

1. What are the necessary characteristics of deadlocks? Explain deadlock prevention and avoidance methods.

→ The necessary characteristics of deadlock are:-

a. Mutual exclusion

→ At least one resource must be held in a non-shareable mode, meaning only one process can use it at a time.

b. Hold and wait

→ A process must be holding at least one resource and waiting to acquire additional resources held by other processes.

c. No preemption

→ Resources cannot be forcibly taken away from a process; they can only be released voluntarily.

d. Circular wait

→ There must exist a circular chain of two or more processes, each waiting for a resource held by the next process in the chain.

• Deadlock prevention

a. Mutual Exclusion

→ Ensure that resources that can lead to deadlock, especially those susceptible to mutual exclusion, are designed to allow multiple processes to share them.

b. Hold and Wait

→ A process must request and be allocated all its resources before it begins execution or release all resources if it cannot be allocated all at once. This ensures a process does not hold resources while waiting for others.

c. No preemption

→ If a process cannot get all its requested resources at once, it releases all the resources it currently holds, preventing the scenario where a resource cannot be forcibly forcibly taken from a process.

d. Circular wait

→ Impose a total ordering of resource types and require that each process requests resources in an increasing order. This way, a circular wait is not possible.

• Deadlock Avoidance.

a. Banker's Algorithm

→ The system keeps track of the maximum and currently allocated resources and calculates the remaining available resources. A process can only request resources if the system is left in a safe state after the allocation.

b. Resource Allocation Graph

→ Dynamically analyze the resource allocation state to determine if granting a resource request will result in a cycle in the Resource allocation graph. If no cycle is present, allocate the resource, otherwise, the request is delayed.

c. Timeouts and Killing Processes

→ Set timeouts for resource requests. If a process cannot acquire resources within a specified time it gets terminated or rolled back to safe state.

2. Draw a Gantt chart and find average turnaround time and waiting time of the following process applying FCFS, SJRF and round robin (with quantum = 3) scheduling algorithm.

Process	A	B	C	D	E
Arrival time (sec)	0	3	4	6	10
Burst time	6	3	6	4	2

For FCFS
Gantt chart

A	B	C	D	E
0(B,C,D)6	9(E)	15	19	21

Process	TAT = ET - AT	WT = TAT - BT
A	6	0
B	6	3
C	11	5
D	13	9
E	11	9
ΣTAT	47	$\Sigma WT = 26$

Average turnaround time = $\frac{47}{5} = 9.4$

Average waiting time = $\frac{26}{5} = 5.2 \text{ sec}$

For SRTF
Gantt chart

A	B	D	E	D	C
0(B,C,D)6	9(E)	10	12	15	21

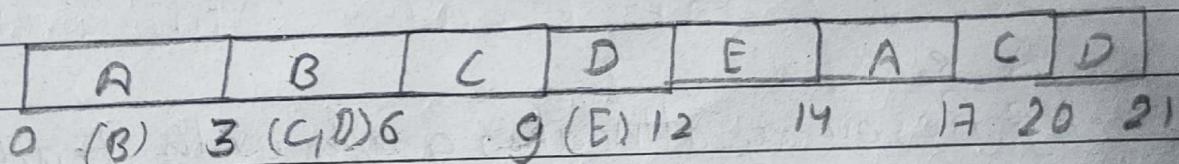
Process	$TAT = ET - AT$	$WT = TAT - BT$
A	6	0
B	6	3
C	17	11
D	9	5
E	2	0
$n=5$	$\sum TAT = 40$	$\sum WT = 19$

Average turnaround time = $\frac{40}{5} = 8 \text{ sec}$

Average waiting time = $\frac{19}{5} = 3.3 \text{ sec.}$

- Round Robin
quantum = 3

Gantt chart :-



Process	$TAT = ET - AT$	$WT = TAT - BT$
A	17	11
B	3	0
C	16	10
D	15	11
E	4	2
$n=5$	$\sum TAT = 55$	$\sum WT = 34$

Average turn around time = $\frac{55}{5} = 11$ sec

Average waiting time = $\frac{37}{5} = 6.8$ sec

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Q1. ~~Why is deadlock a state more critical than starvation?~~ Describe resource allocation graph with a deadlock, with a cycle but no deadlock.

