Chapter 8: Switching

Outline

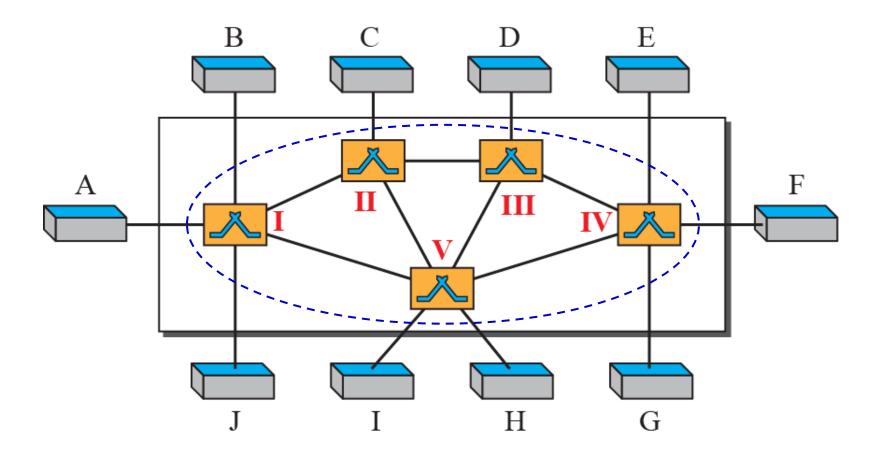
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8-1 INTRODUCTION

A <u>network</u> is a set of <u>connected devices</u>. Whenever we have multiple devices, we have the problem of how to <u>connect</u> them to make <u>one-to-one communication</u> possible. The solution is <u>switching</u>.

A switched network consists of a <u>series of</u> <u>interlinked nodes</u> called <u>switches</u>. Switches are devices capable of creating <u>temporary</u> <u>connections</u> between two or more devices linked to the switch.

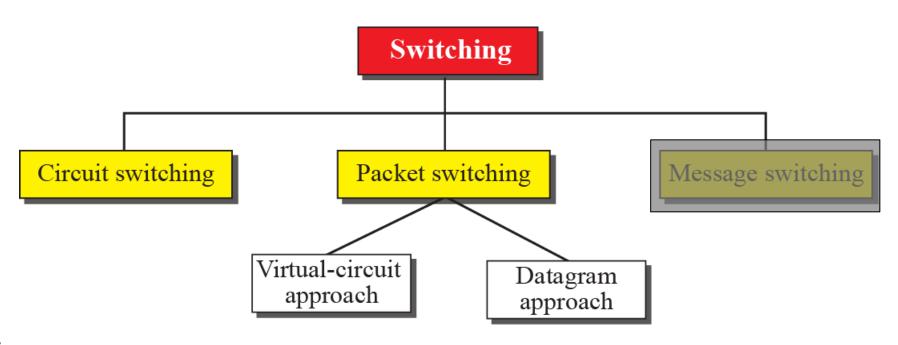
Figure 8.1: Switched network



Switches I, II, III, IV, V are <u>devices</u> capable of creating <u>temporary connections</u> between <u>two or more devices</u> linked to the switch.

8.1.1 Three Methods of Switching

Traditionally, there are three methods of switching: circuit switching, packet switching and message switching. Packet switching can further be divided into two subcategories: datagram approach and virtual-circuit approach.



8.1.2 Switching and TCP/IP Layers

Circuit switching and packet switching are commonly used today. Message switching has been phased out in general communications but still has networking applications.

Switching can happen at <u>several layers</u> of the TCP/IP protocol suite:

