

# EXAM 351 MIDTERM

1) 30 40 10 20

a) 6

Bubblesort comparison: 1+1+1+

40

1+1

b) 7

Insertion sort with a sentinel: 1+1+1+1

1+1

1

c) 4

Mergesort:

30 40 10 20

10 20

1+1+1+1

1

1

d) 5

Heap

40 30

30

e) 3

40 10

20

30

10

20

30

10

1+1+1+1+1+1

40

30

20

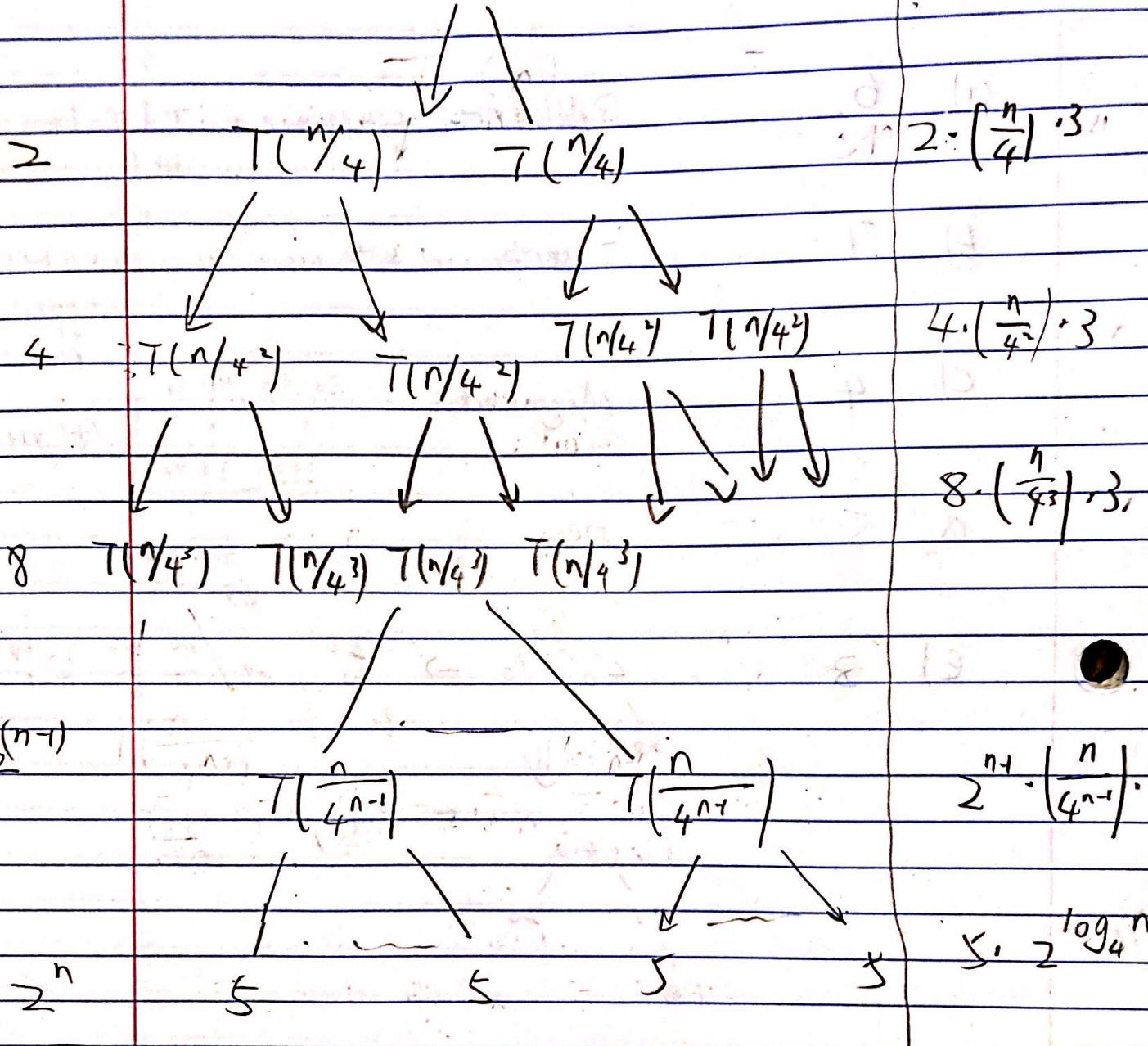
10

Quicksort

30 20 10 20 1+ 10 20

3.

