

并发容器类Map

JDK源码学习方法

逻辑思维能力是梳理学习方法的基础。养成线性思维：两个或者多个概念，像一条线串起来。

1. 演绎推导

因果推理。因为JAVA中网络编程只提供了BIO和NIO两种方式，所以一切框架中，涉及到网络处理的，都可以用着两个知识点去探究原理。

2. 归纳总结

可能正确的猜想。线上10台服务器，有三台总是每天会自动重启，收集相关信息后，发现是运维在修改监控系统配置的时候，漏掉了提高这三台机器的重启阈值。

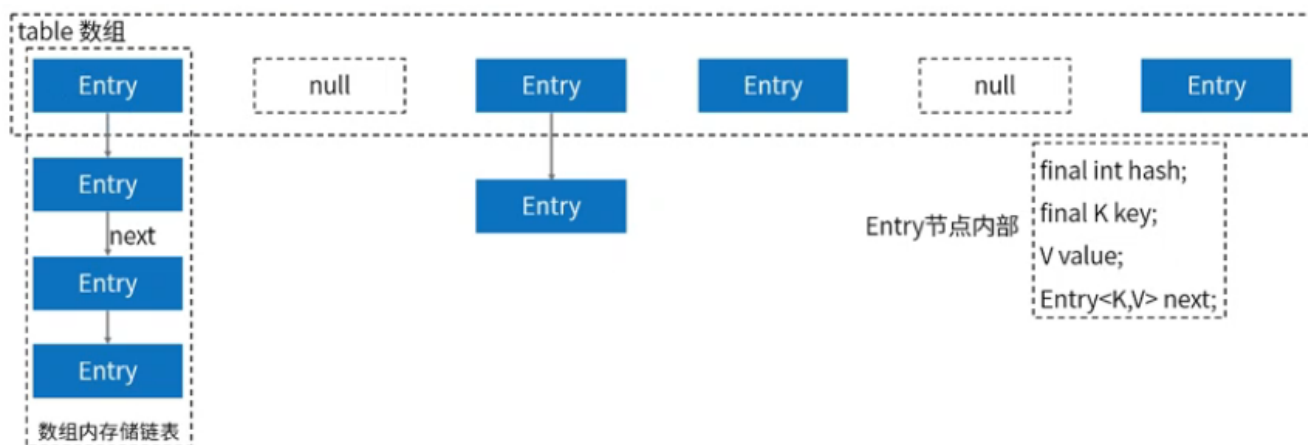
3. 类比法

集群就好像是马在拉车，一匹马拉不动的时候，就使用多匹马去拉。分布式的概念，就像是理发的过程中，洗头发和剪头发是不同的人负责的。

HashMap-非线程安全-JDK1.7

1. 存储结构

如下图所示：HashMap内部的存储结构是数组加链表，Hash不冲突的元素放在数组不同的位置，冲突的元素组织成链表，采用的是头插法。



JAVA8更新：链表中元素超过一个数量后，转变为红黑树结构



```
transient Entry<K,V>[] table;
```

Entry中的主要属性：

```
final K key;  
V value;  
Entry<K,V> next;  
int hash;
```

2. 初始化

```
static final int MAXIMUM_CAPACITY = 1 << 30;  
static final int DEFAULT_INITIAL_CAPACITY = 16;  
static final float DEFAULT_LOAD_FACTOR = 0.75f;  
// 元素数量  
transient int size;  
// 扩容的阈值，等于capacity*loadFactor  
int threshold;  
// 负载因子  
final float loadFactor;  
  
public HashMap() {  
    this(DEFAULT_INITIAL_CAPACITY, DEFAULT_LOAD_FACTOR);  
}  
  
public HashMap(int initialCapacity, float loadFactor) {  
    if (initialCapacity < 0)  
        throw new IllegalArgumentException("Illegal initial capacity: " +  
            initialCapacity);  
  
    if (initialCapacity > MAXIMUM_CAPACITY)  
        initialCapacity = MAXIMUM_CAPACITY;  
    if (loadFactor <= 0 || Float.isNaN(loadFactor))  
        throw new IllegalArgumentException("Illegal load factor: " +  
            loadFactor);  
  
    // 找到小于给定初始容量最接近2的次方的值  
    int capacity = 1;  
    while (capacity < initialCapacity)  
        capacity <<= 1;  
  
    this.loadFactor = loadFactor;  
    threshold = (int) Math.min(capacity * loadFactor, MAXIMUM_CAPACITY + 1);  
    table = new Entry[capacity];  
    useAltoHashing = sun.misc.VM.isBooted() &&  
        (capacity >= Holder.ALTERNATIVE_HASHING_THRESHOLD);  
    init();  
}
```

