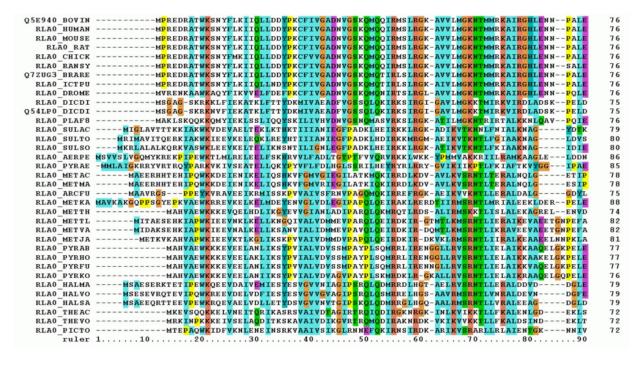
TOPIC 4: Sequence alignment

Learning Outcomes

- Be able to define the two main methods of alignment.
- Understand the two main algorithms for NGS alignment, including strengths and weaknesses.
- Be able to read SAM format

Sequence alignment

 Sequence alignment is a way of arranging the sequences of DNA, RNA, or protein to identify regions of similarity that may be a consequence of functional, structural, or evolutionary relationships between the sequences.



A multiple alignment of protein sequences

Pairwise alignment

- Alignment of two sequences is a relatively straightforward computational problem, but...
 - there are many possible alignments
 - there can be a very large reference
- NOTE: Two sequences can always be aligned and there can be more than one optimal solution

Methods of alignment

- By hand
- Mathematical approach
 - -Dynamic programming (slow, but optimal)
- Heuristic methods (fast, but approximate)
 - -BLAST, short read aligners

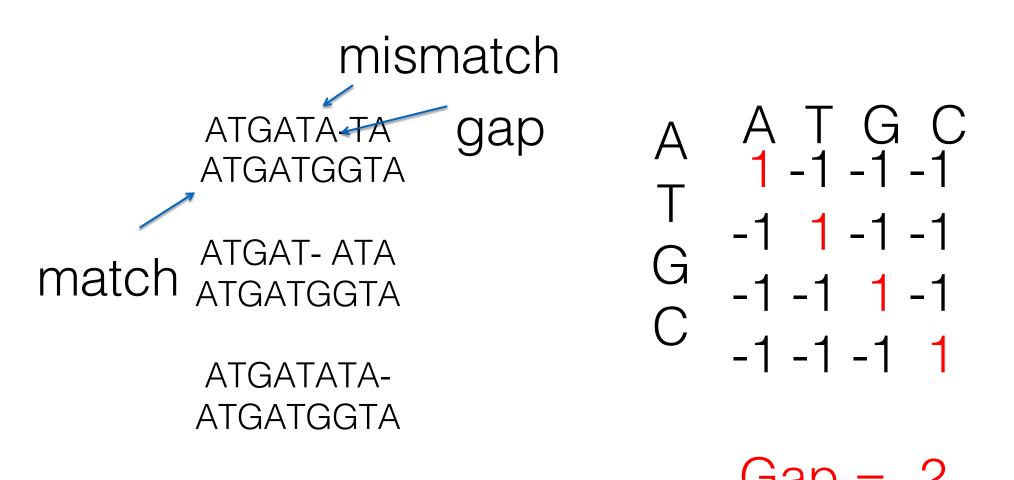
Think-pair-share

How could you quantify the quality of these DNA alignments?

Example: ATGATA-TA ATGATGGTA

ATGAT- ATA ATGATGGTA

ATGATATA-ATGATGGTA scoring matrix: a table of values that describes the probability of a biologically meaningful amino-acid or nucleotide residue-pair occurring in an alignment.



Scoring methods

- Scoring systems:
 - Each symbol pairing is assigned a numerical value, based on a symbol comparison table.
 - nucleotides
 - amino acids (PAM, BLOSUM)
- Gap penalties:
 - Opening: The cost of introducing a gap.
 - Extension: The cost to elongate a gap.

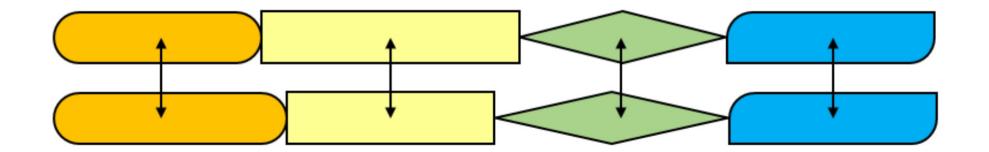
```
BLOSUM62
                           Positive for chemically similar substitution
0 -3 -3 -3
   1 0 0 -3
                              Common amino acids have low weights
                                    Rare amino acids have high weights
-3 -3 -4 -4 -2 -2 -3 -2 -2 -3 -2 -3 -1
ARNDCQEGHILKM
```

Gap penalties

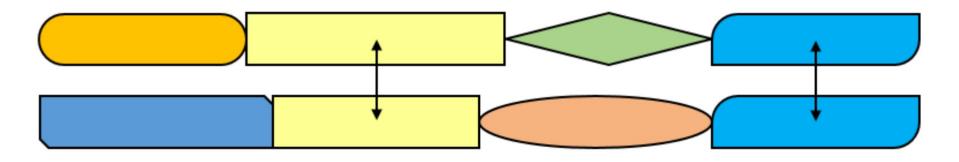
- Too little gap penalty gives nonsense nonhomologous alignments.
- Gaps are common, so too high gap penalty removes real alignments.
- "Affine" gap penalty has a large penalty to introduce a gap and a smaller penalty to extend one.

