

# SCC120 Fundamentals of Computer Science

## Unit 7: Trees (Terminology)



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# Overview

- Concept of a Tree
- Examples of Trees
- Terminology/Definitions

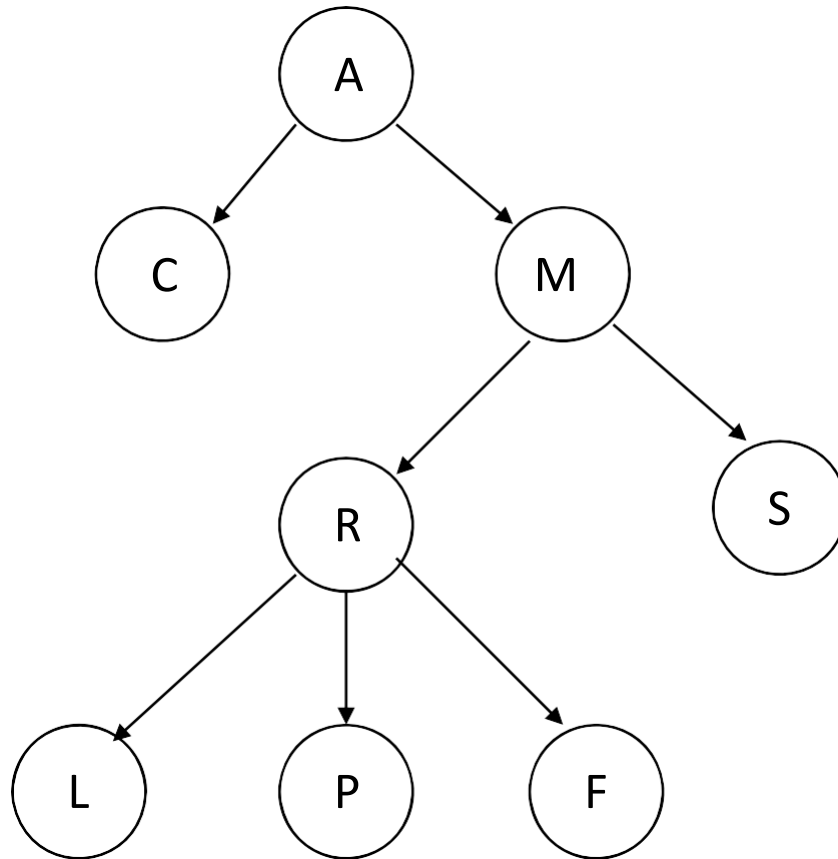
# A Tree

(we've seen this when we looked at graphs)

- A certain engineering company A is divided into a consultancy division C and a manufacturing division M
- The manufacturing division M is divided into a railway section R and a marine engine section S
- Section R is divided into three departments, building locomotives (L), passenger coaches (P), and freight wagons (F)

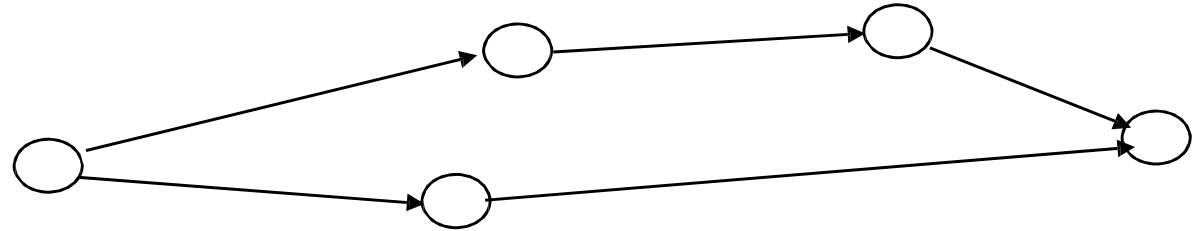


# A Tree

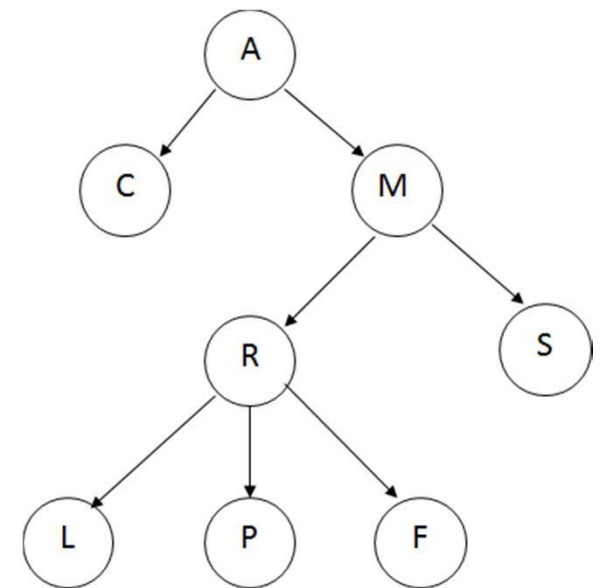


# Graphs vs. Trees

- A directed or undirected graph:
  - has a set of **nodes**
  - and a set of **edges** connecting these nodes
  - may have loops

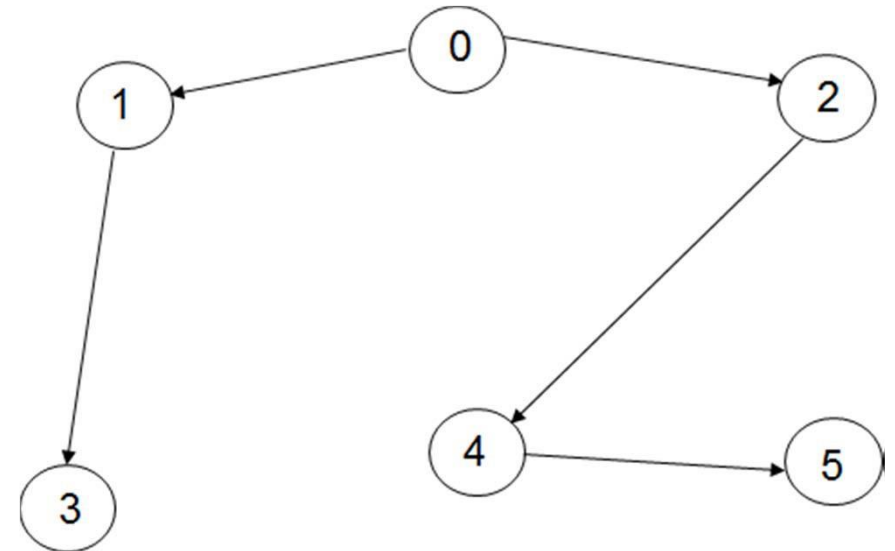
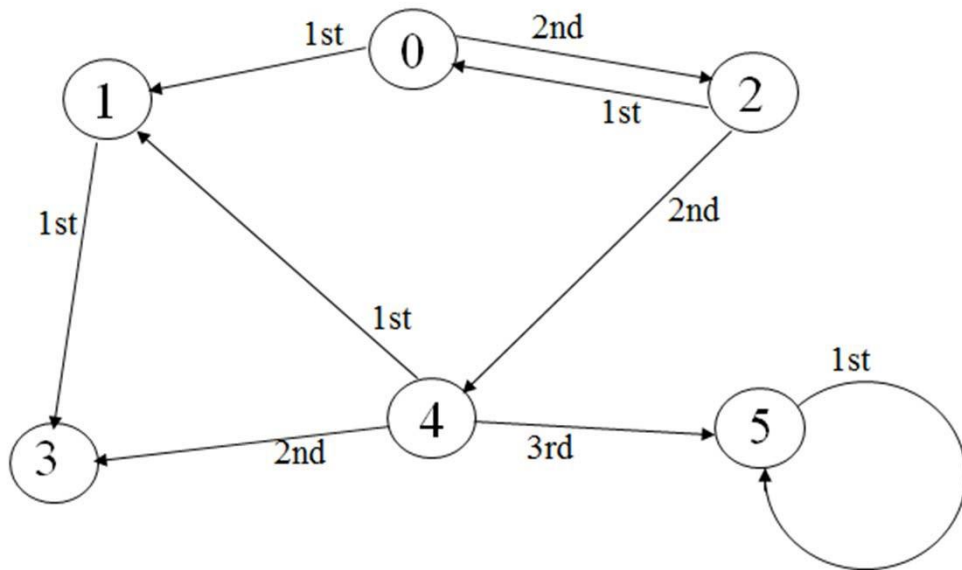


