

PROBLEMS FOR THE CATEGORY THEORY READING COURSE, 2017

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1. ASSIGNMENT 1, DUE END OF WEEK 4

1.1. Functors.

- (1) Leinster 1.2.21 (functors preserve isomorphisms)
- (2) Leinster 1.2.27, 1.2.28b (full, faithful)

1.2. Natural transformations.

- (1) In \mathbf{fdVec} , show that the functors id and $**$ are naturally isomorphic.
- (2) Show the the vertical composition of two natural transformations is in fact a natural transformation.
- (3) Prove carefully that the horizontal composition of two natural transformations is again a natural transformation.
- (4) Show that a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is part of an equivalence of categories if and only if it is *fully faithful* and *essentially surjective*. Clearly state where you are using the axiom of choice, or add hypotheses so it is unnecessary.

1.3. Universal properties.

- (1) Prove that two initial objects in a category are isomorphic.
- (2) For each of the following categories, decide whether there is an initial, final, and/or zero object, and if so, describe them: \mathbf{FinSet} , \mathbf{fdVec} , \mathbf{Top} , \mathbf{Top}_* (pointed topological spaces), field extensions of a fixed field F , \mathbf{Graphs} (your answer may depend on which class of graphs you consider), $\mathbf{Semigroups}$, \mathbf{Groups} .
- (3) Describe the product of two objects as the terminal object in some category.
- (4) Describe the tensor product of two vectors spaces as the initial object in some category.
- (5) Describe both the product and coproduct in the following categories: \mathbf{FinSet} , \mathbf{Top} , \mathbf{Top}_* , $\mathbf{AbGroup}$, \mathbf{Group} , \mathbf{Graphs} .

1.4. Adjunctions.

- (1) Consider the forgetful functor from abelian groups to groups. What is its left adjoint?
- (2) In the category of finite dimensional vector spaces, show that $- \otimes V$ is biadjoint to $- \otimes V^*$.
- (3) Prove that the ‘hom-set isomorphism’ and ‘unit/counit’ definitions of an adjunction are equivalent.

2. ASSIGNMENT 2, DUE AT THE END OF WEEK 7:

2.1. Idempotent completion. The ‘idempotent completion’ $\text{Kar}(C)$ (also called the ‘Karoubi envelope’) is defined as follows:

$$\begin{aligned}\text{Obj Kar}(C) &= \{(X \in \text{Obj}(C), p : X \rightarrow X) \mid p^2 = p\} \\ \text{Kar}(C)((X, p) \rightarrow (X', p')) &= \{f \in C(X \rightarrow X') \mid fp = f = p'f\}.\end{aligned}$$

- (1) Let primeVec denote the full subcategory of Vec consisting of vector spaces with prime dimensions. Show carefully that $\text{Kar}(\text{primeVec}) \cong \text{Vec}$.
- (2) Construct an equivalence $\iota_C : \text{Kar}(\text{Kar}(C)) \cong \text{Kar}(C)$.
- (3) Show that there is a fully faithful functor $C \rightarrow \text{Kar}(C)$ given by $X \mapsto (X, 1_X)$.
- (4) Note that given a functor $F : C \rightarrow \mathcal{D}$, there is functor $\text{Kar}(F) : \text{Kar}(C) \rightarrow \text{Kar}(\mathcal{D})$ given by

$$\begin{aligned}\text{Kar}(F)(X, p) &= (F(X), F(p)) \\ \text{Kar}(F)(f : (X, p) \rightarrow (X', p')) &= F(f).\end{aligned}$$

(You might say that Kar is a 2-functor from CAT to itself – what does this mean at the level of natural transformations?) Given a functor $F : \text{Kar}(C) \rightarrow \text{Kar}(D)$, show that it is determined up to natural isomorphism by its restriction $F|_C : C \rightarrow \text{Kar}(D)$, by showing $F \cong (\iota_D \circ \text{Kar}(F|_C))$.

- (5) (Not on the problem set: there is a forgetful 2-functor $\text{CAT} \rightarrow \text{SemiCat}$, the 2-category of ‘semicategories’ (categories without identities) and their functors and natural transformations. The idempotent completion gives a 2-functor $\text{SemiCat} \rightarrow \text{CAT}$. Are they an adjoint pair?)