Dynamic programming

- Dynamic programming is a general programming technique.
- It structures a large search space into a succession of stages
 - The initial stage contains trivial solutions to sub-problems
 - Each partial solution in a later stage can be calculated by recurring a fixed number of partial solutions in an earlier stage
 - The final stage contains the overall solution



Global Alignment



Local Alignment

Global vs Local alignments

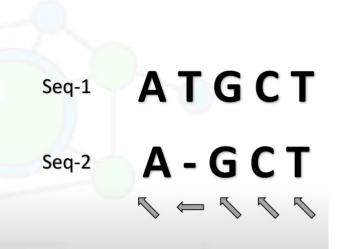
- Global alignment algorithms start at the beginning of two sequences and add gaps to each until the end of one is reached (Needleman-Wunsch).
- Local alignment algorithms finds the region (or regions) of highest similarity between two sequences and build the alignment outward from there (Smith-Waterman).
- Demo of the two algorithms: https://gtuckerkellogg.github.io/pairwise/demo/
- Video demonstrating exactly how to solve the Needleman-Wunsch alignment by hand https://www.youtube.com/watch?v=ipp-pNRIp4g

Basic principles of dynamic programming

There are too many comparisons to try them all so instead:

- Build alignment path matrix
- Stepwise calculation of score values
- Backtracking (evaluation of optimal path)

		Α	T	G	С	Т
	0	-2	-4	-6	-8	-10
Α	-2	1	-1	-3	-5	-7
G	-4	-1	0	0	-2	-4
С	-6	-3	-2	-1	1	-1
Т	-8	-5	-2	-3	-1	2



Basic Local Alignment Search Tool (BLAST)

Used for finding alignments between a small number of sequences and a large database

query

vs database

BLAST SEARCH							
Paste your sequence in FASTA format in the field provided							
>sp P07550 ADRB2_HUMAN Beta-2 adrenergic receptor OS=Homo sapiens MGQPGNGSAFLLAPNGSHAPDHDVTQQRDEVWVVGMGIVMSLIVLAIVFGNVLVITAIAK FERLQTVINYFITSLACADLVMGLAVVPFGAAHILMKMWTFGNFWCEFWTSIDVLCVTAS IETLCVIAVDRYFAITSPFKYQSLLTKNKARVIILMVWIVSGLTSFLPIQMHWYRATHQE AINCYANETCCDFFTNQAYAIASSIVSFYVPLVIMVFVYSRVFQEAKRQLQKIDKSEGRF HVQNLSQVEQDGRTGHGLRRSSKFCLKEHKALKTLGIIMGTFTLCWLPFFIVNIVHVIQD NLIRKEVYILLNWIGYVNSGFNPLIYCRSPDFRIAFQELLCLRRSSLKAYGNGYSSNGNT GEQSGYHVEQEKENKLLCEDLPGTEDFVGHQGTVPSDNIDSPGRNCSTNDSLL							
Specify an e-value cutoff level to use in your query: 1.0							
Submit query Clear fields							

Examples: nr = non-redundant Swiss Prot Model organisms (TAIR)

Your favorite sequences

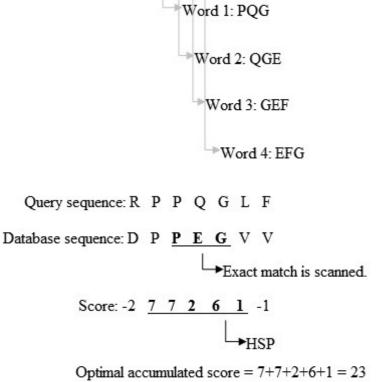
http://www.ncbi.nlm.nih.gov/books/NBK21097/

BLAST

BLAST searches for regions of similarity and doesn't try to align the entire sequence (local alignment). A global alignment aligns two sequences over the entire length

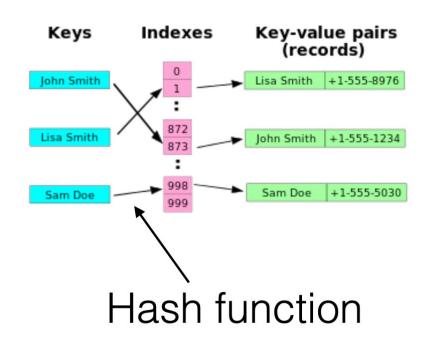
Designed to identify homologous sequences.

Hashed seed-extend algorithm



Hashed seed-extend algorithm

- A "hash" is a structure used in computer programming
- It is a way of storing information in a look-up table
- Allows efficient searching



Step 1. Divide up your query sequence into words (typically 11 for DNA and 3 for proteins)

query sequence list of words LIAWHCMPNAAA LIA IAW AWH WHC

Step 2. Find sequences in the database that have matches to query words with a score greater than your threshold

query sequence LIAWHCMPNAAA

query word WHC

```
database sequence 2 ....WIC... database sequence 3 ....AIC...
```

What is the score for each of these alignments (use the BLOSUM62 scoring matrix)?

