# **Network: Internet Protocol (2)**

Jae Hyeon Kim

### Reference

William Stalling, Data and Computer Communications 10/E, Prentice Hall

#### Subnets and Subnet Masks

- Aims to save the IP addresses by sharing an address with a set of physical networks
- Allows arbitrary complexity within an organization
  - Insulate overall Internet from growth of network numbers and routing complexity
  - Site looks to rest of internet like single network
- Each LAN assigned subnet number
  - Host portion of address partitioned into subnet # and host #
  - Local routers route within subnetted network
- Subnet mask indicates which bits are subnet number and which are host number

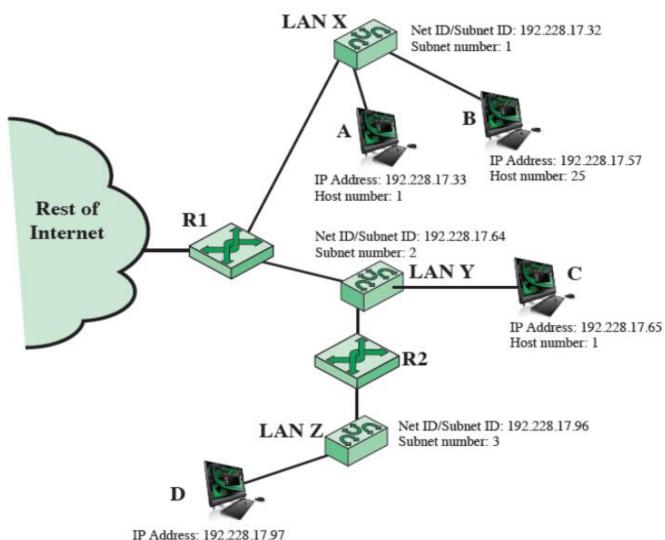
### IP Addresses and Subnet Masks

*		
φ	Binary Representation	Dotted Decimal∘
■IP address。	11000000.11100100.00010001.00111001	192.228.17.57
Subnet mask-	11111111111111111111111111111111000000	255.255.255.224
Bitwise AND of address and mask (resultant network/subnet number)	11000000.11100100.00010001.001000000	192.228.17.32
Subnet number.	11000000.11100100.00010001.001	1₽
Host number	00000000.000000000.0000000000011001	250

#### (b) Default subnet masks

φ	Binary Representation	Dotted Decimal∘
Class A default mask	11111111.00000000.000000000.00000000000	255.0.0.00
Example Class A mask	111111111111000000.00000000000000000000	255.192.0.00
Class B default mask	11111111111111111111000000000.000000000	255.255.0.0
Example Class B mask	1111111111111111111111000.0000000000	255.255.248.0
Class C default mask	1111111111111111111111111110000000000	255. 255. 255.00
Example Class C mask	1111111111111111111111111111111000	255.255.255.252

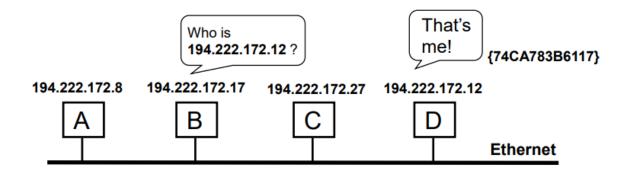
### An Example : Routing Using Subnets



Host number: 1

### Mapping IP Addresses to the DL

- Consider an Ethernet(IEEE802.3) LAN running IP
  - Recall data link layer has it's own 48-bit addresses
  - Network layer provide it's own 32-bit IP address space
  - Data link layer knows nothing about IP addresses
- How do these two sets of addresses get mapped?
  - ARP(Address Resolution Protocol, RFC 826) build a query message and broadcast it
  - All hosts in the network receive it and one says "that's me!"



