

TOA Assignment 3.

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Am1) Union:

	0	1
$\{ \begin{matrix} z_1 \\ z_2 \\ z_3 \end{matrix} \}$	$\begin{matrix} z_2 & 11 \\ z_1 & 11 \\ z_5 & 1 \end{matrix}$	$\begin{matrix} z_3 & 11 \\ z_4 & 1 \\ z_1 & 11 \end{matrix}$
$\{ \begin{matrix} z_4 \\ z_5 \\ z_6 \end{matrix} \}$	$\begin{matrix} z_6 & 11 \\ z_3 & 11 \\ z_4 & 1 \end{matrix}$	$\begin{matrix} z_2 & 11 \\ z_5 & 1 \\ z_6 & 11 \end{matrix}$

	0	1
$\begin{matrix} z_1 & 3 & 111 \\ z_2 & 3 & 111 \\ z_3 & 3 & 111 \\ z_6 & 3 & 111 \end{matrix}$	$\begin{matrix} z_2 & 111 \\ z_1 & 111 \\ z_5 & 1 \\ z_4 & 111 \end{matrix}$	$\begin{matrix} z_3 & 111 \\ z_4 & 111 \\ z_1 & 111 \\ z_6 & 111 \end{matrix}$
$\begin{matrix} z_4 & 3 & 111 \\ z_5 & 3 & 111 \end{matrix}$	$\begin{matrix} z_6 & 111 \\ z_3 & 111 \end{matrix}$	$\begin{matrix} z_2 & 111 \\ z_5 & 1 \end{matrix}$

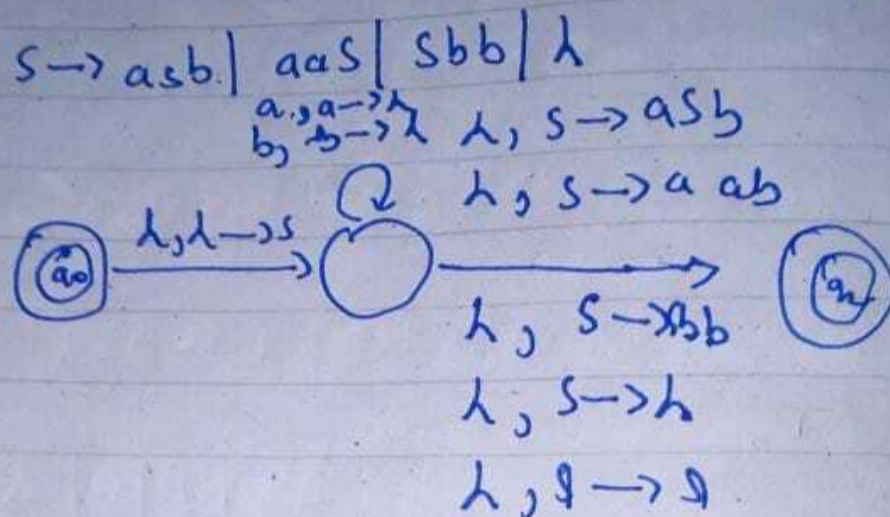
$\begin{matrix} z_1 & 111 \\ z_2 & 111 \\ z_3 & 111 \\ z_6 & 111 \\ z_4 & 111 \\ z_5 & 111 \end{matrix}$	$\begin{matrix} z_2 & 111 \\ z_1 & 111 \\ z_5 & 111 \\ z_4 & 111 \\ z_6 & 111 \\ z_3 & 111 \end{matrix}$	$\begin{matrix} z_3 & 111 \\ z_4 & 111 \\ z_1 & 111 \\ z_6 & 111 \\ z_2 & 111 \\ z_5 & 111 \end{matrix}$
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0 states were merged into one DFA nonreduced form.

① Clavine

$$\begin{array}{c|c|c}
 a_i \text{ III} & s_2, s_1 \text{ II} & s_1 \text{ I} \\
 s_1 \text{ III} & s_2, s_1 \text{ II} & s_1 \text{ II} \\
 s_2, s_1 \text{ III} & s_1 \text{ II} & s_2, s_1 \text{ II}
 \end{array}$$

Ans 2) (a) $\{a^n b^m \mid (n+m) \text{ is even}\}$.