Step 3. If a match is found extend the alignment around it in each direction for as long as the cumulative score is above a threshold (e.g. 0) or if it falls off by more than X from its maximum. (or you get to the end of the sequence)

This extended match is called a High Scoring Pair (HSP)

query sequence

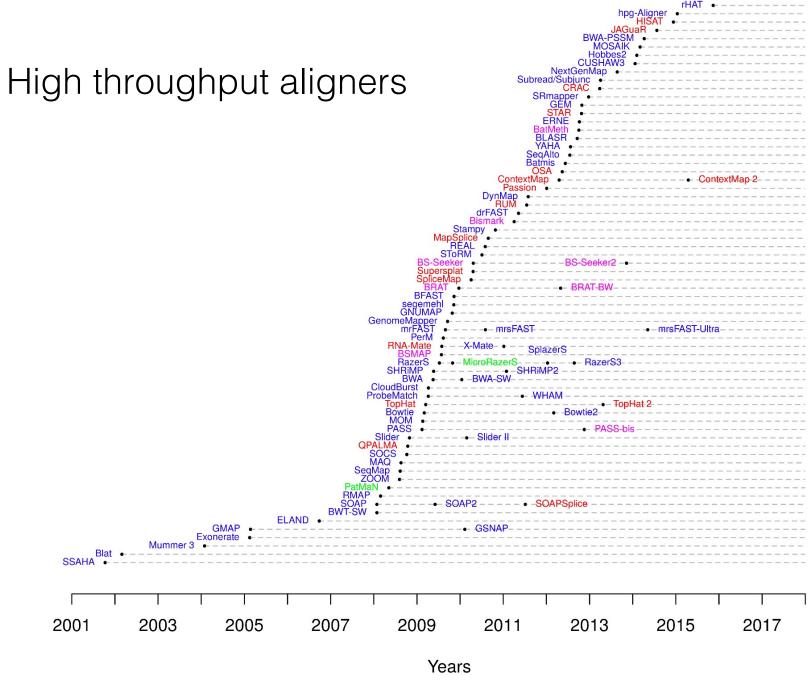
LIAWHCMPNAAA

database sequence 1 IIAWHCMPNDDA



BLAST

- Why not use BLAST for short read data (e.g., aligning short reads to a reference)?
 - Typically takes 0.1 to 1 second to search 1 sequence against a database
 - 60 million reads equates to 70 CPU days

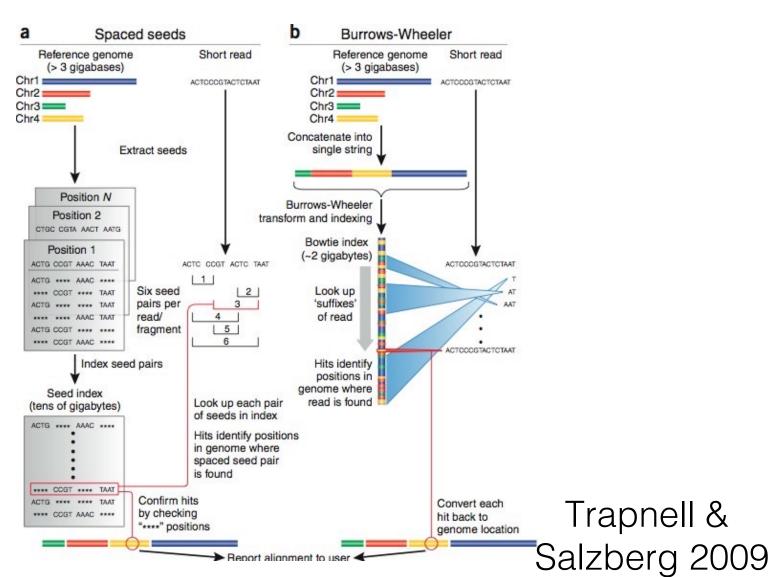


http://www.ebi.ac.uk/~nf/hts_mappers/

Short read alignment is hard

- Billions of short sequences aligned to a very long reference
- Short reads contain less information and are less likely to have a unique mapping location

Approaches to align short reads



Trapnell &

Hashed seed-extend algorithms

- Two step process:
 - Identify a match to the seed sequence in the reference
 - Extend match using sensitive (but slow) Smith-Waterman algorithm

Reference sequence:

...GATCTCGATCGATGATCGTAGGATTGATCAGCTA...

Short read:

TCGATCGATGATCGAAGGATTGATCAG

Reference sequence:

...GATCTCGATCGATGATCGTAGGATTGATCAGCTA...

Short read:

TCGATCGAT GATCGAAGG ATTGATCAG

9bp seed 9bp seed 9bp seed

The algorithm will try to match each seed to the reference. If there is a match with any seed, it performs a local alignment

Reference sequence:

<u>seed</u> ->Extend with Smith-Waterman->

...GATCTCGATCGATGATCGTAGGATTGATCAGCTA...