3. put方法

```
public V put(K key, V value) {  
    if (key == null)
```

```

        return putForNullKey(value);
    // 生成key的hash值
    int hash = hash(key);
    // 根据hash值定位在数组中的位置
    int i = indexFor(hash, table.length);
    for (Entry<K,V> e = table[i]; e != null; e = e.next) {
        Object k;
        // key完全一样则覆盖,返回旧值
        if (e.hash == hash && ((k = e.key) == key || key.equals(k))) {
            V oldValue = e.value;
            e.value = value;
            e.recordAccess(this);
            return oldValue;
        }
    }

    modCount++;
    // 生成新的Entry
    addEntry(hash, key, value, i);
    return null;
}

final int hash(Object k) {
    int h = 0;
    if (useAltHashing) {
        if (k instanceof String) {
            return sun.misc.Hashing.stringHash32((String) k);
        }
        h = hashSeed;
    }

    h ^= k.hashCode();

    h ^= (h >>> 20) ^ (h >>> 12);
    return h ^ (h >>> 7) ^ (h >>> 4);
}

static int indexFor(int h, int length) {
    return h & (length-1);
}

void addEntry(int hash, K key, V value, int bucketIndex) {
    // 当map中元素的数量大于阈值并且下一次放入发生冲突时会resize
    // 即HashMap最多在有27个元素时会发生扩容11 + 16
    if ((size >= threshold) && (null != table[bucketIndex])) {
        resize(2 * table.length);
        hash = (null != key) ? hash(key) : 0;
        bucketIndex = indexFor(hash, table.length);
    }

    createEntry(hash, key, value, bucketIndex);
}

```

```

void createEntry(int hash, K key, V value, int bucketIndex) {
    Entry<K,V> e = table[bucketIndex];
    // 头插法
    table[bucketIndex] = new Entry<>(hash, key, value, e);
    size++;
}

```

4. get方法

```

public V get(Object key) {
    if (key == null)
        return getForNullKey();
    Entry<K,V> entry = getEntry(key);

    return null == entry ? null : entry.getValue();
}

final Entry<K,V> getEntry(Object key) {
    int hash = (key == null) ? 0 : hash(key);
    // 根据hash值确定在数组中的下标
    for (Entry<K,V> e = table[indexFor(hash, table.length)];
        e != null;
        e = e.next) {
        Object k;
        // 如果在链表中的第一个节点找到就返回
        if (e.hash == hash &&
            ((k = e.key) == key || (key != null && key.equals(k))))
            return e;
    }
    return null;
}

```

HashMap-非线程安全-JDK1.8

1. 存储结构

1.8中在冲突不大的情况下还是数组加链表结构，在冲突严重时会链表会转换成红黑树。

```

transient Node<K,V>[] table;

// Node的结构
final int hash;
final K key;
V value;
Node<K,V> next;

```

2. 初始化

```

// 在1.7的基础上增加了如下几个值

// 链表要转换成红黑树的大小阈值
static final int TREEIFY_THRESHOLD = 8;