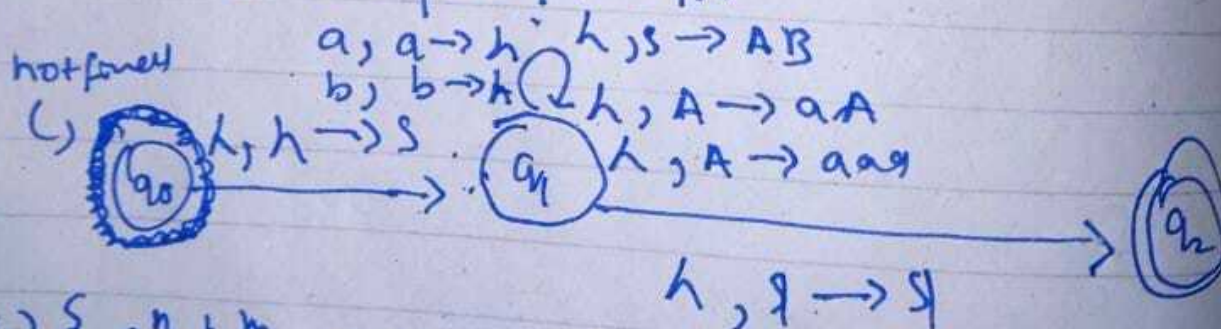


(b) $\{a^n b^m \mid n \geq 4, m \leq 3\}$.

$S \rightarrow AB$

$A \rightarrow aA \mid aaaa$

$B \rightarrow b \mid bb \mid bbb \mid \lambda$



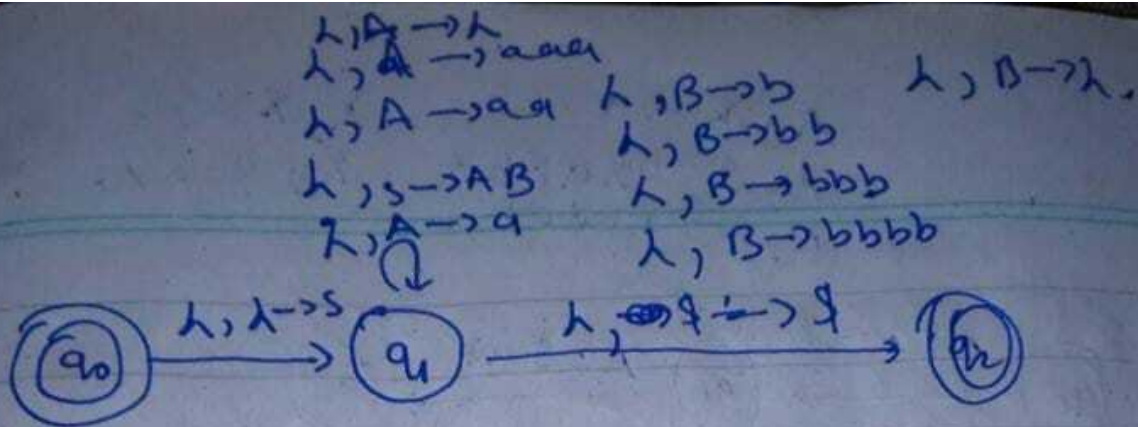
(c) $\{a^n b^m \mid n \leq 4, m \leq 4\}$.

$S \rightarrow AB$

$A \rightarrow a \mid aa \mid aaa \mid \lambda$

$B \rightarrow b \mid bb \mid bbb \mid bbbb \mid \lambda$

aaaa



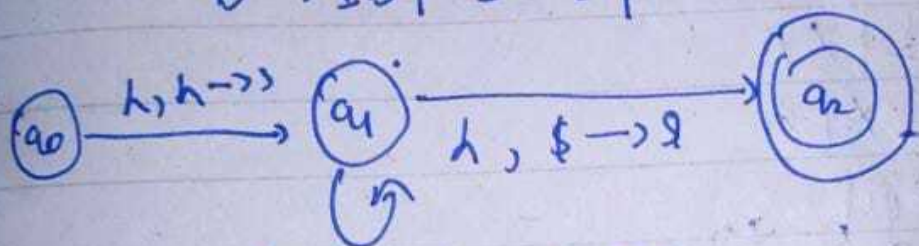
(d) having exactly one pair at consecutive zeros.

$$(1 + 01)^* 00 (1 + 10)^*$$

$$S \rightarrow A 0 0 B$$

$$A \rightarrow 1A \mid 01A \mid \lambda$$

$$B \rightarrow 1B \mid 10B \mid \lambda$$



$$\begin{aligned} \lambda, S &\rightarrow A 0 0 B & \lambda, B &\rightarrow 1 \\ \lambda, A &\rightarrow 1A & \lambda, B &\rightarrow 10 \\ \lambda, A &\rightarrow 01A & \lambda, B &\rightarrow \lambda \\ \lambda, A &\rightarrow \lambda \end{aligned}$$

(c) all strings that contain at least one occurrence of each symbol.