TCGATCGATGATCGAAGGATTGATCAG

Short read:

TCGATCGAT GATCGAAGG ATTGATCAG

9bp seed 9bp seed 9bp seed

Here there is a match with at least one seed

Reference sequence:

...GATCTCGATCGATGATCGTAGGATTGATCAGCTA...

Short read:

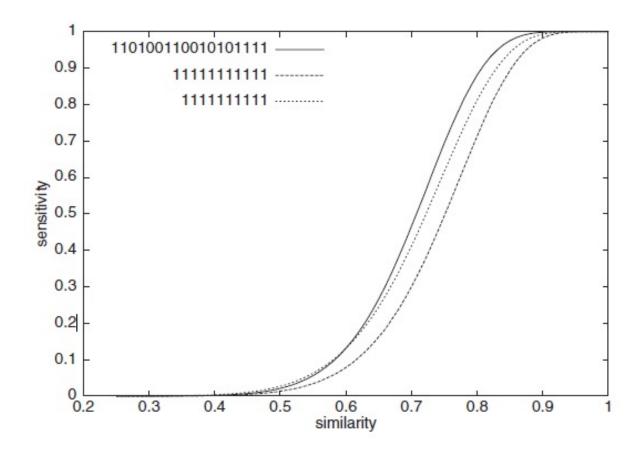
TAGATCGAT GATCGAAGG ATTGAGCAG

9bp seed 9bp seed 9bp seed

With three sequencing errors/SNPs, there can be no matches

Spaced seeds

To increase sensitivity we can used spaced-seeds:



Spaced seeds

To increase sensitivity we can used spaced-seeds:

111111111 GATAGCTAGCTAAT AGCTAGCTA

Consecutive seed template with length 9bp

Reference

Query

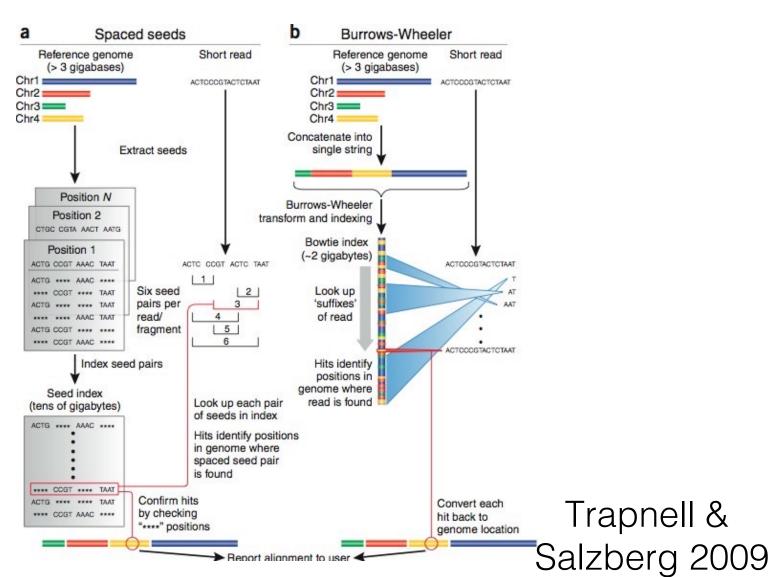
10101101011011
GATAGCTAGCTAAT
GATAGCGAGCTAAT

