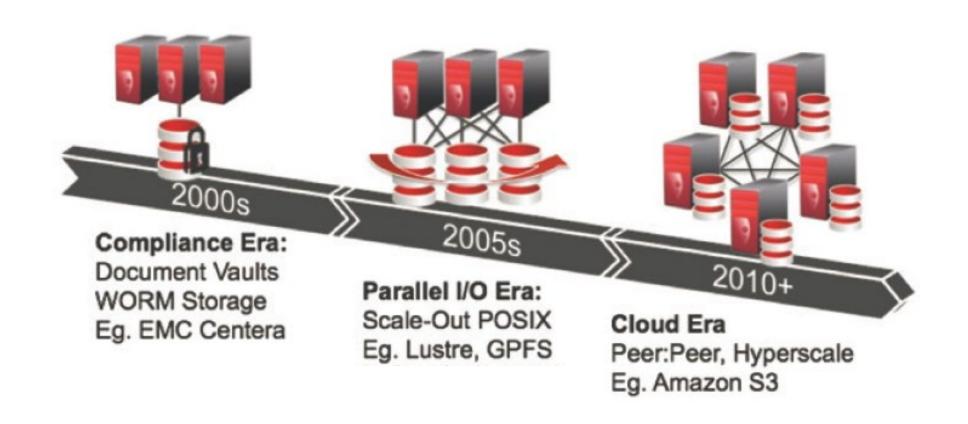
对象存储 Object-Based Storage

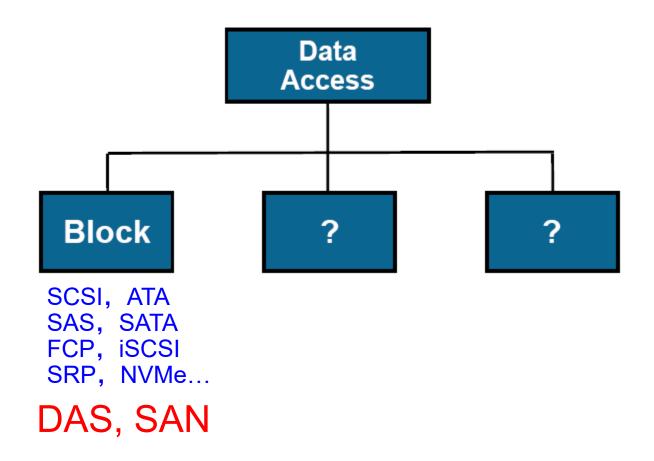


History of Object Storage



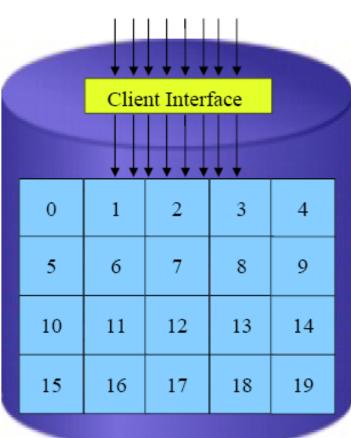
The Data Access Taxonomy

◆ The Block Paradigm



The Block Paradigm

SCSI, ATA, SAS, SATA, FCP, iSCSI, SRP, NVMe...

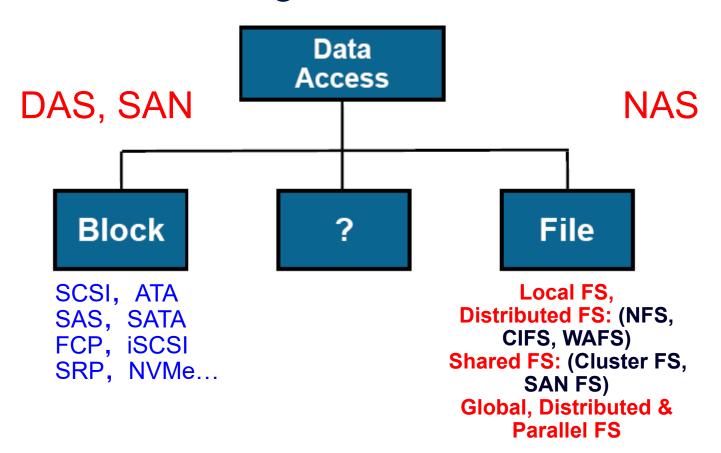


LBA: Logical Block Address

Physical Blocks: e.g. 512 bytes

The Data Access Taxonomy

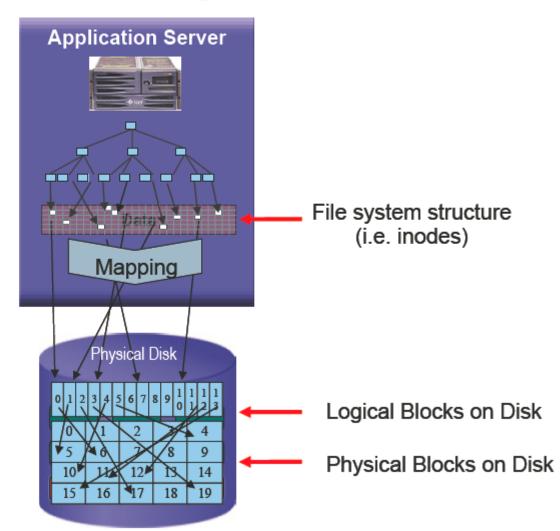
◆ The File Paradigm



Local File Systems

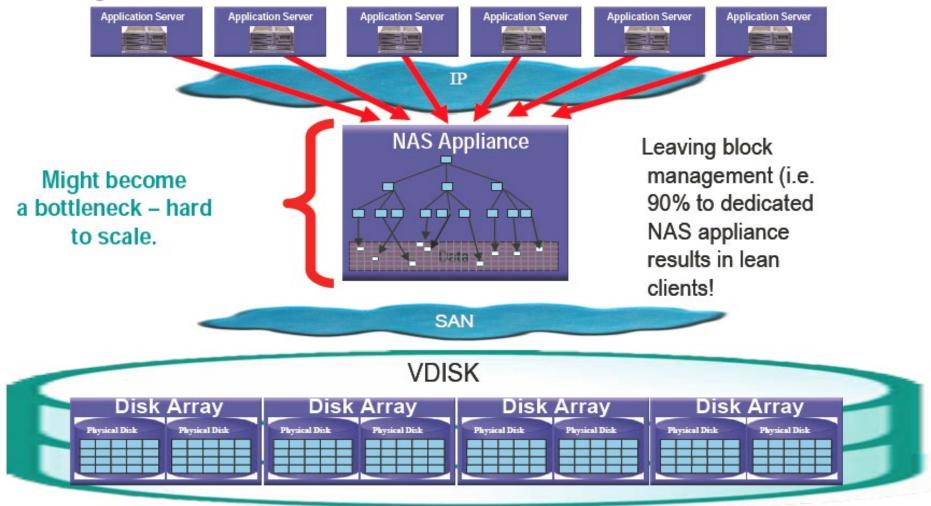
- file/directory management(~10% of workload)
- block/sector management (~90% of workload)