### Address Resolution Protocol (ARP)

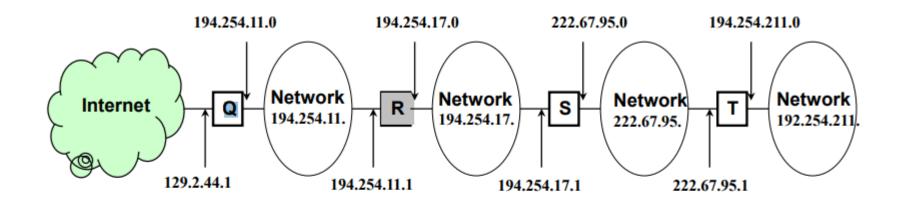
- ARP is a part of the physical network system
  - But it is not a part of the IP
- ARP is a low-level protocol that hides the underlying network physical addressing
  - And permits a machine to assign an arbitrary IP address
- Now, the broadcasting is too expensive, how can it be solved?
  - When a host receives an ARP reply, it saves the sender's IP address and its physical address in its cache for successive lookups
- Is it be possible more refinement?
  - The sender's IP-to-physical address binding is included in every ARP broadcast; receivers update the binding in their cache

### Table Driven IP Routing

- IP routing employs an routing table on host and router
  - The routing table contains information about the possible destinations and how to reach them
  - IP consults the table to decide where to send the datagram
- Then, what information should be kept in routing tables?
  - Minimal information principle : keep network prefix only
  - Information hiding principle: specifies one step along the path from the router to a destination
  - Default routing : if no route appears in the table, the routing routines send the datagram to a default router

## Table Driven IP Routing (An Example)

To reach hosts



Route to

on network	this address	
194.254.11	Deliver Directly	
194.254.17	Deliver Directly	
222.67.95	194.254.17.1	
192.254.211	194.254.17.1	
Default	194.254.11.0	

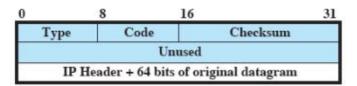
### Routing Protocols in IPv4

- IP routing is based on the destination network ID alone?
  - All IP traffic for a given network tales the same path regardless to the delay or throughput of physical network
  - Only the final router can determine if the destination exists or is operational, the router only can report the delivery to the sender
  - Each router routes traffic independently someone should find out if two-way communication is always possible
- IP routing selects the next hop to be sent the datagram
  - IP simply passes the datagram and the next hop IP address to the network interface software (so-called network driver)
  - The driver software responsible to bind the next hop IP address to a physical address, forms a frame, and sends it

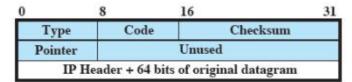
### Internet Control Message Protocol (ICMP)

- RFC 792
- Provides a means for transferring message from routers and other hosts to a host
- Mainly provides feedback about problems
  - Datagram cannot reach its destination
  - Router does not have buffer capacity to forward
  - Router can send traffic on a shorter route
- Encapsulated in IP datagram
  - Hence net reliable

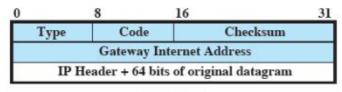
## ICMP Message Format



(a) Destination Unreachable; Time Exceeded; Source Quench



(b) Parameter Problem



(c) Redirect

0		8	16	31
	Type	Code	Checksum	
	Identifier		Sequence Number	9
		Optio	nal data	

(d) Echo, Echo Reply

8	16	31
Code	Checksum	
ntifier	Sequence Number	
Originate	Timestamp	
	ntifier	Code Checksum

(e) Timestamp

8	16	31		
Cod	le Chec	ksum		
Identifier	Sequence	Number		
Originate Timestamp				
Receive Timestamp				
Tran	smit Timestamp			
	Identifier Orig Rec	Code Chec Identifier Sequence Originate Timestamp		

(f) Timestamp Reply



(g) Address Mask Request

0		8	16 3	31
	Type	Code	Checksum	
	Identifier		Sequence Number	
		Addres	ss Mask	

(h) Address Mask Reply

## Why Change IP?

- Address space exhaustion
  - Two level addressing (network and host) wastes space
  - Network addresses used even if not connected to Internet
  - Growth of networks and the Internet
  - Single address per host
- Requirement for new types of service
  - Address configuration
  - Routing flexibility
  - Traffic support
- Security
- mobility

