

(S 726 : HWI)

(a) Prove that all trails not of the form  $x \rightarrow \dots \leftarrow y$  are blocked by  $z$ .

(b) We can prove the above state using d-separation rule.  
The other possible trails are:

$$x \rightarrow \dots z \rightarrow y$$

$$x \leftarrow z \dots \leftarrow y$$

$$x \leftarrow \dots z \rightarrow y$$

In all these 3 cases,  $x$  &  $y$  are blocked by  $z$  because of d-separation rule.

(c) Prove that all trails of the form  $x \rightarrow w, \dots \leftarrow y$  with exactly one V-node is blocked.

(d) Using d-separation principle

$x \rightarrow w, \leftarrow y$   $x$  &  $y$  will be blocked if  $w$  is not observed

If the trail is like:

$$x \rightarrow w_1 \dots \rightarrow w_n \leftarrow y \rightarrow \text{Exactly one V-node}$$

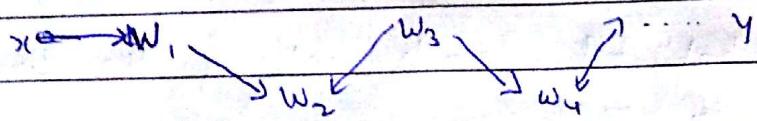
$z$  lies in this trail  $\Rightarrow$  There will be loop in the BN

~~z~~ ~~This will be~~  $\Rightarrow$  It is not a V-node

Hence, the trail of these forms with exactly one V-node is blocked. (We say this using d-separation rule as shown in part (a))

(e) Prove that all trails of the form  $x \rightarrow w, \dots \leftarrow y$  with more than one V-node is also blocked.

(c) Consider a general trail<sup>1</sup> having multiple v-node



Suppose for any v-node, z is the part in that v-node. This will result in the formation of loop ( $\because z = Pa(z) \cup Pa(y)$ ). But formation of loop is not allowed in the BN.

Hence, we cannot observe "z" for a valid BN and using d-separation principle we claim that ~~the~~ trails of the above type would also be blocked.

3 (a) ii)  $A \perp\!\!\!\perp G | F$

Correct

Reason:

node B & C are v-node

~~We are not observing B & C hence their path A to G is~~

Consider the following path  $A \rightarrow E \rightarrow F \rightarrow G$

Using d-separation rule, we claim that

$A \perp\!\!\!\perp G | F$

(ii)  $A \perp\!\!\!\perp D$

Correct

Reason:

We see that all the trails b/w A & D have a v-node. If we do not observe this v-node Then A & D are independent.

(iii)  $B \perp\!\!\!\perp G | C, F$

Wrong

Consider the trail ~~BA~~  $B - A - C - D - G$

So, using d-separation rule B & C cannot be independent if you "observe" C (since it is a v-node)

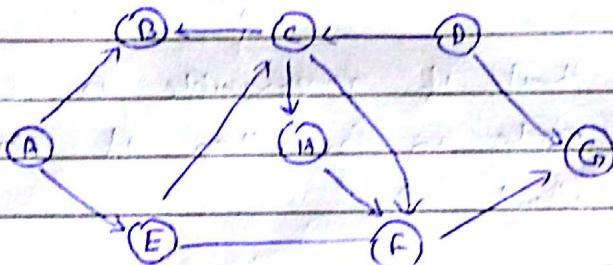
(iv)  $F \perp\!\!\!\perp D \mid C$

Wrong

Consider the trail  $F - E - C - D$

Here, C is a v-node & you are observing C. Hence using d-separation rule F & D become dependent given C.

(b)



We can go to H from G using the two paths  
Path 1: G → D → C → ... H

D blocks the trail b/w H & G (using D-separation)  
 $D \in S$

Path 2

$G \rightarrow F \rightarrow \dots H$

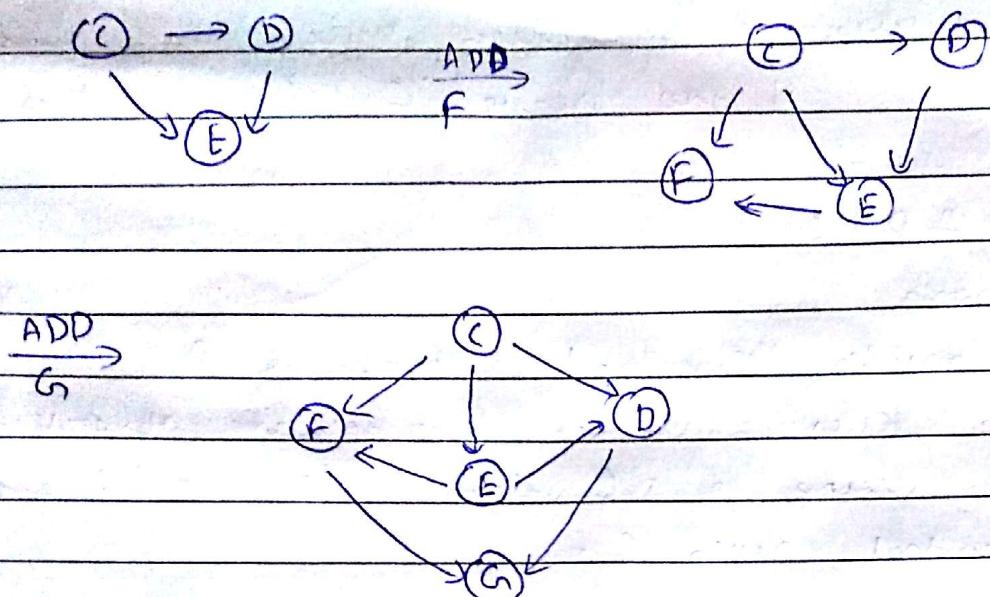
F also blocks the trail b/w G & H. Moreover there is no v-node in the path.