One more level of indirection



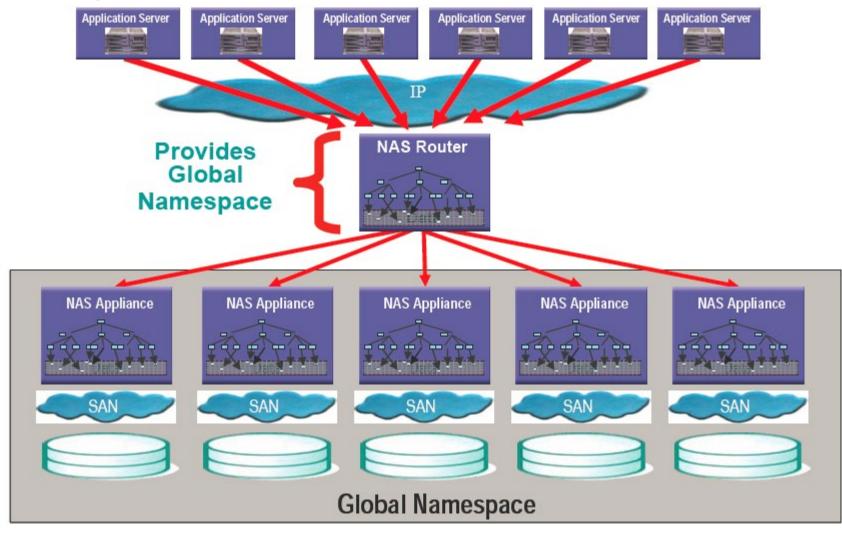
Distributed File Systems

e.g. NAS with NFS,CIFS Protocol



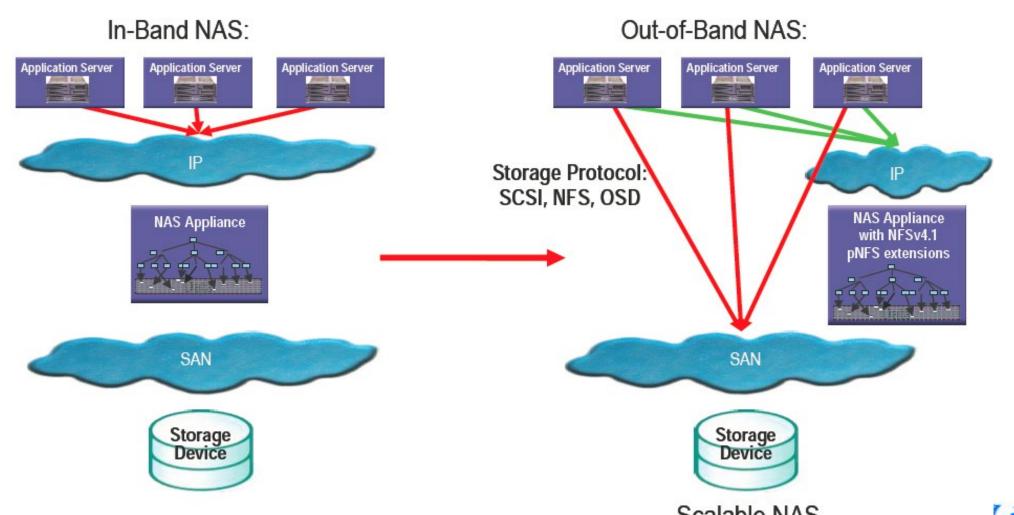
NAS Aggregation/Virtualization

Global Namespace



NAS Cluster

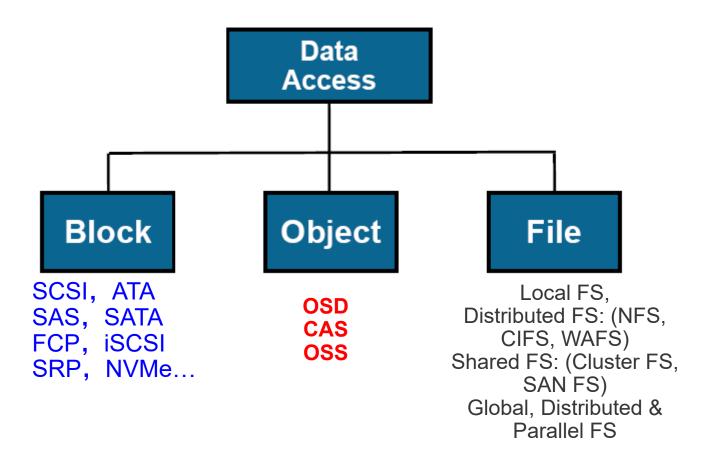
Loosely Coupled NAS: Global Namespace with NFSv4.1 and pNFS



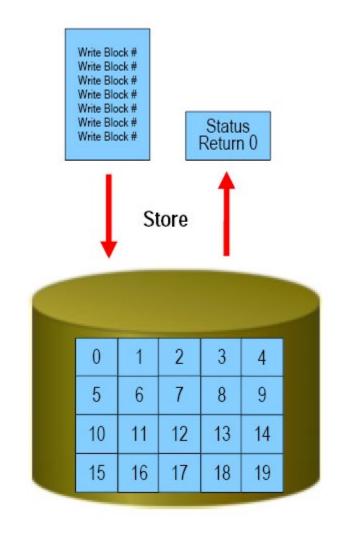
Scalable NAS Loosely Coupled NAS Cluster 存储及