o Physical layer: <u>circuit switching</u>

o Data link layer: packet (frame/cell) switching

o Network layer: <u>packet switching</u>

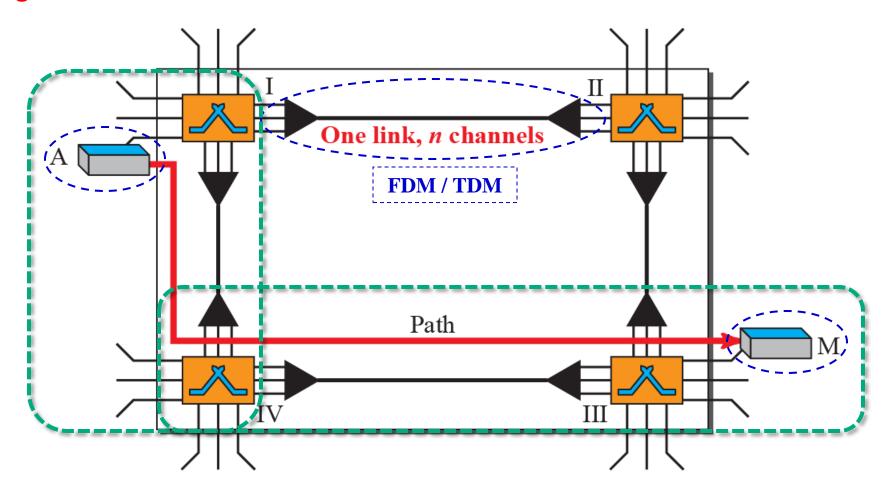
o Application layer: message switching

8-2 CIRCUIT-SWITCHED NETWORKS

A circuit-switched network consists of a <u>set of</u> <u>switches</u> connected by <u>physical links</u>, in which each <u>link</u> is divided into <u>n channels</u>.

A <u>connection</u> between two stations is a <u>dedicated path</u> made of one or more <u>links</u>. Each <u>connection</u> uses a <u>dedicated channel</u> on each link. Each link is normally divided into n channels by using <u>FDM</u> or <u>TDM</u>.

Figure 8.3: A trivial circuit-switched network



In circuit switching, the resources need to be reserved during the <u>setup phase</u>. The resources remain dedicated for the entire duration of the <u>data transfer phase</u> until the <u>teardown phase</u>.

8.2.1 Three Phases

The actual communication in a circuit-switched network requires three phases: (1) connection setup, (2) data transfer and (3) connection teardown:

- (1) Connection Setup: Before the entities can communicate, a dedicated circuit needs to be established. Note that end-to-end addressing (e.g., telephone numbers) is required for creating a connection between the two end systems.
- (2) Data Transfer: After the establishment of the dedicated circuit, the two entities can <u>transfer data</u>.
- (3) Connection Teardown: When one of the entities needs to disconnect, a signal is sent to each switch to <u>release</u> the <u>resources</u>.

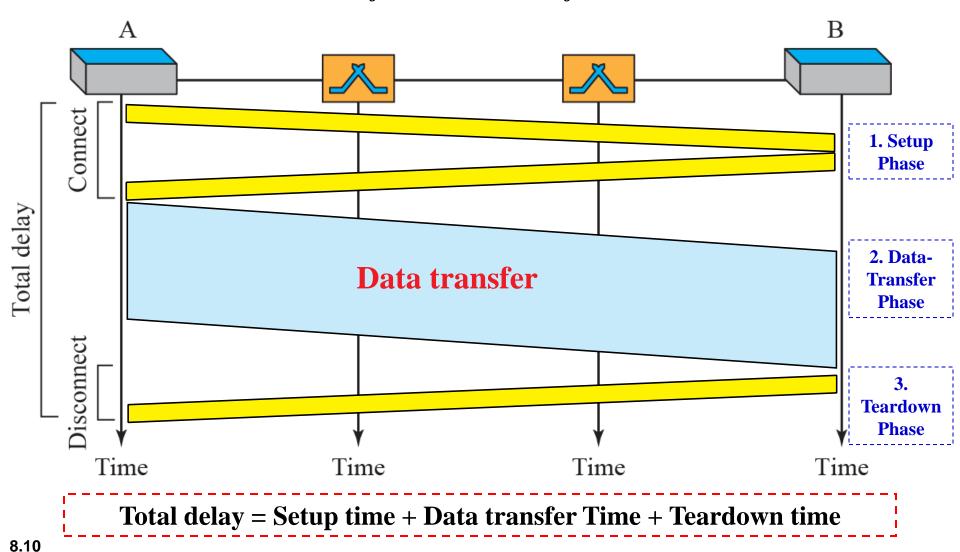
8.2.2 Efficiency

Efficiency: It <u>can be argued</u> that circuit-switched networks are <u>not as efficient</u> as packet- or message-switched networks since <u>resources</u> are <u>allocated</u> for the <u>entire duration</u> of the connection; and <u>unavailable</u> to other connections. Let's consider the following <u>2 scenarios for comparison</u>:

- (a) In a <u>telephone</u> network, people <u>normally terminate</u> the communication when they have finished their conversation.
- (b) In a <u>computer</u> network, a computer <u>can remain</u> <u>connected</u> to another computer even if there is no activity for a long time → other connections are deprived.

