## Numerical Analysis: Midterm

(30 marks, only the 3 best questions count)

## Urbain Vaes

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**Question 1** (Floating point arithmetic, 10 marks). True or false? (+1/0/-1)

- 1. Let  $(\bullet)_2$  denote binary representation. It holds that  $(0.1011)_2 + (0.0101)_2 = 1$ .
- **2.** Let  $(\bullet)_3$  denote base 3 representation. It holds that  $(1000)_3 \times (0.002)_3 = 2$ .
- **3.** A natural number with binary representation  $(b_4b_3b_2b_1b_0)_2$  is even if and only if  $b_0 = 0$ .
- 4. In Julia, Float64(.4) == Float32(.4) evaluates to true.
- **5.** Machine addition  $\widehat{+}$  is a commutative operation. More precisely, given any two double-precision floating point numbers  $x \in \mathbf{F}_{64}$  and  $y \in \mathbf{F}_{64}$ , it holds that x + y = y + x.
- **6.** Let  $\mathbf{F}_{32}$  and  $\mathbf{F}_{64}$  denote respectively the sets of single and double precision floating point numbers. It holds that  $\mathbf{F}_{32} \subset \mathbf{F}_{64}$ .
- 7. The machine epsilon of a floating point format is the smallest strictly positive number that can be represented exactly in the format.
- 8. Let  $\mathbf{F}_{64}$  denote the set of double precision floating point numbers. For any  $x \in \mathbf{R}$  such that  $x \in \mathbf{F}_{64}$ , it holds that  $x + 1 \in \mathbf{F}_{64}$ .
- **9.** Let  $a_i \in \{0,1\}$  for  $i \in \{1,2,3\}$ . If  $(a_1a_2a_3)_2$  is a multiple of 3, then  $(a_1a_2a_3)_4$  is a multiple of 6. Here  $(\bullet)_4$  denotes base 4 representation.
- **10.** Let  $f: \mathbf{R} \to \mathbf{R}$  denote the function that maps  $x \in \mathbf{R}$  to the number of double precision floating point numbers contained in the interval [x-1,x+1]. Then f is a decreasing function of x.
- 11. Let  $n \in \mathbb{N}$ . The number of bits in the binary representation of n is less than or equal to 4 times the number of digits in the decimal representation of n.
- **12.** It holds that  $(0.\overline{2200})_3 = (0.9)_{10}$ .
- **13.** Let  $p \in \mathbb{N}$ . The set  $\{(b_0.b_1b_2...b_{p-1})_2 : b_i \in \{0,1\}\}$  contains  $2^p$  distinct real numbers.

Solution. The correct answers are the following:

- 1. True
- 2. True
- 3. True
- **4.** False, because the binary representation of 0.4 is infinite.
- 5. True, because

$$x + \hat{y} = f(x + y) = f(y + x) = \hat{y} + x,$$

where fl is the rounding operator.

- 6. True
- **7.** False. The smallest number that can be represented in a format is  $2^{E_{\min}-(p-1)}$ , and the machine epsilon is  $2^{-(p-1)}$ .
- 8. False, otherwise there would be infinitely many numbers in the set  $\mathbf{F}_{64}$ .
- **9.** False. For example,  $(110)_2 = 6$  and  $(110)_4 = 20$ .
- 10. False since  $\lim_{x\to-\infty} f(x) = 0$  and f(0) > 0.
- 11. True. Indeed, let d denote the number of digits in the decimal representation of n. Then  $n \leq 10^d 1$ . With 4d bits, all the numbers up to  $2^{4d} 1$  can be represented, and since  $2^{4d} 1 = 16^d 1 \geq 10^d 1$ , the statement is true.
- 12. True because

$$(0.\overline{2200})_3 = (0.2200)_3 \left(1 + 3^{-4} + (3^{-4})^2 + (3^{-4})^3 + \cdots \right) = \left(\frac{2}{3} + \frac{2}{9}\right) \frac{1}{1 - 3^{-4}} = \frac{8}{9} \frac{81}{80} = \frac{9}{10}.$$

13. True because there are  $2^p$  choices for the bits, and distinct sets of bits correspond to distinct real numbers.

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**Question 2** (Interpolation and approximation, 10 marks). Throughout this exercise, we assume that  $x_0 < \ldots < x_n$  are distinct values and that  $u: \mathbf{R} \to \mathbf{R}$  is a smooth function. The notation  $\mathbf{P}(n)$  denotes the set of polynomials of degree less than or equal to n.