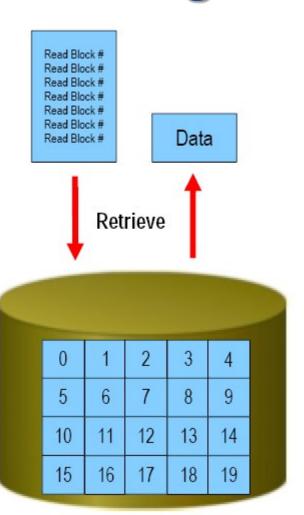
The Data Access Taxonomy

◆ The Object Paradigm

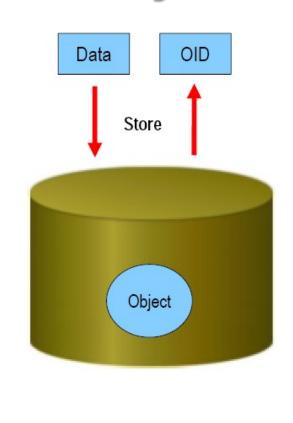


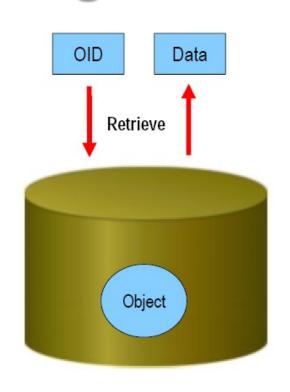
The Old Block Paradigm

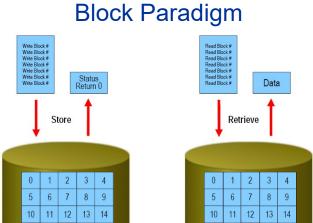




The New Object Paradigm

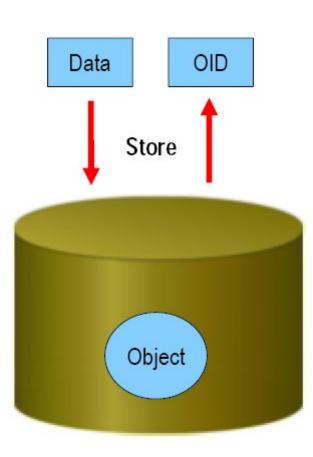






The New Object Paradigm

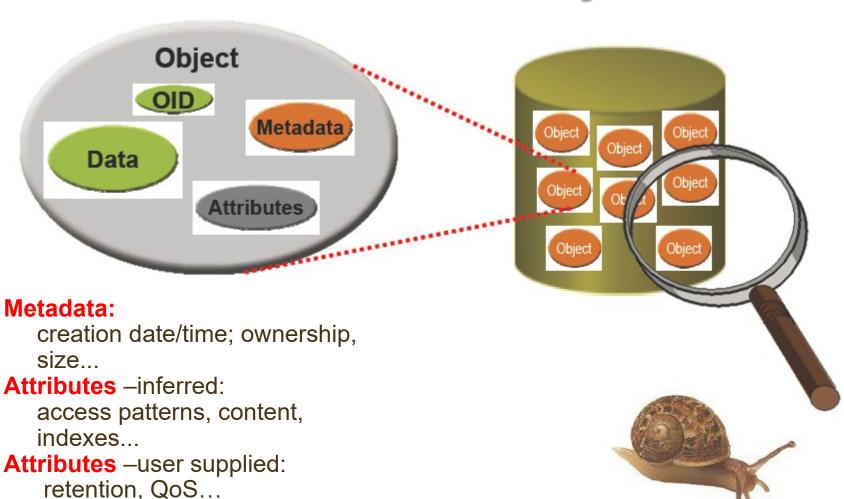
- •WRITE 26,763 Bytes
- •QoS= High
- Description = "X-Ray"
- •Retention = 50 years
- •Access Key = *&^%#
- •Data Payload......



Object Storage Responsibilities:

- Space Management
- Access Control (Identity Mgmt)
- QoS Management
- Cache, Backup
- Policy Migration,
 Retention

