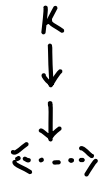


- A knight in chess can move any vertical/horizontal direction two spaces, followed by one space perpendicular to its 2 space motion



- Translate this to a phone pad, given the length of digit is  $n$  ( $1 \leq n \leq 5000$ ) how many unique digits can it produce?

$n=3$

1 2 3  
4 5 6  
7 8 9  
- 0 -

1 2 3  
4 5 6  
7 8 9  
0

Starting @ 1: 1 8 1

1 8 3

1 6 7

1 6 1

1 6 1 6, 1 6 1 8, 1 6 7 2, 1 6 7 6

1 8 1 8, 1 8 1 6, 1 8 3 8, 1 8 3 4

$$\text{unique } 1_n = 2^{n-1}$$

Starting @ 2:

$n=1, l=1$

$n=2, l=4$

2 7 2, 2 9 2, 2 7 6, 2 9 4

$n=3, l=10$

2 7 2 7, 2 7 2 9, 2 7 6 1, 2 7 6 7, 2 7 6 0

2 9 2 9, 2 9 2 7, 2 9 4 3, 2 9 4 9, 2 9 4 0

$n=4$

For every move after 1,  
it can choose 2 paths  
1, 2, 4, 8, 16

There might be a solution  
involving the summation of equations  
too hard to figure out.

Realistic: DFS

Ex.	1 → (6, 8)	f(n: int)
	2 → (9, 7)	for i = 0 → 10
	3 → (4, 8)	ans = unique(n-1, i)
	4 → (0, 3, 9)	
	5 → ( )	unique(n-left: int, value: int, {ans: int})
	6 → (0, 1, 7)	if n-left < 1
	7 → (6, 2)	return 1
	8 → (1, 3)	for values → m[value]
	9 → (4, 2)	X for Key → values
	0 → (4, 6)	<del>return</del> unique(n-1, Key)
		ans +=
		return ans

This is a brute force but can be improved by "caching" values we have seen before.

cache: map<pair<int, int>, long long>

Instead of ans += ...

cache[pair(n-1, Key)] = unique...

ans += cache[pair(n-1, Key)]

Check if cache contains a value b4 calculating.

This solution greatly speeds up algorithm without much manipulation to the function.