

NP#3

7. T & F, CNF

a) $(A \vee \bar{B}) (B \vee \bar{C}) (\bar{A} \vee C)$

b) 2-CNF no var occurs twice

$(A \vee B) (\underline{B \vee \bar{C}}) (\underline{\bar{A} \vee C}) (\underline{B \vee \bar{C}}) (A \vee C)$

c) $(A \vee \textcircled{B} \vee C) (A \vee \textcircled{B} \vee \bar{C}) (A \vee \bar{B} \vee \textcircled{C}) (\bar{A} \vee \textcircled{B} \vee \bar{C})$

$(A \vee \textcircled{B} \vee \bar{D}) (\bar{A} \vee \textcircled{B} \vee \bar{D}) (A \vee \bar{B} \vee \textcircled{D}) (\bar{A} \vee \textcircled{B} \vee D)$

$(A \vee \textcircled{C} \vee D) (A \vee \textcircled{C} \vee \bar{D}) (\bar{A} \vee \textcircled{C} \vee D) (\bar{A} \vee \bar{C} \vee D)$

$(\textcircled{B} \vee C \vee \textcircled{D}) (\textcircled{B} \vee \bar{C} \vee D) (\bar{B} \vee \textcircled{C} \vee D) (\bar{B} \vee \bar{C} \vee \textcircled{D})$

A, D = False

B, C = TRUE

\Rightarrow TRUE

$$n \Rightarrow (A/V./W) \quad n=200$$

8.

a) If the formula is satisfiable,
at least one set of values to the boolean
is required which makes the set TRUE

$$\text{Time: } O(2^n) \quad \underline{2^1 \rightarrow 2^2 \rightarrow 2^3}$$

b) To prove you are right \Rightarrow UN SATISFIABLE

ONLY TIME YOU SHOW ME, Therefore
We are only showing the false values in the
boolean formula.

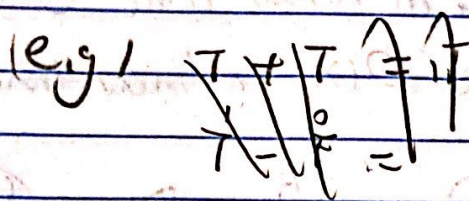
$$\text{Time: } O(n^2)$$

a) Boolean circuit \rightarrow satisfying (SAT)

We make true as 1 false as 0, now it is binary

To make the program tell the difference of (SAT)

We include the addition and minus into our algorithm



+ is and

- is or

* is negative/not

$$T + T = T$$

$$T - T = F$$

$$F - T = T$$

With the algorithm I can tell the

results from the addition and minus and the product

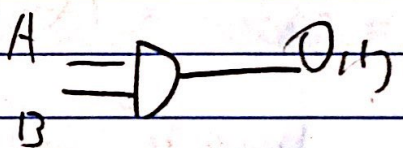
the reason I didn't illustrate my algorithm is because

it is not said in the question itself to write one.

b) $O(m)$ How fast determine satisfiable?

$\left\{ \begin{array}{l} \text{Number of input \& output} \\ \text{No. gates.} \end{array} \right.$

Since circuit can be done in linear time $O(m)$ with size X , we have $O(X)$ in linear time

e.g. $A \wedge B \Rightarrow O(1)$ 

With the implementation depending on gates & circuits the size of n will be conducted in same time $\Theta(n)$

c) $O(m^r)$

$$T : \Rightarrow O(m^r)$$