

A Distributed Augmenting Path Approach for the Bottleneck Assignment Problem
Transactions on Automatic Control
Additional Material

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May 31, 2022

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Function PruneBAP with AugDFS, from agent's perspective

Input: Agent i , initial matched task m_i , agent's edge set \mathcal{E}_i , weight of edges \mathcal{W}_i .

Output: Final matched task m_i .

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1: matching_exists  $\leftarrow$  True
2:  $\bar{\mathcal{E}}_i \leftarrow \mathcal{E}_i$ 
3: while matching_exists do
4:   Find edge with largest weight in  $\bar{\mathcal{E}}_i$ 
5:   Let the edge and the weight be a tuple  $(e_i, w_i)$ 
6:   for  $d \in D$  do
7:     for  $k \in N(\mathcal{G}_C, i)$  do ▷ Neighbours of  $i$ 
8:       Collect  $(e_k, w_k)$ 
9:     end for
10:    Set the new  $(e_i, w_i)$  to be the one containing the largest weight from the collect tuples
11:  end for
12:   $f_i = \text{False}$  ▷  $i$  is unexplored
13:   $\nu_i \leftarrow m_i$ 
14:  search_complete  $\leftarrow$  False
15:  Let  $(\bar{i}, \bar{j}) = e_i$ 
16:   $t \leftarrow \bar{j}$ 
17:   $FILO \leftarrow t$ 
18:  while  $\neg \text{search\_complete}$  do
19:    Check existence of edges  $(i, t) \in \bar{\mathcal{E}}_i$ 
20:    If edge exists, set tuple  $out_i = (m_i, w(i, m_i))$ , or  $out_i = (\hat{b}, \text{tiebreaking identifier})$  if  $i$  is
    free, or else set it to  $out_i = \emptyset$ 
21:    for  $d \in D$  do
22:      for  $k \in N(\mathcal{G}_C, i)$  do
23:        Collect  $out_k$ 
24:      end for
25:      Set the new  $out_i$  to be the one with smallest weight from the collect tuples, or  $out_i =$ 
       $(\hat{b}, \text{tiebreaking identifier})$  if one of the collected tuples shows a free agent exists.
26:    end for
27:    if  $out_i = \emptyset$  and  $t = \bar{j}$  then ▷ No remaining agents
28:      search_complete  $\leftarrow$  True
29:      matching_exists  $\leftarrow$  False
30:    else if  $out_i = \emptyset$  and  $t \neq \bar{j}$  then ▷  $t$  has no children
31:      Check and remove last element in  $FILO, t^*$ 
32:       $t \leftarrow t^{**}$ , where  $t^{**}$  is the new last element in  $FILO$  after removal of  $t^*$ 
33:      if  $m_i == t^*$  then
34:         $\nu_i \leftarrow t^*$ 
35:      end if
36:    else if  $out_i = (\hat{b}, \text{tiebreaking identifier})$  then ▷ Free agent found
37:      if  $i$  matches the tiebreaking identifier then
38:         $\nu_i \leftarrow t$ 
39:      end if
40:      search_complete  $\leftarrow$  True
41:    else ▷ Explore next agent
42:      if  $i$  matches the tiebreaking identifier then
43:         $\nu_i \leftarrow t$ 
44:         $f_i \leftarrow \text{True}$  ▷ Mark  $i$  as explored
45:      end if
46:       $t \leftarrow m_k$  from the current saved  $out_i = (m_k, w(k, m_k))$ 
47:      Append  $t$  to  $FILO$ 
48:    end if
49:  end while
50:   $m_i = \nu_i$ 
51: end while
52: return  $m_i$ 

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