- 1. (4 marks) Are the following statements true or false? (+1/0/-1)
  - There exists a unique polynomial  $p \in \mathbf{P}(n)$  such that

$$\forall i \in \{0, \dots, n\}, \qquad p(x_i) = u(x_i). \tag{1}$$

• Assume that  $p \in \mathbf{P}(n)$  is such that (1) is satisfied. Then there is a constant  $K \in \mathbf{R}$  independent of x such that

$$\forall x \in \mathbf{R}, \quad u(x) - p(x) = K(x - x_0) \dots (x - x_n).$$

- Assume that  $p \in \mathbf{P}(n)$  is such that (1) is satisfied. Then p is of degree exactly n.
- If  $x_0, \ldots, x_n$  are the roots of the Chebyshev polynomial of degree n, then

$$\sup_{x \in \mathbf{R}} \left| (x - x_0) \dots (x - x_n) \right| \leqslant \frac{\pi}{2^n}.$$

• The function  $S \colon \mathbf{N} \to \mathbf{R}$  given by

$$S(n) = \sum_{i=1}^{n} (i + i^{2} + i^{3} + i^{4})$$

is a polynomial of degree 5. (More precisely, there exists a polynomial of degree 5, say q, such that S(n) = q(n) for all  $n \in \mathbb{N}$ .)

Solution. The correct answers are the following:

• True. Indeed assume that p and q both satisfy (1). Then  $p-q \in \mathbf{P}(n)$  and

$$\forall i \in \{0, \dots, n\}, \qquad (p - q)(x_i) = 0.$$

Therefore p-q has at least n+1 roots which, given that p-q if of degree at most n, is possible only if p-q=0.

• False, because if it were true, then it would hold that

$$u(x) = p(x) + K(x - x_0) \dots (x - x_n),$$

implying that u is a polynomial of degree n+1. Therefore, the equation cannot be true for a general smooth function u.

- False. The statement is not true in general since, if (for example) u is the function everywhere equal to zero, then the only  $p \in \mathbf{P}(n)$  that satisfies (1) is p = 0, which is not a polynomial of degree exactly n.
- False, because the supremum on the left-hand side is equal to  $\infty$  as

$$\lim_{x \to \infty} \left| (x - x_0) \dots (x - x_n) \right| = \infty.$$

• True.

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**2.** For  $i \in \{0, ..., n\}$ , let  $u_i = u(x_i)$ , and let  $m \le n$  be a given natural number. We wish to fit the data  $(x_0, u_0), ..., (x_n, u_n)$  with a function  $\widehat{u} \colon \mathbf{R} \to \mathbf{R}$  of the form

$$\widehat{u}(x) = \alpha_0 + \alpha_1 x + \ldots + \alpha_m x^m.$$

Specifically, we wish to find coefficients  $\boldsymbol{\alpha} = (\alpha_0, \dots, \alpha_m)^T$  such that the error

$$J(\boldsymbol{\alpha}) := \frac{1}{2} \sum_{i=0}^{n} |u_i - \widehat{u}(x_i)|^2$$

is minimized. Throughout this exercise, we use the notations

$$A \begin{pmatrix} 1 & x_0 & \dots & x_0^m \\ \vdots & \vdots & & \vdots \\ 1 & x_n & \dots & x_n^m \end{pmatrix}, \qquad \boldsymbol{b} := \begin{pmatrix} u_0 \\ \vdots \\ u_n \end{pmatrix}$$

• (3 marks) Show that  $J(\alpha)$  may be rewritten as

$$J(\boldsymbol{\alpha}) = \frac{1}{2} (\mathsf{A}\boldsymbol{\alpha} - \boldsymbol{b})^T (\mathsf{A}\boldsymbol{\alpha} - \boldsymbol{b}).$$

• (2 marks) Prove that if  $\alpha_* \in \mathbb{R}^{m+1}$  is a minimizer of J, then

$$\mathsf{A}^T \mathsf{A} \boldsymbol{\alpha}_* = \mathsf{A}^T \boldsymbol{b}. \tag{2}$$

• (1 mark) Find a solution to (2) in terms of  $u_0, \ldots, u_n$  and n when m = 0. Explain.

Solution.