```

```

// 红黑树转换回链表的大小阈值
static final int UNTREEIFY_THRESHOLD = 6;
// 链表转红黑树的数组大小阈值
static final int MIN_TREEIFY_CAPACITY = 64;

public HashMap(int initialCapacity) {
    this(initialCapacity, DEFAULT_LOAD_FACTOR);
}

public HashMap(int initialCapacity, float loadFactor) {
    if (initialCapacity < 0)
        throw new IllegalArgumentException("Illegal initial capacity: " +
            initialCapacity);

    if (initialCapacity > MAXIMUM_CAPACITY)
        initialCapacity = MAXIMUM_CAPACITY;
    if (loadFactor <= 0 || Float.isNaN(loadFactor))
        throw new IllegalArgumentException("Illegal load factor: " +
            loadFactor);

    this.loadFactor = loadFactor;
    this.threshold = tableSizeFor(initialCapacity);
}

static final int tableSizeFor(int cap) {
    int n = cap - 1;
    n |= n >>> 1;
    n |= n >>> 2;
    n |= n >>> 4;
    n |= n >>> 8;
    n |= n >>> 16;
    return (n < 0) ? 1 : (n >= MAXIMUM_CAPACITY) ? MAXIMUM_CAPACITY : n + 1;
}

```

3. put方法

```

public V put(K key, V value) {
    return putVal(hash(key), key, value, false, true);
}

final V putVal(int hash, K key, V value, boolean onlyIfAbsent,
               boolean evict) {
    Node<K,V>[] tab; Node<K,V> p; int n, i;
    // 初始化
    if ((tab = table) == null || (n = tab.length) == 0)
        n = (tab = resize()).length;
    // 不冲突, 直接放
    if ((p = tab[i = (n - 1) & hash]) == null)
        tab[i] = newNode(hash, key, value, null);
    else {
        Node<K,V> e; K k;
        // key值完全一样则覆盖
        if (p.hash == hash &&
            ((k = p.key) == key || (key != null && key.equals(k))))
            e = p;
    }
}

```

```

        else if (p instanceof TreeNode)
            e = ((TreeNode<K,V>)p).putTreeVal(this, tab, hash, key, value);
        else {
            for (int binCount = 0; ; ++binCount) {
                // 找到链表尾部
                if ((e = p.next) == null) {
                    // hash冲突插入链表, 尾插法
                    p.next = newNode(hash, key, value, null);
                    // 如果链表中元素多余8个转红黑树
                    if (binCount >= TREEIFY_THRESHOLD - 1) // -1 for 1st
                        treeifyBin(tab, hash);
                    break;
                }
                if (e.hash == hash &&
                    ((k = e.key) == key || (key != null && key.equals(k))))
                    break;
                // 更新循环变量
                p = e;
            }
            if (e != null) { // existing mapping for key
                V oldValue = e.value;
                if (!onlyIfAbsent || oldValue == null)
                    e.value = value;
                afterNodeAccess(e);
                return oldValue;
            }
        }
        ++modCount;
        // 大于阈值直接扩容
        if (++size > threshold)
            resize();
        afterNodeInsertion(evict);
        return null;
    }
}

```

```

final void treeifyBin(Node<K,V>[] tab, int hash) {
    int n, index; Node<K,V> e;
    // 如果转换红黑树时, 数组小组64则扩容
    // 此时发生冲突的主要原因是数组太小
    if (tab == null || (n = tab.length) < MIN_TREEIFY_CAPACITY)
        resize();
    else if ((e = tab[index = (n - 1) & hash]) != null) {
        TreeNode<K,V> hd = null, tl = null;
        do {
            TreeNode<K,V> p = replacementTreeNode(e, null);
            if (tl == null)
                hd = p;
            else {
                p.prev = tl;
                tl.next = p;
            }
        } while ((e = e.next) != null);
        tl = p;
    }
}

```

```

        } while ((e = e.next) != null);
        if ((tab[index] = hd) != null)
            hd.treeify(tab);
    }
}