$F \in S$

Hence set  $S = \{D, F\}$

(c) Start adding nodes in topological order & connect the edges

$\textcircled{C} \rightarrow \textcircled{D}$

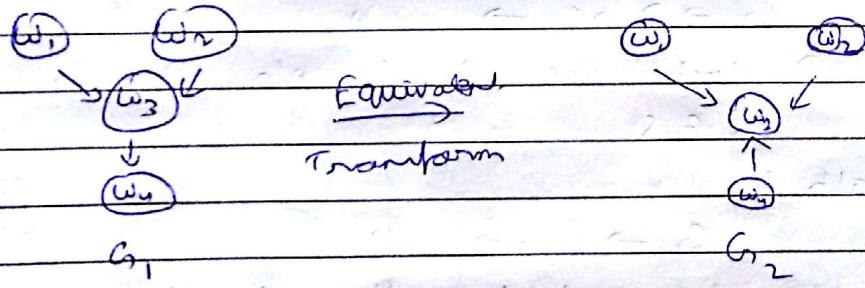


4. (a) Prove that if a variable is not a descendent of a V-node in  $G_1$ , then it cannot be a descendent in equivalent  $G_2$ .

Consider a V-node

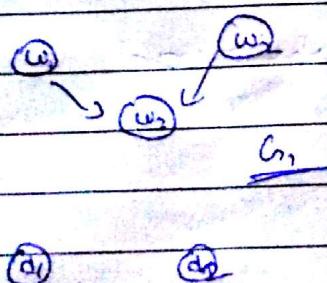
Case 1:

$G_1$ : Assume V-node has a descendent



In  $G_2$ , the transformation is possible but  $w_4$  has become parent. V-node is destroyed

Case 2: node



$w_3$  has not is descendent  
if  $d_1$  &  $d_2$  are non connected  
If we transform this  
graph to equivalent  
graph  $G_2$ .  $b_2$

cannot have a descendent because of the property of equivalent graph that they should have same skeleton.

- (b) In equivalent graph  $G_1$  &  $G_2$ , v-nodes are preserved.

Hence if - in  $G_1$ , ~~there~~ is a node  $\hat{a}$  v-node, then after transformation, that node cannot become v-node.

Consider the following 3 type of trails

$x \dots \leftarrow z \leftarrow \dots y$

$x \dots \rightarrow z \rightarrow \dots y$

$x \dots \leftarrow z \rightarrow \dots y$

~~Free~~  $x$  &  $y$  are  
d-separate

(equivalent transform)

If we transform<sup>T</sup> these type of trail we get ~~sometimes~~ the following 3 types of trail are possible.

$x \dots \leftarrow z \leftarrow y$

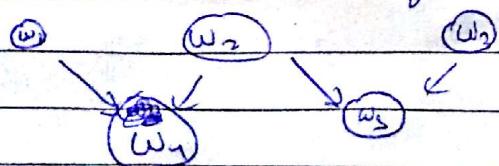
$x \dots \rightarrow z \rightarrow \dots y$

$x \leftarrow z \rightarrow y$

The resultant transform will also ~~to~~ satisfy d-separate rule.

We already know that v-nodes are preserved.

2. Consider a BN having v-node for ex

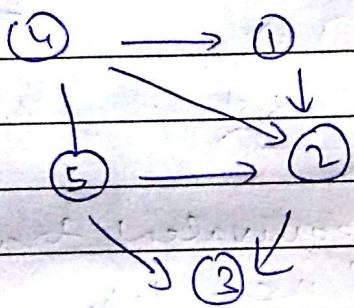


This type of graph will have  $O(n)$  edges

Do the transformation in the following way,

- (1) Pick the V-node one by one in topological order
  - (2) After all the V-node gets exchanged with the other node topological edge with adversarial ordering of vertices
- This transformation will cause  $O(n^2)$

For ex -> After transformation the previous graph would be



This holds for a general graph having V-nodes.