- A tree:
  - also has a set of nodes and edges
  - there are **no loops**
  - well-structured in terms of **levels** of a tree

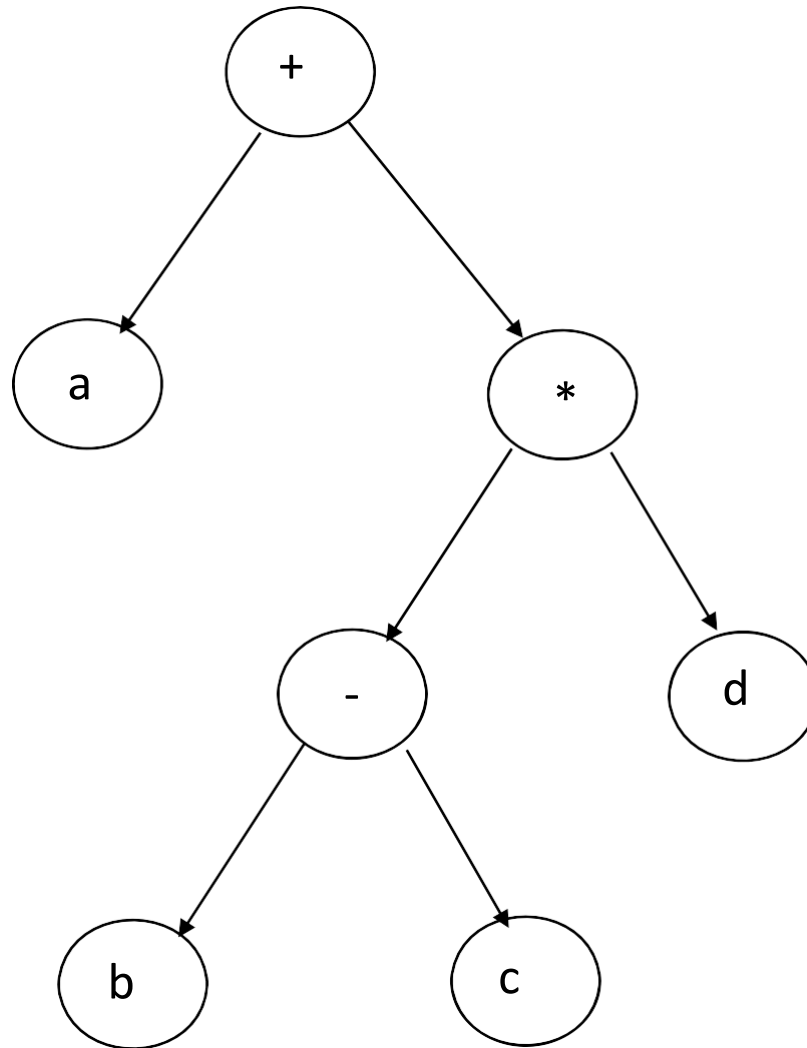


# Graphs vs. Trees

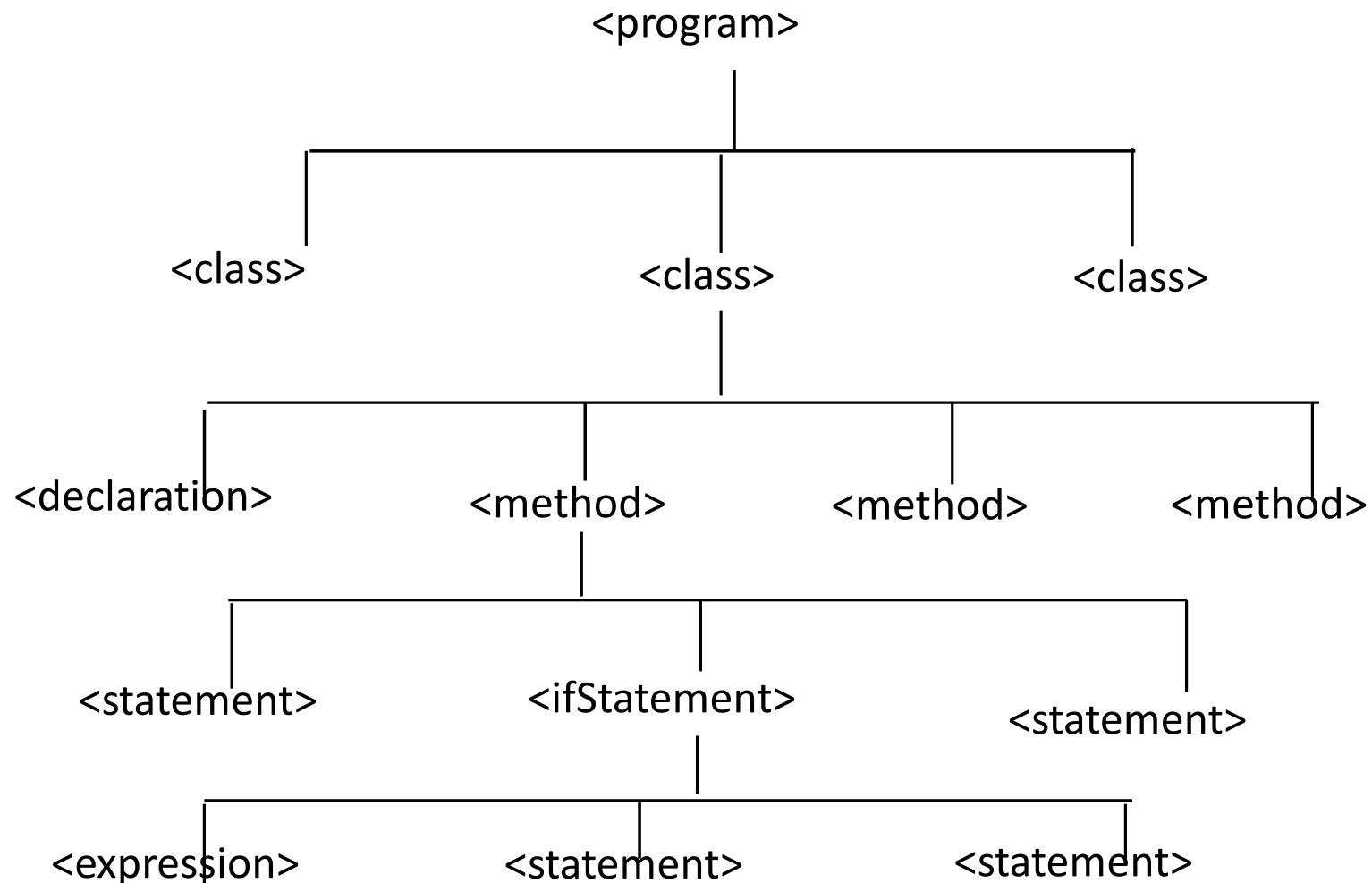
- If we take a connected graph and delete edges until no loops remain, but the structure is still connected, we obtain a **spanning tree** of the graph