2.2. Limits.

- (1) Recall that given a functor $F : \mathcal{J} \rightarrow C$, the *limit* of F , written $\lim_{\mathcal{J}} F$ is a terminal object in the category of cones over F . Explicitly, a cone consists of
 - (a) an object $X \in \text{Obj } C$,
 - (b) for each $j \in \text{Obj } \mathcal{J}$, a map $f_j : X \rightarrow F(j)$,
 - (c) such that for any $g : j \rightarrow j'$, $F(g) \circ f_j = f_{j'}$.

Another way of saying this is that the category of cones is the comma category $(J \downarrow F)$, where here we interpret \mathcal{J} as a functor $C \rightarrow \text{Fun}(\mathcal{J} \rightarrow C)$ by $c \mapsto (j \mapsto c)$ and we interpret F as a functor $1 \rightarrow \text{Fun}(\mathcal{J} \rightarrow C)$ by $1 \mapsto (j \mapsto F(j))$. Explain carefully why these are talking about the same thing!

2.3. The Yoneda embedding.

- (1) Prove the Yoneda lemma.
- (2) Explain why $\text{Fun}(C^{\text{op}} \rightarrow \mathcal{D}) = \text{Fun}(C \rightarrow \mathcal{D}^{\text{op}})$. Prove that $\text{Set}^{\text{op}} \not\cong \text{Set}$, but that $\text{fdVec}^{\text{op}} \cong \text{fdVec}$.

3. ASSIGNMENT 3, DUE MAY 12:

- (1) Prove that every monoidal category is monoidally equivalent to a strict monoidal category.
 (Hint: given a monoidal category C , define a new monoidal category $\text{List}C$, whose objects are lists of objects in C . In $\text{List}C$, tensor product of objects is concatenation of lists, and the tensor unit is the empty list. There should be a functor $\text{List}C \rightarrow C$ defined by sending a list to the tensor product of the elements of that list. Your job is to describe what happens at the level of morphisms, and check that everything works.)
- (2) Let C be a monoidal category. We say a ‘monoid object’ in C (or, as we gain confidence, just a monoid in C) is a tuple $(A \in \text{Obj } C, \iota : 1 \rightarrow A, m : A \otimes A \rightarrow A)$ satisfying some conditions. Look up, or work out, what these conditions should be. (Hint: look at chapter 7 of Etingof, Gelaki, Nikshych, and Ostrik’s book [Tensor Categories](#).) You should be able to show that a monoid object in Vec is what is usually called an associative unital algebra.
 - (a) A ‘module object’ for a monoid object $A \in C$ is a tuple $(M, \triangleright : A \otimes M \rightarrow M)$ satisfying an appropriate condition (what is it?). A morphism f between module objects M and M' is a morphism between the underlying objects, such that $f \circ \triangleright_M = \triangleright_{M'} \circ (1_A \otimes f)$. Draw the string diagram corresponding to this axiom. Define composition of module morphisms, by imitating the definition for modules over a ring. Show that modules for a fixed monoid object form a category.
- (3) Show that $\text{Rep}G$, for G a finite group, forms a monoidal category.
- (4) If $\text{Rep}G \cong \text{Rep}H$, as categories, are G and H isomorphic? (Hint: no, give a counterexample — any pair of non-isomorphic groups with the same number of irreducible representations (equivalently, the same number of conjugacy classes) will do.) What about if $\text{Rep}G \cong \text{Rep}H$ as monoidal categories, and moreover this equivalence is compatible with the forgetful functors to Vec ? (Hint: think about the monoidal automorphisms of the forgetful functor.)

4. ASSIGNMENT 4, DUE MAY 26:

4.1. Abelian categories.

- (1) Show if $F : C \rightarrow D$ is a fully faithful functor, then $f \in C(X \rightarrow Y)$ is a monomorphism if and only if $F(f)$ is. (Similarly for epimorphisms.)
- (2) Use this, and the Yoneda embedding, to show that in a rigid abelian category, the functor $- \otimes X$ is exact. (See Proposition 2.1.8 of Bakalov-Kirillov, attached, if you need some help; they don't explain how they are using Yoneda, however! You may cite the literature for any lemmas you like.)

