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# PRACTICAL1: MATRIX MULTIPLICATION

## INPUT:

```
print("Enter order of 1st matrix:")
a,b = list(map(int,input().split()))

print("Enter Row wise values")

mat1 = []

for i in range(a) :
    print("Enter row",i,"value:")

    row = list(map(int,input().split()))
    mat1.append(row)

print("Enter order of 2nd matrix:")

d,c = list(map(int,input().split()))

print("Enter Row wise values")

mat2 = []

for j in range(d) :
    print("Enter row",j,"value:")

    row = list(map(int,input().split()))
    mat2.append(row)

print("Matrix 1:",mat1)
print("Matrix 2:",mat2)

resultant = []

for i in range(a):
    row = []
    for j in range(c):
        row.append(0)

    resultant.append(row)
```

```
print("Matrix Multiplication: ")
for i in range(a):
    for j in range(c):
        for k in range(b) :
            resultant[i][j] += mat1[i][k] * mat2[k][j]

for row in resultant:
    print(row)
```

## OUTPUT:

```
Enter order of 1st matrix:
Enter row 0 value:
1 3 3
Enter row 1 value:
2 1 9
Enter row 2 value:
4 1 9
Matrix 1: [[1, 2, 3], [1, 2, 3], [1, 2, 3]]
Matrix 2: [[1, 3, 3], [2, 1, 9], [4, 1, 9]]
Matrix Multiplication:
[17, 8, 48]
[17, 8, 48]
[17, 8, 48]
```

# PRACTICAL 2: QUICK SORT

## INPUT:

```
def partition(array, low, high):

    pivot = array[high]

    i = low - 1

    for j in range(low, high):
        if array[j] <= pivot:
            i = i + 1

            (array[i], array[j]) = (array[j], array[i])

    (array[i + 1], array[high]) = (array[high], array[i + 1])

    return i + 1

def quickSort(array, low, high):
    if low < high:

        pi = partition(array, low, high)
        quickSort(array, low, pi - 1)
        quickSort(array, pi + 1, high)

data = [2, 1, 9, 4, 5, 9, -2]
print("Unsorted Array")
print(data)

size = len(data)

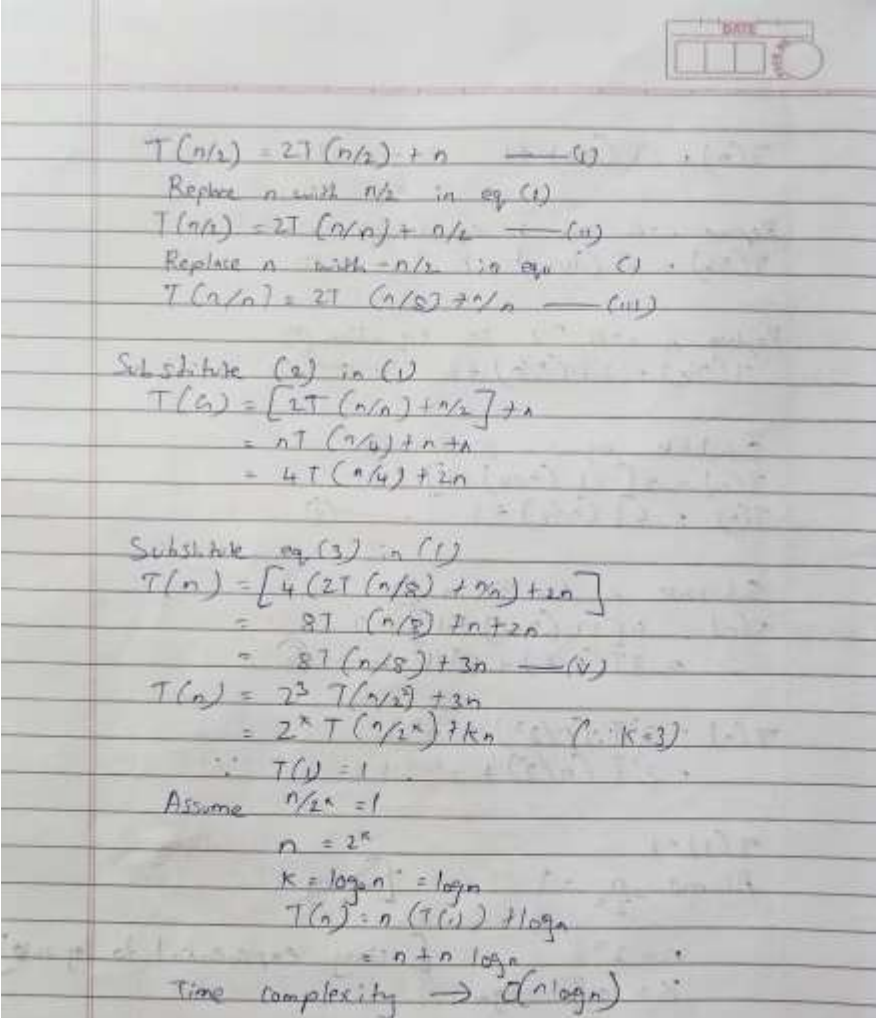
quickSort(data, 0, size - 1)

print('Sorted Array in Ascending Order:')
print(data)
```