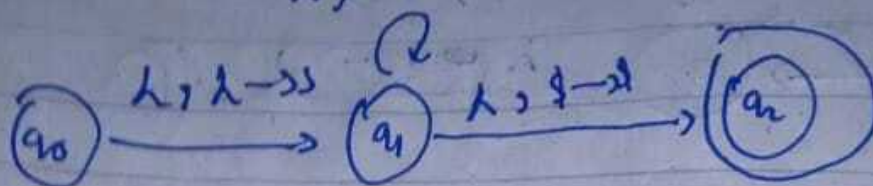
$$RE: (a+b)^* a (a+b)^* b (a+b)^* + (a+b)^* b (a+b)^* a$$

$$(a+b)^* + (a+b)^*$$

$$S \rightarrow x a x b x \mid x b x a x$$

$$x \rightarrow a x \mid b x \mid \lambda$$

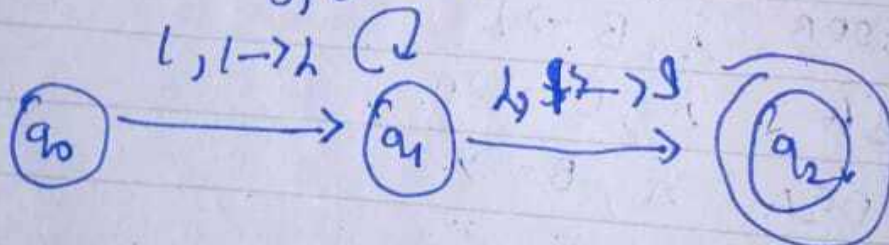
$\lambda, x \rightarrow a$
 $\lambda, y \rightarrow b$
 $\lambda, s \rightarrow xaxbx$
 $\lambda, s \rightarrow xbxax$



(f) All strings not ending in 01

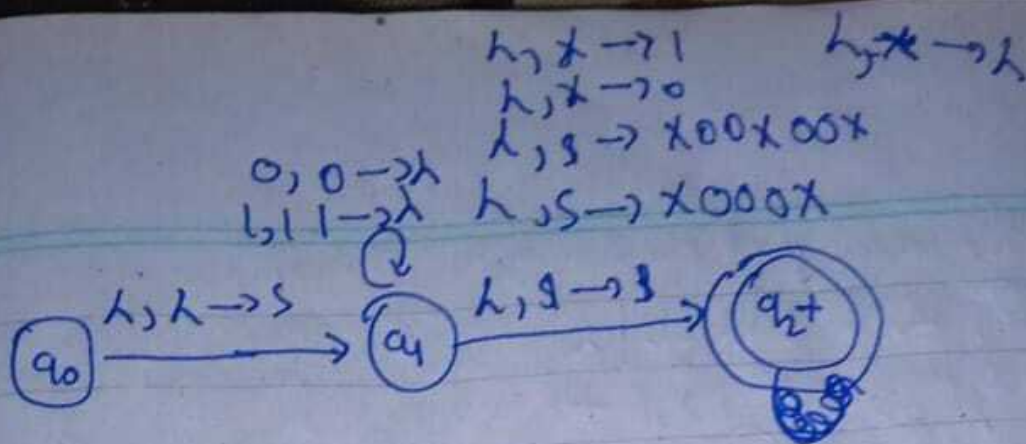
R.E $\lambda + 0 + 1 + (0+1)^* (00+01+11)$

$S \rightarrow \lambda | 0 | 1 | xA$
 $x \rightarrow 0x | 1x | \lambda$
 $A \rightarrow 00 | 01 | 11$
 $0, 0 \rightarrow \lambda$

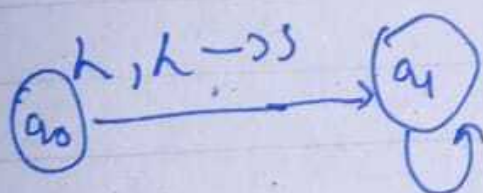


(g) all strings having at least two occurrences of substring 00

$R.E (1+0)^* 00 (1+0)^* 00 (1+0)^* + (1+0)^* 000 (1+0)^*$
 $S \rightarrow x00x00x | x000x$
 $x \rightarrow 1x | 0x | \lambda$



