

Theoretical Exercise Sheet 7

Deadline Friday, June 17, 23:59

About the submission of this sheet.

- You might submit the solutions to exercises in groups of up to 3 students.
- All students of a group need to be in the same tutorial.
- Write the names of **all** students of your group on your solution.
- Hand in the solution **in CMS** and use “Team Groupings”.
 - Go to your personal page in CMS. Here you find the entry “Teams”.
 - When you click “Create team”, you get an invite code.
 - Please share this with your team mates, who need to click on “Join team” and enter the code.

1. (16 points) (4 + 4 + 4 + 4)

1. For each of the following sentences, write down a **predicate logic formula that has the same meaning**. You may only use the following predicates: Tutor(x), Student(x), Assistant(x), Likes(x, y), GivesGoodGradesTo(x, y), GoodStudent(x)

- (i) If a tutor gives good grades to every student, every student likes them.
- (ii) Someone is a good student if and only if there is a tutor that gives them good grades but does not give good grades to every student.

2. For each of the following predicate logic formulas, write down an **English sentence with the same meaning**.

- (i) $\exists x[\text{Student}(x) \wedge \neg \exists y[\text{Tutor}(y) \wedge \text{Likes}(x, y)]]$
- (ii) $\forall x \forall y[\text{Assistant}(x) \wedge \text{Tutor}(y) \wedge \neg \text{Likes}(x, y)] \rightarrow \exists z[\text{GoodStudent}(z) \wedge \text{Student}(z) \wedge \neg \text{GivesGoodGradesTo}(y, z)]$

2. (22 points) (5 + 8 + 9) Transform the following predicate logic formulas into **Clausal Normal Form**. Write down the results of all intermediate steps, specifying which steps you are applying and giving the intermediate results.

Note: Simplify the formulas where possible.

- (a) $\varphi_1 = \forall x[(P(x) \wedge \exists y Q(y)) \rightarrow \exists y \neg R(y, x)]$
- (b) $\varphi_2 = \neg \exists x[\exists y[(\neg P(x) \wedge \neg R(y)) \rightarrow Q(x, y)] \rightarrow \exists y(P(x) \wedge Q(x, y) \wedge \forall z R(z))]$
- (c) $\varphi_3 = \neg \exists x \forall y \forall z[P(x, y) \leftrightarrow (\neg Q(x, z) \vee R(y))]$

3. (21 points) (5 + 6 + 2 + 8)

Run the **unification** algorithm from the lecture on each of the given sets. Assume w, x, y, z to be variables and a, b, c to be constants. For each step, write down the **set to unify**, the **disagreement set**, and the **updated set of substitutions**, i.e., $T_0, D_0, s_1, T_1, D_1, \dots$ until the algorithm terminates. In particular, if the algorithm does not output a failure, list the *entire* content of the substitutions s_1, \dots, s_n . In any case, state the **output** of the algorithm. If multiple substitutions are possible, choose **lexicographically** (e.g., $\frac{x}{f(a)}$ before $\frac{x}{f(y)}$).

- (a) $\{P(a, b, c), P(x, y, f(x))\}$
- (b) $\{P(g(a), y, f(z), h(x, b)), P(g(x), b, f(b), h(a, z))\}$
- (c) $\{P(x, y), P(g(a, x), a)\}$
- (d) $\{P(a, f(b, g(x)), g(z)), P(z, f(x, g(b)), w), P(a, y, g(a))\}$

4. (23 points) (3 + 8 + 12) In this exercise you will see whether Rapunzel is saved by the prince, or cursed by the sorcerer. For this consider the following statements and the set $\theta^* = \{\varphi_1, \varphi_2, \varphi_3, \varphi_4, \varphi_5\}$ of corresponding predicate logic formulas in Skolem normal form:

- 1. Rapunzel lives in a tower.
 $\varphi_1 = \text{lives}(\text{Rapunzel}, \text{tower})$
- 2. If the Prince does not save Rapunzel, then she can't escape the tower.
 $\varphi_2 = \text{saves}(\text{Prince}, \text{Rapunzel}) \vee \neg \text{escapes}(\text{Rapunzel}, \text{tower})$
- 3. Everyone who lives in the tower and does not escape will be cursed by the sorcerer.
 $\varphi_3 = \forall x [\neg \text{lives}(x, \text{tower}) \vee \text{escapes}(x, \text{tower}) \vee \text{curse}(\text{sorcerer}, x)]$
- 4. The Prince does not save anyone that lives in a tower.
 $\varphi_4 = \forall x [\neg \text{saves}(\text{Prince}, x) \vee \neg \text{lives}(x, \text{tower})]$
- 5. The sorcerer does not curse Rapunzel.
 $\varphi_5 = \neg \text{curse}(\text{sorcerer}, \text{Rapunzel})$

- (a) Write down $CF(\theta^*)$ and the resulting Herbrand universe $HU(\theta^*)$.
- (b) Write down the Herbrand expansion $HE(\theta^*)$. Bring each of the formulas in $HE(\theta^*)$ into CNF, resulting in a set of clauses Δ .
- (c) Use **propositional resolution** to prove that Δ is unsatisfiable. Assuming that statements 1.–4. are true, will Prince save Rapunzel from the sorcerer?

5. (10 points) Consider the following set of predicate logic formulas in Skolem Normal Form:

1. Every university mascot is an animal.
 $\varphi_1 = \forall x [\neg Mascot(x) \vee Animal(x)]$
2. Steffi the owl or Kasper the clown are university mascots.
 $\varphi_2 = Mascot(Steffi) \vee Mascot(Kasper)$
3. Every animal is liked by some student.
 $\varphi_3 = \forall x [\neg Animal(x) \vee (Student(f(x)) \wedge Likes(f(x), x))]$
4. No student likes Kasper the clown.
 $\varphi_4 = \forall x [\neg Student(x) \vee \neg Likes(x, Kasper)]$

We want to prove by contradiction that Steffi is an animal. Therefore, we add to our set of clauses:

5. Steffi the owl is not an animal.
 $\varphi_5 = \neg Animal(Steffi)$

These clauses give us Δ :

$$\begin{aligned}
 \varphi_1 &= \{\neg Mascot(x), Animal(x)\} \\
 \varphi_2 &= \{Mascot(Steffi), Mascot(Kasper)\} \\
 \varphi_{3,1} &= \{\neg Animal(y), Student(f(y))\} \\
 \varphi_{3,2} &= \{\neg Animal(z), Likes(f(z), z)\} \\
 \varphi_4 &= \{\neg Student(w), \neg Likes(w, Kasper)\} \\
 \varphi_5 &= \{\neg Animal(Steffi)\}
 \end{aligned}$$

Use **binary PL1 resolution** to show that Δ is unsatisfiable. Also specify the unifiers when applying a resolution step.

Note: If you need to use a clause more than once, make a copy of the clause and rename the variables to avoid name collisions.

Depict the resolution process in form of a tree for easier readability. Label edges with the corresponding unifying substitution