Consecutive seed template with weight 9bp

Reference

Query

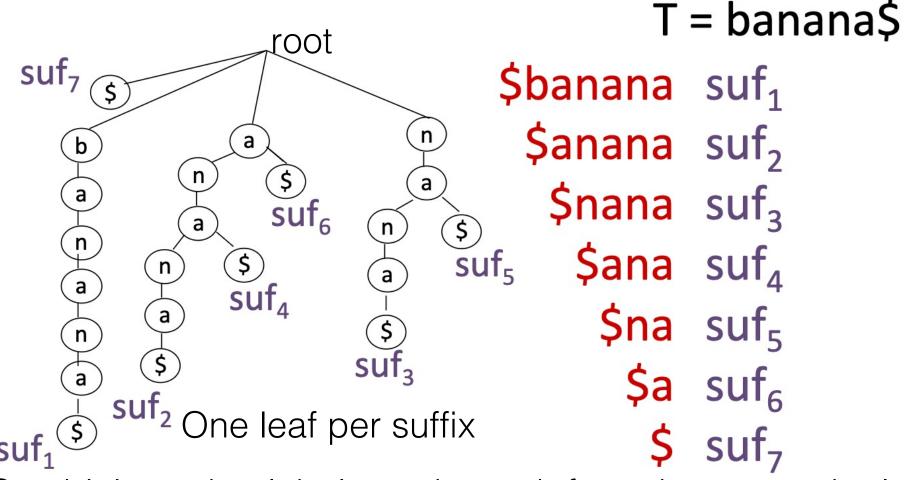
Approaches to align short reads



Trapnell &

Suffix-Trie

A data structure that contains all suffixes and their locations in the text



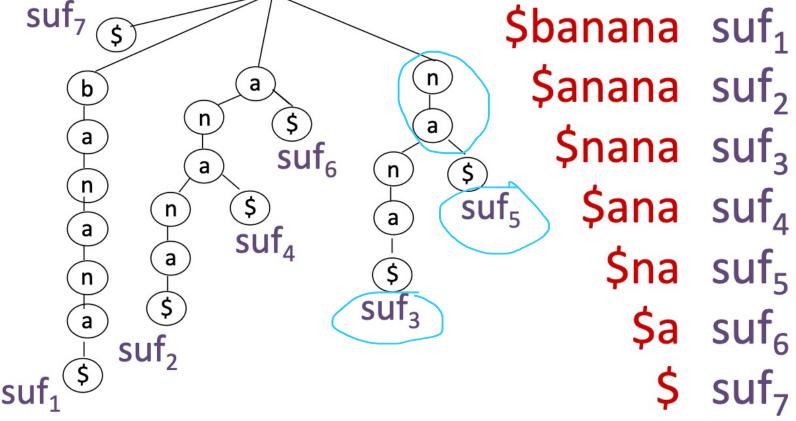
Combining edge labels on the path from the root to the leaf *i* spells out the suffix that begins at position *i*

Suffix-Trie

Search "na"

A data structure that contains all suffixes and their locations in the text

T = banana\$



Suffix tries allow particularly fast implementations of many important string operations (eg. search for substrings quickly)

Try out this website https://visualgo.net/en/suffixtree

Suffix-Prefix Trie

- A family of methods which uses a Trie structure to search a reference sequence (e.g. Bowtie, BWA, SOAP2)
- Trie data structure which stores the suffixes (i.e. ends of a sequence)
- Key advantage over hashed algorithms:
 - Alignment of multiple copies of an identical sequence in the reference only needs to be done once
 - Use of an FM-Index to store Trie can drastically reduce storage/memory requirements (e.g. Human genome can be stored in 2Gb of RAM)
 - Burrows Wheeler Transform to perform fast lookups

More information about the BWT: https://www.youtube.com/watch?v=4n7NPk5lwbl

Burrows-Wheeler Transform (BWT)

- Encodes data so that it is easier to compress
- BWT be reversed to recover the original word

\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana	a
1	anana\$	a\$	5	3	a \$banan	n
2	nana\$	ana\$	3	6	ana\$ban	n
3	ana\$	anana\$	1	2	anana\$b	b
4	na\$	banana\$	0	5	banana\$	\$
5	a\$	na\$	4	1	na\$ bana	a
6	\$	nana\$	2	0	nana\$ba	a

- Encodes data so that it is easier to compress
- Can be reversed to recover the original word

\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana	a
1	anana\$	a\$	5	3	a\$ banan	n
2	nana\$	ana\$	3	6	ana\$ban	n
3	ana\$	anana\$	1	2	anana\$b	b
4	na\$	banana\$	0	5	banana\$	\$
5	a\$	na\$	4	1	na\$ bana	a
6	\$	nana\$	2	0	nana\$ba	a

- Encodes data so that it is easier to compress
- Can be reversed to recover the original word

\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana	a
1	anana\$	₁ a\$	5	3	a\$ banan₁	n
2	nana\$	2ana\$	3	6	ana\$ban ₂	n
3	ana\$	₃anana\$	1	2	anana\$b₁	b
4	na\$	₁banana\$	0	5	banana\$	\$
5	a\$	₁ na\$	4	1	na\$ bana ₂	a
6	\$	2nana\$	2	0	nana\$ba ₃	a

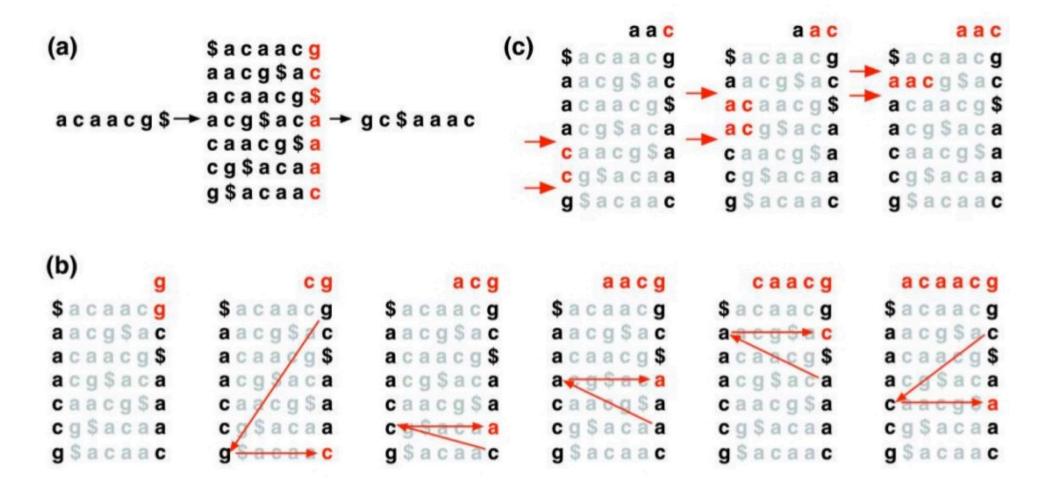
- Encodes data so that it is easier to compress
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\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana ₁	a
1	anana\$	₁ a\$	5	3	a\$banan ₁	n
2	nana\$	2ana\$	3	6	ana\$ban ₂	n
3	ana\$	₃anana\$	1	2	anana\$b₁	b
4	na\$	₁banana\$	0	5	banana\$	\$
5	a\$	₁ na\$	4	1	na \$bana ₂	a
6	\$	2nana\$	2	0	nana\$ba ₃	a