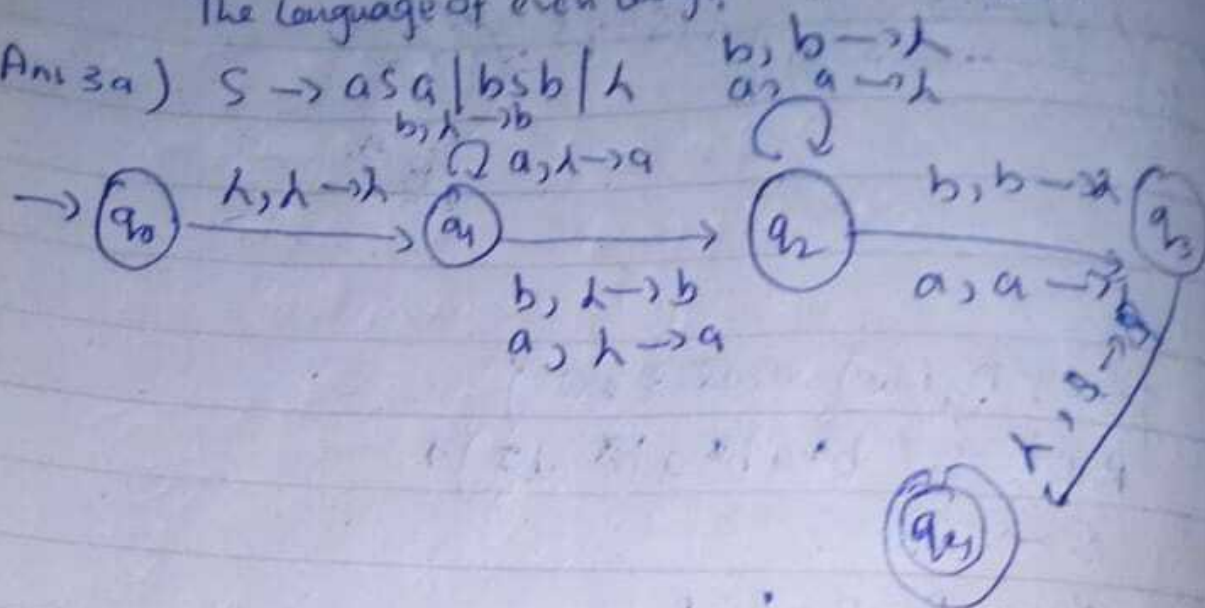
(h) $\{ w : n_a(w) \bmod 3 = 0 \}$
 R.E $\leq b(b^*ab^*ab^*ab^*)^*$
 $S \rightarrow SX$
 $X \rightarrow BaBaBaB$
 $B \rightarrow bB | h$