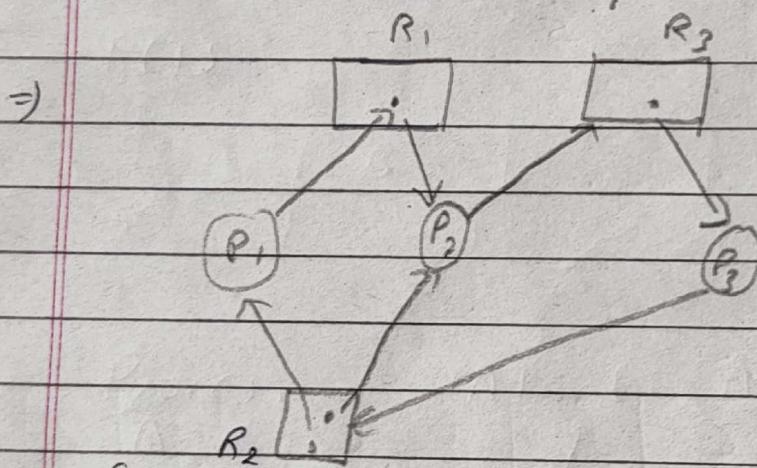


fig: Resource allocation graph with deadlock.

Here cycle is created between R_2, P_2, R_3, P_3 and R_2 where R denote resources and P denote processes.

In this figure, Resource (R_2) is allocated to process (P_2) and P_2 requests resource (R_3), also R_3 is allocated to P_3 and P_3 requests resource

R_2 . At this condition cycle is created and all the processes hold their resources until completion and hence deadlock occurs.

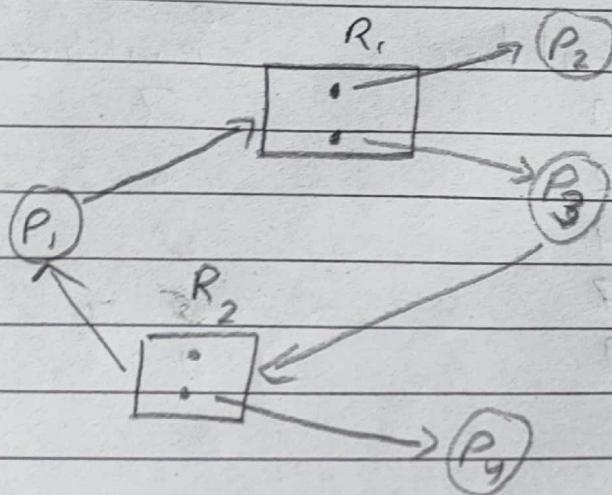


Fig: - Resource allocation graph with cycle but no deadlock.

Here cycle is created between process (P_1, P_3) and Resources (R_2, R_1)

Here, Resource (R_2) is allocated to process (P_1) and P_1 requests for Resource (R_1). But, R_1 can be held by process P_2 and Process P_1 can be executed.

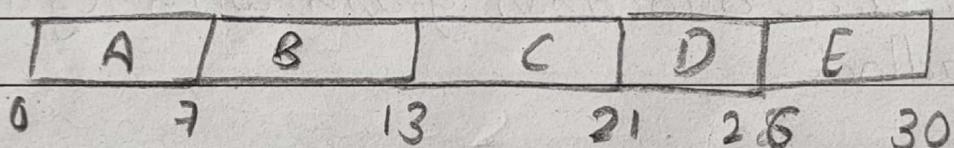
Now Process P_1 holds resources (R_1, R_2). Process (P_2) wants request for resource R_2 and Resource (R_1) is allocated to P_3 so, P_3 can be executed either released by P_1 or P_2 after completion. In this way, deadlock does not occur.

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2. From the following set of information. find the average waiting time and average turn around time using FCFS, SJF, RR (Quantum = 3) and HRRN.

