

# Dynamic Programming: Introduction

(Chapter 8)

# Welcome, dear old friends:

## The Fibonacci numbers

- 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, ...
- Each number is the sum of the previous two:
  - $\text{fib}(0) = 1$
  - $\text{fib}(1) = 1$
  - $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$
- How many can we compute?

# DEMO

THE BASIC ALGORITHM:

**fib (n):**

if  $n < 2$

return  $n$  //i.e. fib(0) is 0 and fib(1) is 1

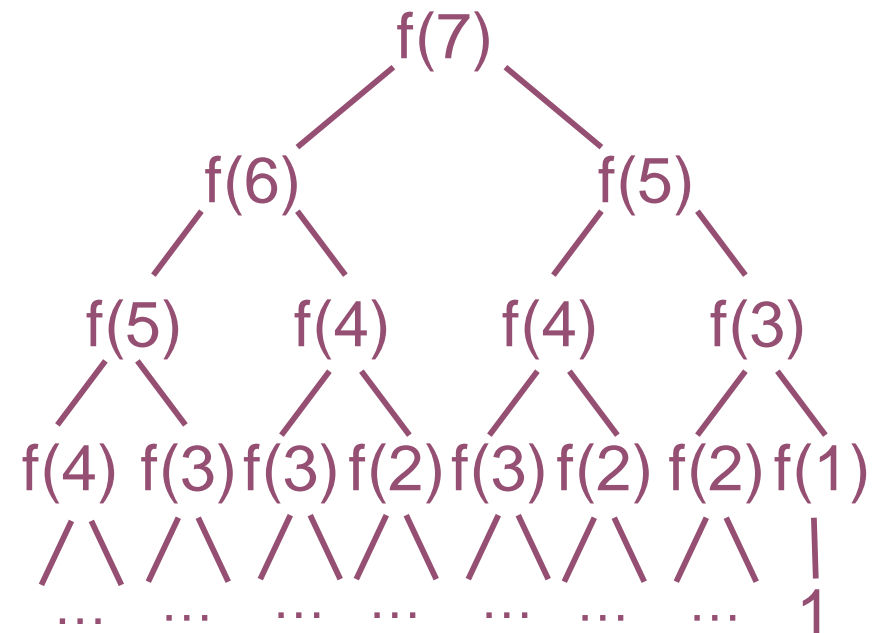
else

return  $\text{fib}(n-1) + \text{fib}(n-2)$

# Fibonacci numbers: Why you so slow?

Execution tree:

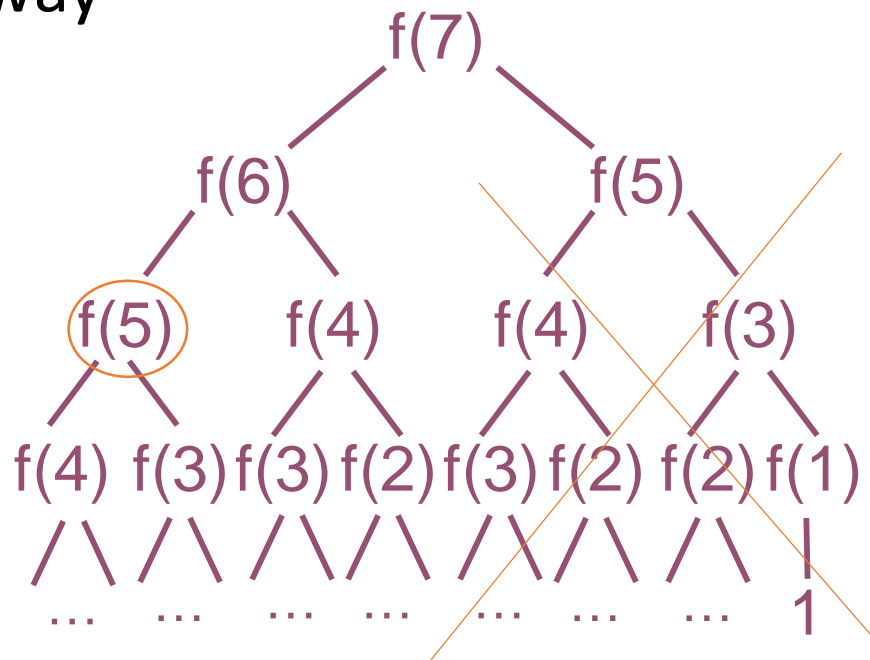
```
fib (n):  
  if n < 2  
    return n;  
  else  
    return = fib(n-1) + fib(n-2)
```



$F(n)$  takes exponential time to compute.