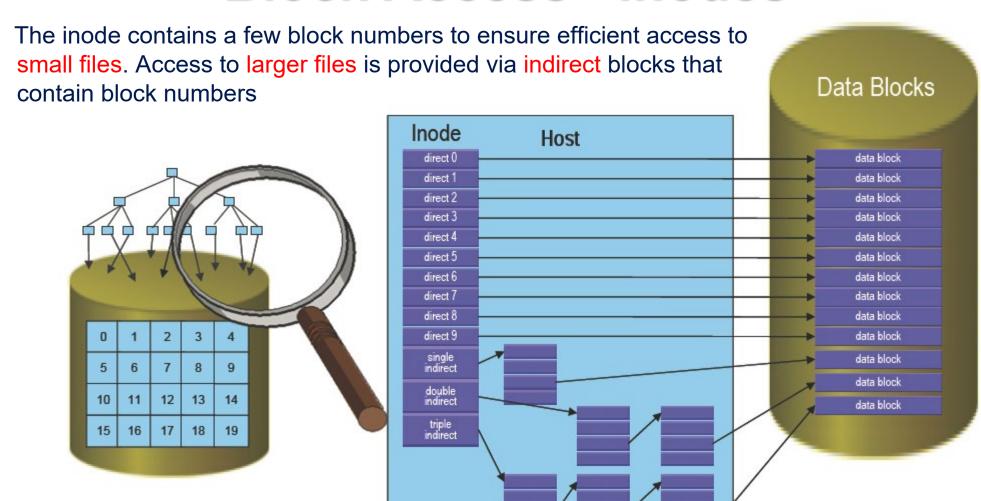
Self-Contained Objects



言息存储及应用实验室

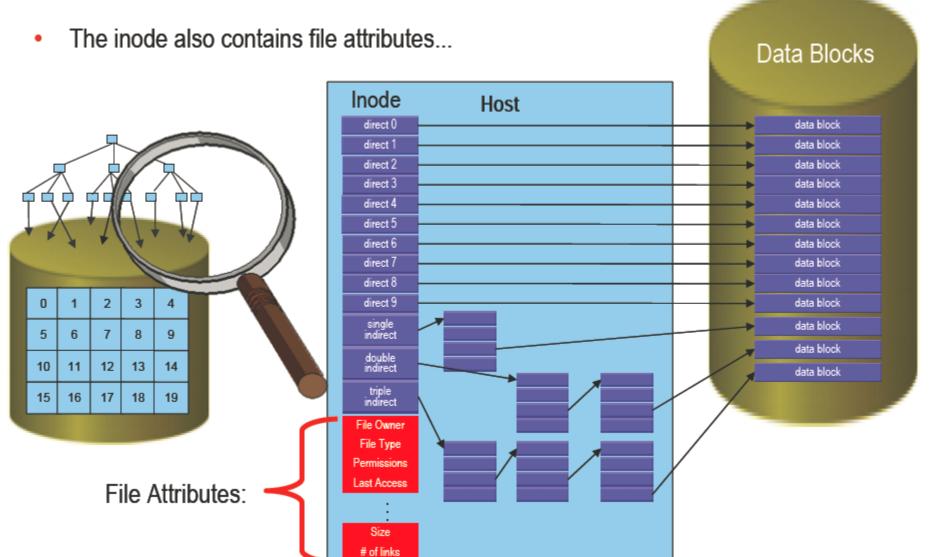
self-contained snail

Block Access - Inodes



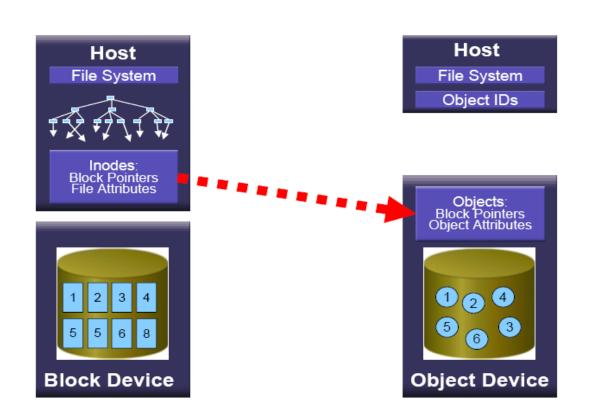
包存储及用实验室

Block Access - Inodes



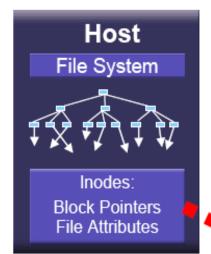
息存储及 用实验室

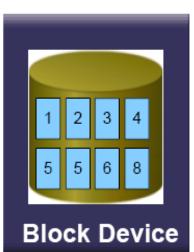
Inodes vs. Objects



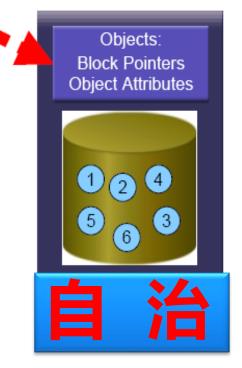
- ➤ Abstract some of the lower layers of storage away from the administrators and applications.
- Inclusion of rich custom metadata within the Object.
 - ✓ specific information(from user or app.)
 for better indexing purposes
 - ✓ Support data-management policies
 - ✓ Centralize management of storage across many individual nodes and clusters
 - ✓ Optimize metadata storage and caching/indexing independently from the data storage

Object Autonomy







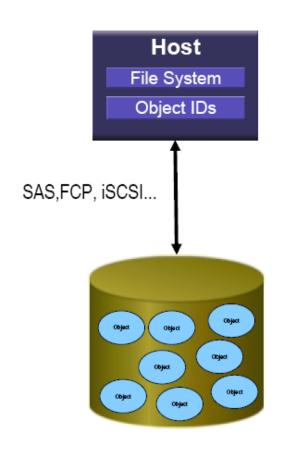


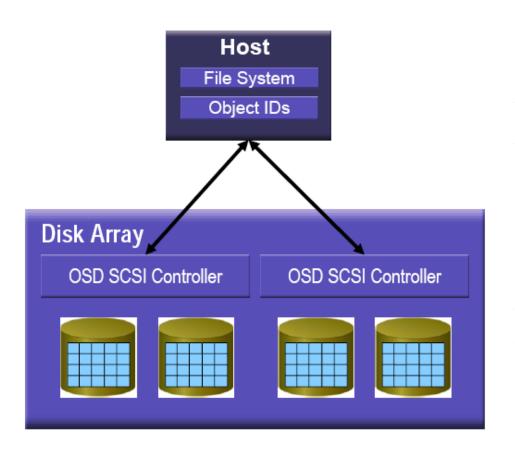
Storage becomes autonomous:

- -capacity planning
- -load balancing
- -backup
- -QoS, SLAs
- –understand data/object grouping
- -aggressive pre-fetching
- -thin provisioning
- -search
- -compression/de-duplication/encryption
- -strong security
- -compliance/retention/secure delete
- -availability/replication
- -audit
 - .

 - .