Process	Arrival time	Service time (burst time)
A	0	7
B	2	6
C	4	8
D	7	5
E	9	4

- FCFS



Process	TAT = ET - AT	WT = TAT - BT
A	7	0
B	11	5
C	17	9
D	19	14
E	21	17

$n=5$ $\sum TAT = 75$ $\sum WT = 45$

- Average turnaround time = $\frac{75}{5} = 15$

- Average waiting time = $\frac{45}{5} = 9$

• SJF

	A	D	E	B	C
o (B,C,D) \neq (E)	7	12	16	22	30

Process	TAT = ET - AT	WT = TAT - BT
A	7	0
B	20	14
C	26	18
D	5	0
E	7	3
n = 5	$\sum TAT = 65$	$\sum WT = 35$

• ~~Total~~ Average Turnaround time = $\frac{\sum TAT}{n} = \frac{65}{5} = 13$

• Average waiting time = $\frac{\sum WT}{n} = \frac{35}{5} = 7$

• RR

quantum = 3

A	B	C	D	E	A	B	C	D	E
0	3	6	9	12	15	18	21	24	26

A	C
27	28

Process	$TAT = ET - AT$	$WT = TAT - BT$
A	28	21
B	19	13
C	26	18
D	19	14
E	18	14
$n = 5$	$\sum TAT = 110$	$\sum WT = 80$

Average turnaround time $= \frac{110}{5} = 22$

Average waiting time $= \frac{80}{5} = 16$

3. Consider the following page reference string S: 1, 2, 3, 4, 2, 1, 5, 6, 2, 1, 2, 3, 1, 7, 6, 3, 2, 1, 1, 2, 3, 6.

How many page faults would occur for each of the following page replacement algorithms assuming 3 pages a frame? In each case calculate fault ratio.

- i) LRU page replacement
- ii) FIFO page replacement
- iii) Second chance page replacement

① Second chance page replacement

Page :

Date 2 / 1 /

1	2	3	4	2	5
0	0	3	3	3	1
2	2	2	2	2	2
1	1	1	4	4	4
Miss	Miss	Miss	Miss	Hit	Miss

6	2	1	2	3	7	6	3
1	2	2	2	2	2	6	6
6	6	6	6	3	3	3	3
5	5	1	1	1	7	7	7
Miss	Miss	Miss	Hit	Miss	Miss	Miss	Hit

2	1	2	3	6
6	1	1	1	6
3	1	3	1	3
2	2	2	1	2
Miss	Miss	Hit	Hit	Miss

Total No. of page fault = 15

ii) LRU page replacement

1	2	3	4	2	1	5	6
1	2	3	3	3	1	1	1
	1	2	2	2	2	2	6
Miss	Miss	Miss	Miss	Hit	Miss	Miss	Miss
2	2	2	2	2	6	6	6
6	6	6	3	3	3	3	3
5	1	1	1	7	7	7	2
Miss	Miss	Hit	Miss	Miss	Miss	Hit	Miss

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1	2	3	6
6	6	3	3
1	1	1	6
2	2	2	2

Miss Hit Miss Miss

Total no. of page faults = 16

iii) ~~FIFO~~ page replacement

1	2	3	4	2	1	5	6	2	1
				3	3	3	5	5	5
	2	2	2	2	1	1	1	2	2
1	1	1	4	4	4	4	6	6	6

Miss Miss Miss Miss Hit Miss Miss Miss MBS Miss

2	3	7	6	3	2	1	3	2	6
1	1	1	6	6	6	6	3	3	3
2	2	7	7	7	7	1	1	1	1
6	3	3	3	3	2	2	2	2	6

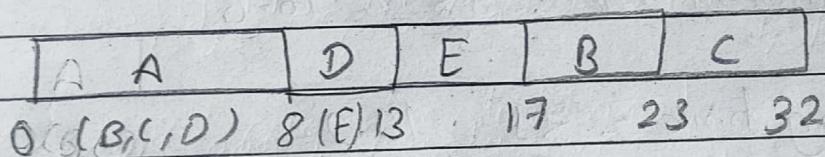
Hit Miss Miss Miss Hit Miss Miss Miss Hit Miss

Total no. of page faults = 16

1 Given the following set of information, what is the average waiting time and average turn-around time using SJF Preemptive, FCFS, RR (Quantum = 3) and HRRN.