# Space-time trade-off

- Augment the algorithm by *remembering* the results you get along the way

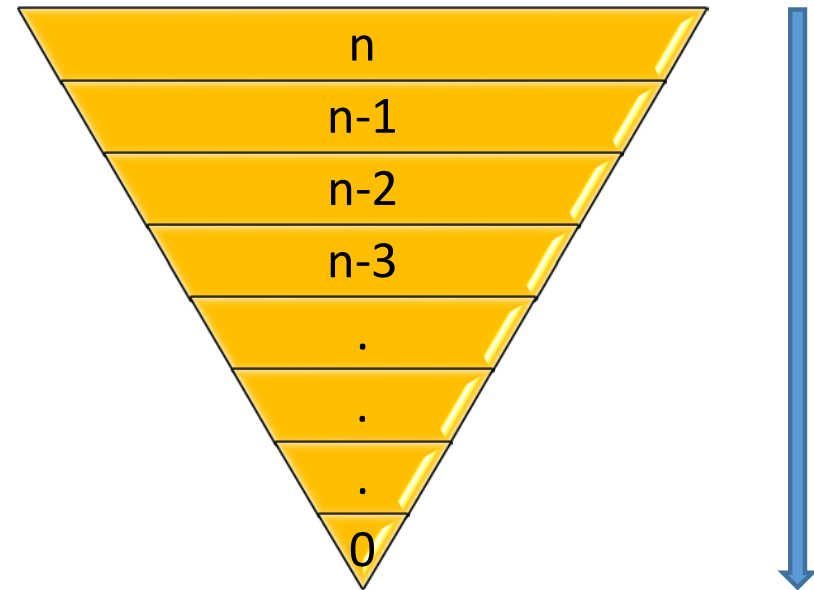


memo

|  |  |  |  |  |   |  |  |
|--|--|--|--|--|---|--|--|
|  |  |  |  |  | 5 |  |  |
|--|--|--|--|--|---|--|--|

# Fibs, top-down

```
fib (n) {  
    if memo[n] exists, return it  
    if n < 2  
        return n  
    else  
        f = fib(n-1) + fib(n-2)  
        memo[n] = f  
    return f  
}
```



top-down (Recursive)

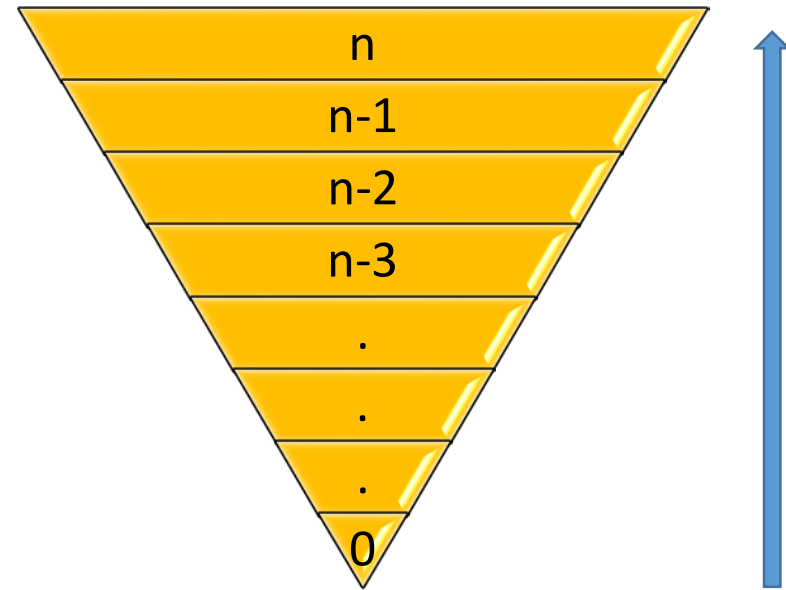
|      |   |   |   |     |                 |                 |               |
|------|---|---|---|-----|-----------------|-----------------|---------------|
| memo | 0 | 1 | 1 | ... | <i>fib(n-2)</i> | <i>fib(n-1)</i> | <i>fib(n)</i> |
|------|---|---|---|-----|-----------------|-----------------|---------------|

Efficiency:

- time:  $O(n)$
- space: Needs an array size  $O(n)$

# Fibs, bottom-up

```
fib (n) {  
    memo[0] = 0;  
    memo[1] = 1;  
    for i ← 2 to n do  
        memo [i] = memo[i-1] + memo[i-2]  
    return memo[n]  
}
```



bottom-up

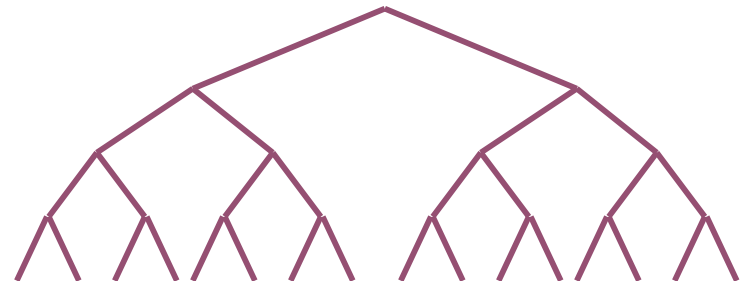
|      |   |   |   |       |                 |                 |               |
|------|---|---|---|-------|-----------------|-----------------|---------------|
| memo | 0 | 1 | 1 | . . . | <i>fib(n-2)</i> | <i>fib(n-1)</i> | <i>fib(n)</i> |
|------|---|---|---|-------|-----------------|-----------------|---------------|

Efficiency:

- time:  $O(n)$
- space: Needs an array size  $O(n)$

# Dynamic programming

- Key point: *remembering* recursively-defined solutions to sub-problems and using them to solve the problem
- Very much like divide-and-conquer ... but store the solutions to sub-problems for possible reuse.
- A good idea if many of the sub-problems are repeats





# Dynamic programming overview

- **Step 1:**
  - Decompose problem into simpler sub-problems
- **Step 2:**
  - Express solution in terms of sub-problems
- **Step 3:**
  - Use table to compute optimal value bottom-up
- **Step 4:**
  - Find optimal solution based on steps 1-3

# Dynamic programming examples

- Fibonacci numbers
- Robot Coin Collecting
- Transitive Closure (Warshall)
- All Pairs Shortest Paths (Floyd)