

$$\forall xs \in L(X): \forall n \in \mathbb{N}: |xs| + n = \text{foldL}(\underbrace{(\lambda (x,s) \in X \times \mathbb{N}. s+1)}_f, \underbrace{0}_n, xs)$$

Beweis durch strukturelle Induktion über $xs \in L(X)$.

1. Fall: $xs = \text{nil}$:

Sei $n \in \mathbb{N}$ beliebig

$$\begin{aligned} \text{foldL}(f, 0, xs) &= \text{foldL}(f, 0, \text{nil}) && | \text{ } xs = \text{nil} \\ &= n && | \text{ Def foldL} \\ &= 0 + n && | \text{ Arithmetik} \\ &= |\text{nil}| + n && | \text{ Def } | _ | \\ &= |xs| + n && | \text{ } xs = \text{nil} \end{aligned}$$

2. Fall: $xs = x :: xr$:

$$\text{Induktionsannahme: } \forall n \in \mathbb{N}: \text{foldL}(f, 0, xr) = |xr| + n$$

Sei $n \in \mathbb{N}$ beliebig.

$$\begin{aligned} \text{foldL}(f, 0, xs) &= \text{foldL}(f, 0, x :: xr) && | \text{ } xs = x :: xr \\ &= \text{foldL}(f, \underbrace{f(x, 0)}_{\substack{\text{Ein mögl. } n \\ \downarrow}}, xr) && | \text{ Def foldL} \\ &= |xr| + f(x, n) && | \text{ Induktion für } xr \\ &= |xr| + (n+1) && | \text{ Def } f \\ &= \underbrace{|xr| + 1}_\uparrow + n && | \text{ Mathe} \\ &= |x :: xr| + n && | \text{ Def } | _ | \\ &= |xs| + n && | \text{ } xs = x :: xr \end{aligned}$$

□

Somit gilt mit $n=0$ die geforderte Aussage.