# Examples of Trees: (Arithmetic) Expressions



# Examples of Trees: A Parse Tree



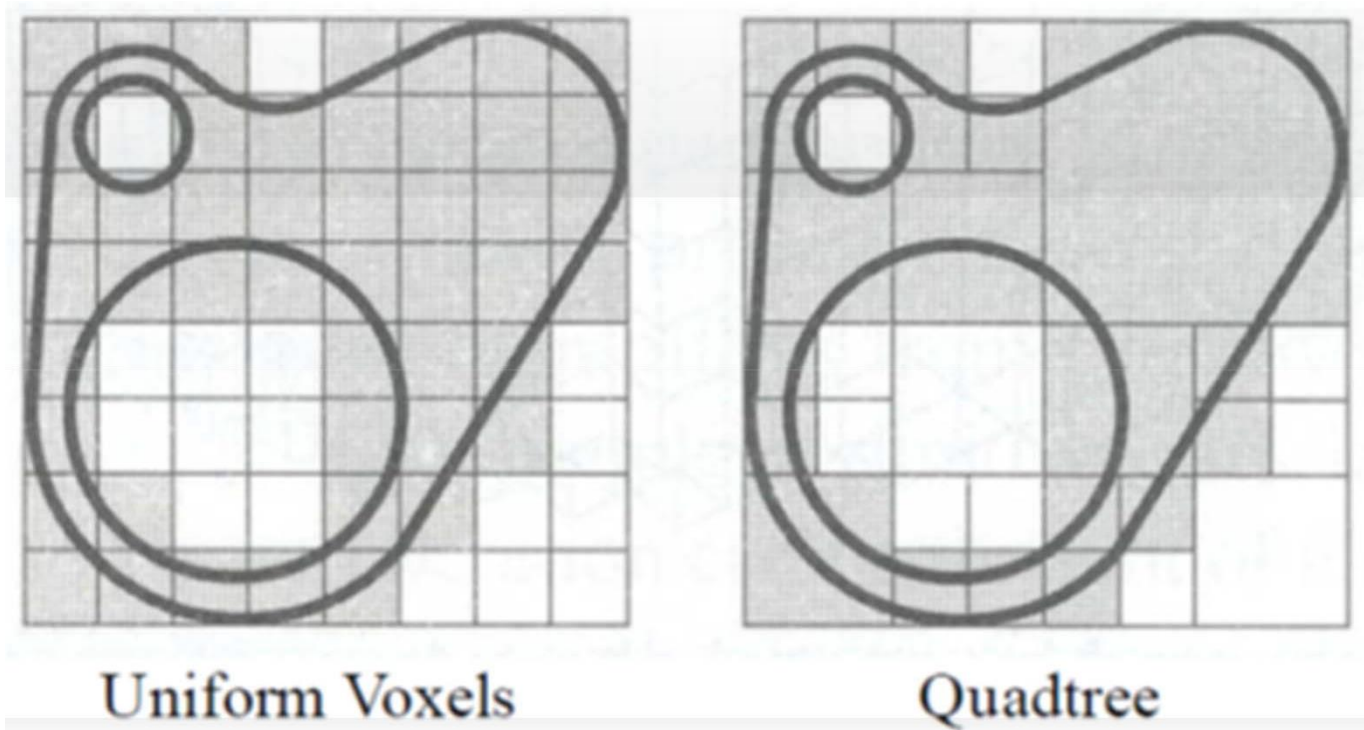


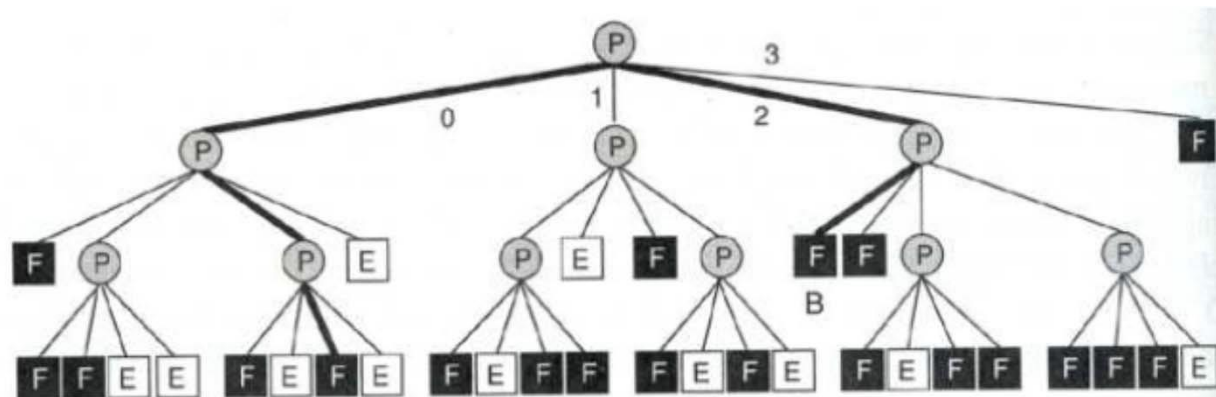
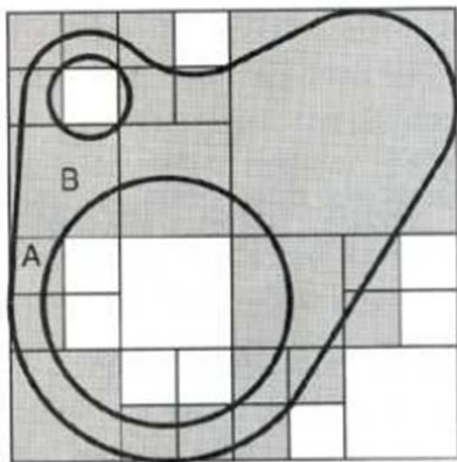
# Examples of Trees: A Document

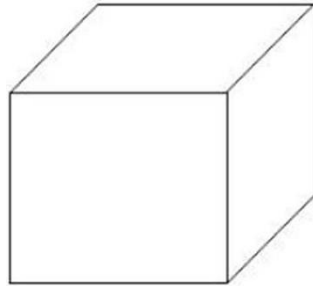
- A book, article, etc.
  - document = {frontMatter, body, backmatter}
  - body = {section, section, section ...}
  - section = {chapter, chapter, chapter, ... }
  - chapter = {title, para, para, para, ...}
  - para = {sentence, sentence, sentence, ...}
  - sentence = {word, word, word, ..., punctuation}



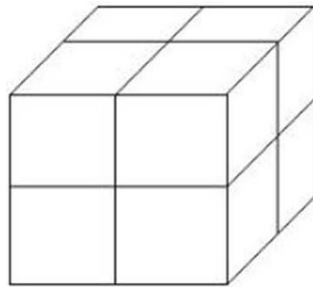
# Examples of Trees: QuadTrees



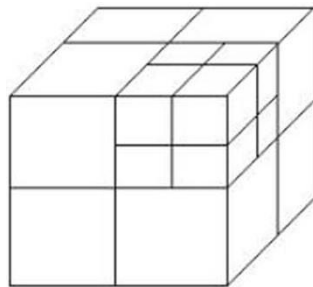
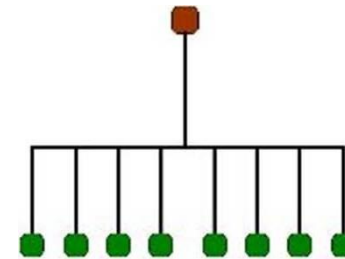