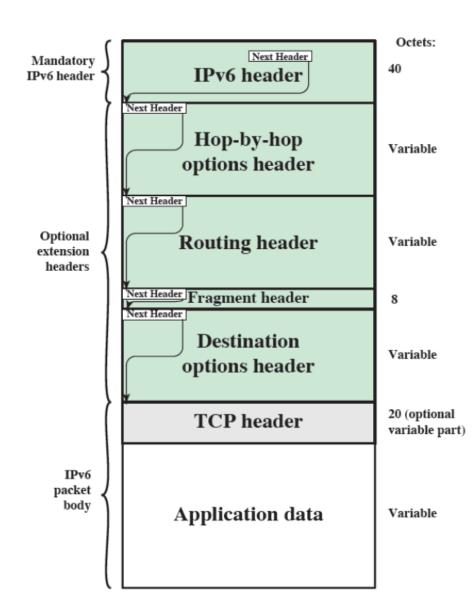
#### IPv6 RFCs

- IP versions
  - IP v1-3 defined and replaced
  - IP v4 current version
  - IP v5 streams protocol
  - IP v6 replacement for IP v4 : during development it was called IPng(IP Next Generation)
- RFC 1752 recommendations for the IPng
  - Requirements
  - Addressing, routing security issues
- RFC 2460 overall specification
- RFC 4291 addressing structure

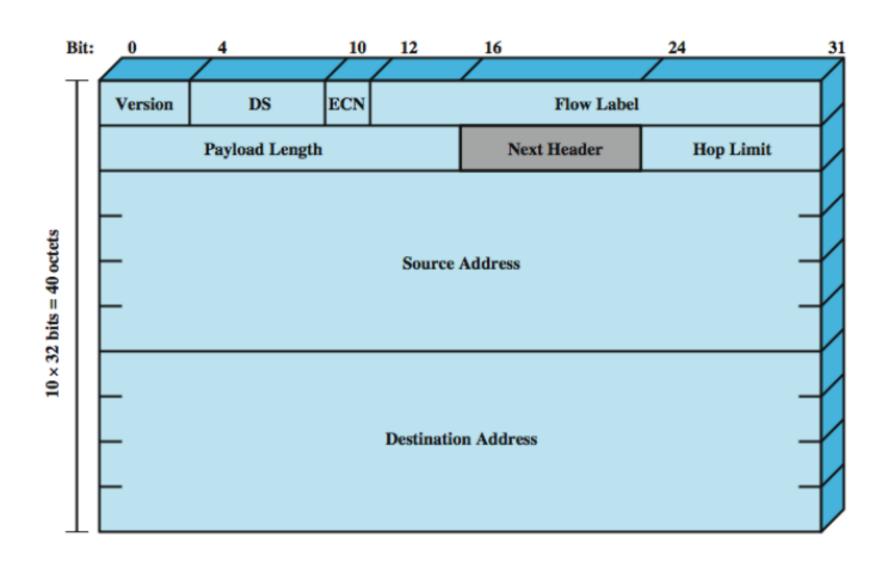
#### IPv6 Enhancements

- Expanded address space
  - 128 bit
- Improved option mechanism
  - Separate optional headers between IPv6 header and transport layer header (most are not examined by intermediate routes)
- Dynamic address assignment
- Increased addressing flexibility
  - Anycast delivered to one of a set of nodes
- Support for resource allocation
  - Labeling of packets to particular traffic flow

#### - IPv6 Packet with Extension Headers



### **IPv6** Header

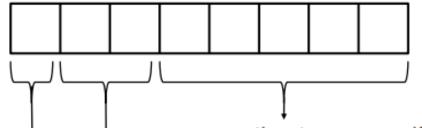


#### IPv6 Flow Label

- Related sequence of packets
  - Needs a special handling for a specific packet
  - Identified by source and destination address + flow label
- Router treats flow as sharing attributes
  - E.g. path, resource allocation, discard requirements, security
- May treat flows differently
  - Buffer sizes, different forwarding precedence, different QoS
- Flow requirements are defined prior to flow start and a unique flow level is assigned to the flow
  - Alternative to including all info in every header

