

SOLUTION CSI2110

Université d'Ottawa
Faculté de génie

École de science informatique
et de génie électrique



uOttawa

University of Ottawa
Faculty of Engineering

School of Electrical Engineering
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L'Université canadienne
Canada's university

CSI2110

Data Structures and Algorithms

Midterm Examination

Length of Examination: 2 hours

Professor: L. Moura

Oct 19, 2014, 15:00

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You'll get marks for your best 50 points out of the 60 points available.
Feel free to answer all 60 points.

Last name:

First name:

Student number:

Signature:

Closed Book.

Please answer in the space provided (in this questionnaire).

No calculators or electronic devices are allowed.

At the end of the exam, when time is up:

- Stop working and turn your exam upside down.
- Remain silent.
- Do not move or speak until *all* exams have been picked up, and a TA or a Professor gives the go-ahead.

Page	Marks of each page
PAGE 2	out of 5
PAGE 3	out of 7
PAGE 4	out of 12
PAGE 5	out of 5
PAGE 6	out of 11
PAGE 7	out of 6
PAGE 8	out of 2
PAGE 9	out of 4
PAGE 10	out of 8
TOTAL	out of 60

NOTE: In all questions where a big-Oh characterization is asked, please give the best possible one.

Question 1 [2 points] What is the worst-case running time of the following algorithms (in big-Oh notation) for input being a bi-dimensional $n \times n$ array A , for an arbitrary integer $n \geq 200$?

```
Algorithm ALG1( $A, n$ ) {
    tot  $\leftarrow 1$ ;
    for  $i \leftarrow 0$  to  $n - 1$  do
        for  $j \leftarrow 0$  to  $100$  do
            for  $k \leftarrow 0$  to  $j * j$  do
                tot  $\leftarrow$  tot +  $A[i][j] * k$ ;
    return tot;
}
```

- a) $O(1)$ b) $O(n)$ c) $O(n^2)$ d) $O(n^3)$ e) $O(n^4)$

REASON :

Loop 1 runs n times
 Loop 2 runs 100 times
 Loop 3 runs ≤ 10000 times

```
Algorithm ALG2( $A, n$ ) {
    tot  $\leftarrow 1$ ;
    for  $i \leftarrow 0$  to  $n - 1$  do
        for  $j \leftarrow 0$  to  $100$  do
            for  $k \leftarrow 0$  to  $i * i$  do
                tot  $\leftarrow$  tot +  $A[i][j] * k$ ;
    return tot;
}
```

- a) $O(1)$ b) $O(n)$ c) $O(n^2)$ d) $O(n^3)$ e) $O(n^4)$

REASON :

Loop 1 runs n times
 Loop 2 runs 100 times
 Loop 3 runs $< n^2$ times

Question 2 [3 points] Fill the blanks below:

- $7\sqrt{n} + 10 \log_2 n^2$ is $O(\sqrt{n})$
- $(\sum_{i=0}^{100} (i \cdot n^3))$ is $O(n^3)$
- $n^4 + 10n^2$ is $O(n^4)$

Question 3 [4 points]

Consider the following part of a recursive program in Java describing the following method of a general tree (not necessarily binary) that stores integer numbers:

```
public int Calc(Position<int> p) {  
    int a=p.getElement();  
    for(Position<int> c: children(p))  
        a = a + Calc(c);  
    return a;  
}
```

- A. What is the result of `Calc(v)`, where v is the root of the tree?

THE SUM OF ALL ELEMENTS STORED IN
THE NODES OF THE TREE.

- B. What is the worst-case running time in big-Oh notation of `Calc(v)`, where v is the root of a tree with n elements?

$\Theta(n)$ (EACH NODE IS VISITED EXACTLY ONCE)

Question 4 [3 points] Prove that $f(n) = n^2 + \log_2 n^4$ is $O(n^2)$ by using the definition of big-Oh.

$$\begin{aligned} f(n) &= n^2 + \log_2 n^4 = n^2 + 4 \log_2 n \\ &\leq n^2 + 4 n^2 = 5 n^2, \\ &\text{for all } n \geq 1. \end{aligned}$$

Question 5 [3 points] Complete the following sentences with the tightest big-Oh estimate for the worst-case running time of the corresponding algorithms. The number of elements in each data structure is an arbitrary number n .