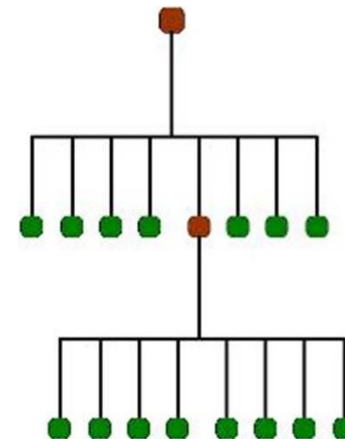
root node

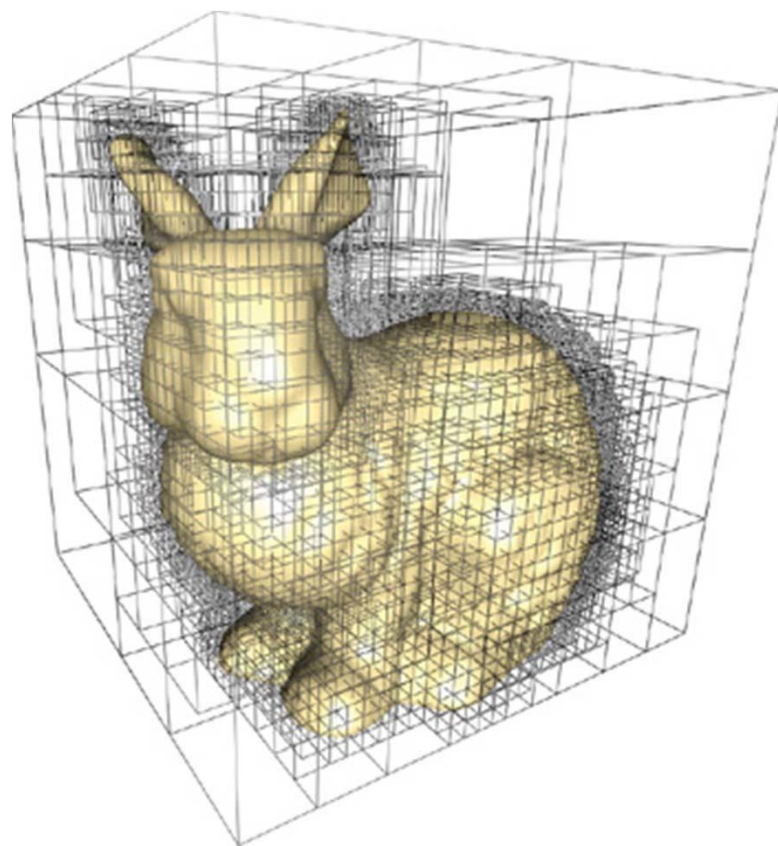
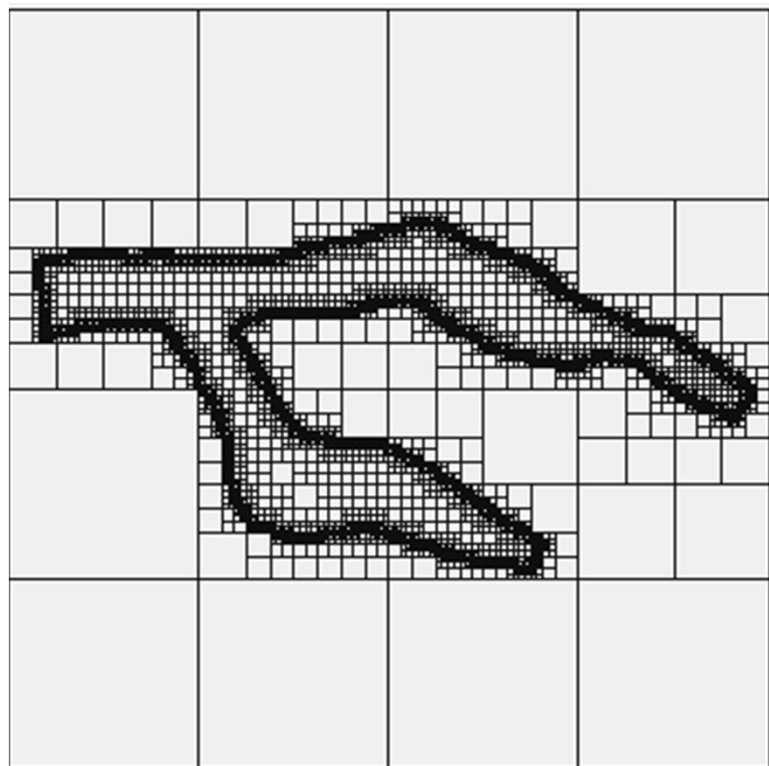


1 level



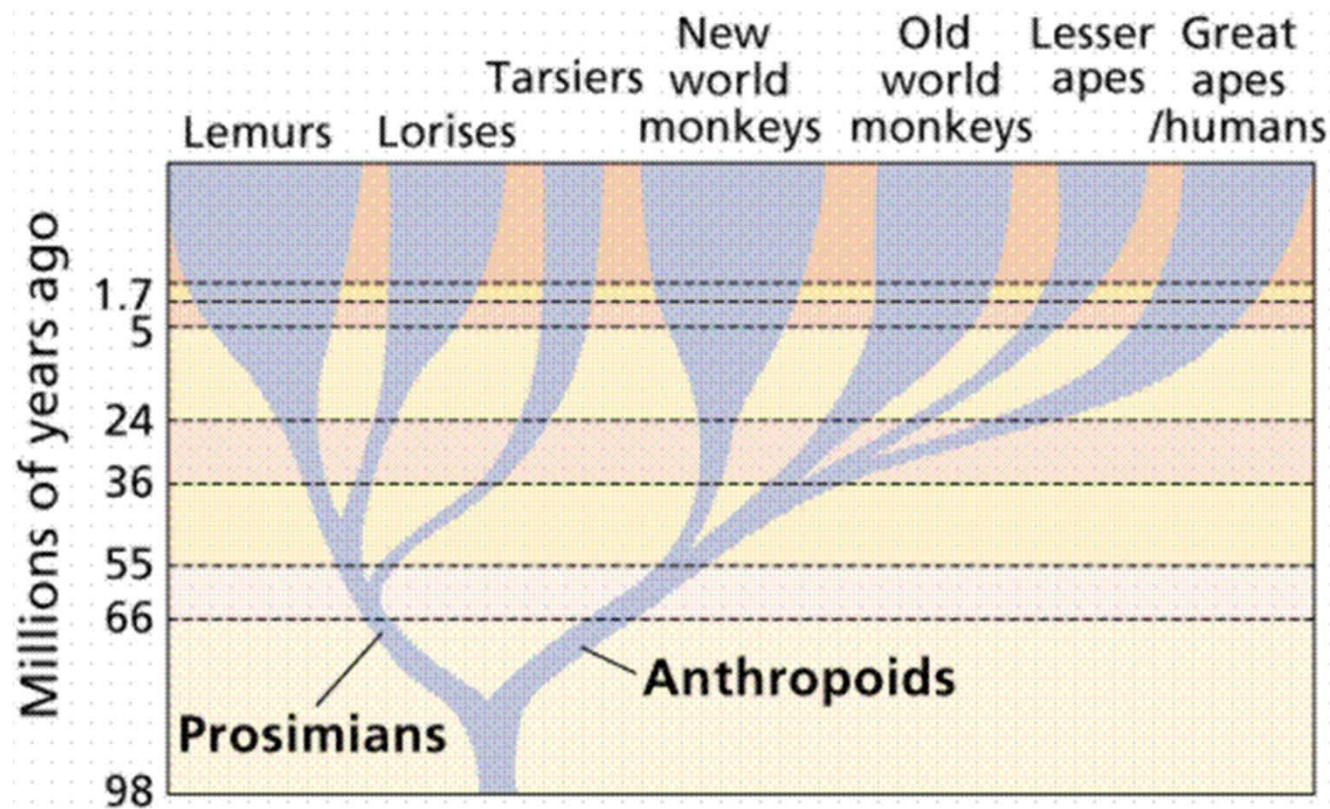
2 levels





# Examples of Trees: Evolutionary Tree

Example of an evolutionary tree (conventionally with the root at the bottom)