Process	Arrival Time	Service Time (Burst Time)
A	0	8
B	2	6
C	4	9
D	7	5
E	9	4

- SJF preemphire



Process	TAT = ET - AT	WT = TAT - RT
A	8	0
B	21	15
C	28	19
D	6	1
E	8	4

- Average Turnaround time = $\frac{71}{5} = 14.2$

- Average waiting time = $\frac{39}{5} = 7.8$

→ FCFS

	A	B	C	D	E	
	0	8	14	23	28	32

Process

$$TAT = ET - AT$$

$$WT = TAT - BT$$

A

8

0

B

12

6

C

19

10

D

21

16

E

23

19

n=5

$$\sum TAT = 83$$

$$\sum WT = 51$$

• Average Turnaround time = $\frac{83}{5} = 16.6$

• Average waiting time = $\frac{51}{5} = 10.2$

→ RR

Quantum=3

	A	B	C	D	E	A	B	C	D	E	A	C	
	0	3	6	9	12	15	18	21	24	26	27	29	32

Process

$$TAT = ET - AT$$

$$WT = TAT - BT$$

A

29

21

B

19

13

C

28

19

D

19

14

E

18

14

$$\sum TAT = 113$$

$$\sum WT = 81$$

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• Average waiting time = $\frac{\sum TAT}{n} = \frac{113}{5} = 22.6$

• Average waiting time = $\frac{\sum WWT}{n} = \frac{81}{5} = 16.2$

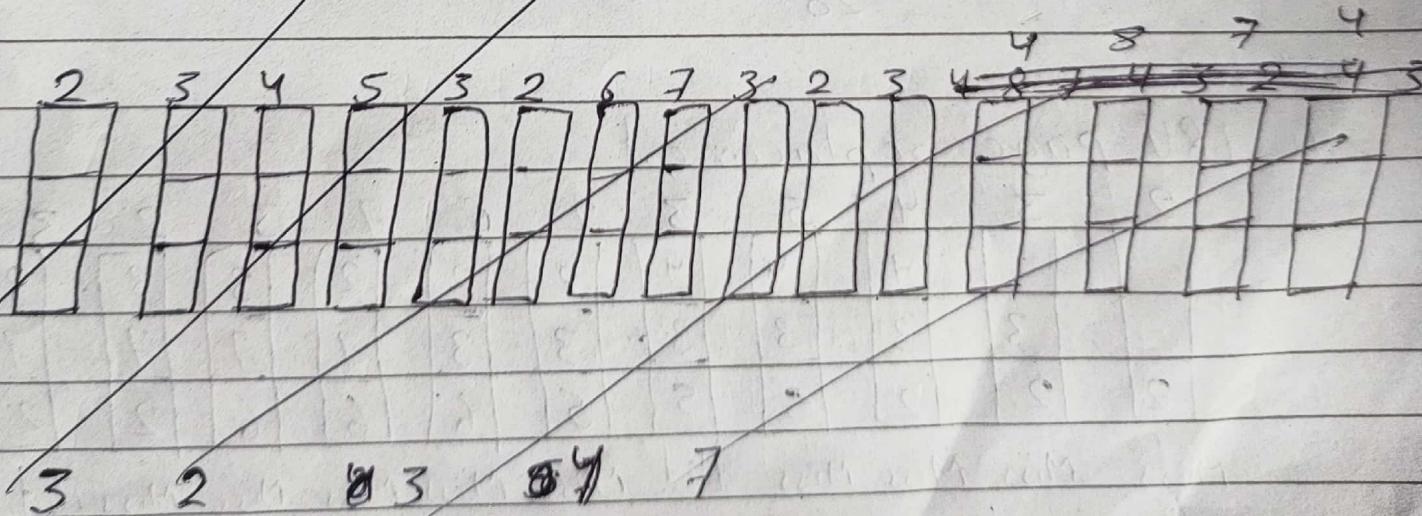
2. Consider the following page reference strings:-

2, 3, 4, 5, 3, 2, 6, 7, 3, 2, 3, 4, 8, 7, 4, 3, 2, 3, 4, 7.

How many page faults would occur for each of the following page replacement algorithms assuming 3 pages a frame? In each case calculate fault ratio.

- Optimal Page replacement (not used for longest period of time)
- LRU page replacement
- FIFO page replacement
- ...