## OUTPUT:

Unsorted Array  
[2, 1, 9, 4, 5, 9, -2]  
Sorted Array in Ascending Order:  
[-2, 1, 2, 4, 5, 9, 9]

## Time Complexity:



Handwritten derivation of the time complexity of Quick Sort:

$$T(n/2) = 2T(n/2) + n \quad \text{--- (i)}$$

Replace  $n$  with  $n/2$  in eq. (i)

$$T(n/2) = 2T(n/2) + n/2 \quad \text{--- (ii)}$$

Replace  $n$  with  $n/2$  in eq. (i)

$$T(n/2) = 2T(n/2) + n/2 \quad \text{--- (iii)}$$

Substitute (ii) in (i)

$$T(n) = [2T(n/2) + n/2] + n$$
$$= nT(n/2) + n + n$$
$$= 4T(n/4) + 2n$$

Substitute eq. (iii) in (i)

$$T(n) = [4(2T(n/2) + n/2) + 2n]$$
$$= 8T(n/2) + 3n \quad \text{--- (iv)}$$
$$T(n) = 2^3 T(n/2^3) + 3n$$
$$= 2^k T(n/2^k) + kn \quad (\because k=3)$$

$\therefore T(1) = 1$

Assume  $n/2^k = 1$

$$n = 2^k$$
$$k = \log_2 n = \log n$$
$$T(n) = n(T(1)) + \log n$$
$$= n + n \log n$$

Time complexity  $\rightarrow O(n \log n)$

# PRACTICAL 3: MERGE SORT

## INPUT:

```
def merge(arr, l, m, r):
    n1 = m - l + 1
    n2 = r - m

    L = [0] * (n1)
    R = [0] * (n2)

    for i in range(0, n1):
        L[i] = arr[l + i]

    for j in range(0, n2):
        R[j] = arr[m + 1 + j]

    i = 0
    j = 0
    k = l
    while i < n1 and j < n2:
        if L[i] <= R[j]:
            arr[k] = L[i]
            i += 1
        else:
            arr[k] = R[j]
            j += 1
        k += 1

    while i < n1:
        arr[k] = L[i]
        i += 1
        k += 1

    while j < n2:
        arr[k] = R[j]
        j += 1
        k += 1

def mergeSort(arr, l, r):
    if l < r:

        m = l+(r-l)//2

        mergeSort(arr, l, m)
        mergeSort(arr, m+1, r)
        merge(arr, l, m, r)

arr = [4, 1, 5, 9, 2, 7]
n = len(arr)
print("Given array is")
for i in range(n):
    print("%d" % arr[i],end=" ")

mergeSort(arr, 0, n-1)
print("\n\nSorted array is")
for i in range(n):
    print("%d" % arr[i],end=" ")
```

## OUTPUT:

```
Given array is
4 1 5 9 2 7

Sorted array is
1 2 4 5 7 9
```

## Time Complexity:

Handwritten derivation of Merge Sort time complexity:

$$T(n) = 2T(n/2) + O(n)$$
$$T(n) = 2T(n/2) + O(n)$$
$$T(n) = 2(2T(n/4) + C \times n/2) + O(n)$$
$$T(n) = n(T(n/4) + 2 \times C \times n)$$
$$T(n) = 2^k \cdot T(n/2^k) + k \times C \times n$$
$$T(n) = NT(1) \neq N \times \log N$$

