Five questions, marked out of a total of $(5+15)+10+(5+10)+10+(10+10) = \mathbf{75}$ **marks**.

Some hints and tips:

- Read the assignment early, so you have time to think about the problems!
- Read **all** the problems and make sure you understand them: sometimes problems can be difficult to understand, but easy once understood.
- You may always assume the solution to a previous part when solving a subsequent part. For instance, if part (a) says provide an algorithm that does X in $O(N)$ time, then in part (b) you may assume that there is some algorithm that does X in $O(N)$ time, and treat it as a black box, *even if you don't know how to solve part (a)*! This means you can skip part (a) and still obtain full credit on part (b), if your solution for (b) is correct.
- Concise answers written in clear English text are perfectly acceptable, as are answers written in *clear* pseudocode.

1. (a) [**5 marks**] *Revision:* Describe how to multiply two n-degree polynomials together in $O(n \log n)$ time, using the Fast Fourier Transform (FFT). You do not need to explain how FFT works – you may treat it as a black box.
Hint: Remember that FFT does not multiply polynomials, what does it do? Refer to the lecture slides.

   (b) In this part we will extend the Fast Fourier Transform (FFT) algorithm described in class to multiply multiple polynomials together (not just two). Suppose you have K polynomials $P_1, \ldots, P_K$ so that

$$\text{degree}(P_1) + \cdots + \text{degree}(P_K) = S$$

   (i) [**5 marks**] Show that you can find the product of these K polynomials in $O(KS \log S)$ time.
Hint: How many points do you need to uniquely determine an S-degree polynomial?

   (ii) [**10 marks**] Show that you can find the product of these K polynomials in $O(S \log S \log K)$ time. Explain why your algorithm has the required time complexity.
Hint: consider using divide-and-conquer!

2. [**10 marks**] You have a set of N coins in a bag, each having a value between 1 and M, where $M \geq N$. Some coins may have the same value. You pick two coins (**without replacement**) and record the sum of their values. Determine what possible sums can be achieved, in $O(M \log M)$ time.

   For example, if there are $N = 3$ coins in the bag with values 1, 4 and 5 (so we could have $M = 5$), then the possible sums are 5, 6 and 9.

   Hint: if the coins have values $v_1, \ldots, v_N$, how might you use the polynomial $x^{v_1} + \cdots + x^{v_N}$? 两个多项式平方 取系数大于1的个数 MlogM

3. Let us define the Fibonacci numbers as $F_0 = 0$, $F_1 = 1$ and $F_n = F_{n-1} + F_{n-2}$ for all $n \geq 2$. Thus, the Fibonacci sequence looks as follows: 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

   (a) [**5 marks**] Show, by induction or otherwise, that

   $$\begin{pmatrix} F_{n+1} & F_n \\ F_n & F_{n-1} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^n$$

   for all integers $n \geq 1$.

   (b) [**10 marks**] Hence or otherwise, give an algorithm that finds F_n in $O(\log n)$ time. *Hint: You may wish to refer to Example 1.5 on page 28 of the "Review Material" found as lecture notes for Topic 0 on the Course Resources page of the course website.*

4. [**10 marks**] You have N items for sale, numbered from 1 to N. Alice is willing to pay $a[i] > 0$ dollars for item i, and Bob is willing to pay $b[i] > 0$ dollars for item i. Alice is willing to buy no more than A of your items, and Bob is willing to buy no more than B of your items. Additionally, you must sell each item to either Alice or Bob, but not both, so you may assume that $N \leq A + B$. Given $N, A, B, a[1..N]$ and $b[1..N]$, show that you can determine the **maximum** total amount of money you can earn in $O(N \log N)$ time.

   Hint: Suppose $N = A + B$ and pretend you wish to sell all items to Alice, but must choose B of them to give to Bob instead. Which ones do you want to give to Bob first? Then extend your approach to handle $N < A + B$ as well.

5. Your army consists of a line of N giants, each with a certain height. You must designate precisely $L \leq N$ of them to be leaders. Leaders must be spaced out across the line; specifically, every pair of leaders must have at least $K \geq 0$ giants standing in between them. Given N, L, K and the heights $H[1..N]$ of the giants in the order that they stand in the line as input, find the *maximum* height of the *shortest* leader among all valid choices of L leaders. We call this the *optimisation* version of the problem.

   For instance, suppose $N = 10, L = 3, K = 2$ and $H = [1, 10, 4, 2, 3, 7, 12, 8, 7, 2]$. Then among the 10 giants, you must choose 3 leaders so that each pair of leaders has at least 2 giants standing in between them. The best choice of leaders has heights 10, 7 and 7, with the shortest leader having height 7. This is the best possible for this case.

   (a) [**10 marks**] In the *decision* version of this problem, we are given an additional integer T as input. Our task is to decide if there exists some valid choice of leaders satisfying the constraints whose shortest leader has height no less than T.

   Give an algorithm that solves the decision version of this problem in $O(N)$ time.

(b) **[10 marks]** Hence, show that you can solve the optimisation version of this problem in $O(N \log N)$ time.