$$1 \quad ?) \quad T(n) = 2T(n/4) + 3n \quad T(1) = 5 \quad 31$$



$$\begin{aligned}
 T(n) &= \sum_{i=1}^{\log_4 n - 1} 2^i \cdot 3 \left(\frac{n}{4^i}\right) = 3n \sum_{i=0}^{\log_4 n - 1} \left(\frac{1}{2}\right)^i = 3n \cdot \left( \frac{\frac{1}{2} \log_4 n - 1}{\frac{1}{2} - 1} \right) \\
 &= 3n \cdot \left( \frac{n \log_4 \left(\frac{1}{2}\right) - 1}{-\frac{1}{2}} \right) = 3n \left( \frac{\frac{1}{2}n - 1}{-\frac{1}{2}} \right) \\
 &= -\frac{\frac{3n}{2}}{2} + \frac{3n}{2} = \boxed{6n - \frac{bn}{\sqrt{n}}} \\
 T_n &= T(0) = \boxed{6n - \frac{bn}{\sqrt{n}}}
 \end{aligned}$$

3)  $R[0], R[1], R[2], \dots, R[m-1]$

$$m = s + 1$$

read(g, h, d), write(g, h, d)

a) Selection sort sends code till  $(i, j, k, n)$ ,

$s$  always works till

for  $i=n$  down to 2 do

if  $i \leq s$  then read ( $l, l, i$ )  $\leftarrow k \leftarrow i$

for  $j=2$  to  $i-1$  do

if  $A[j] > A[k]$  then  $A[j] \leftrightarrow A[k]$

end if

end for  $\leftarrow A[k] \leftrightarrow A[i]$

write ( $l, l, i$ )

else

$R[0] \leftarrow R[\min(s, i)]$

$r \leftarrow \min(s, i)$

while  $r < s$  do

read ( $l, r, s$ )

for  $j=0$  to  $s-1$  do

if  $R[j] > R[j+1]$

end if

end for

write ( $r, 0, s$ )

$[j=0] \leftarrow R[s]$

$r \leftarrow r+s$

end while

read ( $r, l, i-r$ )

for  $j=2$  to  $i-r$  do

if  $R[j] > R[j+1]$

Then  $R[j] \leftarrow R[j+1]$   
 end if  
 end for  
 write ( $r, 0, i-1$ )  
 end if  
 end for

b) case of memory  $f(h)=3$  ~~copy exchange~~

$n > s$   
 $\alpha, B, s, n$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \beta s = B \sum_{i=2}^n s \sum_{k=1}^{i/s} 1 = B \sum_{i=2}^n s \cancel{\frac{i}{s}} - B \sum_{i=2}^n i = \frac{Bn^2}{2}$$

when  $f(h)=3$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \alpha = \alpha \sum_{i=2}^n \sum_{k=1}^{i/s} 1 = \alpha \sum_{i=2}^n i/s = \alpha/s \sum_{i=2}^n i = \frac{\alpha n^2}{6s}$$

When  $f(h)=h$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \alpha ks = \alpha \sum_{i=2}^n s \sum_{k=1}^{i/s} k = \alpha \sum_{i=2}^n s \frac{i^2}{2s^2} = \frac{\alpha}{2s} \sum_{i=2}^n i^2 = \frac{\alpha h^3}{6s}$$