Figure 8.6: Delay in a circuit-switched network

While a circuit-switched network normally has <u>low efficiency</u>, the <u>delay</u> is <u>minimal</u>. During data transfer, the data are not delayed at each switch as the resources are allocated for the duration of the connection.



Example

Let us use a circuit-switched network to connect eight telephones in a small area. Communication is through 4 kHz voice channels. We assume that each link uses FDM to connect a maximum of two voice channels. The bandwidth of each link is then 8 kHz.

Solution

Telephone 1 is connected to telephone 7; 2 to 5; 3 to 8; and 4 to 6. The switch controls the connections as shown below:

Note this is not a fullyconnected setup, e.g., 4 cannot
communicate with 7/8.

Circuit-switched network

Circuit-switched network

5

0
4
kHz kHz

0
4
kHz kHz

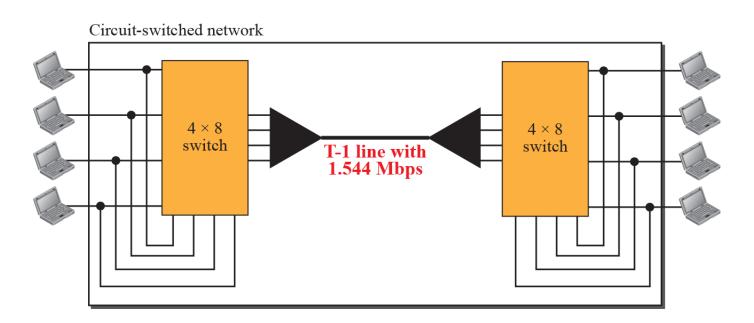
6
kHz kHz

Example

Consider a circuit-switched network that connects computers in two remote offices of a private company. The offices are connected using a T1 line leased from a communication service provider. There are two 4×8 (4 inputs, 8 outputs) switches in this network.

Solution

For each 4 x 8 switch, four output ports are folded back into the input ports to allow communication between computers in the same office. The other four output ports allow communication between the two offices as shown below:



8-3 PACKET SWITCHING

In data communications, we need to send messages from one end system to another. If the message is going to pass through a <u>packet-switched</u> network, the <u>message</u> needs to be divided into <u>packets</u> of fixed or variable size (size of the packet is determined by the network and the governing protocol).

In packet switching, there is no resource reservation for a packet: no reserved bandwidth on the links, no scheduled processing time for each packet. Resources are allocated on-demand on a first-come, first served (FCFS) basis.

8.3.1 Datagram Networks

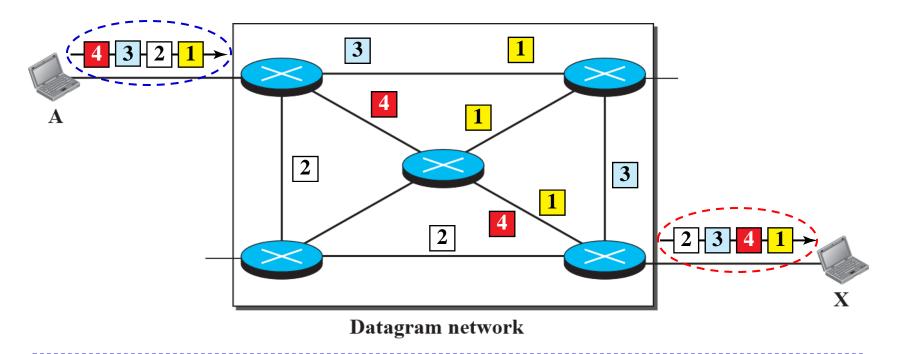
We have two types of packet-switched networks: datagram networks and virtual-circuit networks.

