

GROUP MEMEBERS

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SOLUTIONS:

QUESTION 1:

Step 1: Understand the structure of the tree.

A tree consists of nodes connected by edges. The height $\lfloor h \rfloor$ of a tree is the length path from the root to a leaf. The height $\lfloor h \rfloor$ of a tree in the height $\lfloor f \rfloor$ is achieved when every leaf of the tree is fully populated.

Step 2 (Comparison of the Number of nodes at each level)

Level (Rot) - Level 0 - Rot node is at level 0 and contributes 4 node.

Level 1 - After label on line a maximum of 2 nodes in

level 2 - The level on line a maximum of 4 nodes.

The therefore conclude level 1 has 2 nodes.

Step 3: Maximum number of nodes (Total number of nodes)

Given a height $\lfloor h \rfloor$, the total number of nodes in a tree of height $\lfloor h \rfloor$ is:

$$\lfloor N = 2 + 4 + \dots + 2h. \rfloor$$

This form a geometric series with sum:

$$\lfloor N = \sum_{i=0}^{\lfloor h \rfloor} 2^i = \frac{2^{\lfloor h \rfloor + 1} - 2}{2 - 1} = 2^{\lfloor h \rfloor + 1} - 1. \rfloor$$

QUESTION 2:

1). Depth first search - Is a tree or graph traversal algorithm that explores as far as possible along each branch before backtracking?

Example:

DEE Traversal (Reorder, Rat->Left-> Right):

Chasing from A:

1 Visit A:

2 Go to B - Visit B:

3 Go to E -> Visit E(Backtrack to B):

4 Go to F -> Visit F(Backtrack to B -> A:

5 Go to C -> Visit C(Backtrack to A):

6 Go to D -> Visit D:

7 Go to G -> Visit G(Backtrack to D):

8 Go to H -> Visit H.

Order:

A -> B -> E -> F -> C -> D

→ G -> H

2) Breadth Search

It is a search algorithm that stops at the search of the root and explores all the nodes of the present depth before moving on the nodes of the next depth level.

Example:

Starting from A:

1. Visit A
2. Visit B, C, D
3. Visit E, F, G, H

BFS Order: $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H$

Use case: Shortest path in an unweighted graph.

3) AVL Search

It is performed on AVL Tree, which is a self-balancing binary search tree where the height difference between left and right subtrees at most 1.

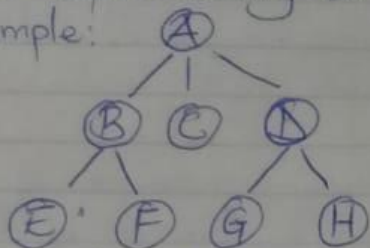
Balance factors for different nodes are:

- 4-2, 1-2, 1-5, 1-1, 1-0, 1-4, 0, and 4-0; hence it is an AVL Tree since all differences are less than or equal to

Use case: Pathfinding in maze.

(i) Breadth-Search \rightarrow It is a search algorithm that starts at the search at the root and explores all nodes at the present depth before moving on the nodes at the next depth level.

Example:



Starting from A:

1. Visit A
2. Visit B, C, D.
3. Visit E, F, G, H

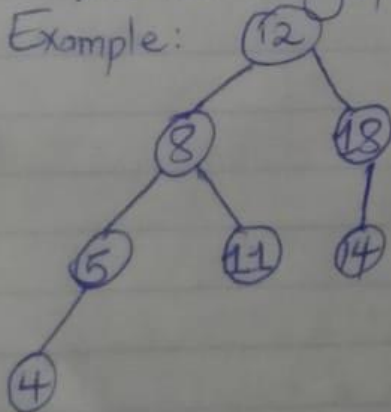
BFS Order: $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F \rightarrow G \rightarrow H$.

Use case: Shortest path in an unweighted graph.

(ii) AVL Search - It is performed on AVL Tree, which is a self-balancing binary search tree where the height difference (balance factor) between left and right subtrees is at most 1.

(3) Tree as a graph data structure.

Example:



Balance factors for different nodes are: 12:1, 18:1, 5:1, 11:0, 14:0, and 4:0; hence it is an AVL Tree since all differences are less than or equal to 1.

(3) Tree as a graph data structure.

A tree is a specific type of graph characterized by being

④ Step 1: Understand the Structure of the Tree

A tree consists of nodes connected by edges. The height h of a tree is the length path from the root to a leaf. The height (h) of a tree is the length of the path from the root to a leaf. The level of the tree is fully populated.