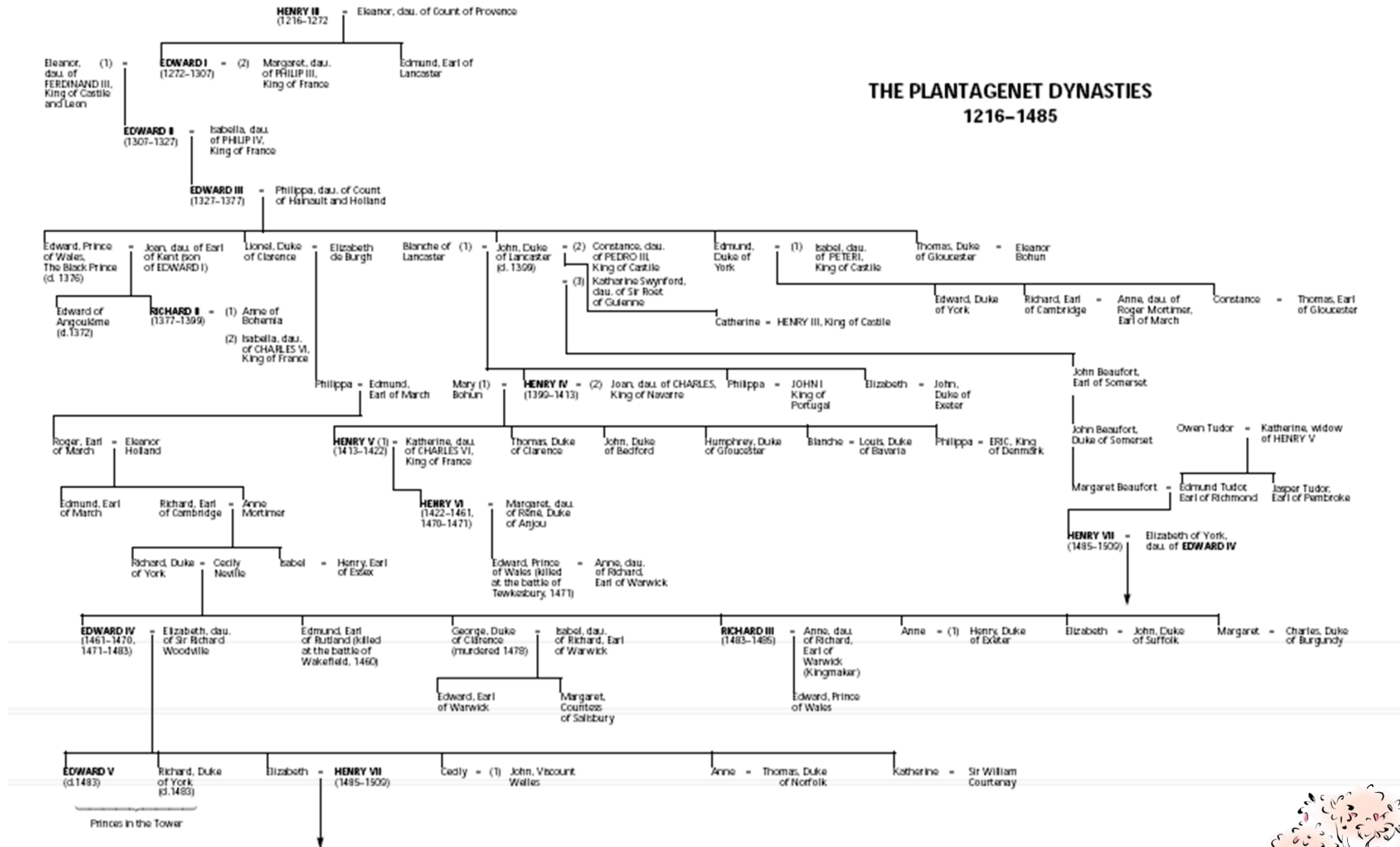


# Examples of Trees: Family Trees

- Simple family trees are trees in the Computer Science sense
  - But marriage of cousins, etc., makes it a graph
  - The line of descent of the English kings and queens is close to being a tree in most representations



# A Family Tree





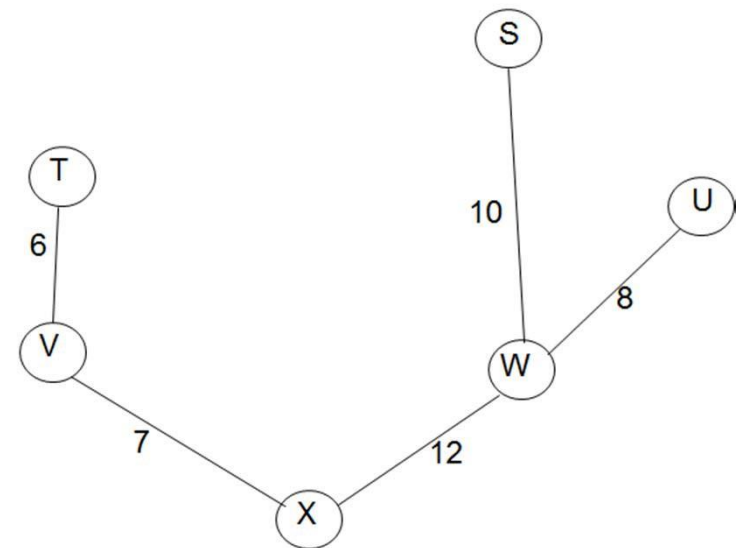
# Overview

- Concept of a Tree
- Examples of Trees
- Terminology/Definitions

# Non-Rooted vs. Rooted Trees

- A non-rooted tree is a connected non-directed graph without any loops

- Example: Minimal Spanning Tree



- A rooted tree has a “root” node, at the “top” level of the tree
  - We focus on rooted trees (i.e. “tree” means rooted tree)

# Definitions of Trees



- A tree is a node pointing to zero or more distinct trees
  - This definition is recursive
  - A tree contains smaller trees within it
- This definition also emphasizes the thought that a rooted tree represents a *hierarchy* of objects
  - Like the structure of a company

# Definitions of Trees

- A tree is a directed graph containing one node of in-degree 0, and zero or more other nodes of in-degree 1
  - This focuses on the nodes of a tree
  - The node of in-degree zero is the *root* node



# Empty Tree

- We can also define a tree which contains no nodes at all
- So we might want our definition on previous slide to start with “A tree is either empty with no nodes and no edges, or ...”

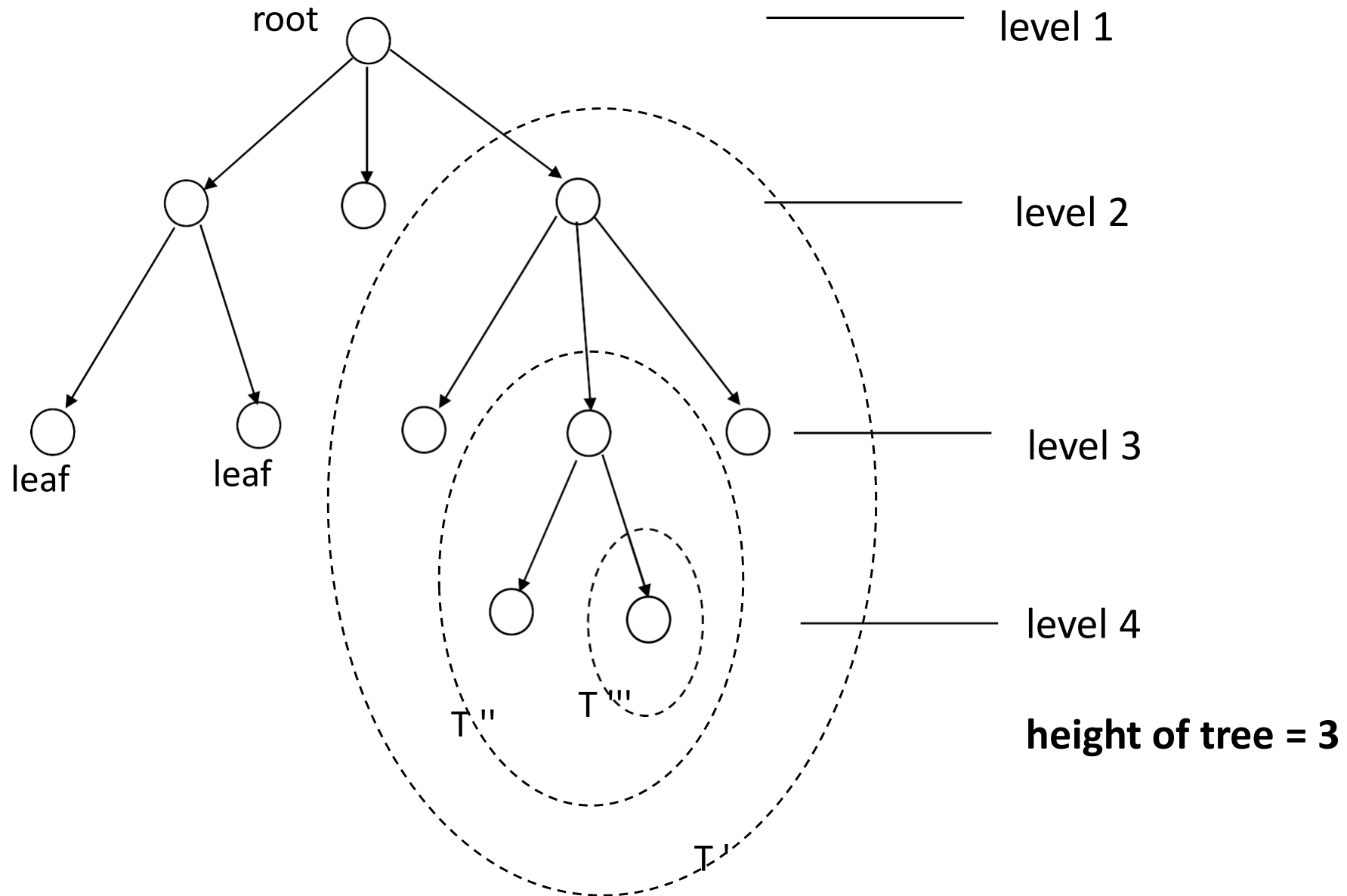


# Terminology

- A node of in-degree zero is called the *root* node of the tree
- Nodes of out-degree zero are called *leaf* nodes
- A node which is neither a root nor a leaf is called an *internal node*



# An Example Tree T



# Comments on the Tree T

- The *levels* of the nodes in the tree defined by  $1 +$  (the number of connections between the node and the root).
- The *height* of a **node** is the number of edges on the longest path between that node and a leaf.
- The *height* of the tree is the height of its root node.
- The *depth* of a node is the number of edges from the tree's root node to the node.