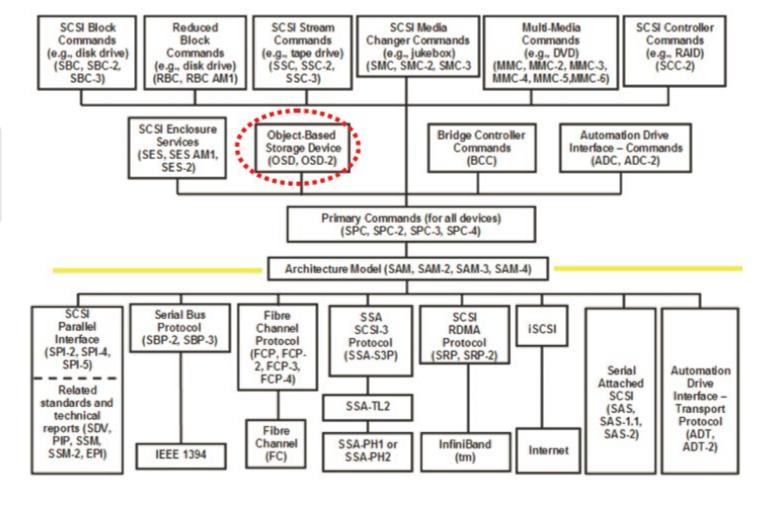
ANSI T10 OSD SCSI Targets





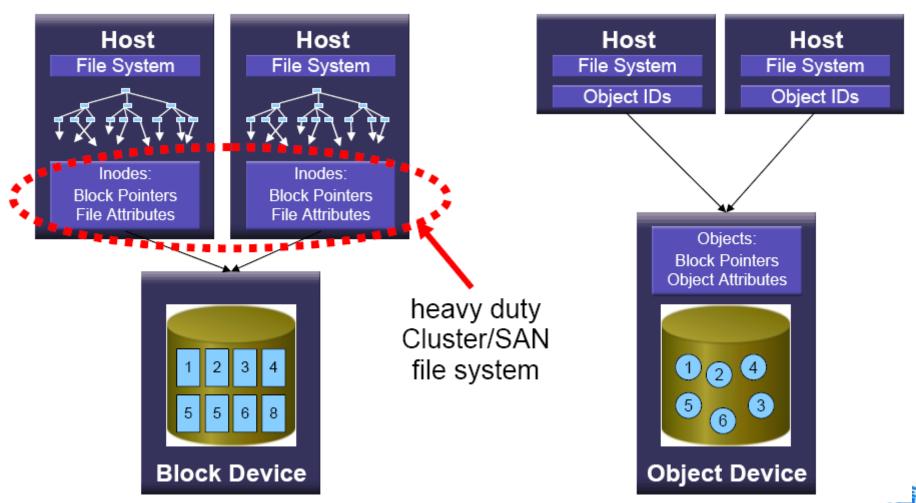
SCSI Standards Architecture

<mark>设备级</mark> 对象存储



Data Sharing

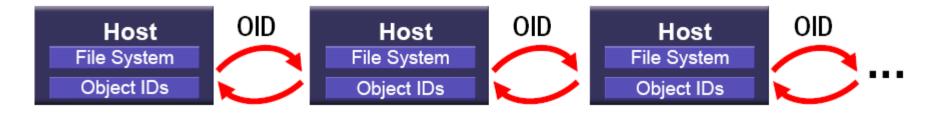
Homogeneous/Heterogeneous

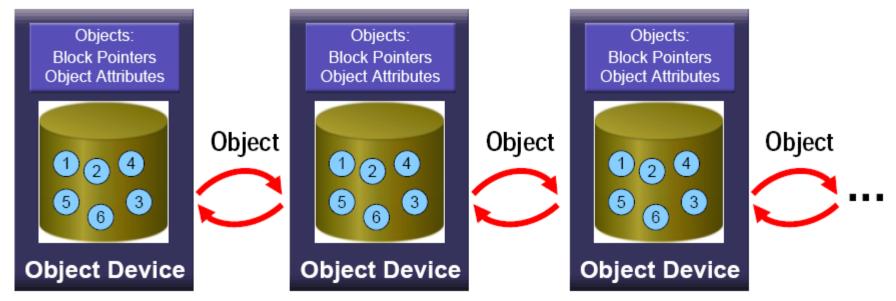


」夏存储及应用实验室

Data Migration - ILM

Homogeneous/Heterogeneous





Additional Layer of Security













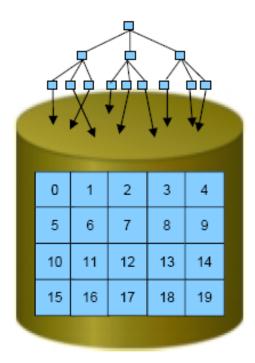
- strong security
 via external service
 - -authentication
 - -authorization

....