## Problem 4.

index (mod 3) 1, 2, 3

insertion

a) Sedg code Not recursive

for  $i=4 \rightarrow n$  doif  $i \equiv 1 \pmod{3}$   
temp  $\leftarrow A[i]$  $j \leftarrow i-1 \quad j \leftarrow i-3$   
while  $j \geq 0$  and  $A[j] > temp$  do begin $A[j+1] \leftarrow A[j]$  $j \leftarrow j-3$ 

end while

 $A[j+3] \leftarrow temp$ 

end if

end for

for  $i=5 \rightarrow n$  doif  $i \equiv 2 \pmod{3}$ temp  $\leftarrow A[i]$  $j \leftarrow i-1 \quad j \leftarrow i-3$ while  $j \geq 0$  and  $A[j] > temp$  do begin $A[j+1] \leftarrow A[j]$  $j \leftarrow j-3$ 

end while

 $A[j+3] \leftarrow temp$ 

end if

end for

for  $i=6 \rightarrow n$  doif  $i \equiv 3 \pmod{3}$ temp  $\leftarrow A[i]$  $j \leftarrow i-1 \quad j \leftarrow i-3$ while  $j \geq 0$  and  $A[j] > temp$  do begin $A[j+1] \leftarrow A[j]$  $j \leftarrow j-3$  $A[j+3] \leftarrow temp$  end while

End if  
 end for  
 for  $i = 7 \rightarrow n$  do  
 if  $(i = 0 \pmod{3})$   
 temp  $\leftarrow A[i]$   
 $j \leftarrow i+1$ ;  $j \leftarrow i-1$   
 while  $j > 0$  and  $A[j] \geq \text{temp}$  do begin  
 $A[j+1] \leftarrow A[j]$   
 $j \leftarrow j-3$   
 end while  
 end if  
 end for

b)  $n=12$  Best case comparison,

$\begin{array}{c} 1 \ 2 \ 3 \ | \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \\ \hline \end{array}$

$\begin{array}{ccccccc} 1 & 4 & 7 & 10 & \text{mod } 3 & ① & -3 \\ 2 & 5 & 8 & 11 & \text{mod } 3 & ② & -3 \\ 3 & 6 & 9 & 12 & \text{mod } 3 & ③ & -3 \end{array}$

$$9 + 11 = \boxed{20}$$

c)  $n=12$  worst case

$$\sum_{i=2}^n \sum_{j=1}^{i-1} 1 = \sum_{i=2}^n (i-1)$$

10 4 1

$\begin{array}{cccccc} 10 & 7 & 4 & 1 & - & \boxed{6} \\ 11 & 8 & 5 & 2 & - & \boxed{6} \\ 12 & 9 & 6 & 3 & - & \boxed{6} \end{array}$

$$= \sum_{i=2}^n i - \frac{1}{2} n(n+1)$$

$$= 3 \left( \frac{n(n+1)}{2} \right)$$

$$3 \cdot \left( \frac{4 \cdot 3}{2} \right) = 3 \cdot 6 = \boxed{18} + \boxed{\cancel{18}} = \boxed{36}$$

Final Answer

d) worst case  $n = 3k$  ( $k \in \mathbb{N}^{>0}$ )

$n \quad n-1 \quad n-2 \quad \dots \quad 2 \quad 1$

$$n \cdot \frac{2n}{3} \cdot \frac{n}{3} \cdot \dots$$

$$n-1 \cdot \frac{2n-1}{3} \cdot \frac{n-1}{3} \cdot \dots$$

sort 1 index comparison  $\frac{2n}{3}$

2 index  $\frac{2n+1}{3}$

3

$\dots \frac{2n+1}{3}$

$$1+2+2+\sum_{i=4}^{\frac{n}{3}} 2i + \sum_{i=5}^{\frac{n}{3}} 2i+1 + \sum_{i=6}^{\frac{n}{3}}$$

5) Assume  $(i, j) \leftarrow \text{MinMax}(A, p, r)$

↓  
smallest      largest  
Array

bhly: 1 compression

a)  $\text{test}(A, p, r) \{$

$B[ ]$  .clearall; // meaning  $B[1] - B[2] - B[n] = 0$

$k = 1$

while  $A.\text{length} \geq 1$

$(i, j) \leftarrow \text{minmax}(A, 1, A.\text{length})$

$B[k] = i$

$B[n-k] = j$

$k++;$

$A[i, j]$  Delete from Array

End while.

if  $A.\text{length} = 1$  then

$B[k] = A[n]$

end if

Return  $B$ ;

}

b) no. comparisons:  $B$  is even  $n/2$

$B$  is odd  $n-1/2$

$$C) \frac{n(\frac{3n}{2}-2)}{2} = \frac{3n^2}{4} - n$$