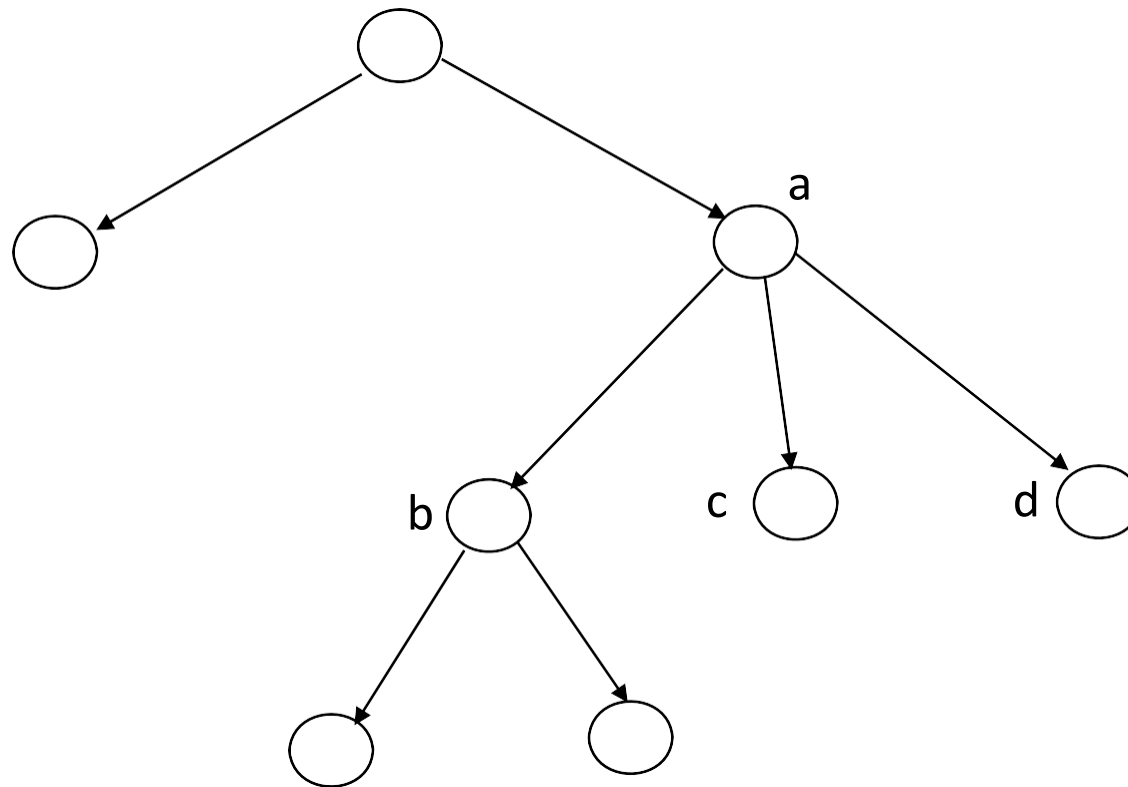


# Comments on the Tree T

- A *subtree* consists of a node X from/within a tree, together with all the nodes and edges below X (if any)
- X is the *root* of the subtree
- In our Tree T:
  - T' is a subtree of T
  - T'' is a subtree of T'
  - T''' is a subtree of T''