Step 2: Computation of the Number of nodes at each level

Level (Root) - Level 0 - Root node is at level 0 and contributes 1 node.

Level 1 - This level can have a maximum of 2 nodes.

Level 2 - This level can have a maximum of 4 nodes.

\therefore We therefore conclude level i has 2^i nodes.

Step 3: Maximum number of nodes (Total number of nodes)

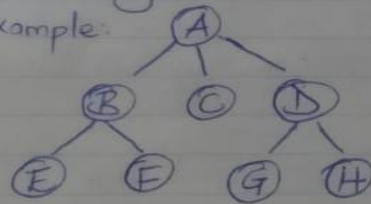
- Given a height h , the total number of nodes in a tree of height h is: $N = 1 + 2 + 4 + \dots + 2^h$.

This forms a geometric series with sum:

$$N = \sum_{i=0}^h 2^i = \frac{2^{h+1} - 1}{2 - 1} = 2^{h+1} - 1.$$

② Depth first search - Is a tree or graph traversal algorithm that explores as far as possible along each branch before backtracking.

Example:



DFS Traversal (Preorder: Root \rightarrow Left \rightarrow Right):

(Starting from A):

1. Visit A.
2. Go to B - Visit B
3. Go to E \rightarrow Visit E (Backtrack to B)
4. Go to F \rightarrow Visit F (Backtrack to B \rightarrow A)
5. Go to C \rightarrow Visit C (Backtrack to A)
6. Go to D \rightarrow Visit D
7. Go to G \rightarrow Visit G (Backtrack to D)
8. Go to H \rightarrow Visit H.

Order: $A \rightarrow B \rightarrow E \rightarrow F \rightarrow C \rightarrow D \rightarrow G \rightarrow H$.

QUESTION 4

```
class Node { int key; Node left, right;
```

```
public Node(int item) {
```

```
    key = item;
```

```
    left = right = null;
```

```
}
```

```
}
```

```
class BinaryTree { Node root;
```

```
BinaryTree() {
```

```
    root = null;
```

```
}
```

```
void insert(int key) {
```

```
    root = insertRec(root, key);
```

```
}
```

```
Node insertRec(Node root, int key) {
```

```
    if (root == null) {
```

```
        root = new Node(key);
```

```
        return root;
```

```
    }
```

```
    if (key < root.key)
```

```
        root.left = insertRec(root.left, key);
    else if (key > root.key)
        root.right = insertRec(root.right, key);
    return root;
}
```

```
void inorder() {
    inorderRec(root);
}
```

```
void inorderRec(Node root) {
    if (root != null) {
        inorderRec(root.left);
        System.out.print(root.key + " ");
        inorderRec(root.right);
    }
}
```

```
Node search(Node root, int key) {
    if (root == null || root.key == key)
        return root;
    if (root.key > key)
        return search(root.left, key);
    return search(root.right, key);
}
```



```
void delete(int key) {  
    root = deleteRec(root, key);  
}
```

```
Node deleteRec(Node root, int key) {  
    if (root == null) return root;  
    if (key < root.key)  
        root.left = deleteRec(root.left, key);  
    else if (key > root.key)  
        root.right = deleteRec(root.right, key);  
    else {  
        if (root.left == null) return root.right;  
        else if (root.right == null) return root.left;  
        root.key = minValue(root.right);  
        root.right = deleteRec(root.right, root.key);  
    }  
    return root;  
}
```

```
int minValue(Node root) {  
    int minv = root.key;  
    while (root.left != null) {  
        minv = root.left.key;  
        root = root.left;  
    }  
    return minv;  
}
```

```
}
```

```
public static void main(String[] args) {
```

```
    BinaryTree tree = new BinaryTree();
```

```
    tree.insert(50);
```

```
    tree.insert(30);
```

```
    tree.insert(70);
```

```
    tree.insert(20);
```

```
    tree.insert(40);
```

```
    tree.insert(60);
```

```
    tree.insert(80);
```

```
    System.out.println("Inorder traversal:");
```

```
    tree.inorder();
```

```
    System.out.println("\nDeleting 40");
```

```
    tree.delete(40);
```

```
    System.out.println("Inorder traversal after deletion:");
```

```
    tree.inorder();
```

```
}
```

```
}
```