• Notice that

$$\mathbf{A}\boldsymbol{\alpha} = \begin{pmatrix} \alpha_0 + \alpha_1 x_0 + \dots + \alpha_m x_0^m \\ \vdots \\ \alpha_0 + \alpha_1 x_n + \dots + \alpha_m x_n^m \end{pmatrix} = \begin{pmatrix} \widehat{u}(x_0) \\ \vdots \\ \widehat{u}(x_n) \end{pmatrix}.$$

Therefore

$$\frac{1}{2} \sum_{i=1}^{n} \left| \widehat{u}(x_i) - u_i \right|^2 = \frac{1}{2} \sum_{i=1}^{n} \left| (\mathsf{A}\boldsymbol{\alpha} - \boldsymbol{b})_i \right|^2 = \frac{1}{2} (\mathsf{A}\boldsymbol{\alpha} - \boldsymbol{b})^T (\mathsf{A}\boldsymbol{\alpha} - \boldsymbol{b})$$

• A necessary condition is that  $\nabla J(\alpha_*) = 0$ . We calculate that

$$\frac{\partial}{\partial x_i} \left( \boldsymbol{b}^T \boldsymbol{x} \right) = \frac{\partial}{\partial x_i} \left( \sum_{j=1}^n b_j x_j \right) = \sum_{j=1}^n b_j \delta_{ij} = b_i.$$

Similarly, for any matrix  $M \in \mathbf{R}^{n \times n}$ , it holds that

$$\frac{\partial}{\partial x_i} \left( \boldsymbol{x}^T \mathsf{M} \boldsymbol{x} \right) = \frac{\partial}{\partial x_i} \left( \sum_{j=1}^n \sum_{k=1}^n m_{jk} x_j x_k \right) = \sum_{j=1}^n \sum_{k=1}^n m_{jk} \frac{\partial}{\partial x_i} (x_j x_k).$$

Applying the formula for the derivative of a product, we obtain

$$\frac{\partial}{\partial_{x_i}} \left( \boldsymbol{x}^T \mathsf{M} \boldsymbol{x} \right) = \sum_{j=1}^n \sum_{k=1}^n m_{jk} \delta_{ij} x_k + m_{jk} x_j \delta_{ik}$$
$$= \sum_{k=1}^n m_{ik} x_k + \sum_{j=1}^n m_{ji} x_j = (\mathsf{M} \boldsymbol{x} + \mathsf{M}^T \boldsymbol{x})_i.$$

Employing these formulae, we calculate that (representing the gradient with a column vector)

$$\nabla_{\boldsymbol{\alpha}} \left( \boldsymbol{b}^T \boldsymbol{\alpha} \right) = \boldsymbol{b}, \qquad \nabla_{\boldsymbol{\alpha}} \left( \boldsymbol{\alpha}^T \mathsf{A}^T \mathsf{A} \boldsymbol{\alpha} \right) = 2 \mathsf{A}^T \mathsf{A} \boldsymbol{\alpha}.$$

It is then simple to conclude.

• In this case  $A^TA = n + 1$  and  $\alpha_*$  is a scalar. The solution is given by

$$\alpha_* = \frac{u_0 + \dots + u_n}{n+1},$$

which is the average of the values  $u_0, \ldots, u_{n+1}$ .

**Question 3** (Numerical integration, 10 marks). The Gauss–Legendre quadrature formula with n nodes is an approximate integration formula of the form

$$I(u) := \int_{-1}^{1} u(x) dx \approx \sum_{i=1}^{n} w_i u(x_i) =: \widehat{I}_n(u),$$
(3)

which is exact when u is a polynomial of degree less than or equal to 2n-1. (Note that the nodes are here numbered starting from 1.)

1. (5 marks) Find the nodes and weights of the Gauss-Legendre rule with n=3 nodes.

Solution. A necessary and sufficient condition in order for (3) to be satisfied for any polynomial  $p \in \mathbf{P}(5)$  is that

$$\int_{-1}^{1} x^{d} dx = \sum_{i=1}^{n} w_{i} x_{i}^{d}, \quad \text{for all } d \in \{0, 1, 2, 3, 4, 5\}.$$

This leads to the following system of equations

$$\begin{cases} 2 = w_1 + w_2 + w_3, \\ 0 = w_1 x_1 + w_2 x_2 + w_3 x_3, \\ \frac{2}{3} = w_1 x_1^2 + w_2 x_2^2 + w_3 x_3^2, \\ 0 = w_1 x_1^3 + w_2 x_2^3 + w_3 x_3^3, \\ \frac{2}{5} = w_1 x_1^4 + w_2 x_2^4 + w_3 x_3^4, \\ 0 = w_1 x_1^5 + w_2 x_2^5 + w_3 x_3^5. \end{cases}$$