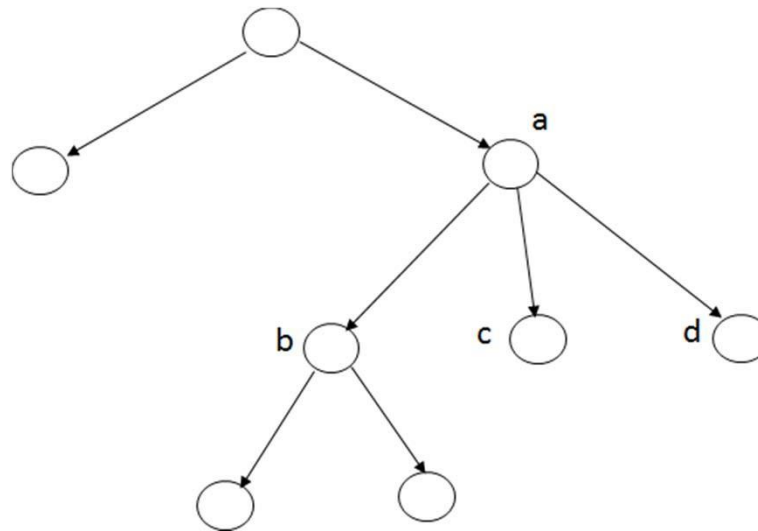


# An Alternative Metaphor

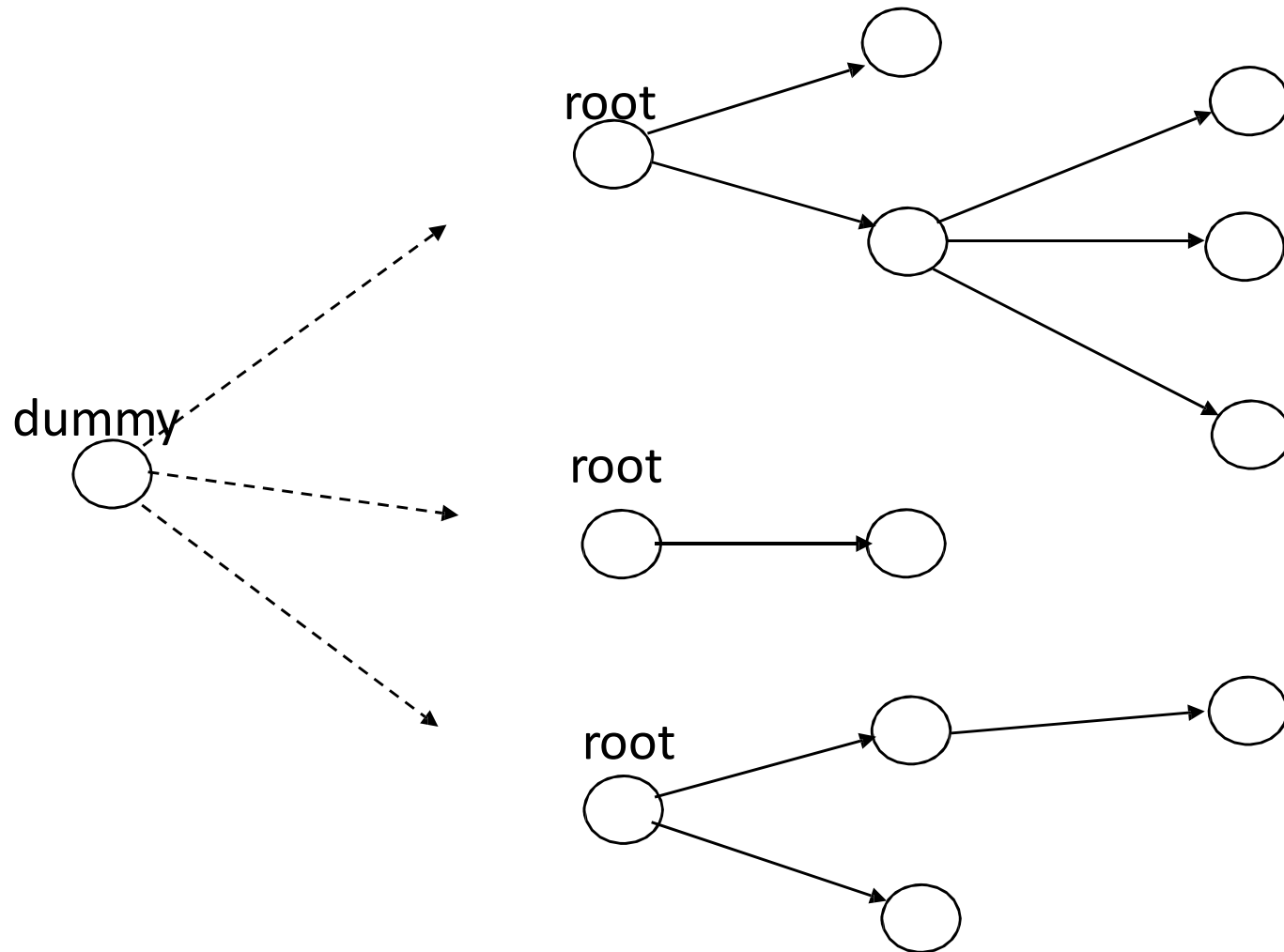


# An Alternative Metaphor

- Instead of metaphors based on real trees (e.g. root and leaf), we can use family metaphors
- Node “a” is the *parent* of nodes “b, c and d”
  - The out-degree of “a” is the number of children it has (here 3)
- Nodes “b, c and d” are the *children* of “a”
- Nodes “b, c and d” are *siblings* of each another



# A Forest of Three Trees



# Formulas for Trees

- For a singly-rooted tree  
num of edges = num of nodes - 1
- For multiple roots  
num of edges = num of nodes - num of roots
- As with a directed graph, the number of nodes in a tree may be called the *weight*



# Restricted Forms of Trees

- Trees may be restricted in various ways
  - We have already seen the restriction to a single root
- These restrictions are generally to make the trees easier to represent and manipulate by computer

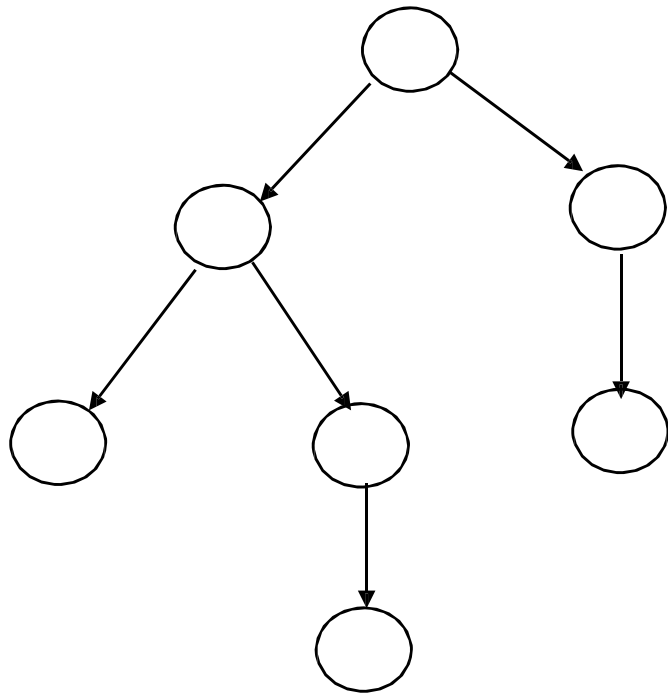


# Node Degree

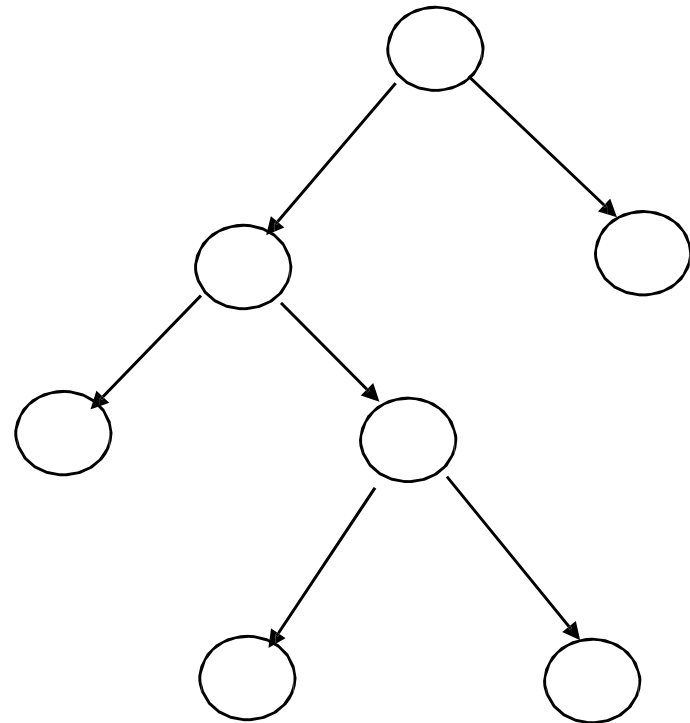
- A tree is said to have *limited out-degree*  $d$  if no node in the tree has an out-degree greater than  $d$
- A tree is said to have *strict out-degree*  $d$  if the out-degree of every node is either 0 or  $d$
- Specifically, for  $d=2$ :
  - In a *binary-limited tree* (or just *binary tree*), the nodes have a maximum out-degree of 2
  - In a *strict binary tree*, the nodes have an out-degree of 0 or 2 (so no out-degrees of 1)



# Node Degree



binary-limited



strict binary

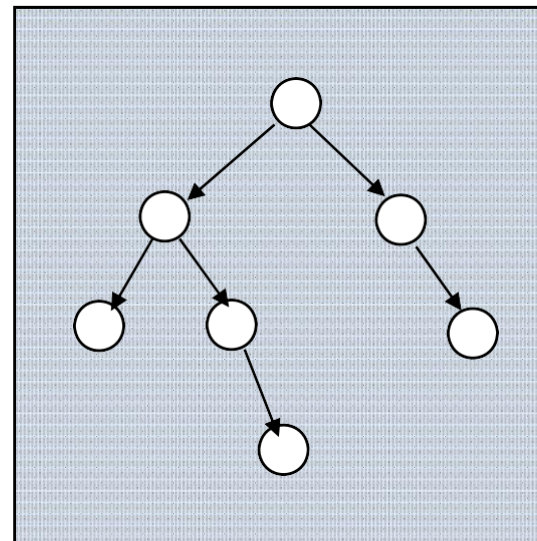
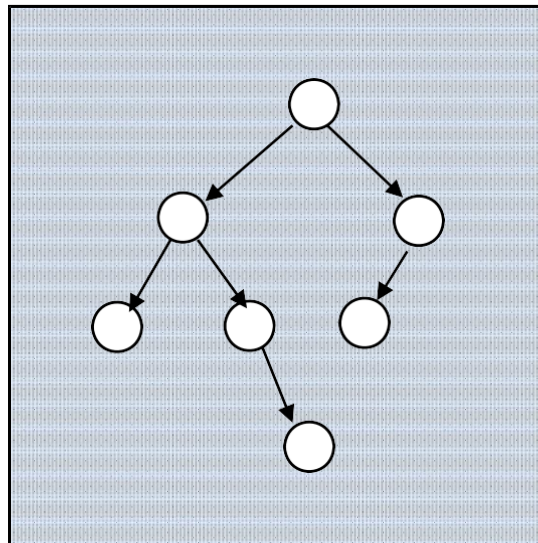
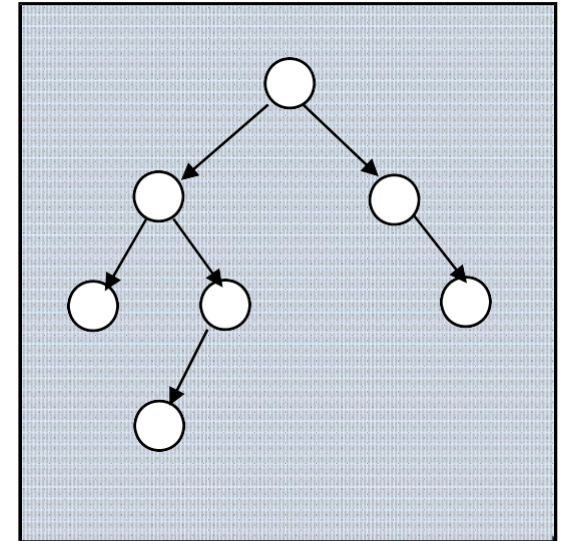
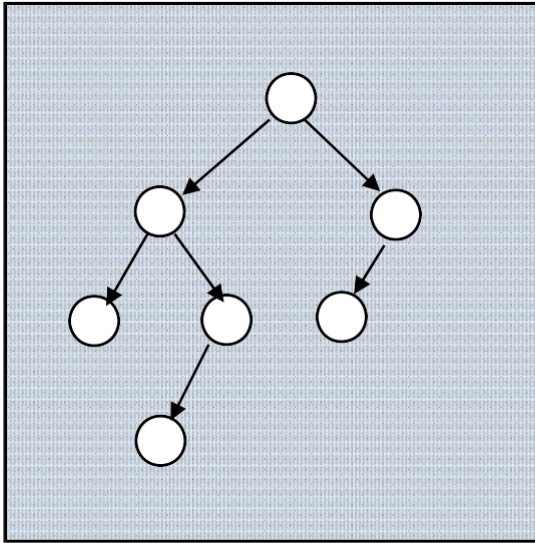