In a datagram network, each <u>packet</u> is treated <u>independently</u> of all others. Even if a packet is part of a multipacket transmission, the network treats it as though it existed alone. Packets in this approach are referred to as <u>datagrams</u>.

Datagram switching is normally done at the <u>network layer</u> (note that the term "switching" is not exclusively referencing only layer 2 "switches").

Figure 8.7: A Datagram network with four switches (routers)

Example of datagram switching to deliver four packets from A to X. All four packets belong to the same message, but may travel different paths from A to X.

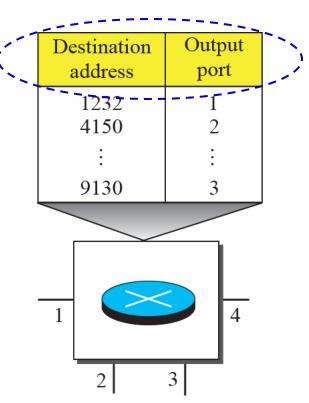


Datagram networks are sometimes referred to as <u>connectionless</u> <u>networks</u>. The term <u>connectionless</u> means that the <u>switch</u> does not keep <u>information</u> about the <u>connection state</u>: each <u>packet</u> is treated the <u>same</u> by a switch regardless of its source or destination.

Figure 8.8: Routing table in a datagram network

The <u>destination address</u> in the header of a packet in a datagram network <u>remains</u> the <u>same</u> during the <u>entire journey</u> of the packet.

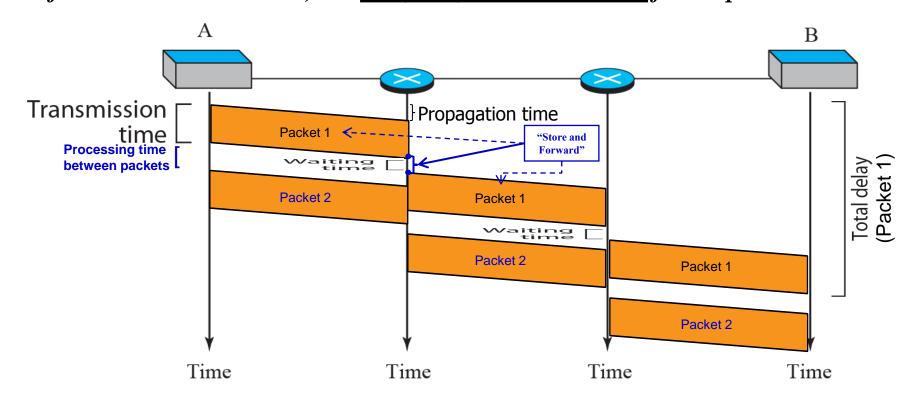
A switch in a datagram network uses a routing table (based on the destination address) where the destination addresses and corresponding forwarding output ports are recorded.



The routing tables are dynamic and updated periodically.

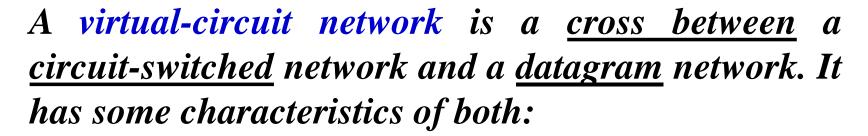
Figure 8.9: Delays in a datagram network

Efficiency of a datagram network is better than that of a circuit-switched network as resources are allocated only when there are packets to be transferred. However, there may be longer delays in a datagram network as each packet may experience a wait time at each switch before being forwarded. In addition, the delay may not be the same for all packets.