Given the symmetry of the problem, it is reasonable to look for a solution of the form

$$(x_1, x_2, x_3, w_1, w_2, w_3) = (-x, 0, x, w_1, w_2, w_1),$$

where only 3 unknown parameters remain. For such a set of parameters, the second, fourth and sixth equations are satisfied, and the other three equations give

$$\begin{cases} 2 = 2w_1 + w_2, \\ \frac{2}{3} = 2w_1 x^2, \\ \frac{2}{5} = 2w_1 x^4. \end{cases}$$

Dividing the third equation by the second, we obtain  $x^2 = 3/5$  and so  $x = \pm \sqrt{\frac{3}{5}}$  (both values lead to the same integration rule in the end). It is then simple to deduce

that  $w_1 = \frac{5}{9}$  and  $w_2 = \frac{8}{9}$ . We have thus derived the formula

$$\int_{-1}^{1}u(x)\approx\frac{5}{9}u\left(-\sqrt{\frac{3}{5}}\right)+\frac{8}{9}u\left(0\right)+\frac{5}{9}u\left(\sqrt{\frac{3}{5}}\right).$$

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**2.** (2 marks) Let  $\{L_0, L_1, \dots\}$  denote orthogonal polynomials for the inner product

$$\langle f, g \rangle := \int_{-1}^{1} f(x)g(x) \, \mathrm{d}x$$

which, in addition, satisfy the following two conditions:

- For all  $i \in \mathbb{N}$ , the polynomial  $L_i$  is of degree i.
- The leading coefficient of  $L_i$ , which multiplies  $x^i$ , is equal to 1.

Calculate  $L_0$ ,  $L_1$ ,  $L_2$  and  $L_3$ . What is the connection between  $L_3$  and the rule found in the first item?

Solution. Clearly  $L_0 = 1$ . Then  $L_1 = x + a_1$  and the requirement that  $\langle L_1, L_0 \rangle = 0$  implies that  $a_1 = 0$ . We then use the ansatz  $L_2 = x^2 + b_2 x + a_2$  for  $L_2$ . The requirement that  $\langle L_2, L_1 \rangle$  leads to  $b_2 = 0$ , and then

$$\langle L_2, L_0 \rangle = \frac{2}{3} + 2a_2,$$

and so  $L_2(x) = x^2 - \frac{1}{3}$ . Finally, for  $L_3$ , we use the ansatz  $L_3 = x^3 + c_3x^2 + b_3x + a_3$ . We calculate

$$\langle L_3, 1 \rangle = \frac{2}{3}c_3 + 2a_3,$$
  
 $\langle L_3, x \rangle = \frac{2}{5} + \frac{2}{3}b_3,$   
 $\langle L_3, x^2 \rangle = \frac{2}{5}c_3 + \frac{2}{3}a_3.$ 

The second equation gives  $b_3 = -\frac{3}{5}$ , and the other two equations lead to  $c_3 = a_3 = 0$ . We conclude that  $L_3(x) = x^3 - \frac{3}{5}x$ . The roots of  $L_3$  are given by  $\left\{-\sqrt{\frac{3}{5}}, 0, \sqrt{\frac{3}{5}}\right\}$ , and they coincide with the nodes of the Gauss–Legendre quadrature with 3 nodes.  $\triangle$ 

- **3.** Assume that  $x_1, \ldots, x_n$  and  $w_1, \ldots, w_n$  are such that (3) is satisfied for all  $u \in \mathbf{P}(2n-1)$ .
  - (2 marks) Show that the weights are given by

$$\forall i \in \{1, ..., n\}, \qquad w_i = \int_{-1}^1 \ell_i(x) \, \mathrm{d}x,$$

where  $\ell_i$  is the Lagrange polynomial

$$\ell_i(x) = \prod_{j \neq i} \frac{x - x_j}{x_i - x_j}.$$

• (1 marks) Show that the weights are all positive:  $w_i > 0$  for all i.

Solution. Since (3) holds true for all  $u \in \mathbf{P}(2n-1)$ , it holds true in particular for the function  $u = \ell_i \in \mathbf{P}(2n-1)$ , which implies that

$$\int_{-1}^{1} \ell_i(x) \, \mathrm{d}x = \sum_{i=1}^{n} w_i \ell_i(x_j) = w_i.$$

Similarly, since (3) holds true also for  $u \in \ell_i^2 \in \mathbf{P}(2n-1)$ , we deduce that

$$\int_{-1}^{1} (\ell_i(x))^2 dx = \sum_{i=1}^{n} w_j (\ell_i(x_j))^2 = w_i.$$

Since the left-hand side is positive, we deduce that  $w_i > 0$ .

