueguards.sample, env)
    else
      return env
    end
  end
end

```

```

        end
    end
    if trueguards.length() > 0
        newenv = helper2(trueguards.sample, env)
        recreate = GCD.new(allguards)
        return helper2(recreate, newenv)
    else
        return env
    end
end
end

def reduce(prog, env)
  case prog
  when GConst
    prog.val
  when GCVar
    env[prog.val]
  when GCOp
    case prog.op
    when :plus
      reduce(prog.val1, env) + reduce(prog.val2, env)
    when :minus
      reduce(prog.val1, env) - reduce(prog.val2, env)
    when :times
      reduce(prog.val1, env) * reduce(prog.val2, env)
    when :div
      reduce(prog.val1, env) / reduce(prog.val2, env)
    end
  when GCComp
    case prog.op
    when :eq
      reduce(prog.val1, env) == reduce(prog.val2, env)
    when :less
      reduce(prog.val1, env) < reduce(prog.val2, env)
    when :greater
      reduce(prog.val1, env) > reduce(prog.val2, env)
    end
  when GCAnd
    reduce(prog.val1, env) && reduce(prog.val2, env)
  when GCOr
    reduce(prog.val1, env) || reduce(prog.val2, env)
  when GCTrue
    true
  when GCFalse
    false
  end
end
end

```

The working of the method is very similar to previous functions we have defined with the only tricky case being the existence of a local variable with the same symbol as a global variable. I have defined the method to prevent clashes and incorrect overrides to occur.