```

4. get方法

```

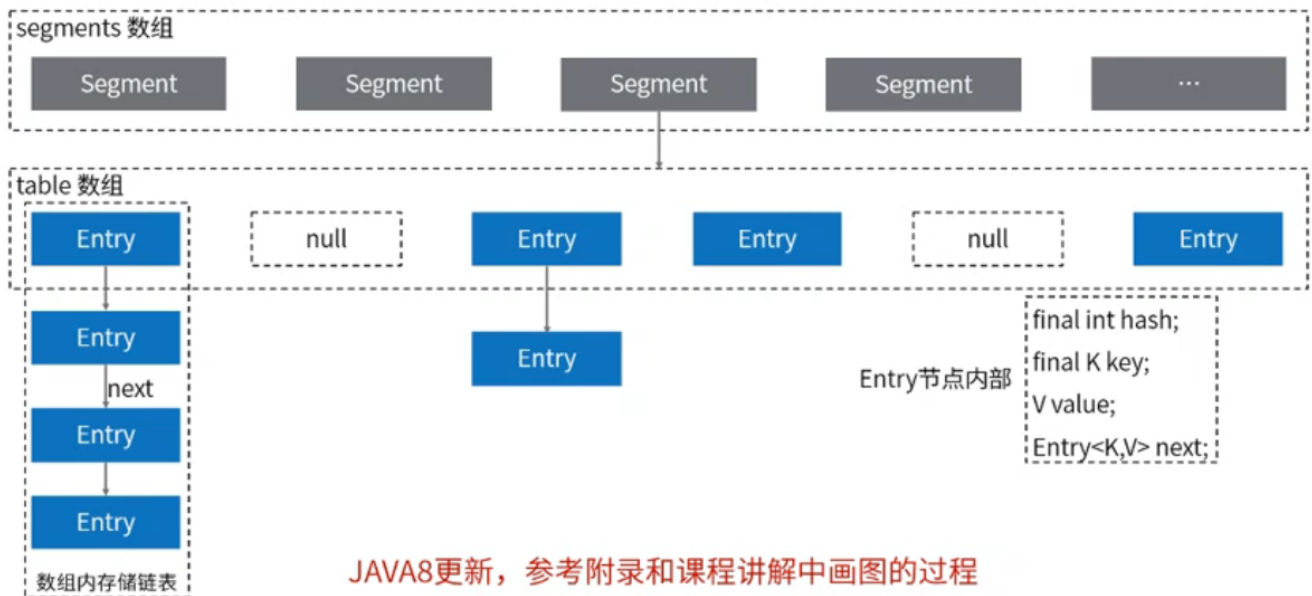
public V get(Object key) {
    Node<K,V> e;
    return (e = getNode(hash(key), key)) == null ? null : e.value;
}

final Node<K,V> getNode(int hash, Object key) {
    Node<K,V>[] tab; Node<K,V> first, e; int n; K k;
    if ((tab = table) != null && (n = tab.length) > 0 &&
        (first = tab[(n - 1) & hash]) != null) {
        // 链表第一个位置即为要找元素
        if (first.hash == hash && // always check first node
            ((k = first.key) == key || (key != null && key.equals(k))))
            return first;
        if ((e = first.next) != null) {
            if (first instanceof TreeNode)
                // 在红黑树中查找
                return ((TreeNode<K,V>)first).getTreeNode(hash, key);
            do {
                // 在链表中查找
                if (e.hash == hash &&
                    ((k = e.key) == key || (key != null && key.equals(k))))
                    return e;
            } while ((e = e.next) != null);
        }
    }
    return null;
}

```

ConcurrentHashMap-线程安全-JDK1.7

HashTable在每个put和get方法加同步锁，效率低，推荐使用ConcurrentHashMap，包含一层segment确保线程安全。



1. 存储结构

`ConcurrentHashMap`采用了分段加锁的机制, 默认16个段, 每个段之间的读写互不影响。

```
final Segment<K,V>[] segments;

// segment其实就是一个锁
static final class Segment<K,V> extends ReentrantLock implements Serializable
// 每个segment里又有一个HashEntry数组
transient volatile HashEntry<K,V>[] table;

// HashEntry结构
final int hash;
final K key;
volatile V value;
volatile HashEntry<K,V> next;
```