**4.** (Bonus +2) Prove the following error estimate: if u is a smooth function, then

$$|I(u) - \widehat{I}_n(u)| \le \frac{C_{2n}}{(2n)!} \int_{-1}^1 (L_n(x))^2 dx, \qquad C_{2n} := \sup_{\xi \in [-1,1]} |u^{(2n)}(\xi)|.$$

**Hint**: You may find it useful to proceed as follows:

• First show that

$$I(u) - \widehat{I}_n(u) = \int_{-1}^1 u(x) - p(x) \, \mathrm{d}x,\tag{4}$$

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for any polynomial  $p \in \mathbf{P}(2n-1)$  such that

$$\forall i \in \{1, \dots, n\}, \qquad p(x_i) = u(x_i). \tag{5}$$

• Notice that equation (4) is true in particular when p is the Hermite interpolation of u at the nodes  $x_1, \ldots, x_n$ . Finally, conclude by using the formula for the interpolation error proved in class: if p is the Hermite interpolant of u at the nodes  $x_1, \ldots, x_n$ , then

$$\forall x \in \mathbf{R}, \quad u(x) - p(x) = \frac{u^{(2n)}(\xi(x))}{(2n)!} (x - x_1)^2 \dots (x - x_n)^2.$$

Solution. Assume that  $p \in \mathbf{P}(2n-1)$  is such that (5) is satisfied. Then by (3) we deduce that

$$\int_{-1}^{1} p(x) dx = \sum_{i=1}^{n} w_i p(x_i) = \sum_{i=1}^{n} w_i u(x_i) = \widehat{I}_n(u).$$

Consequently, we obtain that

$$I(u) - \widehat{I}_n(u) = \int_{-1}^1 u(x) \, dx - \int_{-1}^1 p(x) \, dx = \int_{-1}^1 u(x) - p(x) \, dx.$$

This equation holds true in particular with p being the Hermite interpolation of u at the nodes  $x_1, \ldots, x_n$ . Then, using the formula for the interpolation error, we obtain

$$u(x) - u(x) = \frac{u^{(2n)}(\xi(x))}{(2n)!}(x - x_1)^2 \dots (x - x_n)^2 = \frac{u^{(2n)}(\xi(x))}{(2n)!}(L_n(x))^2.$$

Indeed, as shown in class,  $L_n$  is a polynomial of degree n with single roots at  $x_1, \ldots, x_n$ . Now we conclude by noting that

$$\left| I(u) - \widehat{I}_n(u) \right| = \left| \int_{-1}^1 u(x) - p(x) \, \mathrm{d}x \right| \leqslant \int_{-1}^1 |u(x) - p(x)| \, \mathrm{d}x \leqslant \int_{-1}^1 \frac{C_{2n}}{(2n)!} \left( L_n(x) \right)^2 \, \mathrm{d}x,$$

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which concludes the exercise.

Question 4 (Vector and matrix norms, 10 marks). The 1-norm and the  $\infty$ -norm of a vector  $x \in \mathbb{R}^n$  are defined as follows:

$$\|x\|_1 = |x_1| + \dots + |x_n|$$
 and  $\|x\|_{\infty} = \max\{|x_1|, \dots, |x_n|\}.$ 

These norms both induce a matrix norm through the formula

$$\|A\|_p := \sup \{ \|Ax\|_p : \|x\|_p = 1 \}.$$

Prove, for  $A \in \mathbf{R}^{n \times n}$ , that

1. (10 marks)  $\|A\|_1$  is given by the maximum absolute column sum:

$$\|\mathsf{A}\|_1 = \max_{1 \le j \le n} \sum_{i=1}^n |a_{ij}|. \tag{6}$$

**2.** (Bonus +2)  $\|A\|_{\infty}$  is given by the maximum absolute row sum:

$$\|\mathsf{A}\|_{\infty} = \max_{1 \leqslant i \leqslant n} \sum_{i=1}^{n} |a_{ij}|.$$

**Hint:** In order to prove (6), you may find it useful to proceed as follows:

• Introduce  $j_*$  as the index of the column with maximum absolute sum:

$$j_* = \operatorname*{arg\,max}_{1 \leqslant j \leqslant n} \sum_{i=1}^n |a_{ij}|.$$