$a, a \rightarrow h$ $h, s \rightarrow s$
 $b, b \rightarrow h$

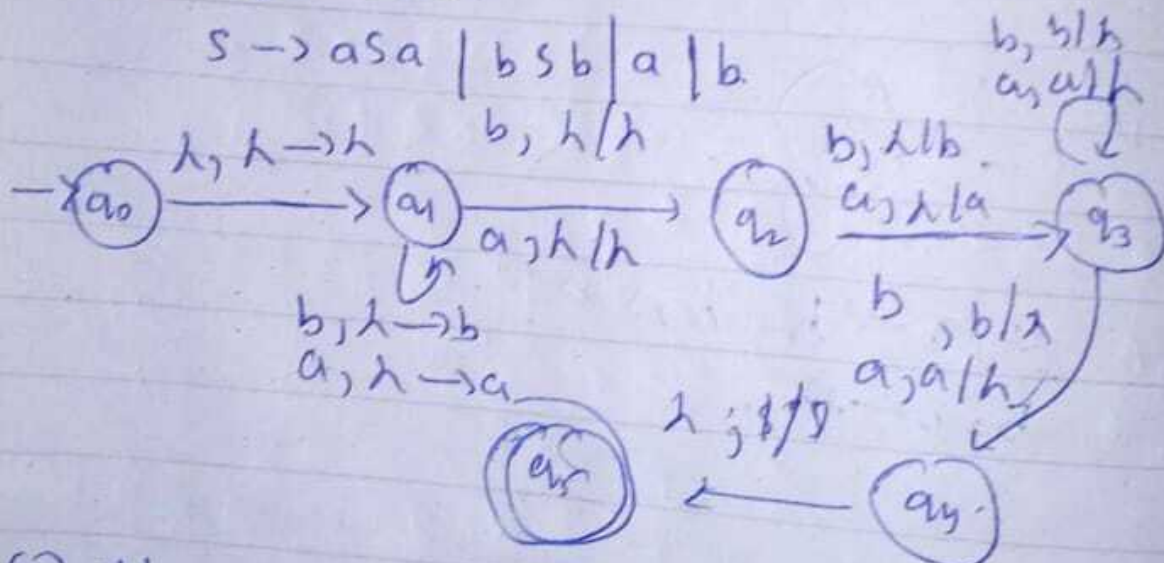
The language of even length palindromes

Ans 3a) $S \rightarrow aSa \mid bSb \mid \lambda$



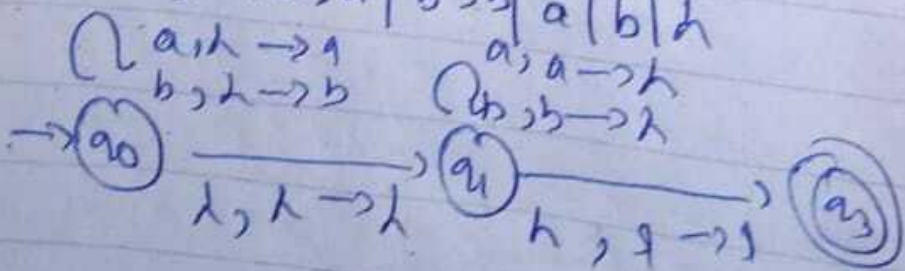
(b) Odd palindrome

$S \rightarrow aSa \mid bSb \mid a \mid b$

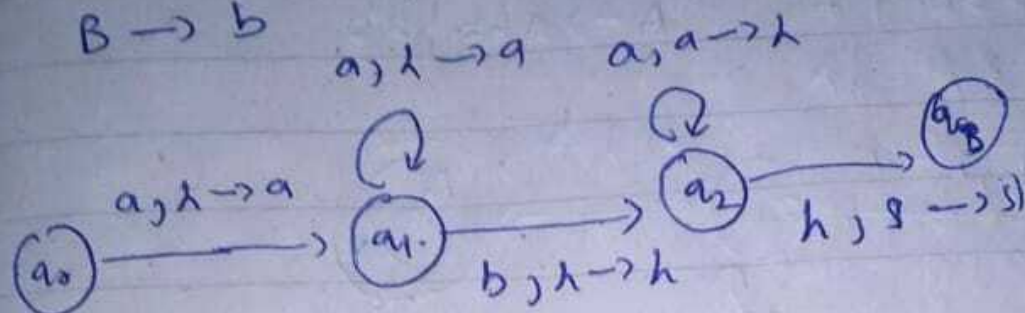


(c) all palindrome

$S \rightarrow aSa \mid bSb \mid a \mid b \mid \lambda$

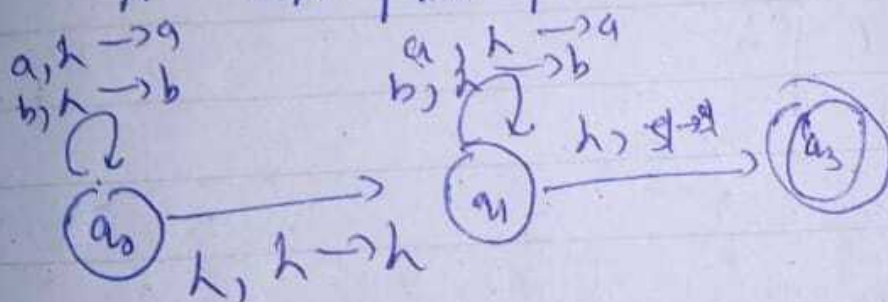


(d) $a^n b a^n$