Time complexity =  $O(N \log N)$

# PRACTICAL 4: LINEAR SEARCH

## INPUT:

```
def linearSearch(array, n, x):  
  
    for i in range(0, n):  
        if (array[i] == x):  
            return i  
    return -1  
  
array = [2, 4, 0, 1, 9]  
x = 1  
n = len(array)  
result = linearSearch(array, n, x)  
if(result == -1):  
    print("Element not found")  
else:  
    print("Element found at index: ", result)
```

## OUTPUT:

```
Element found at index: 3
```

## Time Complexity:

$$\begin{aligned}T(n) &= 2T(n/2) + 1 \\ \text{Replace } n \text{ with } n/2 \\ T(n/2) &= 2T(n/4) + 1 \\ &= 2T(n/8) + 1 \\ &= 2T(n/n) + 1 \\ \text{Replace } n \text{ with } n/n \\ T(n/n) &= 2T(n/8) + 1 \\ T(n) &= 2[2T(n/8) + 1] + 1 \\ &= 4T(n/8) + 3 \\ T(n) &= n(2(1(n/8) + 1) + 3) \\ &= 8T(n/8) + 7 \\ &= 2^3 T(n/2^3) + (2^3 - 1) \\ &= 2^k T(n/2^k) + (2^k - 1) \\ K &= \log n \\ \text{Assume } n/2^k &= 1 \\ T(n) &= 2^k T(1) + (2^k - 1) \\ &= n + (n - 1) \\ &= n + n - 1 \\ T(n) &= 2n - 1 \\ &= O(n)\end{aligned}$$

# PRACTICAL 5

## Time Complexity:

### INPUT:

```
def binary_search(arr, low, high, x):  
  
    if high >= low:  
        mid = (high + low) // 2  
  
        if arr[mid] == x:  
            return mid  
  
        elif arr[mid] > x:  
            return binary_search(arr, low, mid - 1, x)  
  
        else:  
            return binary_search(arr, mid + 1, high, x)  
  
    else:  
        return -1  
  
arr = [ 2, 3, 4, 10, 40 ]  
x = 10  
  
result = binary_search(arr, 0, len(arr)-1, x)  
  
if result != -1:  
    print("Element is present at index", str(result))  
else:  
    print("Element is not present in array")
```

Handwritten derivation of the time complexity for binary search:

$$\begin{aligned} T(N) &= c + T(N/2) \quad \text{--- (I)} \\ T(N/2) &= c + T(N/4) \quad \text{--- (II)} \\ \text{Sub (2) in (I)} \\ T(N) &= T(N/4) + 2c \quad \text{--- (III)} \\ T(N/4) &= c + T(N/8) \quad \text{--- (IV)} \\ \text{Sub (4) in (3)} \\ T(N) &= T(N/8) + 3c \quad \text{--- (V)} \\ T(N) &= T(N/2^k) + kc \quad \text{--- (VI)} \\ k=3 \quad T\left(\frac{N}{2^k}\right) &= T(1) \\ \text{Sub (VI) in (V)} \quad \frac{N}{2^k} &= 1 \\ T(N) &= T\left(\frac{N}{2^{\log_2 N}}\right) + c^{\log_2 N} \quad N = 2^k \\ &= T(1) + c \log_2 N \quad \log_2 N = k \\ &= T(1) + c \log_2 N \\ T(N) &= k + c \log_2 N \\ \text{Time complexity } &O(\log N). \end{aligned}$$

### OUTPUT:

Element found at index: 3

## PRACTICAL NO:-06

**Aim:-**Write a program to implement insertion operation on Binary Search tree and discuss the time complexity of the algorithm. Code:-

```
class Node:
    def __init__(self,index):
        self.left=None
        self.right=None
        self.val=index
```

```
def insert(root,newnode):
    if root is None:
        root=newnode
    else:
        if root.val<=newnode.val:
            if root.right is None:
                root.right=newnode
            else:
                insert(root.right,newnode)
        elif root.val==newnode.val:
            print("Already existing node ",newnode.val)
        else:
            if root.left is None:
                root.left=newnode
            else:
                insert(root.left,newnode)

def inorder(root):
    if root:
        if root.val==None:
```

```
print(end="")
else:
    inorder(root.left)
print(root.val)
inorder(root.right)
```

```
def preorder(root):
    if root:
        if
root.val==None:
    print(end="")
else:
    print(root.val)
    preorder(root.left)
    preorder(root.right)
```

```
def postorder(root):
    if root:
        if root.val==None:
            print(end="")
        else:
            print(root.val)
            postorder(root.left)
            postorder(root.right)
root=Node(100)
insert(root,Node(60))
insert(root,Node(50))
insert(root,Node(90))
```



```
insert(root,Node(40))
insert(root,Node(53))
insert(root, Node(95))
insert(root, Node(75))
```

```
print("inorder Traversal:")
inorder(root)
print("preorder Traversal:")
preorder(root)
print("postorder Traversal:")
postorder(root)
```

```
print("\n")
```

