

Software Engineering

WS 2022/23, Sheet 06



Prof. Dr. Sven Apel
Yannick Lehmen, Maurice Vincon

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Task 1

- a) Explain the concept of *parametric polymorphism*.
- b) **Discussion:** Prepare arguments for and against implementing variability by using parametric polymorphism. You can present your arguments during the discussion in the tutorial sessions.

Solution

- a) *Parametric Polymorphism* is the concept of writing functions or data types of a program generically. This means that the type of the values of data structures or functions can be changed at compile time.

Task 2

Prove the cases 5, 7, 10, and 12 from the *Coarsening and Refinement* slide in the lecture slides. Use the notation that is used to prove the cases 1–4 in the lecture.

Solution

Proof. We have to proof that for any $e, e' \in VExpr$ it holds that $e \equiv e'$ if and only if, for any given configuration $\psi \in \mathbb{B}_F$ with e and e' closed under ψ , $\llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi$ holds. The proof is by structural induction over the terms $e, e' \in VExpr$.

Induction hypothesis: If e and e' both contain no choice expressions, then $e \equiv e'$ if and only if $e = e'$ and therefore also $\forall \psi \in \mathbb{B}_F. \llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi$, as none of the transformation rules are applicable.

Induction step: $e_1, e_2, e_3, e_4 \in VExpr$, $\text{closed}(e_1, \psi)$, $\text{closed}(e_2, \psi)$, $\text{closed}(e_3, \psi)$, $\text{closed}(e_4, \psi)$

Case 5: $e = \text{choice } \langle \phi, e_1, e_2 \rangle < e_3$ $e' = \text{choice } \langle \phi, e_1 < e_3, e_2 < e_3 \rangle$

$$\begin{aligned}\llbracket e \rrbracket_\psi &= \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle < e_3 \rrbracket_\psi = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi < \llbracket e_3 \rrbracket_\psi \\ \llbracket e' \rrbracket_\psi &= \llbracket \text{choice } \langle \phi, e_1 < e_3, e_2 < e_3 \rangle \rrbracket_\psi\end{aligned}$$

if $\psi \Rightarrow \Phi$:

$$\llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi < \llbracket e_3 \rrbracket_\psi = \llbracket e_1 \rrbracket_\psi < \llbracket e_3 \rrbracket_\psi = \llbracket e_1 < e_3 \rrbracket_\psi = \llbracket \text{choice } \langle \phi, e_1 < e_3, e_2 < e_3 \rangle \rrbracket_\psi$$

if $\psi \Rightarrow \neg\Phi$:

$$\llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi < \llbracket e_3 \rrbracket_\psi = \llbracket e_2 \rrbracket_\psi < \llbracket e_3 \rrbracket_\psi = \llbracket e_2 < e_3 \rrbracket_\psi = \llbracket \text{choice } \langle \phi, e_1 < e_3, e_2 < e_3 \rangle \rrbracket_\psi$$

$$\Rightarrow \llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi \xrightarrow{\text{(I.H.)}} e \equiv e'$$

Case 7: $e = \text{if choice } \langle \phi, e_1, e_2 \rangle \text{ then } e_3 \text{ else } e_4$ $e' = \text{choice } \langle \phi, \text{if } e_1 \text{ then } e_3 \text{ else } e_4, \text{if } e_2 \text{ then } e_3 \text{ else } e_4 \rangle$

$$\begin{aligned}\llbracket e \rrbracket_\psi &= \llbracket \text{if choice } \langle \phi, e_1, e_2 \rangle \text{ then } e_3 \text{ else } e_4 \rrbracket_\psi = \text{if } \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ then } \llbracket e_3 \rrbracket_\psi \text{ else } \llbracket e_4 \rrbracket_\psi \\ \llbracket e' \rrbracket_\psi &= \llbracket \text{choice } \langle \phi, \text{if } e_1 \text{ then } e_3 \text{ else } e_4, \text{if } e_2 \text{ then } e_3 \text{ else } e_4 \rangle \rrbracket_\psi\end{aligned}$$

if $\psi \Rightarrow \Phi$:

$$\begin{aligned}\text{if } \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ then } \llbracket e_3 \rrbracket_\psi \text{ else } \llbracket e_4 \rrbracket_\psi &= \text{if } \llbracket e_1 \rrbracket_\psi \text{ then } \llbracket e_3 \rrbracket_\psi \text{ else } \llbracket e_4 \rrbracket_\psi \\ &= \llbracket \text{if } e_1 \text{ then } e_3 \text{ else } e_4 \rrbracket_\psi \\ &= \llbracket \text{choice } \langle \phi, \text{if } e_1 \text{ then } e_3 \text{ else } e_4, \text{if } e_2 \text{ then } e_3 \text{ else } e_4 \rangle \rrbracket_\psi\end{aligned}$$

