

HW 9 Yizhan Ao

1. a) $G: A[1..n, 1..n] \in \{0, 1\}$ list $R \Rightarrow G$

Matrix should use a 2D Array representing Graph

Initialize 2D Array size n^2 n vertices $arr[n][n]$

$i = 0$

if ($i < n$) {

$j = 0$

if ($j < n$) {

if (exist edge in $i \neq j$) {

$arr[i][j] = 1$

} else {

$arr[i][j] = 0$

}

} else {

$j++$

}

} else {

$i++$

}

}

b) initialize $m = \text{arr}[n]$ with size $= n$ }

$i = 0, j = 0$

if ($i < n$) {

$k = \text{empty linked list}$

$\text{arr}[i] = k$

if ($j < n$) {

if (edge V exists in $[i] \rightarrow [j]$) {

Add to list L

$\text{arr}[j]$

} else {

$j++$

}

else {

$i++$

}

}

2. a) If the proposed shortest path

There are two things needed to check

① CYCLE in the tree

② Root vertex reaching all other
Use DFS check $O(m+n)$

Based on above we can use pushing root into stack making it to be visited, iterating root through the element, if anything visited already then, there is a circle we have

b) test (m , adj, visited a)

visited[m] = true

for (each n in adj[n]) do

if (\neg visited[n]) then

test (n , adj, visited)

if $d[n] > d[m] + w(m, n)$ then
return false

end if

else

return false

end if

return true

end test

for ($i = 1 \rightarrow \text{adj.length}$)

visited[i] = false

end for


```

1 if ( test (V, adj, visited) / then }
    for ( j = 1 to visited.length )
        if ( ! visited [j] ) then
            return "No reaching"
        end if
    end for
    else
        return "Tree and tes"
    else
        return "No"
1 end if

```

3. a) Through edge (a,b) which is shorter

We will make the distance from each vertex to b in case that that vertex [i] if $D[i,a] + s < D[i,b]$ then $D[i,b] = D[i,a] + s$. then update the distance between each vertex as much for each vertex [i] and other vertex will be j only if $D[i,b] + D[b,j] < D[i,j]$ then $D[i,j] = D[i,b] + D[b,j]$

b) Graph "G" : Weight Matrix, "M"
Distance Matrix "D" edge(a,b) new edge (a,b)s

Update Distance | G, M, D, (a,b,s) }

$D[a,b] = s$
for (i in G (a,b))

if $(D[i,a] + s < D[i,b])$ {
 $D[i,b] = D[i,a] + s$ }

for (i in G) {

for (j in G) {

if $(D[i,b] + D[b,j] < D[i,j])$ {

$D[i,j] = D[i,b] + D[b,j]$ }

}

}

}