

# Game Theory: Lecture (Week 14)

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## Remaining Course Schedule

Wk	Date	Online (Coursera)	Offline	Notes
14	12/1	1st-half Wk7 video	45min meeting	<b>Quiz 4</b>
15	12/8	2nd-half Wk7 video	75min meeting	
16	12/15	-	-	<b>Final</b>

Two remaining office hours

- ▶ From 3:00 to 4:30PM on Dec 1, 2023
- ▶ (Extended) From 2:00 to 5:00PM on Dec 8, 2023

# Review of Quiz 4

# Cooperative Game Theory

- ▶ So far, we have considered non-cooperative games. Now we switch our gears.
- ▶ Given a set of agents, a cooperative game (or coalitional game) defines how well each group (or coalition) of agents can do for itself.
- ▶ We are not concerned with:
  - ▶ how agents make individual choices within a coalition;
  - ▶ how they coordinate within a coalition.
  - ▶ Thus, we won't talk about individuals' strategies.
- ▶ Instead, we take the payoffs to a coalition as given.
- ▶ Forget about the term “game”. Regard cooperative game as something like an “allocation rule” or “agreeable principle.”

# Coalitional game with transferable utility

## Definition

**A coalitional game with transferable utility** is a pair  $(N, v)$ , where  $N$  is a finite set of players, indexed by  $i$ ; and  $v : 2^N \rightarrow \mathbb{R}$  associates with each coalition  $S \subseteq N$  a real-valued payoff  $v(S)$  that the coalition's members can distribute among themselves. We assume that  $v(\emptyset) = 0$ .

## Definition

A game  $G = (N, v)$  is **superadditive** if for all  $S, T \subset N$ , if  $S \cap T = \emptyset$ , then  $v(S \cup T) \geq v(S) + v(T)$ .

## Coalitional game with transferable utility: Example

Suppose there are three students,  $N \equiv \{1, 2, 3\}$ . They can form a team to work on an assignment.

- ▶ There are eight ( $= 2^3$ ) possible teams:  $\emptyset$ ,  $\{1\}$ ,  $\{2\}$ ,  $\{3\}$ ,  $\{1, 2\}$ ,  $\{1, 3\}$ ,  $\{2, 3\}$ , and  $\{1, 2, 3\}$ .
- ▶ Imagine that each subset has a corresponding value (say, team performance), for example,  $v(\{1\}) = 60$  and  $v(\{1, 3\}) = 75$ .

A coalitional game with transferable utility is 
$$\left( N, \begin{pmatrix} v(\emptyset) \\ v(\{1\}) \\ v(\{2\}) \\ v(\{3\}) \\ v(\{1, 2\}) \\ v(\{1, 3\}) \\ v(\{2, 3\}) \\ v(\{1, 2, 3\}) \end{pmatrix} \right).$$

## Coalitional game with transferable utility: Example

Suppose there are three students,  $N \equiv \{1, 2, 3\}$ . They can form a team to work on an assignment.

- Q1) If  $v(\{1\}) = 60$ ,  $v(\{3\}) = 25$ , and  $v(\{1, 3\}) = 80$ , is the game superadditive?
- Q2) Does  $v(\{1\}) = -5$  violate the superadditivity?
- Q3) Suppose  $v(\emptyset) = v(\{1\}) = v(\{2\}) = v(\{1, 2\}) = 0$ ,  
 $v(\{3\}) = v(\{1, 3\}) = v(\{2, 3\}) = v(\{1, 2, 3\}) = 80$ . Is the game superadditive?

Superadditivity rules out the possibility of 'obviously destructive' coalitions, for example, teams with troll members.

# Shapley value

- ▶ The **Shapley Value** allocates the value of a group according to marginal contribution calculations.
- ▶ Many fair ways of appreciating each member's contribution.  
(Imagine you spent 10 hours for a team project when you are free, and your teammate spent 3 hours while he's hospitalized. Professor told the project scores should be divided within a team. What's the fair way of allocating the scores?)
- ▶  $\psi(N, v) = (\psi_1(N, v), \psi_2(N, v), \dots, \psi_n(N, v))$  is a function that maps a game to each player's value.
- ▶ Shapley value is one way of defining  $\psi(N, v)$ . There could be other ways of allocating value.
- ▶ Let's consider some axioms for  $\psi(N, v)$ .