- Encodes data so that it is easier to compress
- Can be reversed to recover the original word

\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana ₁	a
1	anana\$	₁ a\$	5	3	a\$banan ₁	n
2	nana\$	2ana\$	3	6	ana\$ban ₂	n
3	ana\$	3anana\$	1	2	anana\$b ₁	b
4	na\$	₁banana\$	0	5	banana\$	\$
5	a\$	₁ na\$	4	1	ma\$bana ₂	a
6	\$	2nana\$	2	0	nana\$ba ₃	a

Burrows-Wheeler Transform



Suffix-Prefix Trie

- Less sensitive for sequences that are more divergent from the reference.
 - Sequencing errors
 - Query Reference differences

Comparison

Hash referenced spaced seeds (NextGenMap)

- Requires more RAM
- Runs slower
- Simpler to program
- More sensitive

Suffix/Prefix Trie (BWA)

- Requires less RAM
- Runs much faster
- Complicated to program
- Less sensitive

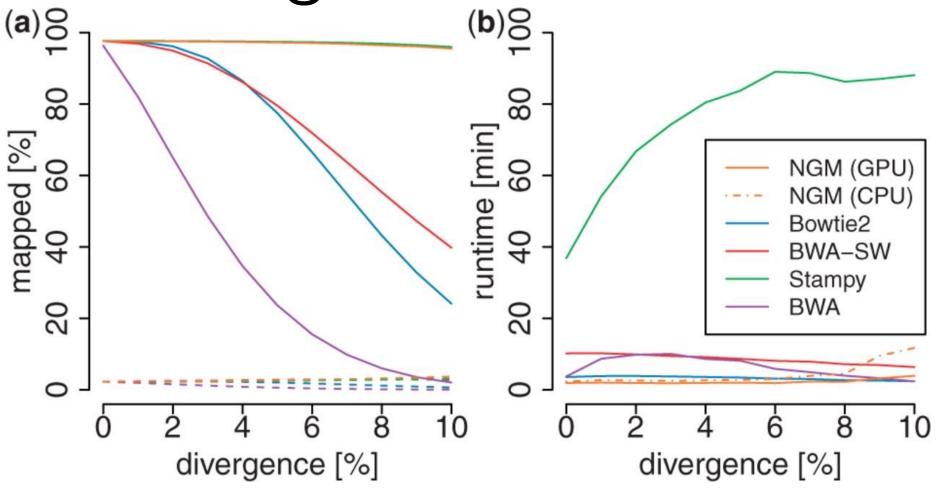
Popular short read aligners

Program	Algorithm	Speed	Accuracy in for divergent sequences	
Bowtie2	Suffix/Prefix	Very fast	Low	
BWA	Suffix/Prefix	Fast	Medium	
Stampy	Hashing ref	Slow	High	
Soap2	Suffix/Prefix	Fast	Low	
Novoalign	Hashing ref	Slow	High	
NextGenMap	Hashing ref	Med	High	

Think-Pair-Share

- Third generation sequencing can produce very long reads (10-50 Kbp), but are very error prone (~5-10% errors)
- Why would suffix-trie based aligners do poorly with this data?

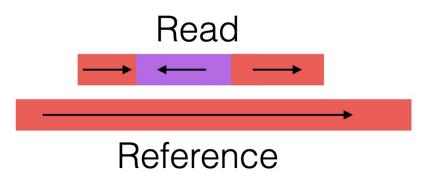
Alignment stats



*From NextGenMap paper

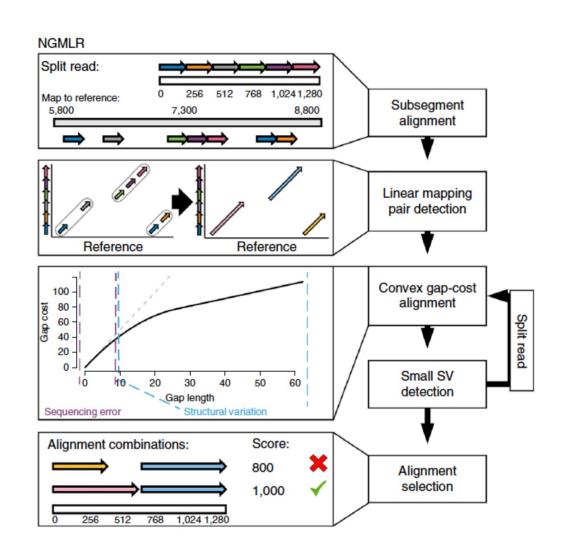
Long read alignment

- Long reads might contain structural variation that makes it hard to form a linear alignment
- For example, a read containing a large inversion would contain 3 linear alignments
- Most long read technologies have higher error rates



Long read alignment

- Find exact matches between read fragment and reference
- 2. Look for chains of matches
- 3. Use local alignment of read to best reference region.