Total delay $(Packet 1) = 3 \times Transmission time + 3 \times Propagation time + 2 \times Waiting time (Note that the transmission, propagation and waiting times may not be the same at each switch)$

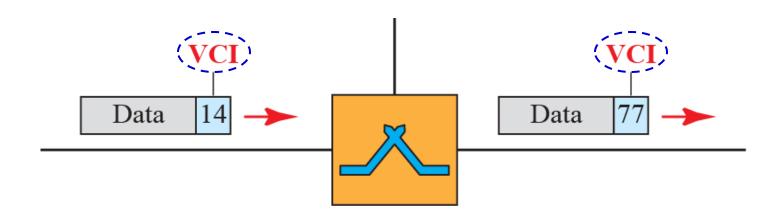
8.3.2 Virtual-Circuit Networks



- 1) There are <u>setup and teardown</u> phases in addition to the <u>data transfer</u> phase (circuit-switched).
- 2) <u>Resources</u> can be <u>allocated</u> during the <u>setup</u> phase (circuit-switched) or <u>on-demand</u> (datagram).
- 3) Data are <u>packetized</u> (datagram) but address in header only has <u>local jurisdiction</u>, i.e., only knows next switch/channel, not end-to-end jurisdiction.
- 4) All packets follow the <u>same path</u> established during the connection (circuit-switched)
- 5) Normally implemented in the <u>data link</u> layer. (Physical layer: circuit-switched; Network layer: datagram)

Figure 8.11: Virtual-circuit identifier

In a virtual-circuit network, two types of <u>addressing</u> are involved: <u>global</u> and <u>local</u>. The <u>global address</u> is used to create a <u>virtual-circuit identifier</u> (VCI), usually a small number, that only has <u>local switch scope</u>, i.e., when a frame arrives at a switch, it has a VCI; when it leaves, it has a different VCI.



As in a circuit-switched network, a source and destination goes through three phases in a virtual-circuit network: <u>setup phase</u>, <u>data transfer phase</u> and <u>teardown phase</u>.

Figure 8.14: Setup request in a virtual-circuit network

In the <u>setup phase</u>, the source and destination use their global addresses to help the switches make table entries for the connection. Two steps are required: (a) the <u>setup request</u> and (b) the acknowledgment.

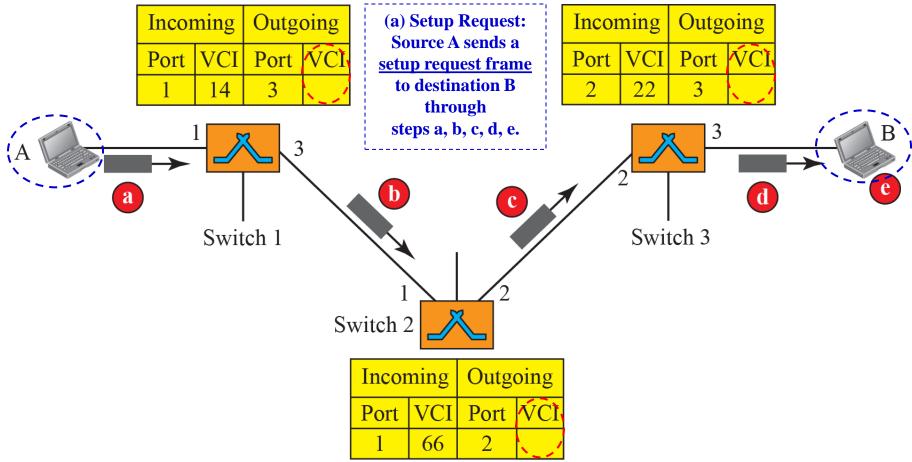


Figure 8.15: Setup acknowledgment in a virtual-circuit network

After the setup request, in the acknowledgment step, a special frame, called the <u>acknowledgment frame</u>, completes the entries in the switching tables. This process creates a <u>virtual circuit</u> between the source and the destination.