 $S \rightarrow aSa \mid \epsilon$
 $B \rightarrow b$



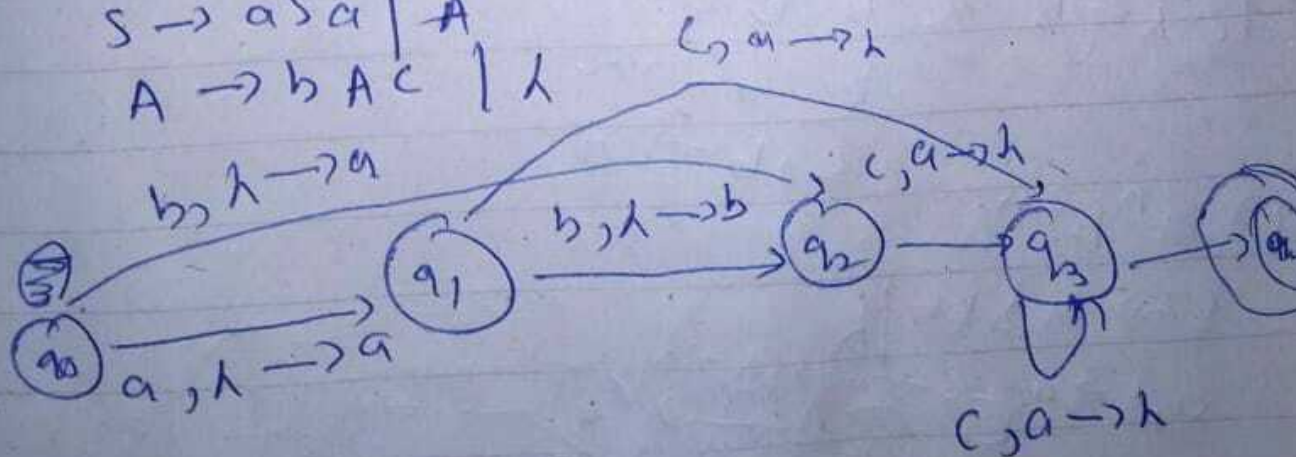
(e) $ww^R, w \in \{a, b\}^*$

$S \rightarrow aAb \mid \epsilon$
 $A \rightarrow aA \mid bA \mid \epsilon$



(f) $a^n b^m c^{n+m}$
 $a^n b^m c^m c^n$

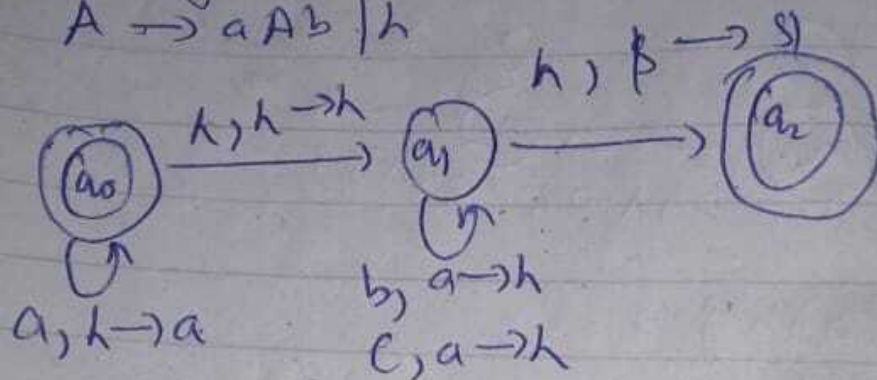
$S \rightarrow aSa \mid A$
 $A \rightarrow bAc \mid \epsilon$



(g) $a^{n+m} b^n c^m$

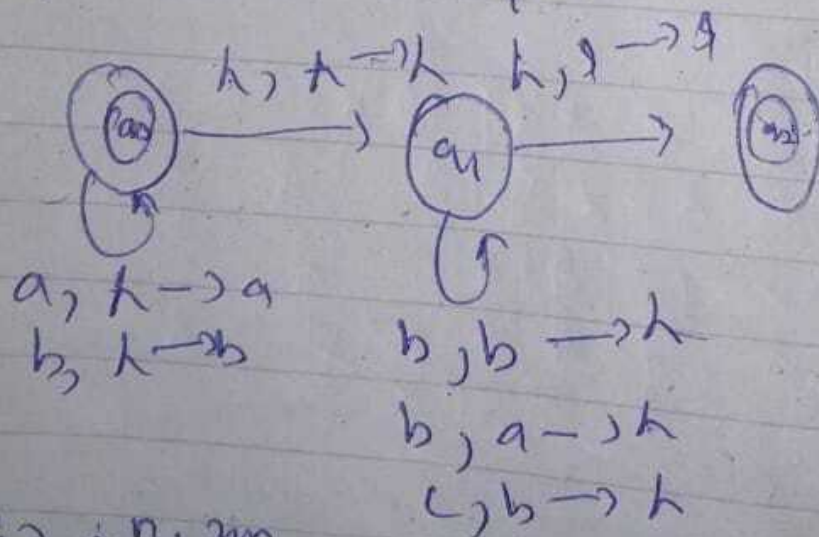
~~$S \rightarrow a^m a^n b^n c^m$~~
 ~~$S \rightarrow aSc$~~ | ~~A~~
 ~~$A \rightarrow aAb$~~ | h