# Shapley value

## Three definitions

- ▶  $i$  and  $j$  are *interchangeable* if  $v(S \cup \{i\}) = v(S \cup \{j\}) \forall S \subset N$ .
- ▶  $i$  is a *dummy player* if  $v(S \cup \{i\}) = v(S) \forall S \subset N$ .
- ▶  $(N, v)$  is *additively separable* to  $(N, v_1)$  and  $(N, v_2)$  if  $v(S) = v_1(S) + v_2(S) \forall S \subset N$ .

## Axiom (Symmetry)

For any  $v$ , if  $i$  and  $j$  are interchangeable, then  $\psi_i(N, v) = \psi_j(N, v)$ .

## Axiom (Dummy Player)

For any  $v$ , if  $i$  is a dummy player, then  $\psi_i(N, v) = 0$ .

## Axiom (Additivity)

If  $(N, v)$  is additively separable to  $(N, v_1)$  and  $(N, v_2)$ , then  $\psi_i(N, v) = \psi_i(N, v_1) + \psi_i(N, v_2)$ .

## Shapley value

Given a coalitional game  $(N, v)$ , the Shapley value of player  $i$  is:

$$\psi_i(N, v) = \sum_{S \subseteq N \setminus \{i\}} \frac{s!(n-s-1)!}{n!} [v(S \cup \{i\}) - v(S)],$$

where  $n = |N|$  and  $s = |S|$ .

It is known that the Shapley value is the unique function  $\psi(N, v)$  that satisfies the Symmetry, Dummy Player and Additivity axioms.

- ▶ Don't be scared of the long mathematical notation!
- ▶ First, list all subsets that do not include  $i$ .
- ▶ Second, for each case, check how much value is added by including  $i$ .  $(v(S \cup \{i\}) - v(S))$
- ▶ Third, calculate the weighted sum of additional values, whose weight is  $\frac{s!(n-s-1)!}{n!}$ .

## Shapley value: Example

Consider again the teamwork example with three students.

Suppose  $v(\emptyset) = 0$ ,  $v(\{1\}) = v(\{2\}) = v(\{3\}) = 20$ ,

$v(\{1, 2\}) = v(\{1, 3\}) = v(\{2, 3\}) = 50$ , and  $v(\{1, 2, 3\}) = 70$ .

What's the Shapley value of player 1?

- ▶ First, list all subsets that exclude 1:  $\emptyset, \{2\}, \{3\}, \{2, 3\}$ .
- ▶ Check the additional value added by player 1.
  - ▶ From  $\emptyset$  to  $\{1\}$ : 20
  - ▶ From  $\{2\}$  to  $\{1, 2\}$ : 30
  - ▶ From  $\{3\}$  to  $\{1, 3\}$ : 30
  - ▶ From  $\{2, 3\}$  to  $\{1, 2, 3\}$ : 20
- ▶ Now calculate the weighted sum:

$$\begin{aligned} & \frac{0!(3-0-1)!}{3!}20 + \frac{1!(3-1-1)!}{3!}30 + \frac{1!(3-1-1)!}{3!}30 + \frac{2!(3-2-1)!}{3!}20 \\ &= \frac{2}{6}20 + \frac{1}{6}30 + \frac{1}{6}30 + \frac{2}{6}20 = \frac{140}{6} = \frac{70}{3} \end{aligned}$$

## Shapley value: Example 2

A game is called a *simple game* if  $v(S) \in \{0, 1\}$  for all  $S \subset N$ .

Suppose  $v(\emptyset) = v(\{1\}) = v(\{2\}) = v(\{3\}) = 0$  and

$v(\{1, 2\}) = v(\{2, 3\}) = v(\{1, 3\}) = v(\{1, 2, 3\}) = 1$ .

(Think a majority voting.) What is the Shapley value of player 1?

- ▶ First, list all subsets that exclude 1:  $\emptyset, \{2\}, \{3\}, \{2, 3\}$ .
- ▶ Check the additional value added by player 1.
  - ▶ From  $\emptyset$  to  $\{1\}$ : 0
  - ▶ From  $\{2\}$  to  $\{1, 2\}$ : 1
  - ▶ From  $\{3\}$  to  $\{1, 3\}$ : 1
  - ▶ From  $\{2, 3\}$  to  $\{1, 2, 3\}$ : 0
- ▶ Now calculate the weighted sum:

$$\frac{0!(3-0-1)!}{3!}0 + \frac{1!(3-1-1)!}{3!}1 + \frac{1!(3-1-1)!}{3!}1 + \frac{2!(3-2-1)!}{3!}0 = \frac{1}{6} + \frac{1}{6} = \frac{1}{3}$$

- ▶ Thus, in a simple game, the Shapley value equals the probability of being pivotal.