### #Time Complexity:-

The time complexity for inserting node **depends on the height of the tree h** , so  $O(h)$

## Output:-

---

inorder Traversal:

40  
50  
53  
60  
75  
90  
95  
100

preorder Traversal:

100  
60  
50  
40  
53  
90  
75  
95

postorder Traversal:

100  
60  
50  
40  
53  
90  
75  
95

## PRACTICAL NO:-07

**Aim:-** Write a python program to delete node from the binary tree from a given data and discuss the time complexity. Code:- class Node: def

```
__init__(self,index):
```

```
    self.left=None
```

```
self.right=None
```

```
self.val=index
```

```
def insert(root,newnode):    if
```

```
root is None:
```

```
root=newnode    else:        if
```

```
root.val<=newnode.val:
```

```
if root.right is None:
```

```
root.right=newnode
```

```
else:
```

```
    insert(root.right,newnode)
```

```
elif root.val==newnode.val:
```

```
    print("Already existing node ",newnode.val)
```

```
else:
```

```
    if root.left is None:
```

```
root.left=newnode
```

```
else:
```

```
    insert(root.left,newnode)
```

```
def inorder(root):    if
```

```
root:        if
```

```
root.val==None:
```

```

print(end="")
else:
    inorder(root.left)
print(root.val)
inorder(root.right)

```

```

def preorder(root):
    if root:
        if
root.val==None:
print(end="")
else:
    print(root.val)
preorder(root.left)
preorder(root.right)

```

```

def postorder(root):
    if root:
        if root.val==None:
print(end="")
else:
    print(root.val)
postorder(root.left)
postorder(root.right)
def
delete(root,node1):
    if root
is None:
    print("Empty
tree")
    elif root.val<node1:
delete(root.right,node1)

```

```
elif root.val>node1:
delete(root.left,node1)
else:
    root.val=None
```

```
root=Node(100)
insert(root,Node(60))
insert(root,Node(50))
insert(root,Node(90))
insert(root,Node(40))
insert(root,Node(53))
insert(root, Node(95))
insert(root, Node(75))
```

```
print("Before Deletion Operation: \n")
print("inorder Traversal:")
inorder(root) print("preorder
Traversal:") preorder(root)
print("postorder Traversal:")
postorder(root)
```

```
print("\n")
```

```
delete(root,40) delete(root,90)
print("After Deletion Operation: \n")
print("inorder Traversal:")
```

```
inorder(root) print("preorder  
Traversal:") preorder(root)  
print("postorder Traversal:")  
postorder(root)
```

### #Time Complexity:-

The time complexity for deleting node , so  $O(n)$

## Output:-

```
>>>
Before Deletion Operation:
inorder Traversal:
40
50
53
60
75
90
95
100
preorder Traversal:
100
60
50
40
53
90
75
95
postorder Traversal:
100
60
50
40
53
90
75
95
After Deletion Operation:
inorder Traversal:
50
53
60
100
preorder Traversal:
100
60
50
53
postorder Traversal:
100
60
50
53
>>>
```

---

## PRACTIACAL NO:-08

**Aim:-** Write a python program to implement Breadth First Traversal of graph.

Code:- def bfs(visited, graph, node):

visited.append(node)

queue.append(node)

while queue:

m=queue.pop(0)

print(m, end=" ")

for neighbour in graph[m]:

if neighbour not in visited:

visited.append(neighbour)

queue.append(neighbour)

graph = {

'A': ['B', 'D', 'G'],

'B': ['A', 'C', 'D'],

'C': ['B', 'E', 'F'],

'D': ['A', 'B', 'E'],

'E': ['D', 'C', 'F', 'G'],

'F': ['C', 'E', 'H'],

'G': ['A', 'E', 'H'],

'H': ['G', 'F'],

}

visited=[] queue=[]



```
print("Following is Breadth-First Search") bfs(visited,graph,'A')
```