# Orientation

- In an *oriented* tree, we think of each node as having  $d$  distinct possible subnode positions
  - None, some or all of these positions may be occupied by a node
  - If only some of the node positions are occupied, the different ways in which subnode positions can be occupied or not are thought of as different trees
  - In other words, the subnodes of a node are *ordered*
- In an oriented *binary* tree, the *left* and *right* subnodes are distinct



# Four Different Oriented Binary Trees



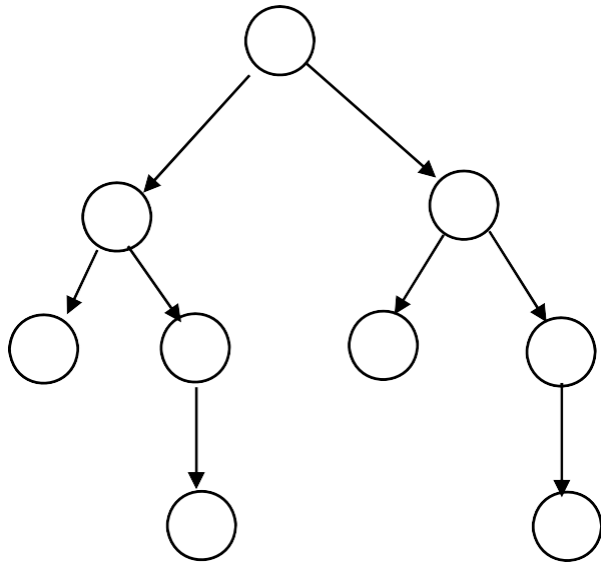
# Balance

- In a balanced tree, the “weights” (numbers of nodes) of all the subtrees of a node are as nearly equal as possible
- In practice, this means that the weights are either equal, or differ by not more than one
  - Other definitions are possible, but this is a reasonable one

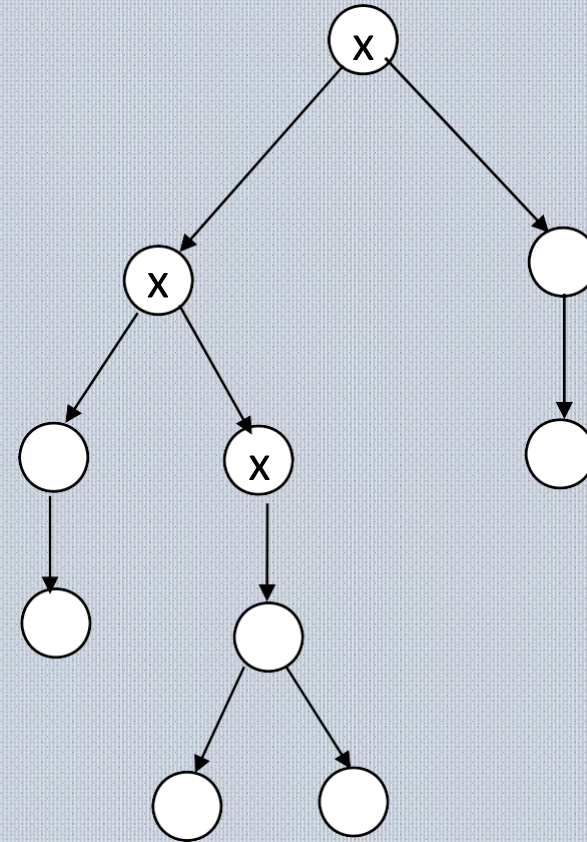




# Balance



balanced



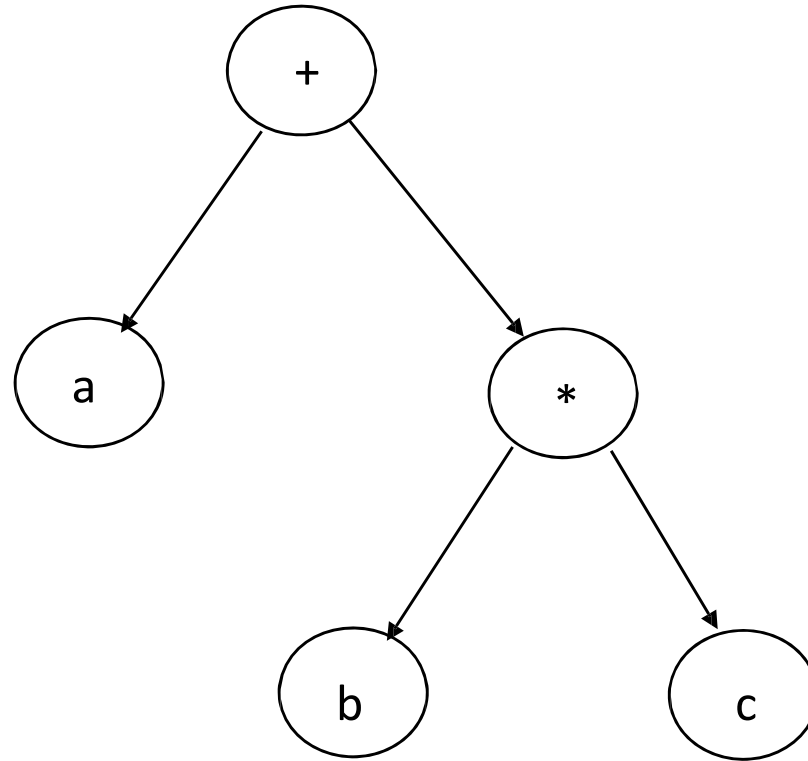
unbalanced at nodes  
marked "x"

# Value Distribution

- In most applications of trees, the nodes (or most of them) hold *values*, while the edges rarely do
- The values may be distributed in different ways
  - All the nodes hold values
  - All the nodes except the root hold values
  - Only the leaf nodes hold values
  - Leaves and non-leaves hold values, but of two distinct types



# Value Distribution



- This tree represents the expression:  $a + (b * c)$
- The leaves hold variables (or their values) and the non-leaves hold operators

# Other Restrictions

- The height of the tree (the number of edges on the longest path between root node and a leaf) may be limited
  - For example, maximum height = 5
- The leaves may all be at the same distance from the root: in this case, we say that the tree is *uniform*



# SCC120 ADT (weeks 7-13)

- Week 7      Abstractions; Set  
                 Stack
- Week 8      Queues Priority  
                 Queues
- Weeks      Graphs (Terminology)  
  9-11      Graphs (Traversals)  
            Graphs (Representations)
- Week 12    Trees (Terminology)
- Week 13