$S \rightarrow aAc$ | h
 $A \rightarrow aAb$ | ab



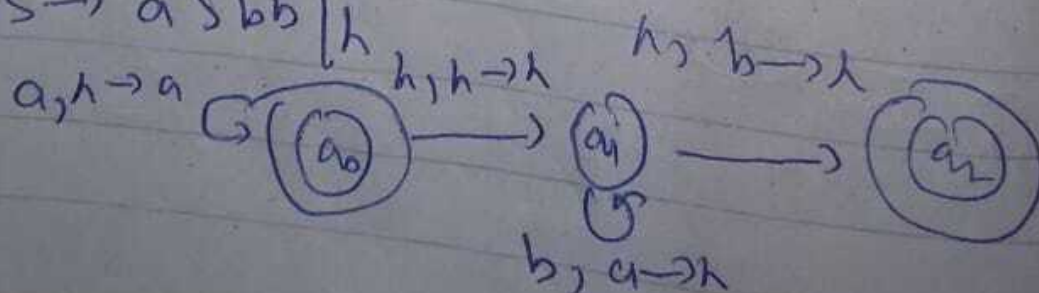
(h) $a^n b^{n+m} c^m$

$S \rightarrow aAc$ | h
 $A \rightarrow bAb$ | bb

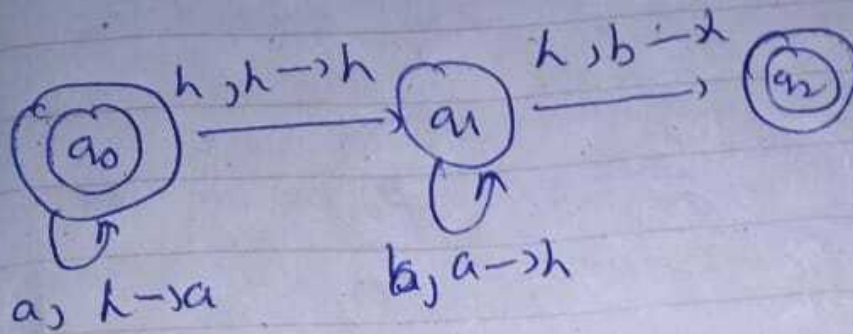


(i) $a^n b^{2m}$

$S \rightarrow aSbb$ | h



(j) $S \rightarrow aa S bbb \mid \epsilon$



Q4) (i) $S \rightarrow abS \mid abA \mid abB$

$A \rightarrow cd$

$B \rightarrow aB$

$C \rightarrow dC$

Remaining C not reachable

$S \rightarrow abS \mid abA \mid abB$

$A \rightarrow cd$

$B \rightarrow aB$

Ans.

(ii) $S \rightarrow ABC | a$

$A \rightarrow b$

$B \rightarrow c$

$C \rightarrow d$

$E \rightarrow e$

$F \rightarrow f$

$g \rightarrow g$

$S \rightarrow ABC | a$

$A \rightarrow b$

$B \rightarrow c$

$C \rightarrow d$

(iii) $S \rightarrow AB | bx$

$A \rightarrow Bad | bSx | c$

$B \rightarrow aSB | bBx$

$x \rightarrow SBD | aBx | cd$

Q3 (iii)

stack operations for

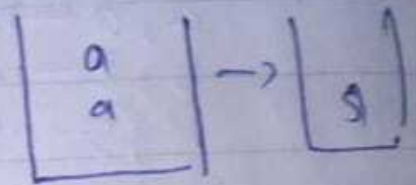
d) let $w = a a b a a$

→ push all a's

→ when b is encountered read it but do not

push (pop anything)

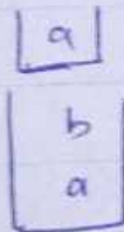
→ pop all a's



(f) let $w = a b c c$

→ push all a's

→ push all b's



→ Read all c's and for every c, pop an a or b out of the stack.

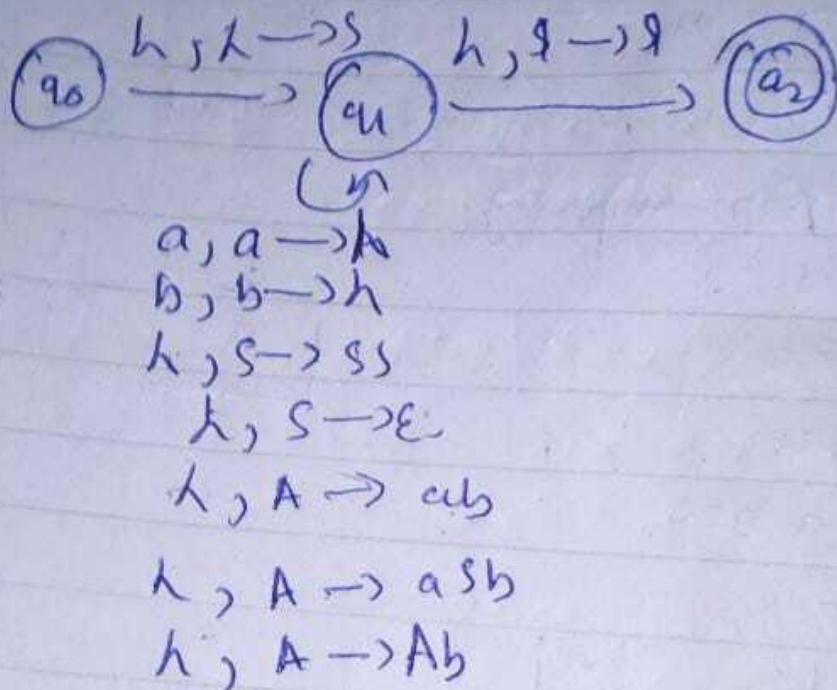
(a) let $w = a b b b b a$

→ Read and push all letters until an assumed midpoint (p.d. a is non-deterministic)