Long read alignment

- Longer reads have more information, but more error.
- Example: NGMLR uses k-mers (short strings of a specific length) to pick region and smith-waterman for exact placement.
- Other programs:
 - KART, BWA-MEM, BLASR, minimap2

Alignment choice

- Speed needed?
- How divergent is sequence from reference? Same species or relative?
- How much variation in your samples?
- Genome size of reference?

Other considerations

- PCR duplicates
- Multi-mapping reads
- Spliced-read mapping

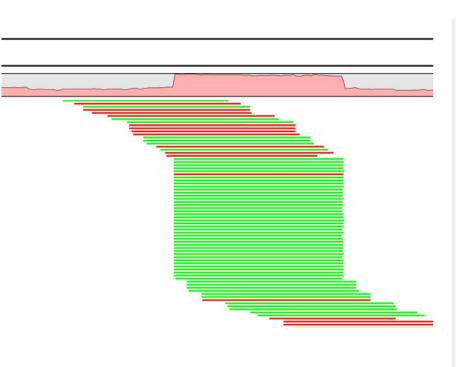
PCR duplicates

Most library preps have at least one PCR amplification step

PCR can introduce errors and then sequencing multiple copies makes it seem like a real SNP

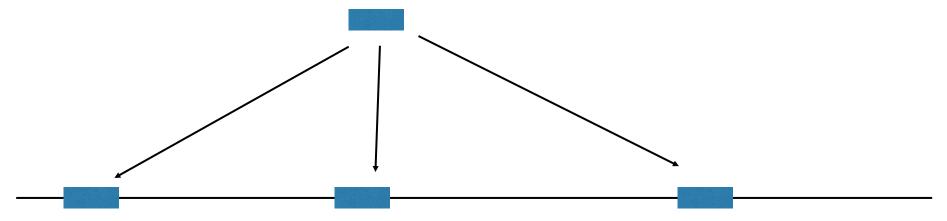
SAMtools and Picard can flag or remove these duplicates based on alignment location

- Samples with same start and stop position are considered duplicates
- Don't flag duplicates for GBS (set start and stop)



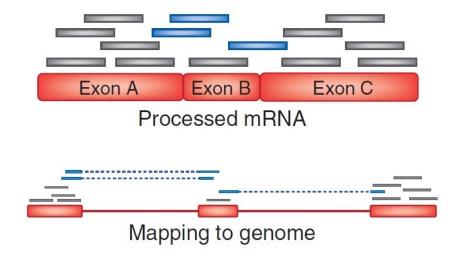
From Qiagen website

Multiple mapping reads



- A single read may occur more than once in a reference genome, due to gene/chromosome duplication or repetitive elements
- Reads may be assigned to one random location
- Affects mapping quality

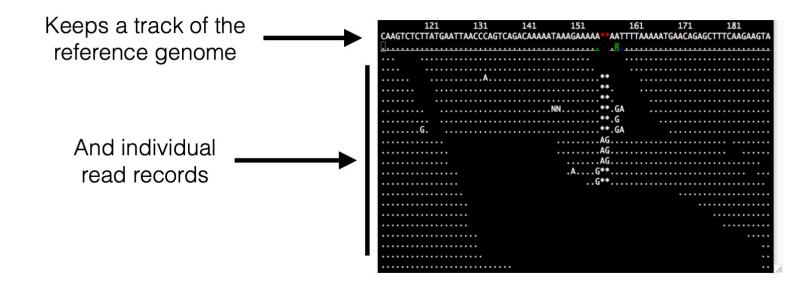
Spliced-read mapping



- Need to account for splicing
- Examples: TopHat, SubRead, Star

SAM (BAM) format

- Sequence Alignment/Map format
 - Universal standard.
 - Generally aligned to reference, but not necessarily
 - Human-readable (SAM) and compressed (BAM) forms



SAM format

Header (lines begin with @)