2. 初始化

```
static final int DEFAULT_INITIAL_CAPACITY = 16;
static final float DEFAULT_LOAD_FACTOR = 0.75f;
static final int DEFAULT_CONCURRENCY_LEVEL = 16;
static final int MAXIMUM_CAPACITY = 1 << 30;
static final int MIN_SEGMENT_TABLE_CAPACITY = 2;
static final int MAX_SEGMENTS = 1 << 16;

public ConcurrentHashMap() {
    this(DEFAULT_INITIAL_CAPACITY, DEFAULT_LOAD_FACTOR, DEFAULT_CONCURRENCY_LEVEL);
}

public ConcurrentHashMap(int initialCapacity, float loadFactor) {
    this(initialCapacity, loadFactor, DEFAULT_CONCURRENCY_LEVEL);
}

public ConcurrentHashMap(int initialCapacity,
```



```

        float loadFactor, int concurrencyLevel) {
    if (!(loadFactor > 0) || initialCapacity < 0 || concurrencyLevel <= 0)
        throw new IllegalArgumentException();
    if (concurrencyLevel > MAX_SEGMENTS)
        concurrencyLevel = MAX_SEGMENTS;
    // 适合参数的2的次方的值
    int sshift = 0;
    int ssize = 1;
    while (ssize < concurrencyLevel) {
        ++sshift;
        ssize <= 1;
    }
    this.segmentShift = 32 - sshift;
    this.segmentMask = ssize - 1;
    if (initialCapacity > MAXIMUM_CAPACITY)
        initialCapacity = MAXIMUM_CAPACITY;
    int c = initialCapacity / ssize;
    if (c * ssize < initialCapacity)
        ++c;
    int cap = MIN_SEGMENT_TABLE_CAPACITY;
    while (cap < c)
        cap <= 1;
    // create segments and segments[0]
    Segment<K,V> s0 =
        new Segment<K,V>(loadFactor, (int)(cap * loadFactor),
                        (HashEntry<K,V>[])new HashEntry[cap]);
    Segment<K,V>[] ss = (Segment<K,V>[])new Segment[ssize];
    UNSAFE.putOrderedObject(ss, SBASE, s0); // ordered write of segments[0]
    this.segments = ss;
}

```

3. put方法

```

public V put(K key, V value) {
    Segment<K,V> s;
    if (value == null)
        throw new NullPointerException();
    int hash = hash(key);
    // 定位在那个segment
    int j = (hash >>> segmentShift) & segmentMask;
    if ((s = (Segment<K,V>)UNSAFE.getObject          // nonvolatile; recheck
         (segments, (j << SSHIFT) + SBASE)) == null) // in ensureSegment
        s = ensureSegment(j);
    return s.put(key, hash, value, false);
}

final V put(K key, int hash, V value, boolean onlyIfAbsent) {
    // put时给每个段加锁
    HashEntry<K,V> node = tryLock() ? null :
        scanAndLockForPut(key, hash, value);
    V oldValue;
    try {
        HashEntry<K,V>[] tab = table;

```

```

        int index = (tab.length - 1) & hash;
        // 链表第一个元素
        HashEntry<K,V> first = entryAt(tab, index);
        for (HashEntry<K,V> e = first;;) {
            // 冲突
            if (e != null) {
                K k;
                // key完全相同覆盖
                if ((k = e.key) == key ||
                    (e.hash == hash && key.equals(k))) {
                    oldValue = e.value;
                    if (!onlyIfAbsent) {
                        e.value = value;
                        ++modCount;
                    }
                    break;
                }
                e = e.next;
            }
            else {
                if (node != null)
                    // 尾插法
                    node.setNext(first);
                else
                    // 第一个元素
                    node = new HashEntry<K,V>(hash, key, value, first);
                int c = count + 1;
                // 超过阈值并且数组大小小于最大值则扩容
                if (c > threshold && tab.length < MAXIMUM_CAPACITY)
                    rehash(node);
                else
                    // 更新segment
                    setEntryAt(tab, index, node);
                ++modCount;
                count = c;
                oldValue = null;
                break;
            }
        }
    } finally {
        unlock();
    }
    return oldValue;
}

```

4. get方法

```

public V get(Object key) {
    Segment<K,V> s; // manually integrate access methods to reduce overhead
    HashEntry<K,V>[] tab;
    int h = hash(key);
    long u = (((h >>> segmentShift) & segmentMask) << SSHIFT) + SBASE;
    // 定位segment

```

```

        if ((s = (Segment<K,V>)UNSAFE.getObjectVolatile(segments, u)) != null &&
            (tab = s.table) != null) {
            // 判断链表中元素
            for (HashEntry<K,V> e = (HashEntry<K,V>) UNSAFE.getObjectVolatile
                (tab, ((long)(((tab.length - 1) & h)) << TSHIFT) + TBASE);
                e != null; e = e.next) {
                K k;
                if ((k = e.key) == key || (e.hash == h && key.equals(k)))
                    return e.value;
            }
        }
        return null;
    }
}