(Check the Coursera videos and problem sets for more examples.)

# Core

- ▶ The Shapley value defined a fair way of dividing the grand coalition's payment among its members.
- ▶ But sometimes smaller coalitions can be more attractive for subsets of the agents, even if they lead to lower value overall.
- ▶ In the prisoner's dilemma, we learned that (C,C) won't be played even if it had greater value overall. Can we consider a concept analogous to Nash equilibrium, say, a subset of the agents that they do not want to deviate from the current coalition?

# Core

## Definition (Core)

A payoff vector  $x$  is in the core of a coalitional game  $(N, v)$  iff

$$\forall S \subseteq N, \sum_{i \in S} x_i \geq v(S) \quad \text{and} \quad \sum_{i \in N} x_i = v(N).$$

The core may not exist. If it exists, it may not be unique. So it is easier to think about the “core set” whose size is 0 (empty), 1 (unique core), or larger.

- ▶ Consider an arbitrary value vector  $x = (x_1, x_2, \dots, x_n)$ .
- 1. List all nonempty subsets of  $N$ .
- 2. List the inequality condition  $(\sum_{i \in S} x_i \geq v(S))$  for each subset.
- 3. Check if there exists  $x$  such that all conditions are satisfied.

## Core: Example

Consider again the teamwork example with three students.

Suppose  $v(\emptyset) = 0$ ,  $v(\{1\}) = v(\{2\}) = v(\{3\}) = 20$ ,

$v(\{1, 2\}) = v(\{1, 3\}) = v(\{2, 3\}) = 50$ , and  $v(\{1, 2, 3\}) = 70$ .

Does the core exist?

1. All nonempty subsets:

$\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}$ .

2. All inequality conditions:

$$\begin{aligned}x_1 &\geq 20, & x_2 &\geq 20, & x_3 &\geq 20 \\x_1 + x_2 &\geq 50, & x_1 + x_3 &\geq 50, & x_2 + x_3 &\geq 50 \\&&&&& x_1 + x_2 + x_3 &= 70\end{aligned}$$

3. You can easily check there is no payoff vector  $x$  that satisfies all inequality conditions.

$x_1 + x_2 \geq 50$ ,  $x_2 + x_3 \geq 50$ , and  $x_1 + x_3 \geq 50$  imply  $x_1 + x_2 + x_3 \geq 75$ .

(If  $v(\{i, j\}) = 40$  instead of 50, the core exists.)

## Core: Example 2

Consider again a simple game with three players.

Suppose  $v(\emptyset) = v(\{1\}) = v(\{2\}) = v(\{3\}) = 0$  and  
 $v(\{1, 2\}) = v(\{2, 3\}) = v(\{1, 3\}) = v(\{1, 2, 3\}) = 1$ .

Does the core exist?

1. All nonempty subsets:

$\{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}$ .

2. All inequality conditions:

$$\begin{aligned}x_1 &\geq 0, & x_2 &\geq 0, & x_3 &\geq 0 \\x_1 + x_2 &\geq 1, & x_1 + x_3 &\geq 1, & x_2 + x_3 &\geq 1 \\&&&& x_1 + x_2 + x_3 &= 1\end{aligned}$$

3. The core doesn't exist.



# Shapley value and Core

## Definition (Convex game)

A game  $(N, v)$  is **convex** if  $v(S \cup T) \geq v(S) + v(T) - v(S \cap T)$   
 $\forall S, T \subset N$ .

- ▶ It is known that every convex game has core, and the Shapley value of a convex game is in the core.
- ▶ Check that the teamwork example used in this note is not a convex game. Hint: Consider  $S = \{1, 2\}$  and  $T = \{2, 3\}$ .  
(If  $v(\{1, 2\})$ ,  $v(\{2, 3\})$ ,  $w(\{1, 3\})$  were to be 40, then it would've been a convex game.)