

3)  $R[0], R[1], R[2], \dots, R[m-1]$

$m = s + 16$

$\text{read}(g, h, d), \text{write}(g, h, d)$

a) Selection sort pseudo code (i, j, k, n),  
S always divides i/

for  $i = n$  down to 2 do

if  $i \leq s$  then  $\text{read}(1, 1, i) \leftarrow k \leftarrow i$

for  $j = 2$  to  $i$  do

if  $A[j] > A[k]$  then  $A[j] \leftrightarrow A[k]$

end if

end for  $\leftarrow A[k] \leftrightarrow A[i]$

$\text{write}(1, 1, i)$

else

$R[0] \leftarrow R[\min(s, i)]$

$r \leftarrow \min(s, i)$

while  $r < i$  do

$\text{read}(1, r, s)$

for  $j = 0$  to  $s-1$  do

if  $R[j] > R[j+1]$

end if

end for

$\text{write}(1, 0, s)$

$j[0] \leftarrow R[s]$

$r \leftarrow r+s$

end while

$\text{read}(r, 1, i-r)$

for  $j = 0$  to  $i-r$  do

if  $R[j] > R[j+1]$



Then  $RC[j] \leftrightarrow RC[j+1]$   
 end if  
 end for  
 write  $(r, 0, i-1)$   
 end if  
 end for

b) case of memory  $f(h)=3$  ~~copy~~ ~~exchange~~  
 $n \gg s$   
 $\propto B, s, n$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \beta s = B \sum_{i=2}^n s \sum_{k=1}^{i/s} 1 = B \sum_{i=2}^n \frac{i}{s} = B \sum_{i=2}^n i = \frac{Bn^2}{2}$$

when  $f(h)=3$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \alpha = \alpha \sum_{i=2}^n \sum_{k=1}^{i/s} 1 = \alpha \sum_{i=2}^n i/s = \alpha/s \sum_{i=2}^n i$$

$$\text{when } f(h)=h \quad = \frac{\alpha n^2}{6s}$$

$$\sum_{i=2}^n \sum_{k=1}^{i/s} \alpha ks = \alpha \sum_{i=2}^n s \sum_{k=1}^{i/s} k = \alpha \sum_{i=2}^n s \frac{i^2}{2s^2} = \frac{\alpha}{2s} \sum_{i=2}^n i^2$$

$$= \frac{\alpha n^3}{6s}$$