- •fine granularity
 - -per object

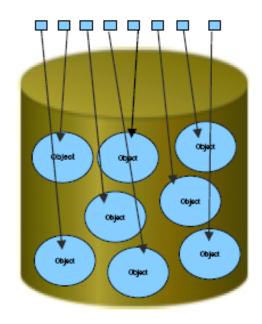
Living in a Flat Namespace

File names / inodes



Traditional Hierarchical

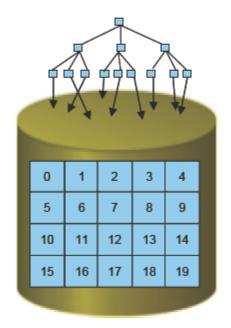
Objects / OIDs



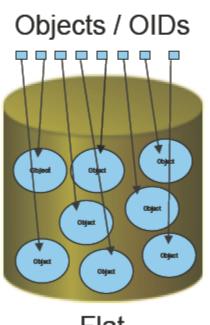
Flat

Virtual View / Virtual File Systems

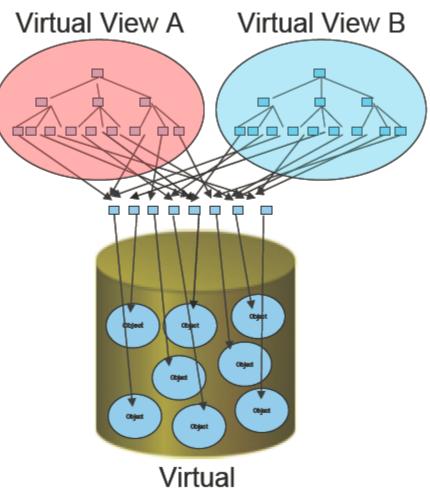
File names / inodes



Traditional

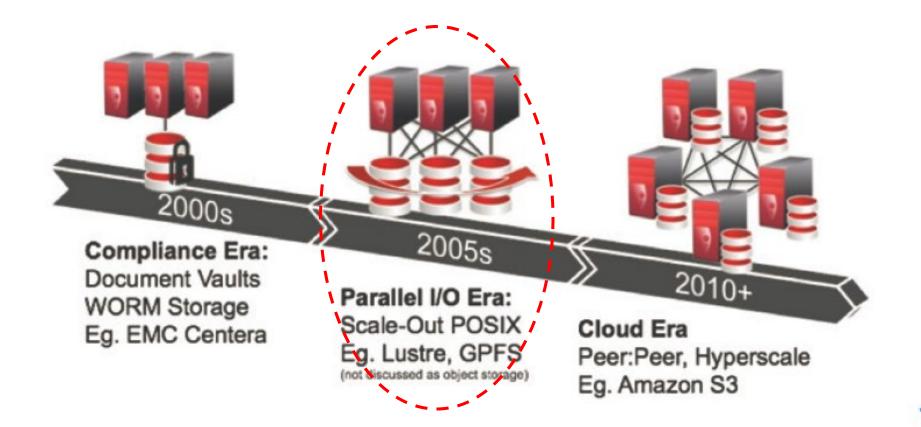


Flat

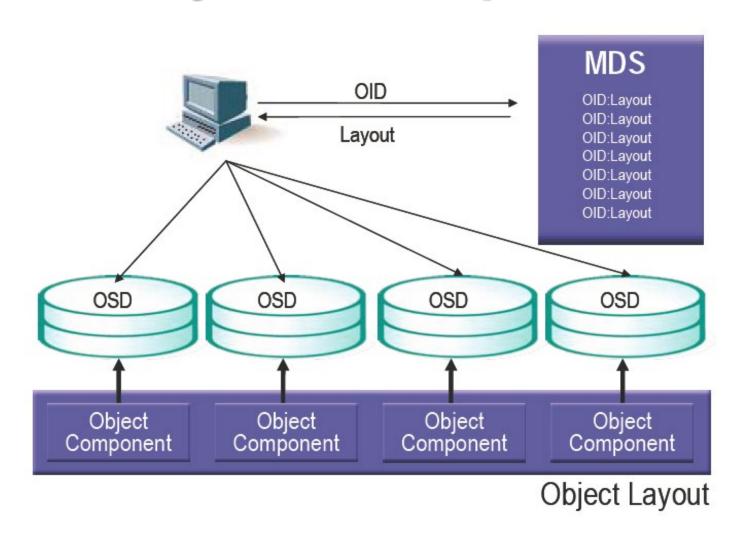


信息存储及 应用实验室

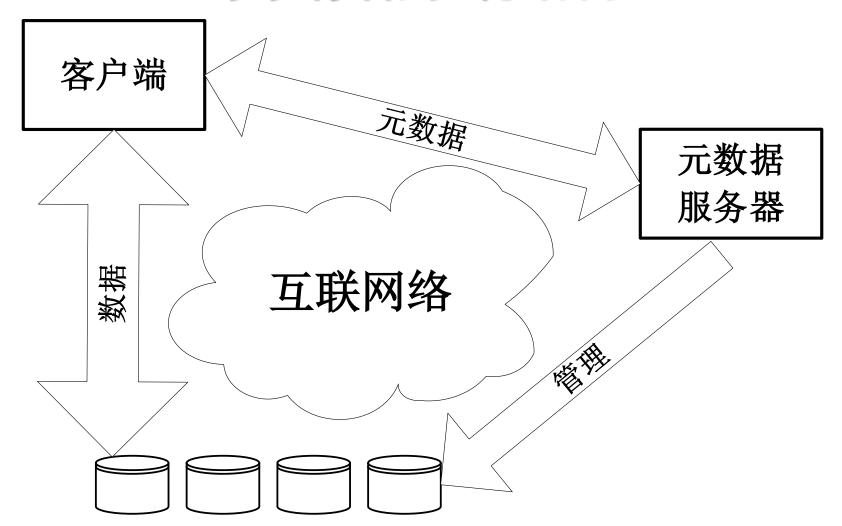
The First Generation of Object Storage (Device-level)



Object Decomposition



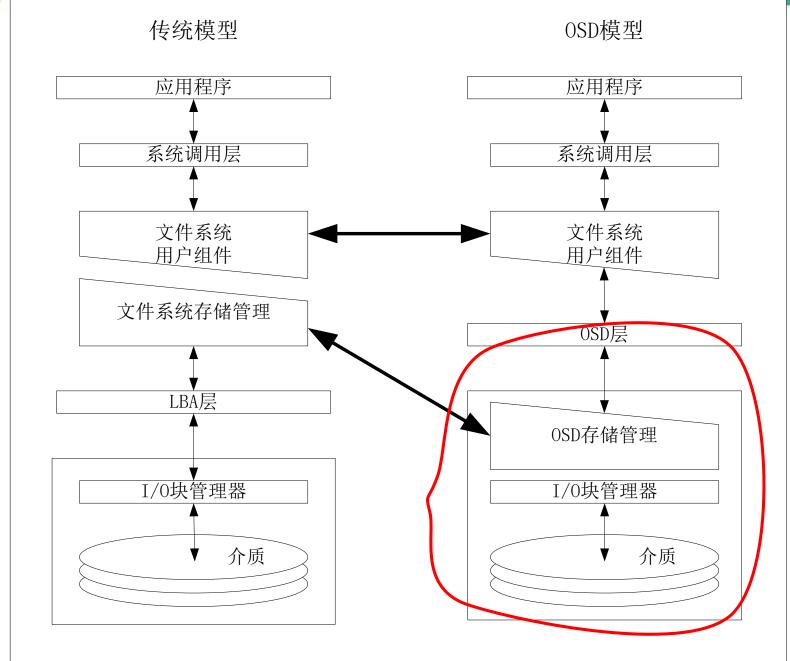
对象存储系统结构



对象存储系统组成

- ◆ 对象(Object)
 - 包含了文件数据以及相关的属性信息,可以进行自我管理
- OSD (Object-based Storage Device)
 - 一个智能存储设备,是Object的集合
- ◆ 文件系统
 - 文件系统运行在客户端上,将应用程序的文件系统请求传输到MDS和OSD上
- ◆ 元数据服务器(Metadata Server, MDS)
 - 系统提供元数据、Cache一致性等服务
- ◆ 网络连接

信息存储理公与社太



对象存储与传统存储的对比

	存储接口	存储系统	优点	缺点
块级存储	块	块存储设备	提供高性能的随机 I/O和数据吞吐率,如: SAN	可扩展性和可管理性较差、价格较高、不能满足成千上万CPU 规模的系统
文件储存	文件	块存储设备 +文件系统	扩展性好、易于管理、 价格便宜,如:NAS	开销高、带宽低、延迟大,不利于高性能集群中应用
对象存储	对象	块存储设备 +文件系统 +定位逻辑 +应用程序	支持高并行性、可伸缩的数据访问,管理性好、安全性高、适合高性能集群使用	处于发展阶段,相应的硬件、软件支持有待进一 步完善

对象存储的特性(总结)

- ◆性能优势
- ◆存储设备的智能化
- ◆数据的共享更容易
- ◆管理更方便
- ◆更好的安全性

是大数据存储理想的选择

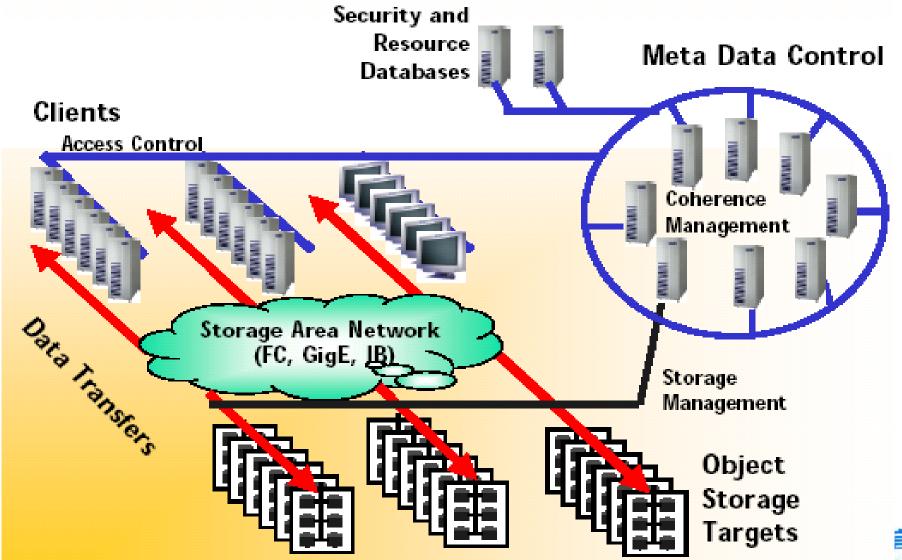
对象存储系统实例: Lustre

- A shared file system for HPC clusters
 - Open Source software (GPL)
 - linux cluster
- Very high metadata and I/O performance
 - 5,000 file creations/sec in 1 dir, 1,000 nodes
 - Single clients up to 290MB/sec.
 - Aggregate up to 11GB/sec
- Scalable to 1,000's of nodes
- In production now on such clusters