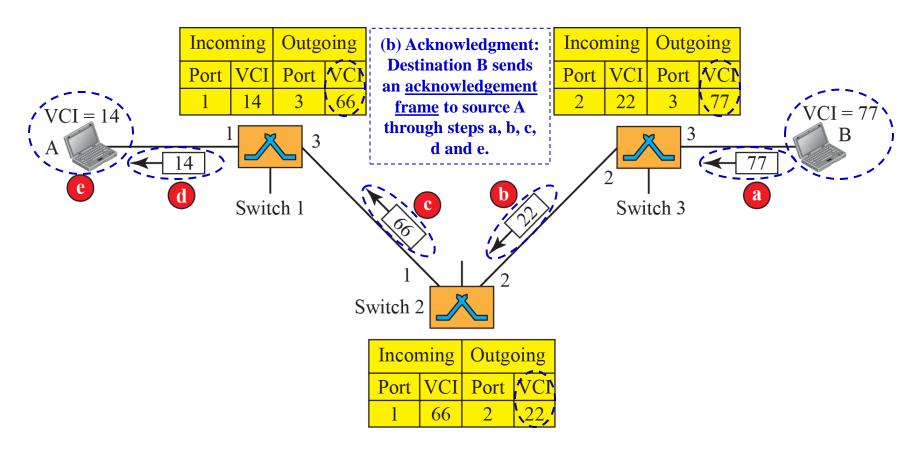
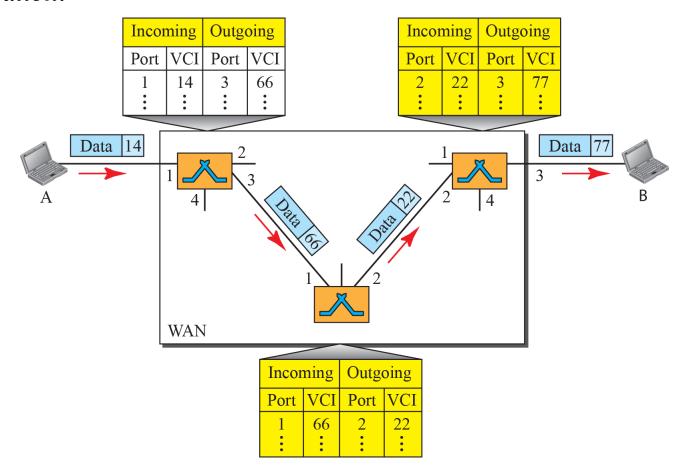


Figure 8.13: Source-to-destination data transfer / Teardown

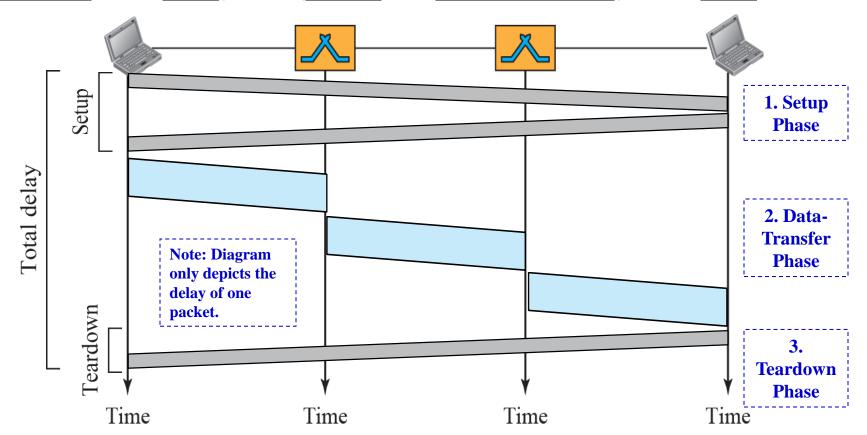
The <u>data transfer phase</u> is active until the source sends all its frames to the destination.



After A sends all its frames to B, it sends a special frame called a <u>teardown</u> request in the <u>teardown phase</u>. B responds with a <u>teardown confirmation</u> and <u>all switches delete</u> the <u>corresponding entries</u> from their tables.

Figure 8.16: Delay in a virtual-circuit network

In virtual-circuit switching, all <u>packets</u> belonging to the <u>same source and</u> <u>destination</u> travel the <u>same path</u>. <u>Resource reservation</u> can be made during the <u>setup phase</u> or on-demand during the <u>data-transfer phase</u>. The <u>delay</u> for <u>each packet</u> is the <u>same</u> for the <u>former</u> and <u>may be different</u> for the <u>latter</u>.



Total delay = Setup time $+ 3 \times Transmission$ time $+ 3 \times Propagation$ time + Teardown time

(Assumption: resource reservation is made during the setup phase, otherwise include additional delay for individual packets at each switch)