• Prove the direction  $\geqslant$  in (6) by finding a vector  $\boldsymbol{x}$  with  $\|\boldsymbol{x}\|_1 = 1$  such that

$$\|\mathsf{A}\boldsymbol{x}\|_1 = \sum_{i=1}^n |a_{ij_*}|.$$

• Prove the direction  $\leq$  in (6) by showing that, for any  $x \in \mathbb{R}^n$  with  $||x||_1 = 1$ ,

$$\|\mathsf{A}\boldsymbol{x}\|_1 \leqslant \sum_{i=1}^n |a_{ij_*}|.$$

Solution.

1. Let  $e_j$  denote the column vector with a 1 at entry j and zero everywhere else. Notice

that  $\|\boldsymbol{e}_j\|_1 = 1$  and

$$\|\mathsf{A}e_{j_*}\|_1 = \sum_{i=1}^n |a_{ij_*}|,$$

and so  $\|\mathbf{A}\|_1 \geqslant \sum_{i=1}^n a_{ij_*}$ . It remains to prove that  $\|\mathbf{A}\|_1 \leqslant \sum_{i=1}^n a_{ij_*}$ . To this end, it is sufficient to show that  $\|\mathbf{A}\boldsymbol{x}\|_1 \leqslant \sum_{i=1}^n a_{ij_*}$  for all  $\boldsymbol{x} \in \mathbf{R}^n$  with  $\|\boldsymbol{x}\|_1 = 1$ . Take  $\boldsymbol{x} \in \mathbf{R}^n$  with  $\|\boldsymbol{x}\|_1 = 1$ . We calculate that

$$\begin{aligned} \|\mathsf{A}\boldsymbol{x}\|_{1} &= \sum_{i=1}^{n} \left| \sum_{j=1}^{n} a_{ij} x_{j} \right| \leqslant \sum_{i=1}^{n} \sum_{j=1}^{n} |a_{ij}| |x_{j}| \\ &= \sum_{j=1}^{n} \left( \sum_{i=1}^{n} |a_{ij}| \right) |x_{j}| \leqslant \sum_{j=1}^{n} \left( \sum_{i=1}^{n} |a_{ij_{*}}| \right) |x_{j}| \\ &= \left( \sum_{i=1}^{n} |a_{ij_{*}}| \right) \sum_{j=1}^{n} |x_{j}| = \left( \sum_{i=1}^{n} |a_{ij_{*}}| \right) \|\boldsymbol{x}\|_{1} = \sum_{i=1}^{n} |a_{ij_{*}}|, \end{aligned}$$

implying that  $\|A\|_1 \leqslant \sum_{i=1}^n |a_{ij_*}|$ .

**2.** Let  $i_*$  denote the index of a row (not necessarily unique) with maximum absolute sum, and let  $\boldsymbol{y}$  be a column vector with entry j equal to  $\operatorname{sign}(a_{i_*j})$ . Then  $\|\boldsymbol{y}\|_{\infty} = 1$  and

$$\|\mathsf{A}\boldsymbol{y}\|_{\infty} = \sum_{j=1}^{n} |a_{i_*j}|,$$

which implies that  $\|A\|_{\infty} \geqslant \sum_{j=1}^{n} |a_{i_*j}|$ . It remains to prove that  $\|A\|_{\infty} \leqslant \sum_{j=1}^{n} |a_{i_*j}|$ . To this end, take  $\boldsymbol{x} \in \mathbf{R}^n$  with  $\|\boldsymbol{x}\|_{\infty} = 1$ . Then for all  $i \in \{1, \dots, n\}$ ,

$$|(\mathbf{A}\boldsymbol{x})_{i}| = \left| \sum_{j=1}^{n} a_{ij} x_{j} \right| \leqslant \sum_{j=1}^{n} |a_{ij}| |x_{j}| \leqslant \left( \sum_{j=1}^{n} |a_{ij}| \right) \max_{1 \leqslant j \leqslant n} |x_{j}|$$

$$= \left( \sum_{j=1}^{n} |a_{ij}| \right) ||\boldsymbol{x}||_{\infty} = \sum_{j=1}^{n} |a_{ij}| \leqslant \sum_{j=1}^{n} |a_{i_{*}j}|,$$

which implies that  $\|A\|_{\infty} \leqslant \sum_{j=1}^{n} |a_{i_*j}|$ .

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