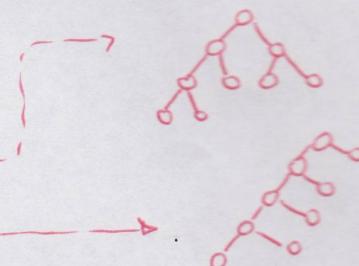
1. Method `insert` of a heap: $O(\log n)$
2. Method `push` of a Stack implemented with an array: $O(1)$
3. Method `push` of a Stack implemented with an extendible array: $O(n)$
4. Method `enqueue` of a Queue implemented with a doubly linked list: $O(1)$
5. Method `removelast` of a Deque (double-ended queue) implemented with a doubly linked list: $O(1)$
6. Algorithm `heapSort` on an array of n elements: $O(n \log n)$

Question 6 [5 points] Short answers:

- The array that represents a **min-heap** is always sorted. [TRUE/FALSE]
- The height of a **full** binary tree of n nodes is $O(\log n)$. [TRUE/FALSE]
- The height of a **complete** binary tree of n nodes is $O(\log n)$. [TRUE/FALSE]
- The **maximum** number of nodes in a binary tree of height h is $2^{h+1} - 1$
- The **minimum** number of nodes in a binary tree of height h is $h + 1$

Question 7 [2 points] Let h_T be the height of a **full binary tree** T with 9 nodes. Let i and e be its number of internal and external nodes, respectively.

- The number of internal nodes of T is $i = 4$
- The number of external nodes of T is $e = 5$
- The minimum possible value for h_T is 3
- The maximum possible value for h_T is 4



Question 8 [2 points] Which of the following data structures can **always** be searched for its minimum element in $O(1)$ time (multiple answers are possible)?

- unsorted array
- sorted array
- sorted doubly-linked list
- min-heap

Question 9 [2 points] Consider a Queue implemented using a circular array, with the contents given below:

Front = 0

Rear = 4

index i	0	1	2	3	4
value $Q[i]$	Z	T	X	Y	

Show the data structure after the following operations executed in the given order:

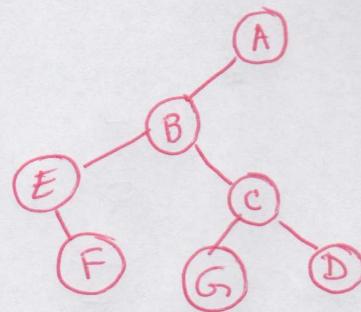
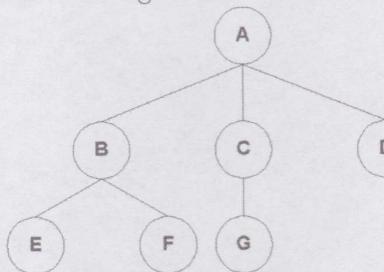
Dequeue(); Enqueue(B); Dequeue(); Enqueue(A);

Front = 2

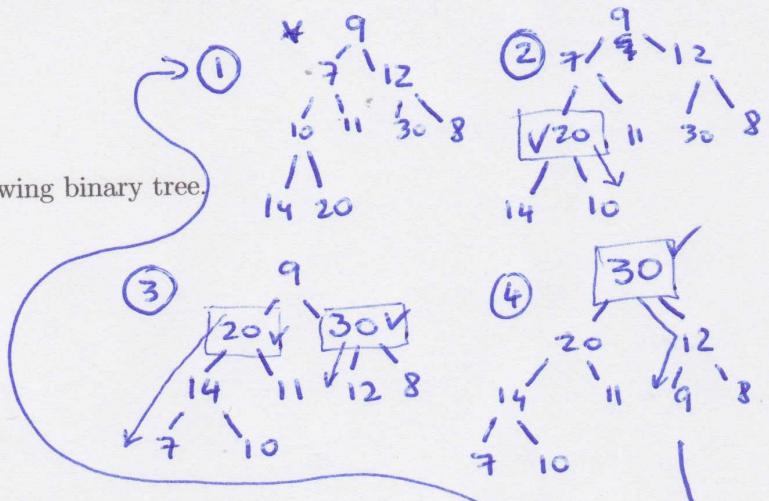
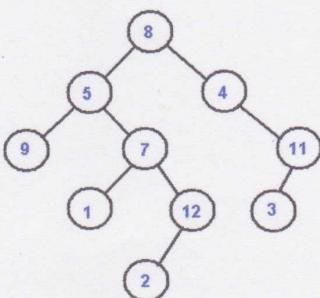
Rear = 1

index i	0	1	2	3	4
value $Q[i]$	A		X	Y	B

Question 10 [3 points] Using the algorithm seen in class, construct the binary tree that corresponds to the general tree below:



Question 11 [3 points] Consider the following binary tree.