Q (a) $S \rightarrow SS | \epsilon$

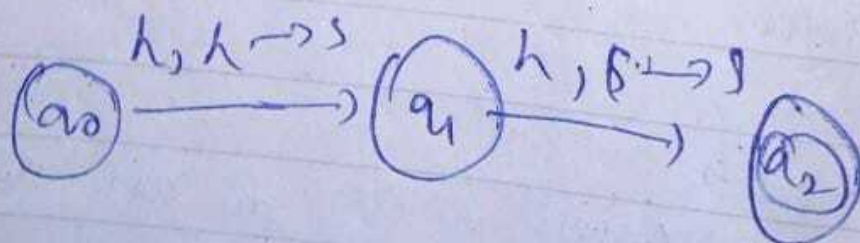
$A \rightarrow aSb | Ab | ab$



(b) $S \rightarrow S+x | x$

$x \rightarrow x*y | y$

$y \rightarrow S$



$h, x \rightarrow x*y$

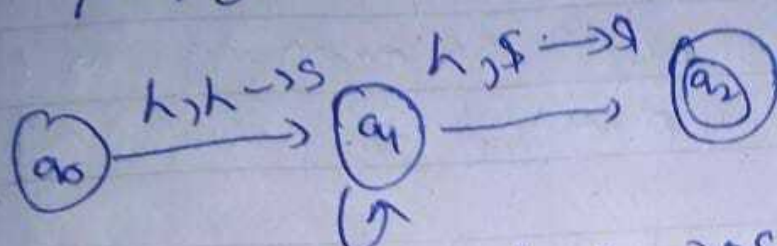
$h, x \rightarrow y$

$h, y \rightarrow S$

$h, S \rightarrow S+x$

$h, S \rightarrow x$

(c) $S \rightarrow OS1 \mid 1S0 \mid \Lambda$
 $S \rightarrow OSX \mid 1SY \mid \Lambda$
 $X \rightarrow 1$
 $Y \rightarrow 0$



$h, Y \rightarrow 0$
 $h, X \rightarrow 1$
 $h, S \rightarrow OSX$
 $h, S \rightarrow 1SY$
 $h, S \rightarrow 1$
 $h, S \rightarrow 0$
 $h, S \rightarrow h$

Qb(i) $S \rightarrow aS \mid \epsilon$

$S \rightarrow ABC$
 $A \rightarrow a$
 $B \rightarrow b$
 $C \rightarrow c$

$w = ABC$
 $w = aBC$
 $w = abc$
 $w = abc$

no Ambiguity as there is only one word in language and only one way to derive it

(#)

(ii) $S \rightarrow as \mid \epsilon$

$w = as$

$w = aas$

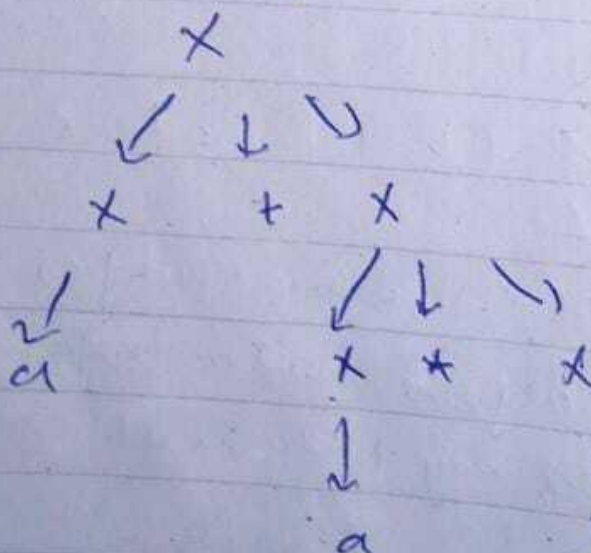
$w = aaas$

$w = aaaa$ no ambiguity as there is only one way to produce every word

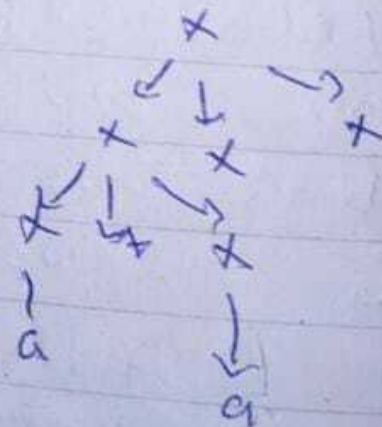
(iii) $X \rightarrow X+X \mid X * X \mid X \mid a$

ambiguous as the example word 'a * a + a' can be derived in multiple ways

Structure 1



Structure 2



~~Two structures for string ambiguity~~
 results