4.2. Braided monoidal categories.

- (1) Explain how any object X in a pivotal braided category gives an oriented link invariant.
- (2) Describe how the Temperley-Lieb category has the structure of a pivotal braided category.
- (3) Calculate the invariant of the trefoil corresponding to the object 1 in Temperley-Lieb. (What is this invariant usually called?)
- (4) Show that in $\text{Kar}(TL)$, we have $(2, 1_2) \cong (2, f^{(2)}) \oplus (0, 1_0)$. (Here $f^{(2)}$ denotes the second Jones-Wenzl idempotent.)
- (5) Calculate the invariant of the unknot corresponding to the object $(2, f^{(2)})$ in (the idempotent completion of) Temperley-Lieb.
- (6) We say a monoid (A, m) in a braided monoidal category is commutative if $m \circ \beta = m$. Define a monoidal structure on the category of modules for a commutative monoid.

LEMMA 2.1.6. *Suppose that V has a dual V^* . Then there exist canonical isomorphisms*

$$(2.1.13) \quad \text{Hom}(U \otimes V, W) = \text{Hom}(U, W \otimes V^*),$$

$$(2.1.14) \quad \text{Hom}(U, V \otimes W) = \text{Hom}(V^* \otimes U, W).$$

PROOF. To $\psi \in \text{Hom}(U \otimes V, W)$ we associate the composition

$$U \xrightarrow{\text{id} \otimes i_V} U \otimes V \otimes V^* \xrightarrow{\psi \otimes \text{id}} W \otimes V^*$$

which is an element of $\text{Hom}(U, W \otimes V^*)$.

Similarly, to $\varphi \in \text{Hom}(U, W \otimes V^*)$ we assign

$$U \otimes V \xrightarrow{\varphi \otimes \text{id}} W \otimes V^* \otimes V \xrightarrow{\text{id} \otimes e_V} W.$$

One can easily check that these two maps are inverse to each other, establishing (2.1.13). The proof of (2.1.14) is similar. \square

In particular, if both U and V have duals, then by Lemma 2.1.6

$$(2.1.15) \quad \text{Hom}(U, V) = \text{Hom}(V^*, U^*) = \text{Hom}(\mathbf{1}, V \otimes U^*).$$

(In the language of abstract nonsense, this means that the category \mathcal{C} has internal Hom's when it has duals.) For $f \in \text{Hom}(U, V)$ its image in $\text{Hom}(V^*, U^*)$ via the isomorphism (2.1.15) will be denoted by f^* .

If the right dual $*$ exists for all objects in \mathcal{C} , then by (2.1.15) it is a contravariant functor, i.e., a functor $\mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$ where \mathcal{C}^{op} is the opposite (or dual) category to \mathcal{C} . (Recall that \mathcal{C}^{op} has the same objects as \mathcal{C} but with all arrows reversed.)

EXERCISE 2.1.7. Show that, in a rigid category, $*$ is an equivalence of categories $\mathcal{C} \rightarrow \mathcal{C}^{\text{op}}$.

Rigidity is a very restrictive requirement. As an illustration, let us prove the following proposition.

PROPOSITION 2.1.8. *In an abelian rigid monoidal category the tensor product functor \otimes is exact, i.e., for any short exact sequence $0 \rightarrow U \rightarrow V \rightarrow W \rightarrow 0$ and an object X , the sequences*

$$0 \rightarrow U \otimes X \rightarrow V \otimes X \rightarrow W \otimes X \rightarrow 0$$

and

$$0 \rightarrow X \otimes U \rightarrow X \otimes V \rightarrow X \otimes W \rightarrow 0$$

are exact.

PROOF. The sequence

$$0 \rightarrow U \otimes X \rightarrow V \otimes X \rightarrow W \otimes X$$

is exact iff

$$0 \rightarrow \text{Hom}(Y, U \otimes X) \rightarrow \text{Hom}(Y, V \otimes X) \rightarrow \text{Hom}(Y, W \otimes X)$$

is exact for any object Y . But by (2.1.12, 2.1.13), $\text{Hom}(Y, U \otimes X) = \text{Hom}(Y \otimes {}^*X, U)$. Since the functor $\text{Hom}(Y, -)$ is left exact, it follows that $- \otimes X$ is left exact. Using Exercise 2.1.7 (or repeating the same argument with duals), we see that it is also right exact. \square

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