List the numbers stored in the nodes in the following order traversals:

- Pre-order: 8, 5, 9, 7, 1, 12, 2, 4, 11, 3
- Post-order: 9, 1, 2, 12, 7, 5, 3, 11, 4, 8
- In-order: 9, 5, 1, 7, 2, 12, 8, 4, 3, 11

Question 12 [4 points] Consider the following complete binary tree using an array representation. This array needs to be transformed into a **max-heap** via the **bottom-up** Heap construction algorithm seen in class.

index i	1	2	3	4	5	6	7	8	9
value $A[i]$	9	7	12	10	11	30	8	14	20

Show the array after *bottom-up* Heap construction algorithm has been applied. (Use back of pages as draft to simulate this algorithm.)

index i	1	2	3	4	5	6	7	8	9
value $A[i]$	30	20	12	14	11	9	8	7	10

Question 13 [4 points] Consider the **in-place** Heapsort algorithm. After the first phase of this algorithm (build heap), the array is as follows:

index i	1	2	3	4	5	6
value $A[i]$	12	11	9	8	5	7

Show the state of the array A after the first 2 `removeMax()` operations in the execution of Heapsort.

index i	1	2	3	4	5	6
$A[i]$ after first <code>removeMax()</code>	11	8	9	7	5	12
$A[i]$ after second <code>removeMax()</code>	9	8	5	7	11	12

$12 \leftrightarrow 7$
 DOWNHEAP 7
 $11 \leftrightarrow 5$
 DOWNHEAP
 HEAP SORTED PART

Question 14 [4 points] We want to sort in increasing order the following sequence S , implemented using an array A , using **in-place Insertion Sort**.

index i	1	2	3	4	5	6	7	8	9	10
value $A[i]$	9	7	3	8	4	12	5	15	2	11

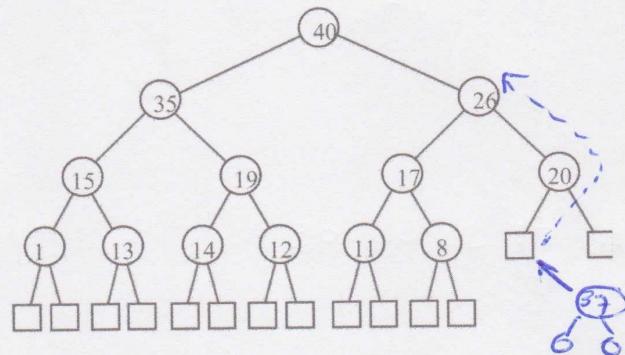
1st Step → 1st element is inserted

2nd Step → 2nd element is inserted

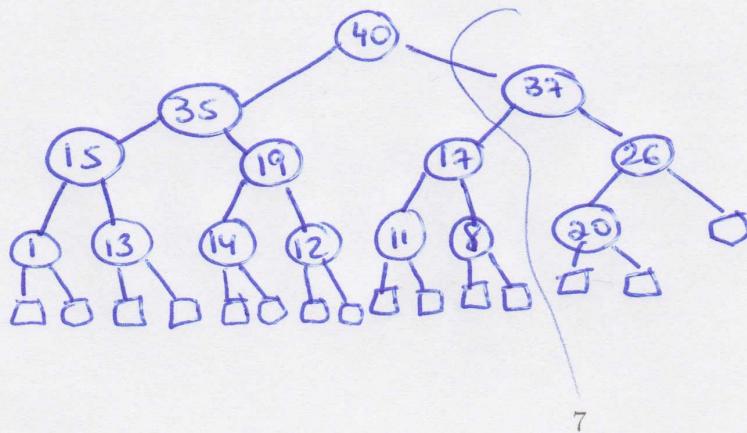
The array will be composed of a sorted part (growing after each step) followed by an unsorted part (shrinking after each step). **Provide the contents** of the array A after the 3rd and 4th steps of the sort (i.e., after inserting the 3rd and the 4th element), and **circle the sorted part** of the array. Of course, you need to first simulate the 1st and 2nd steps using draft space, but they do not need to be shown here.