Lustre Retrospective

- 1999 Initial ideas @CMU
- Seagate: management aspects, prototypes
 - Much survives today
- 2000 National Labs
 - Can Lustre be next generation FS?
 100 GB/sec, trillion files, 10,000's clients, secure, PBs
- ◆ 2002 2003
 - Many partners: Dell, HP, Cray, LNXI, DDN others
 - Production use, 1.0 released
- ◆ April 3rd, 2018, Lustre 2.11.0 released

A Lustre Cluster

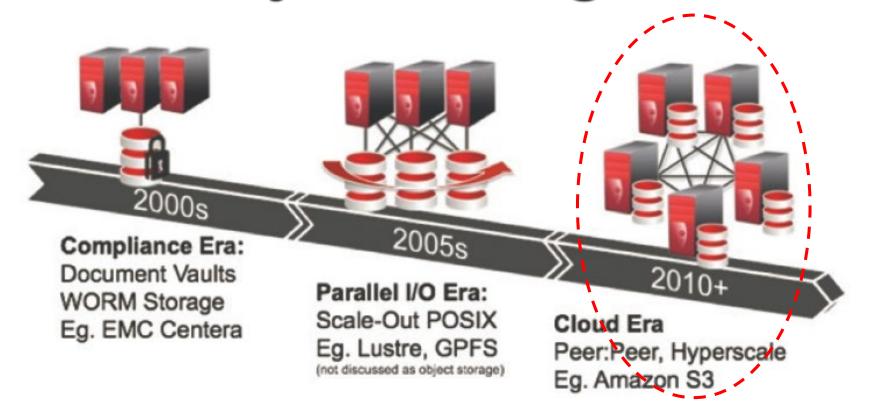


言息存储及 应用实验室

Reference

- ◆ Lustre: A SAN File System for Linux
 - http://lustre.org/documentation/
- ◆ Several presentation materials from Dr. Peter J. Braam

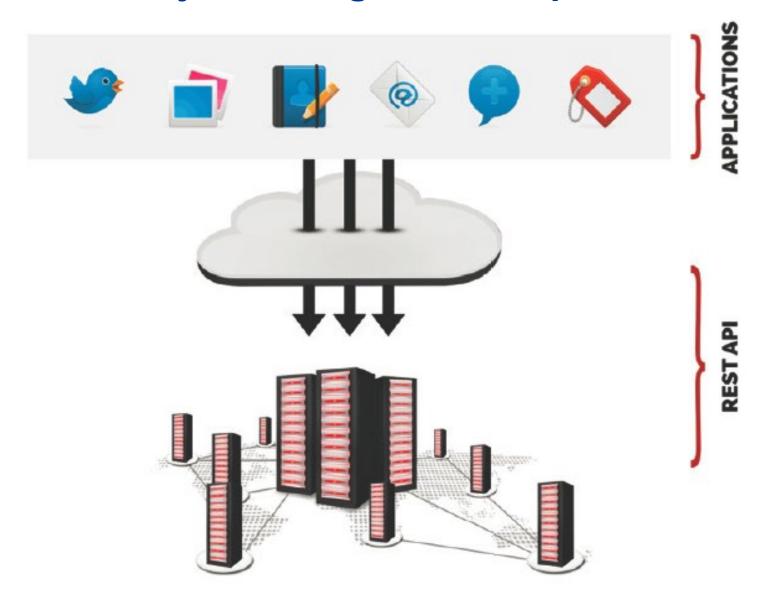
The Current Generation of Object Storage



Cloud Storage, Storage Clouds, Object Storage

- ◆ Cloud Storage is the storage used for Compute Cloud infrastructures
 - Compute Clouds are very IOPS intensive and usually block storage is used in these applications
- ◆ Storage Clouds are "storage in the cloud", whether public or private
 - Storage Clouds are simply storage capacity that is made available through the Internet
 - Most of today's storage clouds use object storage technologies

Scale out object storage with simple REST API



REST API's

- ◆ REST stands for Representational State Transfer
- It is a software architecture that is used for distributed application environments
- ◆ REST API's have become the predominant interface for cloud applications to connect to the cloud
- ◆ For storage-centric cloud applications, a REST API is the interface between the application and the object storage platform
- ◆ PUT GET DELETE...

Current Object Storage Summary

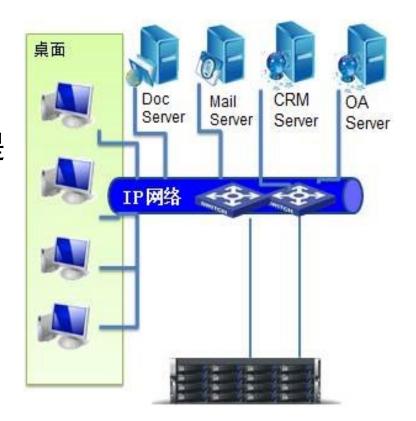
- ◆ Data is stored as objects in one large, scalable pool of storage
- ◆ Objects are stored with metadata information about the object
- ◆ An Object ID is stored, to locate the data
- ◆ REST is the standard interface, simple commands used by applications
- ◆ Objects are immutable; edits are saved as a new object

- ◆ 块存储 (DAS/SAN)
 - ➤FC、iSCSI协议
 - > 专有的系统中, 高读写性能和高可靠性
 - ▶一套存储只服务一个应用系统,例如如交易系统,计费系统。典型行业如金融,制造,能源,电信等



□ 文件存储(NAS)

- ✓ NFS/CIFS协议, IP网络
- ✓兼顾多个应用和更多用户访问,同时提供方便的数据共享手段
- ✓中小企业市场, CRM系统, SCM系统, OA系统等



- ◆ 对象存储(Object)
 - ➤ OSD, HTTP协议
 - ▶互联网或者公网
 - >海量数据,高并发访问
 - ▶ 常见的适配应用如网盘、媒体娱乐, 医疗PACS, 气象, 归档等数据量超大而又相对"冷数据"和非在线处理的应用类型





Throughput easy, latency hard



Throughput is easy



Latency is hard

Throughput is an engineering problem, latency is a physics problem!