if $\psi \Rightarrow \neg\Phi$:

$$\begin{aligned}\text{if } \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ then } \llbracket e_3 \rrbracket_\psi \text{ else } \llbracket e_4 \rrbracket_\psi &= \text{if } \llbracket e_2 \rrbracket_\psi \text{ then } \llbracket e_3 \rrbracket_\psi \text{ else } \llbracket e_4 \rrbracket_\psi \\ &= \llbracket \text{if } e_2 \text{ then } e_3 \text{ else } e_4 \rrbracket_\psi \\ &= \llbracket \text{choice } \langle \phi, \text{if } e_1 \text{ then } e_3 \text{ else } e_4, \text{if } e_2 \text{ then } e_3 \text{ else } e_4 \rangle \rrbracket_\psi\end{aligned}$$

$$\Rightarrow \llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi \xRightarrow{\text{(I.H.)}} e \equiv e'$$

Case 10: $e = (\text{let } x = \text{choice } \langle \phi, e_1, e_2 \rangle \text{ in } e_3) \quad e' = \text{choice } \langle \phi, \text{let } x = e_1 \text{ in } e_3, \text{let } x = e_2 \text{ in } e_3 \rangle$

$$\begin{aligned} \llbracket e \rrbracket_\psi &= \llbracket \text{let } x = \text{choice } \langle \phi, e_1, e_2 \rangle \text{ in } e_3 \rrbracket_\psi = (\text{let } x = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ in } \llbracket e_3 \rrbracket_\psi) \\ \llbracket e' \rrbracket_\psi &= \llbracket \text{choice } \langle \phi, \text{let } x = e_1 \text{ in } e_3, \text{let } x = e_2 \text{ in } e_3 \rangle \rrbracket_\psi \end{aligned}$$

if $\psi \Rightarrow \Phi$:

$$\begin{aligned} (\text{let } x = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ in } \llbracket e_3 \rrbracket_\psi) &= (\text{let } x = \llbracket e_1 \rrbracket_\psi \text{ in } \llbracket e_3 \rrbracket_\psi) \\ &= \llbracket \text{let } x = e_1 \text{ in } e_3 \rrbracket_\psi \\ &= \llbracket \text{choice } \langle \phi, \text{let } x = e_1 \text{ in } e_3, \text{let } x = e_2 \text{ in } e_3 \rangle \rrbracket_\psi \end{aligned}$$

if $\psi \Rightarrow \neg\Phi$:

$$\begin{aligned} (\text{let } x = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi \text{ in } \llbracket e_3 \rrbracket_\psi) &= (\text{let } x = \llbracket e_2 \rrbracket_\psi \text{ in } \llbracket e_3 \rrbracket_\psi) \\ &= \llbracket \text{let } x = e_2 \text{ in } e_3 \rrbracket_\psi \\ &= \llbracket \text{choice } \langle \phi, \text{let } x = e_1 \text{ in } e_3, \text{let } x = e_2 \text{ in } e_3 \rangle \rrbracket_\psi \end{aligned}$$

$$\Rightarrow \llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi \xRightarrow{\text{(I.H.)}} e \equiv e'$$

Case 12: $e = \text{choice } \langle \phi, \text{choice } \langle \phi, e_1, e_2 \rangle, e_2 \rangle \quad e' = \text{choice } \langle \phi, e_1, e_2 \rangle$

if $\psi \Rightarrow \Phi$:

$$\llbracket e \rrbracket_\psi = \llbracket \text{choice } \langle \phi, \text{choice } \langle \phi, e_1, e_2 \rangle, e_2 \rangle \rrbracket_\psi = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi = \llbracket e' \rrbracket_\psi$$

if $\psi \Rightarrow \neg\Phi$:

$$\llbracket e \rrbracket_\psi = \llbracket \text{choice } \langle \phi, \text{choice } \langle \phi, e_1, e_2 \rangle, e_2 \rangle \rrbracket_\psi = \llbracket e_2 \rrbracket_\psi = \llbracket \text{choice } \langle \phi, e_1, e_2 \rangle \rrbracket_\psi = \llbracket e' \rrbracket_\psi$$

$$\Rightarrow \llbracket e \rrbracket_\psi = \llbracket e' \rrbracket_\psi \xRightarrow{\text{(I.H.)}} e \equiv e'$$

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