index i	1	2	3	4	5	6	7	8	9	10
$A[i]$ after step 3	3	7	9	8	4	12	5	15	2	11
$A[i]$ after step 4	3	7	8	9	4	12	5	15	2	11

Question 15 [2 points] Consider the following max-heap:



Show the heap after insertItem(37):



Question 16 [2 points] Consider the abstract data type **Deque** (double-ended queue) with the following methods:

- `void insertFirst(E e)`: insert `e` at the beginning of the Deque.
- `void insertLast(E e)`: insert `e` at the end of the Deque.
- `E removeFirst()` removes and returns the first element of the Deque; an error occurs if the Deque is empty.
- `E removeLast()` removes and returns the last element of the Deque; an error occurs if the Deque is empty.
- `E first()`: returns the first element of the Deque; an error occurs if the Deque is empty.
- `E last()`: returns the last element of the Deque; an error occurs if the Deque is empty.
- `int size()`: returns the number of elements stored in the Deque.
- `boolean isEmpty()`: returns true if and only if the Deque is empty.

Adapt the Deque as a Stack abstract data type and fill in the table below:

or:

Stack method	Deque Method
<code>void push(E e)</code>	<code>InsertLast(e)</code>
<code>E pop()</code>	<code>removeLast()</code>
<code>boolean isEmpty()</code>	<code>isEmpty()</code>
<code>E top()</code>	<code>last()</code>

*insertFirst()
removeFirst()
first()*

Question 17 [4 points]

Consider the **array representation** of a **complete binary tree** where the nodes of the tree are represented by positions $1, 2, \dots, n$ of an array A which stores its elements. Write a method “**void rightmostDescendant(int i)**” that swaps the element stored at position i of the array A with the element of its right-most descendant in the tree.

You may give your algorithm in pseudocode or Java-like program.

```
void rightmostDescendant (int i) {  
    int j = i;  
    while (2*j + 1 <= n) {  
        j = 2*j + 1  
    }  
    temp = A[i];  
    A[i] = A[j];  
    A[j] = temp;  
}
```

Question 18 The following two parts consider a **binary tree** Abstract Data Type that provides the following methods:

method	Description
parent(p)	returns the position of the node that is a parent of node p or null if p is the root
leftChild(p)	returns the position of the left child of p or null if p has no left child.
rightChild(p)	returns the position of the right child of p or null if p has no right child.
isInternal(p)	returns true if and only if p is an internal node.
isExternal(p)	returns true if and only if p is an external node.
size()	returns the number of nodes stored in the tree.

For parts A and B, you may give the required algorithms in pseudocode or Java program excerpt. You may use the methods listed above.

- A. [4 points] Design an algorithm for the following operation of a binary tree T :
`inorderNext(q)`: return the position visited after q in an in-order traversal of T (or `null` if q is the last node to be visited).
- B. [4 points] Write an $O(n)$ algorithm that returns the number of external nodes in a binary tree T with n nodes. Your algorithm will be invoked by calling `numExternal(r)`, where r is the root node of the tree.

Hint: You may implement a recursive algorithm.

```
inorderNext(v) {

    //First try to go right 1 than left all the way
    if(v.right != null){
        next = v.right;
        while (next.left != null)
            next = next.left;
        return next;
    }
    else {
        next = v;
        while(true){ //If you can, go up till you came from the left
            if(next.parent == null) return null;
            if(next.parent.left == next) return next.parent;
            next = next.parent;
        }
    }
```

(blank page to continue answers)

(B)

RECURSIVE PROGRAM:

```
int numExternal (p) {  
    if (p = NULL) then RETURN 0;  
    ELSE if isExternal(p) THEN RETURN 1;  
    ELSE RETURN numExternal(leftChild(p)  
                           + numExternal(rightChild(p))  
}
```

MUST care for Root r:

numExternal(r)