i) Optimal page Replacement



i) Optimal page replacement

2	3	4	5	3	2	6	7	3	2	3	4
3	3	3	3	3	3	3	7	7	7	7	4
2	2	2	5	5	5	6	6	6	2	2	2

Miss Miss Miss Miss Hit Miss Miss Miss Miss Hit Miss

8	7	4	3	2	3	4	7				
3	7	7	7	7	2	2	2	7			
4	4	4	4	4	4	4	4	4			
8	8	8	3	3	3	3	3	3			

Miss Miss Hit Miss Miss Hit Hit Miss

- Total No. of page faults = 15
- Fault ratio = $\frac{15}{20} = 3:4$

ii) LRU page replacement

2	3	4	5	3	2	6	7	3	2	3	4
3	3	3	3	3	3	3	3	7	7	7	4
2	2	2	5	5	5	6	6	6	2	2	2

Miss Miss Miss Miss Hit Miss Miss Miss Miss Hit Miss

8

8	7	4	3	2	3	4	7
3	7	7	8	2	2	2	7
4	4	4	4	4	4	4	4
8	8	8	3	3	3	3	3

Miss Miss Hit Miss Miss Hit Hit Miss

- Total no. of page fault = 15
- Page fault ratio = $\frac{15}{20} = 3:4$

iii FIFO page replacement.

2	3	4	5	3	2	6	7	3	2	3	4
	*	4	4	4	4	6	6	6	2	2	2
	3	3	3	3	2	2	3	3	3	3	3
2	2	2	5	5	5	7	7	7	7	7	4

Miss Miss Miss Miss Hit Miss Miss Miss Miss Hit Miss

8	7	4	3	2	3	4	7
2	7	7	7	7	7	4	4
8	8	8	8	2	2	2	2

Miss Miss Hit Miss Miss Hit Miss Miss

Total no. of page fault = 16

Page Fault ratio = $\frac{16}{20} = 4:5$

3 Five processes and 3 resource types A, B, C and D (Below is the snapshot of the state)

Process	Max	Allocation	Available
P0	A, B, C, D	A, B, C, D	A, B, C, D
P1	P0 6 0 1 2	4 0 0 1	3 2 1 1
P2	P1 2 7 5 0	1 1 0 0	
P3	P2 2 3 5 6	1 2 5 4	
P4	P3 1 6 5 3	0 6 3 3	
P5	P4 1 6 5 6	0 2 1 2	

Is this a safe state? If yes, what is safe sequence?

Page :
Date : / /

1 Consider the deadlock situation that could occur in the dining philosophers problem when the philosopher obtain the chopstick one at a time. Discuss how the four necessary conditions for deadlock indeed hold in this setting. Discuss how deadlock could be avoided by eliminating any one of the four conditions.

⇒ In the dining philosophers problem, a deadlock can occur when each philosopher picks up one chopstick and waits for the other, resulting in a circular waiting scenario. The four necessary condition for deadlock in this situation are:-

- a. Mutual Exclusion:- Philosophers must hold a chopstick exclusively to themselves, creating mutual exclusion.
- b. Hold and wait :- A philosopher can hold one chopstick while waiting for another, contributing to the hold and wait condition.
- c. No preemption:- Once a philosopher holds a chopstick, it cannot be taken away forcibly, adhering to the no preemption condition.
- d. Circular wait :- Philosophers from a circular dependency while waiting for a chopstick, satisfying the circular wait condition.

To avoid deadlock, one of the four conditions must be eliminated. Here are potential solutions:-

1. Mutual Exclusion:- Allow philosophers to share chopsticks.

This eliminates the need for exclusive access and break the mutual exclusion condition.

2. Hold and wait :- Require philosophers to acquire all necessary resources (both chopsticks) simultaneously or release the acquired resources if unable

to get the second chopstick. This way the hold and wait condition is eliminated.

3. No preemptions:- Allow a philosopher to be interrupted and force to release their chopsticks if another philosopher requests them urgently. However this

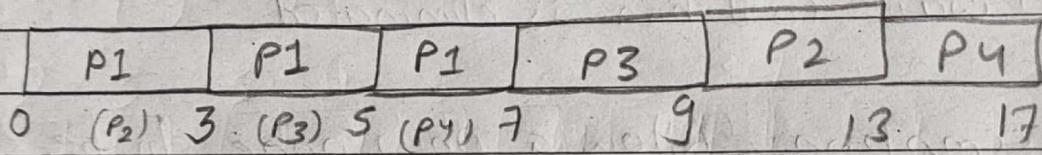
4. Circular wait :- Assign a unique identifier to each chopstick and require philosophers to always pick up chopsticks in a specific order. This breaks the circular dependency.