对象存储与键值存储

◆ 相同点

- ▶对象ID 类似于Key(任意字符串)
- ▶数据可以具有任意大小

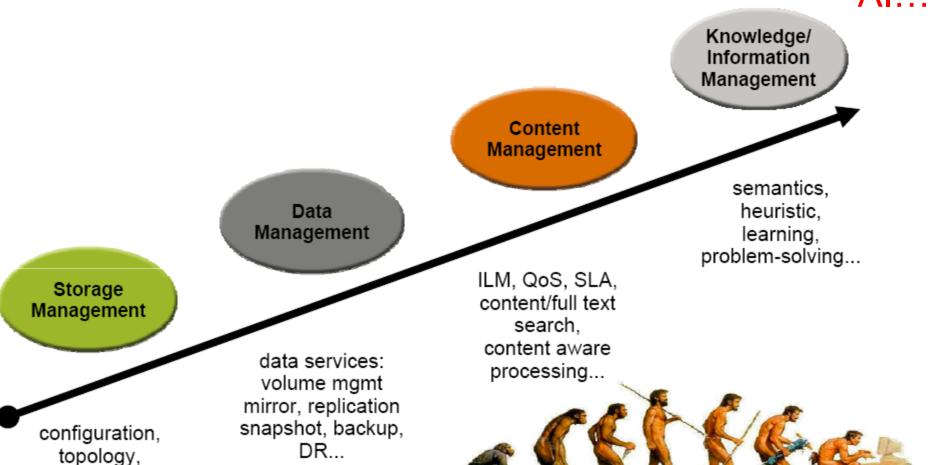
◆ 不同点

- ▶ 对象存储还允许将一组有限的属性(元数据)与每个(数据)对象相关联
- ▶ 对象存储针对大量数据(数百兆甚至千兆字节)进行了优化, 而对于键值存储,其值则相对较小(千字节)
- ▶ 对象存储通常提供较弱的一致性保证,例如最终的一致性,而 键值存储则提供较强的一致性。

RAID...

The Evolution of Data Processing

AI...



信息存储理论与技术

Moving Data to the Processor is Costly

One floating-point calculation



Moving data from DRAM to CPU



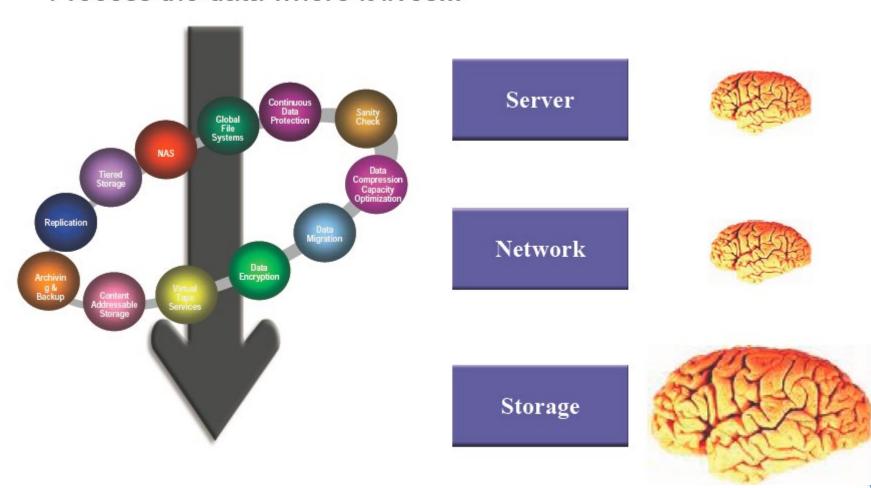
17 picojoules

17,000 picojoules

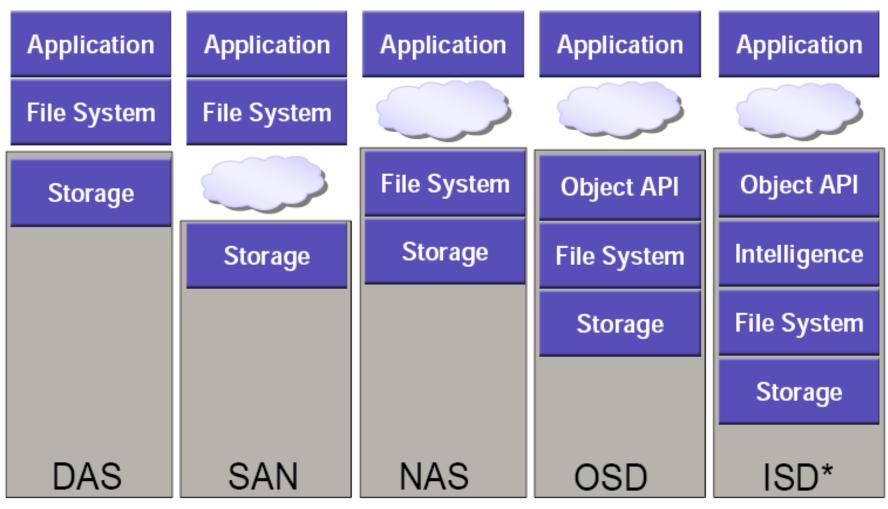
Opportunities for 1000x improvement are increasingly rare

Migration of Storage Application

Process the data where it lives...

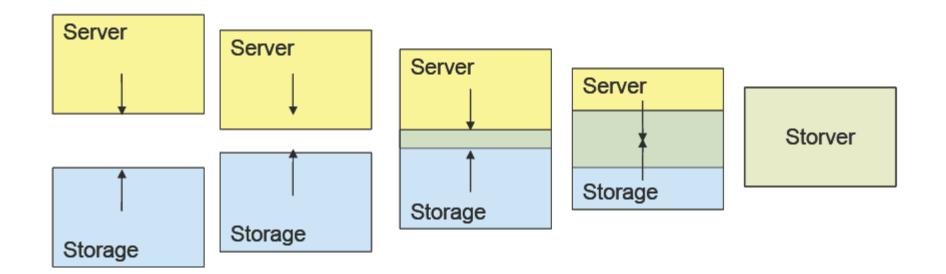


The Evolution of Storage



The Vertical Consolidation

- Storage and server
- Migration of data processing applications
- ♦ No I/O is best I/O



主动对象存储

- ◆传统存储系统被动响应服务请求
- ◆对象具有智能性

智能的系统能够提供主动服务

存储组织结构对比

传统模式

主动对象模式

应用

京统接口 文件系统用户部件 文件系统存储部件 硬盘驱动 总线驱动 适配器驱动

效率低 管理复杂

系统接口 文件系统用户部件 OS 妾口 文件系》存储部件 块I/O管理器

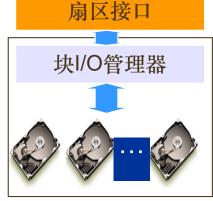
性能高

管理简单

- 含"数据、属性" 和"**操作**"
- 自组织和自管理
- 主动服务

存储系统

主机



- 只含数据
- 主机组织 和管理
- 被动响应

主动对象存储服务机制