VN:1.5 GO:none SO:coordinate

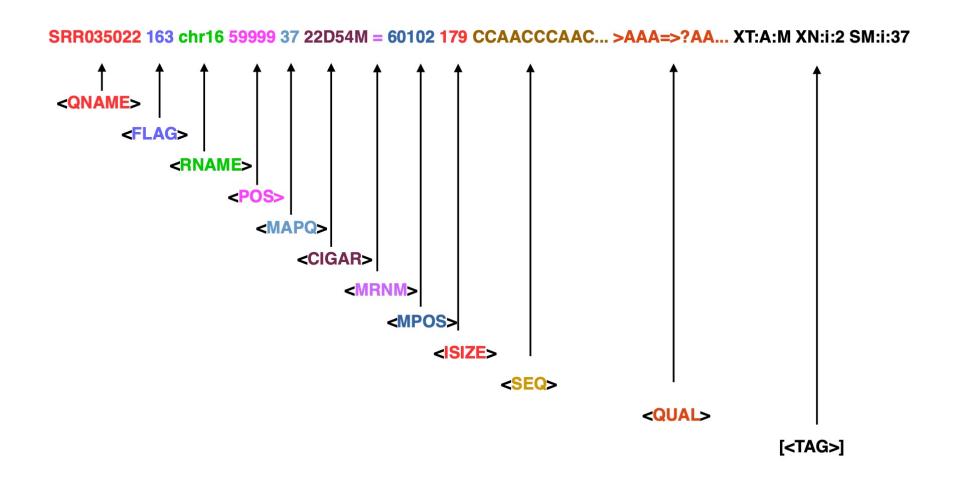
@HD

```
@SQ
      SN:cp ai 88656873
                         LN:151104
                                           Sort order
      SN:mt gi 571031384
@SQ
                         LN:300945
@SQ
      SN:rDNA gi 563582565
                          LN:9814
@SQ
      SN:Ha1 LN:175985764
@SQ
      SN:Ha2 LN:209013747
@SQ
      SN:Ha3 LN:203472901
@SQ
      SN:Ha4 LN:216026857
                                   Reference sequence name (SN) and
@SQ
      SN:Ha5 LN:271056985
@SQ
      SN:Ha6 | N:100519666
@SQ
      SN:Ha7 LN:109221022
                                    length (LN) e.g. Chromosome Ha7,
@SQ
      SN:Ha8 | N:192129815
@SQ
      SN:Ha9 LN:253478808
                                          which is 109,221,022bp long
@SO
      SN:Ha10 I N:327788049
@SQ
      SN:Ha11 LN:208730832
@SQ
      SN:Ha12 LN:208068730
                                       Read group information (@RG)
@SQ
      SN:Ha13 LN:239367298
@SQ
      SN:Ha14 LN:230295834
@SQ
      SN:Ha15 LN:202246870
@SO
      SN:Ha16 I N:226777971
@SQ
      SN:Ha17 LN:267415242
@SO
      SN:Ha0 73Ns
                  LN:359367108
@RG
      ID:HI.2034.006.Index_18.W70_NHK_2013_5 LB:Anomalus PL:ILLUMINA SM:HI.2034.006.Index_18.W70_NHK_2013_5 PU:Anomalus
@PG
      ID:ngm PN:ngm CL:" --affine 0 --argos min score 0 --bam 1 --block multiplier 2 --bs cutoff 6 --bs mapping 0 --cpu threads 11 --dualstrand 1
                 PN:ngm CL:" --affine 0 --argos_min_score 0 --bam 1 --block_multiplier 2 --bs_cutoff 6 --bs_mapping 0 --cpu_threads 11 --
@PG
      ID:ngm.1
@PG
      ID:ngm.2
                 PN:ngm CL:" --affine 0 --argos min score 0 --bam 1 --block multiplier 2 --bs cutoff 6 --bs mapping 0 --cpu threads 11 --
```

Program (@PG) information - what you used to map the reads

SAM format

Alignment lines



SAM format

Alignment lines

SRR035022 163 chr16 59999 37 22D54M = 60102 179 CCAACCCAAC... >AAA=>?AA... XT:A:M XN:i:2 SM:i:37

```
<QNAME> Query name - i.e. the name of the read
```

<FLAG> A combination of bitwise flags that indicate properties of the alignment (complement, strand etc.)

<RNAME> Reference sequence name

<POS> Position in the reference (1-based)

<MAPQ> Mapping quality

<CIGAR> Concise Idiosyncratic Gapped Alignment Report (CIGAR) string - tells you about gaps in the alignment

<MRNM> Read-mate reference sequence

<MPOS> Read-mate position in the reference

<ISIZE> Insert size

<SEQ> The segment's sequence

<QUAL> An ASCII sequence containing quality information for each base in the sequence

[<TAG>] These are optional tags that get added and contain user specified data (For example, SM is the mapping quality of this sequence only - ignoring the read mate)

Mapping Quality

- $MapQ = Qs = -10 log_{10}(P)$
- P = probability that this mapping is NOT the correct one
- MapQ = 0 = equally likely to map somewhere else
- Different programs use different formulas for P
- A value of 255 indicates that the information is not available

Background reading

- Trapnell, C., & Salzberg, S. L. (2009). How to map billions of short reads onto genomes. *Nature* biotechnology, 27(5), 455-457.
- Reinert, K., Langmead, B., Weese, D., & Evers, D.
 J. (2015). Alignment of next-generation sequencing
 reads. *Annual review of genomics and human genetics*, 16, 133-151.

All available on the GitHub

Burrows-Wheeler Algorithm

- Encodes data so that it is easier to compress
- Can be reversed to recover the original word

\overline{i}	Suffix	Sorted Suffix	SA[i]	ISA[i]	Sorted Rotations	BWT[i]
0	banana\$	\$	6	4	\$banana ₁	a
$\bigcirc 1$	anana\$	a\$	5	3	a \$banan₁	n
2	nana\$	ana\$)	3	6	ana\$ban ₂	n
3	ana\$	3 anana\$		2	anana\$b₁	b
4	na\$	banana\$	0	5	banana\$	\$
5	a\$	na\$	-(4)	1	na\$ bana ₂	a
6	\$	2 nana\$	2	0	nana\$ba ₃	a