2. Given the following information, draw a GANTT chart for process scheduling. Preemptive shortest job first and RR (quantum = 2). Also, find average waiting time, average turnaround time and average

response time for all cases

Process	Arrival time	Burst time
P1	0.0	7
P2	3.0	4
P3	5.0	2
P4	6.0	9

- Preemptive shortest job first
GANTT chart



Start time of execution

$$\text{Process} \quad TAT = ET - AT \quad WT = TAT - BT \quad R.T = S.T.E - A.T.$$

P1	7	0	0
P2	10	6	6
P3	4	2	2
P4	11	7	7
	$\sum TAT = 32$	$\sum WT = 15$	$\sum RT = 15$

$$\text{Average turn around time} = \frac{\sum TAT}{n} = \frac{32}{4} = 8$$

$$\text{Average waiting time} = \frac{\sum WT}{n} = \frac{15}{4} = 3.75$$

$$\text{Average response time} = \frac{\sum RT}{n} = \frac{15}{4} = 3.75$$

Round Robin
(Quantum - 2)

GANTT CHART

P ₁	P ₁	P ₂	P ₃	P ₄	P ₁	P ₂	P ₄	P ₁
0	2 (P ₂)	4 (P ₃ , P ₄)	6	8	10	12	14	16

Process	TAT = ET - AT	WAT = TAT - BT	R.T. = ST - AT
P ₁	17	10	0
P ₂	11	7	1
P ₃	3	1	1
P ₄	10	6	2
$\Sigma TAT = 41$	$\Sigma WAT = 24$	$\Sigma R.T. = 4$	

Average Turnaround time = $\frac{\Sigma TAT}{n} = \frac{41}{4} = 10.25$

Average waiting time = $\frac{\Sigma WAT}{n} = \frac{24}{4} = 6$

Average Response time = $\frac{\Sigma R.T.}{n} = \frac{4}{4} = 1$

3 Consider the following page reference strings:
~~2, 3, 3, 4, 5, 6, 5, 7, 1, 2, 5, 8, 6, 4, 1~~. How many page faults ~~would~~ would occur for each of the following page replacement algorithms assuming 3 pages frames?

- i) LRU page replacement
- ii) FIFO page "
- iii) Optimal page "

- LRU page replacement

2	3	3	4	5	6	5	7	1	2	5	8
3	3	3	3	3	8	8	7	7	7	2	2
2	2	2	5	5	5	5	6	6	1	1	8

Miss Miss Hit Miss Miss Miss Hit Miss Miss Miss Miss Miss

6	4	1
5	4	1
6	6	6
8	8	1

Miss Miss Miss

Total no. of page faults = 13

(ii) FIFO

2	3	3	4	5	6	5	7	1	2	5	8	6	4	1
			4	4	4	4	3	7	7	5	5	5	4	4
.	3	3	3	3	6	6	6	6	2	2	8	8	6	6
2	2	2	2	5	5	5	1	1	1	8	8	8	1	1

Miss Miss Hit Miss Miss Miss Hit Miss Miss Miss Miss Miss Miss Miss Miss

Total no. of page fault = 15

(iii) Optimal page replacement

2	3	3	4	5	6	5	7	1	2
			4	4	4	4	7	1	1
.	3	3	3	3	5	5	5	5	2

Miss Miss Hit Miss Miss Miss Hit Miss Miss Miss Miss

5	8	6	4	1
1	1	1	1	1
5	8	8	8	8

Miss Miss Hit Miss Hit

Total no. of page fault = 13

1 Consider the following set of processes with the length of the CPU burst given in milliseconds. Construct the Gantt chart and calculate average waiting time and turnaround time. Also explain which one is the best algorithm.

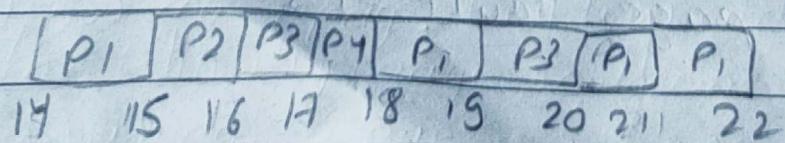
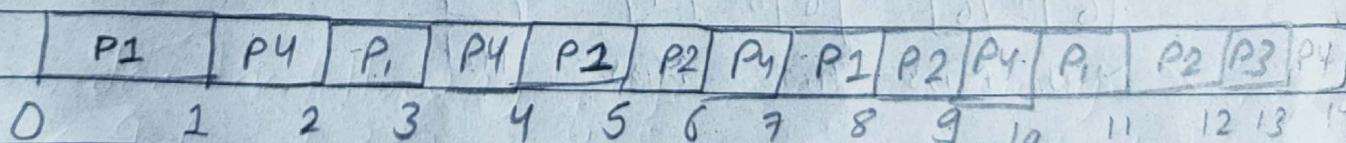
Process	Arrival time	Burst time	Priority
P1	0	9	2
P2	4	4	1
P3	10	3	3
P4	0	6	4