#### Next Header

#### It consists as:



option type: specify a particular option (including TCP/UDP)

indicate the action to be taken by a node that does't recognize this option type

- 00 : skip over this option and continue processing the header
- 01 : discard the packet
- 10 : discard the packet and send an ICMP parameter problem message to the packet source
- 11: discard the packet and send an ICMP parameter problem message to the packet source, for not a multicast address

specifies whether the Option Data field does't change(0), may change (1) on route from source to destination

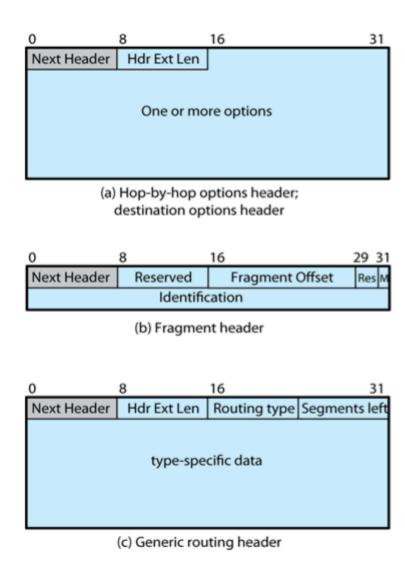
#### IPv6 Addresses

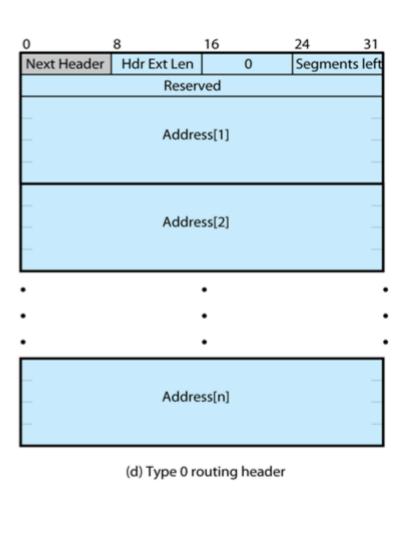
- 128 bits long
- Assigned to an interface rather than a host
  - Single interface may have multiple unicast addresses
- Three types of address
  - Unicast : single interface
  - Anycast: set of interfaces (typically different nodes), delivered to any one interface, usually the "nearest"
  - Multicast: set of interfaces, delivered to all interfaces identified

#### Extension Headers

- Hop-by-hop options
  - Require processing at each router
- Routing options
  - Similar to v4 source routing
- Fragmentation options
  - Only allowed at source, no fragmentation at intermediate routers
- Authentication options
  - Encapsulating security payload
- Destination options
  - Carries optional information for destination node

#### **IPv6 Extension Headers**





### Virtual Private Network (VPN)

- Set of computers interconnected using an unsecure network, as Internet
  - But they provide a secure channel between the organizations, much like as a private network
- Using encryption & special protocols to provide security
  - Eavesdropping (cryptography)
  - Entry point for unauthorized users (authentication)
- Proprietary solutions are problematical
  - Hence, develop the IPSec standard

#### IPSEC

- RFC 1636 (1994) identified security need
- Encryption & authentication to be IPv6
  - But designed also for use with current IPv4
- Applications needing security include:
  - Branch office connectivity
  - Electronic commerce security
- Benefits
  - Provides strong security for external traffic
  - Resistant to bypass
  - Can be transparent to applications as well as end users

#### IPSEC Function

- Authentication header
  - For authentication only
- Encapsulating Security Payload (ESP)
  - For combined authentication/encryption
- A key exchange function
  - Manual or automated
- VPNs usually need combined functions

#### IPSEC Scenario: VPN