```

ConcurrentHashMap-线程安全-JDK1.8

1. 存储结构

和JDK1.8中的HashMap的存储结构一致，兼容1.7的segment，在冲突严重时会转成红黑树。对线程安全处理方式进行了更新，在没有冲突时采用CAS无锁机制，在有冲突时进行加锁。

2. 初始化

```

public ConcurrentHashMap(int initialCapacity,
                        float loadFactor, int concurrencyLevel) {
    if (!(loadFactor > 0.0f) || initialCapacity < 0 || concurrencyLevel <= 0)
        throw new IllegalArgumentException();
    if (initialCapacity < concurrencyLevel) // Use at least as many bins
        initialCapacity = concurrencyLevel; // as estimated threads
    long size = (long)(1.0 + (long)initialCapacity / loadFactor);
    int cap = (size >= (long)MAXIMUM_CAPACITY) ?
        MAXIMUM_CAPACITY : tableSizeFor((int)size);
    this.sizeCtl = cap;
}

```

3. put方法

```

public V put(K key, V value) {
    return putVal(key, value, false);
}

final V putVal(K key, V value, boolean onlyIfAbsent) {
    if (key == null || value == null) throw new NullPointerException();
    int hash = spread(key.hashCode());
    int binCount = 0;
    for (Node<K,V>[] tab = table;;) {
        Node<K,V> f; int n, i, fh;
        // 初始化
        if (tab == null || (n = tab.length) == 0)
            tab = initTable();
        // 没有冲突采用cas机制放值
        else if ((f = tabAt(tab, i = (n - 1) & hash)) == null) {
            if (casTabAt(tab, i, null,

```

```

        new Node<K,V>(hash, key, value, null)))
        break; // no lock when adding to empty bin
    }
    else if ((fh = f.hash) == MOVED)
        tab = helpTransfer(tab, f);
    else {
        V oldVal = null;
        synchronized (f) {
            if (tabAt(tab, i) == f) {
                if (fh >= 0) {
                    binCount = 1;
                    for (Node<K,V> e = f;; ++binCount) {
                        K ek;
                        // key完全相同覆盖
                        if (e.hash == hash &&
                            ((ek = e.key) == key ||
                             (ek != null && key.equals(ek)))) {
                            oldVal = e.val;
                            if (!onlyIfAbsent)
                                e.val = value;
                            break;
                        }
                        Node<K,V> pred = e;
                        // 插入链表尾部
                        if ((e = e.next) == null) {
                            pred.next = new Node<K,V>(hash, key,
                                                         value, null);
                            break;
                        }
                    }
                }
            }
            else if (f instanceof TreeBin) {
                Node<K,V> p;
                binCount = 2;
                if ((p = ((TreeBin<K,V>)f).putTreeVal(hash, key,
                                                         value)) != null) {
                    oldVal = p.val;
                    if (!onlyIfAbsent)
                        p.val = value;
                }
            }
        }
    }
    if (binCount != 0) {
        // 转换为红黑树
        if (binCount >= TREEIFY_THRESHOLD)
            treeifyBin(tab, i);
        if (oldVal != null)
            return oldVal;
        break;
    }
}
}

```

```
        addCount(1L, binCount);
        return null;
    }
```

4. get方法

```
public V get(Object key) {
    Node<K,V>[] tab; Node<K,V> e, p; int n, eh; K ek;
    int h = spread(key.hashCode());
    if ((tab = table) != null && (n = tab.length) > 0 &&
        (e = tabAt(tab, (n - 1) & h)) != null) {
        if ((eh = e.hash) == h) {
            if ((ek = e.key) == key || (ek != null && key.equals(ek)))
                return e.val;
        }
        else if (eh < 0)
            return (p = e.find(h, key)) != null ? p.val : null;
        while ((e = e.next) != null) {
            if (e.hash == h &&
                ((ek = e.key) == key || (ek != null && key.equals(ek))))
                return e.val;
        }
    }
    return null;
}
```