- (i) RR (Quantum = 1ms)
- (ii) Priority scheduling (1 higher priority)
- (iii) Shortest job first (preemptive)

(i) RR

Quantum = 1ms

GANTT chart

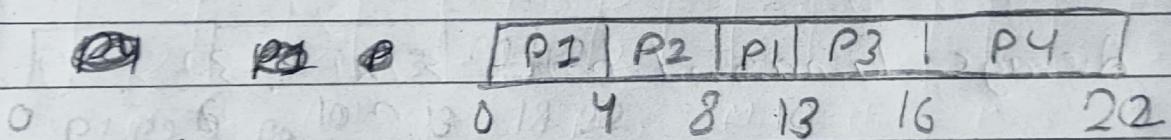


Process	$TAT = ET - AT$	$WT = TAT - BT$
P1	22	13
P2	12	8
P3	10	7
P4	18	12
	$\Sigma TAT = 62$	$\Sigma WT = 40$

Average turnaround time = $\frac{\Sigma TAT}{4} = \frac{62}{4} = 15.5$

Average waiting time = $\frac{\Sigma WT}{4} = \frac{40}{4} = 10$

(ii) Priority scheduling (P higher priority)



Process	$TAT = ET - AT$	$WT = TAT - BT$
P1	13	4
P2	4	0
P3	6	3
P4	22	16
	$\Sigma TAT = 45$	$\Sigma WT = 23$

Average Turnaround time = $\frac{45}{4} = 11.25$

Average waiting time = $\frac{23}{4} = 5.75$

(iii) Shortest job first (preemptive)

P	P4	P2	P3	P1
0	6	10	13	22
P2, P2				

Process	TAT = ET - AT	WT = TAT - BT
P1	22	13
P2	6	2
P3	3	0
P4	6	0
	$\sum TAT = 37$	$\sum WT = 15$

Average Turnaround time = $\frac{\sum TAT}{n} = \frac{37}{4} = 9.25$

Average waiting time = $\frac{15}{4} = 3.75$

2. Consider the following page reference strings 9, 3, 4, 5, 3, 9, 6, 7, 3, 9, 3, 4, 18, 7, 4, 3, 9, 3, 4, 7. How many page faults would occur for each of the following page replacement algorithms assuming 3 pages a frame? In each case calculate fault ratio.

- i) Second chance page replacement

- ii LRU page replacement
 - iii FIFO page replacement

i) Second chance page replacement

9	3	4	5	3	9	6	7	3	9	3	4	8
.			4	4	9	9	9	3	3	3	1	3
	3	3	3	3	1	3	3	7	7	7	7	4
9	9	9	5	5	5	6	6	6	9	9	9	8
M	M	M	M	M	M	M	M	M	M	H	M	M

7	4	3	9	3	4	7
7	3	7	7	7	4	7
4	4	4	9	9	9	9
8	8	3	3	3	3	3

Total no. of page fault = 17

ii) LRU page replacement

9	3	4	5	3	9	6	7	3	9	3	4
	3	3	3	3	3	3	7	7	7	7	4
9	9	9	5	5	5	6	6	6	9	9	9

8	7	4	3	9	3	4	7
3	7	7	7	9	9	9	7
4	7	4	4	4	4	4	4
8	8	8	3	3	3	3	3

M M H M M H H M

Total no. of page faults = 15

ii. FIFO page replacement

9	3	4	5	3	9	6	7	3	9	3
3	3	3	3	3	9	9	6	6	3	3
9	9	9	5	5	5	5	7	7	7	7

M M M M H M M M M M H

4	8	7	4	3	9	3	4	7
9	9	7	7	7	7	7	4	7
3	8	8	8	8	9	9	9	9

M M M H M M H M M M

Total no. of page faults = 16

Page fault ratio = $\frac{16}{20} = 4:5$