## Output:-

```
>>>
Following is Breadth-First Search
A B D G C E H F
\\ \ \ |
```

## PRACTICAL NO:-09

**Aim:-** Write a python program to implement Depth First Traversal of graph.

**Code:-**

```
# Using a Python dictionary to act as an adjacency list graph
```

```
= {
    'A': ['B', 'D', 'G'],
    'B': ['A', 'C', 'D'],
    'C': ['B', 'E', 'F'],
    'D': ['A', 'B', 'E'],
    'E': ['D', 'C', 'F', 'G'],
    'F': ['C', 'E', 'H'],
    'G': ['A', 'E', 'H'],
    'H': ['G', 'F']}
```

```
visited = set() # Set to keep track of visited nodes of graph.
```

```
def dfs(visited, graph, node): #function for dfs
```

```
if node not in visited: print (node,end=" ")
```

```
visited.add(node) for neighbour in
```

```
graph[node]:      dfs(visited, graph,  
neighbour)
```

# Driver Code

```
print("Following is the Depth-First Search") dfs(visited,  
graph, 'A')
```

Output:-

```
>>>  
Following is the Depth-First Search  
A B C E D F H G  
>>> |
```

PRACTIACAL NO:-10

**Aim:-** Write python program for checking whether a given graph g has a simple path from source s to destination d  
**Code:-** graph={'A':['B','C'],

'B':['C','D'],

'C':['D'],

'D':['C'],

'E':['F'],

'F':['C']}

```
def find_all_paths(graph,start,end,path=[]):
```

```
    path=path+[start]#source vertex included in path
```

```
if start==end:    return[path]    if start not in
```

```
graph:
```

```
    print("start vertex is not present in the graph")
```

```
return None    paths=[]
```

```

    for node in graph[start]:#iterating through graph
if node not in path:
    newpaths=find_all_paths(graph,node,end,path)
for newpath in newpaths:

    paths.append(newpath)
return paths

```

```
find_all_paths(graph,'A','D')
```

**Output:-**

```

>>>
>>> find_all_paths(graph, 'A', 'D')
[['A', 'B', 'C', 'D'], ['A', 'B', 'D'], ['A', 'C', 'D']]
>>> |

```

### **PRACTICAL NO:-11**

**Aim:-** Write a python program to implement selection sort algorithm and discuss the time complexity Code:- a=list()

```
n=int(input("Enter the number of elements in the list:-"))
```

```
print("Enter numbers in array") for i in range(n):
```

```
num=input()
```

```
    a.append(int(num))
```

```
print(a)
```

```
for i in range(len(a)):
```

```
    min_index = i    for j in
```

```
range(i+1,len(a)):    if
```

```
a[min_index]>a[j]:
```

```

        min_index=j

    a[min_index],a[i]=a[i],a[min_index]

    print("Iteration :",(i+1))    print(a)

    print("Smallest element is :",a[0])
    print("Largest element is :",a[len(a)-1])

```

## #Time Complexity:-

The time complexity for selection sort algorithm is  **$O(n^2)$  in all three cases**

## Output:-

```

>>>
Enter the number of elements in the list:-5
Enter numbers in array
2
3
8
9
7
[2, 3, 8, 9, 7]
Iteration : 1
[2, 3, 8, 9, 7]
Iteration : 2
[2, 3, 8, 9, 7]
Iteration : 3
[2, 3, 7, 9, 8]
Iteration : 4
[2, 3, 7, 8, 9]
Iteration : 5
[2, 3, 7, 8, 9]
Smallest element is : 2
Largest element is : 9
>>> |

```

## PRACTICAL NO:-12

**Aim:-** Using Tournament method find the second largest number in the given list.

Code:- groups=[]

def largest(list1):

    global groups

    if len(list1)==1:

        #print("\*",list1[0])

    return list1[0]    else:

        left=largest(list1[:len(list1)//2])

        #print("left",left)

    right=largest(list1[len(list1)//2:])

    #print("right",right)

    groups.append((left,right))

    #print("max",max (left,right))    return

    max (left,right)

l1=largest([103,10,9,50,60,30])

print("First Largest is:-",l1) s=[]

for item in groups:    if l1

in item:

        s.append(min(item))

print("Second largest is:-",max(s))

## Output:-

```
111 1111111111
>>>
First Largest is:- 103
Second largest is:- 60
\\
```