

SCC 120: Week 13 Workshop Problems

Question 1. The Fibonacci numbers for $n = 0, 1, \dots$ are defined as follows.

- $fib(0) = 0$;
- $fib(1) = 1$;
- $fib(n) = fib(n-1) + fib(n-2)$;

Consider the following piece of code that calculates the Fibonacci numbers.

```
int fib(int n) {  
    if (n is 0)  
        return 0  
    else if (n is 1)  
        return 1  
    else  
        return fib(n-1) + fib(n-2)  
}
```

What is the worst case complexity of the above code?


Question 2. Write a program to calculate the Fibonacci numbers that runs in linear time in the worst case.

Question 3. Let p and q be propositional variables and \wedge and \neg the usual logical operators. List all formulas entailed by p (that is, if p were true, what other formulas would be true?).

Question 4. What is an approach for determining if a formula is unsatisfiable, that is, inconsistent? What is its lower bound on its worst case complexity? Is this better or worse than satisfiability or validity of a formula?






Question 5. Consider a variant of the subset sum problem called *min-subset-sum*. For min-subset-sum, this is what you must do.

- If there is a subset that sums to 0, you must return a smallest such subset
- Else return *null*.

1. Is min-subset-sum a decision problem? 
2. For the same set can you get different answers?
3. Describe an approach for solving min-subset-sum?
4. Do you think it would be correct to claim that your approach is $\Omega(n)$? How about $\Omega(2^n)$?

Question 6. Use the undecidability of the *Halting* problem to show the undecidability of the *Dead Code* problem.

Question 7. Indicate the true statements

1. NP-Completeness problems are in a sense the hardest problems in NP 
2. If propositional satisfiability (SAT) has a polynomial time solution, then so does Hamiltonian Cycle (HC) 
3. If SAT has a polynomial time solution, then so does linear search 
4. If $P=NP$, then solving any problem is as easy as verifying it 
5. To show HC is NP-Complete, in addition to showing it is in NP, I need to give a reduction from HC to SAT 
6. If a problem is NP-Completeness, that is proof of the intractability of the problem 